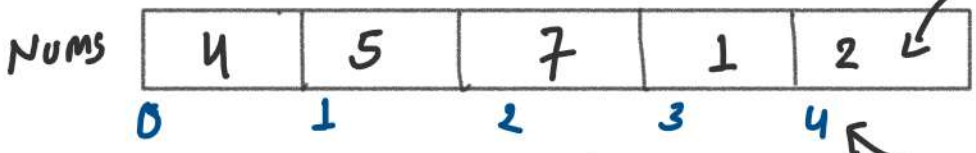
08|01|2024

DYNAMIC PROGRAMMING
CLASS - 2

1. House Robber (Leetcode-198)

Ex:1



Amount of money

Nums[4] = 2

Houli no = 5

i = 4

Find kya karna hai?

↳ Return the maximum amount of money

Case:1
Chori karne wala ghar

↳ 1th
↳ 3th
↳ 4th

amount
4
7
2
Total = 13

Case:2
Chori na karne wala ghar

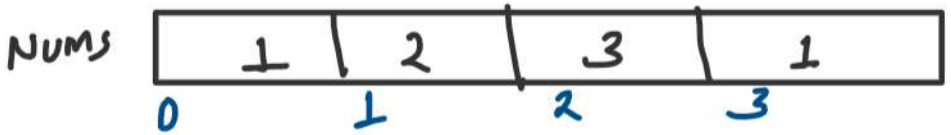
↳ 1th ✗ 0th ✗ 2th
↳ 3th ✗ 2th ✗ 4th

Amount
5
1
Total = 6

maxAmt = 13

Output

Example: 02



Case-1

Chori Hone wala Ghar | Amount

✓ 0th ✗ 1th | 1

✓ 2th ✗ 4th ✗ 3th | 3

4

Case-2

Chori Na Hone wala Ghar | Amount

✓ 1th ✗ 0th ✗ 2th | 2

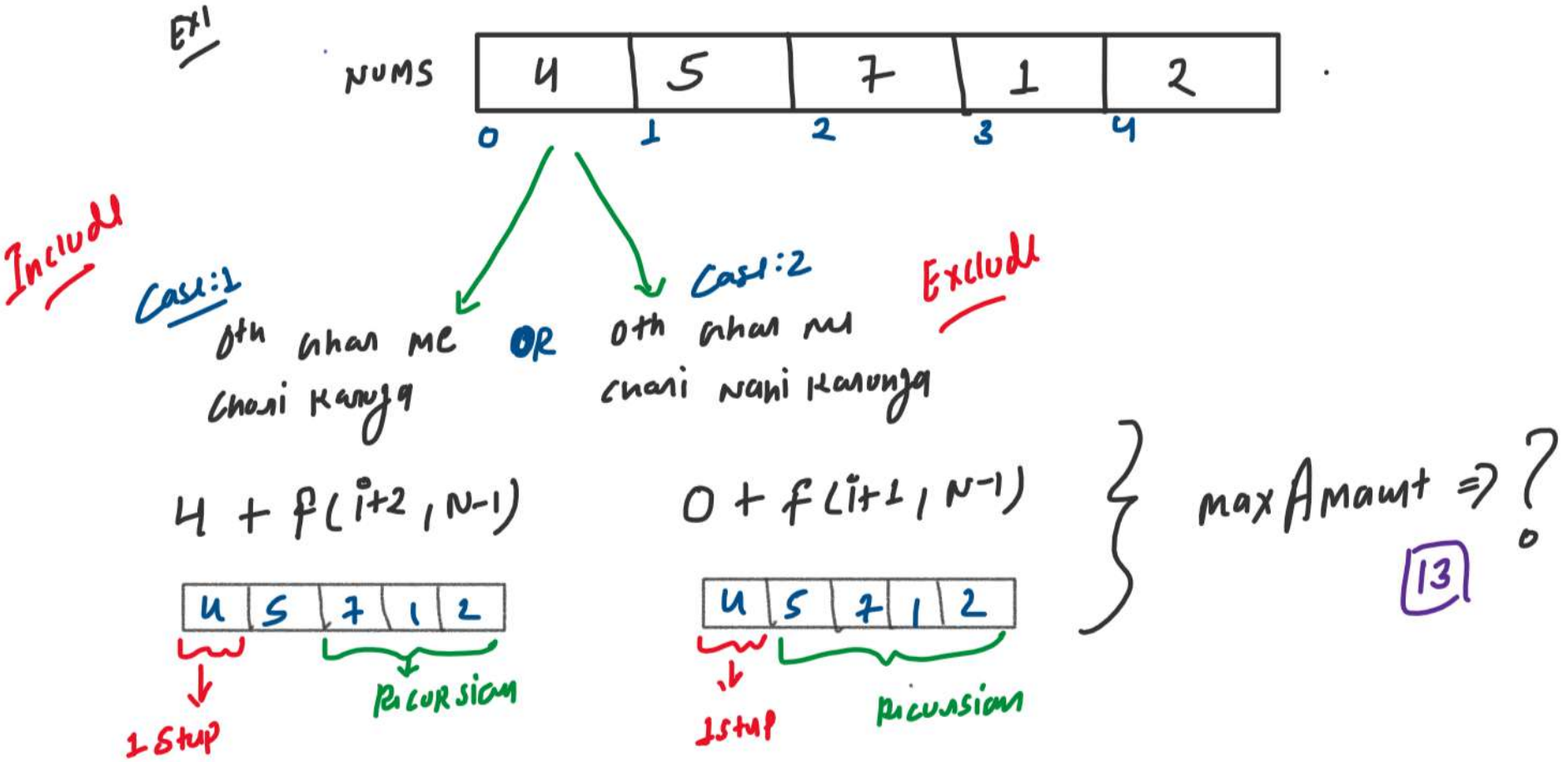
✓ 3th ✗ 2th | 1

3

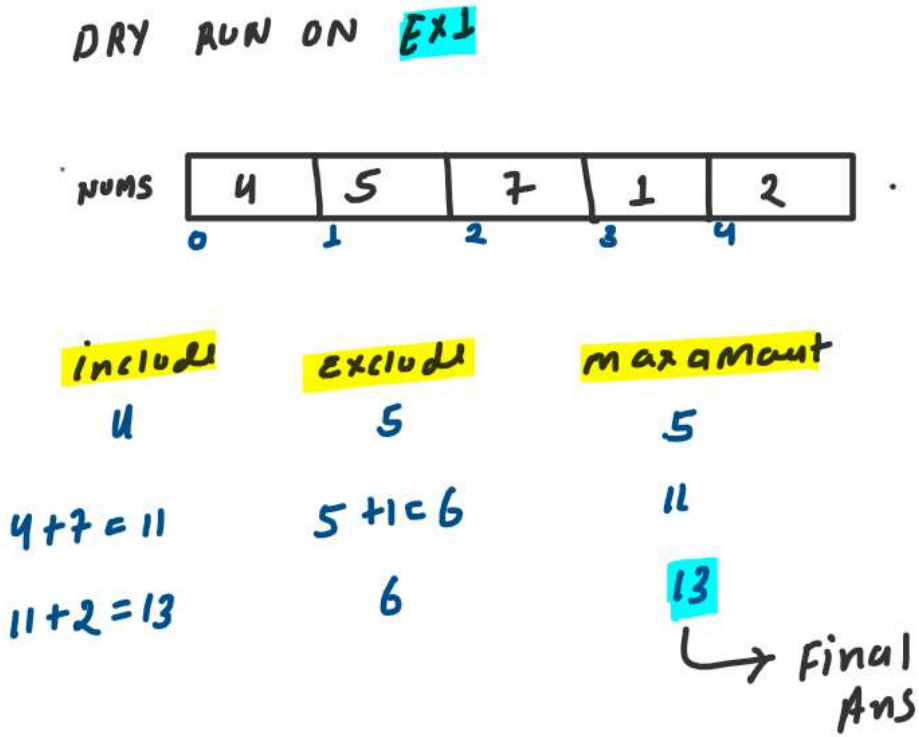
maxAmount = 4

output

Approach 1: Recursion
Inclusive and Exclusive Pattern



```cpp
// Approach 1: Normal Recursion (Inclusive and Exclusive Patterh

class Solution {
public:
    int solveUsingRecursion(vector<int>&nums, int index){
        // Base case
        if(index >= nums.size()){
            return 0;
        }

        // Recursive Relation
        int include = nums[index] + solveUsingRecursion(nums, index+2);
        int exclude = 0 + solveUsingRecursion(nums, index+1);

        // Getting max ammont from both
        int maxAmmount = max(include, exclude);
        return maxAmmount;
    }
    int rob(vector<int>& nums) {
        int index = 0;
        return solveUsingRecursion(nums, index);
    }
};
```



Approach 2: Top Down
Inclusive and Exclusive Pattern

N ↑ 0

TOP DOWN: Traverse from 0 to N

```cpp
// Approach 2: Top Down Approach (Inclusive and Exclusive Pattern

class Solution {
public:
    int solveUsingMemo(vector<int>&nums, int index, vector<int> &dp){
        // Base case
        if(index >= nums.size()){
            return 0;
        }

        // Step 3: if ans max amount already exist then return ans
        if(dp[index] != -1){
            return dp[index];
        }

        // Step 2: store ans max amount and return ans using DP array
        int include = nums[index] + solveUsingMemo(nums, index+2, dp);
        int exclude = 0 + solveUsingMemo(nums, index+1, dp);
        dp[index] = max(include, exclude);
        return dp[index];
    }
    int rob(vector<int>& nums) {
        // Step 1: create DP array
        int n = nums.size();
        vector<int> dp(n, -1);
        int index = 0;
        return solveUsingMemo(nums, index, dp);
    }
};
```

Main whole ARRAY ka ans nahi bata sakta hun to main manually last index ka amount ans main store kar diya hai

Approach 3: Bottom UP
Inclusive and Exclusive Pattern

N ↓ 0

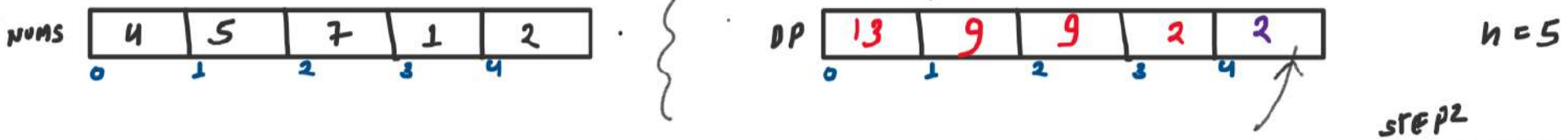
BOTTOM UP: Traverse from N to 0

```cpp
// Approach 3: Bottom-up Approach (Inclusive and Exclusive Pattern

class Solution {
public:
    int solveUsingTabu(vector<int>&nums, int index){
        // Step 1: create DP array
        int n = nums.size();
        vector<int> dp(n, -1);

        // Step 2: fill initial data in DP array according to recursion base case
        dp[n-1] = nums[n-1];

        // Step 3: fill the remaining DP array according to recursion formula/logic
        for(int index = n - 2; index >= 0; index--){
            int tempAns = 0;
            if(index + 2 < n){
                // Corner Case
                tempAns = dp[index+2];
            }
            int include = nums[index] + tempAns;
            int exclude = 0 + dp[index+1];
            dp[index] = max(include, exclude);
        }
        // Return ans
        return dp[0];
    }
    int rob(vector<int>& nums) {
        int index = 0;
        return solveUsingTabu(nums, index);
    }
};
```

Note Return Ans

99% jo index hum pass karte hai ussi index par hamara ans hota hai

Bottom UP DRY RUN



STEP2

```
DP[5-1] = NUM[5-1]
DP[4] = 2
```

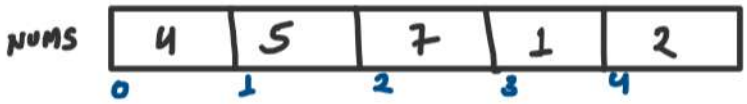
Initially (Jab last me chori karungi T° sirf wahi pass utna hi Amount hoga jitna uski pass rakhi hai)

STEP3

index	Includi	Excludi	DP[Index] max Amount
3	1 + 0 = 1	0 + 2 = 2	2
2	7 + 2 = 9	0 + 2 = 2	9
1	5 + 2 = 9	0 + 9 = 9	9
0	4 + 9 = 13	0 + 9 = 9	13 → Return DP[0]

Approach 4: Space Optimization
Inclusive and Exclusive Pattern

Ex /



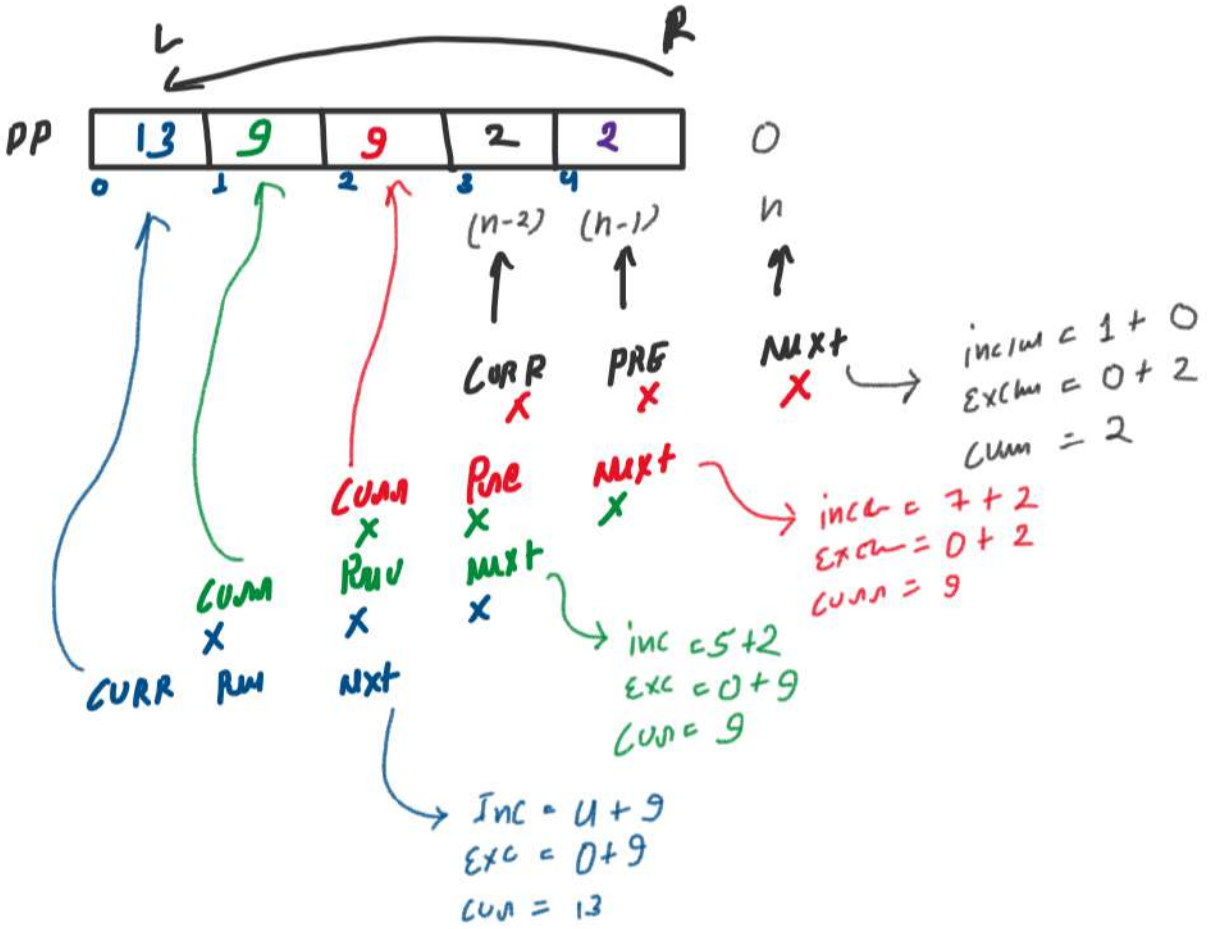
CURR → maxAmount

PREV → nums[n-1]

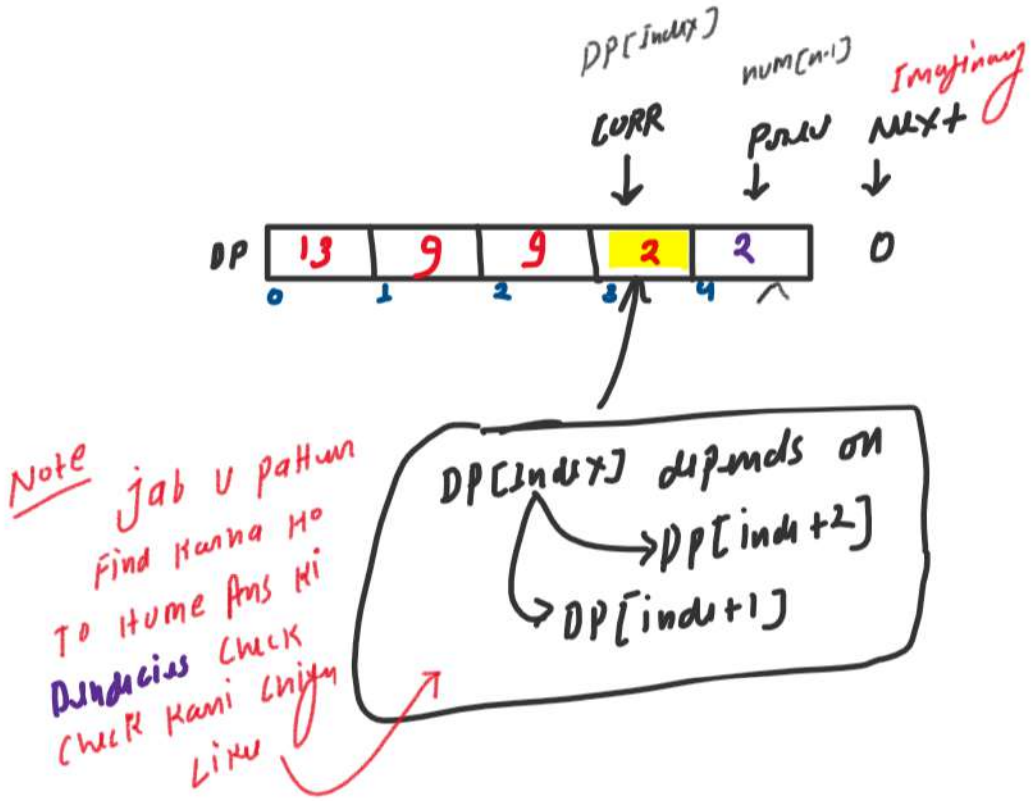
NXT → 0

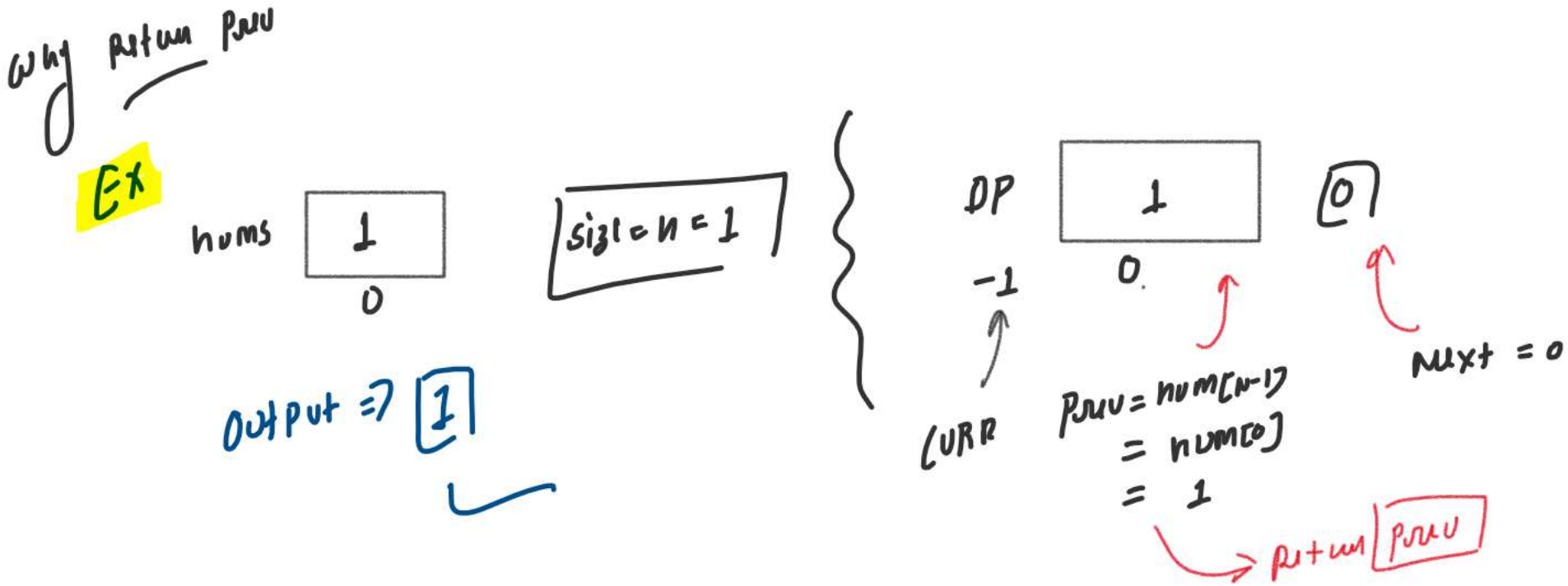
T.C. = O(N)

S.C. = O(1)



```cpp
// Approach 4: Space Optimization Approach (Inclusive and Exclusive Pattern)

class Solution {
public:
    int solveUsingSO(vector<int>&nums, int index){

        int n = nums.size();
        int next = 0;
        int prev = nums[n-1];
        int curr = 0;

        for(int index = n - 2; index >= 0; index--){
            int tempAns = 0;
            if(index + 2 < n){
                // Corner Case
                tempAns = next;
            }
            int include = nums[index] + tempAns;
            int exclude = 0 + prev;
            curr = max(include, exclude);

            // Update krna bhool jata hu
            next = prev;
            prev = curr;
        }
        return prev;
    }
    int rob(vector<int>& nums) {
        int index = 0;
        return solveUsingSO(nums, index);
    }
};
```




Why return prev
Ex
nums
1
0
size = n = 1
Output => 1
DP
1
0
-1
0
CURR
prev = num[n-1]
= num[0]
= 1
next = 0
return prev

2. Coin Change (Leetcode-322)

Example 1:
Input: coins = [1,2,5], amount = 11
Output: 3
Explanation: 11 = 5 + 5 + 1

Example 2:
Input: coins = [2], amount = 3
Output: -1

Example 3:
Input: coins = [1], amount = 0
Output: 0

way 1

Coin = 1 takes 11 times ⇒ 11 × 1 = 11 — 11 coins

Coin = 2 takes 5 times & Coin = 1 takes 1 time ⇒ (5 × 2) + 1 = 11 — 6 coins

Coin = 5 takes 2 times & Coin = 1 takes 1 time ⇒ (2 × 5) + 1 = 11 — 3 coins

Min (11 , 6 , 3) ⇒ 3 coins

Approach 1: Recursion

Explore All Possible Ways Pattern

Ways

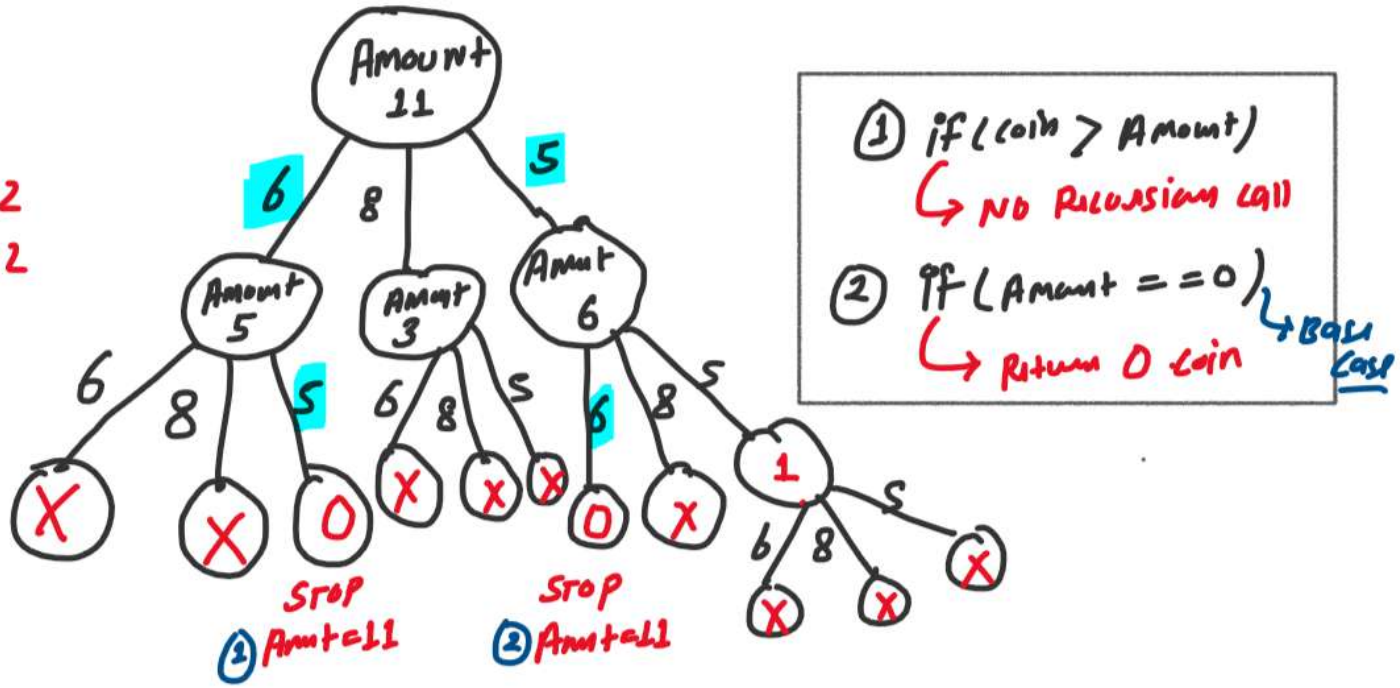
Amount = 11

Coins = 6, 8, 5

Call
① 6 + 5 = 11 } coins = 2
② 5 + 6 = 11 } coins = 2

min(2, 2) = 2

Output



Amount = N = 11

Coins [6 1 8 5]

$f(N) = f(11)$ — 6 coin → $f(5) \Rightarrow \boxed{1} + f(N-6)$ (1 STEP)

$f(N) = f(11)$ — 8 coin → $f(3) \Rightarrow 1 + f(N-8)$

$f(N) = f(11)$ — 5 coin → $f(6) \Rightarrow 1 + f(N-5)$

MIN

{ But AISA Zaroori Nahi Hai Ki Har Bar 3 Types Ka Coin Hi Ho Usse Jada Bhi Ho Sakta Hai To Har Ek Coin Ko Check Karna Ki Let's Loop Ka Use Karke. Ki (Coin <= Amount) Hai Ya Nahi }

DRY RUN

EX

Coins = [3, 2, 5]
0 1 2

Amount = 10

index = 0

mini = 3

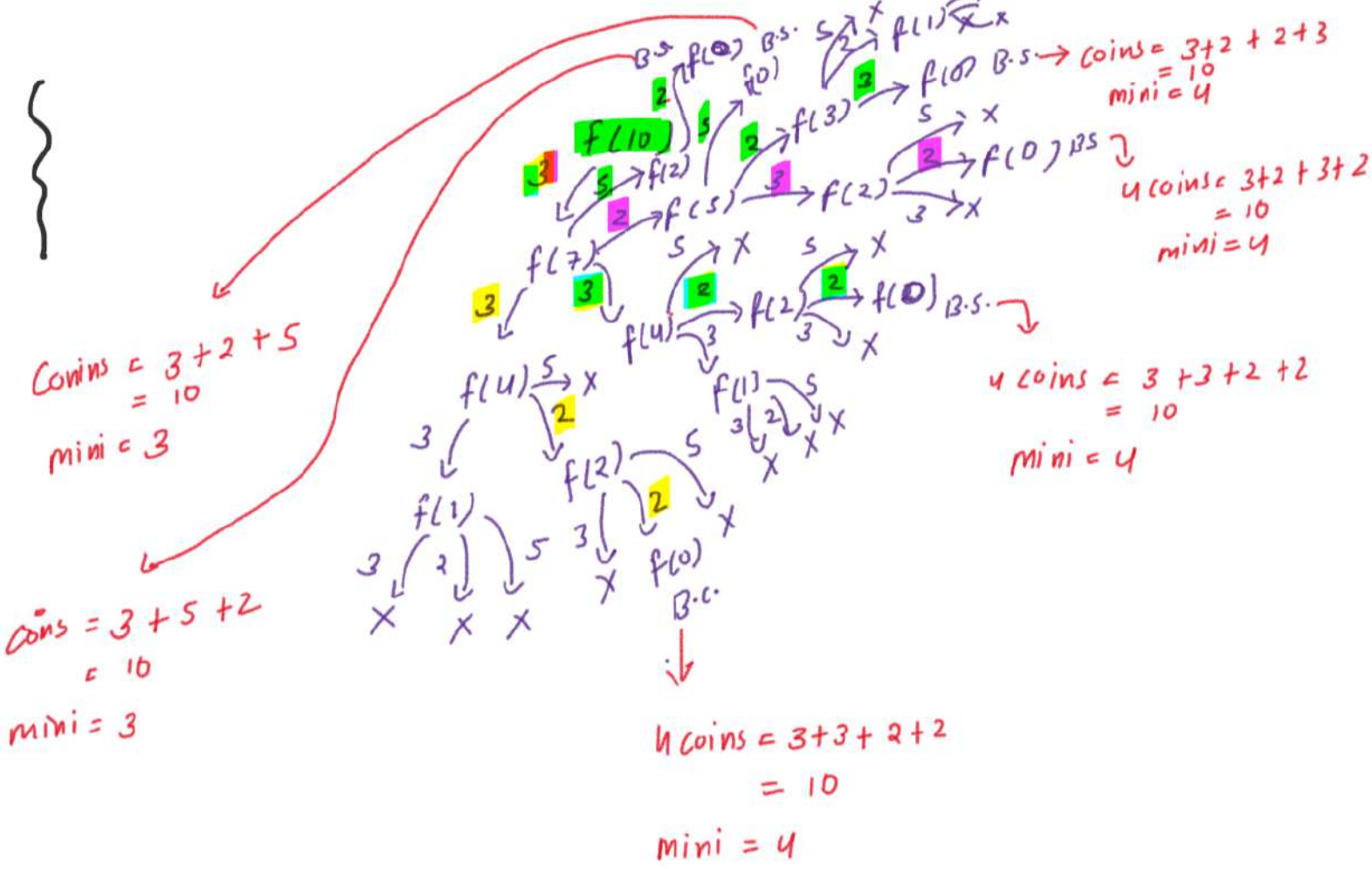
Coins = 3+5+2

Coins = 3 + 2 + 5
= 10

mini = 3

Coins = 3 + 5 + 2
= 10

mini = 3



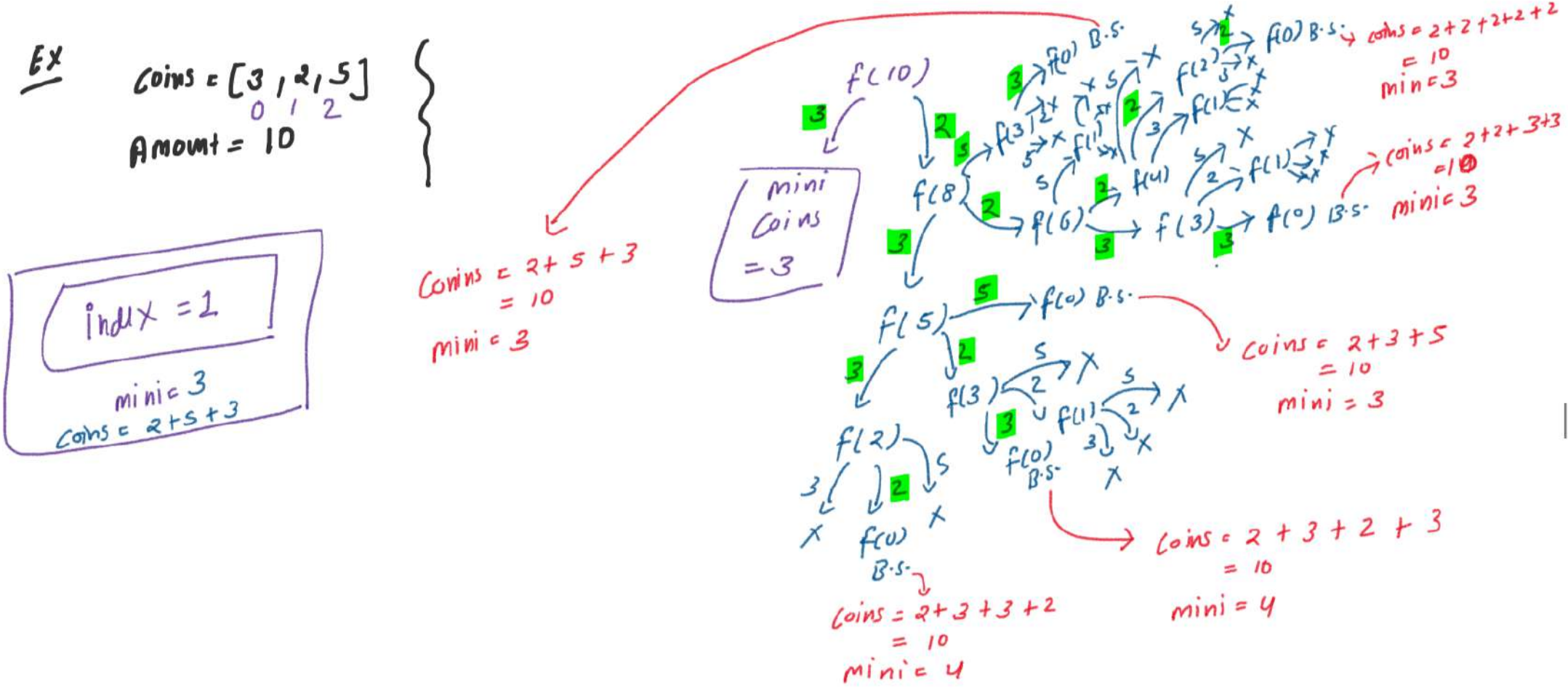
4 coins = 3+2 + 2+3
= 10
mini = 4

4 coins = 3+2 + 3+2
= 10
mini = 4

4 coins = 3 +3 +2 +2
= 10
mini = 4

4 coins = 3+3 + 2+2
= 10

mini = 4


EX
Coins = [3, 2, 5]
0 1 2
Amount = 10
index = 1
mini = 3
Coins = 2+5+3
f(10)
mini Coins = 3
Coins = 2+5+3
= 10
mini = 3
f(8)
f(6)
f(5)
f(3)
f(2)
f(1)
f(0) B.S.
Coins = 2+2+2+2+2
= 10
min = 3
Coins = 2+2+3+3
= 10
minic 3
Coins = 2+3+5
= 10
mini = 3
Coins = 2+3+2+3
= 10
mini = 4
Coins = 2+3+3+2
= 10
mini = 4

Ex

Coins = [3, 2, 5]
0 1 2

Amount = 10

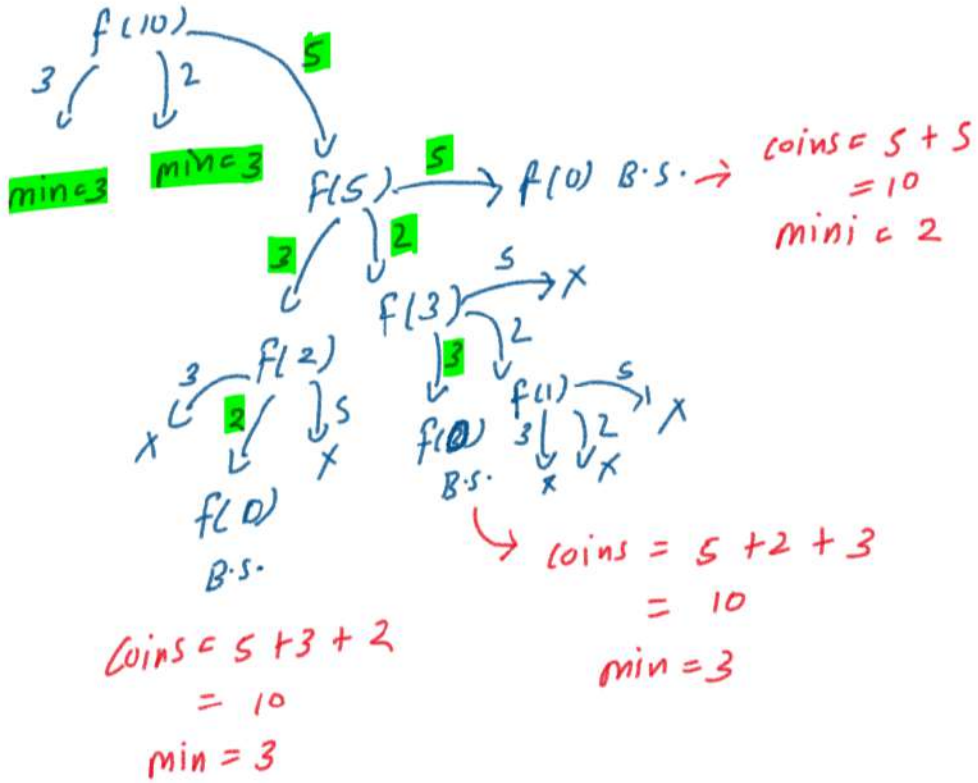
index = 2

minc = 2

Coins = 5+5

Final output

→ 2



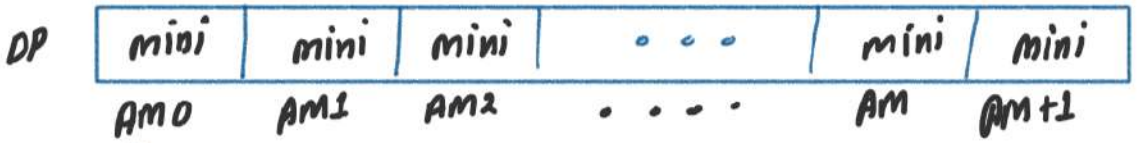
```cpp
// 2. Coin Change (Leetcode-322) Exploring All Possible Ways Pattern

// Approach 1: Normal Recursion Approach

class Solution {
public:
    int solveUsingRec(vector<int>& coins, int amount){
        // Base case
        if(amount == 0){
            return 0;
        }

        // Recursive relation
        int mini = INT_MAX;
        for(int i=0; i<coins.size(); i++){
            int coin = coins[i];

            // Jav amount >= coin se tabhi change kr skte hai
            if(amount - coin >= 0){
                int recKaAns = solveUsingRec(coins, amount - coin);
                // Recursion ka ans (INT_MAX + 1) outOfRange to nhi hai check krlo
                if(recKaAns != INT_MAX ){
                    int ans = 1 + recKaAns;
                    mini = min(mini, ans);
                }
            }
        }
        return mini;
    }

    int coinChange(vector<int>& coins, int amount) {
        int ans = solveUsingRec(coins, amount);
        if(ans == INT_MAX){
            // Invalid ans hai
            return -1;
        }
        else{
            // Valid ans hai
            return ans;
        }
    }
};
```

TLE

Approach 2: Top Down
Explore All Possible Ways Pattern

TOP DOWN: Traverse from amount to 0

STEP 1, 2 (size = Amt + 1)



Kitna Amount ↳ Kam se kam coin Exchange karna

STEP 3 Agar kisi Amount ke minimum coins KO DP me store kar diya hoi TO Dobara store nahi karna. [Return dp[AM]]

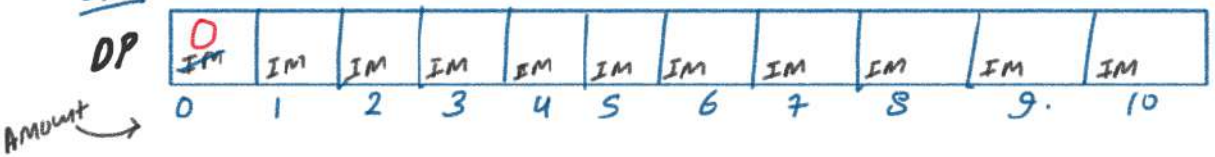
```cpp
// 2. Coin Change (Leetcode-322) Exploring All Possible Ways Pattern

// Approach 2: Top Down Approach

class Solution {
public:
    int solveUsingMemo(vector<int>& coins, int amount, vector<int>&dp){
        // Base case
        if(amount == 0){
            return 0;
        }

        // Step 3: if ans already exist then return ans
        if(dp[amount] != -1){
            return dp[amount];
        }

        // Step 2: store ans and return ans using DP array
        // TOP DOWN: Traverse from amount to 0
        int mini = INT_MAX;
        for(int i=0; i<coins.size(); i++){
            int coin = coins[i];

            // Jav amount >= coin se tabhi change kr skte hai
            if(amount - coin >= 0){
                int recKaAns = solveUsingMemo(coins, amount - coin, dp);
                // Recursion ka ans (INT_MAX + 1) outOfRange to nhi hai check krlo
                if(recKaAns != INT_MAX ){
                    int ans = 1 + recKaAns;
                    mini = min(mini, ans);
                }
            }
        }
        // Store ANS
        dp[amount] = mini;
        // return ANS
        return dp[amount];
    }

    int coinChange(vector<int>& coins, int amount) {
        // Step 1: create DP array
        int n = amount;
        vector<int> dp(n+1, -1);

        int ans = solveUsingMemo(coins, amount, dp);
        if(ans == INT_MAX){
            // Invalid ans hai
            return -1;
        }
        else{
            // Valid ans hai
            return ans;
        }
    }
};
```

Approach 3: Bottom Up
Explore All Possible Ways Pattern

BOTTOM UP: Traverse from 0 to amount

Ex
Coins = [3, 2, 5]
0 1 2
Amount = 10

STEP1



STEP2 Analys- Base case and store initial data to DP

Amount == 0
return 0
} DP[amount] = 0 , when Amnt = 0

STEP3 Fill remaining dp based on Recursive Relation

```cpp
// 2. Coin Change (Leetcode-322) Exploring All Possible Ways Pattern

// Approach 3: Bottom-up

class Solution {
public:
    int solveUsingTabu(vector<int>& coins, int amount) {
        // Step 1: create DP array
        int n = amount;
        vector<int> dp(n+1, INT_MAX);

        // Step 2: fill initial data in DP array according to recursion base case
        dp[0] = 0;

        // Step 3: fill the remaining DP array according to recursion formula/logic
        // BOTTOM UP: Traverse from 0 to amount
        for( int value = 1; value<=amount; value++) {
            int mini = INT_MAX;
            for(int i=0; i<coins.size(); i++) {

                if(value - coins[i] >= 0) {
                    int recKaAns = dp[value - coins[i] ];

                    if(recKaAns != INT_MAX) {

                        int ans = 1 + recKaAns;
                        mini  = min(mini, ans);
                    }
                }
            }
            dp[value] = mini;
        }
        return dp[amount];
    }

    int coinChange(vector<int>& coins, int amount) {
        int ans = solveUsingTabu(coins, amount);
        if(ans == INT_MAX){
            // Invalid ans hai
            return -1;
        }
        else{
            // Valid ans hai
            return ans;
        }
    }
};
```

STEP 3

2 के लिए only 1 coin ki zaroorat Hai.

(Value) DP

0	IM	1	1	2	1	2	2	2	3	2
IM	IM	IM	IM	IM	IM	IM	IM	IM	IM	IM
0	1	2	3	4	5	6	7	8	9	10

Amount →

(2) (3) (2+2) (5) (3+3) (5+2) (5+3) (3+3+3) (5+5)

Ex
Coins = [3, 2, 5]
0 1 2
Amount = 10

value	mini	i	value − coins[i] >= 0	recAns = DP[−]	recAns != IM	ans = 1 + recAns	mini	i++	DP[value] = mini	value++
1	IM	0	1-3 X	—	—	—	—	1		
	IM	1	1-2 X	—	—	—	—	2		
	IM	2	1-5 X	—	—	—	—	3		
	IM	3 stop	—	—	—	—	—	—	IM	2
2	IM	0	2-3 X	—	—	—	—	1		
	IM	1	2-2 ✓	0	✓	1	1	2		
	1	2	2-5 X	—	—	—	—	3		
	1	3 stop	—	—	—	—	—	—	1	3

value	mini	i	value - coins[i] >= 0	recAns = DP[ ]	recAns != IM	ans = 1 + recAns	min	i++	DP[value] = mini	value++
1	IM	0	1-3 X	—	—	—	—	1		
	IM	1	1-2 X	—	—	—	—	2		
	IM	2	1-5 X	—	—	—	—	3		
	IM	3 stop	—	—	—	—	—	—	IM	2
2	IM	0	2-3 X	—	—	—	—	1		
	IM	1	2-2 ✓	0	✓	1	1	2		
	1	2	2-5 X	—	—	—	—	3		
	1	3 stop	—	—	—	—	—	—	1	3
3	IM	0	3-3 ✓	0	✓	1	1	1		
	1	1	3-2 ✓	1	✓	2	1	2		
	1	2	3-5 X	—	—	—	—	3		
	1	3 stop	—	—	—	—	—	—	1	4

DRY RUN AS IT

Approach 4: Optimal Solution

Explore All Possible Ways Pattern

→ (This Approach is not possible)

Why ?

output → DP[Amount] Depends on whole Array

So Yaha Koi pattern ban nahi pa Rha Hai

$$DP[value] = DP[value - coins[i]]$$