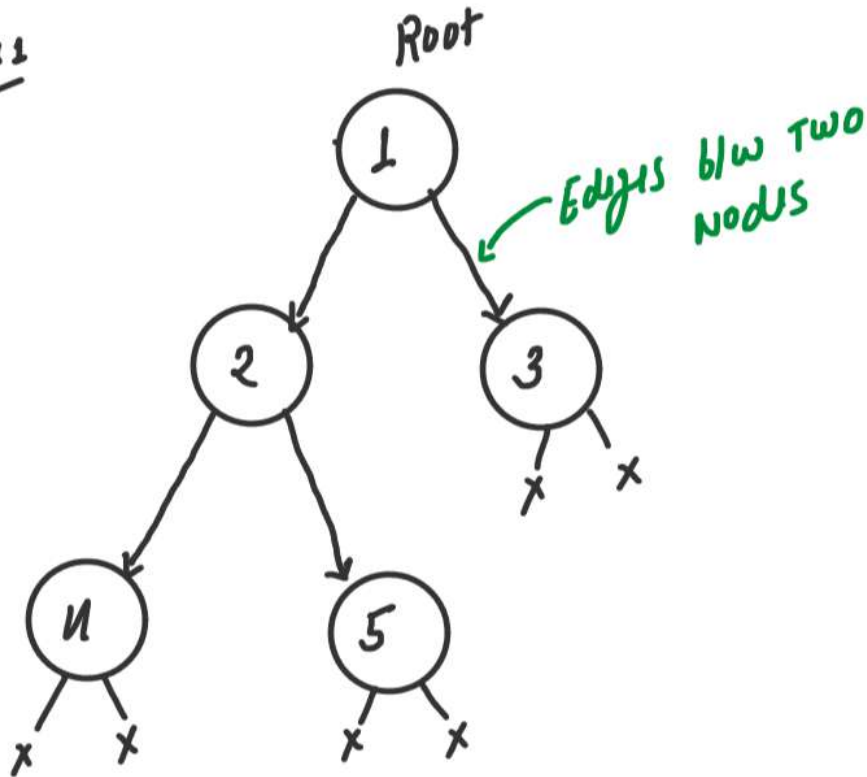
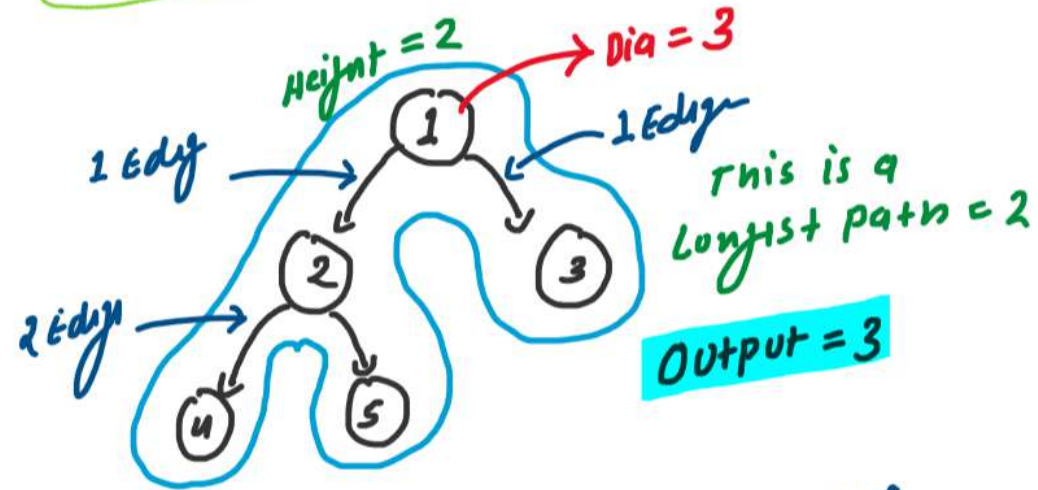


8. Diameter of Binary Tree (Leetcode-543)

Ex 1



Explanation what is Diameter?



$$\begin{aligned}
 \text{Diameter} &= \text{LH Edges} + \text{RH Edges} \\
 &= 2 + 1 \\
 &= 3
 \end{aligned}$$

Ex 1

④  $Dia = 2 + 1 = 3$   
 $LH = 2$  Root  $RH = 1$  *option 3*

②  $Dia = 1 + 1 = 2$   
 $LH = 1$   $RH = 1$

$Dia = 0 + 0 = 0$   
 $LH = 0$   $RH = 0$  ⑤

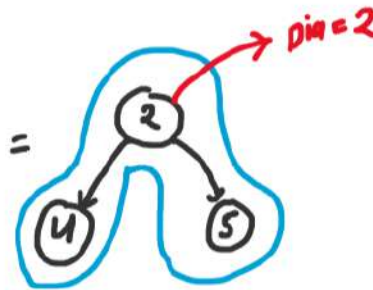
①  $Dia = 0 + 0 = 0$   
 $LH = 0$   $RH = 0$

$Dia = 0 + 0 = 0$   
 $LH = 0$   $RH = 0$  ③

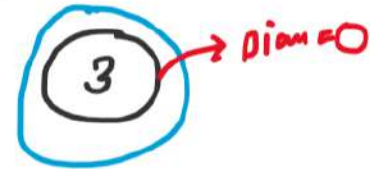
*option 1*

## BRUTE FORCE APPROACH

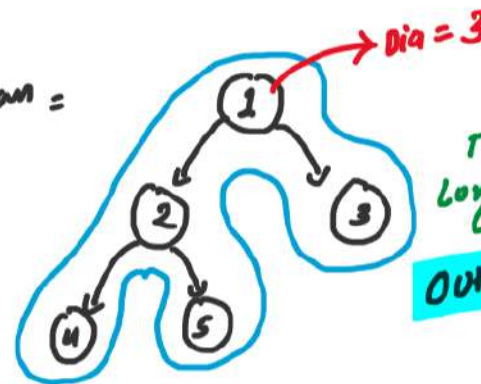
1st option =



2nd option =



3rd option =



This is a longest path

Output = 3

Diameter = max (2, 0, 3)

=> 3  
Ans

```

//  BRUTE FORCE APPROACH (Overhead of the recursive call)
class Solution {
public:
    int height(TreeNode* root){
        if(root == NULL) return 0;

        int LH = height(root->left);
        int RH = height(root->right);
        int finalHeight = max(LH, RH) + 1;
        return finalHeight;
    }

    int diameterOfBinaryTree(TreeNode* root) {
        //base case
        if(root == NULL) {
            return 0;
        }
        int option1 = diameterOfBinaryTree(root->left);
        int option2 = diameterOfBinaryTree(root->right);
        int option3 = height(root->left) + height(root->right);
        int diameter = max(option1, max(option2, option3));
        return diameter;
    }
};

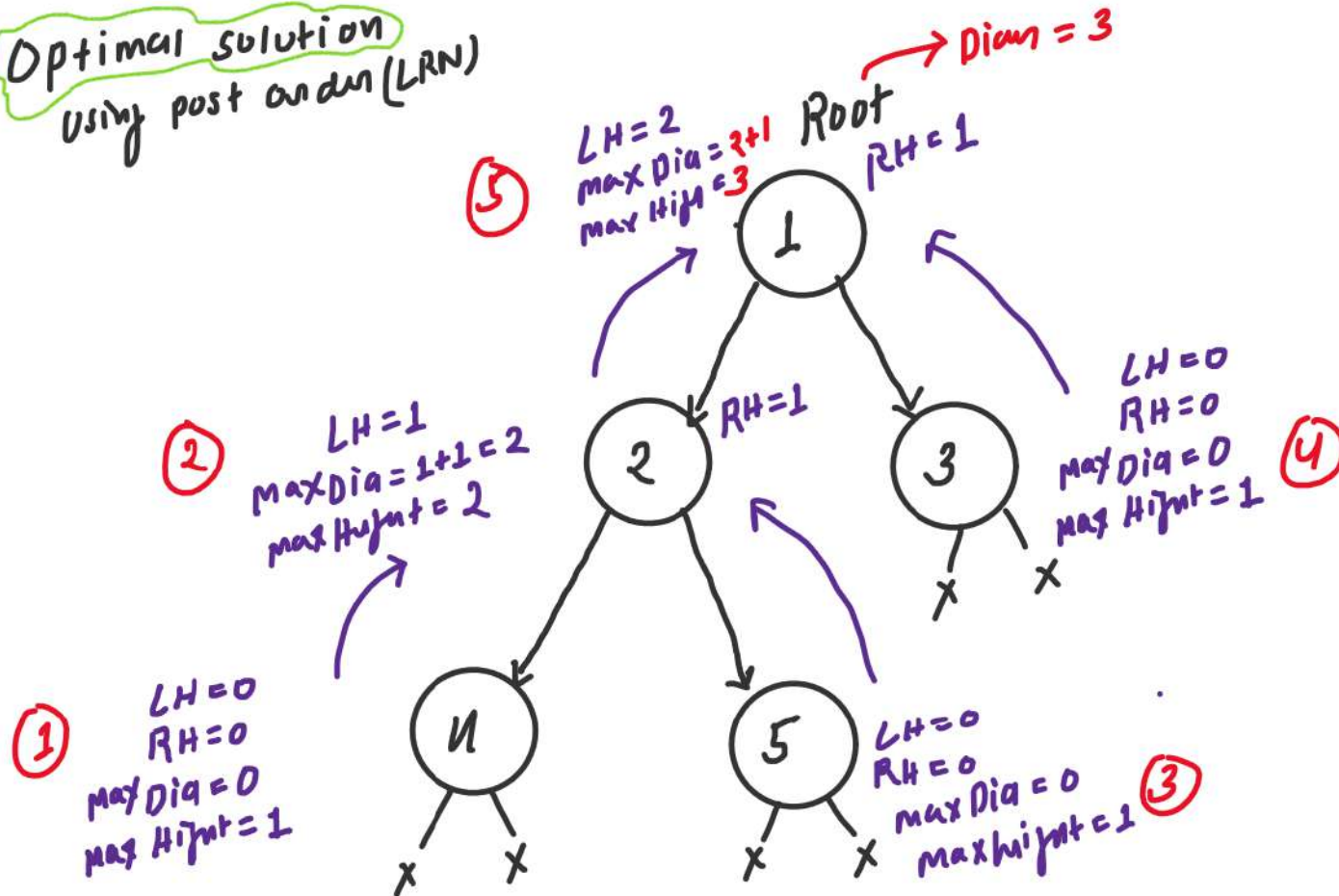
```

**Time complexity:**  $O(N^2)$ , Where N number of nodes

**Why  $O(N^2)$ ?** --> Height function is called by right and left subtree separately for each one node.

**Space complexity:**  $O(N)$  due to the recursive call stack

Optimal solution  
using post order (LRN)



$$\text{Diameter} = \max(\text{Dia}, (\text{LH} + \text{RH}))$$

$$\text{Height} = \max(\text{LH}, \text{RH}) + 1$$

Output

Diameter = 3

post order  
4 → 5 → 3 → 1

```

// ✅ OPTIMAL APPROACH (No overhead of the recursive call)
class Solution {
public:
    int height(TreeNode* root, int &diameter){
        if(root == NULL) return 0;

        int LH = height(root->left, diameter);
        int RH = height(root->right, diameter);
        int maxHeight = max(LH, RH) + 1;
        // Update diameter with longest path of tree
        diameter = max(diameter, (LH+RH));
        return maxHeight;
    }

    int diameterOfBinaryTree(TreeNode* root) {
        //base case
        if(root == NULL) {
            return 0;
        }
        int diameter = 0;
        height(root, diameter);
        return diameter;
    }
};

```

**Time complexity:**  $O(N^2)$ , Where N number of nodes

**Why  $O(N^2)$ ?** --> Height function is not called by right and left subtree separately for each one node.

**Space complexity:**  $O(H)$ , where H is the height of the binary tree