

20/10/2023

QUICK SORT ALGORITHM

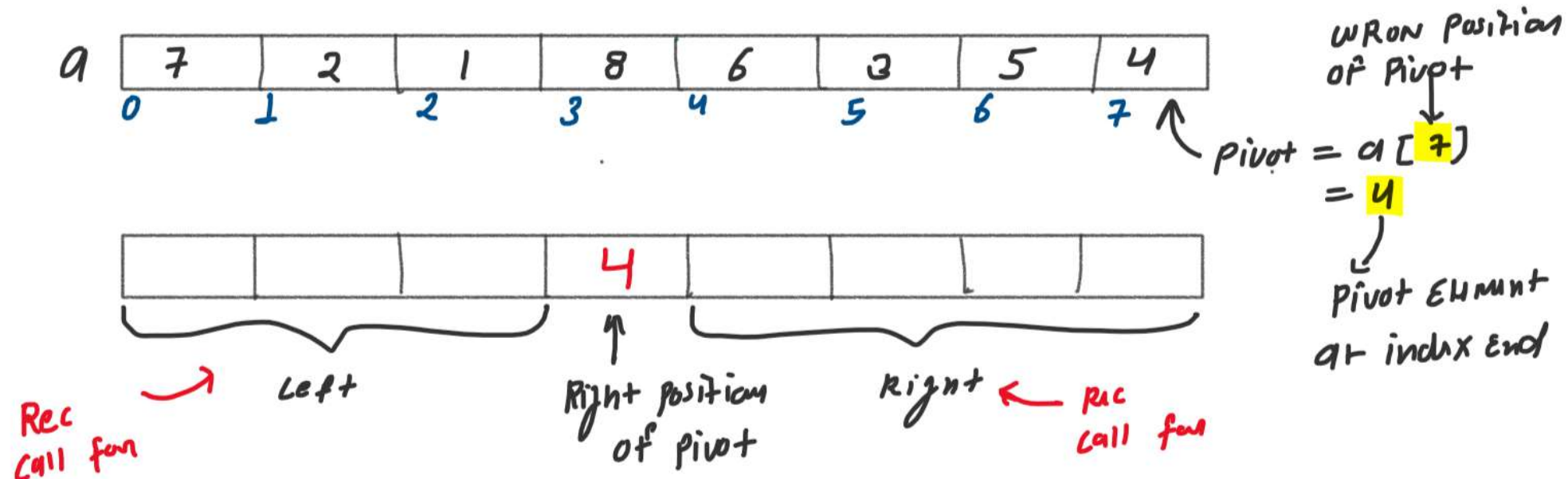
QUICK SORT ALGORITHM

New way to partitioning of array (Ek Step Chalna Mujhe Aata hai)

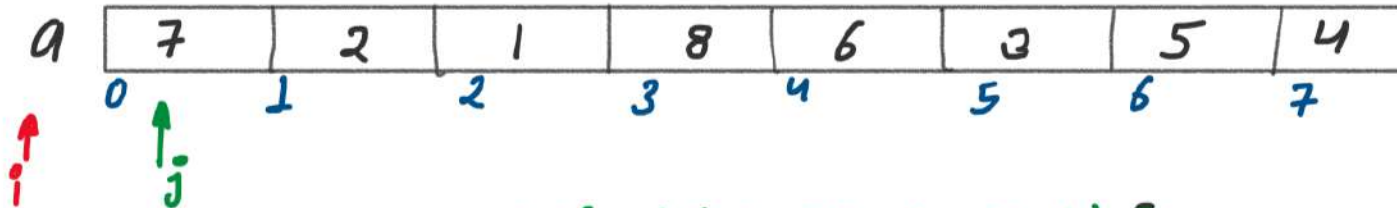
Step 01: Find pivot (It is always end element of array)

- To place pivot such that the element to the **right of pivot** $> a[\text{pivot}]$
- To place pivot such that the element to the **left of pivot** $< a[\text{pivot}]$

Step 02: Now apply recursion for left and right part of pivot (Baki ka recursion sambhal lega)



STEP:01



$start = 0$
 $end = 7$
 $pivot = end$
 $i = start - 1$
 $j = start$

$quickSort(a, start, end) \{$

$while (j < pivot) \{$

$if (a[j] < a[pivot]) \{$

$i++;$

$swap(a[j], a[i]);$

$j++;$

$i++;$

$swap(a[i], a[pivot]);$

$\}$

MUJHE

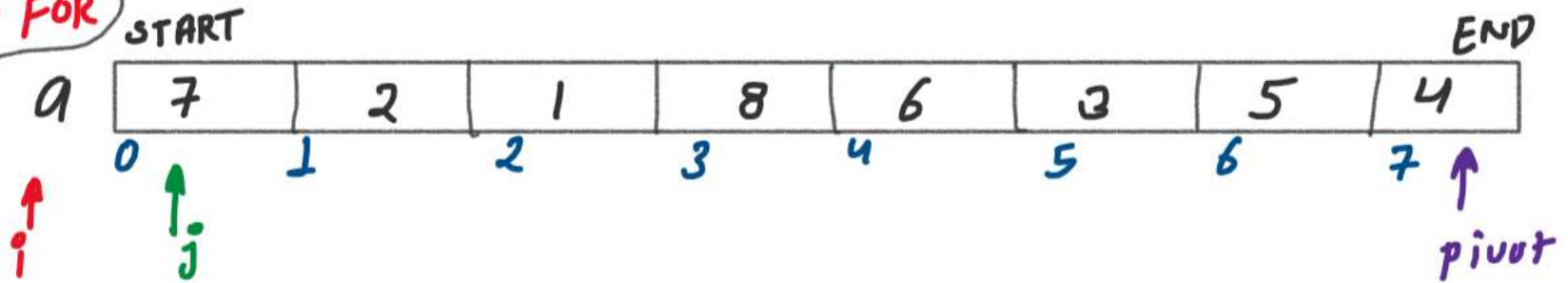
First STEP

Chalna AATA hai
janna par me 4
ko uski **Right**
Position par
RAkhana janta
HU.

DRY RUN FOR

STEP 1

①



start = 0

end = 7

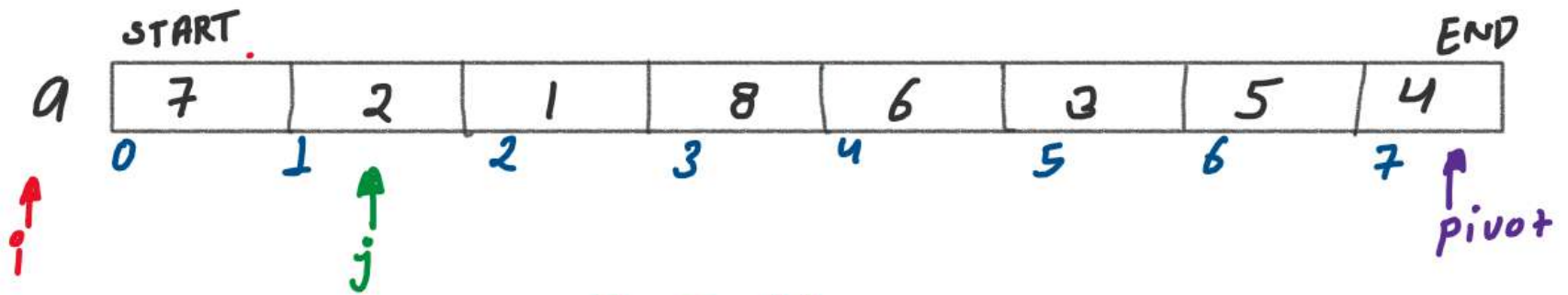
pivot = end

i = start - 1

j = ~~start~~ 1

$a[j] > a[pivot]$
7 > 4
j++

②



Start = 0

End = 7

Pivot = End

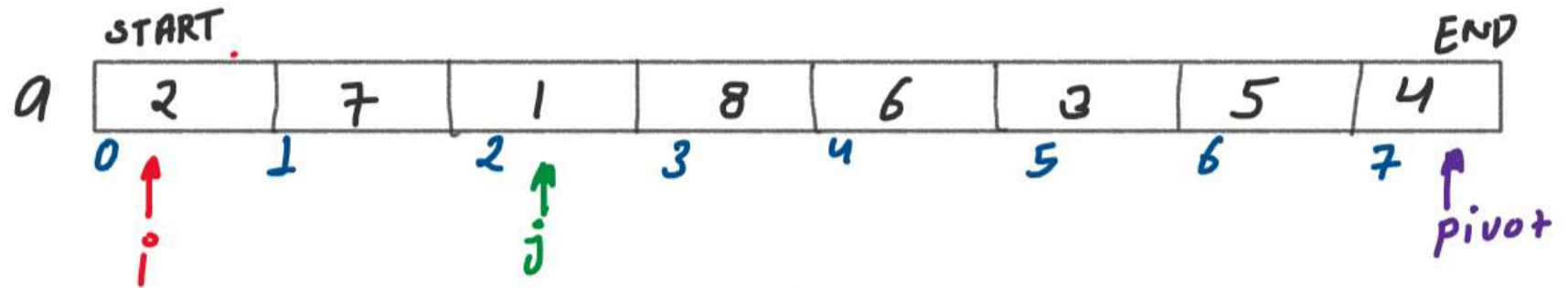
$i = \text{Start} - 1$ 0

$j = 1$ 2

$a[j] < a[\text{Pivot}]$
 $i++$

Swap($a[i]$, $a[j]$)
 $j++$

③



start = 0

end = 7

pivot = end

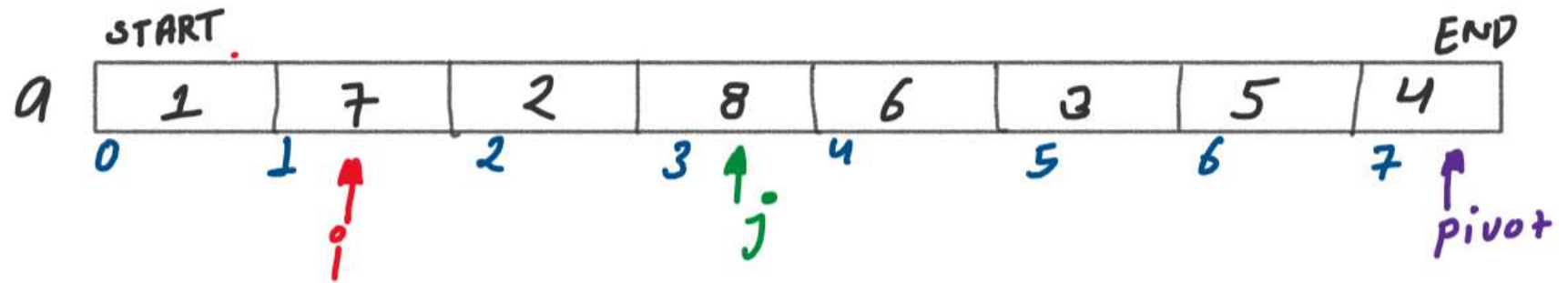
$i = 0$ ~~1~~

$j = 2$ ~~3~~

$1 < 4$
 $a[j] < a[pivot]$
→ $i++$

swap($a[i]$, $a[j]$)
 $j++$ 2 1

4



$start = 0$

$end = 7$

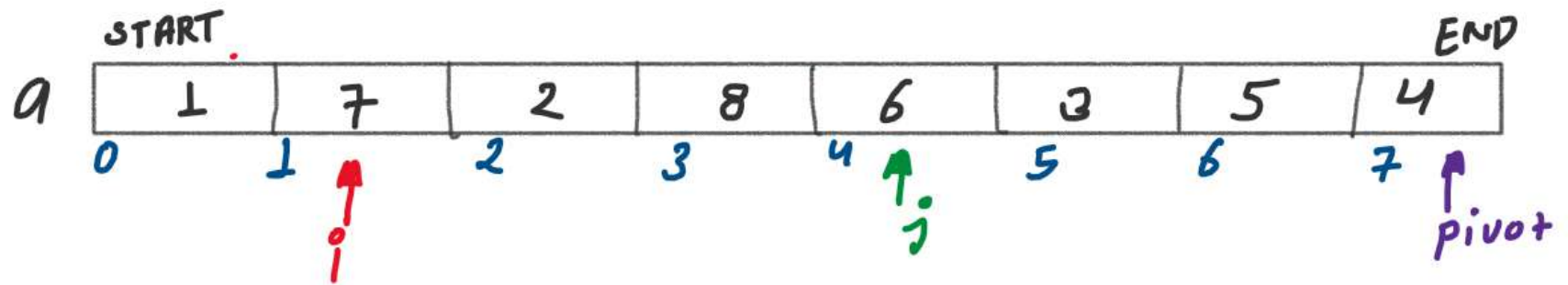
$pivot = end$

$i = 1$

$j = \cancel{3} \ 4$

$a[j] > a[pivot]$
 $8 > 4$
 $j++$

(5)



$\text{start} = 0$

$\text{end} = 7$

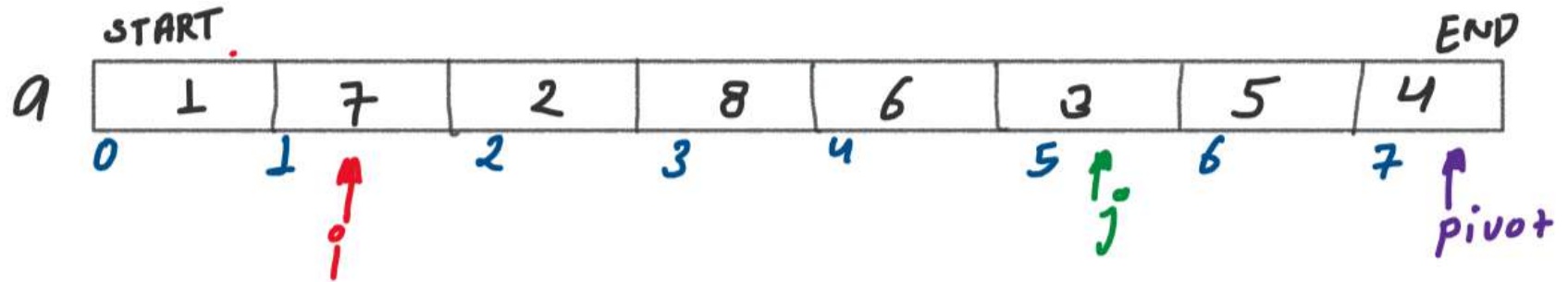
$\text{pivot} = \text{end}$

$i = 1$

$j = \cancel{4} 5$

$a[j] > a[\text{pivot}]$
 $\quad \quad \quad 6 > 4$
 $\quad \quad \quad \downarrow$
 $\quad \quad \quad j++$

⑥



start = 0

end = 7

pivot = end

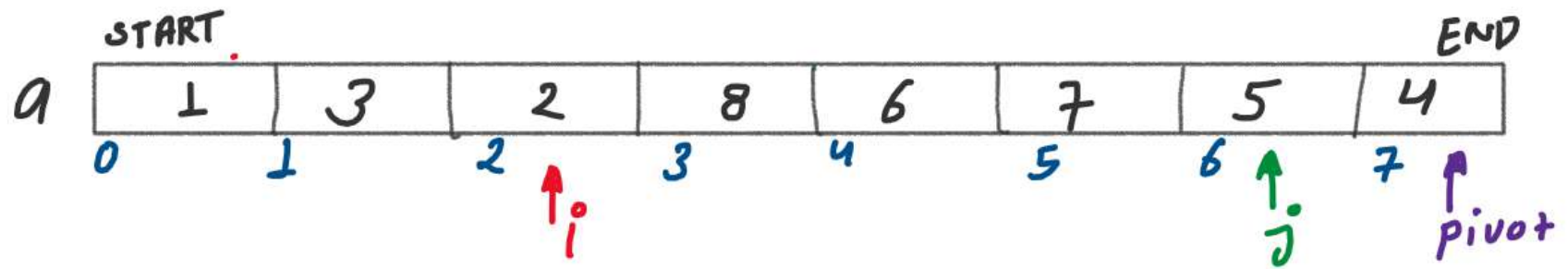
i = ~~1~~ 2

j = ~~5~~ 6

$a[j] < a[pivot]$
 $3 < 4$
→ i++

swapt(a[i], a[j])
j++ 7 3

⑦



$start = 0$

$end = 7$

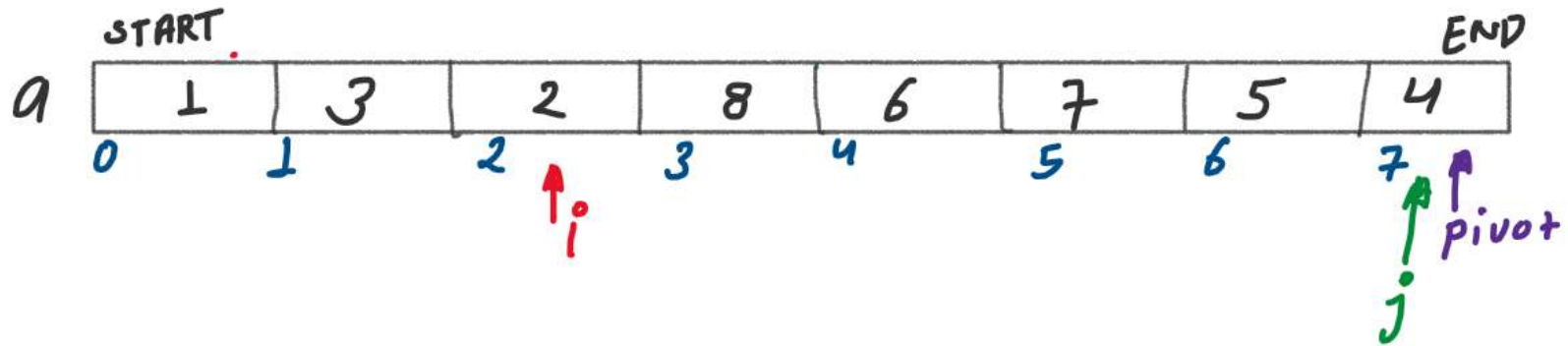
$pivot = end$

$i = 2$

$j = \cancel{6} \ 7$

$a[j] > a[pivot]$
 $\quad \quad \quad \underset{5}{a} \quad \underset{7}{b} \quad \underset{4}{c}$
 $\quad \quad \quad \rightarrow j++$

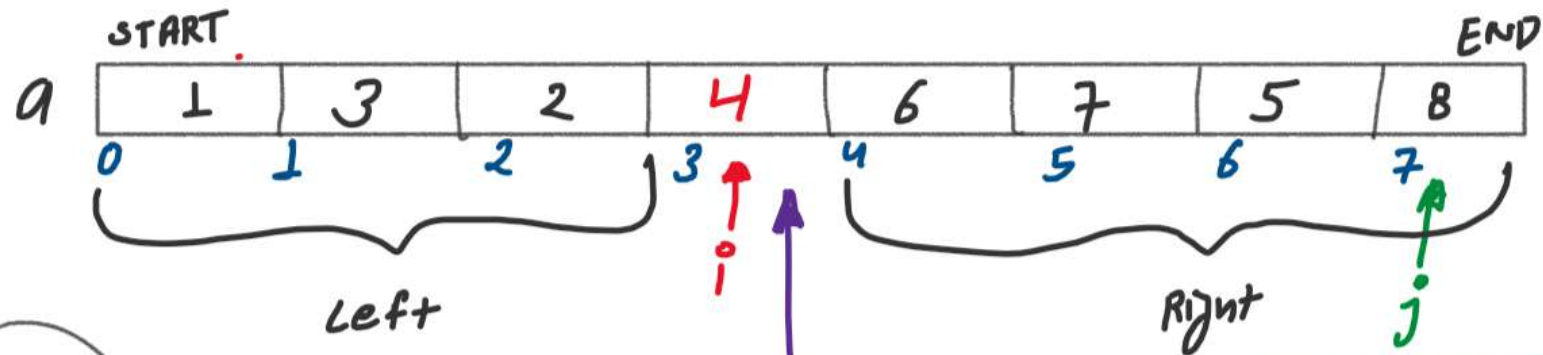
⑧



Start = 0
End = 7
pivot = End
 $i = \cancel{2} \text{ } 3$
 $j = 7$

while($j^0 < \text{pivot}$) { --- }
↓ $7 < 7 \times$ while loop me entry nahi karungi
Outside of while loop
 $i++$
swap($a[i]$, $a[\text{pivot}]$);
 8 4

9



pivot at
right position now

$(1, 3, 2 < 4)$

QuickSort
(a, start, $i-1$)

$(6, 7, 5, 8 > 4)$

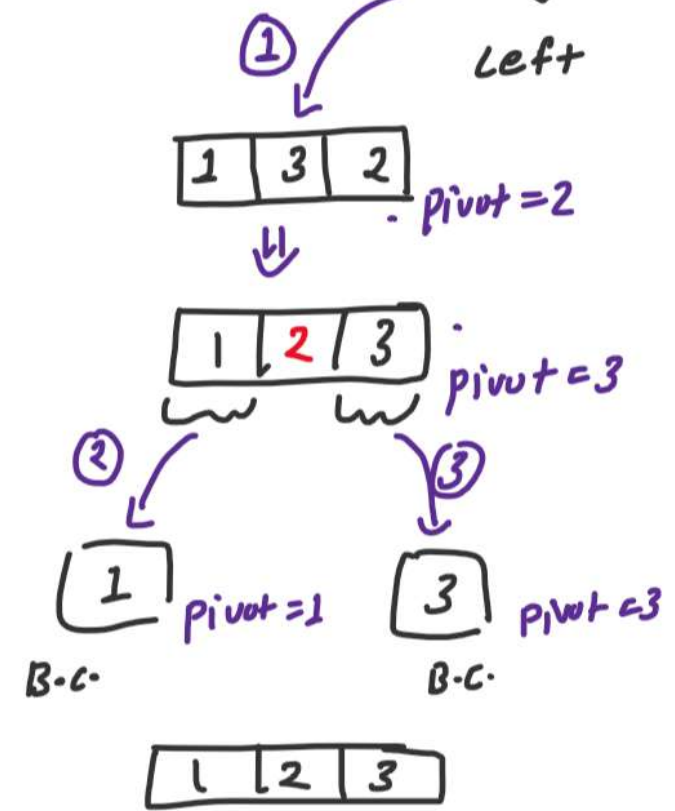
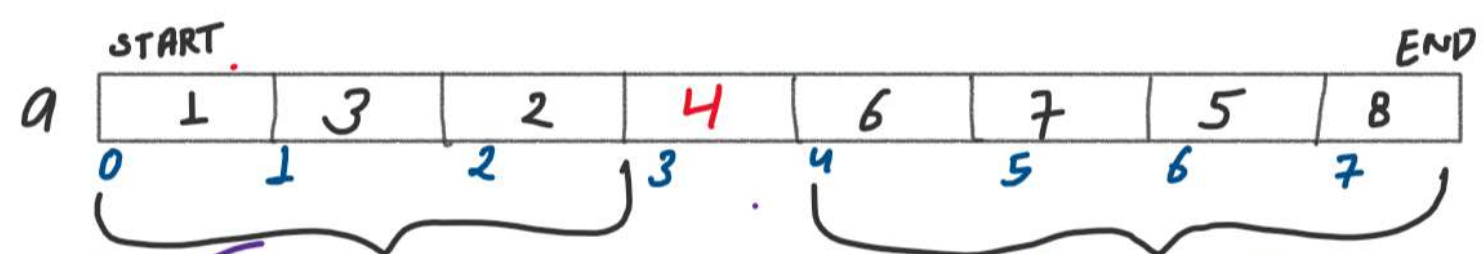
QuickSort
(a, $i+1$, end)

start = 0
end = 7
pivot = end

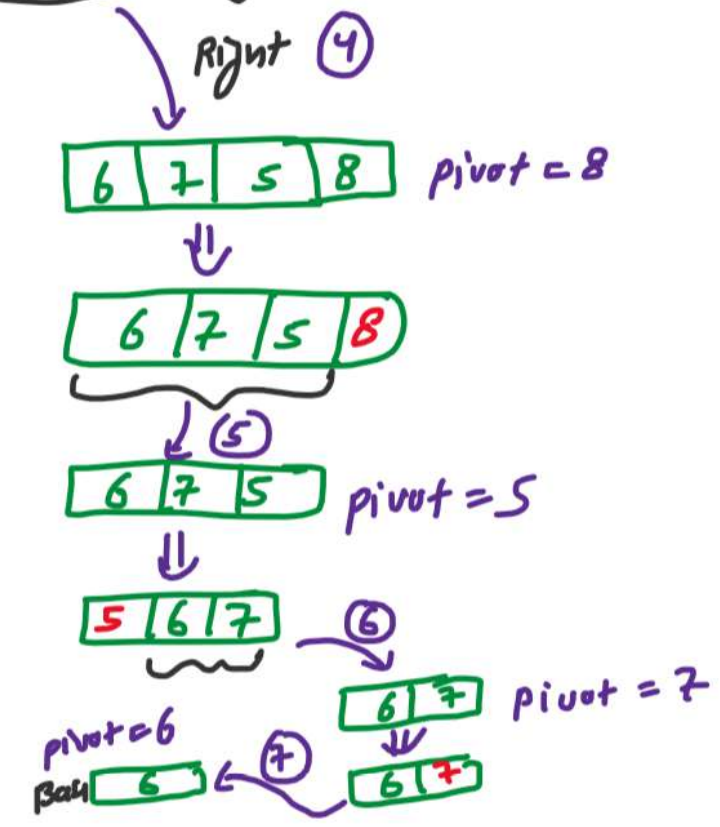
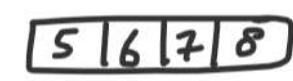
$i = 3$

$j = 7$

DRY RUN
FOR STEP 2



Total
call = 7
= 8-1
= N-1
T.C. = O(N)
S.C. = O(N)



```
//Function to sort an array using quick sort algorithm.
void quickSort(int arr[], int start, int end)
{
```

```
    // Base Case
    if(start >= end) {
        return;
    }
```

```
    int pivot = end;
    int i = start - 1;
    int j = start;

    while(j < pivot){
        if(arr[j] < arr[pivot]){
            i++;
            swap(arr[i], arr[j]);
        }
        j++;
    }
    i++;
    swap(arr[i], arr[pivot]);
```

```
    // Recursive call karlo Left part ke liye
    quickSort(arr, start, i - 1);
    // Recursive call karlo Right part ke liye
    quickSort(arr, i + 1, end);
}
```

STEP 2 T.C. = $O(N)$
S.C. = $O(1)$

→ STEP: 01
T.C. = $O(N^2)$
S.C. = $O(1)$

STEP 1 WORST CASE T.C.

EX

8	7	6	5	4	3	2	1
---	---	---	---	---	---	---	---

T.C. = $O(N^2)$ S.C. = $O(1)$

STEP 2 Recursive call

S.C. = $O(N)$
T.C. = $O(N)$

Total T.C. = $O(N^2) + O(N)$
= $O(N^2)$

Total S.C. = $O(N) + (1)$
= $O(N)$

6	1
(6 7)	2
(6 7 5)	3
(6 7 5 8)	4
(3)	5
(1)	6
(1 3 2)	7
main	

Stack call

} 7 Entry
8-1
⇒ N-1
⇒ $O(N)$
S.C.