Sail Manual

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Contents

1 Sail syntax

l	::=			Source location
	|			
annot	::=			
id	::=			Identifier
	|	x		
	|	(**deinfix** x)		remove infix status
	|	**bool**	M	Built in type identifiers
	|	**bit**	M	
	|	**unit**	M	
	|	**nat**	M	
	|	**string**	M	
	|	**range**	M	
	|	**atom**	M	
	|	**vector**	M	
	|	**list**	M	
	|	**set**	M	
	|	**reg**	M	
	|	*to_num*	M	Built in function identifiers
	|	*to_vec*	M	
kid	::=			variables with kind, ticked to differntiate from program variables
	|	'x		
base_kind	::=			base kind
	|	**Type**		kind of types
	|	**Nat**		kind of natural number size expressions

	|	**Order**	kind of vector order specifications
	|	**Effect**	kind of effect sets
kind	::=		kinds
	|	$base\_kind_1 \rightarrow \ ... \ \rightarrow base\_kind_n$	
nexp	::=		expression of kind Nat, for vector sizes and origins
	|	*kid*	variable
	|	*num*	constant
	|	$nexp_1 * nexp_2$	product
	|	$nexp_1 + nexp_2$	sum
	|	$nexp_1 - nexp_2$	subtraction, error for nexp1 to be smaller than nexp2
	|	$2 * *nexp$	exponential
	|	**neg** *nexp*	For internal use
	|	(*nexp*) S	
order	::=		vector order specifications, of kind Order
	|	*kid*	variable
	|	**inc**	increasing (little-endian)
	|	**dec**	decreasing (big-endian)
	|	(*order*) S	
base_effect	::=		effect
	|	**rreg**	read register
	|	**wreg**	write register
	|	**rmem**	read memory
	|	**wmem**	write memory
	|	**wmea**	signal effective address for writing memory
	|	**wmv**	write memory, sending only value
	|	**barr**	memory barrier

	depend		dynamic footprint
	undef		undefined-instruction exception
	unspec		unspecified values
	nondet		nondeterminism from intra-instruction parallelism
	escape		Tracking of expressions and functions that might call exit
	lset		Local mutation happend; not user-writable

effect	::=		effect set, of kind Effects
	kid		
	$\{base\_effect_1, .., base\_effect_n\}$		effect set
	pure	M	sugar for empty effect set
	$effect_1 \uplus .. \uplus effect_n$	M	meta operation for combining sets of effects

typ	::=		Type expressions, of kind **Type**
	_		Unspecified type
	id		Defined type
	kid		Type variable
	$typ_1 \rightarrow typ_2$ **effect** *effect*		Function type (first-order only in user code)
	$(typ_1, ..., typ_n)$		Tuple type
	$id\langle typ\_arg_1, .., typ\_arg_n\rangle$		type constructor application
	(typ)	S	
	$[\|nexp\|]$	S	sugar for `range<0, nexp>`
	$[\|nexp : nexp'\|]$	S	sugar for `range< nexp, nexp'>`
	[: *nexp* :]	S	sugar for `atom<nexp>` which is special case of `range<nexp,nexp>`
	$typ[nexp]$	S	sugar for vector indexed by [*nexp*]
	$typ[nexp : nexp']$	S	sugar for vector indexed by [*nexp..nexp'*]
	$typ[nexp <: nexp']$	S	sugar for increasing vector indexed as above
	$typ[nexp :> nexp']$	S	sugar for decreasing vector indexed as above

typ_arg	::=	Type constructor arguments of all kinds

	nexp	
	typ	
	order	
	effect	

n_constraint	::=		constraint over kind **Nat**
	|	$nexp = nexp'$	
	|	$nexp \geq nexp'$	
	|	$nexp \leq nexp'$	
	|	kid **IN** $\{num_1, \ldots, num_n\}$	

kinded_id	::=		optionally kind-annotated identifier
	|	*kid*	identifier
	|	*kind kid*	kind-annotated variable

quant_item	::=		Either a kinded identifier or a nexp constraint for a typquant
	|	*kinded_id*	An optionally kinded identifier
	|	*n_constraint*	A constraint for this type

typquant	::=		type quantifiers and constraints
	|	**forall** $quant\_item_1, \ldots, quant\_item_n$.	
	|		sugar, omitting quantifier and constraints

typschm	::=		type scheme
	|	*typquant typ*	

name_scm_opt	::=		Optional variable-naming-scheme specification for variables of defined type
	|		
	|	[**name** = *regexp*]	

type_def	::=		Type definition body
	|	**typedef** *id name_scm_opt* = *typschm*	
			type abbreviation
	|	**typedef** *id name_scm_opt* = **const struct** *typquant*{*typ*$_1$ *id*$_1$; ... ; *typ*$_n$ *id*$_n$;$^?$}	
			struct type definition
	|	**typedef** *id name_scm_opt* = **const union** *typquant*{*type_union*$_1$; ... ; *type_union*$_n$;$^?$}	
			union type definition
	|	**typedef** *id name_scm_opt* = **enumerate** {*id*$_1$; ... ; *id*$_n$;$^?$}	
			enumeration type definition
	|	**typedef** *id* = **register bits** [*nexp* : *nexp*′]{*index_range*$_1$: *id*$_1$; ... ; *index_range*$_n$: *id*$_n$}	
			register mutable bitfield type definition
type_union	::=		Type union constructors
	|	*id*	
	|	*typ id*	
index_range	::=		index specification, for bitfields in register types
	|	*num*	single index
	|	*num*$_1$..*num*$_2$	index range
	|	*index_range*$_1$, *index_range*$_2$	concatenation of index ranges
lit	::=		Literal constant
	|	()	() : **unit**
	|	**bitzero**	**bitzero** : **bit**
	|	**bitone**	**bitone** : **bit**
	|	**true**	**true** : **bool**
	|	**false**	**false** : **bool**
	|	*num*	natural number constant
	|	*hex*	bit vector constant, C-style
	|	*bin*	bit vector constant, C-style

	undefined	constant representing undefined values
	string	string constant

$;^?$	::=		Optional semi-colon
	|		
	|	;	

pat	::=		Pattern
	|	*lit*	literal constant pattern
	|	_	wildcard
	|	$(pat\ \mathbf{as}\ id)$	named pattern
	|	$(typ)pat$	typed pattern
	|	*id*	identifier
	|	$id(pat_1, .., pat_n)$	union constructor pattern
	|	$\{fpat_1; ...; fpat_n\ ;^?\}$	struct pattern
	|	$[pat_1, .., pat_n]$	vector pattern
	|	$[num_1 = pat_1, .., num_n = pat_n]$	vector pattern (with explicit indices)
	|	$pat_1 : .... : pat_n$	concatenated vector pattern
	|	$(pat_1, ..., pat_n)$	tuple pattern
	|	$[\|\|pat_1, .., pat_n\|\|]$	list pattern
	|	(pat)	S

fpat	::=		Field pattern
	|	$id = pat$	

exp	::=		Expression
	|	$\{exp_1; ...; exp_n\}$	block
	|	**nondet** $\{exp_1; ...; exp_n\}$	nondeterminisitic block, expressions evaluate in an unspecified order, or concurrently
	|	*id*	identifier
	|	*lit*	literal constant

$(typ)exp$		cast
$id(exp_1, .., exp_n)$		function application
$id\ exp$	S	No extra parens needed when exp is a tuple
$exp_1\ id\ exp_2$		infix function application
$(exp_1, ..., exp_n)$		tuple
if exp_1 **then** exp_2 **else** exp_3		conditional
if exp_1 **then** exp_2	S	
foreach $(id$ **from** exp_1 **to** exp_2 **by** exp_3 **in** $order)exp_4$		loop
foreach $(id$ **from** exp_1 **to** exp_2 **by** $exp_3)exp_4$	S	
foreach $(id$ **from** exp_1 **to** $exp_2)exp_3$	S	
foreach $(id$ **from** exp_1 **downto** exp_2 **by** $exp_3)exp_4$	S	
foreach $(id$ **from** exp_1 **downto** $exp_2)exp_3$	S	
$[exp_1, ..., exp_n]$		vector (indexed from 0)
$[num_1 = exp_1, ..., num_n = exp_n\ opt\_default]$		vector (indexed consecutively)
$exp[exp']$		vector access
$exp[exp_1..exp_2]$		subvector extraction
$[exp$ **with** $exp_1 = exp_2]$		vector functional update
$[exp$ **with** $exp_1 : exp_2 = exp_3]$		vector subrange update (with vector)
$exp : exp_2$		vector concatenation
$[\|\|exp_1, .., exp_n\|\|]$		list
$exp_1 :: exp_2$		cons
$\{fexps\}$		struct
$\{exp$ **with** $fexps\}$		functional update of struct
$exp.id$		field projection from struct
switch $exp\{$ **case** $pexp_1 ...$ **case** $pexp_n\}$		pattern matching
$letbind$ **in** exp		let expression
$lexp := exp$		imperative assignment
exit exp		expression to halt all current execution, potentially calling a system, trap, or interrupt handler with
assert (exp, exp')		expression to halt with error, when the first expression is true, reporting the optional string as an err

	|	(exp)	S	
	|	$(annot)exp$		This is an internal cast, generated during type checking that will resolve into a syntactic cast after
	|	$annot$		This is an internal use for passing nexp information to library functions, postponed for constraint solving
	|	$annot, annot'$		This is like the above but the user has specified an implicit parameter for the current function
	|	**comment** $string$		For generated unstructured comments
	|	**comment** exp		For generated structured comments
	|	**let** $lexp = exp$ **in** exp'		This is an internal node for compilation that demonstrates the scope of a local mutable variable
	|	**let** $pat = exp$ **in** exp'		This is an internal node, used to distinguised some introduced lets during processing from original ones
	|	**return** (exp)		For internal use to embed into monad definition
$lexp$	::=			lvalue expression
	|	id		identifier
	|	$id(exp_1, .., exp_n)$		memory write via function call
	|	$id\ exp$	S	
	|	$(typ)id$		
	|	$lexp[exp]$		vector element
	|	$lexp[exp_1..exp_2]$		subvector
	|	$lexp.id$		struct field
$fexp$	::=			Field-expression
	|	$id = exp$		
$fexps$	::=			Field-expression list
	|	$fexp_1; ...; fexp_n\ ;^?$		
$opt\_default$	::=			Optional default value for indexed vectors, to define a defualt value for any unspecified positions in a sparse map
	|			
	|	; **default** $= exp$		
$pexp$	::=			Pattern match

	|	*pat* → *exp*	
tannot_opt	::=		Optional type annotation for functions
	|	*typquant typ*	
rec_opt	::=		Optional recursive annotation for functions
	|		non-recursive
	|	**rec**	recursive
effect_opt	::=		Optional effect annotation for functions
	|		sugar for empty effect set
	|	**effect** *effect*	
funcl	::=		Function clause
	|	*id pat* = *exp*	
fundef	::=		Function definition
	|	**function** *rec_opt tannot_opt effect_opt* $funcl_1$ **and** ... **and** $funcl_n$	
letbind	::=		Let binding
	|	**let** *typschm pat* = *exp*	value binding, explicit type (*pat* must be total)
	|	**let** *pat* = *exp*	value binding, implicit type (*pat* must be total)
val_spec	::=		Value type specification
	|	**val** *typschm id*	
	|	**val extern** *typschm id*	
	|	**val extern** *typschm id* = *string*	Specify the type and id of a function from Lem, where the string must provide an explicit path to the requi
default_spec	::=		Default kinding or typing assumption

| | **default** *base_kind kid*
| | **default Order** *order*
| | **default** *typschm id*

scattered_def ::= Function and type union definitions that can be spread across a file. Each one must end in **id**
| | **scattered function** *rec_opt tannot_opt effect_opt id*
scattered function definition header
| | **function clause** *funcl* scattered function definition clause
| | **scattered typedef** *id name_scm_opt* = **const union** *typquant*
scattered union definition header
| | **union** *id* **member** *type_union* scattered union definition member
| | **end** *id* scattered definition end

reg_id ::=
| | *id*

alias_spec ::= Register alias expression forms. Other than where noted, each id must refer to an unaliased register of type
| | *reg_id*.*id*
| | *reg_id*[*exp*]
| | *reg_id*[*exp*..*exp′*]
| | *reg_id* : *reg_id′*

dec_spec ::= Register declarations
| | **register** *typ id*
| | **register alias** *id* = *alias_spec*
| | **register alias** *typ id* = *alias_spec*

def ::= Top-level definition
| | *type_def* type definition
| | *fundef* function definition

		letbind	value definition
		val_spec	top-level type constraint
		default_spec	default kind and type assumptions
		scattered_def	scattered function and type definition
		dec_spec	register declaration
		dec_comm	generated comments

defs	::=		Definition sequence
		def_1 .. def_n	

2 Sail primitive types and functions

built_in_types ::= Type Kind

	bit : **Typ**
	unit : **Typ**
	forall Nat $'n$. **Nat** $'m$. **range** $<' n,' m >$: **Nat** $\rightarrow$ **Nat** $\rightarrow$ **Typ**
	forall Nat $'n$. **atom** $<' n >$: **Nat** $\rightarrow$ **Typ** — singleton number, instead of range¡'n,'n¿
	forall Nat $'n$, **Nat** $'m$, **Order** $'o$, **Typ** $'t$. **vector** $\langle' n,' m,' o,' t\rangle$: **Nat** $\rightarrow$ **Nat** $\rightarrow$ **Order** $\rightarrow$ **Typ**
	forall Typ $'t$. **register** $\langle' t\rangle$: **Typ** $\rightarrow$ **Typ**
	forall Typ $'t$. **reg** $\langle' t\rangle$: **Typ** $\rightarrow$ **Typ** — internal reference cell
	forall Nat $'n$. **implicit** $<' n >$: **Nat** $\rightarrow$ **Typ** — see Kathy for explanation

built_in_type_abbreviations ::=

	bool $\Rightarrow$ **bit**
	nat $\Rightarrow [\vert 0..pos\_infinity \vert]$
	int $\Rightarrow [\vert neg\_infinity..pos\_infinity \vert]$
	$uint8 \Rightarrow [\vert 0..2 * *8 \vert]$
	$uint16 \Rightarrow [\vert 0..2 * *16 \vert]$
	$uint32 \Rightarrow [\vert 0..2 * *32 \vert]$
	$uint64 \Rightarrow [\vert 0..2 * *32 \vert]$

functions ::= Built-in functions: all have effect pure, all order polymorphic

	val forall Typ $'a.'a \rightarrow$ **unit** : **ignore**
	val $([\vert' n..' m \vert], [\vert' o..' p \vert]) \rightarrow [\vert' n +' o..' m +' p \vert] : +$ — arithmetic addition
	val forall Nat $'n$.(**bit** $['n]$, **bit** $['n]) \rightarrow$ **bit** $['n] : +$ — unsigned vector addition
	val forall Nat $'n$.(**bit** $['n]$, **bit** $['n]) \rightarrow$ (**bit** $['n]$, **bit** , **bit**) : $+$ — unsigned vector addition with overflow, carry out
	val forall Nat $'n$.(**bit** $['n]$, **bit** $['n]) \rightarrow$ **bit** $['n] : +\_s$ — signed vector addition
	val forall Nat $'n$.($bit['n]$, **bit** $['n]) \rightarrow$ (**bit** $['n]$, **bit** , **bit**) : $+\_s$ — signed vector addition with overflow, carry out
	val $([\vert' n..' m \vert], [\vert' o..' p \vert]) \rightarrow [\vert' n -' o..' m -' p \vert] : -$ — arithmetic subtraction
	val forall Nat $'n$.(**bit** $['n]$, **bit** $['n]) \rightarrow$ **bit** $['n] : -$ — unsigned vector subtraction
	val forall Nat $'n$.(**bit** $['n]$, **bit** $['n]) \rightarrow$ (**bit** $['n]$, **bit** , **bit**) : $-$ — unsigned vector subtraction with overflow, carry out

$\mathbf{val}\,\mathbf{forall}\,\mathbf{Nat}\,'n.(\,\mathbf{bit}\,['n],\,\mathbf{bit}\,['n]) \to \mathbf{bit}\,['n] : -\_s$	signed vector subtraction
$\mathbf{val}\,\mathbf{forall}\,\mathbf{Nat}\,'n.(bit['n],\,\mathbf{bit}\,['n]) \to (\,\mathbf{bit}\,['n],\,\mathbf{bit}\,,\,\mathbf{bit}\,) : -\_s$	signed vector subtraction with overflow, carry out
$\mathbf{val}\,([|'n..'m|],[|'o..'p|]) \to [|'n *' o..'m *' p|] : *$	arithmetic multiplication
$\mathbf{val}\,\mathbf{forall}\,\mathbf{Nat}\,'n.(\,\mathbf{bit}\,['n],\,\mathbf{bit}\,['n]) \to \mathbf{bit}\,[2 *' n] : *$	unsigned vector multiplication
$\mathbf{val}\,\mathbf{forall}\,\mathbf{Nat}\,'n.(\,\mathbf{bit}\,['n],\,\mathbf{bit}\,['n]) \to \mathbf{bit}\,[2 *' n] : *\_s$	signed vector multiplication
$\mathbf{val}\,([|'n..'m|],[|1..'p|]) \to [|0..'p-1|] : \mathbf{mod}$	arithmetic modulo
$\mathbf{val}\,\mathbf{forall}\,\mathbf{Nat}\,'n.(\,\mathbf{bit}\,['n],\,\mathbf{bit}\,['n]) \to \mathbf{bit}\,['n] : \mathbf{mod}$	unsigned vector modulo
$\mathbf{val}\,([|'n..'m|],[|1..'p|]) \to [|'q..'r|] : \mathbf{quot}$	arithmetic integer division
$\mathbf{val}\,\mathbf{forall}\,\mathbf{Nat}\,'n,\,\mathbf{Nat}\,'m.(\,\mathbf{bit}\,['n],\,\mathbf{bit}\,['m]) \to \mathbf{bit}\,['n] : \mathbf{quot}$	unsigned vector division
$\mathbf{val}\,\mathbf{forall}\,\mathbf{Nat}\,'n,\,\mathbf{Nat}\,'m.(\,\mathbf{bit}\,['n],\,\mathbf{bit}\,['m]) \to \mathbf{bit}\,['n] : quot\_s$	signed vector division
$\mathbf{val}\,\mathbf{forall}\,\mathbf{Typ}\,'a,\,\mathbf{Nat}\,'n.('a['n] \to [:'\,n\,:]) : \mathbf{length}$	
$\mathbf{val}\,\mathbf{forall}\,\mathbf{Typ}\,'a,\,\mathbf{Nat}\,'n,\,\mathbf{Nat}\,'m,'\,n \leq' m.(\,\mathbf{implicit}\,\langle'm\rangle,'\,a['n]) \to' a['m] : \mathbf{mask}$	reduce size of vector, dropping MSBits. Type system supplies implicit p
$\mathbf{val}\,\mathbf{forall}\,\mathbf{Nat}\,'n.(\,\mathbf{bit}\,['n],\,\mathbf{bit}\,['n]) \to \mathbf{bit}\,:\equiv$	vector equality
$\mathbf{val}\,\mathbf{forall}\,\mathbf{Typ}\,'a,\,\mathbf{Typ}\,'b.('a,'\,b) \to \mathbf{bit}\,:\equiv$	
$\mathbf{val}\,\mathbf{forall}\,\mathbf{Typ}\,'a,\,\mathbf{Typ}\,'b.('a,'\,b) \to \mathbf{bit}\,:!=$	
$\mathbf{val}\,([|'n..'m|],[|'o..'p|]) \to \mathbf{bit}\,:\langle$	
$\mathbf{val}\,\mathbf{forall}\,\mathbf{Nat}\,'n.(\,\mathbf{bit}\,['n],\,\mathbf{bit}\,['n]) \to \mathbf{bit}\,:\langle$	unsigned less than
$\mathbf{val}\,\mathbf{forall}\,\mathbf{Nat}\,'n.(\,\mathbf{bit}\,['n],\,\mathbf{bit}\,['n]) \to \mathbf{bit}\,:< \_s$	
$\mathbf{val}\,([|'n..'m|],[|'o..'p|]) \to \mathbf{bit}\,:\rangle$	
$\mathbf{val}\,\mathbf{forall}\,\mathbf{Nat}\,'n.(\,\mathbf{bit}\,['n],\,\mathbf{bit}\,['n]) \to \mathbf{bit}\,:\rangle$	unsigned greater than
$\mathbf{val}\,\mathbf{forall}\,\mathbf{Nat}\,'n.(\,\mathbf{bit}\,['n],\,\mathbf{bit}\,['n]) \to \mathbf{bit}\,:> \_s$	
$\mathbf{val}\,([|'n..'m|],[|'o..'p|]) \to \mathbf{bit}\,:\leq$	
$\mathbf{val}\,\mathbf{forall}\,\mathbf{Nat}\,'n.(\,\mathbf{bit}\,['n],\,\mathbf{bit}\,['n]) \to \mathbf{bit}\,:\leq$	unsigned less than or eq
$\mathbf{val}\,\mathbf{forall}\,\mathbf{Nat}\,'n.(\,\mathbf{bit}\,['n],\,\mathbf{bit}\,['n]) \to \mathbf{bit}\,:<= \_s$	
$\mathbf{val}\,([|'n..'m|],[|'o..'p|]) \to \mathbf{bit}\,:\geq$	
$\mathbf{val}\,\mathbf{forall}\,\mathbf{Nat}\,'n.(\,\mathbf{bit}\,['n],\,\mathbf{bit}\,['n]) \to \mathbf{bit}\,:\geq$	unsigned greater than or eq
$\mathbf{val}\,\mathbf{forall}\,\mathbf{Nat}\,'n.(\,\mathbf{bit}\,['n],\,\mathbf{bit}\,['n]) \to \mathbf{bit}\,:>= \_s$	
$\mathbf{val}\,\mathbf{bit} \to \mathbf{bit}\,:$	bit negation
$\mathbf{val}\,\mathbf{forall}\,\mathbf{Nat}\,'n.\,\mathbf{bit}\,['n] \to \mathbf{bit}\,['n] :$	bitwise negation

	$\mathbf{val}\,(\,\mathbf{bit}\,,\,\mathbf{bit}\,)\rightarrow\,\mathbf{bit}\,:\mid$	bitwise or
|	$\mathbf{val}\,\mathbf{forall}\,\mathbf{Nat}\,'n.(\,\mathbf{bit}\,['n],\,\mathbf{bit}\,['n])\rightarrow\,\mathbf{bit}\,['n]:\mid$	
|	$\mathbf{val}\,(\,\mathbf{bit}\,,\,\mathbf{bit}\,)\rightarrow\,\mathbf{bit}\,:\&$	bitwise and
|	$\mathbf{val}\,\mathbf{forall}\,\mathbf{Nat}\,'n.(\,\mathbf{bit}\,['n],\,\mathbf{bit}\,['n])\rightarrow\,\mathbf{bit}\,['n]:\&$	
|	$\mathbf{val}\,(\,\mathbf{bit}\,,\,\mathbf{bit}\,)\rightarrow\,\mathbf{bit}\,:\uparrow$	bitwise xor
|	$\mathbf{val}\,\mathbf{forall}\,\mathbf{Nat}\,'n.(\,\mathbf{bit}\,['n],\,\mathbf{bit}\,['n])\rightarrow\,\mathbf{bit}\,['n]:\uparrow$	
|	$\mathbf{val}\,\mathbf{forall}\,\mathbf{Nat}\,'n.(\,\mathbf{bit}\,,[\|'n\|])\rightarrow\,\mathbf{bit}\,['n]:\uparrow\uparrow$	duplicate bit into a vector
|	$\mathbf{val}\,\mathbf{forall}\,\mathbf{Nat}\,'n,\,\mathbf{Nat}\,'m,'\,m\leq'\,n.(\,\mathbf{bit}\,['n],[\|'m\|])\rightarrow\,\mathbf{bit}\,['n]:<<$	left shift
|	$\mathbf{val}\,\mathbf{forall}\,\mathbf{Nat}\,'n,\,\mathbf{Nat}\,'m,'\,m\leq'\,n.(\,\mathbf{bit}\,['n],[\|'m\|])\rightarrow\,\mathbf{bit}\,['n]:>>$	right shift
|	$\mathbf{val}\,\mathbf{forall}\,\mathbf{Nat}\,'n,\,\mathbf{Nat}\,'m,'\,m\leq'\,n.(\,\mathbf{bit}\,['n],[\|'m\|])\rightarrow\,\mathbf{bit}\,['n]:<<<$	rotate

functions_with_coercions ::=

|	$\mathbf{val}\,\mathbf{forall}\,\mathbf{Nat}\,'n.(bit['n],\,\mathbf{bit}\,['n])\rightarrow[\|2\,{*}{*}'\,n\|]:+$
|	$\mathbf{val}\,\mathbf{forall}\,\mathbf{Nat}\,'n,\,\mathbf{Nat}\,'o,\,\mathbf{Nat}\,'p.(\,\mathbf{bit}\,['n],[\|'o..'p\|])\rightarrow\,\mathbf{bit}\,['n]:+$
|	$\mathbf{val}\,\mathbf{forall}\,\mathbf{Nat}\,'n,\,\mathbf{Nat}\,'o,\,\mathbf{Nat}\,'p.([\|'o..'p\|],\,\mathbf{bit}\,['n])\rightarrow\,\mathbf{bit}\,['n]:+$
|	$\mathbf{val}\,\mathbf{forall}\,\mathbf{Nat}\,'n,\,\mathbf{Nat}\,'o,\,\mathbf{Nat}\,'p.(\,\mathbf{bit}\,['n],[\|'o..'p\|])\rightarrow[\|'o..'p+2*{*}'n\|]:+$
|	$\mathbf{val}\,\mathbf{forall}\,\mathbf{Nat}\,'n(\,\mathbf{bit}\,['n],\,\mathbf{bit}\,)\rightarrow\,\mathbf{bit}\,['n]:+$
|	$\mathbf{val}\,\mathbf{forall}\,\mathbf{Nat}\,'n(\,\mathbf{bit}\,,\,\mathbf{bit}\,['n])\rightarrow\,\mathbf{bit}\,['n]:+$
|	$\mathbf{val}\,\mathbf{forall}\,\mathbf{Nat}\,'n.(bit['n],\,\mathbf{bit}\,['n])\rightarrow[\|2\,{*}{*}'\,n\|]:+\_s$
|	$\mathbf{val}\,\mathbf{forall}\,\mathbf{Nat}\,'n,\,\mathbf{Nat}\,'o,\,\mathbf{Nat}\,'p.(\,\mathbf{bit}\,['n],[\|'o..'p\|])\rightarrow\,\mathbf{bit}\,['n]:+\_s$
|	$\mathbf{val}\,\mathbf{forall}\,\mathbf{Nat}\,'n,\,\mathbf{Nat}\,'o,\,\mathbf{Nat}\,'p.([\|'o..'p\|],\,\mathbf{bit}\,['n])\rightarrow\,\mathbf{bit}\,['n]:+\_s$
|	$\mathbf{val}\,\mathbf{forall}\,\mathbf{Nat}\,'n,\,\mathbf{Nat}\,'o,\,\mathbf{Nat}\,'p.(\,\mathbf{bit}\,['n],[\|'o..'p\|])\rightarrow[\|'o..'p+2*{*}'n\|]:+\_s$
|	$\mathbf{val}\,\mathbf{forall}\,\mathbf{Nat}\,'n(\,\mathbf{bit}\,['n],\,\mathbf{bit}\,)\rightarrow\,\mathbf{bit}\,['n]:+\_s$
|	$\mathbf{val}\,\mathbf{forall}\,\mathbf{Nat}\,'n(\,\mathbf{bit}\,,\,\mathbf{bit}\,['n])\rightarrow\,\mathbf{bit}\,['n]:+\_s$
|	$\mathbf{val}\,\mathbf{forall}\,\mathbf{Nat}\,'n,\,\mathbf{Nat}\,'o,\,\mathbf{Nat}\,'p.(\,\mathbf{bit}\,['n],[\|'o..'p\|])\rightarrow\,\mathbf{bit}\,['n]:-$
|	$\mathbf{val}\,\mathbf{forall}\,\mathbf{Nat}\,'n,\,\mathbf{Nat}\,'o,\,\mathbf{Nat}\,'p.([\|'o..'p\|],\,\mathbf{bit}\,['n])\rightarrow\,\mathbf{bit}\,['n]:-$
|	$\mathbf{val}\,\mathbf{forall}\,\mathbf{Nat}\,'n,\,\mathbf{Nat}\,'o,\,\mathbf{Nat}\,'p.(\,\mathbf{bit}\,['n],[\|'o..'p\|])\rightarrow[\|'o..'p+2*{*}'n\|]:-$

3 Sail type system

3.1 Internal type syntax

k	::=		Internal kinds
	|	K_Typ	
	|	K_Nat	
	|	K_Ord	
	|	K_Efct	
	|	$K\_Lam(k_0 .. k_n \to k')$	
	|	K_infer	Representing an unknown kind, inferred by context

$t,\ u$	::=			Internal types
	|	x		
	|	$'x$		
	|	$t_1 \to t_2$ *effect*		
	|	$(t_1, ..., t_n)$		
	|	$x\langle t\_args\rangle$		
	|	$t \mapsto t_1$		
	|	**register** $\langle t\_arg\rangle$	S	
	|	**range** $\langle ne\ ne'\rangle$	S	
	|	**atom** $\langle ne\rangle$	S	
	|	**vector** $\langle ne\ ne'\ order\ t\rangle$	S	
	|	**list** $\langle t\rangle$	S	
	|	**reg** $\langle t\rangle$	S	
	|	**implicit** $\langle ne\rangle$	S	
	|	**bit**	S	
	|	*string*	S	
	|	**unit**	S	
	|	$t[t\_arg_1/tid_1 .. t\_arg_n/tid_n]$	M	

optx	::=		
	|	x	
	|		
tag	::=		Data indicating where the identifier arises and thus information necessary in compilation
	|	**None**	
	|	**Intro**	Denotes an assignment and lexp that introduces a binding
	|	**Set**	Denotes an expression that mutates a local variable
	|	**Global**	Globally let-bound or enumeration based value/variable
	|	**Ctor**	Data constructor from a type union
	|	**Extern** *optx*	External function, specied only with a val statement
	|	**Default**	Type has come from default declaration, identifier may not be bound locally
	|	**Spec**	
	|	**Enum** *num*	
	|	**Alias**	
	|	*Unknown_pathoptx*	Tag to distinguish an unknown path from a non-analysis non deterministic path
ne	::=		internal numeric expressions
	|	$'x$	
	|	*num*	
	|	**infinity**	
	|	$ne_1 * ne_2$	
	|	$ne_1 + ... + ne_n$	
	|	$ne_1 - ne_2$	
	|	$2 ** ne$	
	|	$(-ne)$	
	|	**zero**	S
	|	**one**	S
	|	**bitlength** (bin)	M
	|	**bitlength** (hex)	M

	$\mid$	**count** $(num_0 \ldots num_i)$	M
	$\mid$	**length** $(pat_1 \ldots pat_n)$	M
	$\mid$	**length** $(exp_1 \ldots exp_n)$	M

t_arg	::=			Argument to type constructors
	$\mid$	t		
	$\mid$	ne		
	$\mid$	*effect*		
	$\mid$	*order*		
	$\mid$	**fresh**	M	

t_args	::=		Arguments to type constructors
	$\mid$	$t\_arg_1 \ldots t\_arg_n$	

nec	::=		Numeric expression constraints
	$\mid$	$ne \leq ne'$	
	$\mid$	$ne = ne'$	
	$\mid$	$ne \geq ne'$	
	$\mid$	$'x$ **IN** $\{num_1, \ldots, num_n\}$	
	$\mid$	$nec_0 \,..\, nec_n \rightarrow nec'_0 \ldots nec'_m$	
	$\mid$	$nec_0 \ldots nec_n$	

Σ^{N}	::=			nexp constraint lists
	$\mid$	$\{nec_1, \ldots, nec_n\}$		
	$\mid$	${\Sigma^{\mathrm{N}}}_1 \uplus \ldots \uplus {\Sigma^{\mathrm{N}}}_n$	M	
	$\mid$	**consistent_increase** $ne_1\ ne'_1 \ldots ne_n\ ne'_n$	M	Generates constraints from pairs of constraints, where the first of each pair is always larger than the su
	$\mid$	**consistent_decrease** $ne_1\ ne'_1 \ldots ne_n\ ne'_n$	M	Generates constraints from pairs of constraints, where the first of each pair is always smaller than the
	$\mid$	**resolve** (Σ^{N})		

E^{D}	::=		Environments storing top level information, such as defined abbreviations, records, enumerations, and kin

	|	$\langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}} \rangle$	
	|	ϵ	
	|	$E^{\mathrm{D}} \uplus E^{\mathrm{D}\prime}$	
kinf	::=		Whether a kind is default or from a local binding
	|	k	
	|	k **default**	
tid	::=		A type identifier or type variable
	|	*id*	
	|	*kid*	
E^{K}	::=		Kind environments
	|	$\{tid_1 \mapsto kinf_1, .., tid_n \mapsto kinf_n\}$	
	|	${E^{\mathrm{K}}}_1 \uplus .. \uplus {E^{\mathrm{K}}}_n$	M In a unioning kinf, k default u k results in k (i.e. the default is locally forgotten)
	|	$E^{\mathrm{K}} \setminus {E^{\mathrm{K}}}_1 .. {E^{\mathrm{K}}}_n$	M
tinf	::=		Type variables, type, and constraints, bound to an identifier
	|	t	
	|	$E^{\mathrm{K}}, \Sigma^{\mathrm{N}}, tag, t$	
E^{A}	::=		
	|	$\{tid_1 \mapsto tinf_1, .., tid_n \mapsto tinf_n\}$	
	|	$E^{\mathrm{A}}_1 \uplus .. \uplus E^{\mathrm{A}}_n$	
field_typs	::=		Record fields
	|	$id_1 : t_1, .., id_n : t_n$	
E^{R}	::=		Record environments
	|	$\{\{field\_typs_1\} \mapsto tinf_1, .., \{field\_typs_n\} \mapsto tinf_n\}$	

	|	$E^{\mathrm{R}}_1 \uplus .. \uplus E^{\mathrm{R}}_n$	M	
enumerate_map	::=		
	|	$\{num_1 \mapsto id_1 \, ... \, num_n \mapsto id_n\}$		
E^{E}	::=			Enumeration environments
	|	$\{t_1 \mapsto enumerate\_map_1, .., t_n \mapsto enumerate\_map_n\}$		
	|	$E^{\mathrm{E}}_1 \uplus .. \uplus E^{\mathrm{E}}_n$		
E^{T}	::=			Type environments
	|	$\{id_1 \mapsto tinf_1, .., id_n \mapsto tinf_n\}$		
	|	$\{id \mapsto \textbf{overload}\ tinf\ conformsto : tinf_1, ..., tinf_n\}$		
	|	$(E^{\mathrm{T}}{}_1 \uplus .... \uplus E^{\mathrm{T}}{}_n)$	M	
	|	$\uplus E^{\mathrm{T}}{}_1 .. E^{\mathrm{T}}{}_n$	M	
	|	$E^{\mathrm{T}} \setminus id_1 .. id_n$	M	
	|	$(E^{\mathrm{T}}{}_1 \cap .... \cap E^{\mathrm{T}}{}_n)$	M	
	|	$\cap E^{\mathrm{T}}{}_1 .. E^{\mathrm{T}}{}_n$	M	
ts	::=			
	|	$t_1, .., t_n$		
E	::=			Definition environment and lexical environment
	|	$\langle E^{\mathrm{T}}, E^{\mathrm{D}} \rangle$		
	|	ϵ	M	
	|	$E \uplus E'$		
I	::=			Information given by type checking an expression
	|	$\langle \Sigma^{\mathrm{N}}, effect \rangle$		
	|	I_ϵ		Empty constraints, effect
	|	$I_1 \uplus I_2$		

	$I_1 \uplus .. \uplus I_n$	Unions the constraints and effect
formula	::=	
	judgement	
	$formula_1 \quad .. \quad formula_n$	
	$E^{\mathrm{K}}(tid) \triangleright kinf$	Kind lookup
	$E^{\mathrm{A}}(tid) \triangleright tinf$	
	$E^{\mathrm{T}}(id) \triangleright tinf$	Type lookup
	$E^{\mathrm{T}}(id) \triangleright \mathbf{overload}\ tinf : tinf_1 ... tinf_n$	Overloaded type lookup
	$E^{\mathrm{K}}(tid) < -\vert k$	Update the kind associated with id to k
	$E^{\mathrm{R}}(id_0 .. id_n) \triangleright t, ts$	Record lookup
	$E^{\mathrm{R}}(t) \triangleright id_0 : t_0 .. id_n : t_n$	Record looup by type
	$E^{\mathrm{E}}(t) \triangleright enumerate\_map$	Enumeration lookup by type
	$\mathbf{dom}\,(E^{\mathrm{T}}{}_1) \cap \mathbf{dom}\,(E^{\mathrm{T}}{}_2) = \emptyset$	
	$\mathbf{dom}\,(E^{\mathrm{K}}{}_1) \cap \mathbf{dom}\,(E^{\mathrm{K}}{}_2) = \emptyset$	
	$\mathbf{disjoint\,doms}\,(E^{\mathrm{T}}{}_1, ..., E^{\mathrm{T}}{}_n)$	Pairwise disjoint domains
	$id \notin \mathbf{dom}\,(E^{\mathrm{K}})$	
	$id \notin \mathbf{dom}\,(E^{\mathrm{T}})$	
	$id_0 : t_0 .. id_n : t_n \subset id'_0 : t'_0 .. id'_i : t'_i$	
	$num_1 < ... < num_n$	
	$num_1 > ... > num_n$	
	$exp_1 \equiv exp_2$	
	$E^{\mathrm{K}}{}_1 \equiv E^{\mathrm{K}}{}_2$	
	$E^{\mathrm{K}}{}_1 \approx E^{\mathrm{K}}{}_2$	
	$E^{\mathrm{T}}{}_1 \equiv E^{\mathrm{T}}{}_2$	
	$E^{\mathrm{R}}_1 \equiv E^{\mathrm{R}}_2$	
	$E^{\mathrm{E}}_1 \equiv E^{\mathrm{E}}_2$	
	$E^{\mathrm{D}}{}_1 \equiv E^{\mathrm{D}}{}_2$	
	$E_1 \equiv E_2$	
	$\Sigma^{\mathrm{N}}{}_1 \equiv \Sigma^{\mathrm{N}}{}_2$	

| $id \equiv' id$
| $x_1 \neq x_2$
| $lit_1 \neq lit_2$
| $I_1 \equiv I_2$
| $effect_1 \equiv effect_2$
| $t_1 \equiv t_2$
| $ne \equiv ne'$
| $kid \equiv fresh\_kid(E^{\mathrm{D}})$

3.2 Type relations

$\boxed{E^{\mathrm{K}} \vdash_t t\,\mathbf{ok}}$ Well-formed types

$$\frac{E^{\mathrm{K}}('x) \triangleright K\_Typ}{E^{\mathrm{K}} \vdash_t 'x\,\mathbf{ok}}\ \textsc{check\_t\_var}$$

$$\frac{E^{\mathrm{K}}('x) \triangleright K\_infer \quad E^{\mathrm{K}}('x) < -|K\_Typ}{E^{\mathrm{K}} \vdash_t 'x\,\mathbf{ok}}\ \textsc{check\_t\_varInfer}$$

$$\frac{E^{\mathrm{K}} \vdash_t t_1\,\mathbf{ok} \quad E^{\mathrm{K}} \vdash_t t_2\,\mathbf{ok} \quad E^{\mathrm{K}} \vdash_e effect\,\mathbf{ok}}{E^{\mathrm{K}} \vdash_t t_1 \to t_2\ effect\,\mathbf{ok}}\ \textsc{check\_t\_fn}$$

$$\frac{E^{\mathrm{K}} \vdash_t t_1\,\mathbf{ok} \quad .... \quad E^{\mathrm{K}} \vdash_t t_n\,\mathbf{ok}}{E^{\mathrm{K}} \vdash_t (t_1, ..., t_n)\,\mathbf{ok}}\ \textsc{check\_t\_tup}$$

$$\frac{E^{\mathrm{K}}(x) \triangleright K\_Lam(k_1 .. k_n \to K\_Typ) \quad E^{\mathrm{K}}, k_1 \vdash t\_arg_1\,\mathbf{ok} \quad .. \quad E^{\mathrm{K}}, k_n \vdash t\_arg_n\,\mathbf{ok}}{E^{\mathrm{K}} \vdash_t x\langle t\_arg_1 .. t\_arg_n\rangle\,\mathbf{ok}}\ \textsc{check\_t\_app}$$

$\boxed{E^{\mathrm{K}} \vdash_e effect\,\mathbf{ok}}$ Well-formed effects

$$\frac{E^{\mathrm{K}}('x) \triangleright K\_Efct}{E^{\mathrm{K}} \vdash_e 'x\,\mathbf{ok}}\ \textsc{check\_ef\_var}$$

$$\frac{E^{\mathrm{K}}('x) \triangleright K\_infer \quad E^{\mathrm{K}}('x) < -|K\_Efct}{E^{\mathrm{K}} \vdash_e {'x}\,\mathbf{ok}}\ \textsc{check\_ef\_varInfer}$$

$$\frac{}{E^{\mathrm{K}} \vdash_e \{base\_effect_1, .., base\_effect_n\}\,\mathbf{ok}}\ \textsc{check\_ef\_set}$$

$\boxed{E^{\mathrm{K}} \vdash_n ne\,\mathbf{ok}}$ Well-formed numeric expressions

$$\frac{E^{\mathrm{K}}('x) \triangleright K\_Nat}{E^{\mathrm{K}} \vdash_n {'x}\,\mathbf{ok}}\ \textsc{check\_n\_var}$$

$$\frac{E^{\mathrm{K}}('x) \triangleright K\_infer \quad E^{\mathrm{K}}('x) < -|K\_Nat}{E^{\mathrm{K}} \vdash_n {'x}\,\mathbf{ok}}\ \textsc{check\_n\_varInfer}$$

$$\frac{}{E^{\mathrm{K}} \vdash_n num\,\mathbf{ok}}\ \textsc{check\_n\_num}$$

$$\frac{E^{\mathrm{K}} \vdash_n ne_1\,\mathbf{ok} \quad E^{\mathrm{K}} \vdash_n ne_2\,\mathbf{ok}}{E^{\mathrm{K}} \vdash_n ne_1 + ne_2\,\mathbf{ok}}\ \textsc{check\_n\_sum}$$

$$\frac{E^{\mathrm{K}} \vdash_n ne_1\,\mathbf{ok} \quad E^{\mathrm{K}} \vdash_n ne_2\,\mathbf{ok}}{E^{\mathrm{K}} \vdash_n ne_1 * ne_2\,\mathbf{ok}}\ \textsc{check\_n\_mult}$$

$$\frac{E^{\mathrm{K}} \vdash_n ne\,\mathbf{ok}}{E^{\mathrm{K}} \vdash_n 2 ** ne\,\mathbf{ok}}\ \textsc{check\_n\_exp}$$

$\boxed{E^{\mathrm{K}} \vdash_o order\,\mathbf{ok}}$ Well-formed numeric expressions

$$\frac{E^{\mathrm{K}}('x) \triangleright K\_Ord}{E^{\mathrm{K}} \vdash_o {'x}\,\mathbf{ok}}\ \textsc{check\_ord\_var}$$

$$\frac{E^{\mathrm{K}}('x) \triangleright K\_infer \quad E^{\mathrm{K}}('x) < -|K\_Ord}{E^{\mathrm{K}} \vdash_o {'x}\,\mathbf{ok}}\ \textsc{check\_ord\_varInfer}$$

$\boxed{E^{\mathrm{K}}, k \vdash t\_arg\,\mathbf{ok}}$ Well-formed type arguments kind check matching the application type variable

$$\frac{E^{\mathrm{K}} \vdash_t t\ \mathbf{ok}}{E^{\mathrm{K}}, K\_Typ \vdash t\ \mathbf{ok}}\ \textsc{check\_targs\_typ}$$

$$\frac{E^{\mathrm{K}} \vdash_e effect\ \mathbf{ok}}{E^{\mathrm{K}}, K\_Efct \vdash effect\ \mathbf{ok}}\ \textsc{check\_targs\_eff}$$

$$\frac{E^{\mathrm{K}} \vdash_n ne\ \mathbf{ok}}{E^{\mathrm{K}}, K\_Nat \vdash ne\ \mathbf{ok}}\ \textsc{check\_targs\_nat}$$

$$\frac{E^{\mathrm{K}} \vdash_o order\ \mathbf{ok}}{E^{\mathrm{K}}, K\_Ord \vdash order\ \mathbf{ok}}\ \textsc{check\_targs\_ord}$$

$\boxed{E^{\mathrm{K}} \vdash kind \rightsquigarrow k}$

$$\frac{}{E^{\mathrm{K}} \vdash \mathbf{Type} \rightsquigarrow K\_Typ}\ \textsc{convert\_kind\_Typ}$$

$\boxed{E^{\mathrm{D}} \vdash quant\_item \rightsquigarrow {E^{\mathrm{K}}}_1, \Sigma^{\mathrm{N}}}$ Convert source quantifiers to kind environments and constraints

$$\frac{E^{\mathrm{K}} \vdash kind \rightsquigarrow k}{\langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}} \rangle \vdash kind\ 'x \rightsquigarrow \{'x \mapsto k\}, \{\,\}}\ \textsc{convert\_quants\_kind}$$

$$\frac{E^{\mathrm{K}}('x) \triangleright k}{\langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}} \rangle \vdash 'x \rightsquigarrow \{'x \mapsto k\}, \{\,\}}\ \textsc{convert\_quants\_nokind}$$

$$\frac{\vdash nexp_1 \rightsquigarrow ne_1 \quad \vdash nexp_2 \rightsquigarrow ne_2}{E^{\mathrm{D}} \vdash nexp_1 = nexp_2 \rightsquigarrow \{\,\}, \{ne_1 = ne_2\}}\ \textsc{convert\_quants\_eq}$$

$$\frac{\vdash nexp_1 \rightsquigarrow ne_1 \quad \vdash nexp_2 \rightsquigarrow ne_2}{E^{\mathrm{D}} \vdash nexp_1 \geq nexp_2 \rightsquigarrow \{\,\}, \{ne_1 \geq ne_2\}}\ \textsc{convert\_quants\_gteq}$$

$$\frac{\vdash nexp_1 \rightsquigarrow ne_1 \quad \vdash nexp_2 \rightsquigarrow ne_2}{E^{\mathrm{D}} \vdash nexp_1 \leq nexp_2 \rightsquigarrow \{\,\}, \{ne_1 \leq ne_2\}}\ \textsc{convert\_quants\_lteq}$$

$$\frac{}{E^{\mathrm{D}} \vdash \text{'}x\,\mathbf{IN}\,\{num_1,\ ...\ , num_n\} \rightsquigarrow \{\,\}, \{\text{'}x\,\mathbf{IN}\,\{num_1,\ ...\ , num_n\}\}}\ \textsc{convert\_quants\_in}$$

$\boxed{E^{\mathrm{D}} \vdash typschm \rightsquigarrow t, E^{\mathrm{K}}, \Sigma^{\mathrm{N}}}$ Convert source types with typeschemes to internal types and kind environments

$$\frac{E^{\mathrm{D}} \vdash typ \rightsquigarrow t}{E^{\mathrm{D}} \vdash\ typ \rightsquigarrow t, \{\,\}, \{\,\}}\ \textsc{convert\_typschm\_noquant}$$

$$\frac{\begin{array}{l} E^{\mathrm{D}} \vdash quant\_item_1 \rightsquigarrow {E^{\mathrm{K}}}_1, {\Sigma^{\mathrm{N}}}_1 \quad ... \quad E^{\mathrm{D}} \vdash quant\_item_n \rightsquigarrow {E^{\mathrm{K}}}_n, {\Sigma^{\mathrm{N}}}_n \\ E^{\mathrm{K}} \equiv {E^{\mathrm{K}}}_1 \uplus\ ...\ \uplus {E^{\mathrm{K}}}_n \\ E^{\mathrm{D}} \uplus \langle E^{\mathrm{K}}, \{\,\}, \{\,\}, \{\,\}\rangle \vdash typ \rightsquigarrow t \end{array}}{E^{\mathrm{D}} \vdash \mathbf{forall}\, quant\_item_1,\ ...\ , quant\_item_n.\ typ \rightsquigarrow t, E^{\mathrm{K}}, {\Sigma^{\mathrm{N}}}_1 \uplus\ ...\ \uplus {\Sigma^{\mathrm{N}}}_n}\ \textsc{convert\_typschm\_quant}$$

$\boxed{E^{\mathrm{D}} \vdash typ \rightsquigarrow t}$ Convert source types to internal types

$$\frac{E^{\mathrm{K}}(\text{'}x) \triangleright K\_Typ}{\langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle \vdash \text{'}x \rightsquigarrow \text{'}x}\ \textsc{convert\_typ\_var}$$

$$\frac{E^{\mathrm{K}}(x) \triangleright K\_Typ}{\langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle \vdash x \rightsquigarrow x}\ \textsc{convert\_typ\_id}$$

$$\frac{\begin{array}{c} E^{\mathrm{D}} \vdash typ_1 \rightsquigarrow t_1 \\ E^{\mathrm{D}} \vdash typ_2 \rightsquigarrow t_2 \end{array}}{E^{\mathrm{D}} \vdash typ_1 \rightarrow typ_2\, \mathbf{effect}\ effect \rightsquigarrow t_1 \rightarrow t_2\ effect}\ \textsc{convert\_typ\_fn}$$

$$\frac{E^{\mathrm{D}} \vdash typ_1 \rightsquigarrow t_1 \quad .... \quad E^{\mathrm{D}} \vdash typ_n \rightsquigarrow t_n}{E^{\mathrm{D}} \vdash (typ_1,\ ....\ , typ_n) \rightsquigarrow (t_1,\ ....\ , t_n)}\ \textsc{convert\_typ\_tup}$$

$$\frac{\begin{array}{l} E^{\mathrm{K}}(x) \triangleright K\_Lam(k_1\ ..\ k_n \rightarrow K\_Typ) \\ \langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle, k_1 \vdash typ\_arg_1 \rightsquigarrow t\_arg_1 \quad ..\quad \langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle, k_n \vdash typ\_arg_n \rightsquigarrow t\_arg_n \end{array}}{\langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle \vdash x\langle typ\_arg_1,\ ..\ , typ\_arg_n\rangle \rightsquigarrow x\langle t\_arg_1\ ..\ t\_arg_n\rangle}\ \textsc{convert\_typ\_app}$$

$\boxed{E^{\mathrm{D}}, k \vdash typ\_arg \rightsquigarrow t\_arg}$ Convert source type arguments to internals

$$\frac{E^{\mathrm{D}} \vdash typ \rightsquigarrow t}{E^{\mathrm{D}}, K\_Typ \vdash typ \rightsquigarrow t}\ \textsc{convert\_targ\_typ}$$

$\boxed{\vdash nexp \rightsquigarrow ne}$ Convert and normalize numeric expressions

$$\frac{}{\vdash {'x} \rightsquigarrow {'x}}\ \textsc{convert\_nexp\_var}$$

$$\frac{}{\vdash num \rightsquigarrow num}\ \textsc{convert\_nexp\_num}$$

$$\frac{\vdash nexp_1 \rightsquigarrow ne_1 \quad \vdash nexp_2 \rightsquigarrow ne_2}{\vdash nexp_1 * nexp_2 \rightsquigarrow ne_1 * ne_2}\ \textsc{convert\_nexp\_mult}$$

$$\frac{\vdash nexp_1 \rightsquigarrow ne_1 \quad \vdash nexp_2 \rightsquigarrow ne_2}{\vdash nexp_1 + nexp_2 \rightsquigarrow ne_1 + ne_2}\ \textsc{convert\_nexp\_add}$$

$$\frac{\vdash nexp \rightsquigarrow ne}{\vdash 2 ** nexp \rightsquigarrow 2 ** ne}\ \textsc{convert\_nexp\_exp}$$

$\boxed{E^{\mathrm{D}} \vdash t \approx t'}$

$$\frac{E^{\mathrm{K}} \vdash_t t\ \mathbf{ok}}{\langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}} \rangle \vdash t \approx t}\ \textsc{conforms\_to\_refl}$$

$$\frac{E^{\mathrm{D}} \vdash t_1 \approx t_2 \quad E^{\mathrm{D}} \vdash t_2 \approx t_3}{E^{\mathrm{D}} \vdash t_1 \approx t_3}\ \textsc{conforms\_to\_trans}$$

$$\frac{}{E^{\mathrm{D}} \vdash {'x} \approx t}\ \textsc{conforms\_to\_var}$$

$$\frac{}{E^{\mathrm{D}} \vdash t \approx {'x}}\ \textsc{conforms\_to\_var2}$$

$$\frac{E^{\mathrm{A}}(x) \triangleright u \quad \langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}} \rangle \vdash u \approx t}{\langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}} \rangle \vdash x \approx t}\ \textsc{conforms\_to\_abbrev}$$

$$\frac{E^{\mathrm{A}}(x) \triangleright u \quad \langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}} \rangle \vdash t \approx u}{\langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}} \rangle \vdash t \approx x}\ \textsc{conforms\_to\_abbrev2}$$

$$\frac{E^{\mathrm{D}} \vdash t_1 \approx u_1 \quad .... \quad E^{\mathrm{D}} \vdash t_n \approx u_n}{E^{\mathrm{D}} \vdash (t_1, \ldots, t_n) \approx (u_1, \ldots, u_n)} \quad \textsc{conforms\_to\_tup}$$

$$\frac{\begin{array}{l} E^{\mathrm{K}}(x) \triangleright K\_Lam(k_1 .. k_n \to K\_Typ) \\ \langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}} \rangle, k_1 \vdash t\_arg_1 \approx t\_arg'_1 \quad .. \quad \langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}} \rangle, k_n \vdash t\_arg_n \approx t\_arg'_n \end{array}}{\langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}} \rangle \vdash x\langle t\_arg_1 .. t\_arg_n \rangle \approx x\langle t\_arg'_1 .. t\_arg'_n \rangle} \quad \textsc{conforms\_to\_app}$$

$$\frac{\begin{array}{l} x' \neq x \\ E^{\mathrm{A}}(x') \triangleright \{tid_1 \mapsto kinf_1, .., tid_m \mapsto kinf_m\}, \Sigma^{\mathrm{N}}, tag, u \\ \langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}} \rangle \vdash x\langle t\_arg_1 .. t\_arg_n \rangle \approx u[t\_arg'_1/tid_1 .. t\_arg'_m/tid_m] \end{array}}{\langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}} \rangle \vdash x\langle t\_arg_1 .. t\_arg_n \rangle \approx x'\langle t\_arg'_1 .. t\_arg'_m \rangle} \quad \textsc{conforms\_to\_appAbbrev}$$

$$\frac{\begin{array}{l} x' \neq x \\ E^{\mathrm{A}}(x') \triangleright \{tid_1 \mapsto kinf_1, .., tid_n \mapsto kinf_n\}, \Sigma^{\mathrm{N}}, tag, u \\ \langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}} \rangle \vdash u[t\_arg_1/tid_1 .. t\_arg_n/tid_n] \approx x\langle t\_arg'_1 .. t\_arg'_m \rangle \end{array}}{\langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}} \rangle \vdash x'\langle t\_arg_1 .. t\_arg_n \rangle \approx x\langle t\_arg'_1 .. t\_arg'_m \rangle} \quad \textsc{conforms\_to\_appAbbrev2}$$

$$\frac{E^{\mathrm{D}} \vdash t \approx u}{E^{\mathrm{D}} \vdash \mathbf{register}\, \langle t \rangle \approx u} \quad \textsc{conforms\_to\_register}$$

$\boxed{E^{\mathrm{D}}, k \vdash t\_arg \approx t\_arg'}$

$$\frac{E^{\mathrm{D}} \vdash t \approx t'}{E^{\mathrm{D}}, K\_Typ \vdash t \approx t'} \quad \textsc{targconforms\_typ}$$

$$\frac{}{E^{\mathrm{D}}, K\_Nat \vdash ne \approx ne'} \quad \textsc{targconforms\_nexp}$$

$\boxed{\sigma_{conformsto(t,t')}(tinflist) \triangleright tinflist'}$

$$\frac{\begin{array}{c} E^{\mathrm{D}} \vdash t_i \approx t'_i \\ E^{\mathrm{D}} \vdash t'_j \approx t_j \\ \sigma_{\mathbf{full}(t_i,t_j)}(tinf_0 .. tinf_m\ tinf'_0 .. tinf'_n) \triangleright \epsilon \end{array}}{\sigma_{\mathbf{full}(t_i,t_j)}(tinf_0 .. tinf_m\ E^{\mathrm{K}}, \Sigma^{\mathrm{N}}, tag, t'_i \to t'_j\ effect\ tinf'_0 .. tinf'_n) \triangleright E^{\mathrm{K}}, \Sigma^{\mathrm{N}}, tag, t'_i \to t'_j} \quad \textsc{so\_full}$$

$$\frac{\begin{array}{c} E^{\mathrm{D}} \vdash t_i \approx t'_i \\ \sigma_{\mathbf{parm}(t_i,t_j)}(tinf_0 .. tinf_m) \triangleright \epsilon \end{array}}{\sigma_{\mathbf{parm}(t_i,t_j)}(tinf_0 .. tinf_m\ E^{\mathrm{K}}, \Sigma^{\mathrm{N}}, tag, t'_i \to t\ effect\ tinf'_0 .. tinf'_n) \triangleright E^{\mathrm{K}}, \Sigma^{\mathrm{N}}, tag, t'_i \to t} \quad \textsc{so\_parm}$$

$\boxed{E^{\mathrm{D}} \vdash t \precapprox t', \Sigma^{\mathrm{N}}}$

$$\frac{E^{\mathrm{K}} \vdash_t t\ \mathbf{ok}}{\langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle \vdash t \precapprox t, \{\,\}}\ \textsc{consistent\_typ\_refl}$$

$$\frac{E^{\mathrm{D}} \vdash t_1 \precapprox t_2, \Sigma^{\mathrm{N}}{}_1 \quad E^{\mathrm{D}} \vdash t_2 \precapprox t_3, \Sigma^{\mathrm{N}}{}_2}{E^{\mathrm{D}} \vdash t_1 \precapprox t_3, \Sigma^{\mathrm{N}}{}_1 \uplus \Sigma^{\mathrm{N}}{}_2}\ \textsc{consistent\_typ\_trans}$$

$$\frac{E^{\mathrm{A}}(x) \triangleright \{\,\}, \Sigma^{\mathrm{N}}{}_1, tag, u \quad \langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle \vdash u \precapprox t, \Sigma^{\mathrm{N}}}{\langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle \vdash x \precapprox t, \Sigma^{\mathrm{N}} \uplus \Sigma^{\mathrm{N}}{}_1}\ \textsc{consistent\_typ\_abbrev}$$

$$\frac{E^{\mathrm{A}}(x) \triangleright \{\,\}, \Sigma^{\mathrm{N}}{}_1, tag, u \quad \langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle \vdash t \precapprox u, \Sigma^{\mathrm{N}}}{\langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle \vdash t \precapprox x, \Sigma^{\mathrm{N}} \uplus \Sigma^{\mathrm{N}}{}_1}\ \textsc{consistent\_typ\_abbrev2}$$

$$\frac{}{E^{\mathrm{D}} \vdash \text{'}x \precapprox t, \{\,\}}\ \textsc{consistent\_typ\_var}$$

$$\frac{}{E^{\mathrm{D}} \vdash t \precapprox \text{'}x, \{\,\}}\ \textsc{consistent\_typ\_var2}$$

$$\frac{E^{\mathrm{D}} \vdash t_1 \precapprox u_1, \Sigma^{\mathrm{N}}{}_1 \quad .... \quad E^{\mathrm{D}} \vdash t_n \precapprox u_n, \Sigma^{\mathrm{N}}{}_n}{E^{\mathrm{D}} \vdash (t_1, \ldots., t_n) \precapprox (u_1, \ldots., u_n), \Sigma^{\mathrm{N}}{}_1 \uplus \ldots. \uplus \Sigma^{\mathrm{N}}{}_n}\ \textsc{consistent\_typ\_tup}$$

$$\frac{}{E^{\mathrm{D}} \vdash \mathbf{range}\,\langle ne_1\ ne_2\rangle \precapprox \mathbf{range}\,\langle ne_3\ ne_4\rangle, \{ne_3 \leq ne_1, ne_2 \leq ne_4\}}\ \textsc{consistent\_typ\_range}$$

$$\frac{}{E^{\mathrm{D}} \vdash \mathbf{atom}\,\langle ne\rangle \precapprox \mathbf{range}\,\langle ne_1\ ne_2\rangle, \{ne_1 \leq ne, ne \leq ne_2\}}\ \textsc{consistent\_typ\_atomRange}$$

$$\frac{}{E^{\mathrm{D}} \vdash \mathbf{atom}\,\langle ne_1\rangle \precapprox \mathbf{atom}\,\langle ne_2\rangle, \{ne_1 = ne_2\}}\ \textsc{consistent\_typ\_atom}$$

$$\frac{}{E^{\mathrm{D}} \vdash \mathbf{range}\,\langle ne_1\ ne_2\rangle \precapprox \mathbf{atom}\,\langle \text{'}x\rangle, \{ne_1 \leq \text{'}x, \text{'}x \leq ne_2\}}\ \textsc{consistent\_typ\_rangeAtom}$$

$$\frac{E^{\mathrm{D}} \vdash t \precapprox t', \Sigma^{\mathrm{N}}}{E^{\mathrm{D}} \vdash \mathbf{vector}\,\langle ne_1\ ne_2\ order\ t\rangle \precapprox \mathbf{vector}\,\langle ne_3\ ne_4\ order\ t'\rangle, \{ne_1 = ne_3, ne_2 = ne_3\} \uplus \Sigma^{\mathrm{N}}}\ \textsc{consistent\_typ\_vector}$$

$$\frac{E^{\mathrm{K}}(x) \triangleright K\_Lam(k_1 \,..\, k_n \to K\_Typ) \quad \langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle, k_1 \vdash t\_arg_1 \gtrapprox t\_arg'_1, \Sigma^{\mathrm{N}}{}_1 \quad .. \quad \langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle, k_n \vdash t\_arg_n \gtrapprox t\_arg'_n, \Sigma^{\mathrm{N}}{}_n}{\langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle \vdash x\langle t\_arg_1 \,..\, t\_arg_n\rangle \gtrapprox x\langle t\_arg'_1 \,..\, t\_arg'_n\rangle, \Sigma^{\mathrm{N}}{}_1 \uplus \,..\, \uplus \Sigma^{\mathrm{N}}{}_n} \text{ CONSISTENT\_TYP\_APP}$$

$$\frac{x' \neq x \quad E^{\mathrm{A}}(x') \triangleright \{tid_1 \mapsto kinf_1, \,..\, tid_m \mapsto kinf_m\}, \Sigma^{\mathrm{N}}, tag, u \quad \langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle \vdash x\langle t\_arg_1 \,..\, t\_arg_n\rangle \gtrapprox u[t\_arg'_1/tid_1 \,..\, t\_arg'_m/tid_m], \Sigma^{\mathrm{N}}{}_2}{\langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle \vdash x\langle t\_arg_1 \,..\, t\_arg_n\rangle \gtrapprox x'\langle t\_arg'_1 \,..\, t\_arg'_m\rangle, \Sigma^{\mathrm{N}} \uplus \Sigma^{\mathrm{N}}{}_2} \text{ CONSISTENT\_TYP\_APPABBREV}$$

$$\frac{x' \neq x \quad E^{\mathrm{A}}(x') \triangleright \{tid_1 \mapsto kinf_1, \,..\, tid_m \mapsto kinf_m\}, \Sigma^{\mathrm{N}}, tag, u \quad \langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle \vdash u[t\_arg'_1/tid_1 \,..\, t\_arg'_m/tid_m] \gtrapprox x\langle t\_arg_1 \,..\, t\_arg_n\rangle, \Sigma^{\mathrm{N}}{}_2}{\langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle \vdash x'\langle t\_arg'_1 \,..\, t\_arg'_m\rangle \gtrapprox x\langle t\_arg_1 \,..\, t\_arg_n\rangle, \Sigma^{\mathrm{N}} \uplus \Sigma^{\mathrm{N}}{}_2} \text{ CONSISTENT\_TYP\_APPABBREV2}$$

$\boxed{E^{\mathrm{D}}, k \vdash t\_arg \gtrapprox t\_arg', \Sigma^{\mathrm{N}}}$

$$\frac{E^{\mathrm{D}} \vdash t \gtrapprox t', \Sigma^{\mathrm{N}}}{E^{\mathrm{D}}, K\_Typ \vdash t \gtrapprox t', \Sigma^{\mathrm{N}}} \text{ TARG\_CONSISTENT\_TYP}$$

$$\frac{}{E^{\mathrm{D}}, K\_Nat \vdash ne \gtrapprox ne', \{ne = ne'\}} \text{ TARG\_CONSISTENT\_NEXP}$$

$\boxed{E^{\mathrm{D}}, t' \vdash exp : t \triangleright t'', exp', \Sigma^{\mathrm{N}}, effect}$

$$\frac{E^{\mathrm{D}}, u_1 \vdash id_1 : t_1 \triangleright u_1, exp_1, \Sigma^{\mathrm{N}}{}_1, effect_1 \quad .. \quad E^{\mathrm{D}}, u_n \vdash id_n : t_n \triangleright u_n, exp_n, \Sigma^{\mathrm{N}}{}_n, effect_n \quad exp' \equiv \mathbf{switch}\ exp\{\ \mathbf{case}\ (id_1, \,..\, , id_n) \to (exp_1, \,..\, , exp_n)\}}{E^{\mathrm{D}}, (u_1, \,..\, , u_n) \vdash exp : (t_1, \,..\, , t_n) \triangleright (u_1, \,..\, , u_n), exp', \Sigma^{\mathrm{N}}{}_1 \uplus \,..\, \uplus \Sigma^{\mathrm{N}}{}_n, \mathbf{pure}} \text{ COERCE\_TYP\_TUPLE}$$

$$\frac{E^{\mathrm{D}} \vdash u \gtrapprox t, \Sigma^{\mathrm{N}} \quad exp' \equiv (annot)\, exp}{E^{\mathrm{D}}, \mathbf{vector}\ \langle ne_1\ ne_2\ order\ t\rangle \vdash exp : \mathbf{vector}\ \langle ne_3\ ne_4\ order\ u\rangle \triangleright \mathbf{vector}\ \langle ne_3\ ne_4\ order\ t\rangle, exp', \Sigma^{\mathrm{N}} \uplus \{ne_2 = ne_4\}, \mathbf{pure}} \text{ COERCE\_TYP\_VECTORUPDATESTART}$$

$$\frac{E^{\mathrm{D}} \vdash u \gtrapprox \mathbf{bit}, \Sigma^{\mathrm{N}} \quad exp' \equiv to\_num\ exp}{E^{\mathrm{D}}, \mathbf{range}\ \langle ne_1\ ne_2\rangle \vdash exp : \mathbf{vector}\ \langle ne_3\ ne_4\ order\ u\rangle \triangleright \mathbf{range}\ \langle ne_1\ ne_2\rangle, exp', \Sigma^{\mathrm{N}} \uplus \{ne_1 = \mathbf{zero}, ne_2 \geq 2 ** ne_4\}, \mathbf{pure}} \text{ COERCE\_TYP\_TONUM}$$

$$\frac{exp' \equiv to\_vec\ exp}{E^{\mathrm{D}}, \mathbf{vector}\ \langle ne_1\ ne_2\ order\ \mathbf{bit}\rangle \vdash exp : \mathbf{range}\ \langle ne_3\ ne_4\rangle \triangleright \mathbf{vector}\ \langle ne_1\ ne_2\ order\ \mathbf{bit}\rangle, exp', \{ne_3 = \mathbf{zero}, ne_4 \leq 2 ** ne_2\}, \mathbf{pure}} \text{ COERCE\_TYP\_FROMNUM}$$

$$\frac{\begin{array}{c} E^{\mathrm{D}} \vdash typ \rightsquigarrow t \\ exp' \equiv (typ)\, exp \\ E^{\mathrm{D}}, u \vdash exp' : t \triangleright t', exp'', \Sigma^{\mathrm{N}}, \mathbf{pure} \end{array}}{E^{\mathrm{D}}, u \vdash exp : \mathbf{register}\,\langle t \rangle \triangleright t', exp'', \Sigma^{\mathrm{N}}, \{\mathbf{rreg}\}} \quad \textsc{coerce\_typ\_readReg}$$

$$\frac{exp' \equiv exp[numZero]}{E^{\mathrm{D}}, \mathbf{bit} \vdash exp : \mathbf{vector}\,\langle ne_1\ ne_2\ order\ \mathbf{bit} \rangle \triangleright \mathbf{bit}, exp', \{ne_1 = \mathbf{one}\}, \mathbf{pure}} \quad \textsc{coerce\_typ\_accessVecBit}$$

$$\frac{\begin{array}{l} E^{\mathrm{D}} \vdash \mathbf{range}\,\langle \mathbf{zero}\,\mathbf{one} \rangle \precsim\!\approx \mathbf{range}\,\langle ne_1\ ne_2 \rangle, \Sigma^{\mathrm{N}} \\ exp' \equiv \mathbf{switch}\ exp \{\ \mathbf{case}\,\mathbf{bitzero} \rightarrow numZero\ \mathbf{case}\,\mathbf{bitone} \rightarrow numOne\ \} \end{array}}{E^{\mathrm{D}}, \mathbf{range}\,\langle ne_1\ ne_2 \rangle \vdash exp : \mathbf{bit} \triangleright \mathbf{range}\,\langle ne_1\ ne_2 \rangle, exp', \Sigma^{\mathrm{N}}, \mathbf{pure}} \quad \textsc{coerce\_typ\_bitToNum}$$

$$\frac{\begin{array}{l} E^{\mathrm{D}} \vdash \mathbf{range}\,\langle ne_1\ ne_2 \rangle \precsim\!\approx \mathbf{range}\,\langle \mathbf{zero}\,\mathbf{one} \rangle, \Sigma^{\mathrm{N}} \\ exp' \equiv \mathbf{switch}\ exp \{\ \mathbf{case}\ numZero \rightarrow \mathbf{bitzero}\ \mathbf{case}\ numOne \rightarrow \mathbf{bitone}\ \} \end{array}}{E^{\mathrm{D}}, \mathbf{bit} \vdash \mathbf{range} : \mathbf{range}\,\langle ne_1\ ne_2 \rangle \triangleright \mathbf{bit}, exp', \Sigma^{\mathrm{N}}, \mathbf{pure}} \quad \textsc{coerce\_typ\_numToBit}$$

$$\frac{\begin{array}{l} E^{\mathrm{E}}(x) \triangleright \{\ \overline{num_i \mapsto id_i}^{\,i}\ \} \\ exp' \equiv \mathbf{switch}\ exp \{\ \overline{\mathbf{case}\ num_i \rightarrow id_i}^{\,i}\ \} \\ ne_3 \equiv \mathbf{count}\ (\ \overline{num_i}^{\,i}\ ) \end{array}}{\langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}} \rangle, x \vdash exp : \mathbf{range}\,\langle ne_1\ ne_2 \rangle \triangleright x, exp', \{ne_1 \leq \mathbf{zero}, ne_2 \leq ne_3\}, \mathbf{pure}} \quad \textsc{coerce\_typ\_toEnumerate}$$

$$\frac{\begin{array}{l} E^{\mathrm{E}}(x) \triangleright \{\ \overline{num_i \mapsto id_i}^{\,i}\ \} \\ exp' \equiv \mathbf{switch}\ exp \{\ \overline{\mathbf{case}\ id_i \rightarrow num_i}^{\,i}\ \} \\ ne_3 \equiv \mathbf{count}\ (\ \overline{num_i}^{\,i}\ ) \\ \langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}} \rangle \vdash \mathbf{range}\,\langle \mathbf{zero}\ ne_3 \rangle \precsim\!\approx \mathbf{range}\,\langle ne_1\ ne_2 \rangle, \Sigma^{\mathrm{N}} \end{array}}{\langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}} \rangle, \mathbf{range}\,\langle ne_1\ ne_2 \rangle \vdash exp : x \triangleright \mathbf{range}\,\langle \mathbf{zero}\ ne_3 \rangle, exp', \Sigma^{\mathrm{N}}, \mathbf{pure}} \quad \textsc{coerce\_typ\_fromEnumerate}$$

$$\frac{E^{\mathrm{D}} \vdash t \precsim\!\approx u, \Sigma^{\mathrm{N}}}{E^{\mathrm{D}}, u \vdash exp : t \triangleright t, exp, \Sigma^{\mathrm{N}}, \mathbf{pure}} \quad \textsc{coerce\_typ\_eq}$$

$\boxed{t \vdash lit : t' \Rightarrow exp, \Sigma^{\mathrm{N}}}$ Typing literal constants, coercing to expected type t

$$\frac{}{\mathbf{range}\,\langle ne\ ne' \rangle \vdash num : \mathbf{atom}\,\langle num \rangle \Rightarrow num, \{ne \leq num, num \leq ne'\}} \quad \textsc{check\_lit\_num}$$

$$\frac{}{\mathbf{vector}\,\langle ne\ ne'\ order\ \mathbf{bit} \rangle \vdash num : \mathbf{atom}\,\langle num \rangle \Rightarrow to\_vec\ num, \{num + \mathbf{one} \leq 2 ** ne'\}} \quad \textsc{check\_lit\_numToVec}$$

$$\frac{}{\mathbf{bit} \vdash \mathit{numZero} : \mathbf{atom}\,\langle \mathbf{zero} \rangle \Rightarrow \mathbf{bitzero}, \{\,\}} \quad \text{CHECK\_LIT\_NUMBITZERO}$$

$$\frac{}{\mathbf{bit} \vdash \mathit{numOne} : \mathbf{atom}\,\langle \mathbf{one} \rangle \Rightarrow \mathbf{bitone}, \{\,\}} \quad \text{CHECK\_LIT\_NUMBITONE}$$

$$\frac{}{\mathit{string} \vdash \mathit{string} : \mathit{string} \Rightarrow \mathit{string}, \{\,\}} \quad \text{CHECK\_LIT\_STRING}$$

$$\frac{ne \equiv \mathbf{bitlength}\,(\mathit{hex})}{\mathbf{vector}\,\langle \mathit{ne}_1\ \mathit{ne}_2\ \mathit{order}\ \mathbf{bit} \rangle \vdash \mathit{hex} : \mathbf{vector}\,\langle \mathit{ne}_1\ \mathit{ne}\ \mathit{order}\ \mathbf{bit} \rangle \Rightarrow \mathit{hex}, \{ne = \mathit{ne}_2\}} \quad \text{CHECK\_LIT\_HEX}$$

$$\frac{ne \equiv \mathbf{bitlength}\,(\mathit{bin})}{\mathbf{vector}\,\langle \mathit{ne}_1\ \mathit{ne}_2\ \mathit{order}\ \mathbf{bit} \rangle \vdash \mathit{bin} : \mathbf{vector}\,\langle \mathit{ne}_1\ \mathit{ne}\ \mathit{order}\ \mathbf{bit} \rangle \Rightarrow \mathit{bin}, \{ne = \mathit{ne}_2\}} \quad \text{CHECK\_LIT\_BIN}$$

$$\frac{}{\mathbf{unit} \vdash () : \mathbf{unit} \Rightarrow \mathbf{unit}, \{\,\}} \quad \text{CHECK\_LIT\_UNIT}$$

$$\frac{}{\mathbf{bit} \vdash \mathbf{bitzero} : \mathbf{bit} \Rightarrow \mathbf{bitzero}, \{\,\}} \quad \text{CHECK\_LIT\_BITZERO}$$

$$\frac{}{\mathbf{bit} \vdash \mathbf{bitone} : \mathbf{bit} \Rightarrow \mathbf{bitzero}, \{\,\}} \quad \text{CHECK\_LIT\_BITONE}$$

$$\frac{}{t \vdash \mathbf{undefined} : t \Rightarrow \mathbf{undefined}, \{\,\}} \quad \text{CHECK\_LIT\_UNDEF}$$

$\boxed{E, t \vdash \mathit{pat} : t' \triangleright \mathit{pat}', E^{\mathrm{T}}, \Sigma^{\mathrm{N}}}$ Typing patterns, building their binding environment

$$\frac{\begin{array}{l}\mathit{lit} \neq \mathbf{undefined} \\ t \vdash \mathit{lit} : u \Rightarrow \mathit{lit}', \Sigma^{\mathrm{N}} \\ E^{\mathrm{D}} \vdash u \precsim t, \Sigma^{\mathrm{N}\prime}\end{array}}{\langle E^{\mathrm{T}}, E^{\mathrm{D}} \rangle, t \vdash \mathit{lit} : u \triangleright \mathit{lit}', \{\,\}, \Sigma^{\mathrm{N}} \uplus \Sigma^{\mathrm{N}\prime}} \quad \text{CHECK\_PAT\_LIT}$$

$$\frac{}{E, t \vdash \_ : t \triangleright \_, \{\,\}, \{\,\}} \quad \text{CHECK\_PAT\_WILD}$$

$$\frac{\begin{array}{l}E, t \vdash \mathit{pat} : u \triangleright \mathit{pat}', {E^{\mathrm{T}}}_1, \Sigma^{\mathrm{N}} \\ \mathit{id} \notin \mathbf{dom}\,({E^{\mathrm{T}}}_1)\end{array}}{E, t \vdash (\mathit{pat}\ \mathbf{as}\ \mathit{id}) : u \triangleright (\mathit{pat}'\ \mathbf{as}\ \mathit{id}), ({E^{\mathrm{T}}}_1 \uplus \{\mathit{id} \mapsto t\}), \Sigma^{\mathrm{N}}} \quad \text{CHECK\_PAT\_AS}$$

$$\frac{\begin{array}{l}\langle E^{\mathrm{T}}, E^{\mathrm{D}}\rangle, t' \vdash pat : t \triangleright pat', E^{\mathrm{T}}{}_1, \Sigma^{\mathrm{N}} \\ E^{\mathrm{T}}(id) \triangleright \{\,\}, \{\,\}, \mathbf{Default}, t' \\ E^{\mathrm{D}} \vdash t' \precapprox u, \Sigma^{\mathrm{N}\prime}\end{array}}{\langle E^{\mathrm{T}}, E^{\mathrm{D}}\rangle, u \vdash (pat\ \mathbf{as}\ id) : t \triangleright (pat'\ \mathbf{as}\ id), (E^{\mathrm{T}}{}_1 \uplus \{id \mapsto t'\}), \Sigma^{\mathrm{N}} \uplus \Sigma^{\mathrm{N}\prime}}\ \textsc{check\_pat\_asDefault}$$

$$\frac{\begin{array}{l}E^{\mathrm{D}} \vdash typ \rightsquigarrow t \\ \langle E^{\mathrm{T}}, E^{\mathrm{D}}\rangle, t \vdash pat : t \triangleright pat', E^{\mathrm{T}}{}_1, \Sigma^{\mathrm{N}}\end{array}}{\langle E^{\mathrm{T}}, E^{\mathrm{D}}\rangle, u \vdash (typ)pat : t \triangleright pat', E^{\mathrm{T}}{}_1, \Sigma^{\mathrm{N}}}\ \textsc{check\_pat\_typ}$$

$$\frac{\begin{array}{l}E^{\mathrm{T}}(id) \triangleright \{tid_1 \mapsto kinf_1, \ldots, tid_m \mapsto kinf_m\}, \Sigma^{\mathrm{N}}, \mathbf{Ctor}, (u'_1, \ldots, u'_n) \to x\langle t\_arg_1 \ldots t\_arg_m\rangle\ \mathbf{pure} \\ (u_1, \ldots, u_n) \to x\langle t\_args'\rangle\ \mathbf{pure} \equiv (u'_1, \ldots, u'_n) \to x\langle t\_args\rangle\ \mathbf{pure}[t\_arg_1/tid_1 \ldots t\_arg_m/tid_m] \\ \langle E^{\mathrm{T}}, E^{\mathrm{D}}\rangle, u_1 \vdash pat_1 : t_1 \triangleright pat'_1, E^{\mathrm{T}}{}_1, \Sigma^{\mathrm{N}}{}_1 \quad \ldots \quad \langle E^{\mathrm{T}}, E^{\mathrm{D}}\rangle, u_n \vdash pat_n : t_n \triangleright pat'_n, E^{\mathrm{T}}{}_n, \Sigma^{\mathrm{N}}{}_n \\ \mathbf{disjoint\ doms}\,(E^{\mathrm{T}}{}_1, \ldots, E^{\mathrm{T}}{}_n) \\ E^{\mathrm{D}} \vdash x\langle t\_args'\rangle \precapprox t, \Sigma^{\mathrm{N}}\end{array}}{\langle E^{\mathrm{T}}, E^{\mathrm{D}}\rangle, t \vdash id(pat_1, \ldots, pat_n) : x\langle t\_args'\rangle \triangleright id(pat'_1, \ldots, pat'_n), \uplus E^{\mathrm{T}}{}_1 \ldots E^{\mathrm{T}}{}_n, \Sigma^{\mathrm{N}} \uplus \Sigma^{\mathrm{N}}{}_1 \uplus \ldots \uplus \Sigma^{\mathrm{N}}{}_n}\ \textsc{check\_pat\_constr}$$

$$\frac{\begin{array}{l}E^{\mathrm{T}}(id) \triangleright \{tid_1 \mapsto kinf_1, \ldots, tid_m \mapsto kinf_m\}, \Sigma^{\mathrm{N}}, \mathbf{Ctor}, \mathbf{unit} \to x\langle t\_arg_1 \ldots t\_arg_m\rangle\ \mathbf{pure} \\ \mathbf{unit} \to x\langle t\_args'\rangle\ \mathbf{pure} \equiv \mathbf{unit} \to x\langle t\_args\rangle\ \mathbf{pure}[t\_arg_1/tid_1 \ldots t\_arg_m/tid_m] \\ E^{\mathrm{D}} \vdash x\langle t\_args'\rangle \precapprox t, \Sigma^{\mathrm{N}}\end{array}}{\langle E^{\mathrm{T}}, E^{\mathrm{D}}\rangle, t \vdash id : t \triangleright id, \{\,\}, \Sigma^{\mathrm{N}}}\ \textsc{check\_pat\_identConstr}$$

$$\frac{\begin{array}{l}E^{\mathrm{T}}(id) \triangleright \{\,\}, \{\,\}, \mathbf{Default}, t \\ E^{\mathrm{D}} \vdash t \precapprox u, \Sigma^{\mathrm{N}}\end{array}}{\langle E^{\mathrm{T}}, E^{\mathrm{D}}\rangle, u \vdash id : t \triangleright id, (E^{\mathrm{T}} \uplus \{id \mapsto t\}), \Sigma^{\mathrm{N}}}\ \textsc{check\_pat\_varDefault}$$

$$\frac{}{\langle E^{\mathrm{T}}, E^{\mathrm{D}}\rangle, t \vdash id : t \triangleright id, (E^{\mathrm{T}} \uplus \{id \mapsto t\}), \{\,\}}\ \textsc{check\_pat\_var}$$

$$\frac{\begin{array}{l}E^{\mathrm{R}}(\,\overline{id_i}^{\,i}\,) \triangleright x\langle t\_args\rangle, (\,\overline{t_i}^{\,i}\,) \\ \overline{\langle E^{\mathrm{T}}, \langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle\rangle, t_i \vdash pat_i : u_i \triangleright pat'_i, E^{\mathrm{T}}{}_i, \Sigma^{\mathrm{N}}{}_i}^{\,i} \\ \mathbf{disjoint\ doms}\,(\,\overline{E^{\mathrm{T}}{}_i}^{\,i}\,) \\ \langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle \vdash x\langle t\_args\rangle \precapprox t, \Sigma^{\mathrm{N}}\end{array}}{\langle E^{\mathrm{T}}, \langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle\rangle, t \vdash \{\,\overline{id_i = pat_i}^{\,i}\,;^{?}\} : x\langle t\_args\rangle \triangleright \{\,\overline{id_i = pat'_i}^{\,i}\,;^{?}\}, \uplus\,\overline{E^{\mathrm{T}}{}_i}^{\,i}, \Sigma^{\mathrm{N}} \uplus\ \overline{\Sigma^{\mathrm{N}}{}_i}^{\,i}}\ \textsc{check\_pat\_record}$$

$$\frac{\begin{array}{l}\langle E^{\mathrm{T}}, E^{\mathrm{D}}\rangle, t \vdash pat_1 : u_1 \triangleright pat'_1, {E^{\mathrm{T}}}_1, {\Sigma^{\mathrm{N}}}_1 \quad ... \quad \langle E^{\mathrm{T}}, E^{\mathrm{D}}\rangle, t \vdash pat_n : u_n \triangleright pat'_n, {E^{\mathrm{T}}}_n, {\Sigma^{\mathrm{N}}}_n \\ \mathbf{disjoint\ doms}\,({E^{\mathrm{T}}}_1, \ ..., {E^{\mathrm{T}}}_n) \\ E^{\mathrm{D}} \vdash u_1 \precapprox t, {\Sigma^{\mathrm{N}'}}_1 \quad ... \quad E^{\mathrm{D}} \vdash u_n \precapprox t, {\Sigma^{\mathrm{N}'}}_n \\ ne_4 \equiv \mathbf{length}\,(pat_1\ ...\ pat_n) \\ \Sigma^{\mathrm{N}} \equiv {\Sigma^{\mathrm{N}}}_1 \uplus ... \uplus {\Sigma^{\mathrm{N}}}_n \\ \Sigma^{\mathrm{N}'} \equiv {\Sigma^{\mathrm{N}'}}_1 \uplus ... \uplus {\Sigma^{\mathrm{N}'}}_n\end{array}}{\langle E^{\mathrm{T}}, E^{\mathrm{D}}\rangle, \mathbf{vector}\,\langle ne_1\ ne_2\ order\ t\rangle \vdash [pat_1, ..., pat_n] : \mathbf{vector}\,\langle ne_3\ ne_4\ order\ u\rangle \triangleright [pat'_1, ..., pat'_n], ({E^{\mathrm{T}}}_1 \uplus ... \uplus {E^{\mathrm{T}}}_n), \Sigma^{\mathrm{N}} \uplus \Sigma^{\mathrm{N}'} \uplus \{ne_2 = ne_4\}}\ \text{CHECK\_PAT\_VECTOR}$$

$$\frac{\begin{array}{l}\langle E^{\mathrm{T}}, E^{\mathrm{D}}\rangle, t \vdash pat_1 : u_1 \triangleright pat'_1, {E^{\mathrm{T}}}_1, {\Sigma^{\mathrm{N}}}_1 \quad ... \quad \langle E^{\mathrm{T}}, E^{\mathrm{D}}\rangle, t \vdash pat_n : u_n \triangleright pat'_n, {E^{\mathrm{T}}}_n, {\Sigma^{\mathrm{N}}}_n \\ E^{\mathrm{D}} \vdash u_1 \precapprox t, {\Sigma^{\mathrm{N}'}}_1 \quad ... \quad E^{\mathrm{D}} \vdash u_n \precapprox t, {\Sigma^{\mathrm{N}'}}_n \\ ne_4 \equiv \mathbf{length}\,(pat_1\ ...\ pat_n) \\ \mathbf{disjoint\ doms}\,({E^{\mathrm{T}}}_1, \ ..., {E^{\mathrm{T}}}_n) \\ num_1 < ... < num_n \\ \Sigma^{\mathrm{N}} \equiv {\Sigma^{\mathrm{N}}}_1 \uplus ... \uplus {\Sigma^{\mathrm{N}}}_n \\ \Sigma^{\mathrm{N}'} \equiv {\Sigma^{\mathrm{N}'}}_1 \uplus ... \uplus {\Sigma^{\mathrm{N}'}}_n\end{array}}{\langle E^{\mathrm{T}}, E^{\mathrm{D}}\rangle, \mathbf{vector}\,\langle ne_1\ ne_2\ \mathbf{inc}\ t\rangle \vdash [num_1 = pat_1, ..., num_n = pat_n] : \mathbf{vector}\,\langle ne_3\ ne_4\ \mathbf{inc}\ t\rangle \triangleright [num_1 = pat'_1, ..., num_n = pat'_n], ({E^{\mathrm{T}}}_1 \uplus ... \uplus {E^{\mathrm{T}}}_n), \{ne_1 \leq num_1, ne_2 \geq ne_4\} \uplus$$

$$\frac{\begin{array}{l}\langle E^{\mathrm{T}}, E^{\mathrm{D}}\rangle, t \vdash pat_1 : u_1 \triangleright pat'_1, {E^{\mathrm{T}}}_1, {\Sigma^{\mathrm{N}}}_1 \quad ... \quad \langle E^{\mathrm{T}}, E^{\mathrm{D}}\rangle, t \vdash pat_n : u_n \triangleright pat'_n, {E^{\mathrm{T}}}_n, {\Sigma^{\mathrm{N}}}_n \\ E^{\mathrm{D}} \vdash u_1 \precapprox t, {\Sigma^{\mathrm{N}'}}_1 \quad ... \quad E^{\mathrm{D}} \vdash u_n \precapprox t, {\Sigma^{\mathrm{N}'}}_n \\ ne_4 \equiv \mathbf{length}\,(pat_1\ ...\ pat_n) \\ \mathbf{disjoint\ doms}\,({E^{\mathrm{T}}}_1, \ ..., {E^{\mathrm{T}}}_n) \\ num_1 > ... > num_n \\ \Sigma^{\mathrm{N}} \equiv {\Sigma^{\mathrm{N}}}_1 \uplus ... \uplus {\Sigma^{\mathrm{N}}}_n \\ \Sigma^{\mathrm{N}'} \equiv {\Sigma^{\mathrm{N}'}}_1 \uplus ... \uplus {\Sigma^{\mathrm{N}'}}_n\end{array}}{\langle E^{\mathrm{T}}, E^{\mathrm{D}}\rangle, \mathbf{vector}\,\langle ne_1\ ne_2\ \mathbf{dec}\ t\rangle \vdash [num_1 = pat_1, ..., num_n = pat_n] : \mathbf{vector}\,\langle ne_3\ ne_4\ \mathbf{dec}\ t\rangle \triangleright [num_1 = pat'_1, ..., num_n = pat'_n], ({E^{\mathrm{T}}}_1 \uplus ... \uplus {E^{\mathrm{T}}}_n), \{ne_1 \geq num_1, ne_2 \geq ne_4\} \uplus$$

$$\frac{\begin{array}{l}\langle E^{\mathrm{T}}, E^{\mathrm{D}}\rangle, \mathbf{vector}\,\langle ne''_1\ ne'''_1\ order\ t\rangle \vdash pat_1 : \mathbf{vector}\,\langle ne''_1\ ne'_1\ order\ u_1\rangle \triangleright pat'_1, {E^{\mathrm{T}}}_1, {\Sigma^{\mathrm{N}}}_1 \quad ... \quad \langle E^{\mathrm{T}}, E^{\mathrm{D}}\rangle, \mathbf{vector}\,\langle ne''_n\ ne'''_n\ order\ t\rangle \vdash pat_1 : \mathbf{vector}\,\langle ne''_n\ ne'_n\ order\ u_1\rangle \triangleright pat'_n, E \\ E^{\mathrm{D}} \vdash u_1 \precapprox t, {\Sigma^{\mathrm{N}'}}_1 \quad ... \quad E^{\mathrm{D}} \vdash u_n \precapprox t, {\Sigma^{\mathrm{N}'}}_n \\ \mathbf{disjoint\ doms}\,({E^{\mathrm{T}}}_1, \ ..., {E^{\mathrm{T}}}_n) \\ \Sigma^{\mathrm{N}} \equiv {\Sigma^{\mathrm{N}}}_1 \uplus ... \uplus {\Sigma^{\mathrm{N}}}_n \\ \Sigma^{\mathrm{N}'} \equiv {\Sigma^{\mathrm{N}'}}_1 \uplus ... \uplus {\Sigma^{\mathrm{N}'}}_n\end{array}}{\langle E^{\mathrm{T}}, E^{\mathrm{D}}\rangle, \mathbf{vector}\,\langle ne_1\ ne_2\ order\ t\rangle \vdash pat_1 : ... : pat_n : \mathbf{vector}\,\langle ne_1\ ne_4\ order\ t\rangle \triangleright pat'_1 : ... : pat'_n, ({E^{\mathrm{T}}}_1 \uplus ... \uplus {E^{\mathrm{T}}}_n), \{ne'_1 + ... + ne'_n \leq ne_2\} \uplus \Sigma^{\mathrm{N}} \uplus \Sigma^{\mathrm{N}'}$$

$$\frac{\begin{array}{l} E, t_1 \vdash pat_1 : u_1 \triangleright pat'_1, {E^{\mathrm{T}}}_1, {\Sigma^{\mathrm{N}}}_1 \quad .... \quad E, t_n \vdash pat_n : u_n \triangleright pat'_n, {E^{\mathrm{T}}}_n, {\Sigma^{\mathrm{N}}}_n \\ \mathbf{disjoint\ doms}\,({E^{\mathrm{T}}}_1, \ldots, {E^{\mathrm{T}}}_n) \end{array}}{E, (t_1, \ldots, t_n) \vdash (pat_1, \ldots, pat_n) : (u_1, \ldots, u_n) \triangleright (pat'_1, \ldots, pat'_n), ({E^{\mathrm{T}}}_1 \uplus \ldots \uplus {E^{\mathrm{T}}}_n), {\Sigma^{\mathrm{N}}}_1 \uplus \ldots \uplus {\Sigma^{\mathrm{N}}}_n} \text{ CHECK\_PAT\_TUP}$$

$$\frac{\begin{array}{l} \langle E^{\mathrm{T}}, E^{\mathrm{D}} \rangle, t \vdash pat_1 : u_1 \triangleright pat'_1, {E^{\mathrm{T}}}_1, {\Sigma^{\mathrm{N}}}_1 \quad .. \quad \langle E^{\mathrm{T}}, E^{\mathrm{D}} \rangle, t \vdash pat_n : u_n \triangleright pat'_n, {E^{\mathrm{T}}}_n, {\Sigma^{\mathrm{N}}}_n \\ \mathbf{disjoint\ doms}\,({E^{\mathrm{T}}}_1, .., {E^{\mathrm{T}}}_n) \\ E^{\mathrm{D}} \vdash u_1 \precapprox t, {\Sigma^{\mathrm{N}}}'_1 \quad .. \quad E^{\mathrm{D}} \vdash u_n \precapprox t, {\Sigma^{\mathrm{N}}}'_n \\ \mathbf{disjoint\ doms}\,({E^{\mathrm{T}}}_1, .., {E^{\mathrm{T}}}_n) \\ \Sigma^{\mathrm{N}} \equiv {\Sigma^{\mathrm{N}}}_1 \uplus .. \uplus {\Sigma^{\mathrm{N}}}_n \\ \Sigma^{\mathrm{N}\prime} \equiv {\Sigma^{\mathrm{N}}}'_1 \uplus .. \uplus {\Sigma^{\mathrm{N}}}'_n \end{array}}{\langle E^{\mathrm{T}}, E^{\mathrm{D}} \rangle, \mathbf{list}\,\langle t \rangle \vdash [||pat_1, .., pat_n||] : \mathbf{list}\,\langle t \rangle \triangleright [||pat'_1, .., pat'_n||], ({E^{\mathrm{T}}}_1 \uplus .. \uplus {E^{\mathrm{T}}}_n), \Sigma^{\mathrm{N}} \uplus \Sigma^{\mathrm{N}\prime}} \text{ CHECK\_PAT\_LIST}$$

$\boxed{E, t \vdash exp : t' \triangleright exp', I, E^{\mathrm{T}}}$ Typing expressions, collecting nexp constraints, effects, and new bindings

$$\frac{\begin{array}{l} E^{\mathrm{T}}(id) \triangleright \{tid_0 \mapsto kinf_0, .., tid_n \mapsto kinf_n\}, \{\,\}, \mathbf{Ctor}, \mathbf{unit} \rightarrow x\langle t\_args \rangle\,\mathbf{pure} \\ u \equiv x\langle t\_args \rangle[t\_arg_0/tid_0 .. t\_arg_n/tid_n] \\ E^{\mathrm{D}} \vdash u \precapprox t, \Sigma^{\mathrm{N}} \end{array}}{\langle E^{\mathrm{T}}, E^{\mathrm{D}} \rangle, t \vdash id : x \triangleright id, \langle \Sigma^{\mathrm{N}}, \mathbf{pure} \rangle, \{\,\}} \text{ CHECK\_EXP\_UNARYCTOR}$$

$$\frac{\begin{array}{l} E^{\mathrm{T}}(id) \triangleright \{\,\}, \{\,\}, tag, u \\ E^{\mathrm{D}}, t \vdash id : u \triangleright t', exp, \Sigma^{\mathrm{N}}, effect \end{array}}{\langle E^{\mathrm{T}}, E^{\mathrm{D}} \rangle, t \vdash id : u \triangleright id, \langle \Sigma^{\mathrm{N}}, effect \rangle, \{\,\}} \text{ CHECK\_EXP\_LOCALVAR}$$

$$\frac{\begin{array}{l} E^{\mathrm{T}}(id) \triangleright \{tid_1 \mapsto kinf_1, .., tid_n \mapsto kinf_n\}, \Sigma^{\mathrm{N}}, tag, u' \\ u \equiv u'[t\_arg_1/tid_1 .. t\_arg_n/tid_n] \\ E^{\mathrm{D}}, t \vdash id : u \triangleright t', exp, \Sigma^{\mathrm{N}\prime}, effect \end{array}}{\langle E^{\mathrm{T}}, E^{\mathrm{D}} \rangle, t \vdash id : u \triangleright id, \langle \Sigma^{\mathrm{N}} \uplus \Sigma^{\mathrm{N}\prime}, effect \rangle, \{\,\}} \text{ CHECK\_EXP\_OTHERVAR}$$

$$\frac{\begin{array}{l} E^{\mathrm{T}}(id) \triangleright \{tid_0 \mapsto kinf_0, .., tid_n \mapsto kinf_n\}, \{\,\}, \mathbf{Ctor}, t'' \rightarrow x\langle t\_args \rangle\,\mathbf{pure} \\ t' \rightarrow u\,\mathbf{pure} \equiv t'' \rightarrow x\langle t\_args \rangle\,\mathbf{pure}[t\_arg_0/tid_0 .. t\_arg_n/tid_n] \\ E^{\mathrm{D}} \vdash u \precapprox t, \Sigma^{\mathrm{N}} \\ \langle E^{\mathrm{T}}, E^{\mathrm{D}} \rangle, t' \vdash exp : u' \triangleright exp, \langle \Sigma^{\mathrm{N}\prime}, effect \rangle, E^{\mathrm{T}\prime} \end{array}}{\langle E^{\mathrm{T}}, E^{\mathrm{D}} \rangle, t \vdash id(exp) : t \triangleright id(exp'), \langle \Sigma^{\mathrm{N}} \uplus \Sigma^{\mathrm{N}}, effect \rangle, \{\,\}} \text{ CHECK\_EXP\_CTOR}$$

$$\frac{\begin{array}{l} E^{\mathrm{T}}(id) \triangleright \{tid_0 \mapsto kinf_0, .., tid_n \mapsto kinf_n\}, \Sigma^{\mathrm{N}}, tag, u \\ u[t\_arg_0/tid_0 .. t\_arg_n/tid_n] \equiv u_i \to u_j \; effect \\ u_i \equiv (\textbf{implicit}\ \langle ne \rangle, t_0, .., t_m) \\ \langle E^{\mathrm{T}}, E^{\mathrm{D}} \rangle, (t_0, .., t_m) \vdash (exp_0, .., exp_m) : u_i' \triangleright (exp_0', .., exp_m'), I, E^{\mathrm{T}\prime} \\ E^{\mathrm{D}}, t \vdash id(annot, exp_0', .., exp_m') : u_j \triangleright u_j', exp'', \Sigma^{\mathrm{N}\prime}, effect' \end{array}}{\langle E^{\mathrm{T}}, E^{\mathrm{D}} \rangle, t \vdash id(exp_0, .., exp_m) : u_j \triangleright exp'', I \uplus \langle \Sigma^{\mathrm{N}}, effect \rangle \uplus \langle \Sigma^{\mathrm{N}\prime}, effect' \rangle, E^{\mathrm{T}}}\ \textsc{check\_exp\_appImplicit}$$

$$\frac{\begin{array}{l} E^{\mathrm{T}}(id) \triangleright \{tid_0 \mapsto kinf_0, .., tid_n \mapsto kinf_n\}, \Sigma^{\mathrm{N}}, tag, u \\ u[t\_arg_0/tid_0 .. t\_arg_n/tid_n] \equiv u_i \to u_j \; effect \\ \langle E^{\mathrm{T}}, E^{\mathrm{D}} \rangle, u_i \vdash exp : u_i' \triangleright exp', I, E^{\mathrm{T}\prime} \\ E^{\mathrm{D}}, t \vdash id(exp') : u_j \triangleright u_j', exp'', \Sigma^{\mathrm{N}\prime}, effect' \end{array}}{\langle E^{\mathrm{T}}, E^{\mathrm{D}} \rangle, t \vdash id(exp) : u_j \triangleright exp'', I \uplus \langle \Sigma^{\mathrm{N}}, effect \rangle \uplus \langle \Sigma^{\mathrm{N}\prime}, effect' \rangle, E^{\mathrm{T}}}\ \textsc{check\_exp\_app}$$

$$\frac{\begin{array}{l} E^{\mathrm{T}}(id) \triangleright \textbf{overload}\,\{tid_0 \mapsto kinf_0, .., tid_n \mapsto kinf_n\}, \Sigma^{\mathrm{N}}, tag, u : tinf_1 \ldots tinf_n \\ u[t\_arg_0/tid_0 .. t\_arg_n/tid_n] \equiv u_i \to u_j \; effect \\ \langle E^{\mathrm{T}}, E^{\mathrm{D}} \rangle, u_i \vdash exp : u_i' \triangleright exp', I, E^{\mathrm{T}\prime} \\ \texttt{<<no parses (char 3): sel***ect (conformsto(ui', t)) of tinf1 ... tinfn gives tinf >>} \\ \langle (\{id \mapsto tinf\} \uplus E^{\mathrm{T}}), E^{\mathrm{D}} \rangle, t \vdash id(exp) : t' \triangleright exp'', I', E^{\mathrm{T}\prime\prime} \end{array}}{\langle E^{\mathrm{T}}, E^{\mathrm{D}} \rangle, t \vdash id(exp) : u_j \triangleright exp'', I \uplus I' \uplus \langle \Sigma^{\mathrm{N}}, effect \rangle \uplus \langle \Sigma^{\mathrm{N}\prime}, effect' \rangle, E^{\mathrm{T}}}\ \textsc{check\_exp\_appOverload}$$

$$\frac{\begin{array}{l} E^{\mathrm{T}}(id) \triangleright \{tid_0 \mapsto kinf_0, .., tid_n \mapsto kinf_n\}, \Sigma^{\mathrm{N}}, tag, u \\ u[t\_arg_0/tid_0 .. t\_arg_n/tid_n] \equiv u_i \to u_j \; effect \\ \langle E^{\mathrm{T}}, E^{\mathrm{D}} \rangle, u_i \vdash (exp_1, exp_2) : u_i' \triangleright (exp_1', exp_2'), I, E^{\mathrm{T}\prime} \\ E^{\mathrm{D}}, t \vdash exp_1'\ id\ exp_2' : u_j \triangleright u_j', exp, \Sigma^{\mathrm{N}\prime}, effect' \end{array}}{\langle E^{\mathrm{T}}, E^{\mathrm{D}} \rangle, t \vdash exp_1\ id\ exp_2 : t \triangleright exp, I \uplus \langle \Sigma^{\mathrm{N}}, effect \rangle \uplus \langle \Sigma^{\mathrm{N}\prime}, effect' \rangle, E^{\mathrm{T}}}\ \textsc{check\_exp\_infix\_app}$$

$$\frac{\begin{array}{l} E^{\mathrm{T}}(id) \triangleright \textbf{overload}\,\{tid_0 \mapsto kinf_0, .., tid_n \mapsto kinf_n\}, \Sigma^{\mathrm{N}}, tag, u : tinf_1 \ldots tinf_n \\ u[t\_arg_0/tid_0 .. t\_arg_n/tid_n] \equiv u_i \to u_j \; effect \\ \langle E^{\mathrm{T}}, E^{\mathrm{D}} \rangle, u_i \vdash (exp_1, exp_2) : u_i' \triangleright (exp_1', exp_2'), I, E^{\mathrm{T}\prime} \\ \texttt{<<no parses (char 3): sel***ect (conformsto(ui', t)) of tinf1 ... tinfn gives tinf >>} \\ \langle (\{id \mapsto tinf\} \uplus E^{\mathrm{T}}), E^{\mathrm{D}} \rangle, t \vdash exp_1\ id\ exp_2 : t' \triangleright exp, I', E^{\mathrm{T}\prime\prime} \end{array}}{\langle E^{\mathrm{T}}, E^{\mathrm{D}} \rangle, t \vdash exp_1\ id\ exp_2 : t \triangleright exp, I \uplus I \uplus \langle \Sigma^{\mathrm{N}}, effect \rangle \uplus \langle \Sigma^{\mathrm{N}\prime}, effect' \rangle, E^{\mathrm{T}}}\ \textsc{check\_exp\_infix\_appOverload}$$

$$\frac{\begin{array}{l} E^{\mathrm{R}}(\,\overline{id_i}^{\,i}\,) \triangleright x\langle t\_args\rangle, \overline{t_i}^{\,i} \\ \overline{\langle E^{\mathrm{T}}, \langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle\rangle, t_i \vdash exp_i : u_i \triangleright exp_i', \langle \Sigma^{\mathrm{N}}{}_i, effect_i\rangle, E^{\mathrm{T}}}^{\,i} \\ \overline{\langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle \vdash u_i \lessapprox t_i, \Sigma^{\mathrm{N}\prime}{}_i}^{\,i} \\ \Sigma^{\mathrm{N}} \equiv \uplus\, \overline{\Sigma^{\mathrm{N}}{}_i}^{\,i} \\ \Sigma^{\mathrm{N}\prime} \equiv \uplus\, \overline{\Sigma^{\mathrm{N}\prime}_i}^{\,i} \end{array}}{\langle E^{\mathrm{T}}, \langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle\rangle, t \vdash \{\overline{id_i = exp_i}^{\,i}\,;^{?}\} : x\langle t\_args\rangle \triangleright \{\overline{id_i = exp_i'}^{\,i}\,;^{?}\}, \uplus\langle \Sigma^{\mathrm{N}} \uplus \Sigma^{\mathrm{N}\prime}, \uplus\, \overline{effect_i}^{\,i}\rangle, \{\,\}} \text{ CHECK\_EXP\_RECORD}$$

$$\frac{\begin{array}{l} \langle E^{\mathrm{T}}, \langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle\rangle, t \vdash exp : x\langle t\_args\rangle \triangleright exp', I, E^{\mathrm{T}} \\ E^{\mathrm{R}}(x\langle t\_args\rangle) \triangleright \overline{id_n' : t_n'}^{\,n} \\ \overline{\langle E^{\mathrm{T}}, \langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle\rangle, t_i \vdash exp_i : u_i \triangleright exp_i', I_i, E^{\mathrm{T}}}^{\,i} \\ \overline{id_i : t_i}^{\,i} \subset \overline{id_n' : t_n'}^{\,n} \\ \overline{\langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle \vdash u_i \lessapprox t_i, \Sigma^{\mathrm{N}\prime}{}_i}^{\,i} \end{array}}{\langle E^{\mathrm{T}}, \langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle\rangle, t \vdash \{exp\ \mathbf{with}\ \overline{id_i = exp_i}^{\,i}\,;^{?}\} : x\langle t\_args\rangle \triangleright \{exp'\ \mathbf{with}\ \overline{id_i = exp_i'}^{\,i}\,\}, I \uplus \overline{I_i}^{\,i}, E^{\mathrm{T}}} \text{ CHECK\_EXP\_RECUP}$$

$$\frac{\begin{array}{l} \langle E^{\mathrm{T}}, E^{\mathrm{D}}\rangle, t \vdash exp_1 : u_1 \triangleright exp_1', I_1, E^{\mathrm{T}\prime} \quad ... \quad \langle E^{\mathrm{T}}, E^{\mathrm{D}}\rangle, t \vdash exp_n : u_n \triangleright exp_n', I_n, E^{\mathrm{T}\prime} \\ E^{\mathrm{D}} \vdash u_1 \lessapprox t, \Sigma^{\mathrm{N}}{}_1 \quad ... \quad E^{\mathrm{D}} \vdash u_n \lessapprox t, \Sigma^{\mathrm{N}}{}_n \\ \mathbf{length}\,(exp_1\ ...\ exp_n) \equiv ne \\ \Sigma^{\mathrm{N}} \equiv \{ne = ne_2\} \uplus \Sigma^{\mathrm{N}}{}_1 \uplus\ ...\ \uplus \Sigma^{\mathrm{N}}{}_n \end{array}}{E, \mathbf{vector}\,\langle ne_1\ ne_2\ order\ t\rangle \vdash [exp_1, \ ..., exp_n] : \mathbf{vector}\,\langle ne_1\ num\ order\ t\rangle \triangleright [exp_1', \ ..., exp_n'], \langle \Sigma^{\mathrm{N}}, \mathbf{pure}\rangle \uplus I_1 \uplus\ ...\ \uplus I_n, E^{\mathrm{T}}} \text{ CHECK\_EXP\_VECTOR}$$

$$\frac{\begin{array}{l} E, \mathbf{vector}\,\langle ne\ ne'\ order\ t\rangle \vdash exp_1 : \mathbf{vector}\,\langle ne_1\ ne_1'\ \mathbf{inc}\ u\rangle \triangleright exp_1', I_1, E^{\mathrm{T}} \\ E, \mathbf{range}\,\langle ne_2\ ne_2'\rangle \vdash exp_2 : \mathbf{range}\,\langle ne_3\ ne_3'\rangle \triangleright exp_2', I_2, E^{\mathrm{T}} \end{array}}{E, t \vdash exp_1[exp_2] : u \triangleright exp_1'[exp_2'], I_1 \uplus I_2 \uplus \langle\{ne_1 \leq ne_3, ne_3 + ne_3' \leq ne_1 + ne_1'\}, \mathbf{pure}\rangle, E^{\mathrm{T}}} \text{ CHECK\_EXP\_VECTORGETINC}$$

$$\frac{\begin{array}{l} E, \mathbf{vector}\,\langle ne\ ne'\ order\ t\rangle \vdash exp_1 : \mathbf{vector}\,\langle ne_1\ ne_1'\ \mathbf{dec}\ u\rangle \triangleright exp_1', I_1, E^{\mathrm{T}} \\ E, \mathbf{range}\,\langle ne_2\ ne_2'\rangle \vdash exp_2 : \mathbf{range}\,\langle ne_3\ ne_3'\rangle \triangleright exp_2', I_2, E^{\mathrm{T}} \end{array}}{E, t \vdash exp_1[exp_2] : u \triangleright exp_1'[exp_2'], I_1 \uplus I_2 \uplus \langle\{ne_1 \geq ne_3, ne_3 + (-ne_3') \leq ne_1 + (-ne_1')\}, \mathbf{pure}\rangle, E^{\mathrm{T}}} \text{ CHECK\_EXP\_VECTORGETDEC}$$

$$\frac{\begin{array}{l} E, \mathbf{vector}\,\langle ne_1\ ne_1'\ \mathbf{inc}\ t\rangle \vdash exp_1 : \mathbf{vector}\,\langle ne_2\ ne_2'\ \mathbf{inc}\ u\rangle \triangleright exp_1', I_1, E^{\mathrm{T}} \\ E, \mathbf{range}\,\langle ne_3\ ne_3'\rangle \vdash exp_2 : \mathbf{range}\,\langle ne_4\ ne_4'\rangle \triangleright exp_2', I_2, E^{\mathrm{T}} \\ E, \mathbf{range}\,\langle ne_5\ ne_5'\rangle \vdash exp_3 : \mathbf{range}\,\langle ne_6\ ne_6'\rangle \triangleright exp_3', I_3, E^{\mathrm{T}} \end{array}}{E, \mathbf{vector}\,\langle ne\ ne'\ \mathbf{inc}\ t\rangle \vdash exp_1[exp_2..exp_3] : \mathbf{vector}\,\langle ne_7\ ne_7'\ \mathbf{inc}\ u\rangle \triangleright exp_1'[exp_2' : exp_3'], I_1 \uplus I_2 \uplus I_3 \uplus \langle\{ne \geq ne_4, ne \leq ne_4', ne' \leq ne_4 + ne_6', ne_4 \leq ne_2, ne_4 + ne_6' \leq ne_2'\}, \mathbf{pure}$$

$$\frac{\begin{array}{l}E, \mathbf{vector}\,\langle ne_1\ ne_1'\ \mathbf{dec}\ t\rangle \vdash exp_1 : \mathbf{vector}\,\langle ne_2\ ne_2'\ \mathbf{dec}\ u\rangle \triangleright\ exp_1', I_1, E^{\mathrm{T}} \\ E, \mathbf{range}\,\langle ne_3\ ne_3'\rangle \vdash exp_2 : \mathbf{range}\,\langle ne_4\ ne_4'\rangle \triangleright\ exp_2', I_2, E^{\mathrm{T}} \\ E, \mathbf{range}\,\langle ne_5\ ne_5'\rangle \vdash exp_3 : \mathbf{range}\,\langle ne_6\ ne_6'\rangle \triangleright\ exp_3', I_3, E^{\mathrm{T}}\end{array}}{E, \mathbf{vector}\,\langle ne\ ne'\ \mathbf{dec}\ t\rangle \vdash exp_1[exp_2..exp_3] : \mathbf{vector}\,\langle ne_7\ ne_7'\ \mathbf{dec}\ u\rangle \triangleright\ exp_1'[exp_2' : exp_3'], I_1 \uplus I_2 \uplus I_3 \uplus \langle\{ne \le ne_4, ne \ge ne_4', ne' \le ne_6' + (-ne_4), ne_4' \ge ne_2, ne_6' + (-ne_4) \le ne$$

$$\frac{\begin{array}{l}E, \mathbf{vector}\,\langle ne\ ne'\ \mathbf{inc}\ t\rangle \vdash exp : \mathbf{vector}\,\langle ne_1\ ne_2\ \mathbf{inc}\ u\rangle \triangleright\ exp', I, E^{\mathrm{T}} \\ E, \mathbf{range}\,\langle ne_1'\ ne_2'\rangle \vdash exp_1 : \mathbf{range}\,\langle ne_3\ ne_4\rangle \triangleright\ exp_1', I_1, E^{\mathrm{T}} \\ E, t \vdash exp_2 : u \triangleright\ exp_2', I_2, E^{\mathrm{T}}\end{array}}{E, \mathbf{vector}\,\langle ne\ ne'\ \mathbf{inc}\ t\rangle \vdash [exp\ \mathbf{with}\ exp_1 = exp_2] : \mathbf{vector}\,\langle ne_1\ ne_2\ \mathbf{inc}\ u\rangle \triangleright\ [exp'\ \mathbf{with}\ exp_1' = exp_2'], I \uplus I_1 \uplus I_2 \uplus \langle\{ne_1 \le ne_3, ne_2 \ge ne_4\}, \mathbf{pure}\rangle, E^{\mathrm{T}}}\ \textsc{check\_exp\_vectoru}$$

$$\frac{\begin{array}{l}E, \mathbf{vector}\,\langle ne\ ne'\ \mathbf{dec}\ t\rangle \vdash exp : \mathbf{vector}\,\langle ne_1\ ne_2\ \mathbf{dec}\ u\rangle \triangleright\ exp', I, E^{\mathrm{T}} \\ E, \mathbf{range}\,\langle ne_1'\ ne_2'\rangle \vdash exp_1 : \mathbf{range}\,\langle ne_3\ ne_4\rangle \triangleright\ exp_1', I_1, E^{\mathrm{T}} \\ E, t \vdash exp_2 : u \triangleright\ exp_2', I_2, E^{\mathrm{T}}\end{array}}{E, \mathbf{vector}\,\langle ne\ ne'\ \mathbf{dec}\ t\rangle \vdash [exp\ \mathbf{with}\ exp_1 = exp_2] : \mathbf{vector}\,\langle ne_1\ ne_2\ \mathbf{dec}\ u\rangle \triangleright\ [exp'\ \mathbf{with}\ exp_1' = exp_2'], I \uplus I_1 \uplus I_2 \uplus \langle\{ne_1 \ge ne_3, ne_2 \ge ne_4\}, \mathbf{pure}\rangle, E^{\mathrm{T}}}\ \textsc{check\_exp\_vector}$$

$$\frac{\begin{array}{l}E, \mathbf{vector}\,\langle ne_1\ ne_2\ order\ t\rangle \vdash exp : \mathbf{vector}\,\langle ne_3\ ne_4\ \mathbf{inc}\ u\rangle \triangleright\ exp', I, E^{\mathrm{T}} \\ E, \mathbf{atom}\,\langle ne_5\rangle \vdash exp_1 : \mathbf{atom}\,\langle ne_6\rangle \triangleright\ exp_1', I_1, E^{\mathrm{T}} \\ E, \mathbf{atom}\,\langle ne_7\rangle \vdash exp_2 : \mathbf{atom}\,\langle ne_8\rangle \triangleright\ exp_2', I_2, E^{\mathrm{T}} \\ E, \mathbf{vector}\,\langle ne_9\ ne_{10}\ \mathbf{inc}\ t\rangle \vdash exp_3 : \mathbf{vector}\,\langle ne_{11}\ ne_{12}\ \mathbf{inc}\ u\rangle \triangleright\ exp_3', I_3, E^{\mathrm{T}} \\ I_4 \equiv \langle\{ne_3 \le ne_5, ne_3 + ne_4 \le ne_7, ne_{12} = ne_8 + (-ne_6), ne_6 + \mathbf{one} \le ne_8\}, \mathbf{pure}\rangle\end{array}}{E, \mathbf{vector}\,\langle ne_1\ ne_2\ order\ t\rangle \vdash [exp\ \mathbf{with}\ exp_1 : exp_2 = exp_3] : \mathbf{vector}\,\langle ne_3\ ne_4\ \mathbf{inc}\ u\rangle \triangleright\ [exp'\ \mathbf{with}\ exp_1' : exp_2' = exp_3'], I \uplus I_1 \uplus I_2 \uplus I_3 \uplus I_4, E^{\mathrm{T}}}\ \textsc{check\_exp\_vecrangeupInc}$$

$$\frac{\begin{array}{l}E, \mathbf{vector}\,\langle ne_1\ ne_2\ order\ t\rangle \vdash exp : \mathbf{vector}\,\langle ne_3\ ne_4\ \mathbf{inc}\ u\rangle \triangleright\ exp', I, E^{\mathrm{T}} \\ E, \mathbf{atom}\,\langle ne_5\rangle \vdash exp_1 : \mathbf{atom}\,\langle ne_6\rangle \triangleright\ exp_1', I_1, E^{\mathrm{T}} \\ E, \mathbf{atom}\,\langle ne_7\rangle \vdash exp_2 : \mathbf{atom}\,\langle ne_8\rangle \triangleright\ exp_2', I_2, E^{\mathrm{T}} \\ E, u \vdash exp_3 : u' \triangleright\ exp_3', I_3, E^{\mathrm{T}} \\ I_4 \equiv \langle\{ne_3 \le ne_5, ne_3 + ne_4 \le ne_7\}, \mathbf{pure}\rangle\end{array}}{E, \mathbf{vector}\,\langle ne_1\ ne_2\ order\ t\rangle \vdash [exp\ \mathbf{with}\ exp_1 : exp_2 = exp_3] : \mathbf{vector}\,\langle ne_3\ ne_4\ \mathbf{inc}\ u\rangle \triangleright\ [exp'\ \mathbf{with}\ exp_1' : exp_2' = exp_3'], I \uplus I_1 \uplus I_2 \uplus I_3 \uplus I_4, E^{\mathrm{T}}}\ \textsc{check\_exp\_vecrangeupvalu}$$

$$\frac{\begin{array}{l}E, \mathbf{vector}\,\langle ne_1\ ne_2\ order\ t\rangle \vdash exp : \mathbf{vector}\,\langle ne_3\ ne_4\ \mathbf{dec}\ u\rangle \triangleright\ exp', I, E^{\mathrm{T}} \\ E, \mathbf{atom}\,\langle ne_5\rangle \vdash exp_1 : \mathbf{atom}\,\langle ne_6\rangle \triangleright\ exp_1', I_1, E^{\mathrm{T}} \\ E, \mathbf{atom}\,\langle ne_7\rangle \vdash exp_2 : \mathbf{atom}\,\langle ne_8\rangle \triangleright\ exp_2', I_2, E^{\mathrm{T}} \\ E, \mathbf{vector}\,\langle ne_9\ ne_{10}\ \mathbf{dec}\ t\rangle \vdash exp_3 : \mathbf{vector}\,\langle ne_{11}\ ne_{12}\ \mathbf{dec}\ u\rangle \triangleright\ exp_3', I_3, E^{\mathrm{T}} \\ I_4 \equiv \langle\{ne_5 \le ne_3, ne_3 + (-ne_4) \le ne_6 + (-ne_8), ne_8 + \mathbf{one} \le ne_6\}, \mathbf{pure}\rangle\end{array}}{E, \mathbf{vector}\,\langle ne_1\ ne_2\ order\ t\rangle \vdash [exp\ \mathbf{with}\ exp_1 : exp_2 = exp_3] : \mathbf{vector}\,\langle ne_3\ ne_4\ \mathbf{dec}\ u\rangle \triangleright\ [exp'\ \mathbf{with}\ exp_1' : exp_2' = exp_3'], I \uplus I_1 \uplus I_2 \uplus I_3 \uplus I_4, E^{\mathrm{T}}}\ \textsc{check\_exp\_vecrangeupDec}$$

$$\frac{\begin{array}{l} E, \mathbf{vector}\,\langle ne_1\ ne_2\ order\ t\rangle \vdash exp : \mathbf{vector}\,\langle ne_3\ ne_4\ \mathbf{dec}\ u\rangle \triangleright exp', I, E^{\mathrm{T}} \\ E, \mathbf{atom}\,\langle ne_5\rangle \vdash exp_1 : \mathbf{atom}\,\langle ne_6\rangle \triangleright exp_1', I_1, E^{\mathrm{T}} \\ E, \mathbf{atom}\,\langle ne_7\rangle \vdash exp_2 : \mathbf{atom}\,\langle ne_8\rangle \triangleright exp_2', I_2, E^{\mathrm{T}} \\ E, u \vdash exp_3 : u' \triangleright exp_3', I_3, E^{\mathrm{T}} \\ I_4 \equiv \langle \{ne_5 \leq ne_3, ne_3 + (-ne_4) \leq ne_6 + (-ne_8), ne_8 + \mathbf{one} \leq ne_6\}, \mathbf{pure}\rangle \end{array}}{E, \mathbf{vector}\,\langle ne_1\ ne_2\ order\ t\rangle \vdash [exp\ \mathbf{with}\ exp_1 : exp_2 = exp_3] : \mathbf{vector}\,\langle ne_3\ ne_4\ \mathbf{dec}\ u\rangle \triangleright [exp'\ \mathbf{with}\ exp_1' : exp_2' = exp_3'], I \uplus I_1 \uplus I_2 \uplus I_3 \uplus I_4, E^{\mathrm{T}}}\ \text{CHECK\_EXP\_VECRANGEUPVAL}$$

$$\frac{\begin{array}{l} E^{\mathrm{R}}(x\langle t\_args\rangle) \triangleright \overline{id_i : t_i}^{\,i}\ id : u\ \overline{id_j' : t_j'}^{\,j} \\ \langle E^{\mathrm{T}}, \langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle\rangle, t'' \vdash exp : x\langle t\_args\rangle \triangleright exp', I, E^{\mathrm{T}} \\ E^{\mathrm{D}}, t \vdash exp'.id : u \triangleright t', exp_1', \Sigma^{\mathrm{N}\prime}, effect \end{array}}{\langle E^{\mathrm{T}}, \langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle\rangle, t \vdash exp.id : u \triangleright exp_1', I \uplus \langle \Sigma^{\mathrm{N}\prime}, effect\rangle, E^{\mathrm{T}}}\ \text{CHECK\_EXP\_FIELD}$$

$$\frac{\begin{array}{l} \langle E^{\mathrm{T}}, E^{\mathrm{D}}\rangle, t'' \vdash exp : u \triangleright exp', I, E^{\mathrm{T}} \\ \overline{\langle E^{\mathrm{T}}, E^{\mathrm{D}}\rangle, u \vdash pat_i : u_i' \triangleright pat_i', {E^{\mathrm{T}}}_i, {\Sigma^{\mathrm{N}}}_i}^{\,i} \\ \overline{\langle (E^{\mathrm{T}} \uplus {E^{\mathrm{T}}}_i), E^{\mathrm{D}}\rangle, t \vdash exp_i : u_i'' \triangleright exp_i', I_i, {E^{\mathrm{T}}}_i'}^{\,i} \end{array}}{\langle E^{\mathrm{T}}, E^{\mathrm{D}}\rangle, t \vdash \mathbf{switch}\ exp\{\ \overline{\mathbf{case}\ pat_i \rightarrow exp_i}^{\,i}\ \} : u \triangleright \mathbf{switch}\ exp'\{\ \overline{\mathbf{case}\ pat_i' \rightarrow exp_i'}^{\,i}\ \}, I \uplus\ \overline{I_i \uplus \langle {\Sigma^{\mathrm{N}}}_i, \mathbf{pure}\rangle}^{\,i}, E^{\mathrm{T}}}\ \text{CHECK\_EXP\_CASE}$$

$$\frac{\begin{array}{l} \langle E^{\mathrm{T}}, E^{\mathrm{D}}\rangle, t'' \vdash exp : u \triangleright exp', I, E^{\mathrm{T}} \\ E^{\mathrm{D}} \vdash typ \rightsquigarrow t' \\ E^{\mathrm{D}}, t' \vdash exp' : u \triangleright u', exp'', \Sigma^{\mathrm{N}}, effect \\ E^{\mathrm{D}}, t \vdash exp'' : t' \triangleright u'', exp''', \Sigma^{\mathrm{N}\prime}, effect' \end{array}}{\langle E^{\mathrm{T}}, E^{\mathrm{D}}\rangle, t \vdash (typ)exp : t \triangleright exp''', I \uplus \langle \Sigma^{\mathrm{N}} \uplus \Sigma^{\mathrm{N}\prime}, effect \uplus effect'\rangle, E^{\mathrm{T}}}\ \text{CHECK\_EXP\_TYPED}$$

$$\frac{\begin{array}{l} \langle E^{\mathrm{T}}, E^{\mathrm{D}}\rangle \vdash letbind \triangleright letbind', {E^{\mathrm{T}}}_1, \Sigma^{\mathrm{N}}, effect, \{\,\} \\ \langle (E^{\mathrm{T}} \uplus {E^{\mathrm{T}}}_1), E^{\mathrm{D}}\rangle, t \vdash exp : u \triangleright exp', I_2, {E^{\mathrm{T}}}_2 \end{array}}{\langle E^{\mathrm{T}}, E^{\mathrm{D}}\rangle, t \vdash letbind\ \mathbf{in}\ exp : t \triangleright letbind'\ \mathbf{in}\ exp', \langle \Sigma^{\mathrm{N}}, effect\rangle \uplus I_2, E^{\mathrm{T}}}\ \text{CHECK\_EXP\_LET}$$

$$\frac{E, t_1 \vdash exp_1 : u_1 \triangleright exp_1', I_1, {E^{\mathrm{T}}}_1 \quad .... \quad E, t_n \vdash exp_n : u_n \triangleright exp_n', I_n, {E^{\mathrm{T}}}_n}{E, (t_1, \ldots., t_n) \vdash (exp_1, \ldots., exp_n) : (u_1, \ldots., u_n) \triangleright (exp_1', \ldots., exp_n'), I_1 \uplus \ldots. \uplus I_n, E^{\mathrm{T}}}\ \text{CHECK\_EXP\_TUP}$$

$$\frac{\begin{array}{l} \langle E^{\mathrm{T}}, E^{\mathrm{D}}\rangle, t \vdash exp_1 : u_1 \triangleright exp_1', I_1, {E^{\mathrm{T}}}_1 \quad .. \quad \langle E^{\mathrm{T}}, E^{\mathrm{D}}\rangle, t \vdash exp_n : u_n \triangleright exp_n', I_n, {E^{\mathrm{T}}}_n \\ E^{\mathrm{D}} \vdash u_1 \lessapprox t, {\Sigma^{\mathrm{N}}}_1 \quad .. \quad E^{\mathrm{D}} \vdash u_n \lessapprox t, {\Sigma^{\mathrm{N}}}_n \end{array}}{\langle E^{\mathrm{T}}, E^{\mathrm{D}}\rangle, \mathbf{list}\,\langle t\rangle \vdash [||exp_1, .., exp_n||] : \mathbf{list}\,\langle u\rangle \triangleright [||exp_1', .., exp_n'||], \langle {\Sigma^{\mathrm{N}}}_1 \uplus .. \uplus {\Sigma^{\mathrm{N}}}_n, \mathbf{pure}\rangle \uplus I_1 \uplus .. \uplus I_n, E^{\mathrm{T}}}\ \text{CHECK\_EXP\_LIST}$$

$$\frac{\begin{array}{c} E, \mathbf{bit} \vdash exp_1 : \mathbf{bit} \triangleright exp_1', I_1, {E^{\mathrm{T}}}' \\ E, t \vdash exp_2 : u_1 \triangleright exp_2', I_2, {E^{\mathrm{T}}}_2 \\ E, t \vdash exp_3 : u_2 \triangleright exp_3', I_3, {E^{\mathrm{T}}}_3 \\ E^{\mathrm{D}} \vdash u_1 \precapprox t, {\Sigma^{\mathrm{N}}}_1 \\ E^{\mathrm{D}} \vdash u_2 \precapprox t, {\Sigma^{\mathrm{N}}}_2 \end{array}}{\langle E^{\mathrm{T}}, E^{\mathrm{D}} \rangle, t \vdash \mathbf{if}\ exp_1\ \mathbf{then}\ exp_2\ \mathbf{else}\ exp_3 : u \triangleright \mathbf{if}\ exp_1'\ \mathbf{then}\ exp_2'\ \mathbf{else}\ exp_3', \langle {\Sigma^{\mathrm{N}}}_1 \uplus {\Sigma^{\mathrm{N}}}_2, \mathbf{pure} \rangle \uplus I_1 \uplus I_2 \uplus I_3, ({E^{\mathrm{T}}}_2 \cap {E^{\mathrm{T}}}_3)} \ \text{CHECK\_EXP\_IF}$$

$$\frac{\begin{array}{c} \langle E^{\mathrm{T}}, E^{\mathrm{D}} \rangle, \mathbf{range}\, \langle ne_1\ ne_2 \rangle \vdash exp_1 : \mathbf{range}\, \langle ne_7\ ne_8 \rangle \triangleright exp_1', I_1, E^{\mathrm{T}} \\ \langle E^{\mathrm{T}}, E^{\mathrm{D}} \rangle, \mathbf{range}\, \langle ne_3\ ne_4 \rangle \vdash exp_2 : \mathbf{range}\, \langle ne_9\ ne_{10} \rangle \triangleright exp_2', I_2, E^{\mathrm{T}} \\ \langle E^{\mathrm{T}}, E^{\mathrm{D}} \rangle, \mathbf{range}\, \langle ne_5\ ne_6 \rangle \vdash exp_3 : \mathbf{range}\, \langle ne_{11}\ ne_{12} \rangle \triangleright exp_3', I_3, E^{\mathrm{T}} \\ \langle (E^{\mathrm{T}} \uplus \{ id \mapsto \mathbf{range}\, \langle ne_1\ ne_4 \rangle \}), E^{\mathrm{D}} \rangle, \mathbf{unit} \vdash exp_4 : t \triangleright exp_4', I_4, {E^{\mathrm{T}}}' \end{array}}{\langle E^{\mathrm{T}}, E^{\mathrm{D}} \rangle, \mathbf{unit} \vdash \mathbf{foreach}\, (id\ \mathbf{from}\ exp_1\ \mathbf{to}\ exp_2\ \mathbf{by}\ exp_3) exp_4 : t \triangleright \mathbf{foreach}\, (id\ \mathbf{from}\ exp_1'\ \mathbf{to}\ exp_2'\ \mathbf{by}\ exp_3') exp_4', I_1 \uplus I_2 \uplus I_3 \uplus I_4 \uplus \langle \{ ne_1 \leq ne_3 + ne_4 \}, \mathbf{pure} \rangle, E^{\mathrm{T}}} \ \text{CHECK\_EXI}$$

$$\frac{\begin{array}{c} E, t \vdash exp_1 : u \triangleright exp_1', I_1, E^{\mathrm{T}} \\ E, \mathbf{list}\, \langle t \rangle \vdash exp_2 : \mathbf{list}\, \langle u \rangle \triangleright exp_2', I_2, E^{\mathrm{T}} \end{array}}{E, \mathbf{list}\, \langle t \rangle \vdash exp_1 :: exp_2 : \mathbf{list}\, \langle u \rangle \triangleright exp_1' :: exp_2', I_1 \uplus I_2, E^{\mathrm{T}}} \ \text{CHECK\_EXP\_CONS}$$

$$\frac{t \vdash lit : u \Rightarrow exp, \Sigma^{\mathrm{N}}}{E, t \vdash lit : u \triangleright exp, \langle \Sigma^{\mathrm{N}}, \mathbf{pure} \rangle, E^{\mathrm{T}}} \ \text{CHECK\_EXP\_LIT}$$

$$\frac{\langle E^{\mathrm{T}}, E^{\mathrm{D}} \rangle, \mathbf{unit} \vdash exp : \mathbf{unit} \triangleright exp', I, {E^{\mathrm{T}}}_1}{\langle E^{\mathrm{T}}, E^{\mathrm{D}} \rangle, \mathbf{unit} \vdash \{ exp \} : \mathbf{unit} \triangleright \{ exp' \}, I, E^{\mathrm{T}}} \ \text{CHECK\_EXP\_BLOCKBASE}$$

$$\frac{\begin{array}{c} \langle E^{\mathrm{T}}, E^{\mathrm{D}} \rangle, \mathbf{unit} \vdash exp : \mathbf{unit} \triangleright exp', I_1, {E^{\mathrm{T}}}_1 \\ \langle (E^{\mathrm{T}} \uplus {E^{\mathrm{T}}}_1), E^{\mathrm{D}} \rangle, \mathbf{unit} \vdash \{\ \overline{exp_i}^{\,i}\ \} : \mathbf{unit} \triangleright \{\ \overline{exp_i'}^{\,i}\ \}, I_2, {E^{\mathrm{T}}}_2 \end{array}}{\langle E^{\mathrm{T}}, E^{\mathrm{D}} \rangle, \mathbf{unit} \vdash \{ exp;\ \overline{exp_i}^{\,i}\ \} : \mathbf{unit} \triangleright \{ exp';\ \overline{exp_i'}^{\,i}\ \}, I_1 \uplus I_2, E^{\mathrm{T}}} \ \text{CHECK\_EXP\_BLOCKREC}$$

$$\frac{\langle E^{\mathrm{T}}, E^{\mathrm{D}} \rangle, \mathbf{unit} \vdash exp : \mathbf{unit} \triangleright exp', I, {E^{\mathrm{T}}}_1}{\langle E^{\mathrm{T}}, E^{\mathrm{D}} \rangle, \mathbf{unit} \vdash \mathbf{nondet}\, \{ exp \} : \mathbf{unit} \triangleright \{ exp' \}, I, E^{\mathrm{T}}} \ \text{CHECK\_EXP\_NONDETBASE}$$

$$\frac{\begin{array}{c} \langle E^{\mathrm{T}}, E^{\mathrm{D}} \rangle, \mathbf{unit} \vdash exp : \mathbf{unit} \triangleright exp', I_1, {E^{\mathrm{T}}}_1 \\ \langle (E^{\mathrm{T}} \uplus {E^{\mathrm{T}}}_1), E^{\mathrm{D}} \rangle, \mathbf{unit} \vdash \mathbf{nondet}\, \{\ \overline{exp_i}^{\,i}\ \} : \mathbf{unit} \triangleright \{\ \overline{exp_i'}^{\,i}\ \}, I_2, {E^{\mathrm{T}}}_2 \end{array}}{\langle E^{\mathrm{T}}, E^{\mathrm{D}} \rangle, \mathbf{unit} \vdash \mathbf{nondet}\, \{ exp;\ \overline{exp_i}^{\,i}\ \} : \mathbf{unit} \triangleright \{ exp';\ \overline{exp_i'}^{\,i}\ \}, I_1 \uplus I_2, E^{\mathrm{T}}} \ \text{CHECK\_EXP\_NONDETREC}$$

$$\frac{\begin{array}{c} E, t \vdash \mathit{exp} : u \triangleright \mathit{exp}', I_1, {E^{\mathrm{T}}}_1 \\ E \vdash \mathit{lexp} : t \triangleright \mathit{lexp}', I_2, {E^{\mathrm{T}}}_2 \end{array}}{E, \mathbf{unit} \vdash \mathit{lexp} := \mathit{exp} : \mathbf{unit} \triangleright \mathit{lexp}' := \mathit{exp}', I \uplus I_2, {E^{\mathrm{T}}}_2} \quad \textsc{check\_exp\_assign}$$

$\boxed{E \vdash \mathit{lexp} : t \triangleright \mathit{lexp}', I, E^{\mathrm{T}}}$ Check the left hand side of an assignment

$$\frac{E^{\mathrm{T}}(id) \triangleright \mathbf{register}\,\langle t \rangle}{\langle E^{\mathrm{T}}, E^{\mathrm{D}} \rangle \vdash \mathit{id} : t \triangleright \mathit{id}, \langle \{\,\}, \{\mathbf{wreg}\} \rangle, E^{\mathrm{T}}} \quad \textsc{check\_lexp\_wreg}$$

$$\frac{E^{\mathrm{T}}(id) \triangleright \mathbf{reg}\,\langle t \rangle}{\langle E^{\mathrm{T}}, E^{\mathrm{D}} \rangle \vdash \mathit{id} : t \triangleright \mathit{id}, I_\epsilon, E^{\mathrm{T}}} \quad \textsc{check\_lexp\_wlocl}$$

$$\frac{E^{\mathrm{T}}(id) \triangleright t}{\langle E^{\mathrm{T}}, E^{\mathrm{D}} \rangle \vdash \mathit{id} : t \triangleright \mathit{id}, I_\epsilon, E^{\mathrm{T}}} \quad \textsc{check\_lexp\_var}$$

$$\frac{\mathit{id} \notin \mathbf{dom}\,(E^{\mathrm{T}})}{\langle E^{\mathrm{T}}, E^{\mathrm{D}} \rangle \vdash \mathit{id} : t \triangleright \mathit{id}, I_\epsilon, \{\mathit{id} \mapsto \mathbf{reg}\,\langle t \rangle\}} \quad \textsc{check\_lexp\_wnew}$$

$$\frac{\begin{array}{c} E^{\mathrm{T}}(id) \triangleright \mathbf{register}\,\langle t \rangle \\ E^{\mathrm{D}} \vdash \mathit{typ} \rightsquigarrow u \\ E^{\mathrm{D}} \vdash u \lessapprox t, \Sigma^{\mathrm{N}} \end{array}}{\langle E^{\mathrm{T}}, E^{\mathrm{D}} \rangle \vdash (\mathit{typ})\mathit{id} : t \triangleright \mathit{id}, \langle \Sigma^{\mathrm{N}}, \{\mathbf{wreg}\} \rangle, E^{\mathrm{T}}} \quad \textsc{check\_lexp\_wregCast}$$

$$\frac{\begin{array}{c} E^{\mathrm{T}}(id) \triangleright \mathbf{reg}\,\langle t \rangle \\ E^{\mathrm{D}} \vdash \mathit{typ} \rightsquigarrow u \\ E^{\mathrm{D}} \vdash u \lessapprox t, \Sigma^{\mathrm{N}} \end{array}}{\langle E^{\mathrm{T}}, E^{\mathrm{D}} \rangle \vdash (\mathit{typ})\mathit{id} : t \triangleright \mathit{id}, \langle \Sigma^{\mathrm{N}}, \mathbf{pure} \rangle, E^{\mathrm{T}}} \quad \textsc{check\_lexp\_wloclCast}$$

$$\frac{\begin{array}{c} E^{\mathrm{T}}(id) \triangleright t \\ E^{\mathrm{D}} \vdash \mathit{typ} \rightsquigarrow u \\ E^{\mathrm{D}} \vdash u \lessapprox t, \Sigma^{\mathrm{N}} \end{array}}{\langle E^{\mathrm{T}}, E^{\mathrm{D}} \rangle \vdash (\mathit{typ})\mathit{id} : t \triangleright \mathit{id}, \langle \Sigma^{\mathrm{N}}, \mathbf{pure} \rangle, E^{\mathrm{T}}} \quad \textsc{check\_lexp\_varCast}$$

$$\frac{\begin{array}{c} \mathit{id} \notin \mathbf{dom}\,(E^{\mathrm{T}}) \\ E^{\mathrm{D}} \vdash \mathit{typ} \rightsquigarrow t \end{array}}{\langle E^{\mathrm{T}}, E^{\mathrm{D}} \rangle \vdash (\mathit{typ})\mathit{id} : t \triangleright \mathit{id}, I_\epsilon, \{\mathit{id} \mapsto \mathbf{reg}\,\langle t \rangle\}} \quad \textsc{check\_lexp\_wnewCast}$$

$$\frac{\begin{array}{l} E^{\mathrm{T}}(id) \triangleright E^{\mathrm{K}}, \Sigma^{\mathrm{N}}, \mathbf{Extern}\,, t_1 \to t\,\{\,\overline{base\_effect_i}^{\,i}\,, \mathbf{wmem},\ \overline{base\_effect'_j}^{\,j}\,\} \\ \langle E^{\mathrm{T}}, E^{\mathrm{D}}\rangle, t_1 \vdash exp : u_1 \triangleright exp', I, {E^{\mathrm{T}}}_1 \end{array}}{\langle E^{\mathrm{T}}, E^{\mathrm{D}}\rangle \vdash id(exp) : t \triangleright id(exp'), I \uplus \langle \Sigma^{\mathrm{N}}, \{\mathbf{wmem}\}\rangle, E^{\mathrm{T}}} \text{ CHECK\_LEXP\_WMEM}$$

$$\frac{\begin{array}{c} E, \mathbf{atom}\,\langle ne\rangle \vdash exp : u \triangleright exp', I_1, E^{\mathrm{T}} \\ E \vdash lexp : \mathbf{vector}\,\langle ne_1\ ne_2\,\mathbf{inc}\ t\rangle \triangleright lexp', I_2, E^{\mathrm{T}} \end{array}}{E \vdash lexp[exp] : t \triangleright lexp'[exp'], I_1 \uplus I_2 \uplus \langle \{ne_1 \le ne, ne_1 + ne_2 \ge ne\}, \mathbf{pure}\rangle, E^{\mathrm{T}}} \text{ CHECK\_LEXP\_WBITINC}$$

$$\frac{\begin{array}{c} E, \mathbf{atom}\,\langle ne\rangle \vdash exp : u \triangleright exp', I_1, E^{\mathrm{T}} \\ E \vdash lexp : \mathbf{vector}\,\langle ne_1\ ne_2\,\mathbf{dec}\ t\rangle \triangleright lexp', I_2, E^{\mathrm{T}} \end{array}}{E \vdash lexp[exp] : t \triangleright lexp'[exp'], I_1 \uplus I_2 \uplus \langle \{ne \le ne_1, ne_1 + (-ne_2) \le ne\}, \mathbf{pure}\rangle, E^{\mathrm{T}}} \text{ CHECK\_LEXP\_WBITDEC}$$

$$\frac{\begin{array}{c} E, \mathbf{atom}\,\langle ne_1\rangle \vdash exp_1 : u_1 \triangleright exp'_1, I_1, E^{\mathrm{T}} \\ E, \mathbf{atom}\,\langle ne_2\rangle \vdash exp_2 : u_2 \triangleright exp'_2, I_2, E^{\mathrm{T}} \\ E \vdash lexp : \mathbf{vector}\,\langle ne_3\ ne_4\,\mathbf{inc}\ t\rangle \triangleright lexp', I_3, E^{\mathrm{T}} \end{array}}{E \vdash lexp[exp_1 : exp_2] : \mathbf{vector}\,\langle ne_1\ ne_2 + (-ne_1)\,\mathbf{inc}\ t\rangle \triangleright lexp'[exp'_1 : exp'_2], I_1 \uplus I_2 \uplus I_3 \uplus \langle \{ne_3 \le ne_1, ne_3 + ne_4 \le ne_2 + (-ne_1)\}, \mathbf{pure}\rangle, E^{\mathrm{T}}} \text{ CHECK\_LEXP\_WSLICEINC}$$

$$\frac{\begin{array}{c} E, \mathbf{atom}\,\langle ne_1\rangle \vdash exp_1 : u_1 \triangleright exp'_1, I_1, E^{\mathrm{T}} \\ E, \mathbf{atom}\,\langle ne_2\rangle \vdash exp_2 : u_2 \triangleright exp'_2, I_2, E^{\mathrm{T}} \\ E \vdash lexp : \mathbf{vector}\,\langle ne_3\ ne_4\,\mathbf{inc}\ t\rangle \triangleright lexp', I_3, E^{\mathrm{T}} \end{array}}{E \vdash lexp[exp_1 : exp_2] : \mathbf{vector}\,\langle ne_1\ ne_2 + (-ne_1)\,\mathbf{inc}\ t\rangle \triangleright lexp'[exp'_1 : exp'_2], I_1 \uplus I_2 \uplus I_3 \uplus \langle \{ne_1 \le ne_3, ne_3 + (-ne_4) \le ne_1 + (-ne_2)\}, \mathbf{pure}\rangle, E^{\mathrm{T}}} \text{ CHECK\_LEXP\_WSLICEDEC}$$

$$\frac{\begin{array}{l} E^{\mathrm{R}}(x\langle t\_args\rangle) \triangleright \overline{id_i : t_i}^{\,i}\ id : t\ \overline{id'_j : t'_j}^{\,j} \\ \langle E^{\mathrm{T}}, \langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle\rangle \vdash lexp : x\langle t\_args\rangle \triangleright lexp', I, E^{\mathrm{T}} \end{array}}{\langle E^{\mathrm{T}}, \langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle\rangle \vdash lexp.id : t \triangleright lexp'.id, I, E^{\mathrm{T}}} \text{ CHECK\_LEXP\_WRECORD}$$

$\boxed{E \vdash letbind \triangleright letbind', E^{\mathrm{T}}, \Sigma^{\mathrm{N}}, effect, E^{\mathrm{K}}}$ Build the environment for a let binding, collecting index constraints

$$\frac{\begin{array}{l} \langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle \vdash typschm \rightsquigarrow t, {E^{\mathrm{K}}}_2, \Sigma^{\mathrm{N}} \\ \langle E^{\mathrm{T}}, \langle E^{\mathrm{K}} \uplus {E^{\mathrm{K}}}_2, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle\rangle, t \vdash pat : u \triangleright pat', {E^{\mathrm{T}}}_1, {\Sigma^{\mathrm{N}}}_1 \\ \langle E^{\mathrm{T}}, \langle E^{\mathrm{K}} \uplus {E^{\mathrm{K}}}_2, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle\rangle, t \vdash exp : u' \triangleright exp', \langle {\Sigma^{\mathrm{N}}}_2, effect\rangle, {E^{\mathrm{T}}}_2 \\ \langle E^{\mathrm{K}} \uplus {E^{\mathrm{K}}}_2, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle \vdash u' \precsim\!\!\approx u, {\Sigma^{\mathrm{N}}}_3 \end{array}}{\langle E^{\mathrm{T}}, \langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle\rangle \vdash \mathbf{let}\ typschm\ pat = exp \triangleright \mathbf{let}\ typschm\ pat' = exp', {E^{\mathrm{T}}}_1, \Sigma^{\mathrm{N}} \uplus {\Sigma^{\mathrm{N}}}_1 \uplus {\Sigma^{\mathrm{N}}}_2 \uplus {\Sigma^{\mathrm{N}}}_3, effect, {E^{\mathrm{K}}}_2} \text{ CHECK\_LETBIND\_VAL\_ANNOT}$$

$$\frac{\begin{array}{l} \langle E^{\mathrm{T}}, E^{\mathrm{D}}\rangle, t \vdash pat : u \triangleright pat', {E^{\mathrm{T}}}_1, {\Sigma^{\mathrm{N}}}_1 \\ \langle (E^{\mathrm{T}} \uplus {E^{\mathrm{T}}}_1), E^{\mathrm{D}}\rangle, u \vdash exp : u' \triangleright exp', \langle {\Sigma^{\mathrm{N}}}_2, effect\rangle, {E^{\mathrm{T}}}_2 \end{array}}{\langle E^{\mathrm{T}}, E^{\mathrm{D}}\rangle \vdash \mathbf{let}\ pat = exp \triangleright \mathbf{let}\ pat' = exp', {E^{\mathrm{T}}}_1, {\Sigma^{\mathrm{N}}}_1 \uplus {\Sigma^{\mathrm{N}}}_2, effect, \{\,\}} \text{ CHECK\_LETBIND\_VAL\_NOANNOT}$$

$\boxed{E^{\mathrm{D}} \vdash \mathit{type\_def} \triangleright E}$ Check a type definition

$$\frac{E^{\mathrm{D}} \vdash \mathit{typschm} \rightsquigarrow t, E^{\mathrm{K}}, \Sigma^{\mathrm{N}}}{E^{\mathrm{D}} \vdash \mathbf{typedef}\ \mathit{id}\ \mathit{name\_scm\_opt} = \mathit{typschm} \triangleright \langle \{\,\}, \langle \{\,\}, \{\mathit{id} \mapsto E^{\mathrm{K}}, \Sigma^{\mathrm{N}}, \mathbf{None}, t\}, \{\,\}, \{\,\}\rangle\rangle}\ \text{CHECK\_TD\_ABBREV}$$

$$\frac{\begin{array}{l} E^{\mathrm{D}} \vdash \mathit{typ}_1 \rightsquigarrow t_1 \quad .. \quad E^{\mathrm{D}} \vdash \mathit{typ}_n \rightsquigarrow t_n \\ E^{\mathrm{R}} \equiv \{\{\mathit{id}_1 : t_1, .., \mathit{id}_n : t_n\} \mapsto x\} \end{array}}{E^{\mathrm{D}} \vdash \mathbf{typedef}\ x\ \mathit{name\_scm\_opt} = \mathbf{const\ struct}\ \{\mathit{typ}_1\ \mathit{id}_1; ..; \mathit{typ}_n\ \mathit{id}_n\ ;^{?}\} \triangleright \langle \{\,\}, \langle \{x \mapsto \mathit{K\_Typ}\}, \{\,\}, E^{\mathrm{R}}, \{\,\}\rangle\rangle}\ \text{CHECK\_TD\_UNQUANT\_RECORD}$$

$$\frac{\begin{array}{l} \overline{\langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle \vdash \mathit{quant\_item}_i \rightsquigarrow {E^{\mathrm{K}}}_i, {\Sigma^{\mathrm{N}}}_i}^{\,i} \\ \langle E^{\mathrm{K}} \uplus \overline{{E^{\mathrm{K}}}_i}^{\,i}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle \vdash \mathit{typ}_1 \rightsquigarrow t_1 \quad .. \quad \langle E^{\mathrm{K}} \uplus \overline{{E^{\mathrm{K}}}_i}^{\,i}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle \vdash \mathit{typ}_n \rightsquigarrow t_n \\ \{x'_1 \mapsto k_1, .., x'_m \mapsto k_m\} \equiv \uplus \overline{{E^{\mathrm{K}}}_i}^{\,i} \\ E^{\mathrm{R}}_1 \equiv \{\{\mathit{id}_1 : t_1, .., \mathit{id}_n : t_n\} \mapsto \{x'_1 \mapsto k_1, .., x'_m \mapsto k_m\}, \uplus \overline{{\Sigma^{\mathrm{N}}}_i}^{\,i}, \mathbf{None}, x\langle x'_1 .. x'_m\rangle\} \\ {E^{\mathrm{K}}}'_1 \equiv \{x \mapsto \mathit{K\_Lam}(k_1 .. k_m \to \mathit{K\_Typ})\} \end{array}}{\langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle \vdash \mathbf{typedef}\ x\ \mathit{name\_scm\_opt} = \mathbf{const\ struct\ forall}\ \overline{\mathit{quant\_item}_i}^{\,i}.\{\mathit{typ}_1\ \mathit{id}_1; ..; \mathit{typ}_n\ \mathit{id}_n\ ;^{?}\} \triangleright \langle \{\,\}, \langle {E^{\mathrm{K}}}', \{\,\}, E^{\mathrm{R}}_1, \{\,\}\rangle\rangle}\ \text{CHECK\_TD\_QUANT\_RECORD}$$

$$\frac{\begin{array}{l} E^{\mathrm{T}} \equiv \{\mathit{id}_1 \mapsto \{\,\}, \{\,\}, \mathbf{Ctor}, t_1 \to x\ \mathbf{pure}, \ldots, \mathit{id}_n \mapsto \{\,\}, \{\,\}, \mathbf{Ctor}, t_n \to x\ \mathbf{pure}\} \\ {E^{\mathrm{K}}}_1 \equiv \{x \mapsto \mathit{K\_Typ}\} \\ \langle E^{\mathrm{K}} \uplus {E^{\mathrm{K}}}_1, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle \vdash \mathit{typ}_1 \rightsquigarrow t_1 \quad \ldots \quad \langle E^{\mathrm{K}} \uplus {E^{\mathrm{K}}}_1, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle \vdash \mathit{typ}_n \rightsquigarrow t_n \end{array}}{\langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle \vdash \mathbf{typedef}\ x\ \mathit{name\_scm\_opt} = \mathbf{const\ union}\ \{\mathit{typ}_1\ \mathit{id}_1; \ldots; \mathit{typ}_n\ \mathit{id}_n\ ;^{?}\} \triangleright \langle E^{\mathrm{T}}, \langle {E^{\mathrm{K}}}_1, \{\,\}, \{\,\}, \{\,\}\rangle\rangle}\ \text{CHECK\_TD\_UNQUANT\_UNION}$$

$$\frac{\begin{array}{l} \overline{\langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle \vdash \mathit{quant\_item}_i \rightsquigarrow {E^{\mathrm{K}}}_i, {\Sigma^{\mathrm{N}}}_i}^{\,i} \\ \{x'_1 \mapsto k_1, \ldots, x'_m \mapsto k_m\} \equiv \uplus \overline{{E^{\mathrm{K}}}_i}^{\,i} \\ {E^{\mathrm{K}}}' \equiv \{x \mapsto \mathit{K\_Lam}(k_1 \ldots k_m \to \mathit{K\_Typ})\} \uplus \overline{{E^{\mathrm{K}}}_i}^{\,i} \\ \langle E^{\mathrm{K}} \uplus {E^{\mathrm{K}}}', E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle \vdash \mathit{typ}_1 \rightsquigarrow t_1 \quad \ldots \quad \langle E^{\mathrm{K}} \uplus {E^{\mathrm{K}}}', E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle \vdash \mathit{typ}_n \rightsquigarrow t_n \\ t \equiv x\langle x'_1 \ldots x'_m\rangle \\ E^{\mathrm{T}} \equiv \{\mathit{id}_1 \mapsto {E^{\mathrm{K}}}', \uplus \overline{{\Sigma^{\mathrm{N}}}_i}^{\,i}, \mathbf{Ctor}, t_1 \to t\ \mathbf{pure}, \ldots, \mathit{id}_n \mapsto {E^{\mathrm{K}}}', \uplus \overline{{\Sigma^{\mathrm{N}}}_i}^{\,i}, \mathbf{Ctor}, t_n \to t\ \mathbf{pure}\} \end{array}}{\langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle \vdash \mathbf{typedef}\ \mathit{id}\ \mathit{name\_scm\_opt} = \mathbf{const\ union\ forall}\ \overline{\mathit{quant\_item}_i}^{\,i}.\{\mathit{typ}_1\ \mathit{id}_1; \ldots; \mathit{typ}_n\ \mathit{id}_n\ ;^{?}\} \triangleright \langle E^{\mathrm{T}}, \langle {E^{\mathrm{K}}}', \{\,\}, \{\,\}, \{\,\}\rangle\rangle}\ \text{CHECK\_TD\_QUANT\_UNION}$$

$$\frac{\begin{array}{l} E^{\mathrm{T}} \equiv \{\mathit{id}_1 \mapsto x, \ldots, \mathit{id}_n \mapsto x\} \\ E^{\mathrm{E}} \equiv \{x \mapsto \{\mathit{num}_1 \mapsto \mathit{id}_1 \ldots \mathit{num}_n \mapsto \mathit{id}_n\}\} \end{array}}{E^{\mathrm{D}} \vdash \mathbf{typedef}\ x\ \mathit{name\_scm\_opt} = \mathbf{enumerate}\ \{\mathit{id}_1; \ldots; \mathit{id}_n\ ;^{?}\} \triangleright \langle E^{\mathrm{T}}, \langle \{\mathit{id} \mapsto \mathit{K\_Typ}\}, \{\,\}, \{\,\}, E^{\mathrm{E}}\rangle\rangle}\ \text{CHECK\_TD\_ENUMERATE}$$

$\boxed{E \vdash \mathit{fundef} \triangleright \mathit{fundef}', E^{\mathrm{T}}, \Sigma^{\mathrm{N}}}$ Check a function definition

$$\frac{\begin{array}{l} E^{\mathrm{T}}(id) \triangleright E^{\mathrm{K}\prime}, \Sigma^{\mathrm{N}\prime}, \mathbf{Global}, t_1 \to t\ effect \\ \overline{E^{\mathrm{D}} \vdash quant\_item_i \rightsquigarrow E^{\mathrm{K}}{}_i, \Sigma^{\mathrm{N}}{}_i}^{\,i} \\ \Sigma^{\mathrm{N}\prime\prime} \equiv \uplus\, \overline{\Sigma^{\mathrm{N}}{}_i}^{\,i} \\ E^{\mathrm{K}\prime} \equiv \overline{E^{\mathrm{K}}{}_i}^{\,i} \\ E^{\mathrm{D}}{}_1 \equiv \langle E^{\mathrm{K}\prime}, \{\,\}, \{\,\}, \{\,\}\rangle \uplus E^{\mathrm{D}} \\ E^{\mathrm{D}}{}_1 \vdash typ \rightsquigarrow u \\ E^{\mathrm{D}}{}_1 \vdash u \precapprox t, \Sigma^{\mathrm{N}}{}_2 \\ \overline{\langle E^{\mathrm{T}}, E^{\mathrm{D}}{}_1\rangle, t_1 \vdash pat_j : u_j \triangleright pat'_j, E^{\mathrm{T}}{}_j, \Sigma^{\mathrm{N}\prime\prime\prime}_j}^{\,j} \\ \overline{\langle (E^{\mathrm{T}} \uplus E^{\mathrm{T}}{}_j), E^{\mathrm{D}}{}_1\rangle, u \vdash exp_j : u' \triangleright exp'_j, \langle \Sigma^{\mathrm{N}\prime\prime\prime\prime}_j, effect'_j\rangle, E^{\mathrm{T}\prime}_j}^{\,j} \\ \Sigma^{\mathrm{N}\prime\prime\prime\prime\prime} \equiv \Sigma^{\mathrm{N}}{}_2 \uplus\, \overline{\Sigma^{\mathrm{N}\prime\prime\prime}_j \uplus \Sigma^{\mathrm{N}\prime\prime\prime\prime}_j}^{\,j} \\ effect \equiv \uplus\, \overline{effect'_j}^{\,j} \\ \Sigma^{\mathrm{N}} \equiv \mathbf{resolve}\,(\Sigma^{\mathrm{N}\prime} \uplus \Sigma^{\mathrm{N}\prime\prime} \uplus \Sigma^{\mathrm{N}\prime\prime\prime\prime\prime}) \end{array}}{\langle E^{\mathrm{T}}, E^{\mathrm{D}}\rangle \vdash \mathbf{function\ rec\ forall}\ \overline{quant\_item_i}^{\,i}\ .\ typ\ \mathbf{effect}\ effect\ \overline{id\ pat_j = exp_j}^{\,j} \triangleright \mathbf{function\ rec\ forall}\ \overline{quant\_item_i}^{\,i}\ .\ typ\ \mathbf{effect}\ effect\ \overline{id\ pat'_j = exp'_j}^{\,j}, E^{\mathrm{T}}, \Sigma^{\mathrm{N}}}\ \textsc{check\_fd\_rec}$$

$$\frac{\begin{array}{l} E^{\mathrm{T}}(id) \triangleright E^{\mathrm{K}\prime}, \Sigma^{\mathrm{N}\prime}, \mathbf{Global}, t_1 \to t\ effect \\ E^{\mathrm{D}} \vdash typ \rightsquigarrow u \\ E^{\mathrm{D}} \vdash u \precapprox t, \Sigma^{\mathrm{N}}{}_2 \\ \overline{\langle E^{\mathrm{T}}, E^{\mathrm{D}}\rangle, t_1 \vdash pat_j : u_j \triangleright pat', E^{\mathrm{T}}{}_j, \Sigma^{\mathrm{N}\prime\prime}_j}^{\,j} \\ \overline{\langle (E^{\mathrm{T}} \uplus E^{\mathrm{T}}{}_j), E^{\mathrm{D}}\rangle, u \vdash exp_j : u'_j \triangleright exp'_j, \langle \Sigma^{\mathrm{N}\prime\prime\prime}_j, effect'_j\rangle, E^{\mathrm{T}\prime}_j}^{\,j} \\ effect \equiv \uplus\, \overline{effect'_j}^{\,j} \\ \Sigma^{\mathrm{N}} \equiv \mathbf{resolve}\,(\Sigma^{\mathrm{N}}{}_2 \uplus \Sigma^{\mathrm{N}\prime} \uplus\, \overline{\Sigma^{\mathrm{N}\prime\prime}_j \uplus \Sigma^{\mathrm{N}\prime\prime\prime}_j}^{\,j}) \end{array}}{\langle E^{\mathrm{T}}, E^{\mathrm{D}}\rangle \vdash \mathbf{function\ rec}\ typ\ \mathbf{effect}\ effect\ \overline{id\ pat_j = exp_j}^{\,j} \triangleright \mathbf{function\ rec}\ typ\ \mathbf{effect}\ effect\ \overline{id\ pat'_j = exp'_j}^{\,j}, E^{\mathrm{T}}, \Sigma^{\mathrm{N}}}\ \textsc{check\_fd\_rec\_function2}$$

$$\frac{\begin{array}{l}\overline{\langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle \vdash \mathit{quant\_item}_i \rightsquigarrow {E^{\mathrm{K}}}_i, {\Sigma^{\mathrm{N}}}_i}^{\,i} \\ \Sigma^{\mathrm{N}\prime} \equiv \uplus\, \overline{{\Sigma^{\mathrm{N}}}_i}^{\,i} \\ E^{\mathrm{K}\prime} \equiv E^{\mathrm{K}} \uplus\, \overline{{E^{\mathrm{K}}}_i}^{\,i} \\ \langle E^{\mathrm{K}\prime}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle \vdash \mathit{typ} \rightsquigarrow t \\ \overline{\langle E^{\mathrm{T}}, \langle E^{\mathrm{K}\prime}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle\rangle, t_1 \vdash \mathit{pat}_j : u_j \triangleright \mathit{pat}'_j, {E^{\mathrm{T}}}_j, {\Sigma^{\mathrm{N}}}''_j}^{\,j} \\ E^{\mathrm{T}\prime} \equiv (E^{\mathrm{T}} \uplus \{\mathit{id} \mapsto E^{\mathrm{K}\prime}, \Sigma^{\mathrm{N}\prime}, \mathbf{Global}, t_1 \to t\ \mathit{effect}\}) \\ \overline{\langle (E^{\mathrm{T}\prime} \uplus {E^{\mathrm{T}}}_j), \langle E^{\mathrm{K}\prime}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle\rangle, t \vdash \mathit{exp}_j : u'_j \triangleright \mathit{exp}'_j, \langle {\Sigma^{\mathrm{N}}}'''_j, \mathit{effect}'_j\rangle, {E^{\mathrm{T}}}'_j}^{\,j} \\ \mathit{effect} \equiv \uplus\, \overline{\mathit{effect}'_j}^{\,j} \\ \Sigma^{\mathrm{N}} \equiv \mathbf{resolve}\,(\Sigma^{\mathrm{N}\prime} \uplus\, \overline{{\Sigma^{\mathrm{N}}}''_j \uplus {\Sigma^{\mathrm{N}}}'''_j}^{\,j}) \end{array}}{\langle E^{\mathrm{T}}, \langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle\rangle \vdash \mathbf{function\ rec\ forall}\ \overline{\mathit{quant\_item}_i}^{\,i}\ .\ \mathit{typ}\ \mathbf{effect}\ \mathit{effect}\ \overline{\mathit{id}\ \mathit{pat}_j = \mathit{exp}_j}^{\,j} \triangleright \mathbf{function\ rec\ forall}\ \overline{\mathit{quant\_item}_i}^{\,i}\ .\ \mathit{typ}\ \mathbf{effect}\ \mathit{effect}\ \overline{\mathit{id}\ \mathit{pat}'_j = \mathit{exp}'_j}^{\,j}, E^{\mathrm{T}\prime}, \Sigma^{\mathrm{N}}}\ \mathrm{C}$$

$$\frac{\begin{array}{l} E^{\mathrm{D}} \vdash \mathit{typ} \rightsquigarrow t \\ \overline{\langle E^{\mathrm{T}}, E^{\mathrm{D}}\rangle, t_1 \vdash \mathit{pat}_j : u_j \triangleright \mathit{pat}'_j, {E^{\mathrm{T}}}_j, {\Sigma^{\mathrm{N}}}'_j}^{\,j} \\ E^{\mathrm{T}\prime} \equiv (E^{\mathrm{T}} \uplus \{\mathit{id} \mapsto \{\,\}, \{\,\}, \mathbf{Global}, t_1 \to t\ \mathit{effect}\}) \\ \overline{\langle (E^{\mathrm{T}\prime} \uplus {E^{\mathrm{T}}}_j), E^{\mathrm{D}}\rangle, t \vdash \mathit{exp}_j : u'_j \triangleright \mathit{exp}'_j, \langle {\Sigma^{\mathrm{N}}}'_j, \mathit{effect}'_j\rangle, {E^{\mathrm{T}}}'_j}^{\,j} \\ \mathit{effect} \equiv \uplus\, \overline{\mathit{effect}'_j}^{\,j} \\ \Sigma^{\mathrm{N}} \equiv \mathbf{resolve}\,(\uplus\, \overline{{\Sigma^{\mathrm{N}}}'_j \uplus {\Sigma^{\mathrm{N}}}''_j}^{\,j}) \end{array}}{\langle E^{\mathrm{T}}, E^{\mathrm{D}}\rangle \vdash \mathbf{function\ rec}\ \mathit{typ}\ \mathbf{effect}\ \mathit{effect}\ \overline{\mathit{id}\ \mathit{pat}_j = \mathit{exp}_j}^{\,j} \triangleright \mathbf{function\ rec}\ \mathit{typ}\ \mathbf{effect}\ \mathit{effect}\ \overline{\mathit{id}\ \mathit{pat}'_j = \mathit{exp}'_j}^{\,j}, E^{\mathrm{T}\prime}, \Sigma^{\mathrm{N}}}\ \textsc{check\_fd\_rec\_function\_no\_spec}2$$

$$\frac{\begin{array}{l} E^{\mathrm{T}}(id) \triangleright E^{\mathrm{K}\prime}, \Sigma^{\mathrm{N}\prime}, \mathbf{Global}, t_1 \to t\ effect \\ \overline{\langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle \vdash quant\_item_i \rightsquigarrow E^{\mathrm{K}}{}_i, \Sigma^{\mathrm{N}}{}_i}^{\,i} \\ \Sigma^{\mathrm{N}\prime\prime} \equiv \uplus\, \overline{\Sigma^{\mathrm{N}}{}_i}^{\,i} \\ E^{\mathrm{K}\prime\prime} \equiv \overline{E^{\mathrm{K}}{}_i}^{\,i} \\ \langle E^{\mathrm{K}\prime\prime} \uplus E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle \vdash typ \rightsquigarrow u \\ \langle E^{\mathrm{K}\prime\prime} \uplus E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle \vdash u \precapprox t, \Sigma^{\mathrm{N}}{}_2 \\ \overline{\langle E^{\mathrm{T}}, \langle E^{\mathrm{K}} \uplus E^{\mathrm{K}\prime\prime}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle\rangle, t_1 \vdash pat_j : u_j \triangleright pat'_j, E^{\mathrm{T}}{}_j, \Sigma^{\mathrm{N}\prime\prime}{}_j}^{\,j} \\ \overline{\langle (E^{\mathrm{T}} \setminus id \uplus E^{\mathrm{T}}{}_j), \langle E^{\mathrm{K}} \uplus E^{\mathrm{K}\prime\prime}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle\rangle, t \vdash exp_j : u'_j \triangleright exp'_j, \langle \Sigma^{\mathrm{N}\prime\prime\prime}{}_j, effect'_j\rangle, E^{\mathrm{T}\prime}{}_j}^{\,j} \\ \Sigma^{\mathrm{N}\prime\prime\prime\prime} \equiv \uplus\, \overline{\Sigma^{\mathrm{N}\prime\prime}{}_j \uplus \Sigma^{\mathrm{N}\prime\prime\prime}{}_j}^{\,j} \\ effect \equiv \uplus\, \overline{effect'_j}^{\,j} \\ \Sigma^{\mathrm{N}} \equiv \mathbf{resolve}\,(\Sigma^{\mathrm{N}\prime} \uplus \Sigma^{\mathrm{N}\prime\prime} \uplus \Sigma^{\mathrm{N}\prime\prime\prime\prime}) \end{array}}{\langle E^{\mathrm{T}}, \langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle\rangle \vdash \mathbf{function\ forall}\ \overline{quant\_item_i}^{\,i}\ .\ typ\ \mathbf{effect}\ effect\ \overline{id\ pat_j = exp_j}^{\,j} \triangleright \mathbf{function\ forall}\ \overline{quant\_item_i}^{\,i}\ .\ typ\ \mathbf{effect}\ effect\ \overline{id\ pat'_j = exp'_j}^{\,j}, E^{\mathrm{T}}, \Sigma^{\mathrm{N}}}\ \textsc{check\_f}$$

$$\frac{\begin{array}{l} E^{\mathrm{T}}(id) \triangleright \{\,\}, \Sigma^{\mathrm{N}}{}_1, \mathbf{Global}, t_1 \to t\ effect \\ E^{\mathrm{D}} \vdash typ \rightsquigarrow u \\ E^{\mathrm{D}} \vdash u \precapprox t, \Sigma^{\mathrm{N}}{}_2 \\ \overline{\langle E^{\mathrm{T}}, E^{\mathrm{D}}\rangle, t_1 \vdash pat_j : u_j \triangleright pat_j, E^{\mathrm{T}}{}_j, \Sigma^{\mathrm{N}\prime}{}_j}^{\,j} \\ \overline{\langle (E^{\mathrm{T}} \setminus id \uplus E^{\mathrm{T}}{}_j), E^{\mathrm{D}}\rangle, u \vdash exp_j : u'_j \triangleright exp'_j, \langle \Sigma^{\mathrm{N}\prime\prime}{}_j, effect'_j\rangle, E^{\mathrm{T}\prime}{}_j}^{\,j} \\ effect \equiv \uplus\, \overline{effect'_j}^{\,j} \\ \Sigma^{\mathrm{N}} \equiv \mathbf{resolve}\,(\Sigma^{\mathrm{N}}{}_1 \uplus \Sigma^{\mathrm{N}}{}_2 \uplus\, \overline{\Sigma^{\mathrm{N}\prime}{}_j \uplus \Sigma^{\mathrm{N}\prime\prime}{}_j}^{\,j}) \end{array}}{\langle E^{\mathrm{T}}, E^{\mathrm{D}}\rangle \vdash \mathbf{function}\ \ typ\ \mathbf{effect}\ effect\ \overline{id\ pat_j = exp_j}^{\,j} \triangleright \mathbf{function}\ \ typ\ \mathbf{effect}\ effect\ \overline{id\ pat'_j = exp'_j}^{\,j}, E^{\mathrm{T}}, \Sigma^{\mathrm{N}}}\ \textsc{check\_fd\_function2}$$

$$\frac{\begin{array}{l}\overline{\langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle \vdash \mathit{quant\_item}_i \rightsquigarrow {E^{\mathrm{K}}}_i, {\Sigma^{\mathrm{N}}}_i}^{\,i} \\ \Sigma^{\mathrm{N}\prime} \equiv \uplus\, \overline{{\Sigma^{\mathrm{N}}}_i}^{\,i} \\ E^{\mathrm{K}\prime\prime} \equiv E^{\mathrm{K}} \uplus\, \overline{{E^{\mathrm{K}}}_i}^{\,i} \\ \langle E^{\mathrm{K}\prime\prime}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle \vdash \mathit{typ} \rightsquigarrow t \\ \overline{\langle E^{\mathrm{T}}, \langle E^{\mathrm{K}\prime\prime}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle\rangle, t_1 \vdash \mathit{pat}_j : u_j \triangleright \mathit{pat}_j, {E^{\mathrm{T}}}_j, {\Sigma^{\mathrm{N}\prime\prime}}_j}^{\,j} \\ E^{\mathrm{T}\prime} \equiv (E^{\mathrm{T}} \uplus \{\mathit{id} \mapsto E^{\mathrm{K}\prime\prime}, \Sigma^{\mathrm{N}\prime}, \mathbf{Global}, t_1 \to t\ \mathit{effect}\}) \\ \overline{\langle (E^{\mathrm{T}} \uplus {E^{\mathrm{T}}}_j), \langle E^{\mathrm{K}\prime\prime}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle\rangle, t \vdash \mathit{exp}_j : u'_j \triangleright \mathit{exp}'_j, \langle {\Sigma^{\mathrm{N}\prime\prime}}_j, \mathit{effect}'_j\rangle, {E^{\mathrm{T}\prime}}_j}^{\,j} \\ \mathit{effect} \equiv \uplus\, \overline{\mathit{effect}'_j}^{\,j} \\ \Sigma^{\mathrm{N}} \equiv \mathbf{resolve}\,(\Sigma^{\mathrm{N}\prime} \uplus\, \overline{{\Sigma^{\mathrm{N}\prime}}_j \uplus {\Sigma^{\mathrm{N}\prime\prime}}_j}^{\,j}) \end{array}}{\langle E^{\mathrm{T}}, \langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle\rangle \vdash \mathbf{function\ forall}\ \overline{\mathit{quant\_item}_i}^{\,i}\ .\ \mathit{typ}\ \mathbf{effect}\ \mathit{effect}\ \overline{\mathit{id}\ \mathit{pat}_j = \mathit{exp}_j}^{\,j} \triangleright \mathbf{function\ forall}\ \overline{\mathit{quant\_item}_i}^{\,i}\ .\ \mathit{typ}\ \mathbf{effect}\ \mathit{effect}\ \overline{\mathit{id}\ \mathit{pat}'_j = \mathit{exp}'_j}^{\,j}, E^{\mathrm{T}\prime}, \Sigma^{\mathrm{N}}}\ \text{CHECK\_}$$

$$\frac{\begin{array}{l} E^{\mathrm{D}} \vdash \mathit{typ} \rightsquigarrow t \\ \overline{\langle E^{\mathrm{T}}, E^{\mathrm{D}}\rangle, t_1 \vdash \mathit{pat}_j : u_j \triangleright \mathit{pat}'_j, {E^{\mathrm{T}}}_j, {\Sigma^{\mathrm{N}\prime}}_j}^{\,j} \\ E^{\mathrm{T}\prime} \equiv (E^{\mathrm{T}} \uplus \{\mathit{id} \mapsto \{\,\}, \Sigma^{\mathrm{N}}, \mathbf{Global}, t_1 \to t\ \mathit{effect}\}) \\ \overline{\langle (E^{\mathrm{T}} \uplus {E^{\mathrm{T}}}_j), E^{\mathrm{D}}\rangle, t \vdash \mathit{exp}_j : u'_j \triangleright \mathit{exp}', \langle {\Sigma^{\mathrm{N}\prime}}_j, \mathit{effect}'_j\rangle, {E^{\mathrm{T}\prime}}_j}^{\,j} \\ \mathit{effect} \equiv \uplus\, \overline{\mathit{effect}'_j}^{\,j} \\ \Sigma^{\mathrm{N}} \equiv \mathbf{resolve}\,(\uplus\, \overline{{\Sigma^{\mathrm{N}\prime}}_j \uplus {\Sigma^{\mathrm{N}\prime\prime}}_j}^{\,j}) \end{array}}{\langle E^{\mathrm{T}}, E^{\mathrm{D}}\rangle \vdash \mathbf{function}\ \ \mathit{typ}\ \mathbf{effect}\ \mathit{effect}\ \overline{\mathit{id}\ \mathit{pat}_j = \mathit{exp}_j}^{\,j} \triangleright \mathbf{function}\ \ \mathit{typ}\ \mathbf{effect}\ \mathit{effect}\ \overline{\mathit{id}\ \mathit{pat}'_j = \mathit{exp}'_j}^{\,j}, E^{\mathrm{T}\prime}, \Sigma^{\mathrm{N}}}\ \text{CHECK\_FD\_FUNCTION\_NO\_SPEC2}$$

$\boxed{E \vdash \mathit{val\_spec} \triangleright E^{\mathrm{T}}}$ Check a value specification

$$\frac{E^{\mathrm{D}} \vdash \mathit{typschm} \rightsquigarrow t, {E^{\mathrm{K}}}_1, \Sigma^{\mathrm{N}}}{\langle E^{\mathrm{T}}, E^{\mathrm{D}}\rangle \vdash \mathbf{val}\ \mathit{typschm}\ \mathit{id} \triangleright \{\mathit{id} \mapsto {E^{\mathrm{K}}}_1, \Sigma^{\mathrm{N}}, \mathbf{Global}, t\}}\ \text{CHECK\_SPEC\_VAL\_SPEC}$$

$$\frac{E^{\mathrm{D}} \vdash \mathit{typschm} \rightsquigarrow t, {E^{\mathrm{K}}}_1, \Sigma^{\mathrm{N}}}{\langle E^{\mathrm{T}}, E^{\mathrm{D}}\rangle \vdash \mathbf{val\ extern}\ \mathit{typschm}\ \mathit{id} = \mathit{string} \triangleright \{\mathit{id} \mapsto {E^{\mathrm{K}}}_1, \Sigma^{\mathrm{N}}, \mathbf{Extern}\,, t\}}\ \text{CHECK\_SPEC\_EXTERN}$$

$\boxed{E^{\mathrm{D}} \vdash \mathit{default\_spec} \triangleright E^{\mathrm{T}}, {E^{\mathrm{K}}}_1}$ Check a default typing specification

$$\frac{E^{\mathrm{K}} \vdash \mathit{base\_kind} \rightsquigarrow k}{\langle E^{\mathrm{K}}, E^{\mathrm{A}}, E^{\mathrm{R}}, E^{\mathrm{E}}\rangle \vdash \mathbf{default}\ \mathit{base\_kind}\ 'x \triangleright \{\,\}, \{'x \mapsto k\ \mathbf{default}\}}\ \text{CHECK\_DEFAULT\_KIND}$$

$$\frac{E^{\mathrm{D}} \vdash \mathit{typschm} \rightsquigarrow t, {E^{\mathrm{K}}}_1, {\Sigma^{\mathrm{N}}}}{E^{\mathrm{D}} \vdash \mathbf{default}\ \mathit{typschm}\ \mathit{id} \triangleright \{\mathit{id} \mapsto {E^{\mathrm{K}}}_1, {\Sigma^{\mathrm{N}}}, \mathbf{Default}, t\}, \{\,\}} \quad \text{CHECK\_DEFAULT\_TYP}$$

$\boxed{E \vdash \mathit{def} \triangleright \mathit{def}', E'}$ Check a definition

$$\frac{E^{\mathrm{D}} \vdash \mathit{type\_def} \triangleright E}{\langle E^{\mathrm{T}}, E^{\mathrm{D}} \rangle \vdash \mathit{type\_def} \triangleright \mathit{type\_def}, \langle E^{\mathrm{T}}, E^{\mathrm{D}} \rangle \uplus E} \quad \text{CHECK\_DEF\_TDEF}$$

$$\frac{E \vdash \mathit{fundef} \triangleright \mathit{fundef}', E^{\mathrm{T}}, {\Sigma^{\mathrm{N}}}}{E \vdash \mathit{fundef} \triangleright \mathit{fundef}', E \uplus \langle E^{\mathrm{T}}, \epsilon \rangle} \quad \text{CHECK\_DEF\_FDEF}$$

$$\frac{\begin{array}{l} E \vdash \mathit{letbind} \triangleright \mathit{letbind}', \{\mathit{id}_1 \mapsto t_1, \ldots, \mathit{id}_n \mapsto t_n\}, {\Sigma^{\mathrm{N}}}, \mathbf{pure}, E^{\mathrm{K}} \\ {\Sigma^{\mathrm{N}}}_1 \equiv \mathbf{resolve}\,({\Sigma^{\mathrm{N}}}) \end{array}}{E \vdash \mathit{letbind} \triangleright \mathit{letbind}', E \uplus \langle \{\mathit{id}_1 \mapsto E^{\mathrm{K}}, {\Sigma^{\mathrm{N}}}, \mathbf{None}, t_1, \ldots, \mathit{id}_n \mapsto E^{\mathrm{K}}, {\Sigma^{\mathrm{N}}}, \mathbf{None}, t_n\}, \epsilon \rangle} \quad \text{CHECK\_DEF\_VDEF}$$

$$\frac{E \vdash \mathit{val\_spec} \triangleright E^{\mathrm{T}}}{E \vdash \mathit{val\_spec} \triangleright \mathit{val\_spec}, E \uplus \langle E^{\mathrm{T}}, \epsilon \rangle} \quad \text{CHECK\_DEF\_VSPEC}$$

$$\frac{E^{\mathrm{D}} \vdash \mathit{default\_spec} \triangleright {E^{\mathrm{T}}}_1, {E^{\mathrm{K}}}_1}{\langle E^{\mathrm{T}}, E^{\mathrm{D}} \rangle \vdash \mathit{default\_spec} \triangleright \mathit{default\_spec}, \langle (E^{\mathrm{T}} \uplus {E^{\mathrm{T}}}_1), E^{\mathrm{D}} \uplus \langle {E^{\mathrm{K}}}_1, \{\,\}, \{\,\}, \{\,\} \rangle \rangle} \quad \text{CHECK\_DEF\_DEFAULT}$$

$$\frac{E^{\mathrm{D}} \vdash \mathit{typ} \rightsquigarrow t}{\langle E^{\mathrm{T}}, E^{\mathrm{D}} \rangle \vdash \mathbf{register}\ \mathit{typ}\ \mathit{id} \triangleright \mathbf{register}\ \mathit{typ}\ \mathit{id}, \langle (E^{\mathrm{T}} \uplus \{\mathit{id} \mapsto \mathbf{register}\ \langle t \rangle\}), E^{\mathrm{D}} \rangle} \quad \text{CHECK\_DEF\_REGISTER}$$

$\boxed{E \vdash \mathit{defs} \triangleright \mathit{defs}', E'}$ Check definitions, potentially given default environment of built-in library

$$\frac{\begin{array}{l} E \vdash \mathit{def} \triangleright \mathit{def}', E_1 \\ E \uplus E_1 \vdash \overline{\mathit{def}_i}^{\,i} \triangleright \overline{\mathit{def}'_i}^{\,i}, E_2 \end{array}}{E \vdash \mathit{def}\ \overline{\mathit{def}_i}^{\,i} \triangleright \mathit{def}'\ \overline{\mathit{def}'_i}^{\,i}, E_2} \quad \text{CHECK\_DEFS\_DEFS}$$

4 Sail operational semantics {TODO}