Microprocessor Lab Report

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1 Assignment 1

1.1 Objective

Find out the sum of the first 30 natural numbers.

1.2 Tool/Experimental setup considered

- Jubin's 8085 Simulator

1.3 Procedure

We know that

$$1 + 2 + 3 + ... + 29 + 30 = \frac{30 \times 29}{2} = 435 = 01D1H$$

This result is not possible to store in a single register, so we need to use register pair to store the result.

1.4 Program

```
;end result is 465, more than 255 hence we need to do extended additions
       MVI D,1E ;setup D, the counter as 30
       MVI C,01 ;setup BC as 1

L1:    DAD B    ;Double add BC to HL
       INX B    ;extended increment BC
       DCR D    ;decrement D
       JNZ L1   ;if D becomes 0, Z flag becomes 0 and we break
       HLT
;ans will be in HL
```

Listing 1: assembly program to find sum of the first 30 natural numbers

1.5 Experimentation

Register	Value	7	6	5	4	3	2	1	0
Accumulator	00	0	0	0	0	0	0	0	0
Register B	00	0	0	0	0	0	0	0	0
Register C	1F	0	0	0	1	1	1	1	1
Register D	00	0	0	0	0	0	0	0	0
Register E	00	0	0	0	0	0	0	0	0
Register H	01	0	0	0	0	0	0	0	1
Register L	D1	1	1	0	1	0	0	0	1
Memory(M)	00	0	0	0	0	0	0	0	0

Figure 1: Register configuration after execution (observe HL)

1.6 Conclusion

We see that after the execution of program, the data stored in HL register pair is 01D1H, which is the hexadecimal value of 435.
Hence the program is working as expected.

2 Assignment 2

2.1 Objective

From an array of 10-byte size integers (unsigned) find out the maximum and minimum.

2.2 Tool/Experimental setup considered

- Jubin's 8085 Simulator

2.3 Procedure

The idea is to linearly iterate through all the values of the arr, and update the register for minimum(C) and maximum(B) values.

2.4 Program

```
;Actual Program
# ORG 5000H
# ARR: DB 5,2,3,4,F,C,7,A,B,1
# ORG 0000
       LXI H,ARR
       MOV B,M  ;B is maximum val
       MOV C,M  ;C is minimum val
       MVI D,0A
;CMP R does (A - R) in background
;If A - R > 0 then Cy = 0, Z = 0
;If A - R = 0 then Cy = 0, Z = 1
;If A - R < 0 then Cy = 1, Z = 1

LP:    MOV A,M
       CMP B
       JC MIN   ;will Jump when Cy = 1, A - B < 0, A < B
       MOV B,A  ;will only happen if A > B

MIN:   CMP C
       JNC SKIP ;will Jump when Cy = 0, A - C > 0, A > C
       MOV C,A  ;will only happen if A < C

SKIP:  INX H
       DCR D
       JNZ LP
       HLT
```

Listing 2: assembly program to find minimum and maximum number in 10-byte unsigned array

2.5 Experimentation

Register	Value	7	6	5	4	3	2	1	0
Accumulator	01	0	0	0	0	0	0	0	1
Register B	0F	0	0	0	0	1	1	1	1
Register C	01	0	0	0	0	0	0	0	1
Register D	00	0	0	0	0	0	0	0	0
Register E	00	0	0	0	0	0	0	0	0
Register H	50	0	1	0	1	0	0	0	0
Register L	0A	0	0	0	0	1	0	1	0
Memory(M)	00	0	0	0	0	0	0	0	0

(a) result for {5,2,3,4,F,C,7,A,B,1}

Register	Value	7	6	5	4	3	2	1	0
Accumulator	06	0	0	0	0	0	1	1	0
Register B	0F	0	0	0	0	1	1	1	1
Register C	06	0	0	0	0	0	1	1	0
Register D	00	0	0	0	0	0	0	0	0
Register E	00	0	0	0	0	0	0	0	0
Register H	50	0	1	0	1	0	0	0	0
Register L	0A	0	0	0	0	1	0	1	0
Memory(M)	00	0	0	0	0	0	0	0	0

(b) result for {F,E,D,C,B,A,9,8,7,6}

Figure 2: Result for different inputs

2.6 Conclusion

We see that that after program execution, B has the maximum value of array, and C has the minimum value of the array.
Hence the program is working as expected.

3 Assignment 3

3.1 Objective

Write a routine that produces a delay. The delay value must be passed to register pair DE.

3.2 Tool/Experimental setup considered

- Jubin's 8085 Simulator

3.3 Procedure

Idea is to assign DE a very big value (say FFFF), and decrement it in a loop till DE becomes 0 to produce delay in execution.

3.4 Program

```
1        LXI D,FFFF
2        CALL DELAY
3        HLT
4  ;delay: this subroutine produces delay
5  ;in: value in DE pair
6  ;out: none
7  ;destroys: A
8
9  DELAY: DCX D    ;doesn't affect any flags, that's why doing OR
10       MOV A,E
11       ORA D    ;will give 0 only when both D and E 00
12       JNZ DELAY
13       RET
```

Listing 3: assembly program to produce delay

3.5 Conclusion

We see that the code runs for sometime, then completes it's execution, signifying that the delay function worked and delayed execution of CPU for some time.
Hence the program is working as expected.

4 Assignment 4

4.1 Objective

Write a subroutine to move a block of bytes from location X to location Y.
Note that the caller would specify

- X, the source address
- Y, the destination address
- Z, the block size

Note that X, Y and Z are 16-bit quantities.

4.2 Tool/Experimental setup considered

- Jubin's 8085 Simulator

4.3 Procedure

Start reading numbers from location X and save them to location Y, after each iteration, update address of X and Y to next byte. Do this Z times and the whole block is copied.

4.4 Program

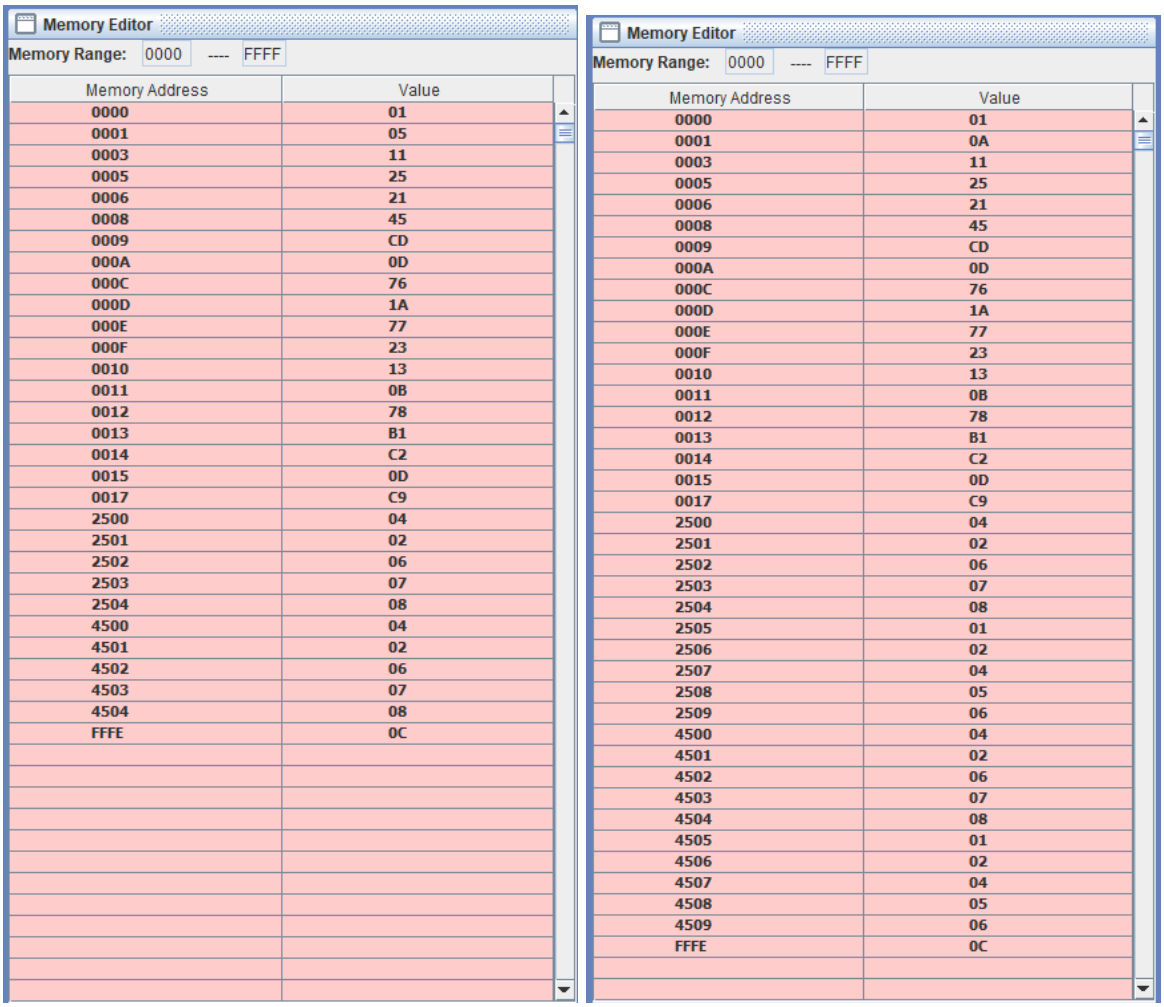
```
;setup of data
#ORG 2500
#ARR: DB 4,2,6,7,8

# DESTLOC EQU 4500
#ORG 0000
;let BC = 5, X = 2500, Y = 4500
;hence add data from 2500 to 2504
       LXI B,0005
       LXI D, ARR
       LXI H,DESTLOC
       CALL MOVE
       HLT
;MOVE: move Z number of bytes from loc (X to X+Z) to loc (Y to Y+Z)
;Z store in BC
;X store in DE
;Y store in HL

MOVE:  LDAX D   ;A = Mem[DE]
       MOV M,A  ;Mem[HL] = A
       INX H    ;HL++
       INX D    ;DE++
       DCX B    ;BC--
;DCX doesn't set flags so do manual check by OR
       MOV A,B
       ORA C
       JNZ MOVE
       RET
```

Listing 4: assembly program to move block

4.5 Experimentation



Memory Editor

Memory Range: 0000 ---- FFFF

Memory Address	Value
0000	01
0001	05
0003	11
0005	25
0006	21
0008	45
0009	CD
000A	0D
000C	76
000D	1A
000E	77
000F	23
0010	13
0011	0B
0012	78
0013	B1
0014	C2
0015	0D
0017	C9
2500	04
2501	02
2502	06
2503	07
2504	08
4500	04
4501	02
4502	06
4503	07
4504	08
FFFE	0C

Memory Editor

Memory Range: 0000 ---- FFFF

Memory Address	Value
0000	01
0001	0A
0003	11
0005	25
0006	21
0008	45
0009	CD
000A	0D
000C	76
000D	1A
000E	77
000F	23
0010	13
0011	0B
0012	78
0013	B1
0014	C2
0015	0D
0017	C9
2500	04
2501	02
2502	06
2503	07
2504	08
2505	01
2506	02
2507	04
2508	05
2509	06
4500	04
4501	02
4502	06
4503	07
4504	08
4505	01
4506	02
4507	04
4508	05
4509	06
FFFE	0C

(a) result for {4,2,6,7,8}

(b) result for {4,2,6,7,8,1,2,4,5,6}

Figure 3: Result for different inputs

4.6 Conclusion

We see that all the data from location 2500(X) to (2500 + Z) has been copied to location 4500(Y) to (4500 + Z) [Z is 5 in 3a and 10 in 3b].
Hence the program is working as expected.

5 Assignment 5

5.1 Objective

Write a function isODD(unsigned n) in assembly that takes an unsigned integer (a byte) and determines if it is odd (returns 1) or 0 if it is even.

5.2 Tool/Experimental setup considered

- Jubin's 8085 Simulator

5.3 Procedure

Odd numbers will always be in the form of $2x + 1$, which means that they will have 1 as their LSB. So we just check if $number \wedge 01$ is 1 or not. If the result is 1, then the number is odd, else it is even.

5.4 Program

```
# NUM EQU 49 ;73 in hex
        MVI B,NUM
        CALL ISODD
        HLT
;isODD(n): function which tell if n is odd or even
;in: n = B
;out: ans in C, if n even C = 0,else 1
;destroys: A
;idea: X AND 01 = 0 if X is even, 1 is X is Odd

ISODD: MOV A,B
        ANI 01    ;A = A and 01
        MOV C,A   ;store result in C
        RET
```

Listing 5: assembly program to check if given number isOdd or not

5.5 Experimentation

Register	Value	7	6	5	4	3	2	1	0
Accumulator	01	0	0	0	0	0	0	0	1
Register B	49	0	1	0	0	1	0	0	1
Register C	01	0	0	0	0	0	0	0	1
Register D	00	0	0	0	0	0	0	0	0
Register E	00	0	0	0	0	0	0	0	0
Register H	00	0	0	0	0	0	0	0	0
Register L	00	0	0	0	0	0	0	0	0
Memory(M)	06	0	0	0	0	0	1	1	0

(a) result for odd number(73)

Register	Value	7	6	5	4	3	2	1	0
Accumulator	00	0	0	0	0	0	0	0	0
Register B	4A	0	1	0	0	1	0	1	0
Register C	00	0	0	0	0	0	0	0	0
Register D	00	0	0	0	0	0	0	0	0
Register E	00	0	0	0	0	0	0	0	0
Register H	00	0	0	0	0	0	0	0	0
Register L	00	0	0	0	0	0	0	0	0
Memory(M)	06	0	0	0	0	0	1	1	0

(b) result for even number(73)

Figure 4: Result for both even and odd numbers

5.6 Conclusion

We see that in case of odd number, C is 1 after execution and in case of even number, C is 0 after execution.
Hence the program is working as expected.

6 Assignment 6

6.1 Objective

Write a function to add two multi-byte numbers stored in location X and Y. The result is stored in X. Pass a parameter Z indicating the no. of bytes to be added.

6.2 Tool/Experimental setup considered

- Jubin's 8085 Simulator

6.3 Procedure

We simulate the default way of adding numbers, we go from right to left, adding (with carry) the numbers and adding it to stack, then we keep popping the elements and save it in X.

6.4 Program

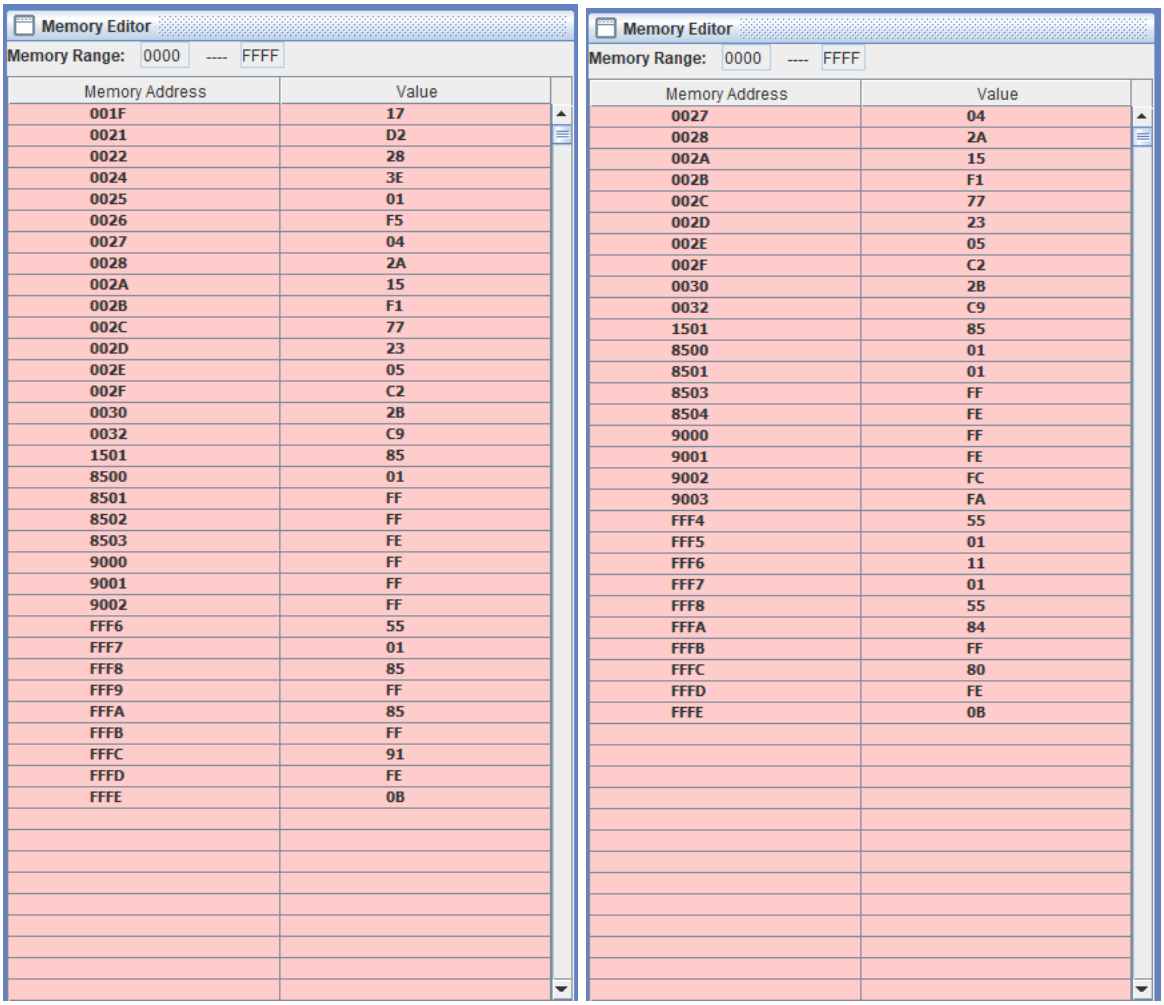
```
# ORG 8500
# NUM1: DB FF,FF,FF
# ORG 9000
# NUM2: DB FF,FF,FF
# NUMBYTE EQU 3
# ORG 0000
       MVI C,NUMBYTE    ;C <- 03
       LXI H,NUM1       ;HL <- 8500
       LXI D,NUM2       ;DE <- 9000
       CALL ADDB
       HLT
;add Z bytes, first number starts from X, other number starts from Y
;Z->C
;X->HL
;Y->DE
;destroys -> B
# SAVELOC EQU 1500
ADDB:  SHLD SAVELOC     ;save start of X to 1500 addr to be read later
       MVI B,00         ;set B as 0, will act as counter afterward
;now we need to put HL and DE pair to back of array
       XCHG             ;swap HL and DE
       DAD B            ;HL = HL + BC [with B = 0] [HL now is DE]
       DCX H
       XCHG             ;swap HL and DE again
       DAD B            ;HL = HL + BC
       DCX H
;now both DE and HL are in the end of their array
LOOP:  LDAX D           ;A -> Mem[DE]
       ADC M            ;A -> A+ Mem[HL] + carry
       PUSH PSW         ;push AF data to SP, first A, then F
       INR B
       DCX H
       DCX D
       DCR C
       JNZ LOOP         ;jump to Loop till C!=0
;now addition done and saved to Stack, need to check if carry exist
       JNC SKIP
;these will execute only when there is a carry
       MVI A,01
       PUSH PSW
       INR B

SKIP:  LHLD SAVELOC     ;read saved LH data to go to start of X

L1:    POP PSW          ;pop SP to AF, first F then A
       MOV M,A
       INX H
       DCR B
       JNZ L1
       RET
```

Listing 6: assembly program to add multi-byte numbers

6.5 Experimentation



(a) result for {FF,FF,FF} + {FF,FF,FF}

(b) result for {01,02,03,04} + {FF,FE,FC,FA}

Figure 5: Result of multi-byte addition (start looking from address 8500)

6.6 Conclusion

We see that the result of multi-byte addition is correct.
Hence the program is working as expected.

7 Assignment 7

7.1 Objective

Write a fast sub-routine to multiply 9 by 15.

7.2 Tool/Experimental setup considered

- Jubin's 8085 Simulator

7.3 Procedure

We use the Shift-and-Add Multiplication to fast multiply 15 and 9, as register size is 8 bits, we can do this multiplication by using a loop which runs 8 times, i.e O(1)
This method is faster than default loop method, which run in O(min(m, n)), where m and n are the numbers that will be multiplied.

7.4 Program

```
1        LXI D,000F    ;15 in hex
2        MVI A,09
3        LXI H,0000    ;result will be in HL
4        MVI B,08      ;8 bit data, 8 rotations to iterate through all the bits of A
5
6  LP:   DAD H         ;HL = HL + HL (multiply by 2)(assume left shift)
7        RAL           ;rotate A left(<-), leftmost value in C flag
8        JNC SKIP
9        DAD D         ;if 1 in A's bit, we add D also
10 SKIP: DCR B
11       JNZ LP
12       SHLD 2500
13       HLT
14 ;ans in HL register
```

Listing 7: assembly program to fast multiply 15 times 9

7.5 Experimentation

Register	Value	7	6	5	4	3	2	1	0
Accumulator	04	0	0	0	0	0	1	0	0
Register B	00	0	0	0	0	0	0	0	0
Register C	00	0	0	0	0	0	0	0	0
Register D	00	0	0	0	0	0	0	0	0
Register E	0F	0	0	0	0	1	1	1	1
Register H	00	0	0	0	0	0	0	0	0
Register L	87	1	0	0	0	0	1	1	1
Memory(M)	00	0	0	0	0	0	0	0	0

Figure 6: Register configuration after execution (look at HL)

7.6 Conclusion

$9 \times 15 = 135 = 0087H$
We see that the result in HL register is same as what we expected.
Hence the program is working as expected.

8 Assignment 8

8.1 Objective

Write a subroutine to sort a 5-element byte array (Any algorithm will do)

8.2 Tool/Experimental setup considered

- Jubin's 8085 Simulator

8.3 Procedure

We use bubble sort algorithm to sort the array.

8.4 Program

```
# ORG 2500
# ARR: DB 4,2,5,6,7
# LEN EQU 5
# ORG 0000

START:      LXI H,ARR   ;start of array
            MVI C,LEN
            DCR C       ;c-- as we will only go through first 4
            MVI D,00    ;d acts a bool, if any swap happened, d = 1

CHECK:      MOV A,M
            INX H
            CMP M       ;compare a[i] with a[i+1]
;if a[i] - a[i+1] > 0 -> Cy = 0, Z = 0 (we swap)
;if a[i] - a[i+1] = 0 -> Cy = 0, Z = 1 (we don't swap)
;if a[i] - a[i+1] < 0 -> Cy = 1, Z = 0 (we don't swap)
            JC NEXTBYTE
            JZ NEXTBYTE
;here swap occurs
            MOV B,M     ;a[i] = A, a[i+1] = B
            MOV M,A     ;a[i+1] = A
            DCX H
            MOV M,B     ;a[i] = B
            INX H
            MVI D,01    ;set flag that swap occurred

NEXTBYTE:   DCR C
            JNZ CHECK
            MOV A,D
            CPI 01      ;compare immediate A and 01
;if A - 01 > 0 -> Cy = 0, Z = 0 -> A(D) is zero, no swap taken place, exit
;if A - 01 = 0 -> Cy = 0, Z = 1 -> A or D is 1, swap taken place redo
            JZ START
            HLT
```

Listing 8: assembly program to bubble sort array

8.5 Experimentation

Memory Editor

Memory Range: 0000 ---- FFFF

Memory Address	Value
0000	21
0002	25
0003	0E
0004	05
0005	0D
0006	16
0008	7E
0009	23
000A	BE
000B	DA
000C	18
000E	CA
000F	18
0011	46
0012	77
0013	2B
0014	70
0015	23
0016	16
0017	01
0018	0D
0019	C2
001A	08
001C	7A
001D	FE
001E	01
001F	CA
0022	76
2500	02
2501	04
2502	05
2503	06
2504	07

Figure 7: Register configuration after execution (look from address 2500)

8.6 Conclusion

We see that after program execution, values of address 2500 - 2504 is sorted in ascending order. Hence the program is working as expected.

9 Assignment 9

9.1 Objective

Write a sub-routine to store all the registers (A, F, B, C, D, E, H, L, I, SPL, SPH, PCL, PC, in that order) starting from location MYREGISTERS.

9.2 Tool/Experimental setup considered

- Jubin's 8085 Simulator

9.3 Program

```
        MVI A,10      ;set A
        LXI B,5092    ;set BC
        LXI D,2794    ;set DE
        LXI H,3792    ;set HL
        SIM           ;set inturrept mask as value of A
        LXI SP,F001
        CALL DEBUG
        HLT

# MYREGISTERS EQU 2000H
DEBUG: PUSH H
        PUSH D
        PUSH B
        PUSH PSW
        LXI H,MYREGISTERS
        POP D
        CALL STORE    ;store AF
        POP D
        CALL STORE    ;store BC
        POP D
        CALL STORE    ;store DE
        POP D
        CALL STORE    ;store HL
        RIM
        MOV M,A       ;store I
        INX H
        XCHG          ;swap HL and DE
        LXI H,0000
        DAD SP        ;now HL <- SP
        XCHG          ;now HL-> save addr, DE<- SP
;problem, SP also has this function call stuff, so we need to remove it's info from DE
;(DE = DE + 2 [remember SP stack works reverse that's why +])
        INX D
        INX D
;now store DE
        CALL STORE
;now we need to store PC, which will be in Stack due to function call
        XCHG          ;swap HL and DE to store HL stuff in DE
        XTHL          ;get PC data from SP (stack mem has garbage now)
        XCHG          ;now HL has save addr, DE has data
        CALL STORE
        XCHG          ;PC data back in HL
        XTHL          ;PC data back in stack
        XCHG          ;save addr back in HL
        RET

;Procedure STORE
;stores data in DE to memory whose address is in HL
STORE: MOV M,D
        INX H
        MOV M,E
        INX H
        RET
```

Listing 9: assembly program to store register configuration

9.4 Experimentation

Memory Editor

Memory Range: 0000 ---- FFFF

Memory Address	Value
0036	42
0038	EB
0039	E3
003A	EB
003B	CD
003C	42
003E	EB
003F	E3
0040	EB
0041	C9
0042	72
0043	23
0044	73
0045	23
0046	C9
2000	10
2002	50
2003	92
2004	27
2005	94
2006	37
2007	92
2009	F0
200A	11
200C	12
EFF7	1E
EFF9	22
EFFB	26
EFFD	3E
EFFF	12

Figure 8: Register configuration after execution (look from address 2000)

9.5 Conclusion

We see that after program execution, all the predefined value of registers are stored in memory starting from address 2000.
Hence the program is working as expected.

10 Assignment 10

10.1 Objective

Implement a POST or power-on-self-test where each RAM location is tested for stuck-at-zero or stuck-at-one fault. In your case the function takes the start address of the RAM block and the block size in bytes. The function sets CY in case of any error (else it is set to 0); HL contains the faulty location and Acc contains 0 for stuck at zero fault and 1 for stuck at one fault. [Note: usually there wont be any error as your RAM is not faulty, so direct checking may not set CY flag]

10.2 Tool/Experimental setup considered

- Jubin's 8085 Simulator

10.3 Procedure

We iteratively go through STARTLOC to (STARTLOC + LEN) and check if any value is 1 or not, if it is, we exit from there setting the C flag.

10.4 Program

```
#STARTLOC EQU 6400
#LEN EQU 05
;let position 6403 have stuck at one fault
        LXI H,6403
        MVI M,01
;let we start out search from 6400 an search 05 addresses
        LXI H,STARTLOC
        MVI B,LEN
        CALL POST
        HLT

POST:   MOV A,M
        CPI 01
;A will be 0 or 1, if 1 we need to exit
;A zero -> A - 1 < 0, Cy =1
;A one -> A - 1 = 0, Z = 1, we exit here
        JZ ERR      ;if 1 found, we exit are set Cy 1 directly
        INX H
        DCR B
        JNZ POST
;we will reach here only if there is no 1 found, so we need to set Cy 0
        JC REVERT   ;if Cy 1 , we revert, else we direct return
        RET
;we reach here when 1 found

ERR:    STC         ;set Cy 1
        RET
;we reach here if no 1 found but Cy 1 due to previous CMP stuff

REVERT: CMC         ;complements Cy, if Cy 1, Cy becomes 0
        RET
```

Listing 10: assembly program for POST for stuck at 1 fault

10.5 Experimentation

Register	Value	7	6	5	4	3	2	1	0
Accumulator	01	0	0	0	0	0	0	0	1
Register B	02	0	0	0	0	0	0	1	0
Register C	00	0	0	0	0	0	0	0	0
Register D	00	0	0	0	0	0	0	0	0
Register E	00	0	0	0	0	0	0	0	0
Register H	64	0	1	1	0	0	1	0	0
Register L	03	0	0	0	0	0	0	1	1
Memory(M)	01	0	0	0	0	0	0	0	1

Resister	Value	S	Z	*	AC	*	P	*	CY
Flag Resister	55	0	1	0	1	0	1	0	1

Type	Value
Stack Pointer(SP)	0000
Memory Pointer (HL)	6403
Program Status Word(PSW)	0155
Program Counter(PC)	000D
Clock Cycle Counter	221
Instruction Counter	29

(a) result of the program

Register	Value	7	6	5	4	3	2	1	0
Accumulator	00	0	0	0	0	0	0	0	0
Register B	00	0	0	0	0	0	0	0	0
Register C	00	0	0	0	0	0	0	0	0
Register D	00	0	0	0	0	0	0	0	0
Register E	00	0	0	0	0	0	0	0	0
Register H	64	0	1	1	0	0	1	0	0
Register L	05	0	0	0	0	0	1	0	1
Memory(M)	00	0	0	0	0	0	0	0	0

Resister	Value	S	Z	*	AC	*	P	*	CY
Flag Resister	54	0	1	0	1	0	1	0	0

Type	Value
Stack Pointer(SP)	0000
Memory Pointer (HL)	6405
Program Status Word(PSW)	0054
Program Counter(PC)	0008
Clock Cycle Counter	266
Instruction Counter	37

(b) result of the program if we comment line 4 and 5

Figure 9: Result of Program Execution

10.6 Conclusion

We see that if there exist a 1 in some memory location, the program sets Cy flag and HL register to stuck address; when there is no 1 in whole array, C is unset
Hence the program is working as expected.

11 Assignment 11

11.1 Objective

Implement a binary search —- the function would take the start address and no. of elements in the array. If successful the function resets CY flag and the HL pair points to the element found else CY is set and the value in HL pair is irrelevant.

11.2 Tool/Experimental setup considered

- Jubin's 8085 Simulator

11.3 Program

```
# ORG 2000H
# ARR : DB 2,4,7,8,10
# ORG 0000H
# N EQU 5
# X EQU 7
            MVI C,N     ;Number of elements in the Array
            DCR C       ;Now we will be using C as the high pointer of Binary Search
            MVI B,00    ;It is the low pointer of Binary Search
            MVI D,X     ;Number we are looking for
            CALL SEARCH
            LXI H,ARR
            ADD L
            MOV L,A
            JNC NAD
            INR H
NAD:        MOV A,D
            CMP M
            JZ ND
            JC ND
            HLT
ND:         CMC         ;Complement the carry flag
            HLT         ;Halt

SEARCH:     MOV A,B
            CMP C       ;Checking if low < high
            RNC         ;Return if not carry
            ADD C       ;Now acc has B+C
            JNC NOCARRY
            CMC

NOCARRY:    RAR         ;Rotate Accumulator right by 1 bit (A = A >> 1) i.e A = (B+C)/2
            MOV E,A
            LXI H,ARR   ;H-L now has 2000h
            ADD L       ;Accumlator now has the offset to the mid point
            MOV L,A     ;H-L now points to the mid
            JNC NOPE
            INR H

NOPE:       MOV A,D     ;Accumulator now has the number we are looking for
            CMP M       ;Compare the number we are looking for with the mid point
            JC LEFT     ;If less than mid point, go left
            JZ ND       ;If equal, return
            MOV B,E     ;E had the previous mid point index
            INR B       ;Now looking for [M+1:]
            JMP SEARCH  ;Recall search

LEFT:       MOV C,E
            DCR C
            JNZ SEARCH
            JZ ND
            RET
;CMP M
;If A less than (R/M), the CY flag is set and Zero flag is reset.
;If A equals to (R/M), the Zero flag is set and CY flag is reset.
;If A greater than (R/M), the CY and Zero flag are reset.
```

Listing 11: assembly program for the implementation of Binary Search

11.4 Experimentation

Register	Value	7	6	5	4	3	2	1	0
Accumulator	07	0	0	0	0	0	1	1	1
Register B	00	0	0	0	0	0	0	0	0
Register C	04	0	0	0	0	0	1	0	0
Register D	07	0	0	0	0	0	1	1	1
Register E	02	0	0	0	0	0	0	1	0
Register H	20	0	0	1	0	0	0	0	0
Register L	02	0	0	0	0	0	0	1	0
Memory(M)	07	0	0	0	0	0	1	1	1

Resister	Value	S	Z	*	AC	*	P	*	CY
Flag Resister	55	0	1	0	1	0	1	0	1

Type	Value
Stack Pointer(SP)	FFFE
Memory Pointer (HL)	2002
Program Status Word(PSW)	0755
Program Counter(PC)	001D
Clock Cycle Counter	148
Instruction Counter	22

(a) result of the program when X is in array

Register	Value	7	6	5	4	3	2	1	0
Accumulator	1A	0	0	0	1	1	0	1	0
Register B	04	0	0	0	0	0	1	0	0
Register C	04	0	0	0	0	0	1	0	0
Register D	1A	0	0	0	1	1	0	1	0
Register E	03	0	0	0	0	0	0	1	1
Register H	20	0	0	1	0	0	0	0	0
Register L	04	0	0	0	0	0	1	0	0
Memory(M)	10	0	0	0	1	0	0	0	0

Resister	Value	S	Z	*	AC	*	P	*	CY
Flag Resister	04	0	0	0	0	0	1	0	0

Type	Value
Stack Pointer(SP)	0000
Memory Pointer (HL)	2004
Program Status Word(PSW)	1A04
Program Counter(PC)	001B
Clock Cycle Counter	682
Instruction Counter	106

(b) result of the program when X is not in array

Figure 10: Result of Program Execution

11.5 Conclusion

We see that when X is in array, Cy flag is set and HL points to the data's address in array; but then X is not present, Cy flag is not set.
Hence the program is working as expected.

12 Assignment 12

12.1 Objective

Suppose that you are reading a bit of an input port (say, PORT 0) to which the output of a function generator, producing a rectangular wave, is connected. Measure the ON of this repetitive rectangular waveform time in terms of millisecond.

12.2 Procedure

1. We first wait in an infinite loop, waiting for the signal to turn 0 (in case we started program when signal was already 1).
2. Then when the signal turns 0, the actual count procedure starts, we again keep waiting in an infinite loop, waiting for the signal to go 1 again.
3. When the signal goes 1 now, we start counting the amount of time the signal is 1 in another infinite loop.
    - The two register pair is too small to count the number of machine cycles it took for the signal to go to 0.
    - so we added a delay (1 ms) after every count, now the count signifies the amount of (1ms) delays it took for signal to go to 0.
4. now the signal is back to 0, and we have the amount of delays it took for signal to go to 0, we write that data to SAVELOC.

12.3 Program

```
# ORG 1000H
;delay -> 1 ms
DELAY:      MVI D,DE
DELAYLoop:  DCR D
            JNZ DELAYLoop
            RET
# ORG 0000
# PORTX EQU 10
INPUT:      IN PORTX        ;input data goes to A
            RAR
            JC INPUT        ;if already high read again
;now input port data 0 (actual low->high timing cont begins now)
L1:         IN PORTX
            RAR
            JNC L1          ;if LOW read again
;now input HIGH -> start counting
            LXI B,0000
COUNT:      INX B           ;6T
            CALL DELAY
            IN PORTX        ;10T
            RAR             ;4T
            JC COUNT        ;10T -> jump till input is 1
;now input is 0 we save counter
;time taken = BC*(30*1/3 + 1ms) = BC*11ms +- 2 ms
# SAVELOC   EQU 5000H
            LXI H,SAVELOC
            MOV M,B
            INX H
            MOV M,C
            HLT
```

Listing 12: assembly program to calculate time period of rectangular wave

12.4 Conclusion

We created a program which will calculate the time period of a rectangular wave as follows

$$T = BC * (30 * 1/3 + 1ms) = BC * 11ms \pm 2ms$$

13 Assignment 13

13.1 Objective

- Using auto vectored input RST 7.5 prepare a scheme to count the number of key-press done at this interrupting input.
- The main routine after initialization of the interrupt mechanism waits in an infinite loop waiting for the key-press.
- On a key-press (that simulates as if you have excited the RST 7.5 input) it increases a counter at a predefined memory location (used to hold the count value).

13.2 Tool/Experimental setup considered

- Jubin's 8085 Simulator

13.3 Procedure

1. We use an RST 7.5 interrupt line to call a procedure every time we get an interrupt.
2. The Problem is that the computer is so fast that it can take in multiple interrupt in a single button press(click).
3. To solve this, we disable the interrupt mechanism inside the Interrupt Service Routine, and also add a small delay to make the computer "wait" for the key-press to end.

13.4 Program

```
# ORG 003CH
            JMP COUNT
# ORG 0000H
            MVI A,0B    ;00001011B
            SIM
            EI
            LXI D,0000

INFLOOP:    JMP INFLOOP
            HLT

COUNT:      DI
            INX D
            CALL DEL125
            EI
            RET

DEL125:     LXI B,3F93
DEL125Loop: DCX B
            MOV A,B
            ORA C
            JNZ DEL125Loop
            RET
```

Listing 13: assembly program to find the number of key presses during the program execution

13.5 Experimentation

Register	Value	7	6	5	4	3	2	1	0
Accumulator	2F	0	0	1	0	1	1	1	1
Register B	2D	0	0	1	0	1	1	0	1
Register C	0B	0	0	0	0	1	0	1	1
Register D	00	0	0	0	0	0	0	0	0
Register E	04	0	0	0	0	0	1	0	0
Register H	00	0	0	0	0	0	0	0	0
Register L	00	0	0	0	0	0	0	0	0
Memory(M)	4F	0	1	0	0	1	1	1	1

Resister	Value	S	Z	*	AC	*	P	*	CY
Flag Resister	00	0	0	0	0	0	0	0	0

Type	Value
Stack Pointer(SP)	0004
Memory Pointer (HL)	0000
Program Status Word(PSW)	2F00
Program Counter(PC)	0054
Clock Cycle Counter	1286871
Instruction Counter	214545

SOD	SID	INTR	TRAP	R7.5	R6.5	R5.5
0	0	0	0	0	0	0

For SIM instruction

SOD	SDE	*	R7.5	MSE	M7.5	M6.5	M5.5
0	0	0	0	0	0	0	0

For RIM instruction

SID	I7.5	I6.5	I5.5	IE	M7.5	M6.5	M5.5
0	0	0	0	0	0	0	0

No. Converter Tool :

Hexadecimal	Decimal	Binary
0	0	0

(a) register config for 4 interrupt clicks

Register	Value	7	6	5	4	3	2	1	0
Accumulator	2E	0	0	1	0	1	1	1	0
Register B	2E	0	0	1	0	1	1	1	0
Register C	40	0	1	0	0	0	0	0	0
Register D	00	0	0	0	0	0	0	0	0
Register E	0C	0	0	0	0	1	1	0	0
Register H	00	0	0	0	0	0	0	0	0
Register L	00	0	0	0	0	0	0	0	0
Memory(M)	4F	0	1	0	0	1	1	1	1

Resister	Value	S	Z	*	AC	*	P	*	CY
Flag Resister	04	0	0	0	0	0	1	0	0

Type	Value
Stack Pointer(SP)	0014
Memory Pointer (HL)	0000
Program Status Word(PSW)	2E04
Program Counter(PC)	0056
Clock Cycle Counter	4406713
Instruction Counter	734651

SOD	SID	INTR	TRAP	R7.5	R6.5	R5.5
0	0	0	0	0	0	0

For SIM instruction

SOD	SDE	*	R7.5	MSE	M7.5	M6.5	M5.5
0	0	0	0	0	0	0	0

For RIM instruction

SID	I7.5	I6.5	I5.5	IE	M7.5	M6.5	M5.5
0	0	0	0	0	0	0	0

No. Converter Tool :

Hexadecimal	Decimal	Binary
0	0	0

(b) register config for 12 interrupt clicks

Figure 11: Result of Program Execution (look at DE register pair)

13.6 Conclusion

We see that the value in DE register pair after program stop matches with the number of clicks made for interrupts.
Hence the program is working as expected.

14 Assignment 14

14.1 Objective

Draw the hardware for implementing RST 7.5 based key-press counting – don't forget to implement h/w based key debouncing problem.

14.2 Procedure

We previously used software method (using delay) to mitigate debouncing, now we will look into the method to reducing debouncing hardware wise.

One of the simplest hardware-based switch debounce solutions employs a resistor-capacitor (RC) network in conjunction with a switch (i.e our key).

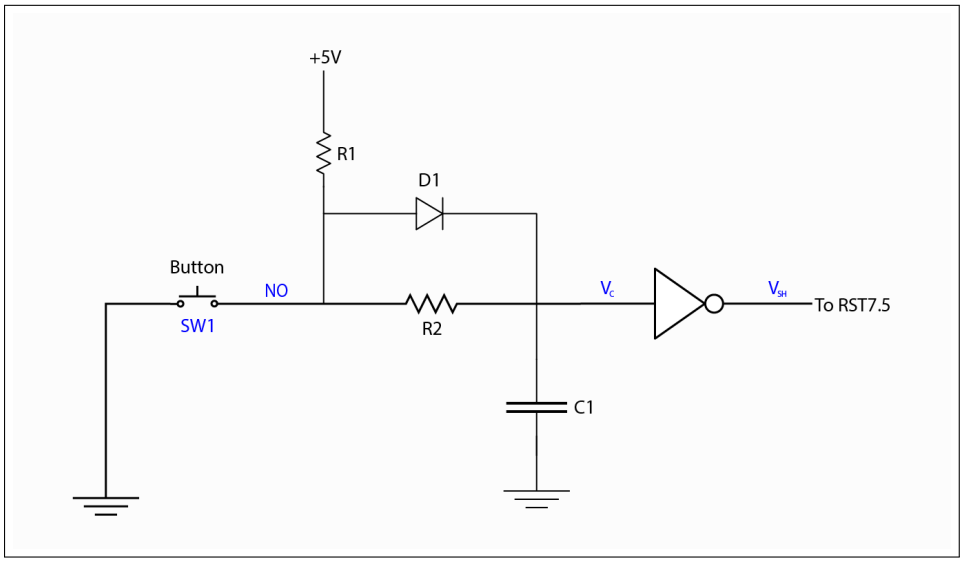
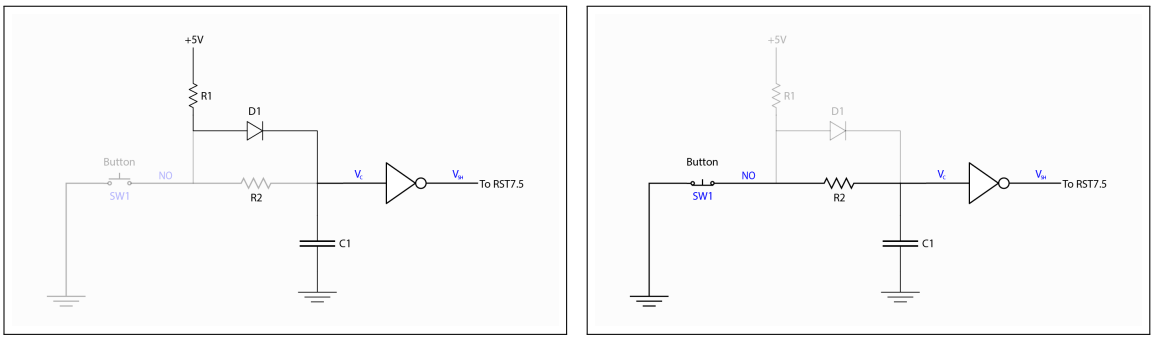


Figure 12: Hardware based Switch Debouncing

When the key is not pressed, there is a way for the power to go from +ve to ground via R1, D1 and C1, which keeps V_c as high, and V_{sh} as low as shown in 13a.

When the key is pressed, the circuit is shorted due to the new circuit and V_c becomes low, making V_{sh} as high as shown in 13b

The capacitor is used to remove the false signals due to debouncing and providing a smoother gradient towards On and Off as shown in 14



(a) Circuit when key is not pressed

(b) Circuit when key is pressed

Figure 13: Hardware based Switch Debouncing

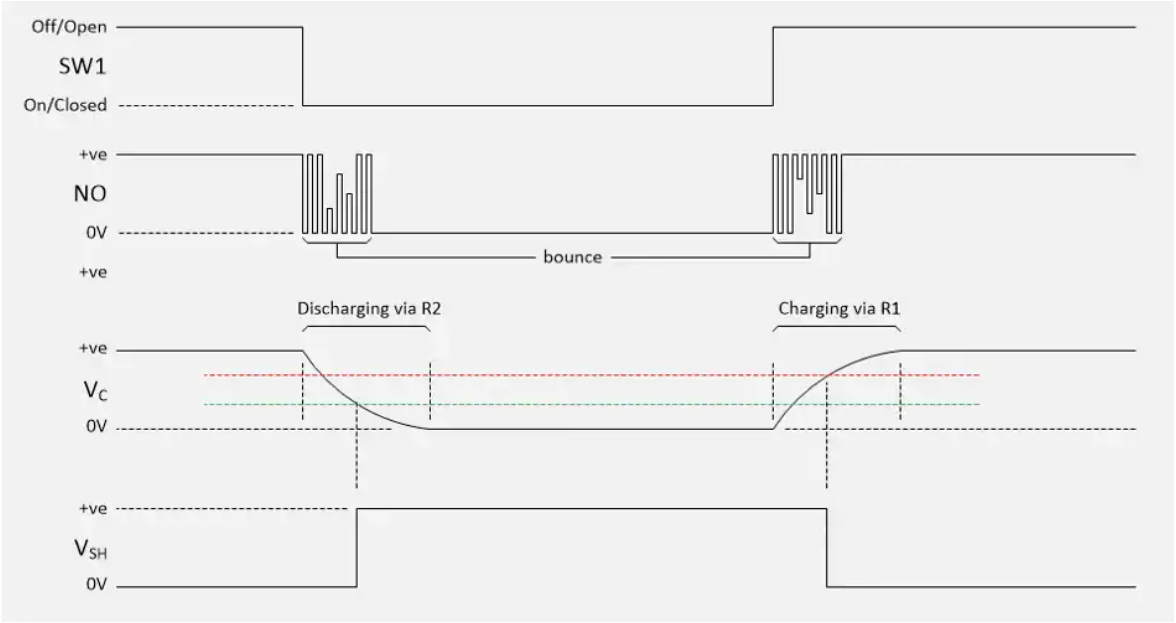


Figure 14: Hardware based Switch Debouncing (Source)

14.3 Conclusion

We discussed hardware based debouncing mechanism to reduce errors while taking input.