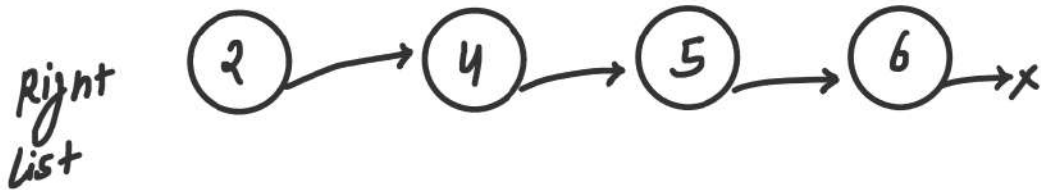
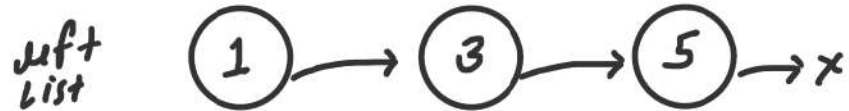


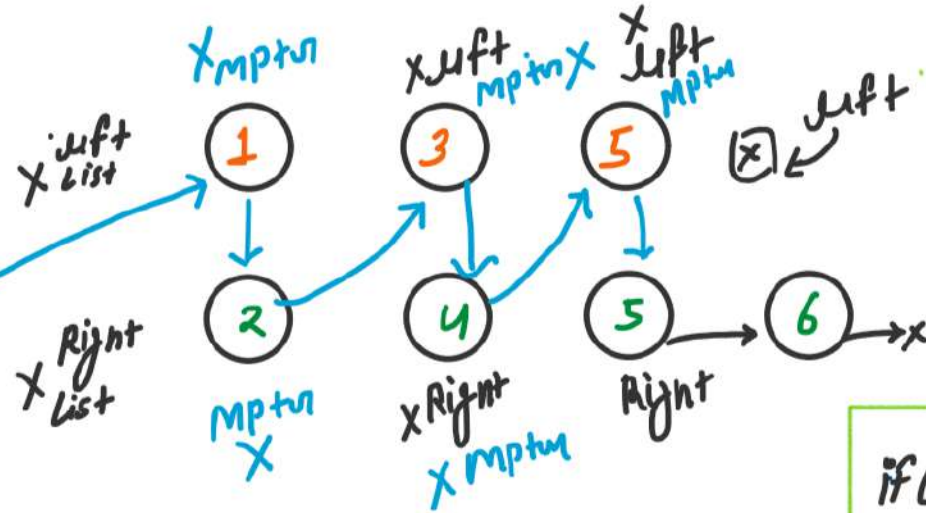
HW 01: Merge Two Sorted Lists (Leetcode-21)

Ex



DRY RUN

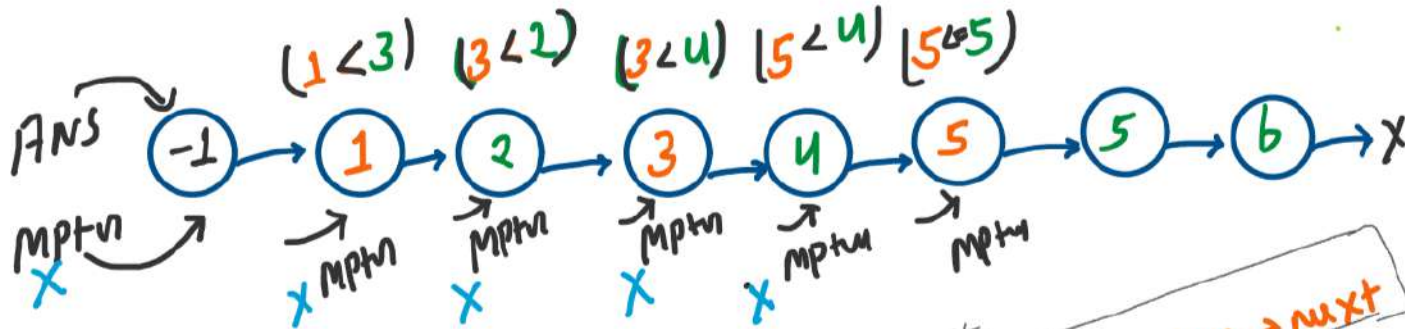
list Ans = -1
 mptn = Ans
 [-1 | •]
 data next



```

if (left) {
  mptn → next = left;
}
if (right) {
  mptn → next = right;
}

```



 return Ans → next

```

if (left → data <= right → data)
{
  mptn → next = left;
  mptn = left;
  left = left → next;
}
else {
  mptn → next = right;
  mptn = right;
  right = right → next;
}

```

```

/**
 * Definition for singly-linked list.
 * struct ListNode {
 *     int val;
 *     ListNode *next;
 *     ListNode() : val(0), next(nullptr) {}
 *     ListNode(int x) : val(x), next(nullptr) {}
 *     ListNode(int x, ListNode *next) : val(x), next(next) {}
 * };
 */
class Solution {
public:
    ListNode* mergeTwoLists(ListNode* left, ListNode* right) {
        if(left == NULL) return right;
        if(right == NULL) return left;

        ListNode* ans = new ListNode(-1);
        ListNode* mptr = ans;

        while(left != NULL && right != NULL){
            if(left->val <= right->val){
                mptr->next = left;
                mptr = left;
                left = left->next;
            }
            else{
                mptr->next = right;
                mptr = right;
                right = right->next;
            }
        }

        if(left != NULL){
            mptr->next = left;
            // mptr = left;
            // left = left->next;
        }

        if(right != NULL){
            mptr->next = right;
            // mptr = right;
            // right = right->next;
        }

        return ans->next;
    }
};

```

Time complexity = $O(N)$

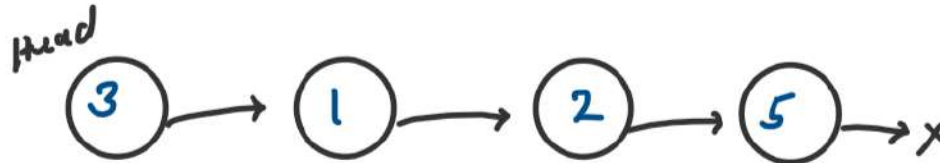
where, N is total numbers of nodes of both list.

Space complexity = $O(1)$

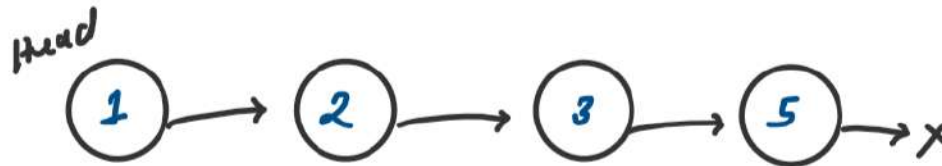
No extra space used by ans and mptr.

HW 02: Sort Lists using Merge Sort (Leetcode-148)

Input



Output



MERGE SORT ALGORITHM

Step 1: Find mid position of the list

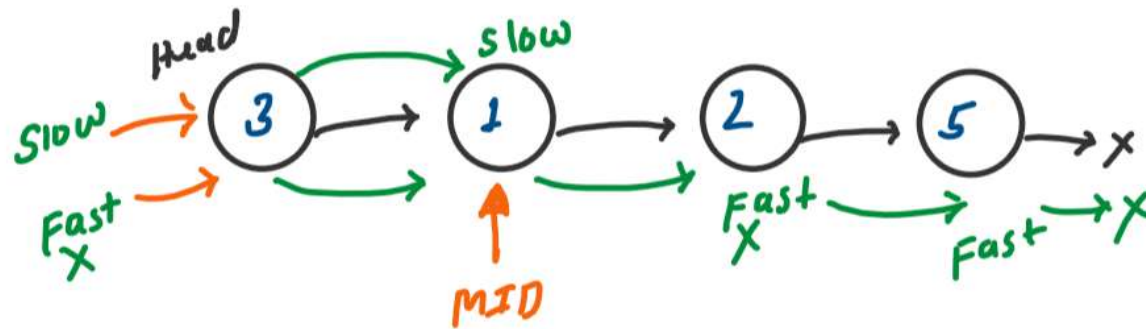
Step 2: Divide list into two half using mid

Step 3: Sort RE

Step 4: Merge both sorted list left and right

Step 1: Find mid position of the List

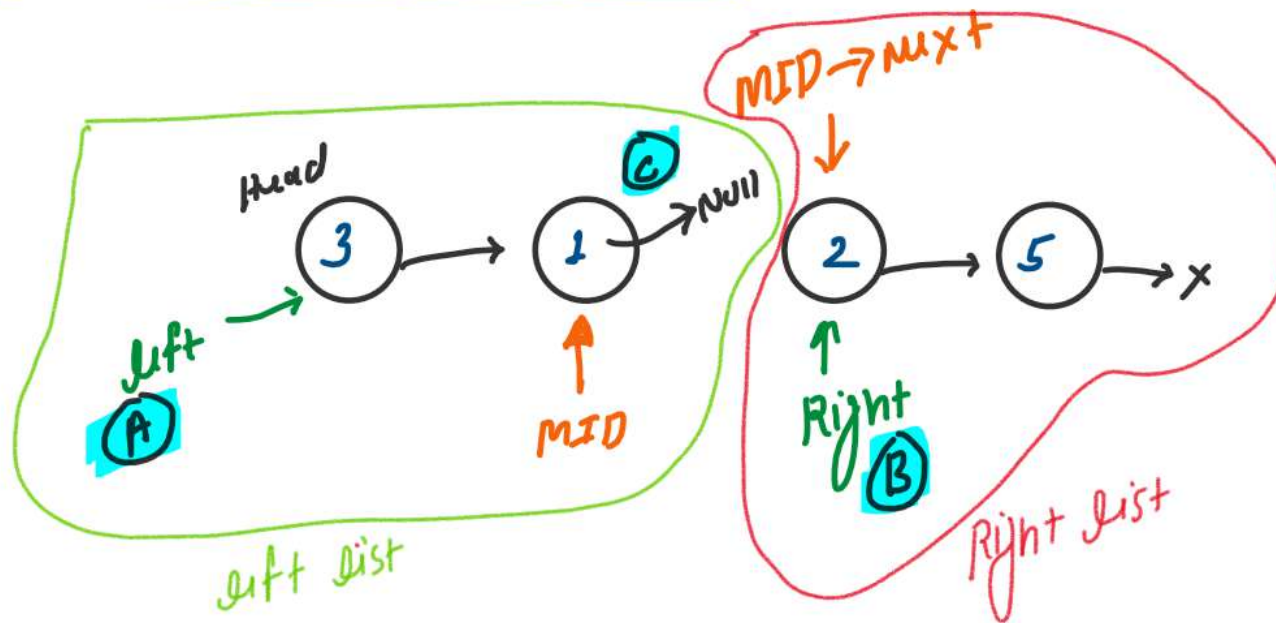
using slow & Fast Algorithm



```
ListNode* getMid(ListNode* head){
    ListNode* slow = head;
    ListNode* fast = head;

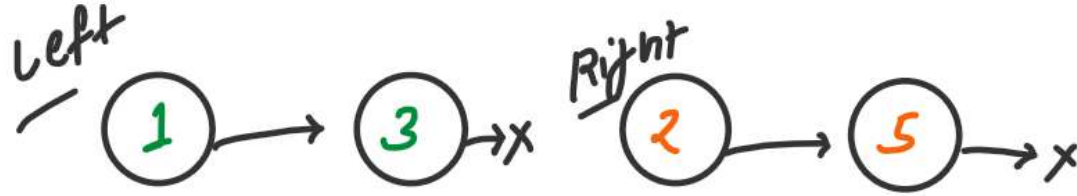
    while(fast->next != NULL){
        fast = fast->next;
        if(fast->next != NULL){
            fast = fast->next;
            slow = slow->next;
        }
    }
    return slow;
}
```

Step 2: Divide list into two half using mid



- (A) $\text{listNode} * \text{Left} = \text{head};$
- (B) $\text{listNode} * \text{Right} = \text{mid} \rightarrow \text{next};$
- (C) $\text{mid} \rightarrow \text{next} = \text{Null};$

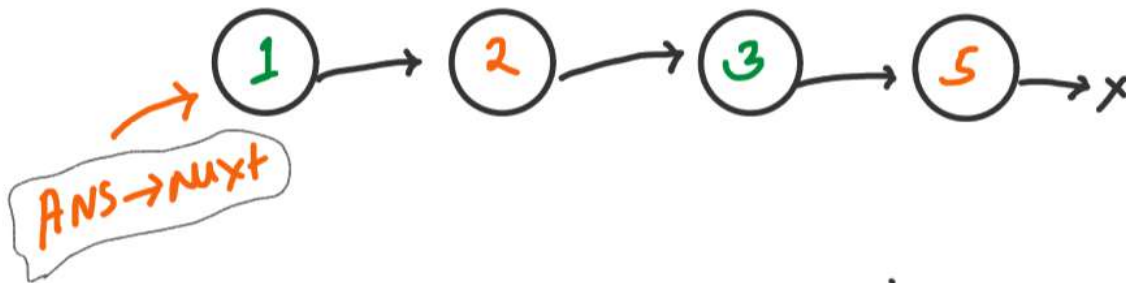
Step 3: Sort both List left and right RE



$left = sortList(left)$

$Right = sortList(Right)$

Step 4: Merge both sorted List Left and right



```
ListNode* merge(ListNode* left, ListNode* right) {  
    if(left == NULL) return right;  
    if(right == NULL) return left;  
  
    ListNode* ans = new ListNode(-1);  
    ListNode* mptr = ans;  
  
    while(left != NULL && right != NULL){  
        if(left->val <= right->val){  
            mptr->next = left;  
            mptr = left;  
            left = left->next;  
        }  
        else{  
            mptr->next = right;  
            mptr = right;  
            right = right->next;  
        }  
    }  
  
    if(left != NULL){  
        mptr->next = left;  
    }  
  
    if(right != NULL){  
        mptr->next = right;  
    }  
  
    return ans->next;  
}
```


COMPLETE CODE

```
// HW 02: Sort Lists using Merge Sort (Leetcode-148)
/**
 * Definition for singly-linked list.
 * struct ListNode {
 *     int val;
 *     ListNode *next;
 *     ListNode() : val(0), next(nullptr) {}
 *     ListNode(int x) : val(x), next(nullptr) {}
 *     ListNode(int x, ListNode *next) : val(x), next(next) {}
 * };
 */
class Solution {
public:
    ListNode* getMid(ListNode* head){...}

    ListNode* merge(ListNode* left, ListNode* right){...}

    ListNode* sortList(ListNode* head) {
        // Base case
        if(head == NULL || head->next == NULL){
            return head;
        }

        // Step 1: Find mid position of the list
        ListNode* mid = getMid(head); ✓

        // Step 2: Divide list into two half using mid
        ListNode* left = head;
        ListNode* right = mid->next;
        mid->next = NULL;

        // Step 3: Sort RE
        left = sortList(left);
        right = sortList(right);

        // Step 4: Merge both sorted list left and right
        ListNode* mergeLR = merge(left, right); ✓
        return mergeLR;
    }
};
```

T.C.

getMid $\Rightarrow T.C. = O(N)$

merge $\Rightarrow T.C. = O(N)$

sortList $\Rightarrow T.C. \Rightarrow O(\log N)$

Overall
T.C.

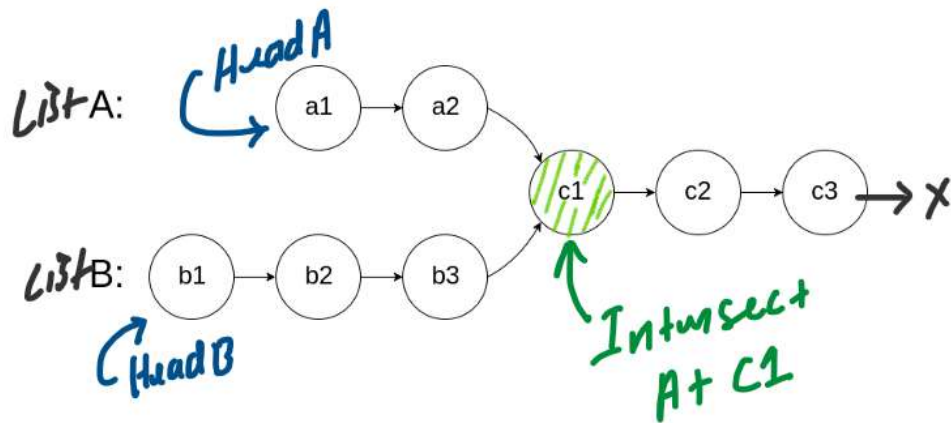
$$\begin{aligned} &= O((O(\log N) + O(N)) * (O(\log N))) \\ &= O((O(N) + O(N)) * (O(\log N))) \\ &= O(O(N) * O(\log N)) \\ &= O(N \log N) \end{aligned}$$

HW 03: Intersection of Two Linked Lists (Leetcode-160)

PROBLEM STATEMENT:

Given the heads of two singly linked-lists headA and headB, return the node at which the two lists intersect. If the two linked lists have no intersection at all, return `NULL`.

For example,
the following two linked lists begin to intersect at node c1:

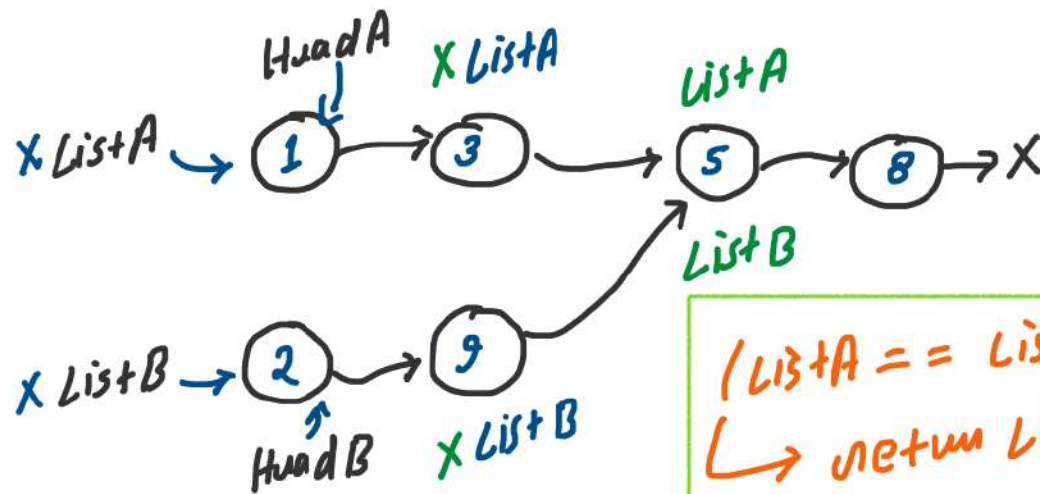


Note that the linked lists must retain their original structure after the function returns.

Ex:1 DRY RUN

Equal Length of List A & List B

List A Length = List B Length = 4



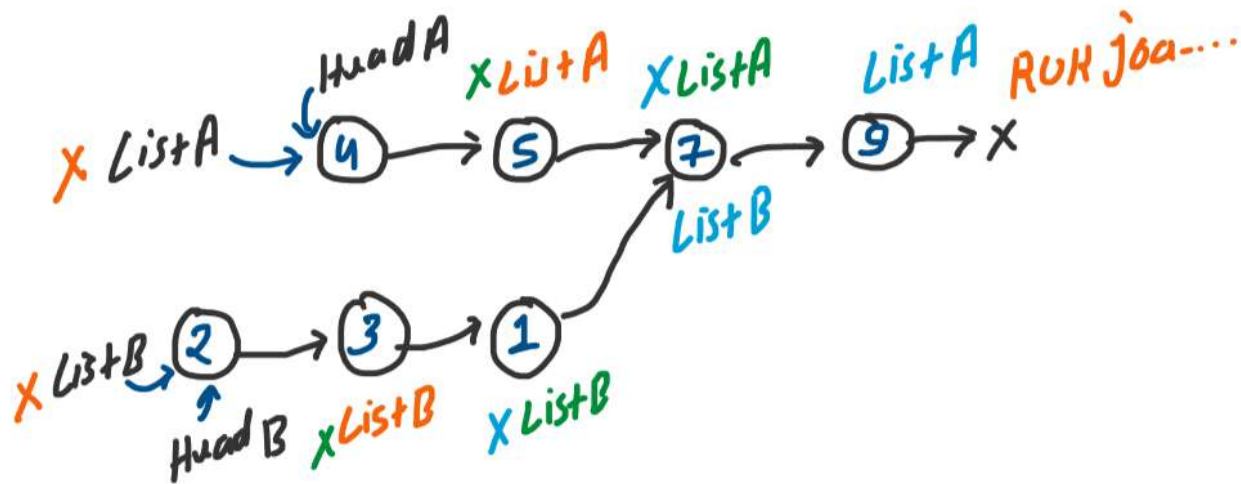
(List A == List B)
↳ return List A
Output = 5

```
ListNode* listA = headA;
ListNode* listB = headB;

while(listA->next != NULL && listB->next != NULL){
    if(listA == listB){
        // Agar listA and listB equal length ki hai
        // iska mtlb wo yahin se intersect Node return kar degi
        return listA;
    }
    listA = listA->next;
    listB = listB->next;
}
```

Ex: 2 DRY RUN

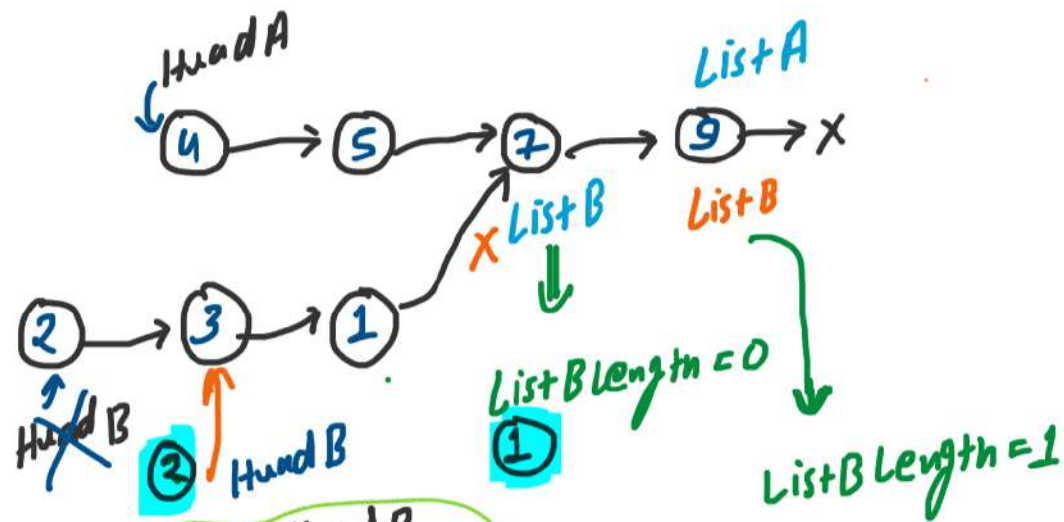
NO Equal Length of List A & List B



Step: 1

```
ListNode* listA = headA;
ListNode* listB = headB;

while(listA->next != NULL && listB->next != NULL){
    if(listA == listB){
        // Agar listA and listB equal length ki hai
        // iska mtlb wo yahin se intersect Node return kar degi
        return listA;
    }
    listA = listA->next;
    listB = listB->next;
}
```



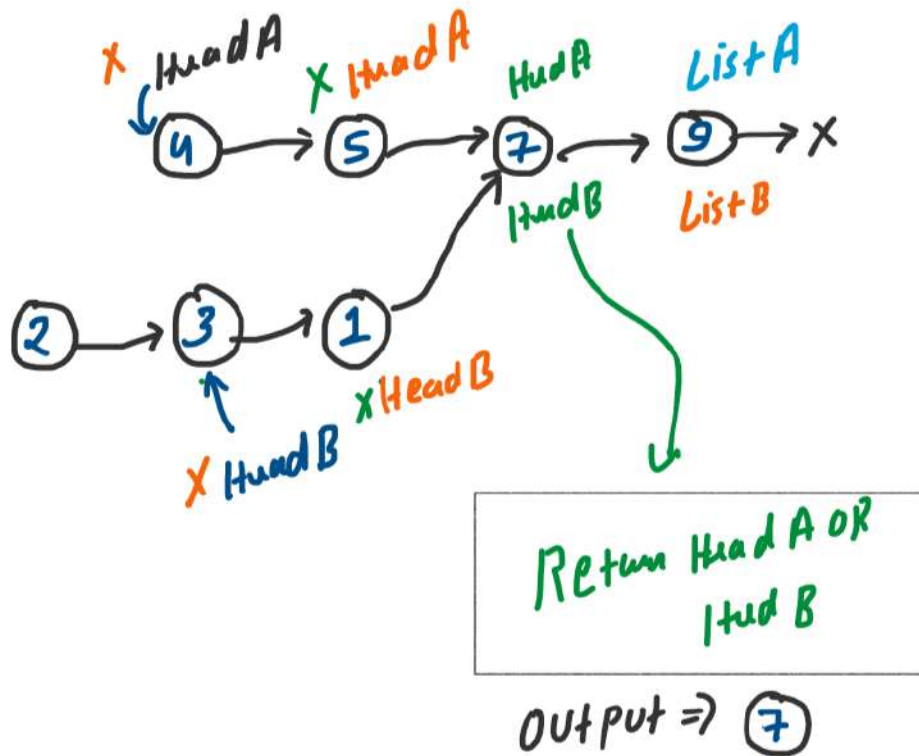
Again Head B

No iss (3) node
par set kardi
TO Intersection
Node mil jayiga

But How

- 1 Find Diff b/w both List
- 2 Set at head at right position.

```
// Me yanha tak tabhi pahuncha hu
// jab listA and listB ki length equal nhi hai
if(listA->next == NULL){
    // First.....
    // ListB is bigger
    // We need to find the length of ListB
    int listBLength = 0;
    while(listB->next != 0){
        listBLength++;
        listB = listB->next;
    }
    // Second.....
    // In starting, Set headB at right node to get the intersection Node
    while(listBLength--){
        headB = headB->next;
    }
}
```

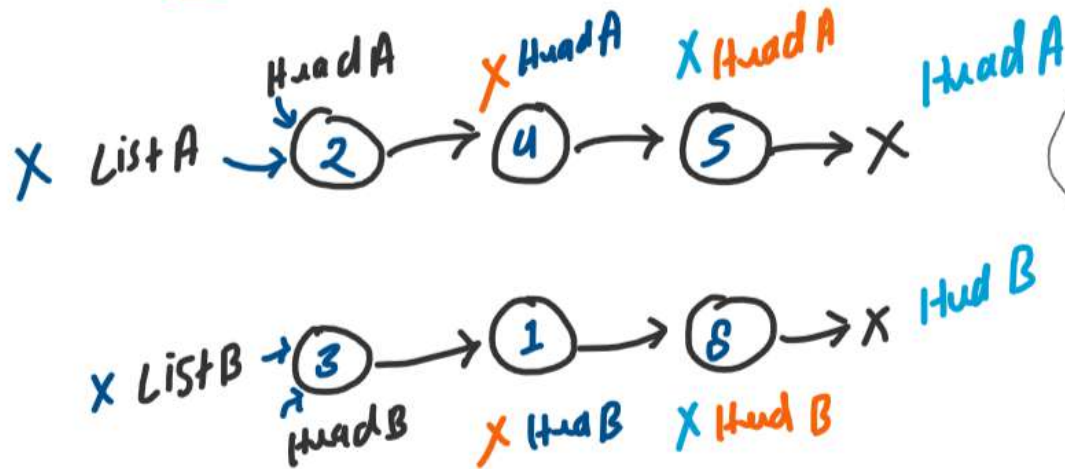


```
// Me yaha tak tabhi pahucha hu jab headA and headB starting me right  
// Node par set ho chuke honge  
// We need to traverse again to get the intersection Node  
while(headA != headB){  
    headA = headA->next;  
    headB = headB->next;  
}  
return headA;
```


Ex:3

DRY RUN

NO Intersection b/w both Lists



```
// Me yaha tak tabhi pahucha hu jab headA and headB starting me right
// Node par set ho chuke honge
// We need to traverse again to get the intersection Node
while(headA != headB){
    headA = headA->next;
    headB = headB->next;
}
return headA;
```

Return Head A
OR
Head B

Output = Null

Complete code

```
// HW 03: Intersection of Two Linked Lists (Leetcode-160)
/**
 * Definition for singly-linked list.
 * struct ListNode {
 *     int val;
 *     ListNode *next;
 *     ListNode(int x) : val(x), next(NULL) {}
 * };
 */
class Solution {
public:
    ListNode *getIntersectionNode(ListNode *headA, ListNode *headB) {
        ListNode* listA = headA;
        ListNode* listB = headB;

        while(listA->next != NULL && listB->next != NULL){
            if(listA == listB){
                // Agar listA and listB equal length ki hai
                // iska mtlb wo yahin se intersect Node return kar degi
                return listA;
            }
            listA = listA->next;
            listB = listB->next;
        }

        // Me yaha tak tabhi pahuncha hu
        // jab listA and listB ki length equal nhi hai
        if(listA->next == NULL){...}

        if(listB->next == NULL){...}

        // Me yaha tak tabhi pahuncha hu jab headA and headB starting me right
        // Node par set ho chuke honge
        // We need to traverse again to get the intersection Node
        while(headA != headB){...}
        return headA;
    }
};
```

```
// Me yaha tak tabhi pahuncha hu
// jab listA and listB ki length equal nhi hai
if(listA->next == NULL){
    // ListB is bigger
    // We need to find the length of ListB

    int listBLength = 0;
    while(listB->next != 0){
        listBLength++;
        listB = listB->next;
    }

    // In starting, Set headB at right node
    // to get the intersection Node
    while(listBLength--){
        headB = headB->next;
    }
}
```

```
if(listB->next == NULL){
    // ListA is bigger
    // We need to find the length of ListA

    int listALength = 0;
    while(listA->next != 0){
        listALength++;
        listA = listA->next;
    }

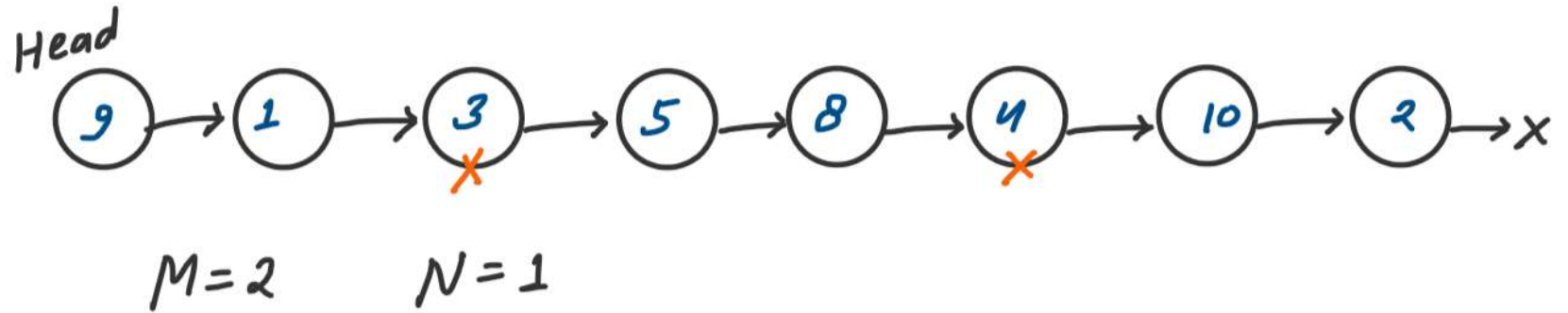
    // In starting, Set headA at right node
    // to get the intersection Node
    while(listALength--){
        headA = headA->next;
    }
}
```

```
while(headA != headB){
    headA = headA->next;
    headB = headB->next;
}
```

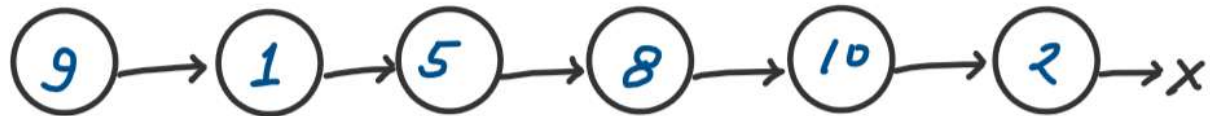
$T.C. \Rightarrow O(N)$
 $S.C. \Rightarrow O(1)$

HW 04: Delete N Nodes after M Nodes (GFG)

Ex 1

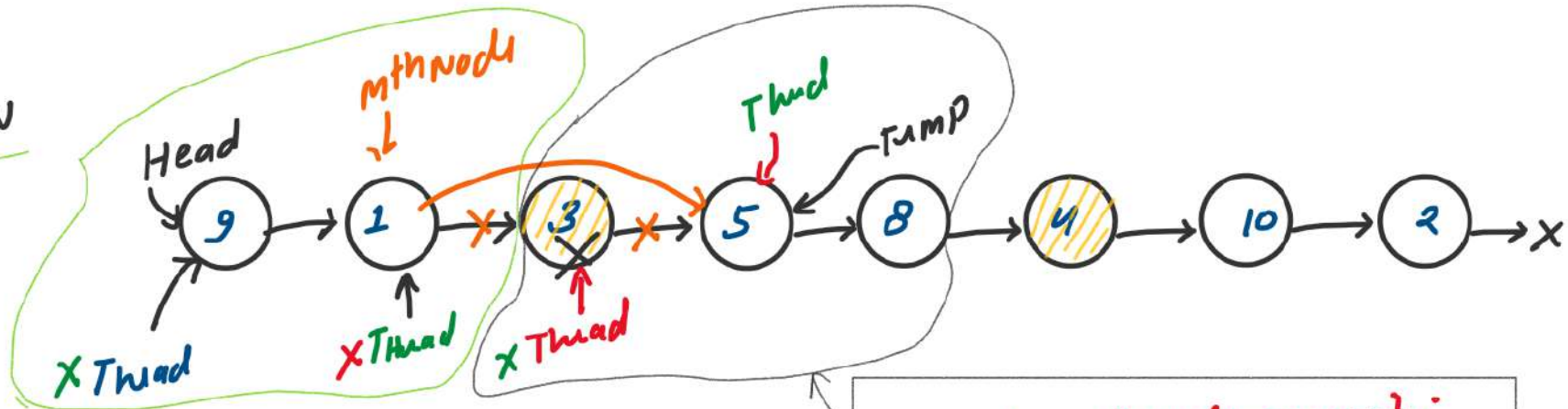


Output



DRY RUN

M=2
N=1



STEP 1

```

Node* THead = Head;
for(int i=0; i<M-1; i++)
{
    THead = THead->next;
}
Node* mthNode = THead;
    
```

STEP 3

fun(THead, M, N) Recursion call

STEP 2

```

THead = mthNode->next;
for(int i=0; i<N; i++)
{
    if(!THead) Break;
    Node* Tmp = THead->next;
    delete THead;
    THead = Tmp;
    mthNode->next = THead;
}
    
```

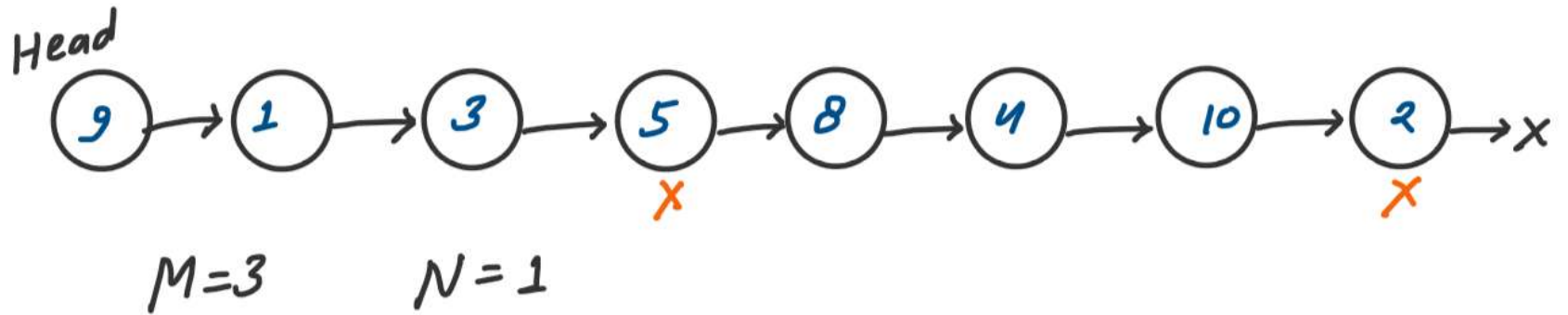
RUN TIME ERROR

→ Nth Node Available nahi Hai

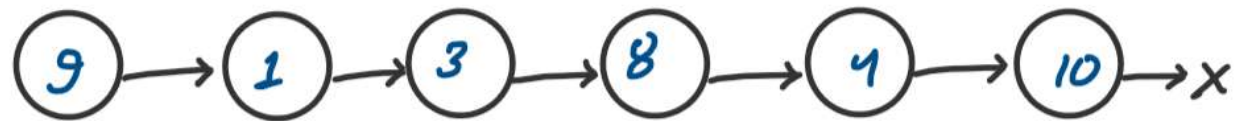
BASE CASE

if(!THead) return

Ex2



Output

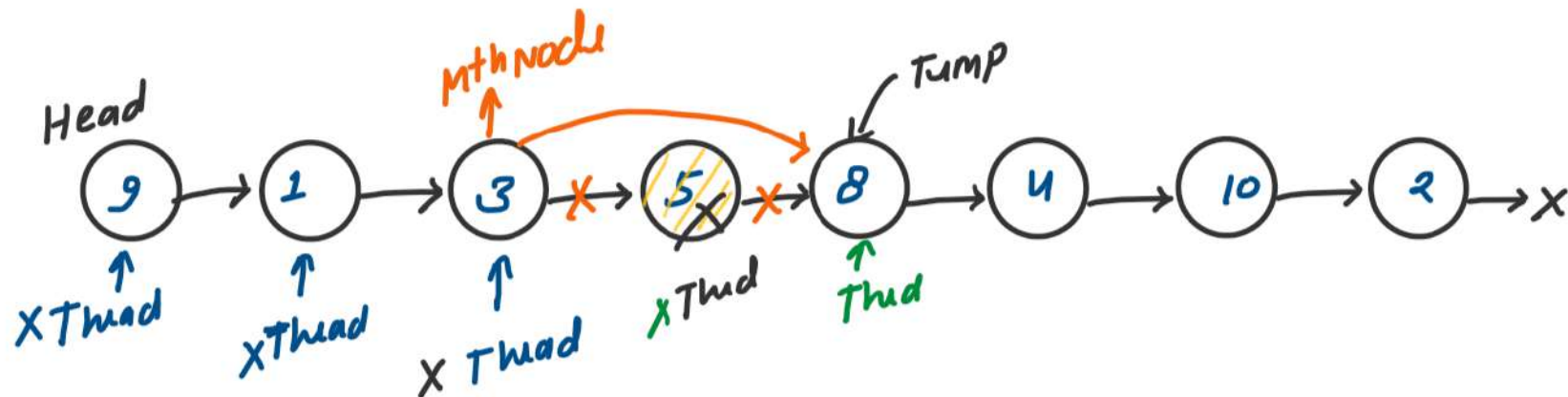


DRY RUN

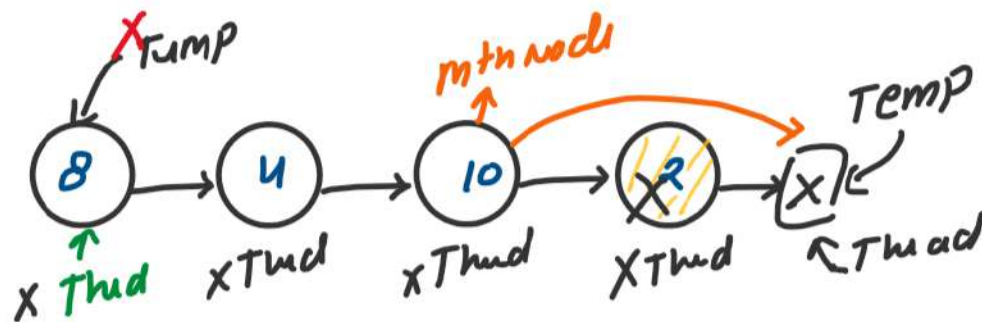
M=3

N=1

Iteration 1



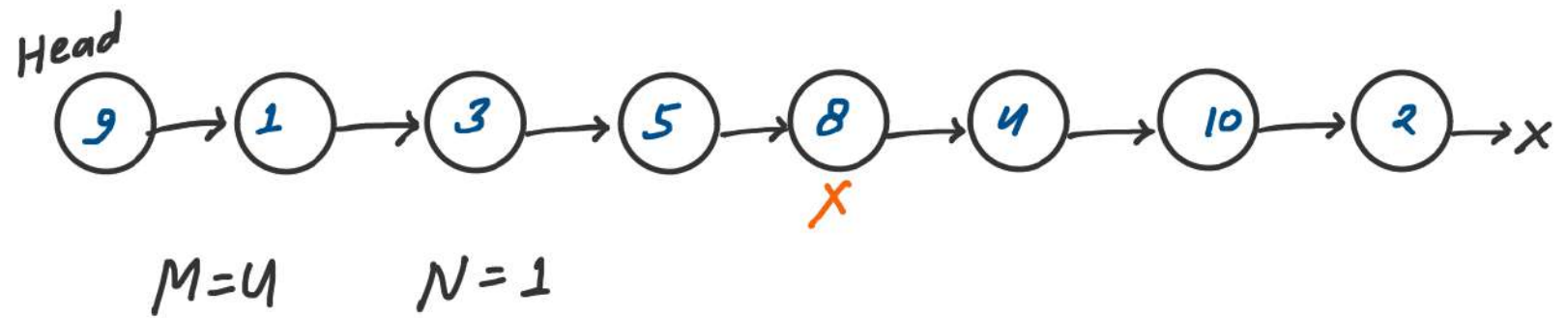
Iteration 2



Temp
[X]
Thud
if (!head) return

Output
(9) → (1) → (3) → (8) → (4) → (10) → (X)

Ex3



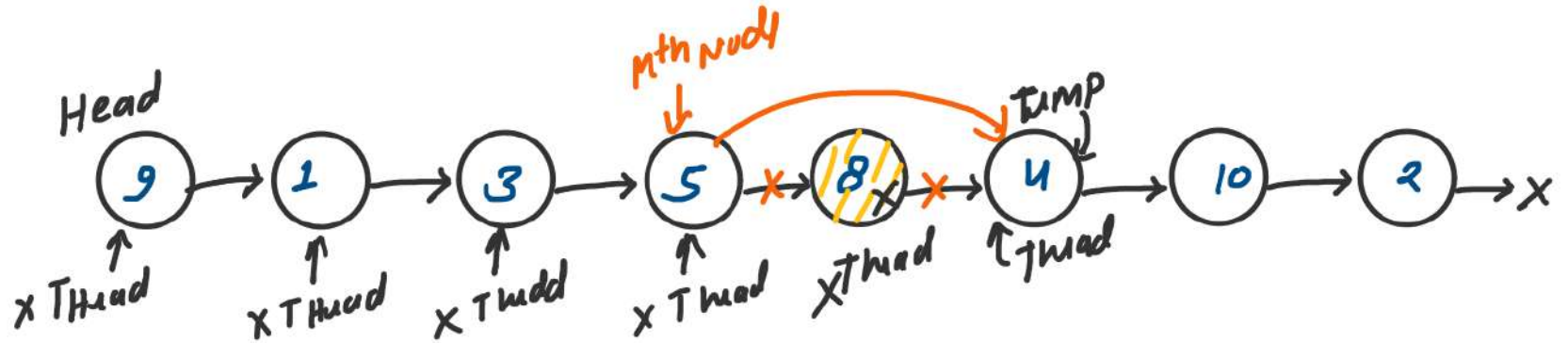
Output



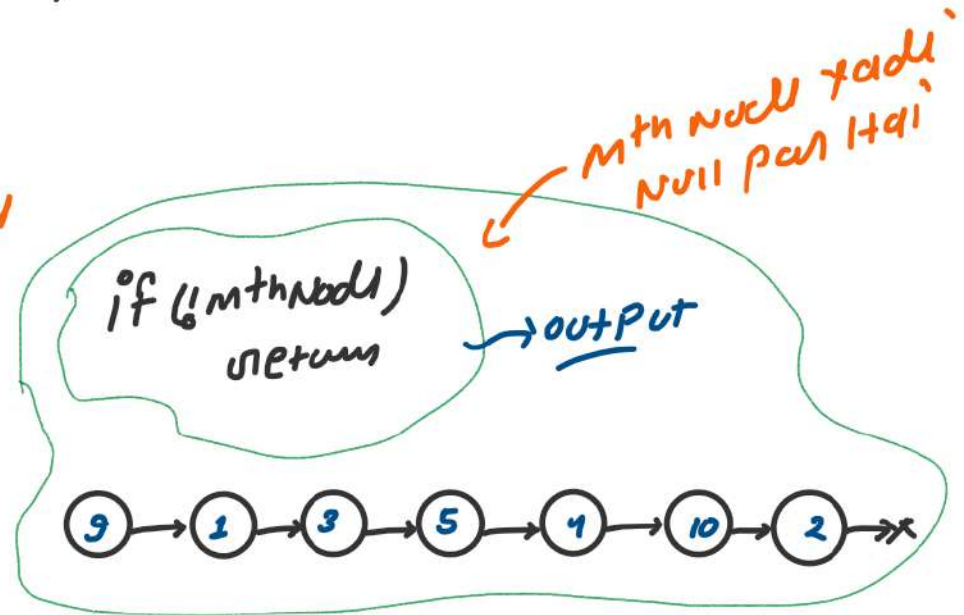
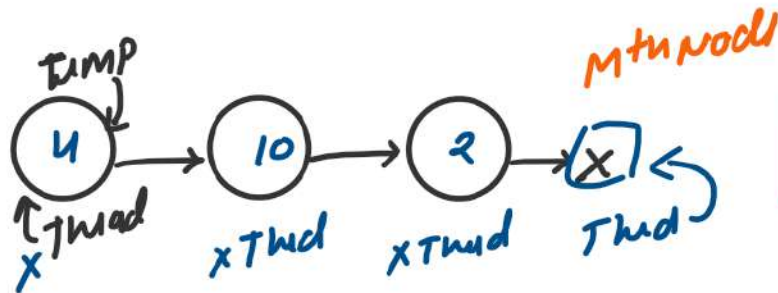
DRY RUN

M=4
N=1

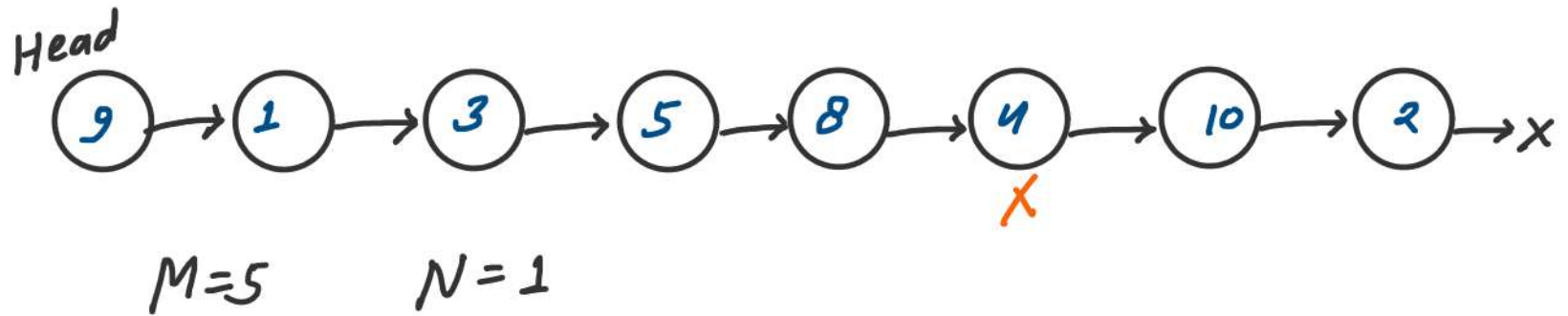
Iteration 1



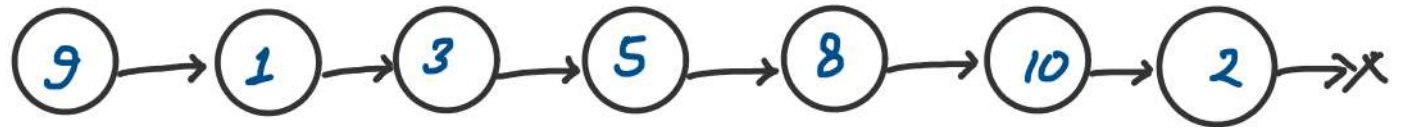
Iter 2



Ex 4



Output

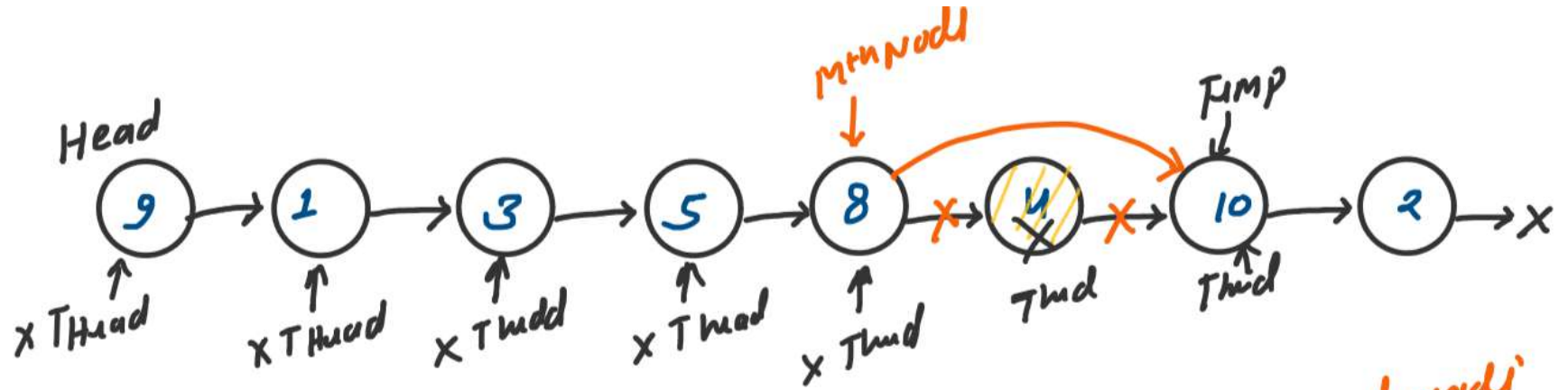


DRY RUN

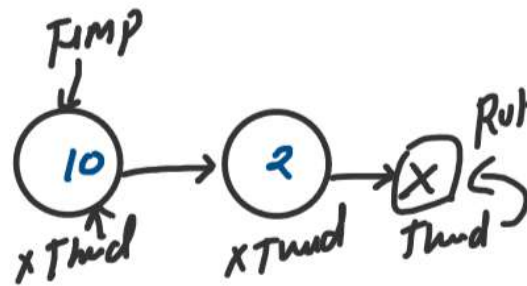
M=5

N=1

Iteration 1



Iteration 2



for (i=3 ---) {
if (!Ttmp) return
}

Ttmp head yadi
Null par hai

output



```

// HW 04: Delete N Nodes after M Nodes (GFG)
class Solution
{
public:
    void linkdelete(struct Node *head, int M, int N)
    {
        // Base case
        if(!head) return;

        // Step 1: Traverse list to M position from 0th to (M-1)
        Node* tempHead = head;
        for(int i=0; i<M-1; i++){
            // Temp Head yadi Null par hai
            if(!tempHead) return;
            tempHead = tempHead->next;
        }
        Node* MthNode = tempHead;

        // Mth Node yadi null par hai
        if(!MthNode) return;

        // Step 2: Delete N node
        tempHead = MthNode->next;
        for(int i=0; i<N; i++){
            // Nth node available nhi hai
            // matlab tempHead null hai
            if(!tempHead) break;

            Node* temp = tempHead->next;
            delete tempHead;
            tempHead = temp;
        }
        MthNode->next = tempHead;

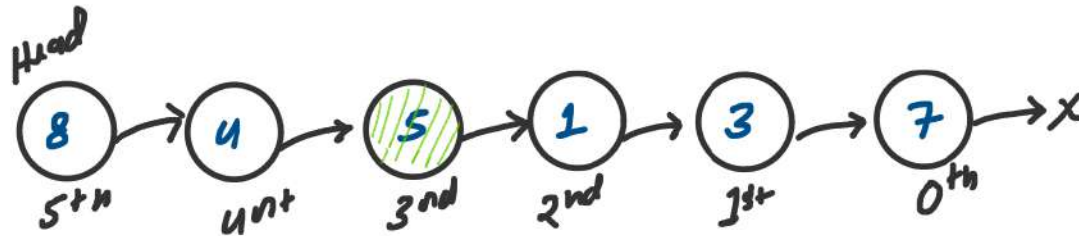
        // Recursive call
        linkdelete(tempHead, M, N);
    }
};

```

$T.C. \Rightarrow O(N)$
 Where N is number of nodes in the list.
 $S.C. \Rightarrow O(1)$

HW 05: Print kth Node from the End (Hacker Rank)

Ex 1



$$pos = kth = 3$$

Output = 5
3rd

RECURSIVE APPROACH

STEP: 1 TRAVERSE the list from Head to Tail

STEP: 2 Again TRAVERSE the list from Tail to [jab tak pos == 0]

STEP: 3 Return Ans = 5

DRY RUN
pos=3

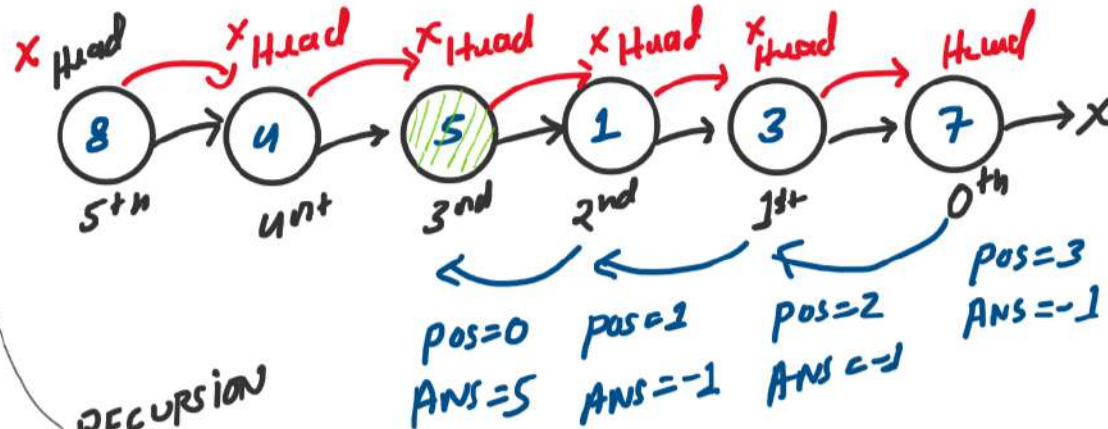
STEP 1

Base case
if (!Head)
return

fun(Head->next, pos, ans)

RECURSION

STEP: 3
return ans



STEP 2

if (pos == 0) {
ans = head->data;
}

pos--;

catch (Galti ki chances hai)

BACKTRACKING

```

// HW 05: Print kth Node from the End (Hackerrank)

/*
 * For your reference:
 *
 * SinglyLinkedListNode {
 *     int data;
 *     SinglyLinkedListNode* next;
 * };
 */

void solve(SinglyLinkedListNode* head, int &pos, int &ans){
    // Base case
    if(head == 0) return;

    // Step 1: traverse list from head to tail
    solve(head->next, pos, ans);

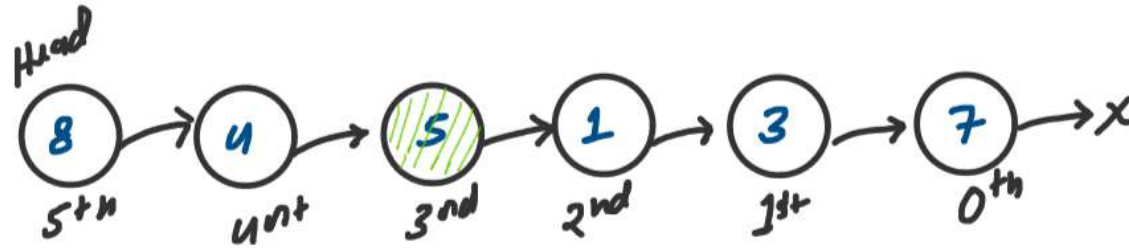
    // Step 2: traverse list from tail to (Jab tak pos == 0)
    if(pos == 0){
        ans = head->data;
    }
    --pos;
}

int getNode(SinglyLinkedListNode* llist, int positionFromTail) {
    // Step 3: return ans
    int ans = -1;
    solve(llist, positionFromTail, ans);
    return ans;
}

```

$T.C. \Rightarrow O(N)$
 Where N is number of nodes in the list.
 $S.C. \Rightarrow O(1)$

Ex 1



$pos = k^{th} = 3$

Output = 5
3rd

ITERATIVE APPROACH

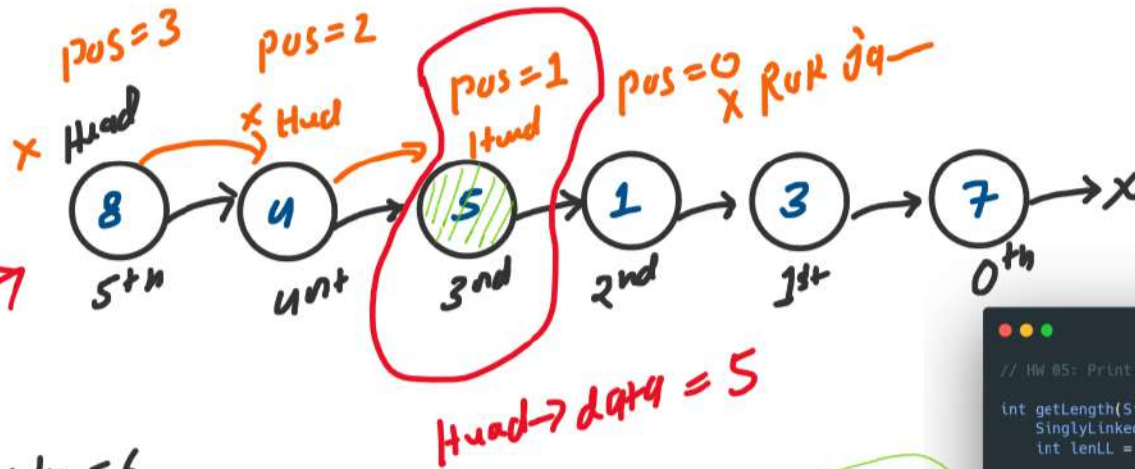
STEP:1 Get length of list

STEP:2 Subtract length - pos

STEP:3 Traverse list from head to [jab tak pos == 0].

DRY RUN

K=3



Step 1 length = 6

Step 2 $pos = length - K = 6 - 3 = 3$

Step 3

we have $1^{st} \text{ node} \rightarrow data$

T.C. $\Rightarrow O(N) + O(N)$
 $= O(N)$
 S.C. $\Rightarrow O(1)$

```
// HW 05: Print kth Node from the End (Hackerrank)

int getLength(SinglyLinkedListNode* head){
    SinglyLinkedListNode* temp = head;
    int lenLL = 0;

    // Base case
    if(head == 0) return lenLL;

    while (temp->next != 0) {
        lenLL++;
        temp = temp->next;
    }

    return lenLL;
}

int getNode(SinglyLinkedListNode* llist, int positionFromTail) {
    // Step 1: get length of list
    int length = getLength(llist);

    // Step 2: subtract pos from length
    int posFromHead = length - positionFromTail;

    // Step 3: traverse list from head to (jab tak pos==0)
    while(posFromHead != 0){
        llist = llist->next;
        posFromHead--;
    }

    return llist->data;
}
```

$\rightarrow O(N)$

$\rightarrow O(N)$

HW 06: Flatten Linked List (GFG)

PROBLEM STATEMENT:

Given a Linked List of size N , where every node represents a **sub-linked-list** and contains two pointers:

(I) a **next** pointer to the next node,

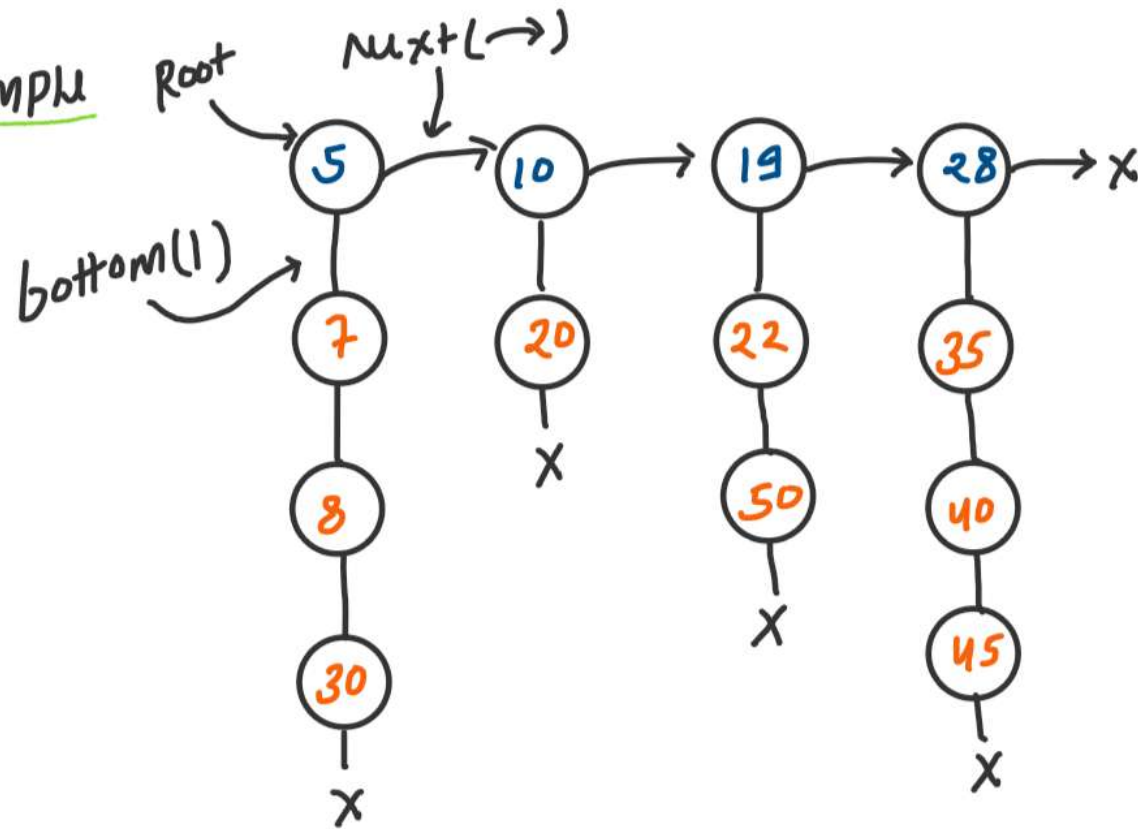
(II) a **bottom** pointer to a linked list where this node is head.

Each of the sub-linked-list is in sorted order.

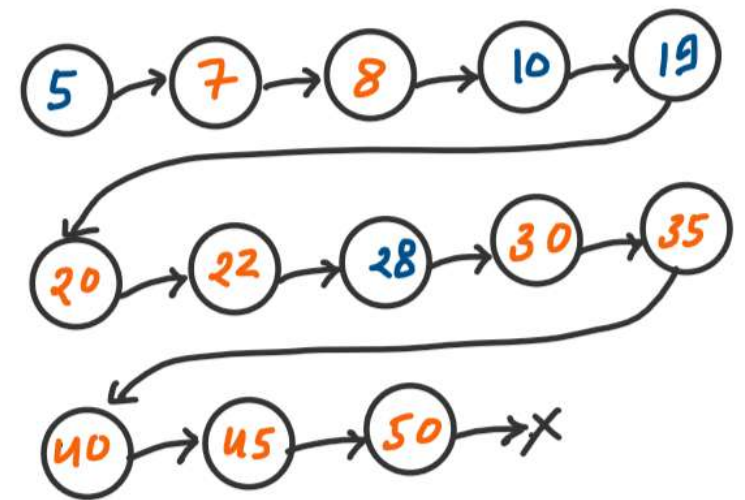
Flatten the Link List such that all the nodes appear in a single level while maintaining the sorted order.

Note: The flattened list will be printed using the bottom pointer instead of the next pointer.

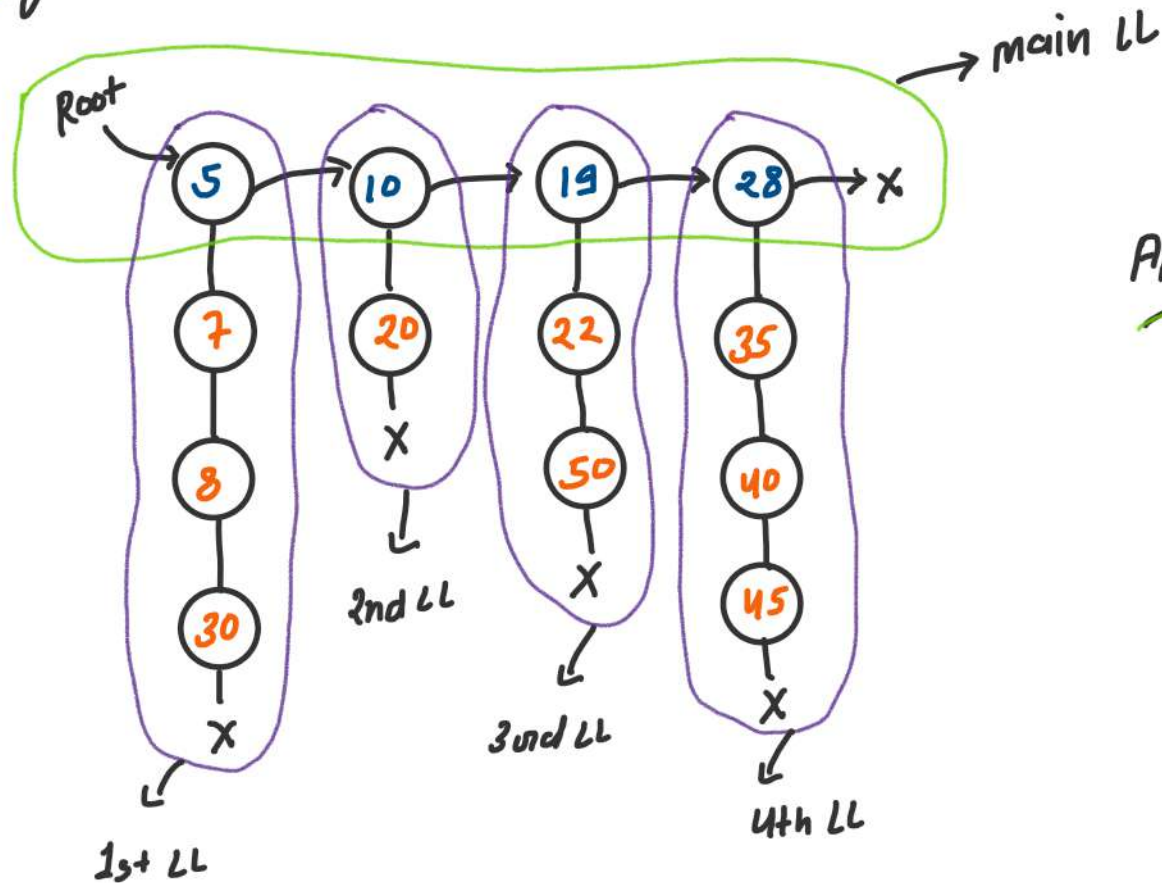
EXAMPLE



Output



Build Logic



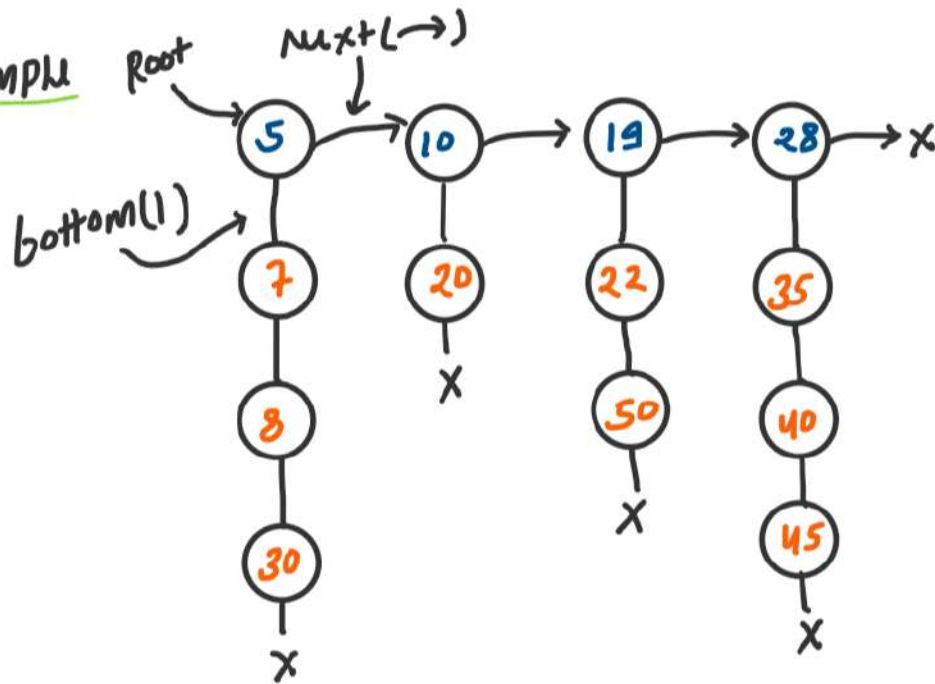
Approach: Recursive Approach

Step 1 merge two sorted LL

Step 2 merge next sorted LL with already two merged LL

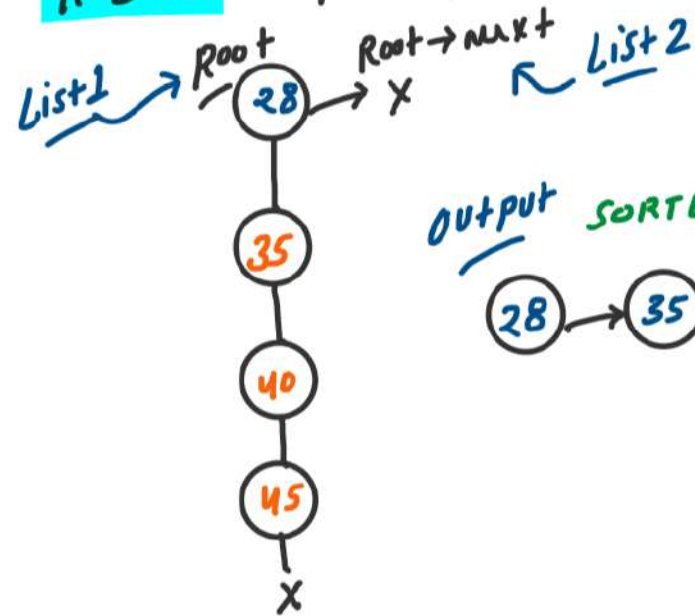
DRY RUN

EXAMPLE



R. Call 1

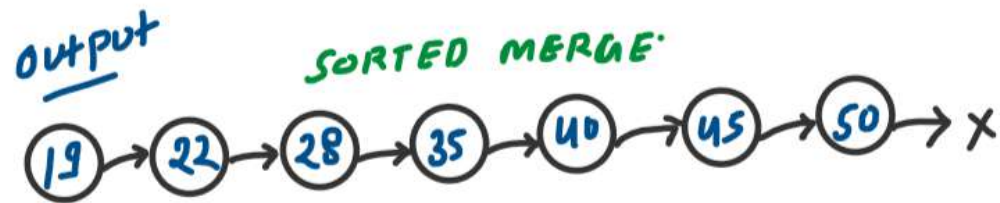
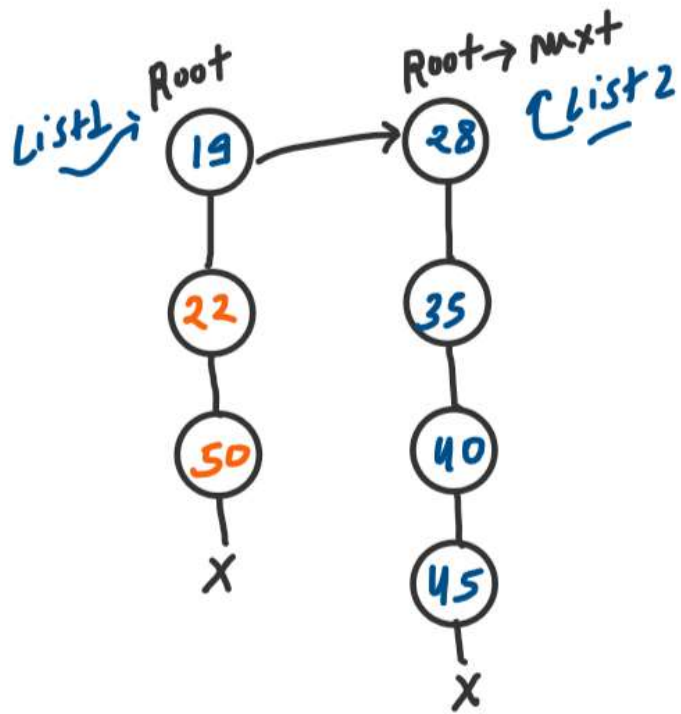
f(L 28, NULL)



Output SORTED MERGE

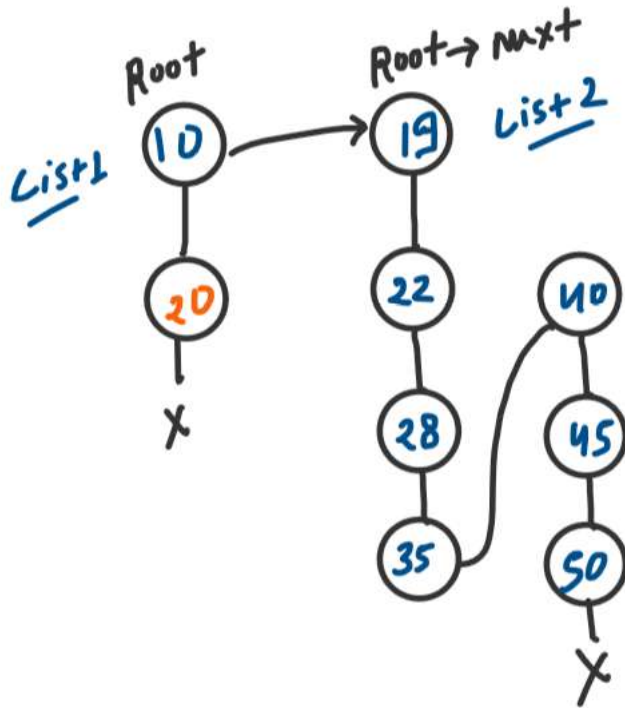


R = Call 2 f(19, 28)



R-Call 3

f(10, 19)

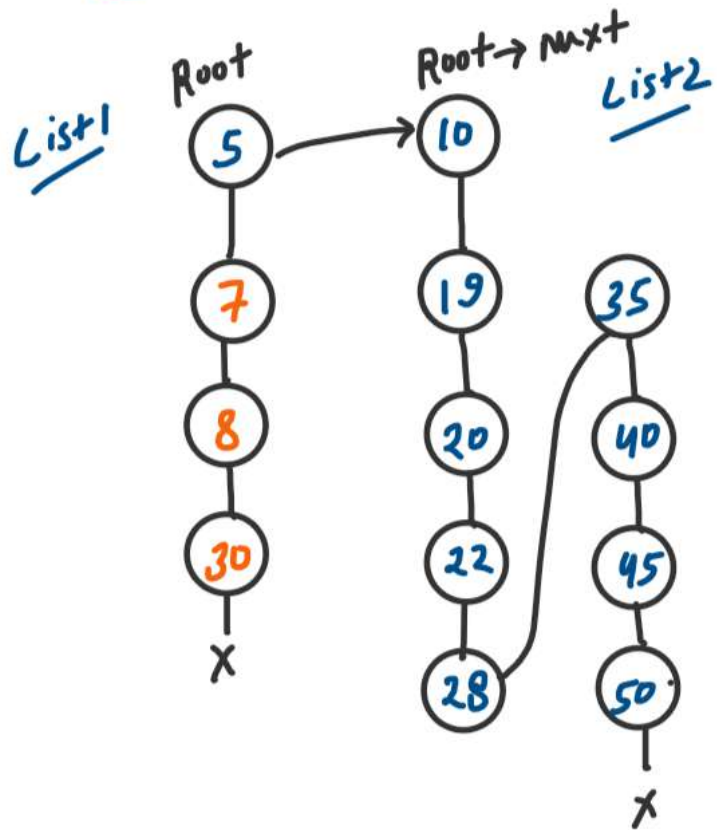


Output

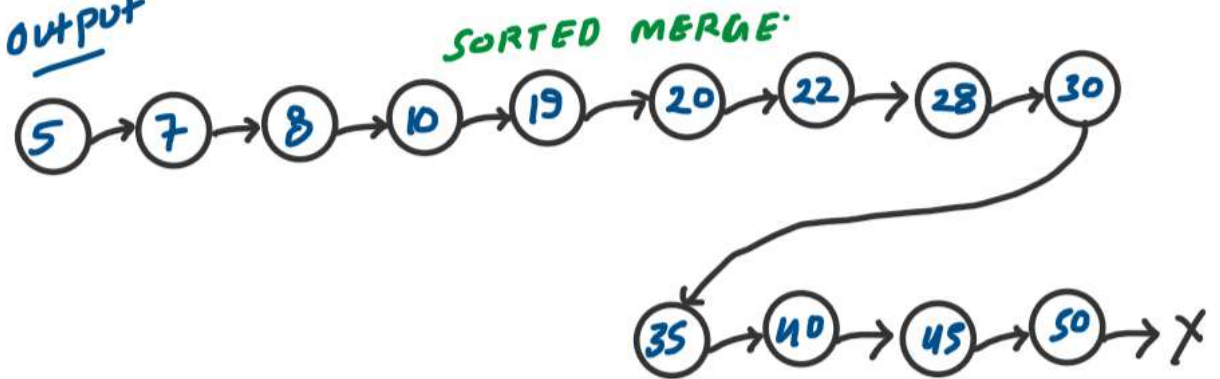
SORTED MERGE



R-Callu f(5, 10)



Final output



ALGORITHM

Main LL → First Two Node pick

Root
List 1

&

Root → next
List 2

MERGE
with Bottom
in single sorted list

```

/* Node structure used in the program
struct Node{
    int data;
    struct Node * next;
    struct Node * bottom;

    Node(int x){
        data = x;
        next = NULL;
        bottom = NULL;
    }
};
*/

```

```

Node* mergeTwoSortedLL(Node* list1, Node* list2){...}

```

```

Node *flatten(Node *root)
{
    // Base case
    if(root == NULL) return NULL;

    Node* mergedLL = mergeTwoSortedLL(root, flatten(root->next));
    return mergedLL;
}

```

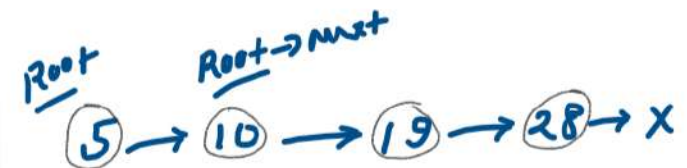
```

Node* mergeTwoSortedLL(Node* list1, Node* list2){
    // Base case
    if(list1 == NULL) return list2;
    if(list2 == NULL) return list1;

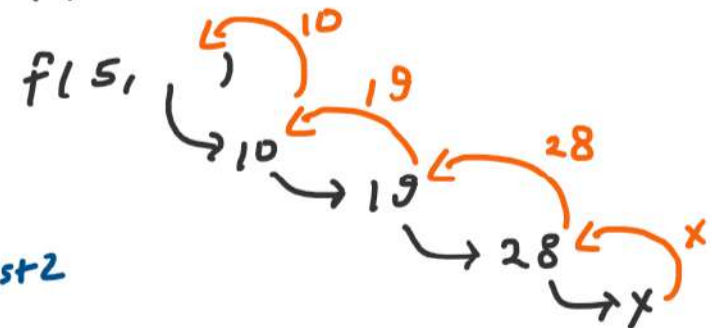
    // Processing
    Node* ans = NULL;
    if(list1->data < list2->data){
        ans = list1;
        list1->bottom = mergeTwoSortedLL(list1->bottom, list2);
    }
    else{
        ans = list2;
        list2->bottom = mergeTwoSortedLL(list1, list2->bottom);
    }

    return ans;
}

```



$f(\text{root}, \text{flat}(\text{root} \rightarrow \text{next}))$

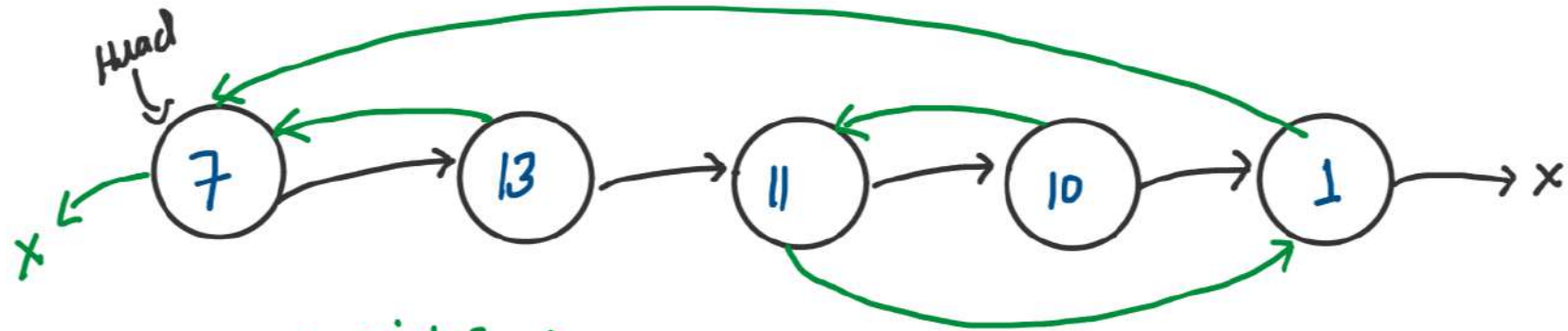


list1 → list2

1st call $f(28, X)$
 2nd call $f(19, 28)$
 3rd call $f(10, 19)$
 4th call $f(5, 10)$

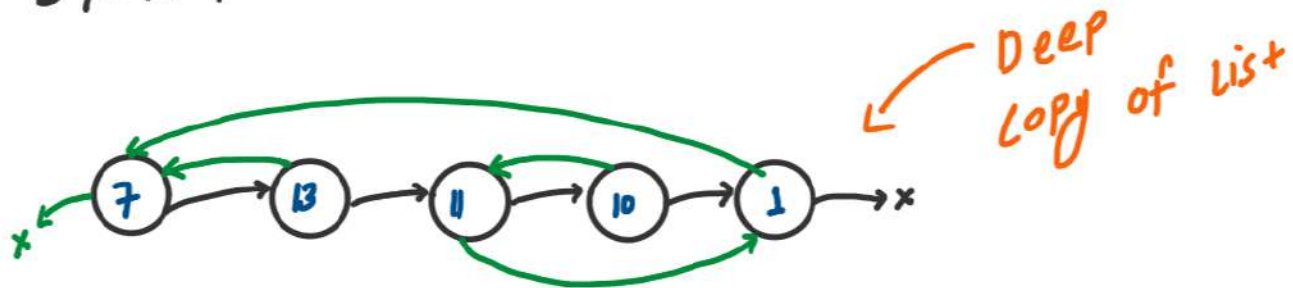
HW 07: Copy List with Random Pointer (Leetcode-138)

Input



- Random pointer →
- Next pointer →

Output

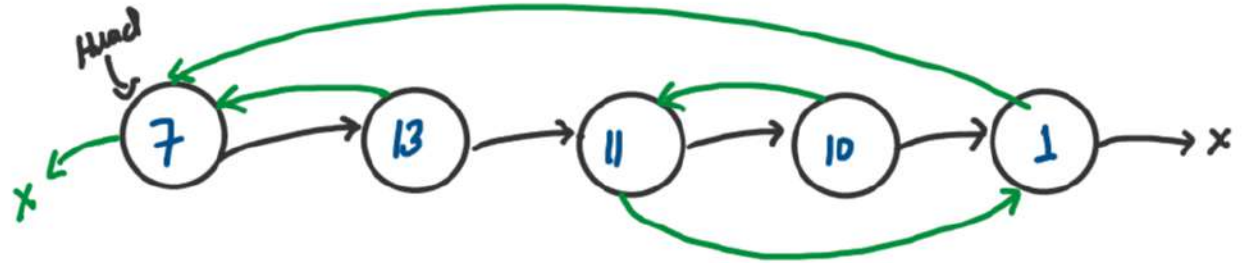


Approach 1 Using map

Step 1 Copy List in map
using Recursion

MP

old ptr	new ptr
7	7
13	13
11	11
10	10
1	1



Base case

if (!head) return null;

Node* newNode = new Node(head->val);

MP[head] = newhead;

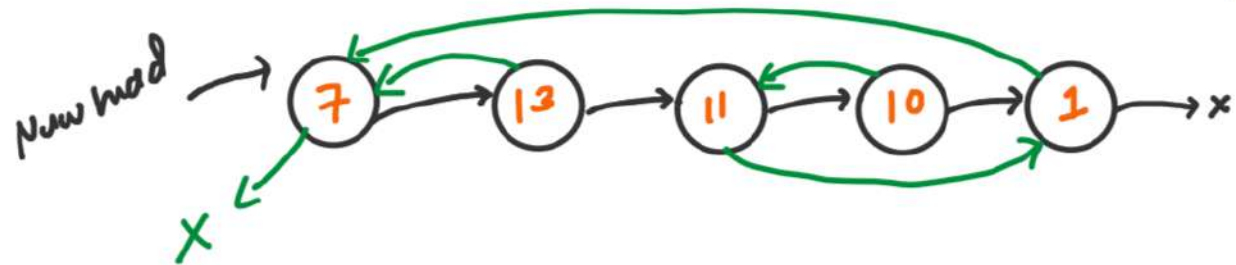
newNode->next = f(head->next, MP);



Step 2 Allocate the Random
pointer

**
Map[old ptr] = new pointer
key of map value of map

```
if (head → Random) {  
    newhead → Random = mp [head → Random];  
}
```



Return newhead output


```
// HW 07: Copy List with Random Pointer (Leetcode-138)
```

```
/*  
Definition for a Node.
```

```
class Node {  
public:  
    int val;  
    Node* next;  
    Node* random;
```

```
    Node(int _val) {  
        val = _val;  
        next = NULL;  
        random = NULL;  
    }  
};
```

```
*/
```

```
class Solution {
```

```
public:
```

```
    Node* solve(Node* head, unordered_map<Node*, Node*> &mp){
```

```
        // Base case  
        if(!head) return NULL;
```

```
        // Step 1: Copy list in map  
        Node* newHead = new Node(head->val);  
        mp[head] = newHead;  
        newHead->next = solve(head->next, mp);
```

```
        // Step 2: Allocate the random pointer  
        if(head->random){  
            newHead->random = mp[head->random];  
        }
```

```
        return newHead;
```

```
    }  
  
    Node* copyRandomList(Node* head) {  
        unordered_map<Node*, Node*> mp;  
        return solve(head, mp);  
    }
```

```
};
```

Time Complexity: $O(N)$,
Where N is number of nodes in list

Space Complexity: $O(N)$,
where N is number of elements (Nodes) stored in map

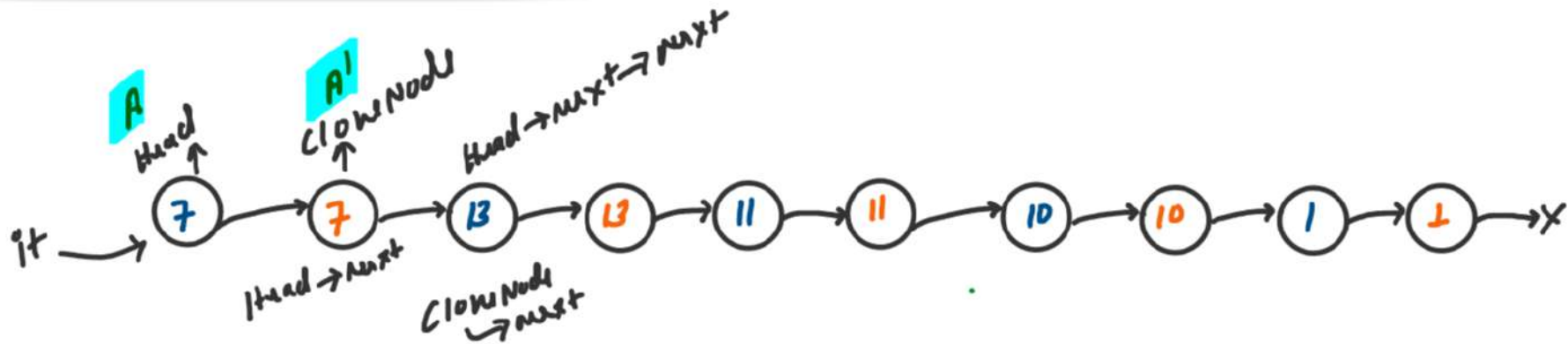
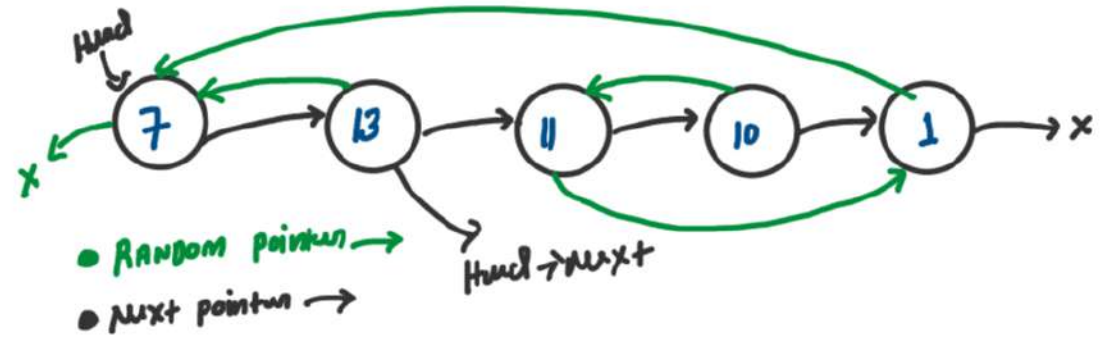
Approach 2 without using map

Step: 1

```
Node* solve(Node* head){
    if(!head) return NULL;

    // Step 1: Clone A->A'
    Node* it = head; // Iterating Over Old Head
    while(it){
        Node* cloneNode = new Node(it->val);
        cloneNode->next = it->next;
        it->next = cloneNode;
        it = cloneNode->next;
    }
}
```

Input



step 2

Assign Random pointer A'
with the help of A

$$i_+ = \text{mod } i$$

whim (it) ε

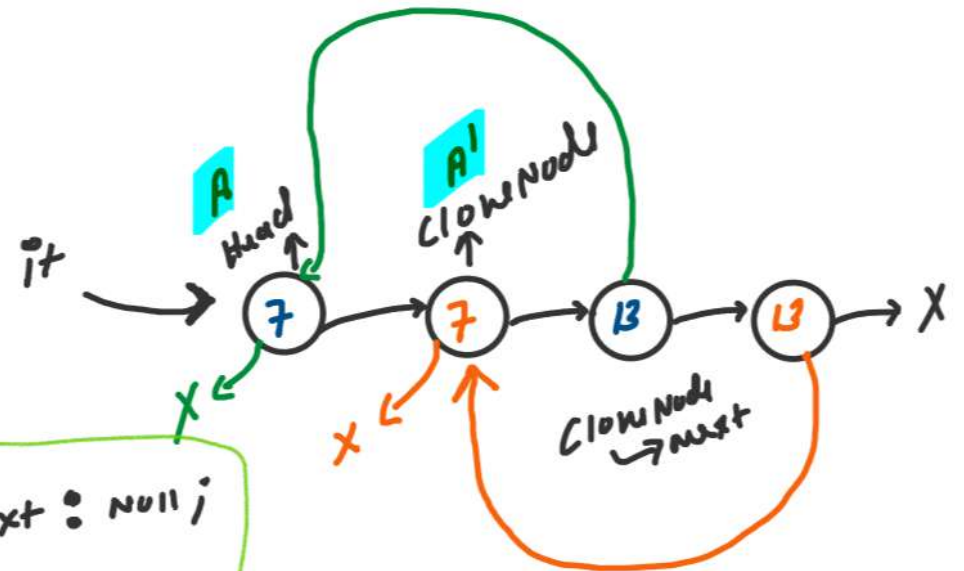
Node* cloneNode = it->next;

Clone Node \rightarrow Random =

it → Random ? it → Random → next : null;

$$i_t = i_t \rightarrow \text{next} \rightarrow \text{next} i$$

3

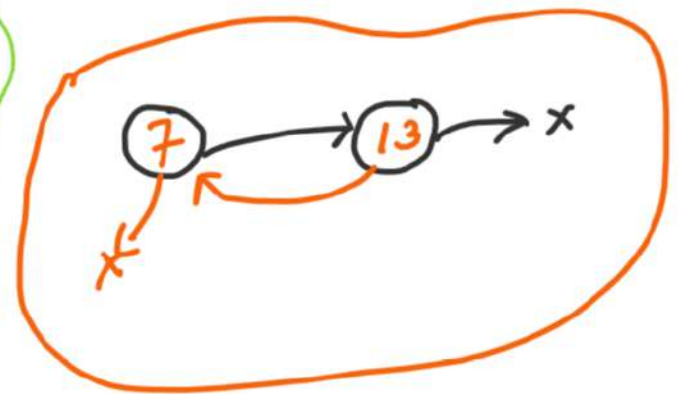
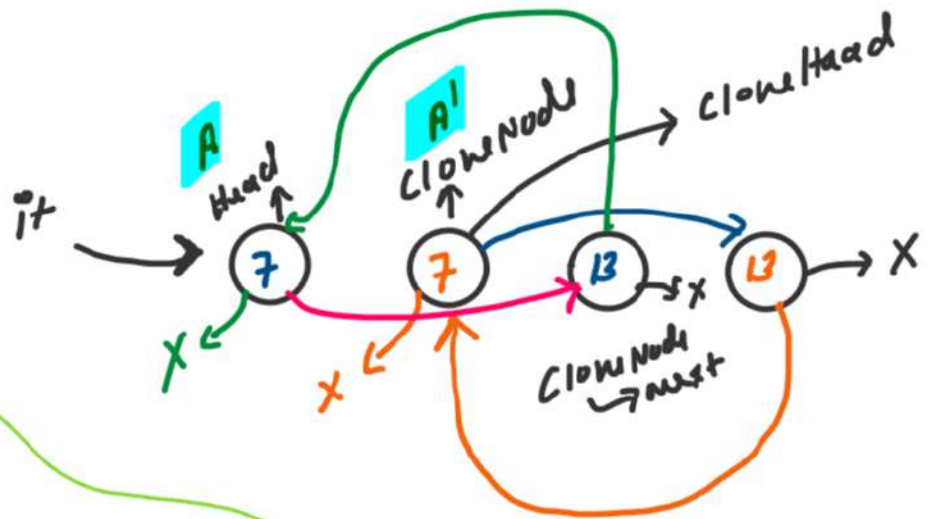


STEP 3 Detach A' from A

```

it = head;
Node* cloneHead = it->next;
while (it) {
    Node* cloneNode = it->next;
    it->next = cloneNode->next;
    if (cloneNode->next) {
        cloneNode->next = cloneNode->next->next;
    }
    it = it->next;
}
return cloneHead;

```



```

class Solution {
public:
    Node* solve(Node* head){
        if(!head) return NULL;

        // Step 1: Clone A->A'
        Node* it = head; // Iterating Over Old Head
        while(it){
            Node* cloneNode = new Node(it->val);
            cloneNode->next = it->next;
            it->next = cloneNode;
            it = cloneNode->next;
        }

        // Step 2: Assign random pointer of A' with the help of A
        it = head;
        while(it){
            Node* cloneNode = it->next;
            cloneNode->random = it->random ? it->random->next : NULL;
            it = cloneNode->next;
        }

        // Step 3: Detach A' from A

    }

    Node* copyRandomList(Node* head) {
        return solve(head);
    }
};

```

Step 3

```

it = head;

// cloneHead is not changed after its initial assignment
Node* cloneHead = it->next;

while(it){
    Node* cloneNode = it->next;
    it->next = cloneNode->next;
    if(cloneNode->next){
        cloneNode->next = cloneNode->next->next;
    }
    it = it->next;
}

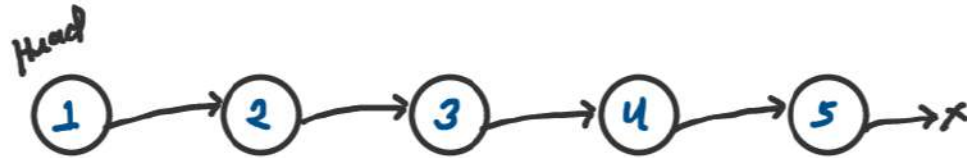
return cloneHead;

```

$T.C. \Rightarrow O(N)$
 $S.C. \Rightarrow O(1)$

HW 08: Rotate List (Leetcode-61)

Ex



$K=1$ $5 \rightarrow 1 \rightarrow 2 \rightarrow 3 \rightarrow 4 \rightarrow x$

$K=2$ $4 \rightarrow 5 \rightarrow 1 \rightarrow 2 \rightarrow 3 \rightarrow x$

$K=3$ $3 \rightarrow 4 \rightarrow 5 \rightarrow 1 \rightarrow 2 \rightarrow x$

$K=4$ $2 \rightarrow 3 \rightarrow 4 \rightarrow 5 \rightarrow 1 \rightarrow x$

$K=5$ $1 \rightarrow 2 \rightarrow 3 \rightarrow 4 \rightarrow 5 \rightarrow x$

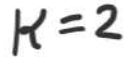
$K=6$

$5 \rightarrow 1 \rightarrow 2 \rightarrow 3 \rightarrow 4 \rightarrow x$

$K=7$

$4 \rightarrow 5 \rightarrow 1 \rightarrow 2 \rightarrow 3 \rightarrow x$

Input


$$\text{dist length} < 5$$
$$\begin{aligned}\text{Actual Rotation} &= K \% \text{ length} \\ &= 2 \% 5 \\ &= 2\end{aligned}$$

3
 new last node posi = $Mn - \text{Actual Rotor} - 1$
 $= 5 - 2 - 1$
 $= 2$

- Step 1 Find length of List
- Step 2 Find Actual Rotation of K
- Step 3 Find new last node position

```
// HW 08: Rotate List (Leetcode-61)

/**
 * Definition for singly-linked list.
 * struct ListNode {
 *     int val;
 *     ListNode *next;
 *     ListNode() : val(0), next(nullptr) {}
 *     ListNode(int x) : val(x), next(nullptr) {}
 *     ListNode(int x, ListNode *next) : val(x), next(next) {}
 * };
 */
class Solution {
public:
    int getLength(ListNode* head){...}

    ListNode* rotateRight(ListNode* head, int k) {
        // Corner Case
        if(!head) return NULL;

        // Step 1: Find length of list
        int len = getLength(head);

        // Step 2: Find actual rotation of k
        int actualRotateK = k % len;

        // Corner case
        if(actualRotateK == 0) return head;

        // Step 3: Find position of lastNewNode
    }
};
```

```
int getLength(ListNode* head){
    ListNode* temp = head;
    int len = 0;

    while(temp){
        len++;
        temp = temp->next;
    }

    return len;
}
```

$T.C. \Rightarrow O(N)$
 $S.C. \Rightarrow O(1)$

```
int newLastNodePos = len - actualRotateK - 1;

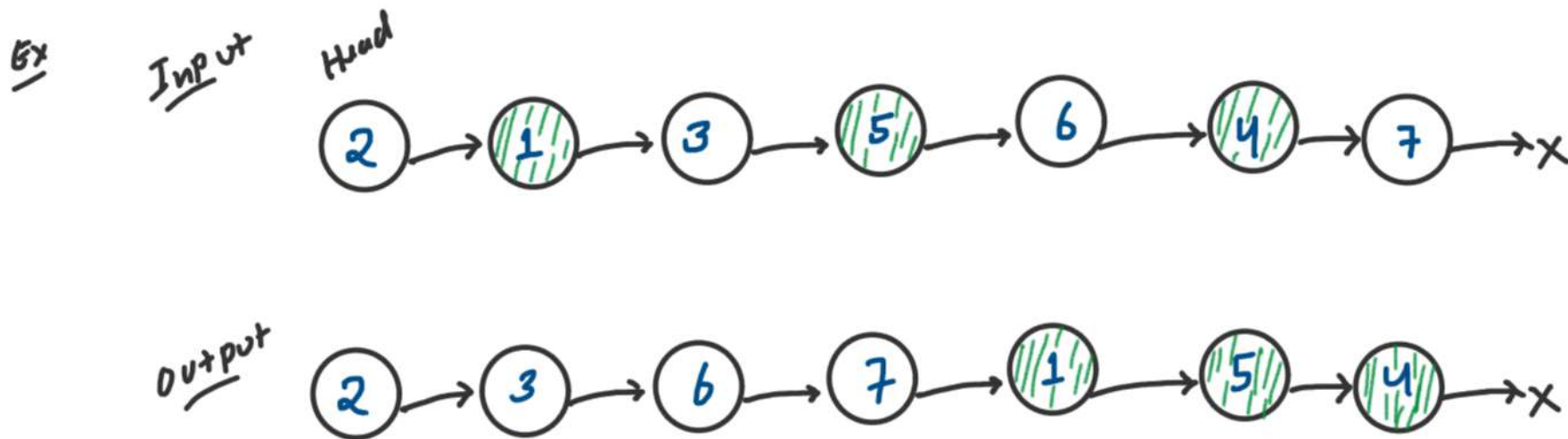
ListNode* newLastNode = head;
for(int i=0; i<newLastNodePos; i++){
    newLastNode = newLastNode->next;
}

// Save newLastNode->next in newHead to track
ListNode* newHead = newLastNode->next;
newLastNode->next = NULL;

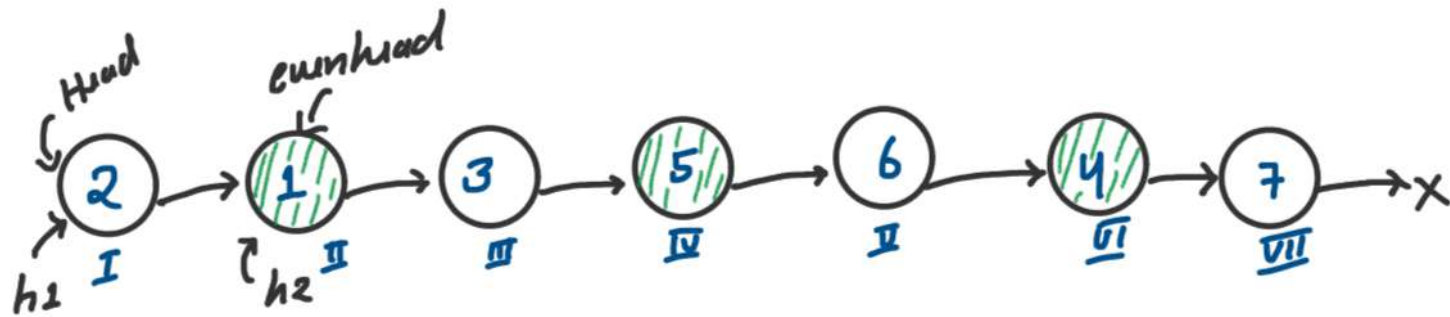
// newHead ka next node yadi null ho jata hai
// to use old Head se meet kara do
ListNode* it = newHead;
while(it->next != NULL){
    it = it->next;
}
it->next = head;

return newHead;
```

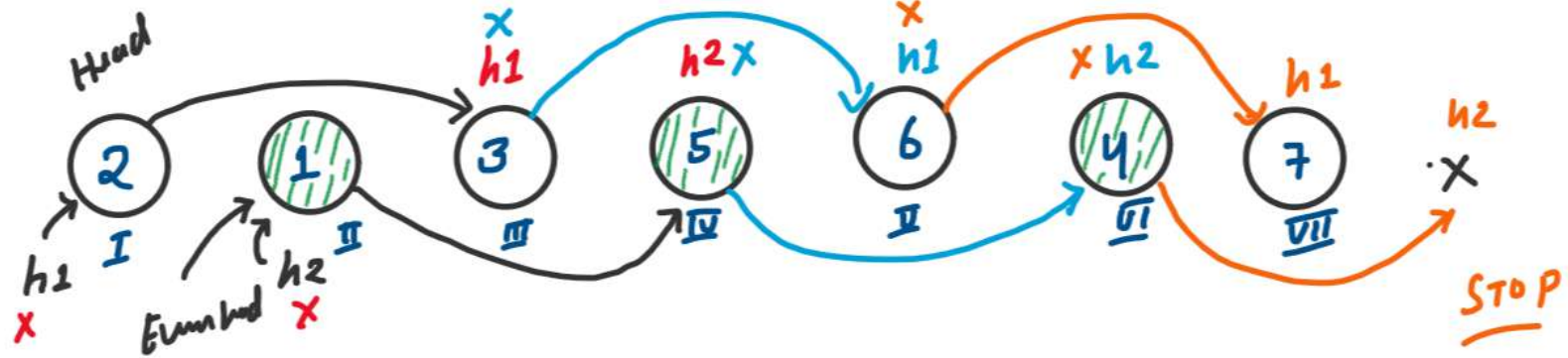
HW 09: Odd Even Linked List (Leetcode-328)



Logic build



Node * h1 = head; ← odd Indexed list
Node * h2 = head → next; ← Even Indexed list
Node * Evenhead = h2; ← save the evenhead to link the odd list

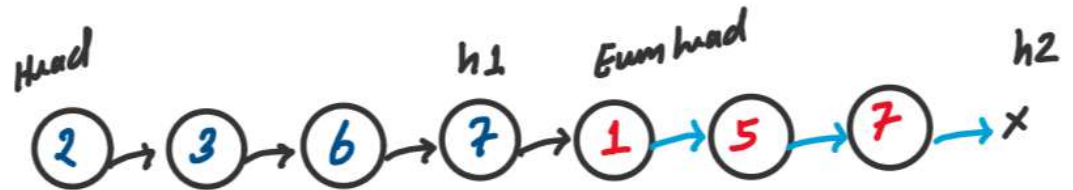


$h1 \rightarrow \text{next} = h2 \rightarrow \text{next};$
 $h2 \rightarrow \text{next} = h2 \rightarrow \text{next} \rightarrow \text{next};$

$h1 = h1 \rightarrow \text{next};$

$h2 = h2 \rightarrow \text{next};$

$h2 == \text{NULL} \text{ \&\& }$
 $h2 \rightarrow \text{next} == \text{NULL}$



Re-group the both list

$\{$
 $h1 \rightarrow \text{next} = \text{evenhead};$
 $\text{return head};$
 $\}$

```

// HW 09: Odd Even Linked List (Leetcode-328)

class Solution {
public:
    ListNode* oddEvenList(ListNode* head) {
        if(head == NULL || head->next == NULL) return head;
        // Odd indexed list
        ListNode* h1 = head;
        // Even indexed list
        ListNode* h2 = head->next;
        // Save h2 for attaching the odd index list
        ListNode* evenHead = h2;

        while(h2 && h2->next){
            h1->next = h2->next;
            h2->next = h2->next->next;

            h1 = h1->next;
            h2 = h2->next;
        }

        // Odd and even indexed list ko regroup krdo
        h1->next = evenHead;
        return head;
    }
};

```

Time complexity: $O(N)$,
Where N is number of nodes of the Linked List

Space complexity: $O(1)$,
Where no extra space used

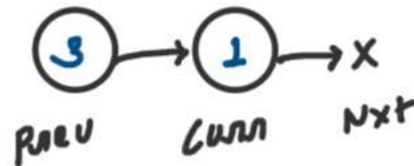
HW 10: Find Minimum and Maximum Number of Nodes
Between Critical Points (Leetcode-2048)

Ex 1

Input



head



Doj Run

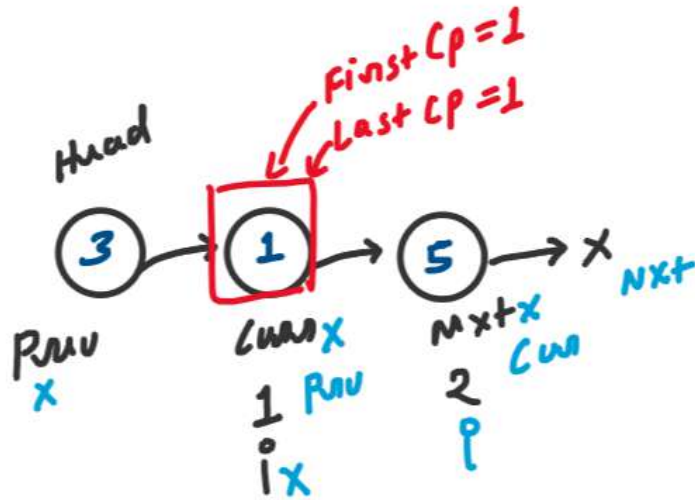
NO critical point exist
Output [-1, -1]

```
vector<int> ans = {-1, -1};  
Node* prev = head;  
if (!prev) return ans;  
Node* curr = head->next;  
if (!curr) return ans;  
Node* next = curr->next;  
if (!next) return ans;
```

Ex2

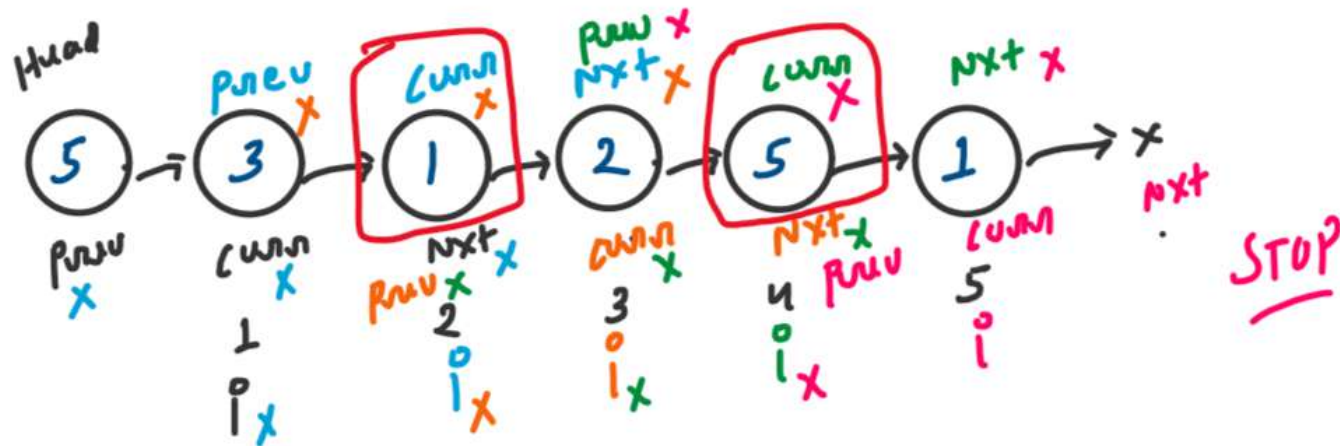


Dry Run



One critical point exist
 $1 < 3$ & $1 < 5$
Output → [-1, -1]

Ex:3



$$\text{minDis} = i - \text{lastCP} \Rightarrow 4 - 2 \Rightarrow 2$$

$$\text{firstCP} = -1 \quad 2 \quad 2$$

$$\text{lastCP} = -1 \quad 2 \quad 4$$

$$\text{maxDis} = \text{lastCP} - \text{firstCP}$$

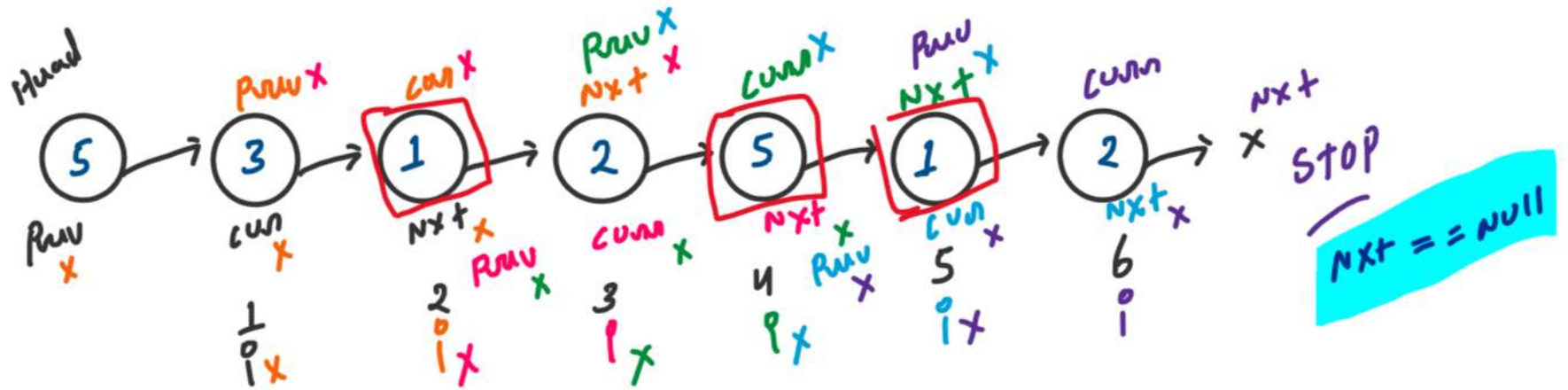
$$= 4 - 2$$

$$= 2$$

Two critical points exist

$$\text{O/P} \Rightarrow [2, 2]$$

Ex: 4



First CP = ~~1~~ 2
 Last CP = ~~1~~ 2 4 5
 minDis = i - last CP $\Rightarrow 4 - 2 = 2$
 maxDis \Rightarrow last CP = first CP
 $\Rightarrow 5 - 2 = 3$

Three critical points exist
 output $\Rightarrow [1, 3]$

```
// HW 10: Find Minimum and Maximum Number of Nodes Between  
Critical Points (Leetcode-2048)
```

```
class Solution {  
public:  
    vector<int> nodesBetweenCriticalPoints(ListNode* head) {  
        // vector initialized with minDis and maxDis  
        vector<int> ans = {-1, -1};  
        ListNode* prev = head;  
        if(!prev) return ans;  
        ListNode* curr = prev->next;  
        if(!curr) return ans;  
        ListNode* nxt = curr->next;  
        if(!nxt) return ans;  
  
        int firstCP = -1;  
        int lastCP = -1;  
        int i = 1;  
        int minDis = INT_MAX;  
  
        while(nxt){...}  
  
        // Only one critical point found condition  
        if(lastCP == firstCP){  
            return ans;  
        }  
        else {  
            ans[0] = minDis;  
            ans[1] = lastCP - firstCP;  
        }  
        return ans;  
    }  
};
```

```
while(nxt){  
    bool isCP = ((curr->val > prev->val && curr->val > nxt->val)  
                || (curr->val < prev->val && curr->val < nxt->val))  
                ? true : false;  
  
    // first critical point condition  
    if(isCP && firstCP == -1){  
        firstCP = i;  
        lastCP = i;  
    }  
    // at least one critical point ke liye condition  
    else if(isCP){  
        minDis = min(minDis, i - lastCP);  
        lastCP = i;  
    }  
    i++;  
    prev = prev->next;  
    curr = curr->next;  
    nxt = nxt->next;  
}
```

Time complexity: $O(N)$,
Where N is number of nodes of the Linked List

Space complexity: $O(1)$,
Where no extra space used

HW 11: Merge Nodes in between Zeros (Leetcode-2181)

Example 1:

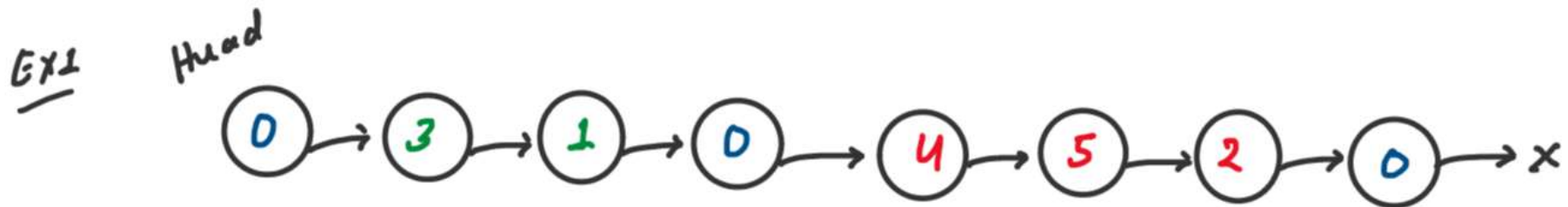
Input: head = [0,3,1,0,4,5,2,0]

Output: [4,11]

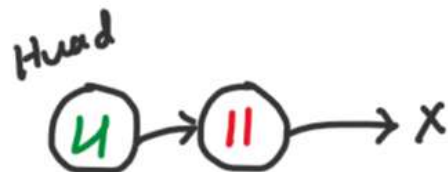
Example 2:

Input: head = [0,1,0,3,0,2,2,0]

Output: [1,3,4]

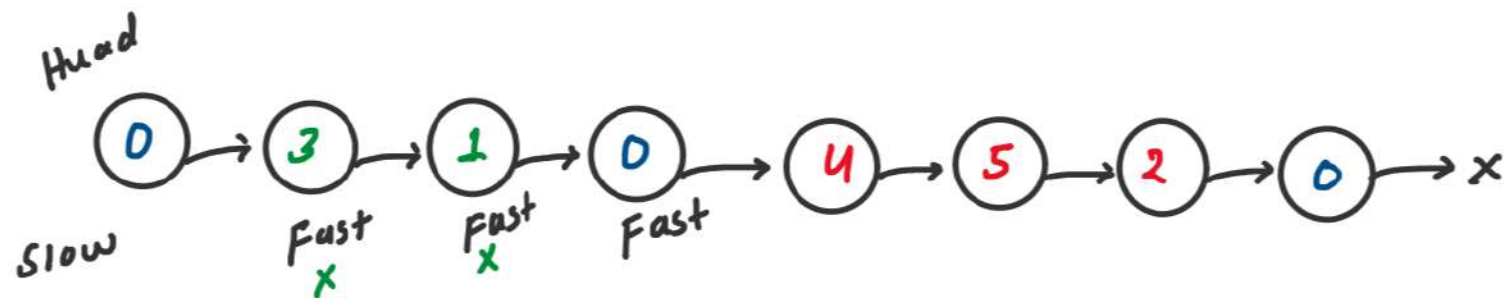


Explain



Output [3+1, 4+5+2]
[4, 11]

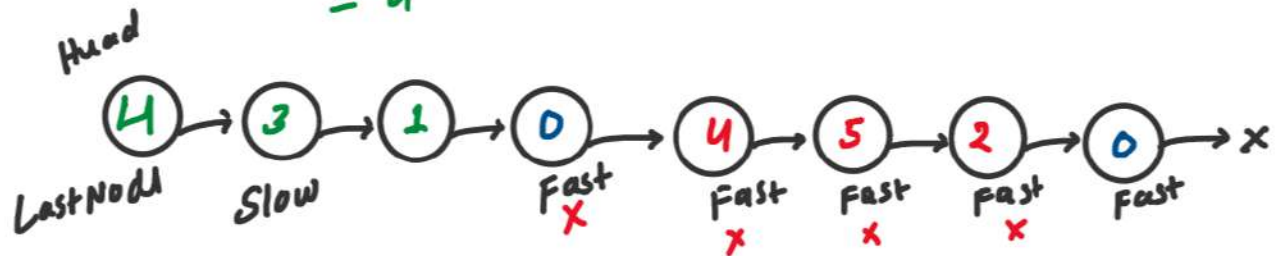
DRY RUN



Node* slow = head;
 Node* fast = head->next;
 Node* newLastNode = NULL;
 int sum = 0

SUM Kab tak karna hai-jab tak fast->val == 0 na ho

$$\text{sum} = 0 + 3 + 1 \\ = 4$$



$$\text{sum} = 0 + 4 + 5 + 2 \\ = 11$$

```

Node* slow = head;
Node* fast = head->next;
Node* newLastNode = NULL;
int sum = 0

```

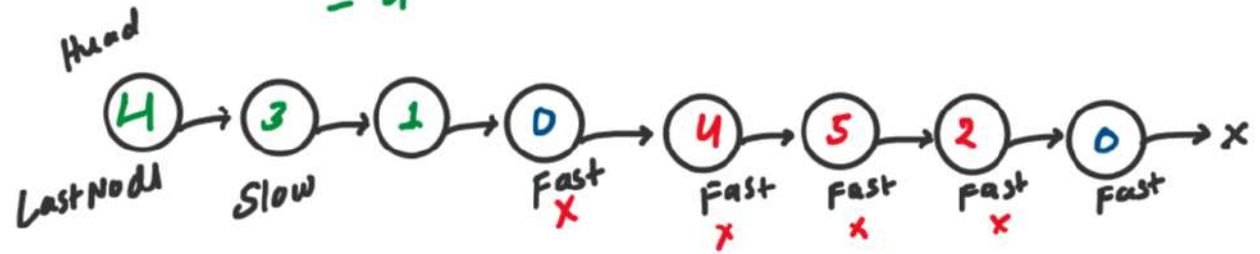
```

while (fast) {
    if (fast->val != 0) {
        sum += fast->val;
    }
    if (fast->val == 0) {
        slow->val = sum;
        lastNode = slow;
        slow = slow->next;
        sum = 0;
    }
    fast = fast->next;
}

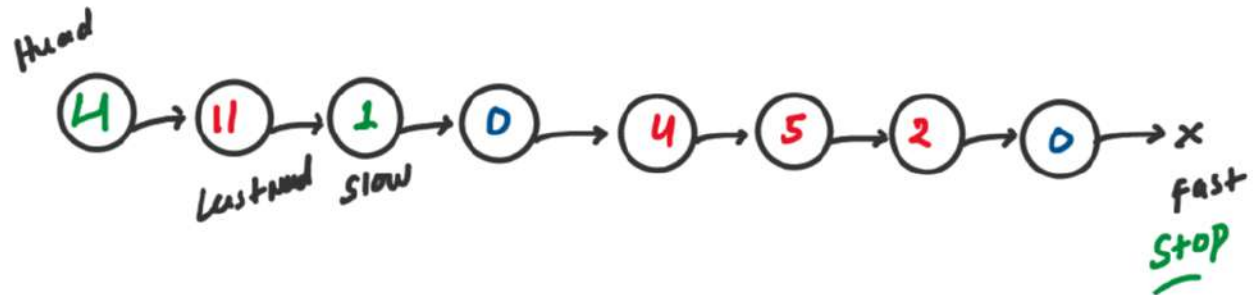
```

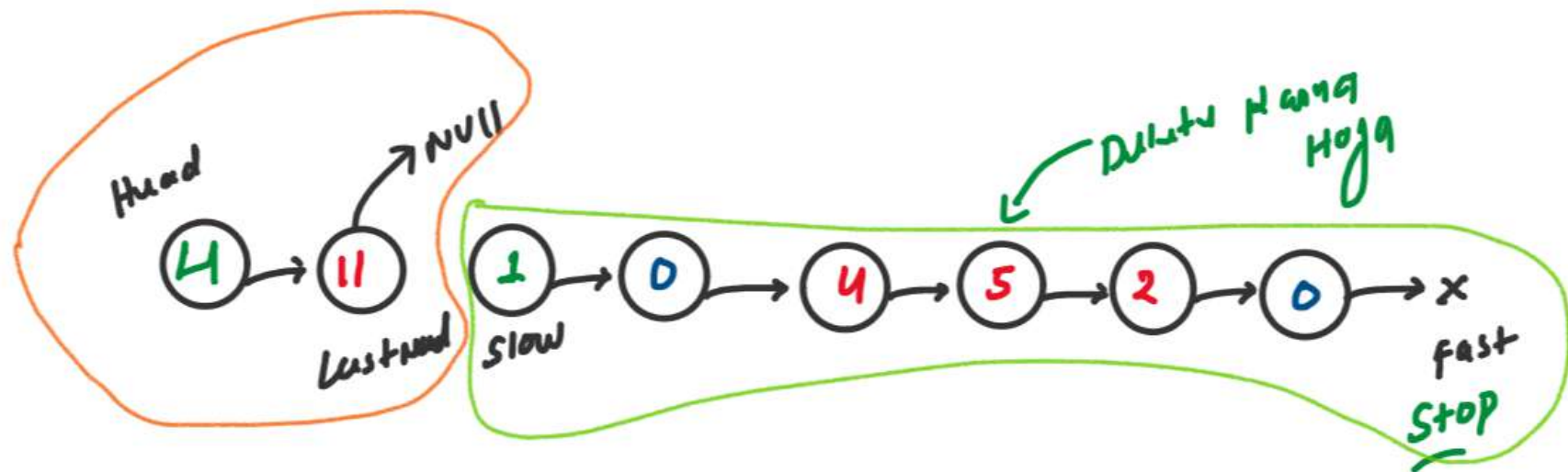
SUM Kab tak karna hai-jab tak fast->val == 0 na ho

$$\text{sum} = 0 + 3 + 1 = 4$$



$$\text{sum} = 0 + 4 + 5 + 2 = 11$$





```

{
    node* temp = slow;
    lastNode->next = NULL;
    delete temp;
    return head;
}

```

```
// HW 11: Merge Nodes in between Zeros (Leetcode-2181)
```

```
class Solution {  
public:  
    ListNode* mergeNodes(ListNode* head) {  
        ListNode* slow = head;  
        ListNode* fast = head->next;  
        ListNode* newLastNode = NULL;  
        int sum = 0;  
  
        while(fast){  
            if(fast->val != 0){  
                sum += fast->val;  
            }  
            else{  
                // fast->val == 0  
                slow->val = sum;  
                newLastNode = slow;  
                slow = slow->next;  
                sum = 0;  
            }  
            fast = fast->next;  
        }  
  
        ListNode* temp = slow;  
  
        // Just formed new list  
        newLastNode->next = NULL;  
  
        // Deleting old list  
        delete temp;  
  
        return head;  
    }  
};
```

Time complexity: $O(N)$,

Where N is number of nodes of the linked list

Space complexity: $O(1)$,

Where no extra space used