

HW 05: Book Allocation Problem**Problem Statement**

You are given an array 'pages' of integer numbers.

In this array, the 'pages[i]' represents the number of pages in the 'i-th' book.

There are 'm' number of students, and the task is to allocate all the books to the students.

Allocate books in a way such that:

- Each student gets at least one book.

- Each book should be allocated to a student.

- Book allocation should be in a contiguous manner.

→ Monotonic ORDER 1 2 3 ... 10 ... ∞

You have to allocate the books to 'm' students such that the maximum number of pages assigned to a student is minimum.

Example 01:

Input: pages[] = { 10, 20, 30, 40 }, n=4 and m = 2

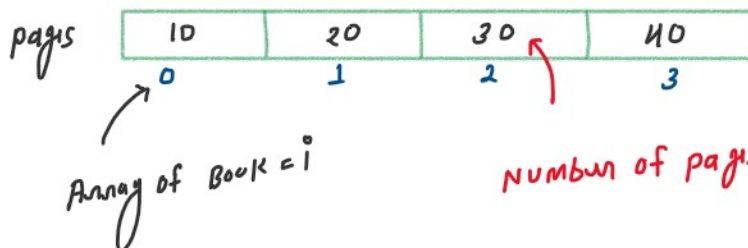
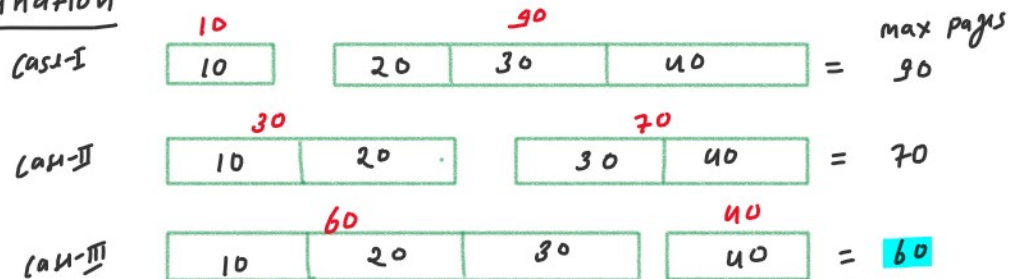
Output: 60

Number of books(n) = 4

Number of students(m) = 2

Number of pages in a book(i) = pages[i]

The minimum of the maximum number of pages assigned = $\min\{90, 70, 60\} = 60$

**Explanation**

Output

{ Minimum pages of maximum pages = 60 }

OPTIMAL APPROACH: Define search space and predicate function

Step 01: find total sum of array to create search space's end point (Maximum number of pages)

Step 02: apply binary search on search space BinarySearch()

Step 03: create predicate function isPossibleSolution()

Highest

pages

10	20	30	40
0	1	2	3

STEP:01

SUM = 100

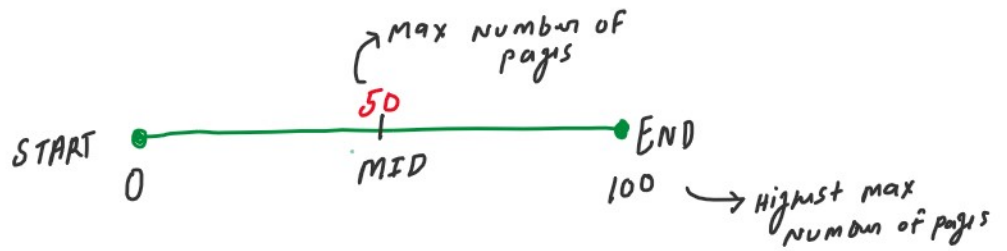
10 + 20 + 30 + 40 = 100

→ Max number of

STEP:02

Iter:01
Sum=100
Start=0
End=100
Mid=50
M=2
Ans=-1

$$\text{Mid} = \frac{0+100}{2} = 50$$



① $\text{mid} \leq \text{sum}$
Find minimum numbers of pages of Max Number of pages till $\text{Case} \leq M$

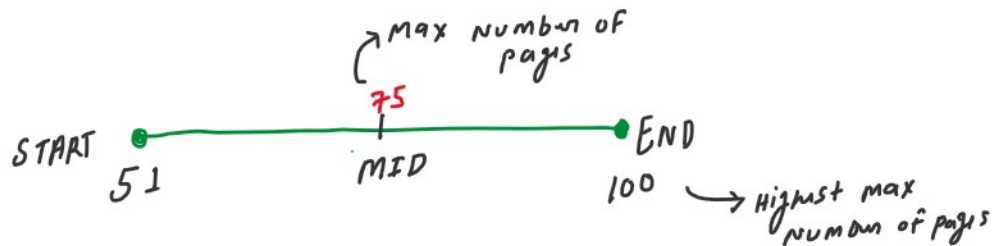
STEP:3 {
Case | $\text{Psum} \leq \text{mid}$
I $\rightarrow 10 + 20 + 30 = 30 \checkmark$
II $\rightarrow 30 + 40 = 30 \checkmark$
III \rightarrow Ruk ja... $\text{Case} \leq M \times$

No possible
 $\text{Start} = \text{mid} + 1$

STEP:02

Iter:02
Sum=100
Start=51
End=100
Mid=75
M=2
Ans=-1

$$\text{Mid} = \frac{51+100}{2} = 75$$



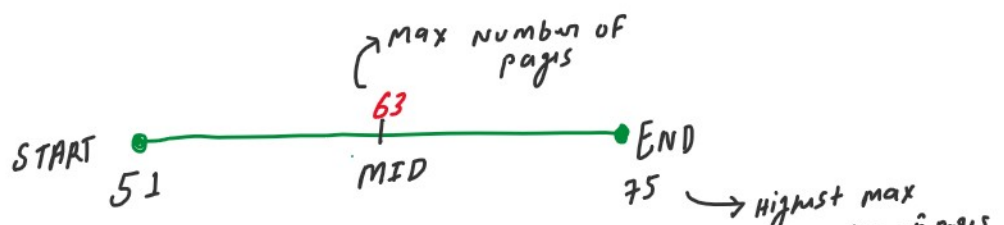
① $\text{mid} \leq \text{sum}$
Find minimum numbers of pages of Max Number of pages till $\text{Case} \leq M$

Case | $\text{Psum} \leq \text{mid}$
I $\rightarrow 10 + 20 + 30 + 40 = 60$
II $\rightarrow 40 = 40$

Possible sol.ⁿ
 $\text{Ans} = \text{mid}$
 $= 75$
 $\text{End} = \text{mid} - 1$

STEP:02

Iter:03



Iter: 03

Sum = 100

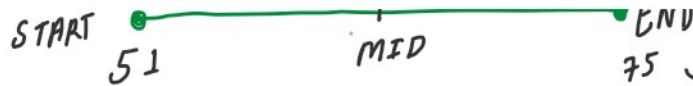
Start = 51

End = 74

mid = 63

M = 2

Ans = 75



$$\text{mid} = \frac{51 + 75}{2} = 63$$

① mid <= sum

Find minimum numbers of pages of max number of pages till $\text{Case} \leq M$

Case	$\text{psum} \leq \text{mid}$	
I	$10 + 20 + 30 + 40$	= 60
II	40	= 40

Possible sol.ⁿ
ANS = mid
= 63
End = mid - 1

STEP: 02

Iter: 04

Sum = 100

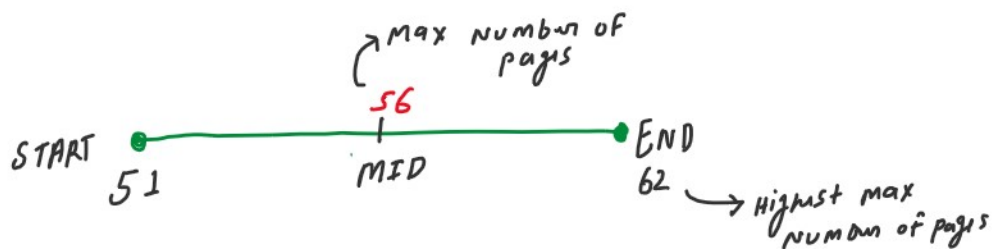
Start = 51

End = 62

mid = 56

M = 2

Ans = 63



$$\text{mid} = \frac{51 + 62}{2} = 56$$

① mid <= sum

Find minimum numbers of pages of max number of pages till $\text{Case} \leq M$

Case	$\text{psum} \leq \text{mid}$	
I	$10 + 20 + 30$	= 30
II	$30 + 40$	= 30
III	END	

NO Possible sol.ⁿ
Start = mid + 1

STEP: 02

Iter: 05

Sum = 100

Start = 57

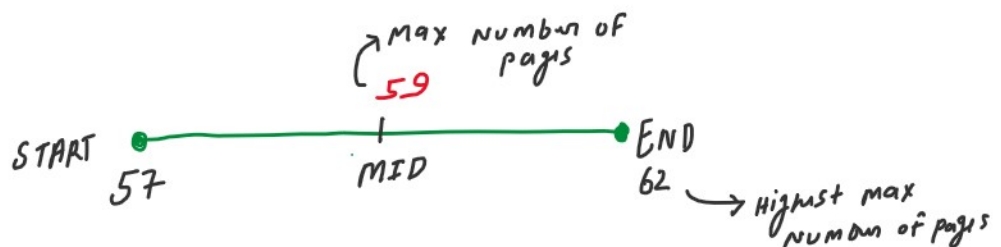
End = 62

mid = 59

$$\text{mid} = \frac{57 + 62}{2} = 59$$

① mid <= sum

Find minimum numbers of pages of



End = 62
 mid = 59
 M = 2
 Ans = 63

$$\text{mid} = \frac{60 + 62}{2} = 59$$

① $\text{mid} < \text{sum}$
 Find minimum numbers of pages of max number of pages till $\text{case} < M$

Case $\text{psum} < \text{mid}$
 I $\rightarrow 10 + 20 + 30 = 36$
 II $\rightarrow 30 + 40 = 30$
 III $\rightarrow \text{END}$
 NO possible sol.ⁿ
 $\text{start} = \text{mid} + 1$

STEP: 02

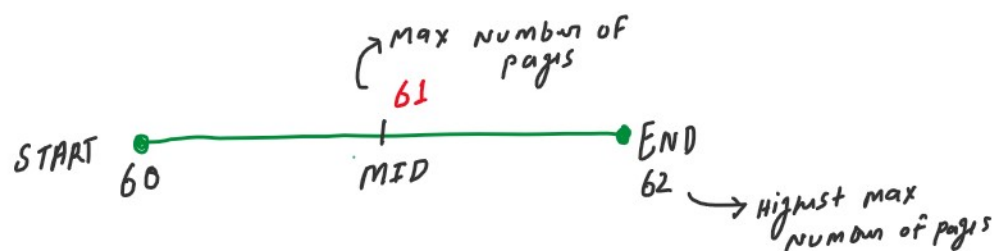
Iter = 6
 Sum = 100
 Start = 60
 End = 62
 mid = 61
 M = 2
 Ans = 63

$$\text{mid} = \frac{60 + 62}{2} = 61$$

① $\text{mid} < \text{sum}$
 Find minimum numbers of pages of max number of pages till $\text{case} < M$

Case $\text{psum} < \text{mid}$
 I $\rightarrow 10 + 20 + 30 = 60$
 II $\rightarrow 40 = 40$

possible sol.ⁿ
 $\text{Ans} = \text{mid} = 61$
 $\text{End} = \text{mid} - 1$

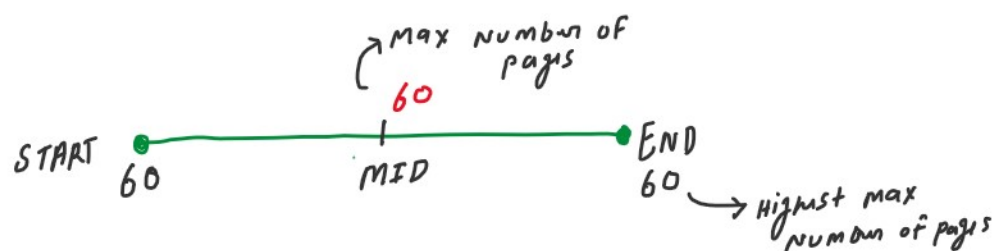


STEP: 02

Iter = 7
 Sum = 100
 Start = 60
 End = 60
 mid = 60
 M = 2
 Ans = 61

$$\text{mid} = \frac{60 + 60}{2} = 60$$

① $\text{mid} < \text{sum}$
 Find minimum numbers of pages of max number of pages till $\text{case} < M$



M = 2
Ans = 61

max number of pages till $Case \leq M$

Case $psum \leq mid$
I $\rightarrow 10 + 20 + 30 = 60$
II $\rightarrow 40 = 40$

Possible sol.ⁿ
 \hookrightarrow Ans = mid
 = 61
End = mid - 1

Iter: 8

Start = 60

End = 59

Ans = 60

START END
60 59

Ruk joo... (Start > END)
Final
Output

Minimum number of maximum number of pages
 $\Rightarrow 60$

```
// HW 05: Book Allocation Problem(GFG and Code studio)
#include <bits/stdc++.h>
bool isPossibleSolution(vector<int> arr, int n, int m, int mid){
    int cases = 1;
    int pageSum = 0;

    int i = 0;
    while(i < n){
        if(pageSum + arr[i] <= mid){
            pageSum += arr[i];
        }
        else{
            cases++;
            if(cases > m || arr[i] > mid){
                return false;
            }
            // reset pageSum
            pageSum = arr[i];
        }
        i++;
    }
    return true;
}
```

PREDICATE
FUNCTION
To 0 \Rightarrow O(N)
N \Rightarrow ARRAY'S SIZE

```
int BinarySearch(vector<int> arr, int n, int m, int end){
    int start = 0;
    int mid = start + (end - start)/2;
    int ans = -1;

    while(start <= end){
        // Step 03: create predicate function isPossibleSolution()
        if(isPossibleSolution(arr, n, m, mid)){
            ans = mid;
            end = mid - 1;
        }
        else{
            start = mid + 1;
        }
        mid = start + (end - start)/2;
    }
    return ans;
}
```

Binary search
T.C. \Rightarrow O(log end)
OR
O(log sum)

```
int allocateBooks(vector<int> arr, int n, int m){
    // Corner Case: 01
    if(m > n){
        return -1;
    }
    // Step 01: find total sum of array to create search space's end point (Maximum number of pages)
    int sum = 0;
    for(int i = 0; i < n; i++){
        sum += arr[i];
    }
    // Step 02: apply binary search on search space BinarySearch()
    int ans = BinarySearch(arr, n, m, sum);
    return ans;
}
```

```
for(int i=0; i<n; i++){  
    sum+=arr[i];  
}  
// Step 02: apply binary search on search space BinarySearch()  
int ans = BinarySearch(arr, n, m, sum);  
return ans;  
}
```

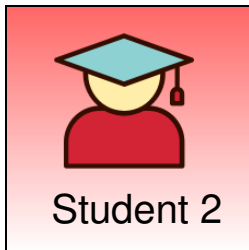
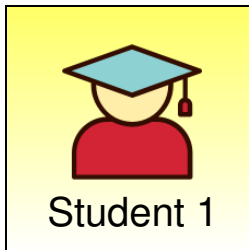
$T.C. \Rightarrow O(N \log \text{sum})$ $S.C. = O(1)$

Pages in books

1	2	3	4
---	---	---	---



4 Books



Possible page...

1	2	3	4
1	2	3	4
1	2	3	4

$$\max (1 + 2 + 3, 4) = 6$$

$$\max (1 + 2, 3 + 4) = 7$$

$$\max (1, 2 + 3 + 4) = 9$$

$$\min (6, 7, 9) = 6$$

```

// HW 05: Book Allocation Problem(GFG and Code studio)
#include <bits/stdc++.h>
bool isPossibleSolution(vector<int> arr, int n, int m, int mid){
    int cases = 1;
    int pageSum = 0;

    int i = 0;
    while(i<n){
        if(pageSum + arr[i] <= mid){
            pageSum += arr[i];
        }
        else{
            cases++;
            if(cases>m || arr[i]>mid){
                return false;
            }
            // reset pageSum
            pageSum = arr[i];
        }
        i++;
    }
    return true;
}

int BinarySearch(vector<int> arr, int n, int m, int end ){
    int start = 0;
    int mid = start + (end - start)/2;
    int ans = -1;

    while(start<=end){
        // Step 03: create predicate function isPossibleSolution()
        if(isPossibleSolution(arr, n, m, mid)){
            ans = mid;
            end = mid - 1;
        }
        else{
            start = mid + 1;
        }
        mid = start + (end - start)/2;
    }
    return ans;
}

int allocateBooks(vector<int> arr, int n, int m) {
    //Corner Case: 01
    if(m>n){
        return -1;
    }
    // Step 01: find total sum of array to create search space's end point (Maximum number of pages)
    int sum = 0;
    for(int i=0; i<n; i++){
        sum+=arr[i];
    }
    // Step 02: apply binary search on search space BinarySearch()
    int ans = BinarySearch(arr, n, m, sum);
    return ans;
}

```