

# ▣ Doubts with Lakshay Kharya [Week 4]

Swap Implement  $\rightarrow (a, b) \rightarrow$  Swap

Three Methods:

① Temp variable.

②  $(+, -)$  Arithmetic

③ XOR.

Arithmetic :-

int a, int b

①  $a = a + b$

②  $b = a - b$

↓

$5 - 3 = 2$

$b = 2$

③  $a = a - b$

$= 5 - 2$

$a = 3$

2 a

3 b

(i)  $a = a + b$

(ii)  $b = a - b$

(iii)  $a = a - b$

XOR :-

(fastest)

int a, int b

↓  
2

↓  
3

$a = a \wedge b = 2 \wedge 3$

$b = a \wedge b = 2 \wedge 3 \wedge 3 = 2$

$a = a \wedge b = 2 \wedge 3 \wedge 2 = 3$

(i)  $a = a \wedge b$

(ii)  $b = a \wedge b$

(iii)  $a = a \wedge b$

Rotate Array

- ① Extra space. ✓
- ② Single rotate  $k$  times. ✓
- ③ Modulo  $((i+k)\%n) \rightarrow O(n)$  ✓
- ④ Reversal method  $\rightarrow O(n)$

④ Reversal Method

$K = 3$

I/P  $\rightarrow 1|2|3|4|5|6|7$ Also

↓

 $7|6|5|4|3|2|1$ 

① Reverse whole Array.

 $5|6|7|4|3|2|1$ 

Rev(0, n-1)

↓

② Rev(0, K-1)

O/P  $\rightarrow 5|6|7|1|2|3|4$ 

③ Rev(K, n-1)

Code.

$K = K \% \text{nums.size}$

what if  $K = \text{nums.size}$   
or  $K > \text{nums.size}$

```
reverse(nums.begin(), nums.end());
reverse(nums.begin(), nums.begin() + K);
reverse(nums.begin() + K, nums.end());
```

 $\rightarrow$  Reverse (STL) $1|2|3|4|5|$ 

↓

nums.begin()

nums.end()

0

 $\rightarrow$ 

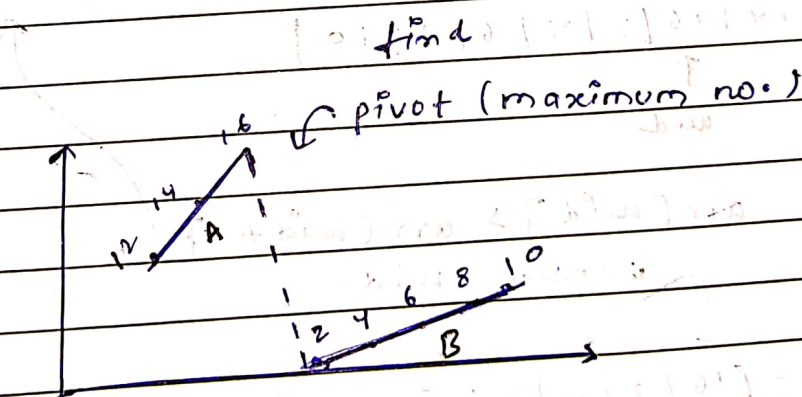
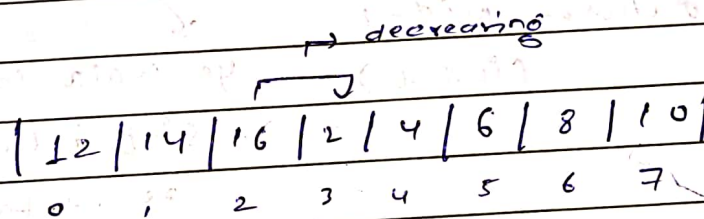
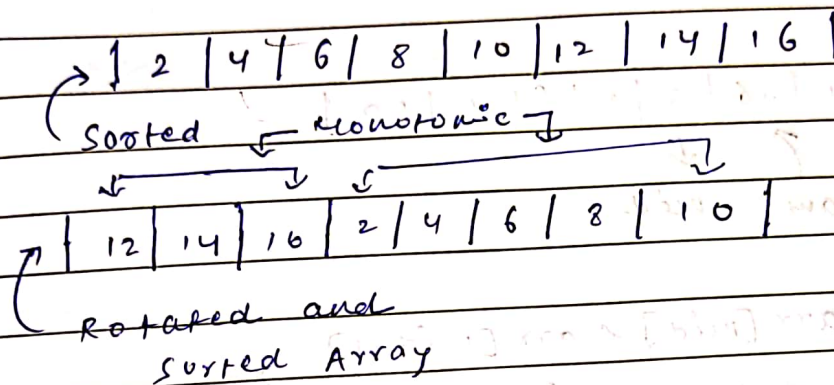
n-1



## ▣ Searching and Sorting - Level 2

→ find pivot element  
↓  
Subquestion.

[ Search in a rotated & Sorted Array ]  
↓  
question



### # Approach - 1

→ Linear Search.  $O(n)$

### # Approach - 2

→ sort  $O(n \log n)$

## #3 Approach -

Binary search (algo)

[ 12 | 14 | 16 | 2 | 4 | 6 | 8 | 10 ]

(A)    ↑    ↑                      (B)

↳ separately handle.

[ 12 | 14 | 16 | 2 | 4 | 6 | 8 | 10 ]

↑    ↓  
ans   mid

 $arr[mid] < arr[mid-1]$ 

return mid-1

(index)

ye wala condition

first index do element

pe lagta hai

ans

[ 12 | 14 | 16 | 2 | 4 | 6 | 8 | 10 ]

↑  
mid

 $arr[mid] > arr[mid+1]$ 

return mid.

[ 12 | 14 | 16 | 2 | 4 | 6 | 8 | 10 ]

(A)

(B)

↳ Agar me is line p note to mai  
pucka 12 k upar hota.

 $arr[start] > arr[mid]$ 

↳ left

 $end = mid - 1;$



else  $s = mid + 1$   $\rightarrow$  Right me jayenge.

Hum pivot index mil gaya hai aur  
ab hum array me search karenge.

| 12 | 14 | 16 | 2 | 4 | 6 | 8 | 10 |

⏟ ⏟  
④ ⓐ

(Binary search) (Binary search)

target = 14

Agar tumhara target 12 se 16 ke beech me  
data hai to ④ me binary search  
Laga lo

Agar tumhara target 2 se 10 ke beech me  
data hai to ② me binary search  
Laga lo.

Edge case. ✗

[1, 3]  $\rightarrow$  Isme code fatega.

↓  
mid arr[mid] < arr[mid-1]

that's why we take condn invalid index.

if (mid < n-1)  
if (mid > 0) } on pivot element  
function.

code // finding pivot element.

```
int pivotIndex(vector<int> &nums) {  
    int n = nums.size();  
    int s = 0;  
    int e = n-1;  
    while (s <= e) {  
        if (s == e) return s;  
        // corner case if the size of array is 1.  
        int mid = (s+e)/2;  
        if (mid < n-1 & nums[mid] > nums[mid+1])  
            return mid;  
        else if (mid > 0 & nums[mid] < nums[mid-1])  
            return mid;  
        else if (nums[mid] < nums[s])  
            e = mid-1;  
        else s = mid+1;  
    }  
    return -1;  
}
```

// Binary search for finding target

```
int binarySearch(vector<int> &nums, int s, int e, int target)  
{  
    while (s <= e)  
    {  
        int mid = s + (e-s)/2;  
        if (nums[mid] == target) return mid;  
        else if (target > nums[mid]) s = mid+1;  
        else e = mid-1;  
    }  
    return -1;  
}
```



// searching in sorted array.

```
int search (vector<int> &nums, int target)
{
    int n = nums.size() - 1;
    int pivotRes = pivotIndex(nums);
    if (target >= nums[0] && target <= nums[pivotRes])
    {
        return binarySearch(nums, 0, pivotRes, target);
    }
    else
    {
        return binarySearch(nums, pivotRes + 1, n, target);
    }
}

return -1;
```

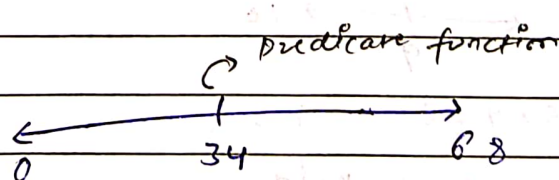
Q. Sqrt(x) (without using STL function)

i/p → x

o/p → Sqrt(x) →  $\sqrt{x}$  ans.

eg: x = 68

o/p  $\sqrt{68} \rightarrow 8$  (nearly)



steps

① search space

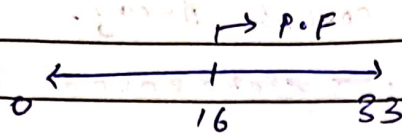
② predicate function

s = 0  
e = 68  
mid = 34

$34 \times 34 = 68 \rightarrow \text{false}$

$34 \times 34 > 68 \rightarrow \text{left me jage.}$

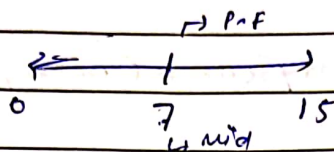
$e = 34 - 1 = 33$



$$\text{mid} = 16$$

$$16 \times 16 = 256 > 68$$

$$e = 16 - 1 = 15$$

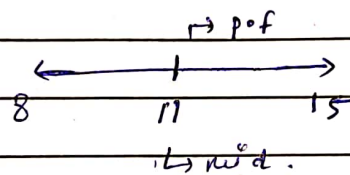


$$7 \times 7 = 49 < 68$$

ans = 7  $\hookrightarrow$  store k or length  
ans

right me jayenge.

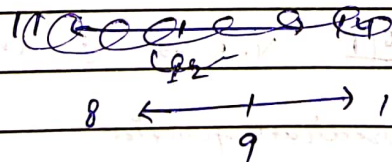
$$s = \text{mid} + 1 = 7 + 1 = 8$$



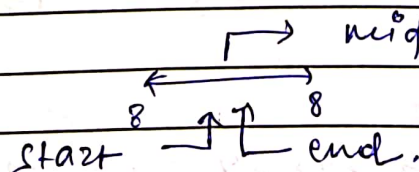
$$11 \times 11 = 121 > 68$$

$$e = \text{mid} - 1$$

$$e = 11 - 1 = 10$$



$$9 \times 9 > 68 \quad e = 8$$



$$8 \times 8 < 68$$

$$\text{ans} = 8$$



steps

- ① search space
- ② Predicate function  $\rightarrow$  True / false
- ③ Store ans.

Code

```
int mySqrt (int x) {  
    int s = 0;  
    int e = x;  
    int ans = -1;  
    while (s <= e) {  
        long long int mid = s + (e - s) / 2;  
        long long int sqrt = mid * mid;  
        if (x == sqrt) return mid;  
        else if (sqrt > x) e = mid - 1;  
        else {  
            ans = mid;  
            s = mid + 1;  
        }  
    }  
    return ans;  
}
```

Time Complexity =  $O(\log n)$   
Space complexity =  $O(1)$