

# Final Exam

Automata Theory Monsoon 2021, IIIT Hyderabad

December 4, 2021

Total Points: 25

Time: 90 mins [75 mins for writing + 15 mins for submission]

**General Instructions:** Upload a scanned PDF of handwritten answers on moodle. No submissions in any form after 10:05 will be evaluated. You should strictly stop writing at 9:45 and start scanning and uploading. DFA is deterministic finite automata. NFA is non-deterministic finite automata. CFL is context free language. If  $S_1$  and  $S_2$  are two strings,  $S_1S_2$  represents their concatenation. All notation is standard from the lectures.

## 0.0.1 Parsing

---

1. [3 points] Comment on the Decidability of the Halting problem for Push-down Automata, compare it to that of the Turing Machine. If they don't match, explain why.
2. [4 points] Let  $R_1, R_2$  be a recursive language, and let  $RE_1, RE_2$  be two recursively enumerable languages. Are the following languages recursive or not? Are they recursively enumerable or not? Justify.

A:  $RE_1 \cup RE_2$

B:  $R_1 \cap RE_2$

C:  $R_1 - R_2 = \{x | (x \in R_1 \text{ and } x \notin R_2) \text{ or } (x \in R_2 \text{ and } x \notin R_1)\}$

3. [4 points] For any language  $L$  over  $\Sigma$ , the prefix closure of  $L$  is defined as:

$$Pre(L) = \{x \in \Sigma^* | \exists y \in \Sigma^* : xy \in L\}$$

Prove that if  $L$  is regular,  $Pre(L)$  is regular.

4. [3 points] Define the language  $EQ_{DFA}$  as follows:

$$EQ_{DFA} = \{\langle M_1, M_2 \rangle | L(M_1) = L(M_2)\}$$

where  $M_1, M_2$  are encodings of DFAs as strings. Prove that  $EQ_{DFA}$  is recursive.

5. [2 points] Which of the following is the correct order and why?

A: Regular  $\subset$  Context Free  $\subset$  Recursive  $\subset$  Recursively Enumerable

B: Regular  $\subseteq$  Context Free  $\subset$  Recursive  $\subseteq$  Recursively Enumerable

C: Regular  $\subset$  Context Free  $\subset$  Recursively Enumerable  $\subset$  Recursive

D: Regular  $\subseteq$  Context Free  $\subseteq$  Recursive  $\subseteq$  Recursively Enumerable

6. [3 points] Use the pumping lemma for CFLs to show that  $L = \{www|w \in \{a,b\}^*\}$  is not context free.
7. [4 points] Let  $L$  be a CFL and  $LEAST(L) = \{y \in L \mid y \text{ is a proper prefix of } x, \forall x \in L\}$ . Show that CFLs are not closed under LEAST operation. (A proper prefix  $X$  of a string  $A$  is a substring such that there exists a non-empty string  $S \in \Sigma^+$  such that  $XS = A$ )
8. [2 points] A. The regular expression  $0^*(10^*)^*$  denotes the same set as
  - (a)  $(1^*0)^*1^*$
  - (b)  $0^+(0^+10)^*$
  - (c) none of theseB. Draw an NFA for the same RegEx.