

# Information-Theoretic Methods in Computer Science (CS1.502)

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Assignment 3

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## Problem 1

### Part 1

We define, for an arbitrary permutation  $\tau$ ,

$$N_1(\sigma, \tau) = |\{1, 2, 3, 4\} - \{\sigma(\tau(1)), \sigma(\tau(2)), \dots, \sigma(\tau(k-1))\}| = |[4] - P(k)|,$$

where  $k = \tau^{-1}(1)$ .

We need to show that for all  $\tau$  such that  $\tau(1) = j$ , we have  $N_1(\sigma, \tau) = 5 - j$ .

We note first that if  $\tau(1) = j$ , then  $j = \tau^{-1}(1) = k$ .

Now, since  $\sigma$  and  $\tau$  are both permutations, their composition  $(\sigma \circ \tau)$  must be one also. Therefore  $P(j)$ , being a permutation of  $[j-1]$ , must have exactly  $(j-1)$  elements.

Furthermore, since  $P(j) \subseteq [4]$ , we can say that

$$|[4] - P(j)| = 4 - |P(j)|.$$

Therefore  $|N_1(\sigma, \tau)| = 4 - (j-1) = 5 - j$ , QED.

### Part 2

We have shown in Part 1 that

$$|N_1(\sigma, \tau)| = 5 - \tau^{-1}(1).$$

Therefore, all  $\tau$  such that  $|N_1(\sigma, \tau)| = j$  are such that  $\tau^{-1}(1) = 5 - j$ , so  $1 = \tau(5 - j)$ .

Therefore, the mapping of  $5 - j$  is fixed, and so it is left to map the set  $[4] - \{5 - j\}$  to  $\{2, 3, 4\}$ . The number of such permutations is  $|S_3| = 3! = 6$ , QED.

## Question 2

Let  $Y \sim \mathcal{N}(\mu, \sigma^2)$ . Since entropy is independent of the mean, we can set  $\mu = 0$  WLOG.

Let  $X \sim P_X$  be an arbitrary r.v. with variance  $\sigma^2$ .

Now, note that

$$\begin{aligned} D(X \parallel Y) &= \int_{\mathbb{R}} p_X(t) \log \frac{p_X(t)}{p_Y(t)} dt \\ &= \int_{\mathbb{R}} p_X(t) \log p_X(t) dt - \int_{\mathbb{R}} p_X(t) \log p_Y(t) dt \\ &= -h(X) - I. \end{aligned}$$

The value of  $I$  can be computed as follows:

$$\begin{aligned} I &= \int_{\mathbb{R}} p_X(t) \log p_Y(t) dt \\ &= \int_{\mathbb{R}} p_X(t) \log \left( \frac{1}{\sqrt{2\pi\sigma^2}} \cdot e^{-\frac{t^2}{2\sigma^2}} \right) dt \\ &= \int_{\mathbb{R}} p_X(t) \log \left( \frac{1}{\sqrt{2\pi\sigma^2}} \right) dt + \int_{\mathbb{R}} p_X(t) \log e^{-\frac{t^2}{2\sigma^2}} dt \\ &= -\frac{1}{2} \log(2\pi\sigma^2) \int_{\mathbb{R}} p_X(t) dt - \frac{\log e}{2\sigma^2} \int_{\mathbb{R}} t^2 p_X(t) dt \\ &= -\frac{1}{2} \log(2\pi\sigma^2) \cdot 1 - \frac{1}{2} \log e \\ &= -\frac{1}{2} \log(2\pi e\sigma^2). \end{aligned}$$

Therefore

$$D(X \parallel Y) = -h(X) + \frac{1}{2} \log(2\pi e\sigma^2).$$

Since KL-divergence is always nonnegative, we can conclude that

$$h(X) \leq \frac{1}{2} \log(2\pi e\sigma^2),$$

QED.

## Question 3

Consider the case where  $P_X$  and  $Q_X$  are binary distributions. Let  $p = P_X(0)$  and  $q = Q_X(0)$ . Furthermore, let  $p = \frac{1}{2}$  and  $q = \frac{1}{2} + \varepsilon$ .

Now, we have that

$$D_2(P_X \parallel Q_X) \geq \frac{k}{\ln 2} \text{TV}(P_X, Q_X)^2.$$

From this, we can say that

$$g_p(q) = D_2(P_X \parallel Q_X) - \frac{k}{\ln 2} \text{TV}(P_X, Q_X)^2 \geq 0,$$

*i.e.*, 0 is a minimum of  $g_p(q)$ , and so  $\frac{d}{dq}g_p(q) = 0$ .

Now,

$$\begin{aligned} g'_p(q) = 0 &\implies \frac{p}{\ln 2} \log \frac{p}{q} + \frac{1-p}{\ln 2} \log \frac{1-p}{1-q} - \frac{k}{\ln 2} (p-q)^2 = 0 \\ &\implies -\frac{p}{q} + (1-p) \frac{1}{1-q} + 2k(p-q) = 0 \\ &\implies k = \frac{1}{2(p-q)} \left( \frac{p}{q} - (1-p) \frac{1}{1-q} \right) \\ &\implies k = \frac{1}{2(p-q)} \left( \frac{p(1-q) - q(1-p)}{q(1-q)} \right) \\ &\implies k = \frac{1}{2(p-q)} \left( \frac{p-q}{q(1-q)} \right) \\ &\implies k = \frac{1}{2q(1-q)} \end{aligned}$$

WLOG, we can say that  $q \leq \frac{1}{2}$ . Therefore,

$$\begin{aligned} q(1-q) &\leq \frac{1}{4} \\ \implies \frac{1}{q(1-q)} &\geq 4 \\ \implies k &\geq 2, \end{aligned}$$

QED.

#### Question 4

We have two probability distributions  $P_X, Q_X$  and a function  $f : \mathcal{X} \rightarrow [0, B]$ . We need to show that

$$|\mathbb{E}_{P_X}[f(X)] - \mathbb{E}_{Q_X}[f(X)]| \leq B \cdot \text{TV}(P_X, Q_X).$$

We can proceed as follows.

$$\begin{aligned}
|\mathbb{E}_{P_X}[f(X)] - \mathbb{E}_{Q_X}[f(X)]| &= \left| \sum_{\mathcal{X}} P_X(x)f(x) - \sum_{\mathcal{X}} Q_X(x)f(x) \right| \\
&= \left| \sum_{\mathcal{X}} f(x)(P_X(x) - Q_X(x)) \right| \\
&= \left| \sum_{\mathcal{X}} (P_X(x) - Q_X(x)) \left( f(x) - \frac{B}{2} \right) + \sum_{\mathcal{X}} (P_X(x) - Q_X(x)) \cdot \frac{B}{2} \right| \\
&= \left| \sum_{\mathcal{X}} (P_X(x) - Q_X(x)) \left( f(x) - \frac{B}{2} \right) + 0 \cdot \frac{B}{2} \right| \\
&= \left| \sum_{P_X(x) > Q_X(x)} (P_X(x) - Q_X(x)) \left( f(x) - \frac{B}{2} \right) \right. \\
&\quad \left. - \sum_{P_X(x) \leq Q_X(x)} (Q_X(x) - P_X(x)) \left( f(x) - \frac{B}{2} \right) \right| \\
&\leq \left| \sum_{P_X(x) > Q_X(x)} (P_X(x) - Q_X(x)) \left( f(x) - \frac{B}{2} \right) \right| \\
&\quad + \left| \sum_{P_X(x) \leq Q_X(x)} (Q_X(x) - P_X(x)) \left( f(x) - \frac{B}{2} \right) \right| \\
&= \left| f(x) - \frac{B}{2} \right| \left| \sum_{\mathcal{X}} P_X(x) - Q_X(x) \right| \\
&\leq \frac{B}{2} \left| \sum_{\mathcal{X}} P_X(x) - Q_X(x) \right| \\
&= B \cdot \text{TV}(P_X, Q_X),
\end{aligned}$$

QED.