Data and Its Applications (CS4.301)

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Relational Data Model (contd.)

Relational Algebra

The relational algebra describes a set of operators used to manipulate relations. There are two kinds of operators: set operators (union, intersection, difference and cartesian product) and database-specific operators (select, project and join).

The select operator preserves the columns (attributes) of its input; it filters only the rows (tuples). If we write $\sigma_c(R(A_1,\ldots,A_n))=R'(A_1,\ldots,A_n)$, then the rows of R' are exactly those rows of R which satisfy the condition c. In SQL, this is expressed as SELECT * FROM WHERE c.

The project operator, conversely, preserves the rows (tuples) and returns only some of columns (attributes). For example, $\pi_{(A_{i_1},\dots,A_{i_k})}(R(A_1,\dots,A_n))=$

 $R'(A_{i_1},\ldots,A_{i_k})$. This is expressed as the SQL command SELECT Ai1, ..., Aik FROM . Note that in the general case, only distinct rows are included in the projection. It remains a proper set, like the input.

We write $\pi_k(R)$ to denote the projection of only the key attribute of R.

The join operator is equivalent to the cartesian product, followed by select. For example, $R \bowtie_c S = \sigma_c(R \times S)$. If the condition c involves equality of some attributes, it is called an equijoin.

For an equijoin, if the condition is A = B, then a natural join (denoted R S) does not include the attribute B.

The set operations of union, intersection and difference require that the operands be *compatible*, *i.e.*, if the operation is performed on $R(A_1, \ldots, A_n)$ and $S(B_1, \ldots, B_n)$, then $\mathrm{dom}(A_i) = \mathrm{dom}(B_i)$ for all i.

The cartesian product combines the tuples of two relations. It returns a relation with all the attributes of both its operands, i.e., $R_1(A_1,\ldots,A_n)\times R_2(B_1,\ldots,B_n)=R(A_1,\ldots,A_n,B_1,\ldots,B_n).$