

18/10/2023

DIVIDE AND CONQUER

CLASS 1

Merge sort

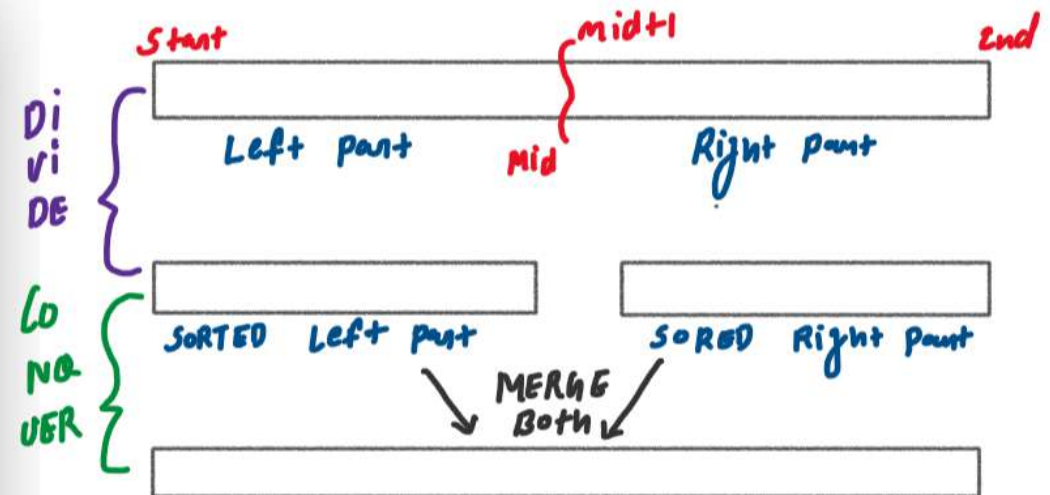
Step 01: Find mid and break the original array into two equal part

Left part [start, mid] and Right part [mid+1, end]

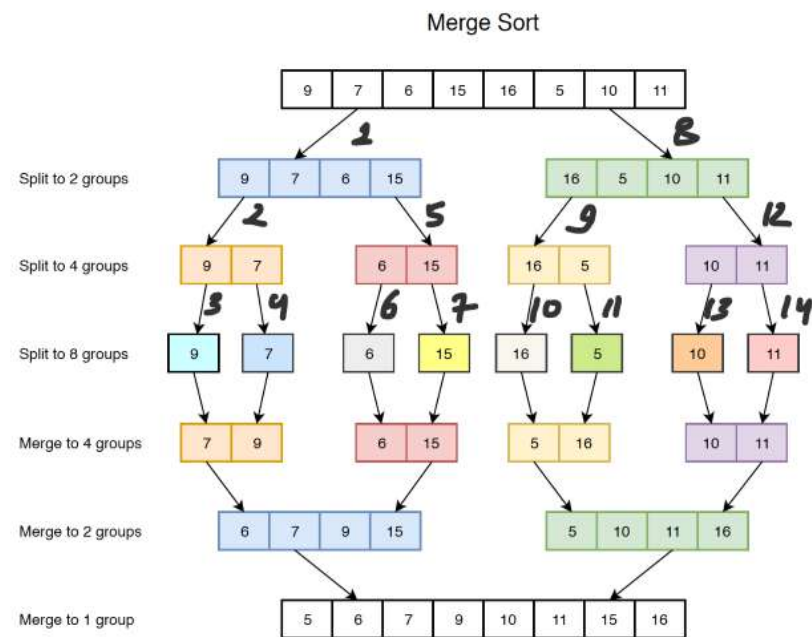
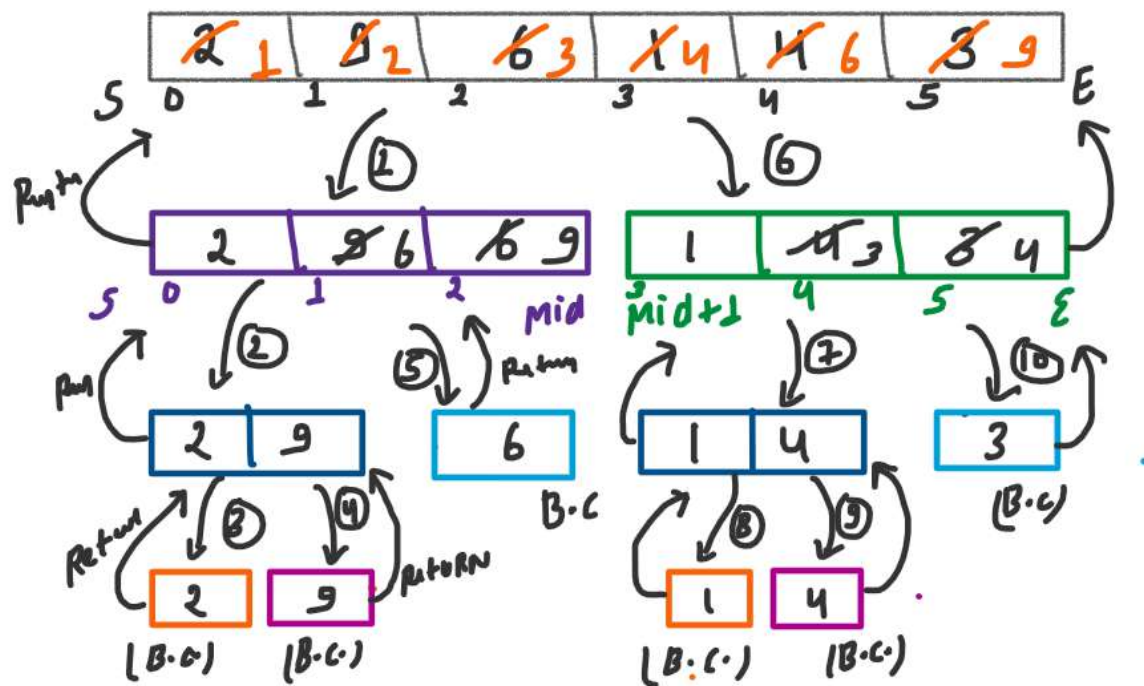
Step 02: Recursion call for sorting left and right part

Step 03: Merge two sorted arrays

```
1 // Merge Sort (Divide and conquer algorithm)
2 void mergeSort(int arr[], int start, int end){
3     // Base Case
4     if(start > end){
5         // Invalid Array
6         return;
7     }
8     if(start == end){
9         // Single element array
10        return;
11    }
12
13    // Step 01: Find mid and break the original array into two equal part
14    int mid = start + (end - start)/2;
15
16    // Step 02: Recursion call for sorting left and right part
17    // Recursive call for left part
18    mergeSort(arr, start, mid);
19    // Recursive call for right part
20    mergeSort(arr, mid+1, end);
21
22    // Step 03: Merge two sorted arrays
23    merge(arr, start, end, mid);
24 }
```



DRY RUN



DYNAMIC MEMORY ALLOCATION

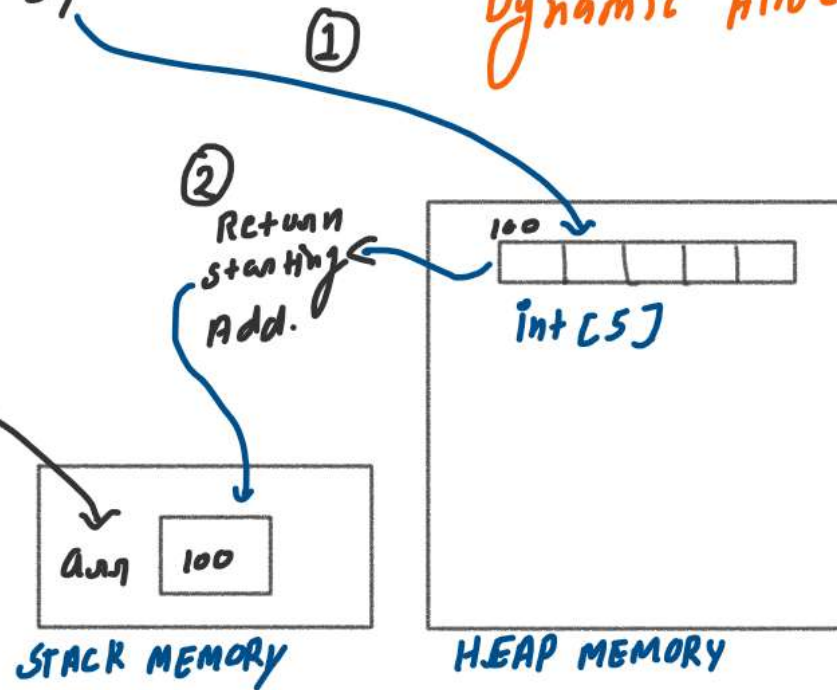
```
int *arr = new int[5];
```

STATIC
Allocation

Dynamic Allocation

😊
deAllocate
karna mat
Bhoolna

```
int arr;
```



MERGE FUNCTION (STEP:03)

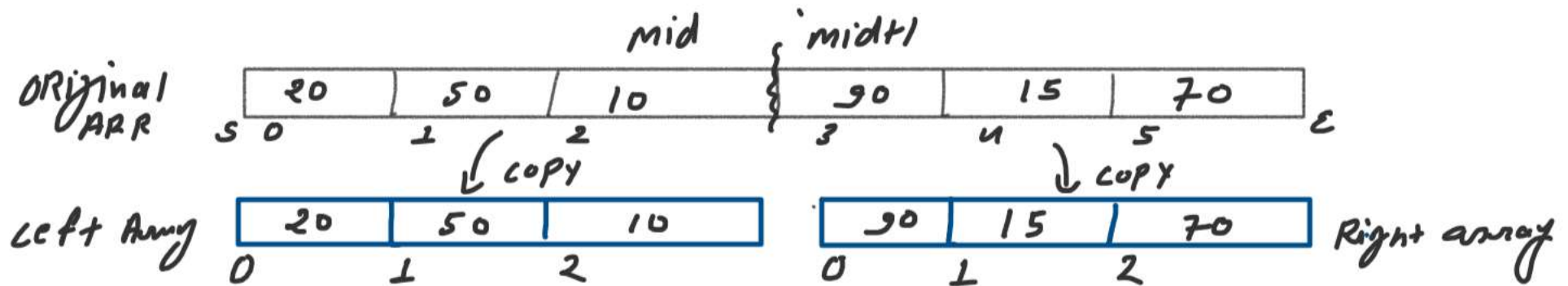
1st step

Step 01: find the length of left and right part array

Step 02: create left and right part array

Step 03: copy value from original array to left and right part array

Step 04: write actual logic to merge left and right sorted array



$$\begin{aligned} \text{lenLeft} &= \text{mid} - s + 1 \\ &= 2 - 0 + 1 \\ &= 3 \end{aligned}$$

$$\begin{aligned} \text{lenRight} &= e - (\text{mid} + 1) + 1 \\ &= 5 - (3) + 1 \\ &= 3 \end{aligned}$$

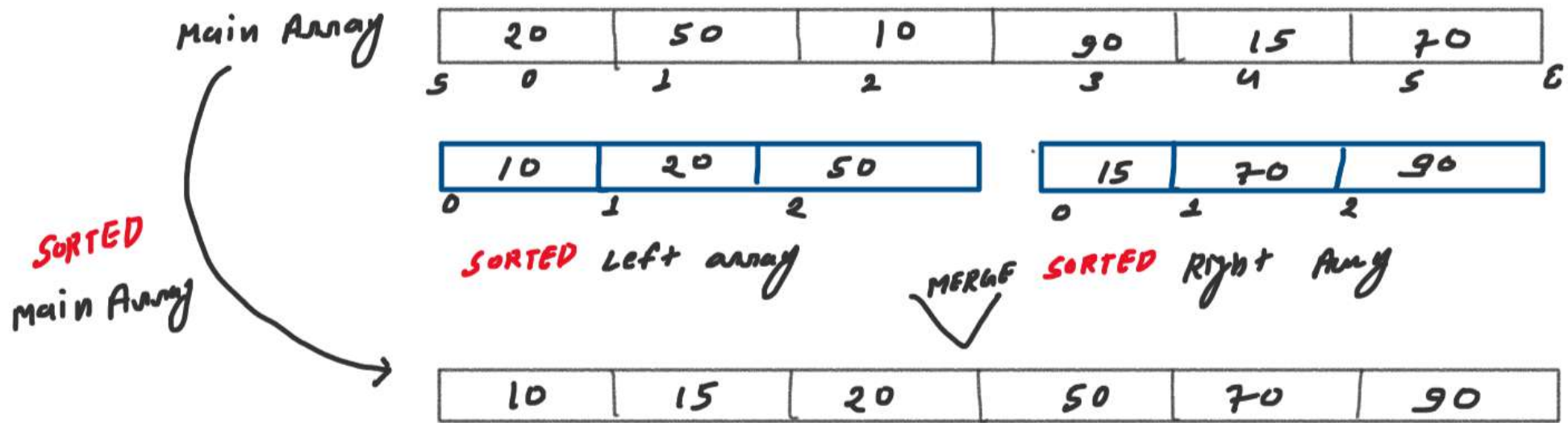
```

1 // Merge Function
2 void merge(int arr[], int start, int end, int mid){
3     // Step 01: find the length of left and right part array
4     int lenLeft = mid - start + 1;
5     int lenRight = end - (mid + 1) + 1;
6
7     // Step 02: create left and right part array
8     int *left = new int[lenLeft];
9     int *right = new int[lenRight];
10
11     // Step 03: copy value from original array to left and right part array
12     int k = start;
13     // copy value from original array to left array
14     for(int i=0; i<lenLeft; i++){
15         left[i] = arr[k];
16         k++;
17     }
18     // copy value from original array to right array
19     for(int i=0; i<lenRight; i++){
20         right[i] = arr[k];
21         k++;
22     }
23
24     // Step 04: write actual logic to merge left and right sorted array
25     mergeTwoSortedArray(arr, start, left, right, lenLeft, lenRight);
26
27     // 😊 De-allocate (Free heap memory from arrays are left and right)
28     delete[] left;
29     delete[] right;
30 }

```

Free memory karna mat
Bhoolna

STEP 4 Merge Two Sorted Array



starting Index of all 3 Array

{

Left Index = 0

Right Index = 0

Main Array Index = 5

}

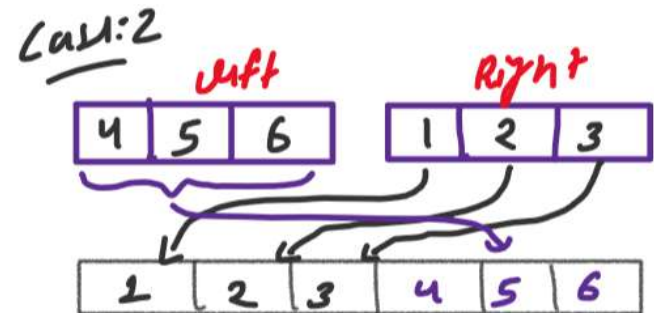
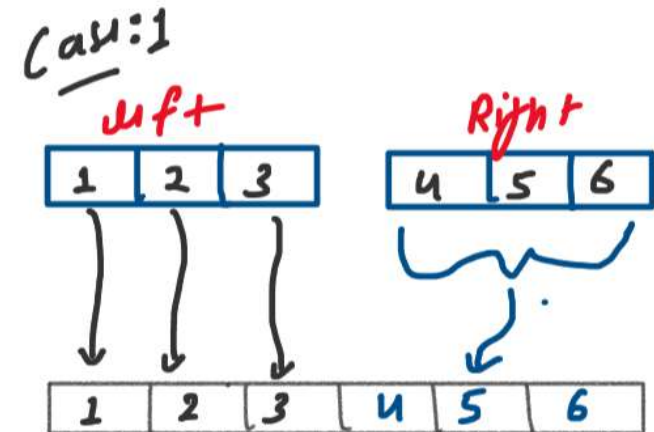
Two Pointers Approach

```

1 // Merge Two Sorted Array Function (Merge two sorted array)
2 void mergeTwoSortedArray(int arr[], int start, int *left, int *right, int lenLeft, int
  lenRight){
3     // left and right part array are already sorted
4     int leftIndex = 0;
5     int rightIndex = 0;
6     int mainArrayIndex = start; // Yeh Catch hai Yanha galti hone ke chance hai
7
8     while(leftIndex < lenLeft && rightIndex < lenRight){
9         if(left[leftIndex] < right[rightIndex]){
10             arr[mainArrayIndex] = left[leftIndex];
11             mainArrayIndex++;
12             leftIndex++;
13         }
14         else{
15             arr[mainArrayIndex] = right[rightIndex];
16             mainArrayIndex++;
17             rightIndex++;
18         }
19     }
20
21     // Case 01: Left array exhaust but right array me abhi bhi element bache huee hai
22     while(rightIndex < lenRight){
23         arr[mainArrayIndex] = right[rightIndex];
24         mainArrayIndex++;
25         rightIndex++;
26     }
27
28     // Case 02: Right array exhaust but left array me abhi bhi element bache huee hai
29     while(leftIndex < lenLeft){
30         arr[mainArrayIndex] = left[leftIndex];
31         mainArrayIndex++;
32         leftIndex++;
33     }

```

3 Catch
1



Time complexity of merge sort

MS() {

Base case $\rightarrow K_1$

MS (Left) $\rightarrow \frac{n}{2}$

MS (Right) $\rightarrow \frac{n}{2}$

merge() $\rightarrow n * K$

}

$$T(N) = K_1 + T(\frac{n}{2}) + T(\frac{n}{2}) + n * K$$

$$2^0 * T(N) = 1K_1 + 2T(\frac{n}{2}) + n * K$$

$$2 * T(\frac{N}{2}) = 2K_1 + 4T(\frac{n}{4}) + 2(\frac{n}{2} * K)$$

$$4 * T(\frac{N}{4}) = 4K_1 + 8T(\frac{n}{8}) + 4(\frac{n}{4} * K)$$

$$\vdots$$

$$2^{a-1} * T(1) = 2^{a-1} K_1$$

} a times

$$T(N) = K(1 + 2 + 4 + \dots + 2^{a-1}) + (a-1)(n * K)$$

$$T(N) = K_1 \underbrace{(1 + 2 + 4 + 8 + \dots + 2^{a-1})}_{\text{G.O.P.}} + (a-1)(n * K)$$

$$S_n = a \times \left(\frac{n^n - 1}{n - 1} \right)$$

Check Binary search
To G. To understand
 $a = \log n$

$$\Rightarrow 1 * (2^a - 1) \Rightarrow 2^a \Rightarrow 2^{\log 2^n} \Rightarrow n$$

$$\begin{aligned} T(N) &= \cancel{K_1 n} + a * n * \cancel{K} \text{ Ignore} \\ &= n + \log n * n \\ &= n \log n \end{aligned}$$

$$O(N \log N)$$

To G. of merge sort