

# BCH Codes

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# BCH Codes

- Discovered by Hocquenghem in 1959 and independently by Bose and Chaudhari in 1960
- Cyclic structure proved by Peterson in 1960
- Decoding algorithms proposed/refined by Peterson, Gorenstein and Zierler, Chien, Forney, Berlekamp, Massey. . .
- We will discuss a subclass of BCH codes — binary primitive BCH codes

# Binary Primitive BCH Codes

For positive integers  $m \geq 3$  and  $t < 2^{m-1}$ , there exists an  $(n, k)$  BCH code with parameters

- $n = 2^m - 1$
- $n - k \leq mt$
- $d_{\min} \geq 2t + 1$

## Definition

Let  $\alpha$  be a primitive element in  $F_{2^m}$ . The generator polynomial  $g(x)$  of the  $t$ -error-correcting BCH code of length  $2^m - 1$  is the least degree polynomial in  $\mathbb{F}_2[x]$  that has

$$\alpha, \alpha^2, \alpha^3, \dots, \alpha^{2t}$$

as its roots.

Let  $\varphi_i(x)$  be the minimal polynomial of  $\alpha^i$ . Then  $g(x)$  is the LCM of  $\varphi_1(x), \varphi_2(x), \dots, \varphi_{2t}(x)$

## Degree of Generator Polynomial

$$g(x) = \text{LCM} \{ \varphi_1(x), \varphi_2(x), \varphi_3(x), \dots, \varphi_{2t}(x) \}$$

- If  $i$  is an even integer, then  $i = i'2^a$  where  $i'$  is odd
- Since  $\alpha^i = (\alpha^{i'})^{2^a}$ ,  $\alpha^i$  and  $\alpha^{i'}$  have the same minimal polynomial
- Every even power of  $\alpha$  has the same minimal polynomial as some previous odd power of  $\alpha$

$$g(x) = \text{LCM} \{ \varphi_1(x), \varphi_3(x), \varphi_5(x), \dots, \varphi_{2t-1}(x) \}$$

# BCH Codes of Length 15

- Let  $\alpha$  be a primitive element of  $F_{16}$  and a root of  $x^4 + x + 1$

Power	Polynomial	Tuple
0	0	(0 0 0 0)
1	1	(1 0 0 0)
$\alpha$	$\alpha$	(0 1 0 0)
$\alpha^2$	$\alpha^2$	(0 0 1 0)
$\alpha^3$	$\alpha^3$	(0 0 0 1)
$\alpha^4$	$1 + \alpha$	(1 1 0 0)
$\alpha^5$	$\alpha + \alpha^2$	(0 1 1 0)
$\alpha^6$	$\alpha^2 + \alpha^3$	(0 0 1 1)
$\alpha^7$	$1 + \alpha + \alpha^3$	(1 1 0 1)
$\alpha^8$	$1 + \alpha^2$	(1 0 1 0)
$\alpha^9$	$\alpha + \alpha^3$	(0 1 0 1)
$\alpha^{10}$	$1 + \alpha + \alpha^2$	(1 1 1 0)
$\alpha^{11}$	$\alpha + \alpha^2 + \alpha^3$	(0 1 1 1)
$\alpha^{12}$	$1 + \alpha + \alpha^2 + \alpha^3$	(1 1 1 1)
$\alpha^{13}$	$1 + \alpha^2 + \alpha^3$	(1 0 1 1)
$\alpha^{14}$	$1 + \alpha^3$	(1 0 0 1)

## BCH Codes of Length 15

- Let  $\alpha$  be a primitive element of  $F_{16}$  and a root of  $x^4 + x + 1$
- The minimal polynomials of  $F_{16}$  are  $\mathbb{F}_2[x]$  factors of  $x^{16} + x$   
$$x^{16} + x = x(x+1)(x^2+x+1)(x^4+x+1)(x^4+x^3+1)(x^4+x^3+x^2+x+1)$$
- A single error correcting BCH code of length 15 has generator polynomial  $g(x) = \varphi_1(x) = x^4 + x + 1$
- A double error correcting BCH code of length 15 has generator polynomial

$$g(x) = \text{LCM}\{\varphi_1(x), \varphi_3(x)\} = (x^4 + x + 1)(x^4 + x^3 + x^2 + x + 1)$$

- The maximum value of  $t$  for a BCH code of length 15 is 7
- What is the generator polynomial for correcting seven errors?

Questions? Takeaways?