

Slime Mold Simulation Tool

Motivation

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graph TD; 0((0)) --- 1((1)); 1 --- 4((4)); 4 --- 5((5)); 5 --- 9((9)); 9 --- 10((10)); 10 --- 11((11)); 11 --- 12((12)); 12 --- 16((16)); 16 --- 17((17)); 17 --- 18((18)); 10 --- 13((13)); 13 --- 14((14)); 14 --- 15((15)); 2((2)); 3((3)); 6((6)); 7((7)); 8((8)); 16((16));
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Figure 3: After 30 time steps, the system converges to the final solution. Edges that do not lie along the shortest path are not visible, and the remaining edges show the calculated shortest path between food sources.

Tool Details

The slime mold algorithm implemented is shown below, presented in [1].

Algorithm 1 Improved *Physarum Ploycephalum* algorithm

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Let  $N$  be the number of nodes in the network;
Let  $L$  be the  $N \times N$  matrix whose  $(i, j)$  entry  $L_{ij}$  is the
length between node  $i$  and  $j$ ;
Let  $s$  and  $e$  be the start and end nodes;
Let  $D_{ij} \in (0, 1]$  for all  $i, j = 1, 2, \dots, N$ ;
Let  $Q_{ij} := 0$  for all  $i, j = 1, 2, \dots, N$ ;
Let  $p_i := 0$  for all  $i = 1, 2, \dots, N$ ;
Set iteration counter  $t = 0$ ;
while termination criteria not met do
    Let  $p_e := 0$  (pressure at ending node  $e$ ) ;
    Calculate the pressure of every node in the network
    solving (5);
    Let  $Q_{ij} := D_{ij}(p_i - p_j)/L_{ij}$  using (1);
    Let  $D_{ij} := \frac{1}{2} \left( \frac{Q_{ij}(p_i - p_j)}{L_{ij}(p_s - p_e)} + D_{ij} \right)$  using (6);
    Let  $t = t + 1$ ;
end

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Equation (5): Fixes the net flow of fluid of all nodes to 0, except for source and sink nodes. The source node will have a net flow of +1 out, while the sink will have a net flow of -1 out (+1 entering).

$$\sum_i \frac{D_{ij}}{L_{ij}} (p_i - p_j) = \begin{cases} +1 & \text{for } j = 1, \\ -1 & \text{for } j = 2, \\ 0 & \text{otherwise.} \end{cases}$$

In the algorithm, D_{ij} corresponds to the diameter of the edge between nodes i and j , L_{ij} is the length of the edge, and Q_{ij} is the fluid flow between the nodes, and p_i is the fluid pressure at node i . Note that all values in D are initialized to 0, while those in L are initialized to 1. Actual values for L are set based on real distances between nodes of the maze.

Solving for node pressures

To solve for the pressure at every node, we use matrix operations. Using Kirchhoff's Current Law, we can create a system of equations. For each node given by index i , in a system with N total nodes the following holds:

$$\sum_{\forall j \in N} ((p_i - p_j) * C_{i,j}) = \begin{cases} 1 & \text{if } i : \text{source node} \\ -1 & \text{if } i : \text{sink node} \\ 0 & \text{otherwise} \end{cases}$$

This equation simply says the sum of all fluid flow (pressure difference * conductance) out of a node will be 0 for all, except for source and sink nodes. The sum can be expanded to the following equation:

$$EQ1: p_i(C_{0,1} + C_{0,2} + \dots + C_{0,N-1}) - p_1 C_{0,1} - p_2 C_{0,2} \dots p_{N-1} C_{0,N-1}$$

Let us represent the sum

$$(C_{0,1} + C_{0,2} + \dots + C_{0,N-1}) = \sum_{\forall j \in N} C_{0,j} = S_i$$

This is the sum of conductances between node i and all other nodes. Note that

$$\forall i \in N: C_{i,i} = 0$$

Then EQ1 can be represented by the inner product:

$$< \{S_i, -C_{0,1}, -C_{0,2}, \dots, -C_{0,N-1}\}, \{p_0, p_1, \dots, p_{N-1}\} >$$

This inner product will equal the sum of flows at node i . The first vector in this inner product can be constructed for every single node in the system in a similar fashion, leading to an $N \times N$ matrix, **A**. Then the entire system can be represented as:

$$A * p = flow$$

Where **p** is the $N \times 1$ vector of pressures we need to solve for, and **flow** is the $N \times 1$ vector of fluid flow sums for each node. Since we already know **flow** and **A**, we can simply solve for **p** with:

$$p = A^{-1} * flow$$

The rest of the algorithm is straightforward and allows us to update the diameters of all edges for each time step.

Future Work

Although the tool works well for paths between just two nodes, it would be interesting to extend it for multiple sources and sinks. Real slime mold is able to solve more complex problems like minimum spanning tree. Although I had some success implementing this using this algorithm, I found that the math does not converge under some circumstances, and decided to leave it out of the final version.

It would also be worthwhile to simulate slime mold's ability to actually explore a space. Based on some literature it seems slime mold mostly explores its surrounding space randomly, then decides to optimize once it finds food sources. [4] contains some interesting findings about how slime mold uses its own slime deposits as external memory, which may be of use for a simulator.

References

- [1] Brabazon, Anthony & McGarraghy, Seán. (2020). Slime mould foraging: An inspiration for algorithmic design. *International Journal of Innovative Computing and Applications*. 11. 30. 10.1504/IJICA.2020.105316.
- [2] Nakagaki, T., Yamada, H. and Toth, A. (2000) 'Maze-solving by an amoeboid organism', *Nature*, Vol. 407, No. 6803, pp.470.
- [3] Mechanism of signal propagation in *P. polycephalum*
Karen Alim, Natalie Andrew, Anne Pringle, Michael P. Brenner *Proceedings of the National Academy of Sciences* May 2017, 114 (20) 5136-5141; DOI: 10.1073/pnas.1618114114
- [4] Slime mold uses an externalized spatial "memory" Chris
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