

# Mathematical Framework to Compute Optimal $x_i$ at Time $t_i$

We are given:

- A total order size  $S$  that must be completed by the end of day.
- The day is split into  $N$  discrete trading intervals.
- For each time step  $t_i$ , we have a temporary impact function  $g_i(x) \approx \beta_i x$ , assumed linear.

## Goal:

Determine the allocation vector  $x = (x_1, x_2, \dots, x_N)$  such that:

- The total shares traded satisfies:  $\sum x_i = S$
- The total temporary impact (slippage) cost:  $\sum \beta_i * x_i$  is minimized.

## Mathematical Formulation (Linear Program):

We set this up as a constrained linear optimization problem:

Minimize:  $\sum \beta_i * x_i$  Subject to:  $\sum x_i = S$   $x_i \geq 0$  for all  $i \in \{1, \dots, N\}$

## Interpretation:

- $\beta_i$ : Cost per share at time  $t_i$  (i.e., slope of the linear temporary impact model at  $t_i$ ).
- $x_i$ : Number of shares to execute at time  $t_i$ .

Since both the objective function and constraints are linear, this problem can be efficiently solved using standard Linear Programming (LP) solvers such as:

- `scipy.optimize.linprog` (Python)
- Gurobi, CPLEX, or other commercial solvers

## Algorithm Sketch:

1. **Input:**  $S$ , vector  $\beta=(\beta_1,\beta_2,\dots,\beta_N)$
2. **Define LP:**
  - Objective:  $\min \beta^T x$
  - Constraints:
    - $Ax=S$  where  $A=[1 \ 1 \ \dots \ 1] \in \mathbb{R}^{1 \times N}$
    - $x \geq 0$
3. **Solve** LP using simplex or interior-point method
4. **Output:** Optimal allocation  $x^*$
- 1.

## Notes:

- The LP ensures the solution is globally optimal due to convexity.
- If integer quantities are required, this becomes an Integer Linear Program (ILP), solvable using branch-and-bound methods.
- If future  $\beta_i$  values are unknown at time  $t_i$ , then a model predictive control (MPC) or online optimization variant would be necessary.