



Introduction - Virtualization

Virtualization

- ▶ Software Virtualization
- ▶ Hardware virtualization

Sources

- ▶ <http://www.slideshare.net/alanmcsweeney/an-introduction-to-server-virtualisation>
- ▶ www.cs.usfca.edu/~cruse/cs686s07/lesson19.ppt
- ▶ <http://www.slideshare.net/ACMBangalore/virtualization-tutorial-at-acm-bangalore-compute-2009>
- ▶ <https://www.slideshare.net/sidhushahbaz/handout2o>
- ▶ <https://www.csa.iisc.ac.in/~vg/teaching/E0-253/slides/VirtualMachines.pptx>



Introduction to Server Virtualization

A loose definition

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Virtualization is a framework or methodology of dividing the resources of a computer into multiple execution environments, by applying one or more concepts or technologies such as hardware and software partitioning, time-sharing, partial or complete machine simulation, emulation, quality of service, and many others.

VM is an isolated duplicate of the physical machine

Without virtualization (bare metal)

Application
1

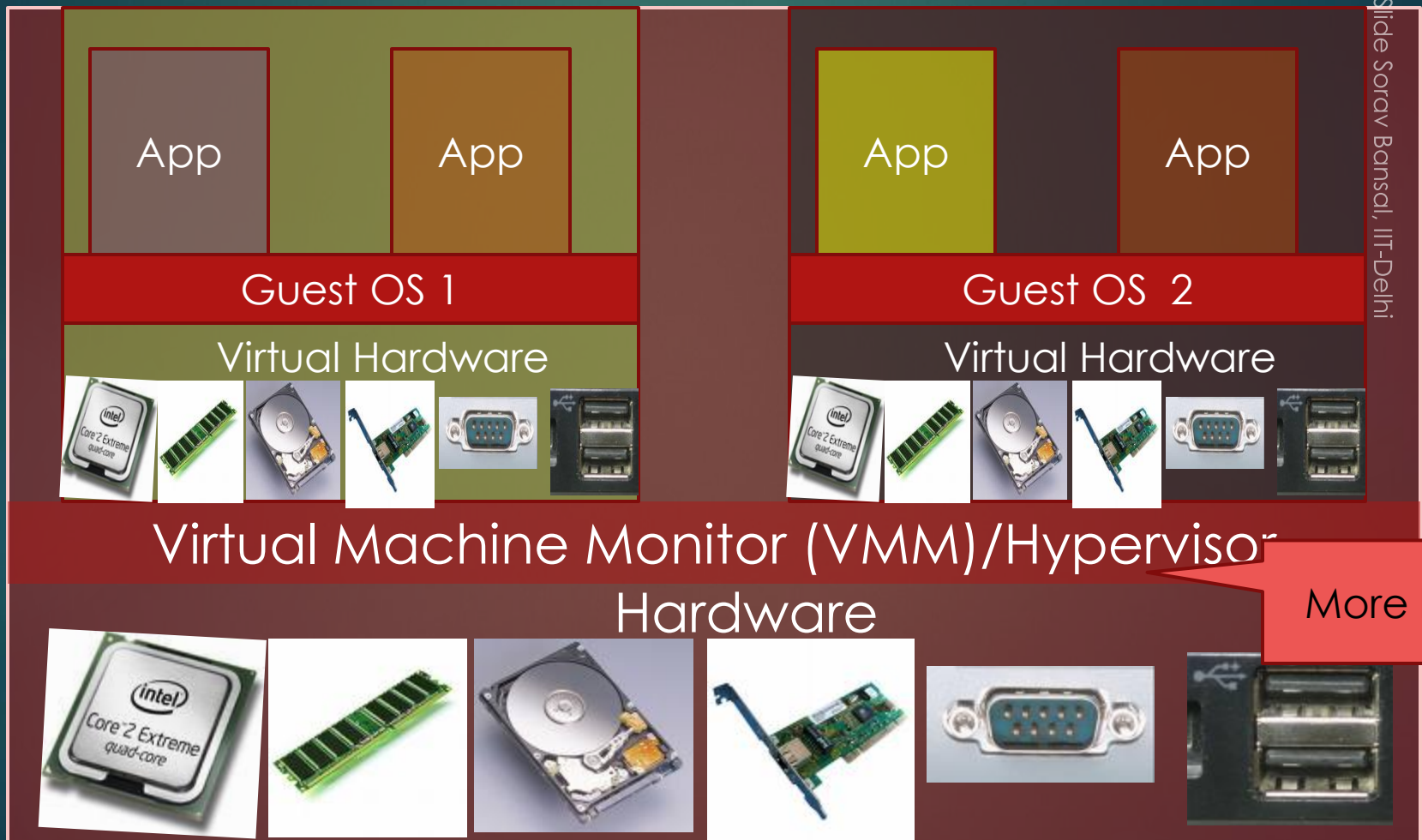
Application
2

OS

Hardware



With virtualization



Virtualization and Cloud Computing

- ▶ A key technology in Cloud Computing
- ▶ Enables different users to share the same physical resources

Some history

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- ▶ An old concept – first virtual machines created on IBM mainframes in early '60s
- ▶ Typically, IBM's virtual machines were identical "copies" of the underlying hardware. Each instance could run its own operating system.
- ▶ Virtualisation formed the basis of "time sharing"



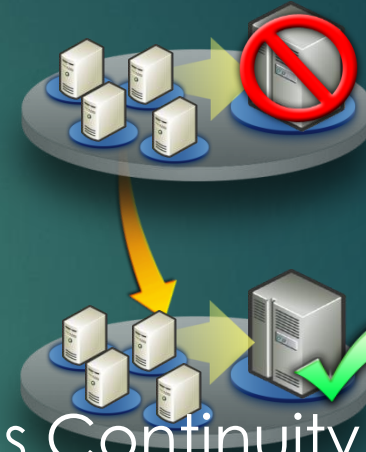
Usage Scenarios for Virtualization

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Consolidation



Business Continuity
Management



Workload Mobility



Development and Test

Why Virtualization

- ▶ Server consolidation
- ▶ Workload mobility
- ▶ Development and test
- ▶ Increase hardware utilization by running multiple applications in isolation on same physical server
- ▶ Move applications from one server to another
- ▶ Easily provision virtual resources for test and dev

Are there other types of VMs?

- ▶ Language VM: Language runtime focused on running a single application
 - ▶ e.g., Java Virtual Machine, Microsoft Common Language runtime, Javascripts
 - ▶ Not in this **course**
- ▶ “Lightweight” VM: Does not run guest OS; But isolates applications from other
 - ▶ e.g., Docker – covered earlier
- ▶ This lecture is only on **System** virtual machine: presents a “copy” of whole machine

How is it implemented?

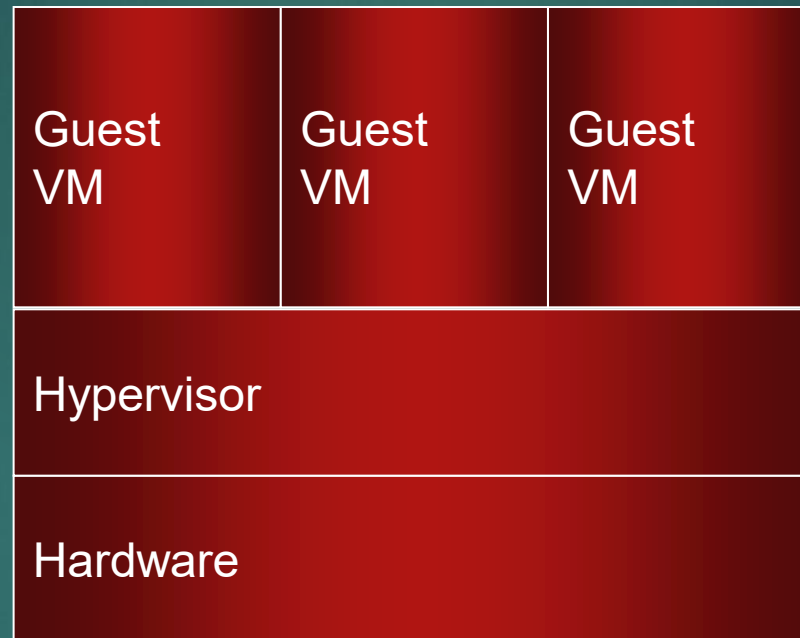
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- ▶ Virtualization usually involves software called a Virtual Machine Monitor (VMM) or hypervisor that creates the illusion of real machines for multiple virtual machines.
- ▶ The hypervisor can run directly on hardware or as an application on a host operating system.

Bare Metal Hypervisors (Type 1)

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IBM CP/CMS

VMware ESX

Windows Virtualisation (2008)

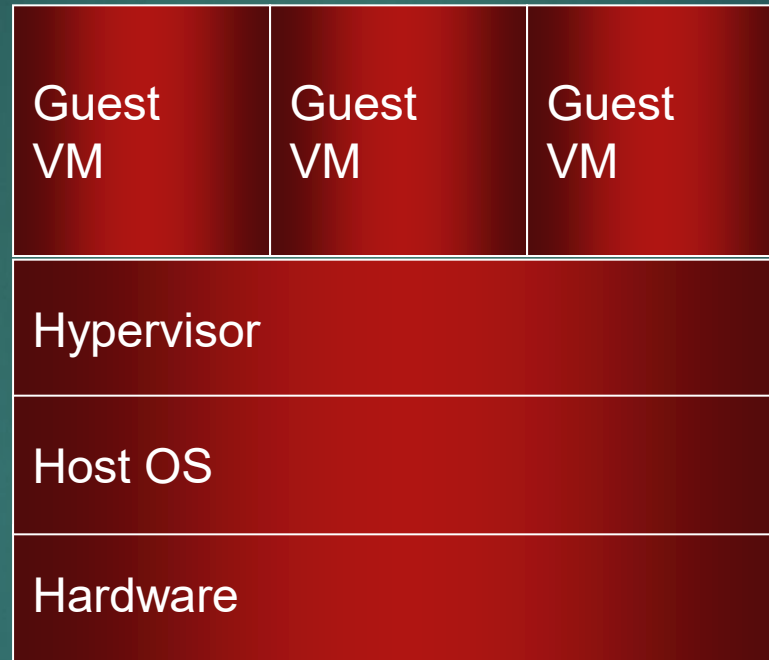
Xen

Virtual Iron

Hosted Hypervisors (Type 2)

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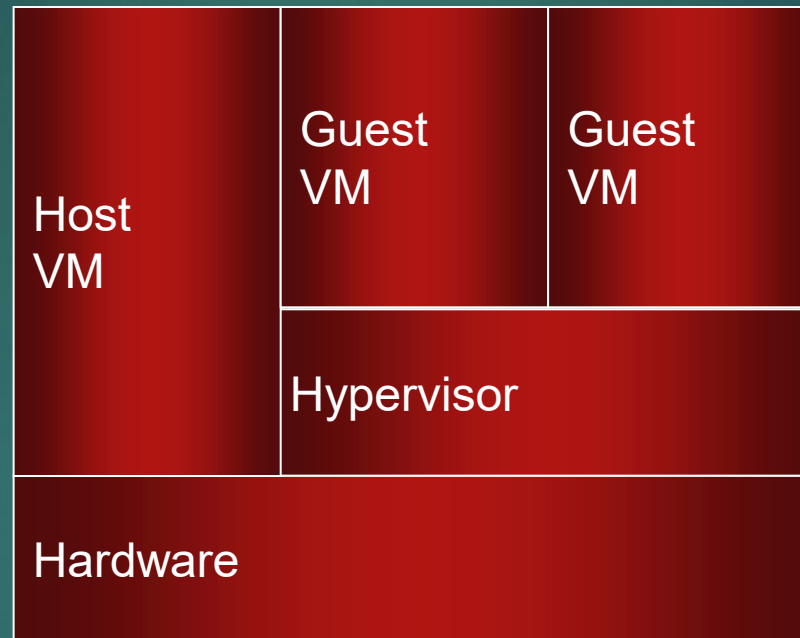


VirtualBox
VMWare Workstation

Hybrid Hypervisors

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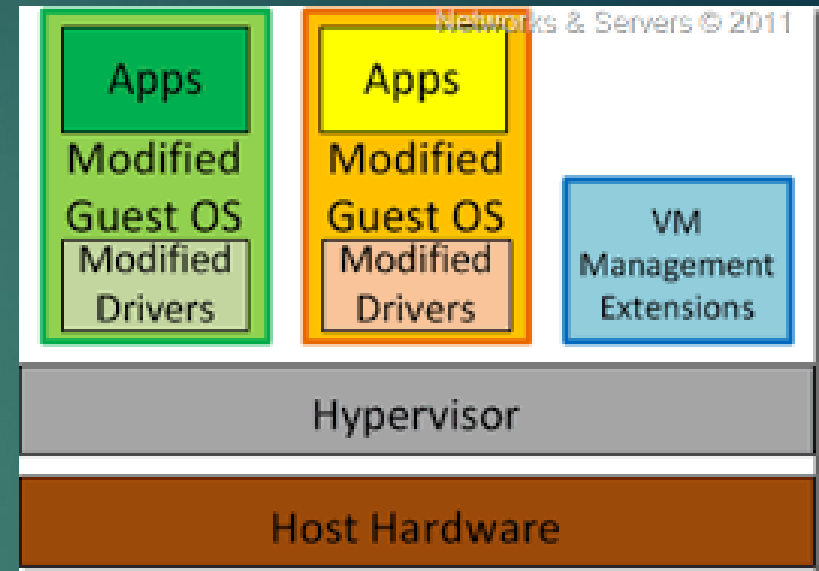
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MS Virtual Server
MS Virtual PC

Types of Virtualization

- ▶ Paravirtualization
 - ▶ Modify OS to run on a hypervisor
 - ▶ Xen, Linux-kvm
 - ▶ Useful if source code of OS is modifiable
 - ▶ IBM: MVS, VM
 - ▶ Linux
 - ▶ Microsoft: Windows
- ▶ Full (transparent) virtualization
 - ▶ OS runs without modification
 - ▶ VMWare
 - ▶ kvm



Types of virtualization


- ▶ Bare-metal vs hosted
- ▶ What is paravirtualization?
- ▶ Pros and cons of paravirtualization
- ▶ Bare-metal – runs on bare machine. Hosted runs on OS
- ▶ Paravirtualization: VMM provides APIs for guest OS
- ▶ Pros: can be more efficient. Cons: Need to modify guest OS

Implementing a VMM

- ▶ In order to virtualize, we must look at following
 - ▶ Instructions -- defined by the ISA (e.g., x86, ARM)
 - ▶ Memory
 - ▶ I/O (network, disk)



- ▶ Let's first look at Instructions



Server Virtualization Techniques: Trap and Emulate Virtualization

VM requirements

- ▶ Guest OS (and process on guest OS) needs to execute on the processor
- ▶ One Guest OS should not interfere with another – in memory and I/O

How to run the VM?

- ▶ The VMM abstracts a CPU

- ▶ Just like CPU reads every instruction, interprets and executes it

- ▶ VMM is responsible for doing the same

- ▶ One way to do it is to

- ▶ Simulate every instruction

- ▶ But this would be very slow

- ▶ Examples: Bochs
(<http://bochs.sourceforge.net/>)

Mov r1, (1000)

VM

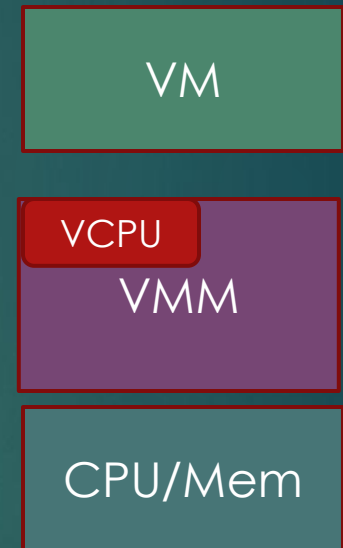
VCPU

VMM

CPU/Mem

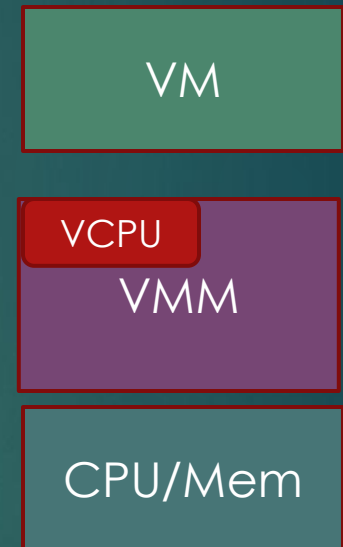
How to run the VM?

- ▶ The VM is also code
- ▶ Why can we not run the VM on the CPU
- ▶ Example
 - ▶ `addl %eax, %ecx`
- ▶ Should be able to run fine on the hardware
- ▶ Can you think of examples of code that cannot run on the hardware?



What cannot run directly?

- ▶ Any code that requires higher privileges to run
- ▶ For example
 - ▶ syscall
 - ▶ `Movl something, %cr3`
 - ▶ I/O

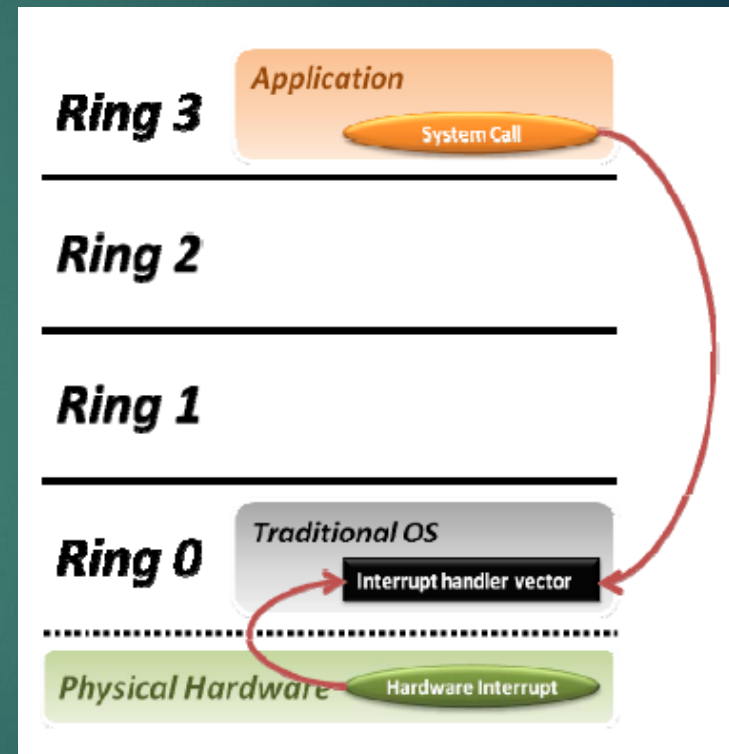


A large red arrow originates from the 'I/O' item in the list and points towards the text box below.

CR3 is the page table.

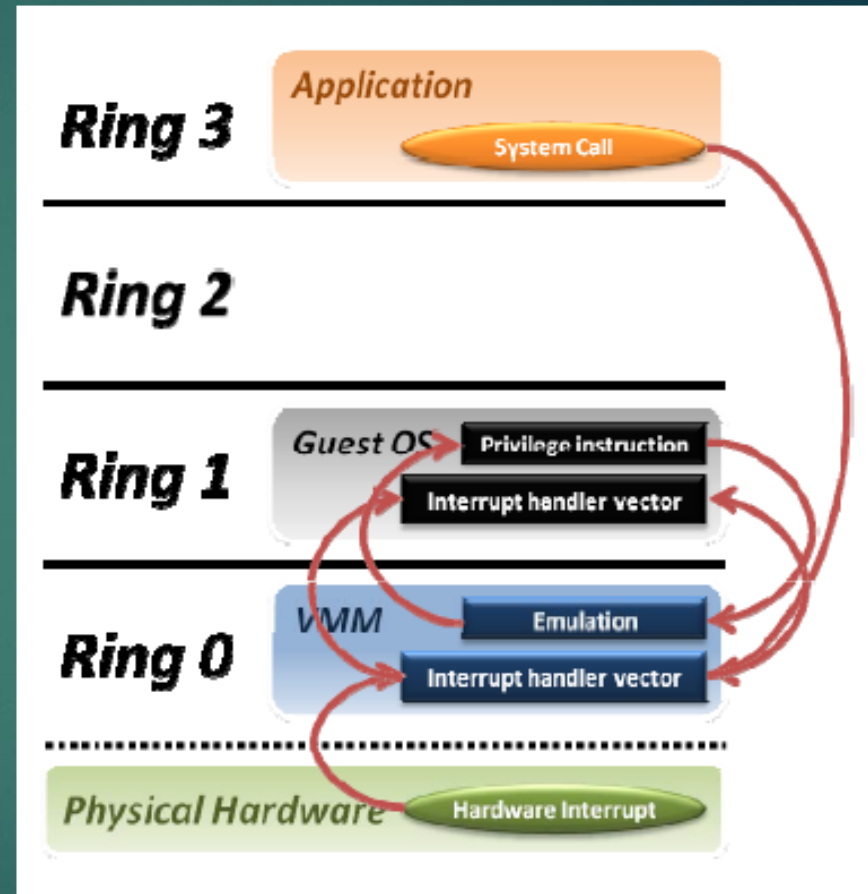
Why cannot it run directly?

- ▶ Privilege Rings in Hardware
- ▶ Typically OS runs in high privilege mode (Ring 0)
- ▶ And application in Ring 3 (less privilege)
- ▶ But with a VMM, the
 - ▶ VMM runs in ring 0
 - ▶ Which ring does Guest OS run in?



Trap and Emulate

- ▶ Run the Guest OS in Ring 1
- ▶ On executing a privileged instruction in Ring 1
 - ▶ Trap into the VMM
 - ▶ VMM emulates the instruction
 - ▶ Does everything that the original instruction was supposed to do but without executing it
- ▶ Then go back to the Guest OS



Trap and Emulate: Earliest Virtualization Technique

Transparent virtualization

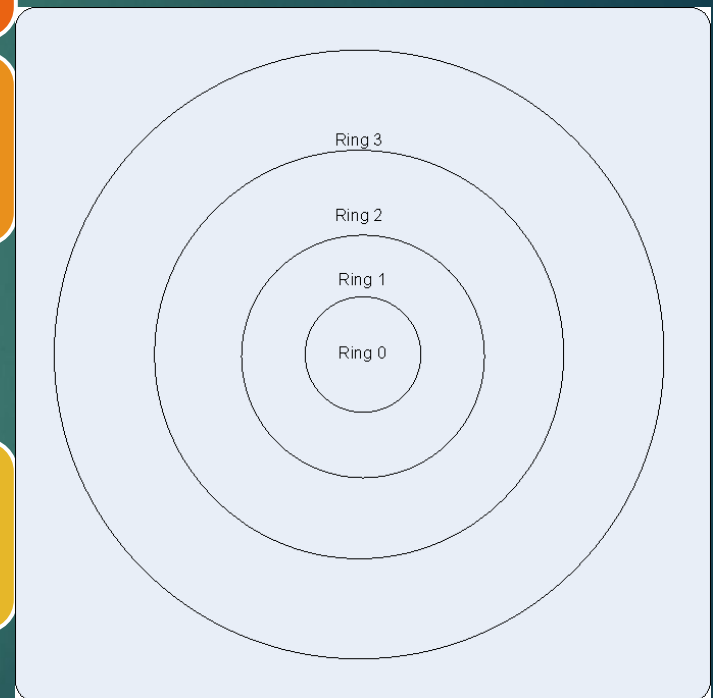
Hypervisor traps privileged instructions and emulates

- Access to physical pages not trapping
- Physical I/O devices privileged
- Control registers instructions could

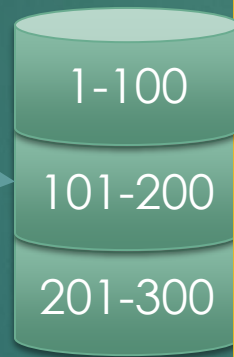
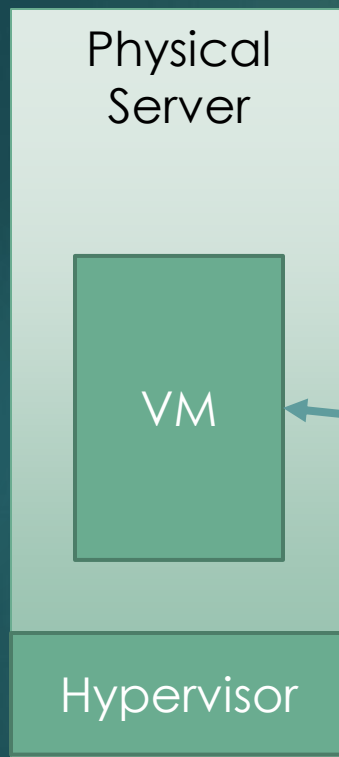
lead to security breach

Handling privileges

- All processors have rings of privilege
- Run hypervisor in highest privilege ring (ring 0)
- Run guest in lower ring

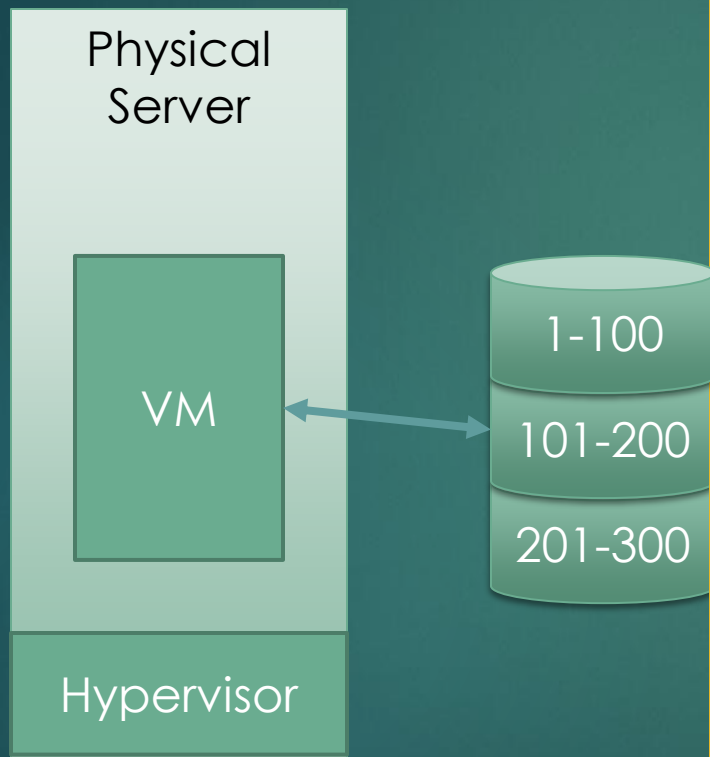


Exercise 1 (10 minutes)



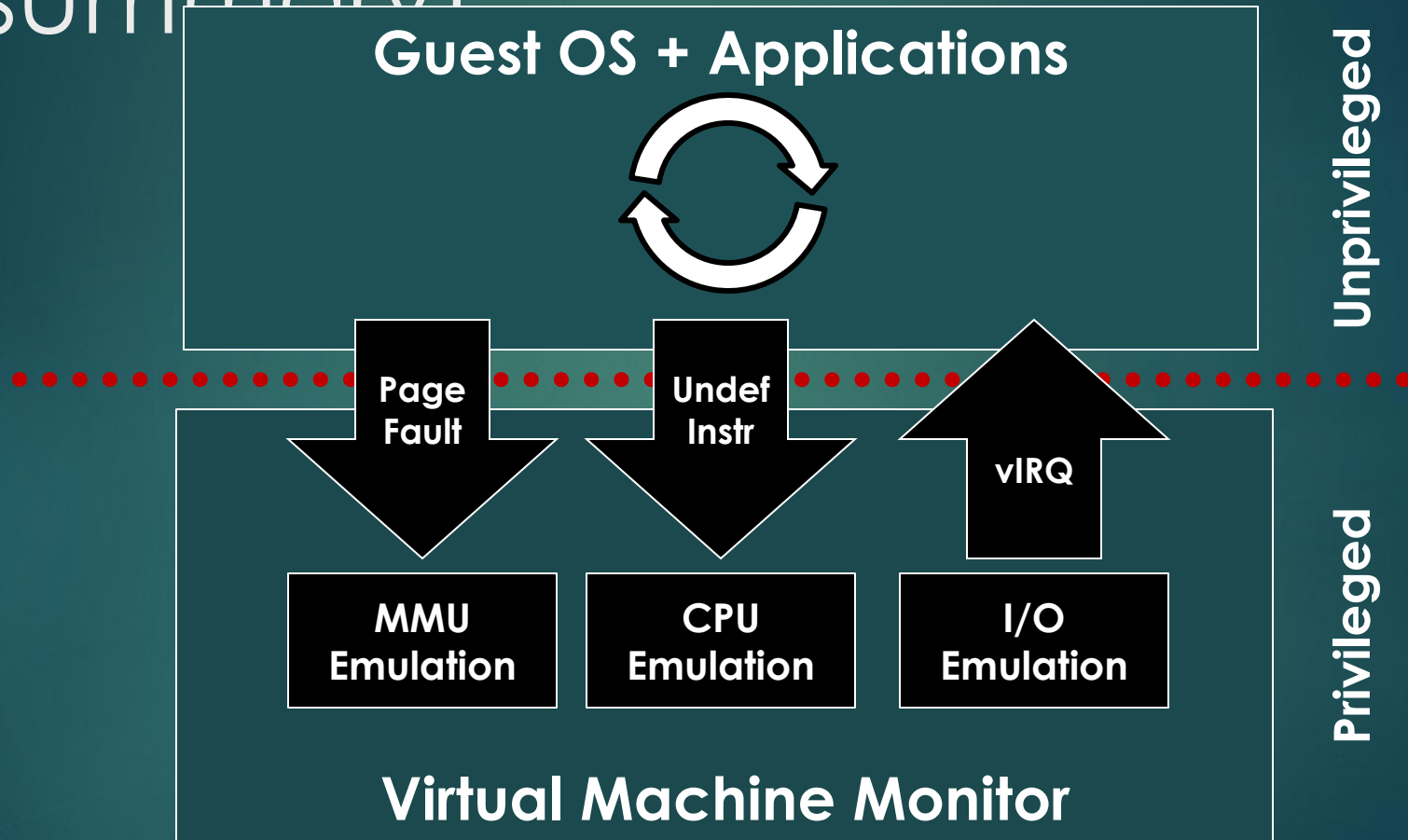
- ▶ Disk has 3 virtual disks – cylinders: (1) 1-100, (2) 101-200, (3) 201-300
- ▶ VM is connected to virtual disk (2)
- ▶ VM executes an instruction to read cylinder 1, block 5 of virtual disk (2)
- ▶ Why is this a privileged instruction?
- ▶ What happens when VM executes instruction?

Exercise 1 solution



- ▶ VM executes an instruction to read cylinder 1, block 5 of virtual disk (2)
- ▶ Why is this a privileged instruction?
 - ▶ Otherwise, VM could read virtual disk (1) or (3). These could belong to other VMs
- ▶ What happens when VM executes instruction?
 - ▶ Hypervisor traps instruction, replaces it with an instruction to read cylinder 101 block 5 of physical disk

Trap and Emulate (summary)



Source: <https://labs.vmware.com/download/45>, Scott Devine, VMWare

Limitations of Trap and Emulate Virtualization

- ▶ Performance Overhead
 - ▶ E.g., trapping privileged instructions
- ▶ Some architectures cannot be virtualized by use of trap and emulate virtualization
 - ▶ Depends on the fact that executing an instruction at a lower privilege than required will cause a **trap** to the VMM.

True or False?

- ▶ Trap and emulate virtualization does not need any special support
- ▶ Trap and emulate virtualization is very efficient
- ▶ False: Trap and emulate virtualization requires CPU to support access rings.
- ▶ False: for many operations, there is a significant overhead