

The ret instruction:

```
ret == popl    %eip
```

the ret instruction accesses the address of the top of the stack from esp register. Call that address as x.

It transfers the content of 4 bytes of the stack top (i.e. $M[x : x + 3]$) to eip register.

```
1) eip <- M[esp : esp + 3]
```

```
2) esp <- esp + 4
```

```
Step 1 + Step 2 == popl %eip
```

```
#-----
```

For learning function definition and function call in systematic manner, we will divide C function in 4 progressive stages of the evolution:

Step 1 : A C function without any parameter, without return value,
without local variables, without nested function calls.

Step 2: A C function with parameters and return value. but no local variables
and no nested function calls

Step 3: A C function

- 1) with parameters
- 2) with return value
- 3) with local variables
- 4) without nested function calls

Step 4: <FULL FLEDGED C FUNCTION>

- 1) function with parameters
- 2) function with return value
- 3) function with local variables
- 4) function with nested calls (having all 4 capacities)

```
#-----
```

Step 1:

step1.c

```
int num1, num2, result;

void test_function(void){
    // global variable manipulation
    num1 = 10;
    num2 = 20;
    result = num1 + num2;
}
```

```
void main(){
    test_function();
}
```

```
.section .bss
    .comm num1, 4, 4
    .comm num2, 4, 4
    .comm result, 4, 4

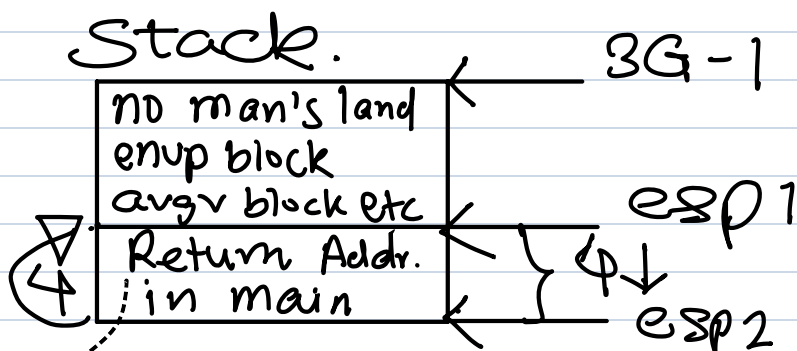
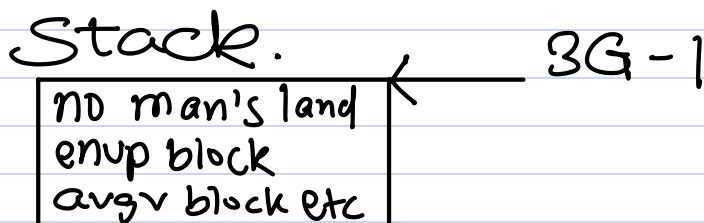
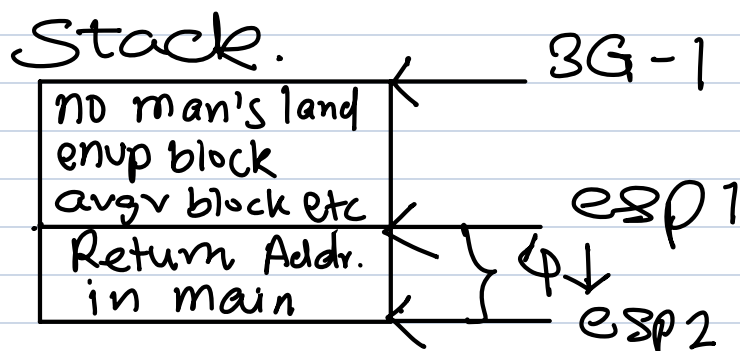
.section .text
.globl main
.type main, @function
```

```
main:
    call test_function --- eip1
    ret
```

```
.globl test_function
.type test_function, @function
```

```
test_function:
    movl $10, num1 --- eip2
    movl $20, num2 --- eip3
    movl num1, %eax --- eip4
    addl num2, %eax --- eip5
    movl %eax, result --- eip6
    ret --- eip7
```

eip8



```

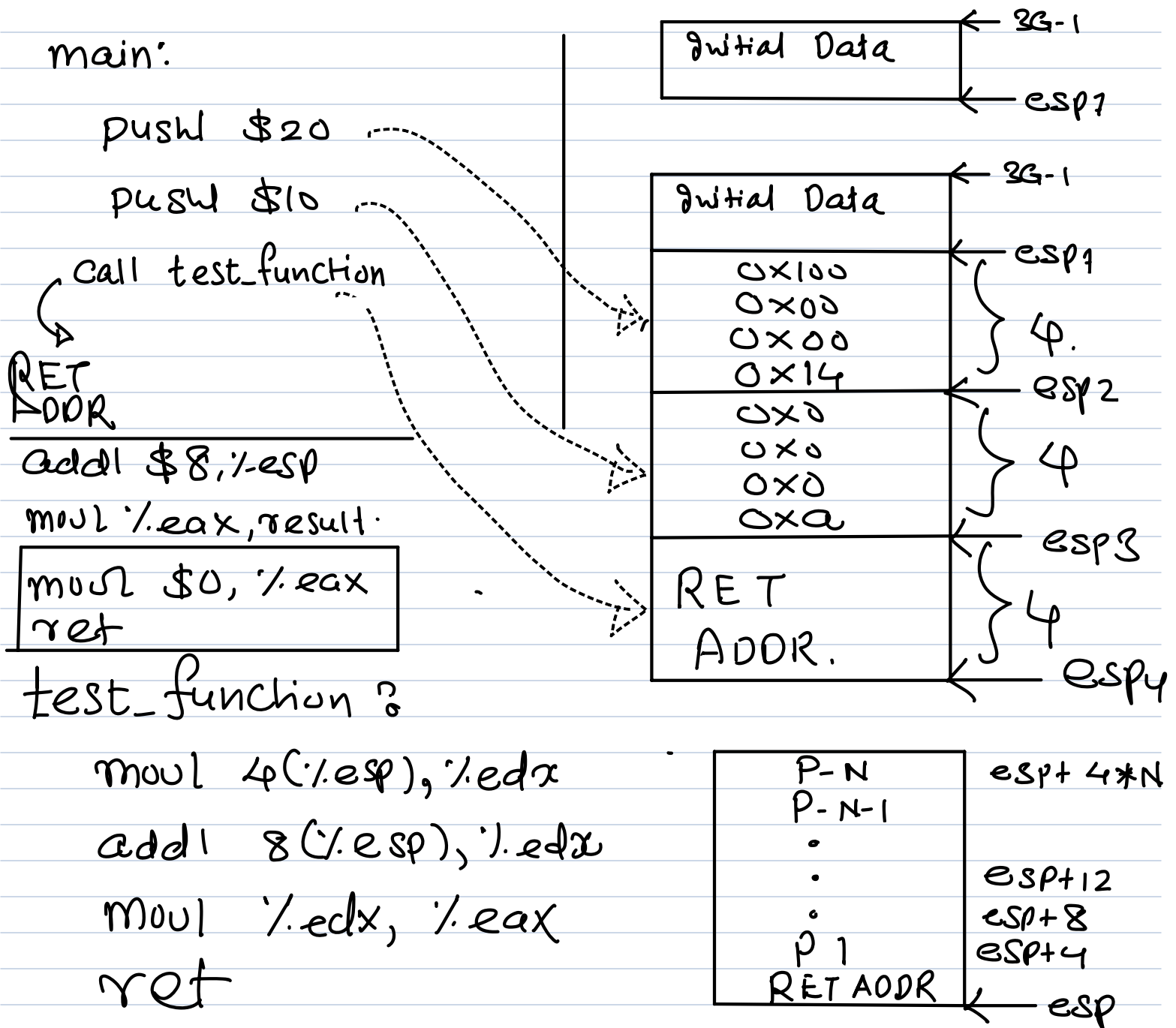
int test_function(int n1, int n2){
    return (n1+n2);
}

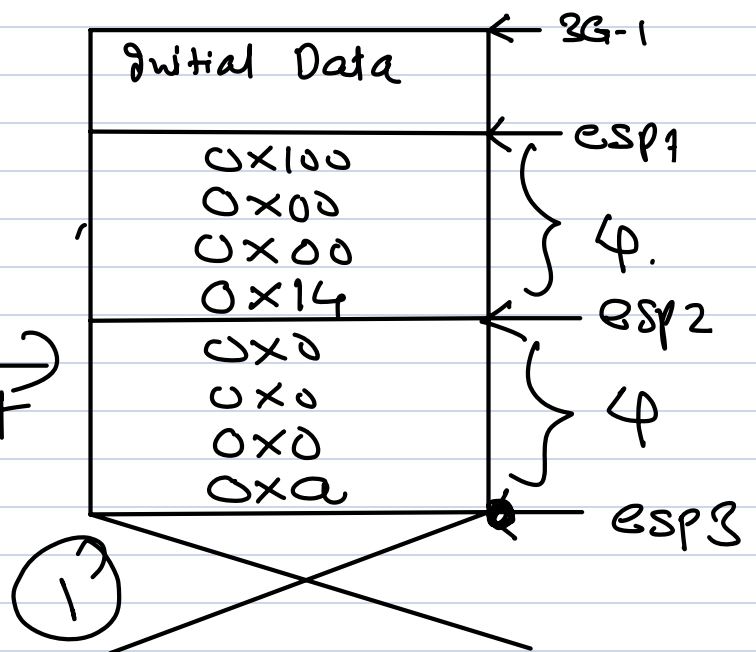
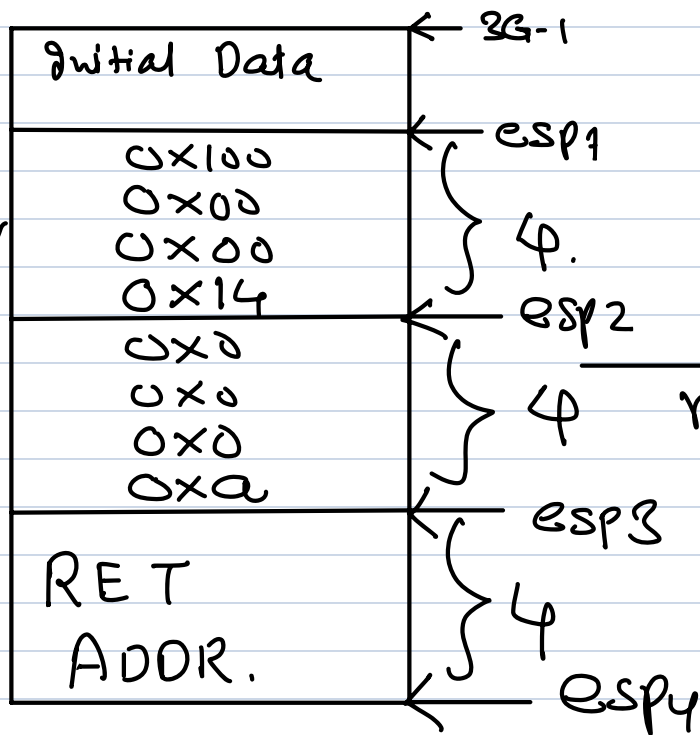
```

```

int main() {
    result = test_function(10, 20);
}

```

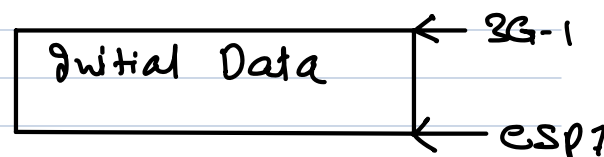
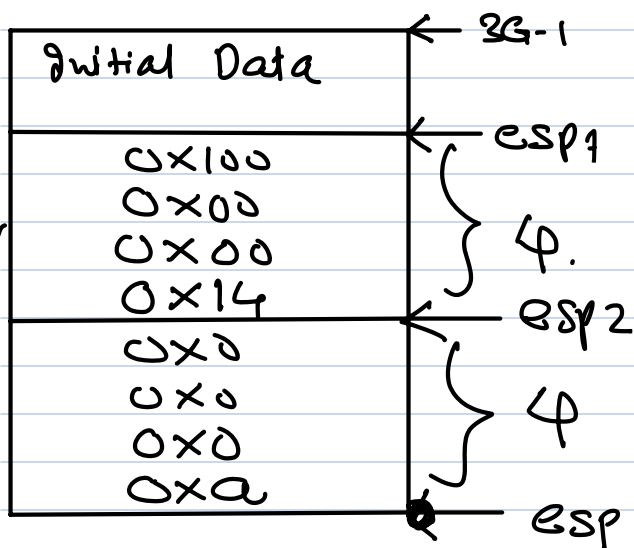




pushl \$20

pushl \$10

(2) call test_function
 RET ADDR. → addl \$8, %esp



addl \$8, %esp

pushl \$20

pushl \$10

call test_function

addl \$8, %esp

movl %eax, result

result = test_function(10, 20);

ret = fun_name($T_1 P_1, \dots, T_n P_n$);

pushl P_n

.

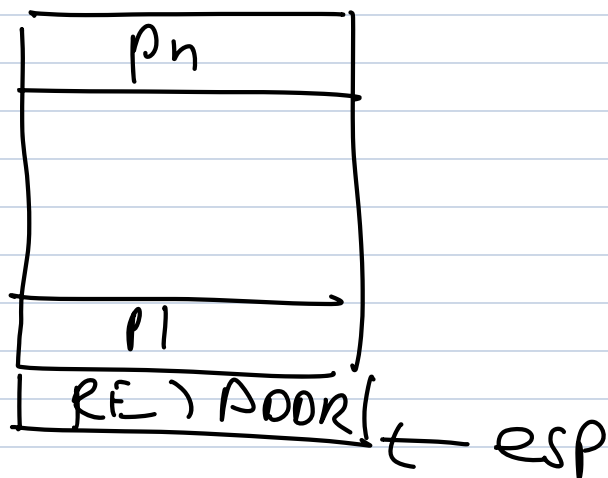
.

pushl P_1

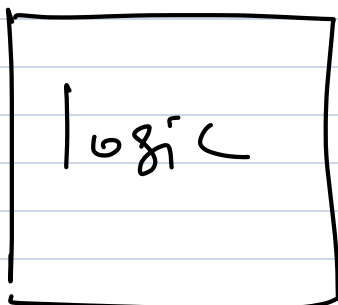
call fun_name

addl \$4n, %esp

movl %eax, ret



fun_name: $4 \times k(\%esp)$ $1 \leq k \leq n$.



movl \rightarrow %eax
ret

XWindow | AS7 | graph \rightarrow adj
 \rightarrow BFS, DFS | RB tree

```
int test_function(int n1, int n2)
{
    int rs1, rs2;
    rs1 = n1 + n2;
    rs2 = n1 - n2;
    return (rs1 * rs2);
}
```

```
int rs;
int main(void) {
    rs = test_function(10, 20);
    return (0);
}
```

• section .bss.

• comm rs, 4, 4.

• section .text.
main:

pushl \$20

pushl \$0

call test_function

R.A. \hookrightarrow addl \$8, %esp

movl %eax, rs

init. data

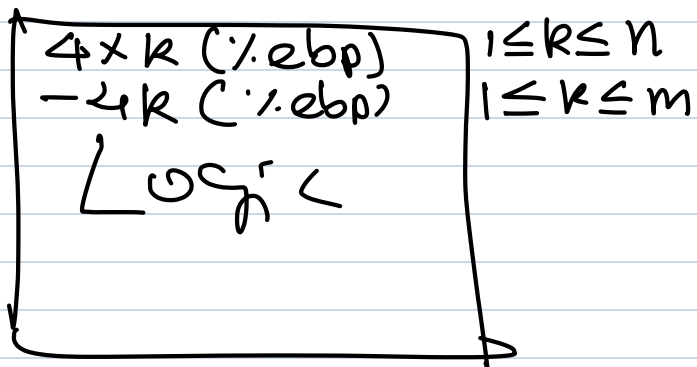
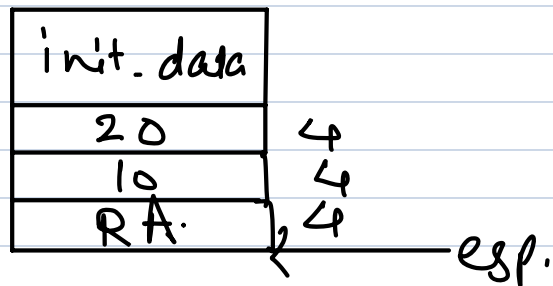
init. data
20
10
RA.

4
4

test_function:

movl %esp, %ebp

subl \$8, %esp.

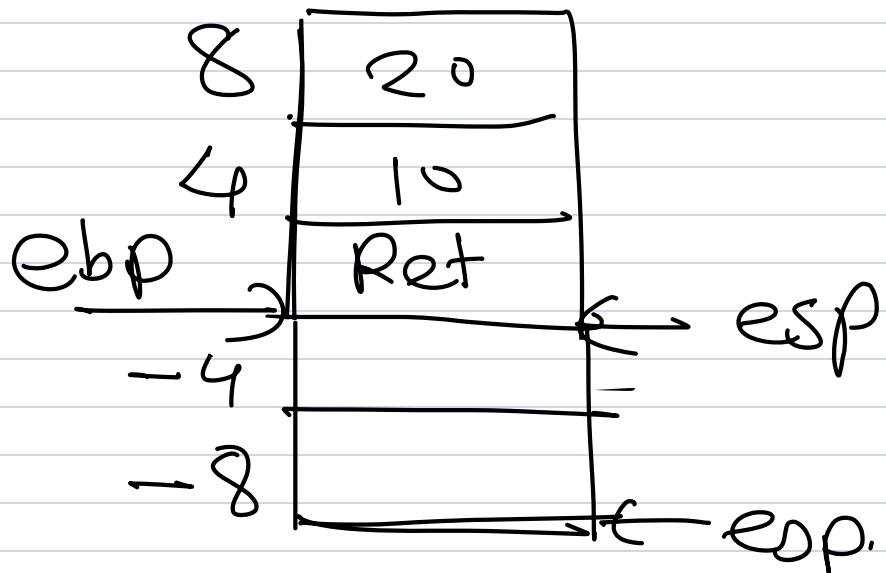
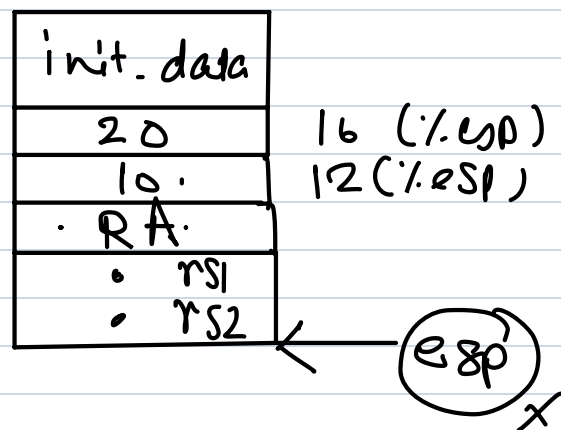


$O(\%esp) \rightarrow R.A.$

$4(\%esp) = p1$

$8(\%esp) = p2$

movl — %eax
 movl %ebp, %esp
 ret



$ret = fun_name(p_1, \dots, p_n);$

$Tr \ fun_name(T, p_1, \dots, p_n)$

$\{ \quad // m \text{ local variables}$

logic.

$\} \quad return \ V_{Tr};$

Caller.

pushl p_n
pushl p_{n-1}

pushl p_1

call fun_name

addl $\$4n, \%esp$

movl $\%eax, ret$

Callee:

$fun_name:$

movl $\%esp, \%ebp$

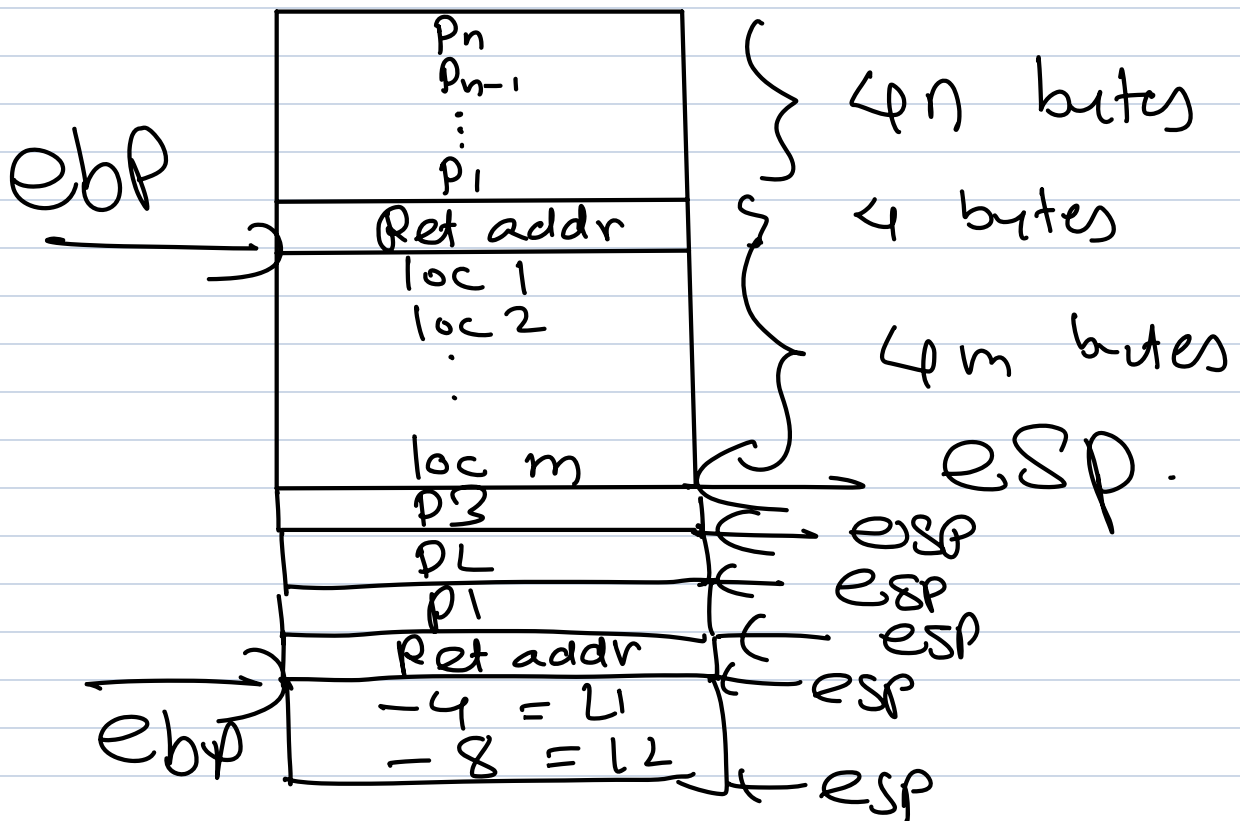
subl $\$4m, \%esp$

$4k(\%ebp) =$
kth Param

$-4k(\%ebp) =$
kth local

logic

movl $_, \%eax$
movl $\%ebp, \%esp$
ret



```

pushl %ebp
movl %esp, %ebp
subl $4m, %esp

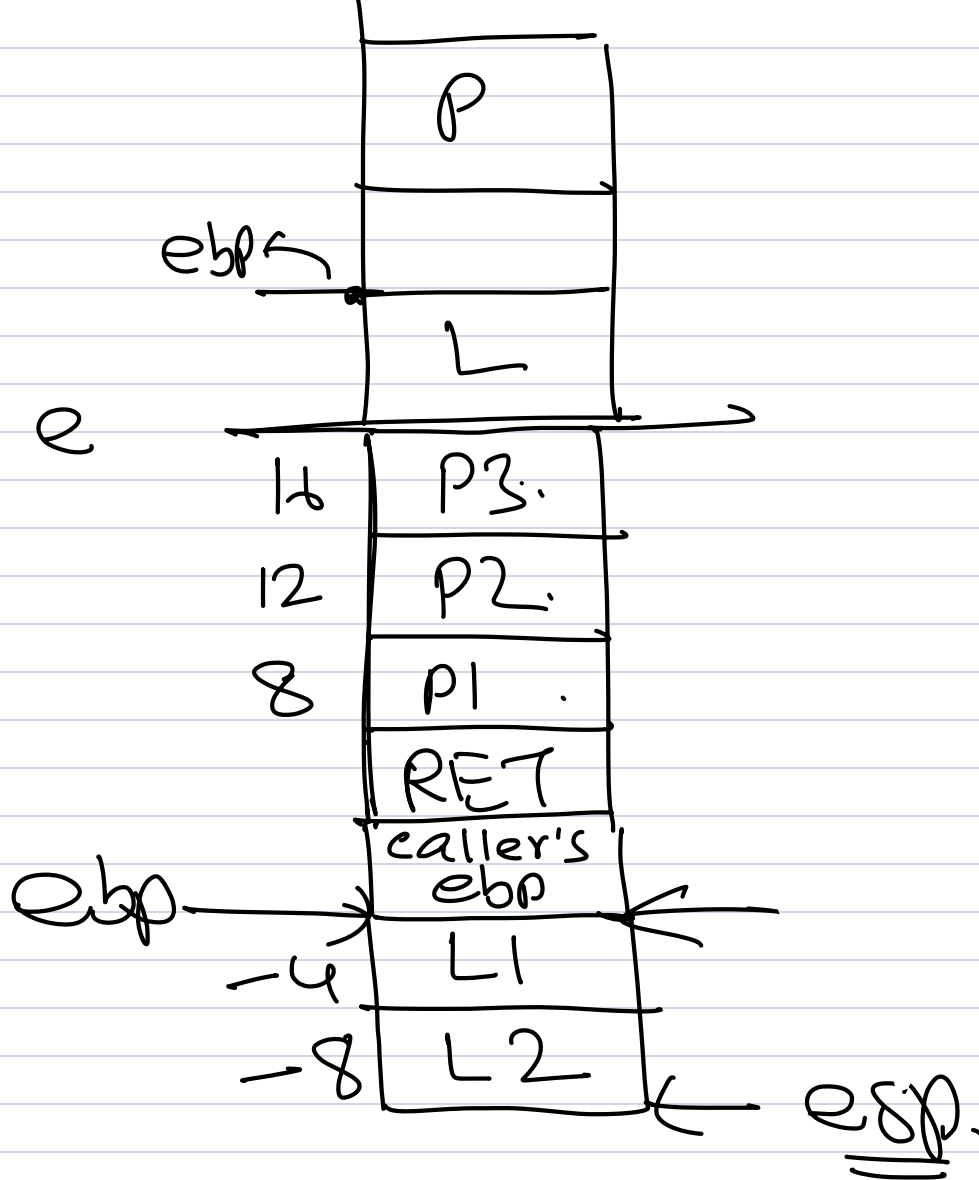
```

$4k + 4(\%ebp) =$
 k^{th} Param
 $-4k(\%ebp) =$
 k^{th} local
 logic

```

movl —, %eax
movl %ebp, %esp
popl %ebp
ret

```



```

int f (int a, int b, int c)
{
    int m, n;
    [ ]
    return (-);
}

```

उपोद्घात = Prologue.

विषयप्रवेश	=	<input type="text"/>
विषयवस्तु,	=	<input type="text"/>

उपसंघर्ष = Epilogue.