**WEEK:: 03** 





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WEEK:: 03 DAY: 01 DATE: 01-05-2023

## TIME & SPACE COMPLEXITY

It is defined as the time taken by the algorithm as our input size.

Time Complexity ⇒ Input Size :: The time complexity depends on the size of our input. Then we can compare our different algorithms to see which one is better.

#### **Time Complexity != Time Taken**

O(Worst Case) :: It indicates the maximum time required by an algorithm for all input values. It represents the worst case of an algorithm's time complexity.

Omega(Best Case) :: It indicates the minimum time required by an algorithm for all input values. It represents the best case of an algorithm's time complexity.

**Theta(Average Case)** ::It indicates the average bound of an algorithm. It represents the average case of an algorithm's time complexity.

#### **Exp :: 01::** Print n natural numbers?

#### Exp :: 02:: Find num in Array {1, 6, 85, 8, 8, 5, }

### **How to choose Algorithm::**

We choose any algorithm according to Time and Complexity which algorithm needs less.

#### **#For Nested Loop:**

#### Sum of n num:

Algorithm 1	Algorithm 2
for(int i=0; i<= n; i++) {    Sum = sum + i; } [execute n time for n output] (n==n)	Sum = n(n+1)/2  [execute one time for any output] (1==n)
O(n)	0

```
for(int i=1; i<=10; i++)
{
    cout<<"Hello CoderArmy";
}

If change (i<=n):: now depend on input
TC = O(n)
```

```
for(int i=1; i<n; i=i*2)
                                 Print input ::
                                 i=1;
                                         "print"
{
                                         "print"
                                 i=2;
                                                      i=i*2
   cout<<"Print";</pre>
                                         "print"
                                 i=4;
}
                                 i=8;
                                        "print"
                                 i=n;
                                         "print" :: i= 1
                                                             2
                                                                    4
                                                                            8 - - -
                                                            2^1
                                                                   2^2
                                                                          2^3
                                                                                    2^k
                                                                                           k=k+1
                                 n = 2^k = \log n = \log 2^k = \log n = k \log 2
                                 k = log n / lg 2 = log_2 n TC = O(log_2 n)
```

```
for(int j=1; i<=n; i++) 

{ for(int j=1; j<=n; j=j+i; cout<<" print"; i=2; j=1 to n; print n/2 [jump 2 number] i=3; j=1 to n; print n/3 [jump 3 number] i=4; j=1 to n; print n/4 [jump 4 number] i=n; j=1 to n; print n/n [jump n number] Sum = (n + n/2 + n/3 + n/4 + ..... + n/n) n[1 + ½ + ½ + ½ + ½ + ½ + ..... + 1/n] = lne n {Harmonic Series} TC = (n \log e n)
```



WEEK:: 03 DAY: 02 DATE: 02-05-2023

# **TIME & SPACE COMPLEXITY PART 2**

Time Complexity: It states that the amount of time taken by an algorithm to run as a function of the length of the input.

```
Choose a better algorithm: 1. O(log10 (x) ) 2. O(\sqrt{n})
Put the value of n = Very high value(small value to high value).
log2N=k \Rightarrow 2 N = k \Rightarrow 2^k = N; || log2 8 = 2^x = 8\Rightarrow x = 3;

n = 10 100 10000
```

```
logn= 1 2 4 It is better.

\sqrt{n} = 3 10 100
```

**Better Time complexity algorithm: Table** 

O(1),  $O(\log n)$ ,  $O(\sqrt{n})$ , O(n),  $O(n \log n)$ ,  $O(n^3)$ ,  $O(2^n)$ , O(n!)

Space Complexity :: It states that the amount of space taken by an algorithm to run a function of the length of input.

- 1. Auxiliary space.
- 2. Total space complexity.

Auxiliary Space	Total space Complexity
It depends on the size of input but does not include the presence of input space which already takes space computer memory.	It depends on both spaces which are obtained from input.
Ignore :: Geven Input size space.	

#### # Print 1 to n::

```
int n;
cin>>n;
for(int i=1; i<=10; i++)
cout<<l<";

Which one takes space::
here, n & i take space.
So, space complexity O(1).
Not depend on input.
```

# Print Array input: {2, 4, 6, 7}; output: {20, 40, 60, 70};

```
#include<iostream>
using namespace std;

int main()
{
   int n;
```

```
cout<<"Enter Input Size: ";</pre>
cin>>n;
                                               Ignore the input loop.
int Input_arr[n];
                                               Time complexity::
cin>>Input_arr[n];
                                               First loop = n
                                               Second loop = n
                                                = n + n = 2n = O(n)
cin>>Input arr[i];
                                               Space Complexity:
int arr[n];
                                               Total space Complexity:
                                               Input loop = 1 + n
arr[i] = Input_arr[i]*10;
                                               Array copy loop = n + 1
                                               Print loop = n + 1
cout<<arr[i]<<" ";
                                               = 1 + n + n + 1 + n + 1 = O(n)
                                               Auxiliary Space = O(n)
```

WEEK:: 03 DAY: 03 DATE: 03-05-2023

## **BINARY SEARCH**

Binary Search is defined as a searching algorithm used in a sorted array by repeatedly dividing the search interval in half.

**Given array**:: arr[] = { 3, 7, 10, 14, 17, 18, 25} **Find**: 7

	3	7	10	14	17	18	25
Index	0	1	2	3	4	5	6

First, We will go to the middle of the array. (0 + 6)/2 = 3

3 7 10 14 17 18 25
--------------------

Check 7 & 14 is equal or not. Then check the left side.



Again go to the middle. And find 7

7

Then it's equal to 7.

#### Example02::

3	10	17	18	24	29	32	37	50
Index 0	1	2	3	4	5	6	7	8

Target: 3

We go to the middle of the array.

Check target & 24 is equal or not. (24>3) find left side

3	10	17	18
---	----	----	----

Again, find the middle point.



Check target & 10 is equal or not. (10>3) find left side

3

Finally we get a target.

### **Binary Search Algorithm:-**

- 1. Find the middle element (Index)
- 2. Check if the value are equal or not

```
a. Both are equal [ target = arr[ mid ] cout << "Find Target";</li>
b. Target value > arr[ mid ] Move to right side
c. Target value < arr[ mid ]</li>
```

Move to left side

#Code::

```
include <iostream>
using namespace std;
int main()
    int nums[8] = \{5, 6, 8, 9, 11, 15, 17, 21\};
    int mid, target = 15, start = 0, end = 7;
    while (start <= end)</pre>
        mid = (start + end) / 2;
        if (nums[mid] == target)
            cout << mid;
        else if (nums[mid] > target)
            end = mid - 1;
            start = mid + 1;
    cout << "-1";
```

**Time Complexity:** 

```
It's going to be half by half. n/2 \rightarrow n/4 \rightarrow n/8 O(log2 n) Space Complexity:
```

O(1) {Auxiliary}, O(n) {Total space complexity}

```
include <iostream>
using namespace std;
int main()
   int v[8] = \{ 1, 2, 2, 2, 4, 5, 8, 11 \}; // arr = v
   int x = 2; // target
   long long mid, start = 0, end = 7, left_index, right_index;
 while (start <= end)</pre>
   {
       mid = (start + end) / 2;
       if (v[mid] == x)
        {
           left_index = mid;
           end = mid - 1;
        else if (v[mid] > x)
            end = mid - 1;
           start = mid + 1;
    }
   start = 0;
    end = 7;
   while (start <= end)</pre>
        mid = (start + end) / 2;
        if (v[mid] == x)
        {
           right index = mid;
           start = mid + 1;
        else if (v[mid] > x)
            end = mid - 1;
            start = mid + 1;
```

```
cout<<"left_index: "<<left_index<<" "<<"right_index:
"<<right_index;
    return 0;
};</pre>
```

## Exp:: 03:: Search Insert Position << LeetCode >>

```
class Solution {
public:
    int searchInsert(vector<int>& arr, int target) {
        int mid, start = 0, end = arr.size() - 1, index;
        while (start<=end)</pre>
        {
            mid = (start + end)/2;
             if (arr[mid] == target)
            return mid;
             else if (arr[mid]<target)</pre>
                 start = mid+1;
                 index = mid+1;
             }
             else
                 end = mid -1;
                 index = mid;
             }
        }
        return index;
    }
};
```

# Exp:: 04:: Peak Index in a Mountain Array << LeetCode >>

#### Own Code ::

```
include <iostream>
using namespace std;
int main()
    int arr[5] = \{5, 6, 5, 4, 1\};
    int start = 0, mid, end = 4;
    while (start <= end)</pre>
        mid = end + (start - end) / 2;
        if (arr[mid] > arr[mid + 1] && arr[mid] > arr[mid - 1])
            cout<< mid;</pre>
        else if (arr[mid] > arr[mid - 1] && arr[mid] < arr[mid + 1])</pre>
            start = mid + 1;
            end = mid - 1;
```

WEEK :: 03 DAY: 04 DATE: 04-05-2023

### BINARY SEARCH IN DETAIL

Explain :: Arr[7] = { 2, 3, 6, 7, 5, 1, 0 }; [ Find index of pick value ]

Arr[7] = { 2, 3, 6, 7, 5, 1, 0 } index = 0 1 2 3 4 5 6

```
Pick element of Arr 7; And index = 3;
Satisfy condition = Arr[i-1] < Arr[i] > Arr[i+1] 6 < 7 > 5

Increasing order : Arr[i-1] < Arr[i] < arr[i+1] = find right side :- start = mid+1

Decreasing order : Arr[i+1] > Arr[i] > arr[i+1] = find left side :- end = mid -1

Code :: Last day Quation
```

Exp:: 01 :: Find missing element

arr	0	1	3	4	5
Index	0	1	2	3	4

1st, go middle of arr = 3 Check [index == element]

Miss match [ index != element ] ,then go left side but store index because if it is a small miss match element.

arr	0	1	2	5	8
Index	0	1	2	3	4

1st, go middle of arr = 3 Check [index == element]

Not Miss match [index == element], then go right side but store index because if it is a small miss match element.

```
#include <iostream>
using namespace std;

int main()
{
   int n;
   cin >> n;
   int arr[n];
```

```
Exp:: 01 :: Find how many times rotate the arr. Arr[ 6 ] = { 4, 5, 6, 7, 2, 3 };

Explain :: [ Pivot Method ]

Coside pivot value = 4;

1st, go mid = 6; compare with pivot value

arr[mid] > pivot value :: store index and go right side :: start = mid +1

arr[mid] < pivot value :: store index and go left side :: start = mid - 1
```

# << Rotation >>

```
//User function template for C++
class Solution{
public:
      int findKRotation(int arr[], int n) {
          // code here
           if (arr[0] < arr[n - 1])</pre>
          return 0;
       int low = 0, high = n - 1;
       while (low < high)</pre>
       {
          int mid = (low + high)/2;
          if (arr[low] <= arr[mid] && arr[high] <= arr[mid])</pre>
              low = mid + 1;
           }
          else
           {
```

```
high = mid;
}
return low;
}
```

```
Exp:: 02 ::Search in a Rotated Array << <u>GEEKSFORGEEKS</u> >>

Arr[ 6 ] = {5, 6, 8, 9, 1, 2}; key =
```

```
Index = 0 1 2 3 4 5

1st :: pivot arr[0] = 5;

2nd :: go mid = 8;
    if( arr[mid] == key)
        Return mid;

3rd:: compare arr[mid] != pivot (arr[0])
        Case 1:: pivot < key end = mid -1;
        Case 2:: arr[mid]] < key
        Start = mid +1

4th :: arr[mid] < pivot
        But pivot < key end = mid -1;
```

**Next Question** 

Case 1:: key<arr[mid] end = mid-1

Other case start = mid +1

WEEK :: 03 DAY: 04 DATE: 06-05-2023

## **BINARY SEARCH IN DETAIL PART II**

```
9
                                          1
                                              2
             5
                  6
                                                       4
                            Target = 3
Index =
                       2
                                 4
                                                        9
             0
                 1
                            3
                                      5
                                          6
```

```
# Find middle (0 + 9)/2
# check key == arr[middle] return middle
# arr[start <= arr[middle]
                            arr sorted
# 5 <= target <= 9 Not alive, move right
# again find middle 2,
# check target alive between 1 & 2, Not - move right
# again find middle check middle==target and print
   1. if(arr[mid] == key)
         return mid;
   2. else if (arr[start] < arr[mid])
           if(key>= arr[start] && key<= arr[mid])
             End = mid - 1;
            else
             start = mid + 1;
   3. else
        {
             If ( key > arr[mid] && key<= arr[end]
            start = mid +1;
       }
```

```
else if(arr[mid]>arr[start])
             {
                  if(arr[mid]>key && arr[start]<=key)</pre>
                  end = mid -1;
                  else
                  start = mid +1;
             }
             else
             {
                  if(arr[mid] < key && key <= arr[end])</pre>
                  start = mid +1;
                  else
                  end = mid -1;
             }
        return index;
    }
};
```

# Allocate minimum number of pages << GEEKSFORGEEKS >>

```
class Solution
    public:
    //Function to find minimum number of pages.
    int findPages(int arr[], int n, int m)
         //code here
        int start=arr[0], end=0, mid;
        if(m>n)
        return -1;
        for(int i=0; i<n; i++)</pre>
             end = end+arr[i];
             if(start<arr[i])</pre>
             start = arr[i];
         }
        int ans;
        while (start<=end)</pre>
             mid =(start+end)/2;
             int sum=0, count=1;
             for (int i=0; i<n; i++)</pre>
                 sum+=arr[i];
                 if (sum>mid)
                  {
                      count++;
                      sum = arr[i];
             }
             if (count<=m)</pre>
```

```
{
     end = mid-1;
     ans = mid;
}
else
     start = mid+1;
}
return ans;
}
```

### Aggressive Cows << <u>GEEKSFORGEEKS</u> >>

```
// User function Template for C++
class Solution {
public:
    int solve(int n, int k, vector<int> &stalls) {
        // Write your code here
        int start =1, end = stalls[0], mid, minimum = stalls[0];
        for(int i=1; i<n; i++)</pre>
            minimum = min(minimum, stalls[i]);
            end = max(end, stalls[i]);
        };
        sort(stalls.begin(), stalls.end());
        end -= minimum;
        int ans;
        while (start<=end)</pre>
            mid =(start+end)/2;
            int lastPos = stalls[0], count = 1;
            for(int i=0; i<n; i++)</pre>
                 if (stalls[i]-lastPos>=mid)
                     count++;
                     lastPos = stalls[i];
                 if (count==k)
                break;
            }
            if (count==k)
                 ans = mid;
                 start = mid +1;
            else
            end = mid-1;
        return ans;
    }
};
```