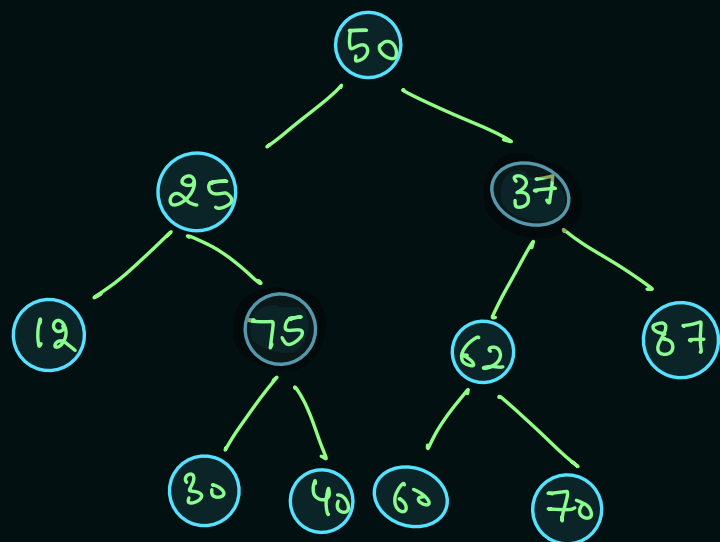


Recover BST :-

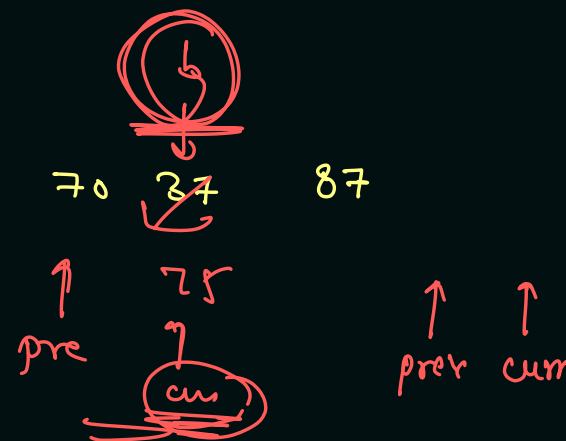


Inorder of BST → 12 25 30 ~~75~~ 40 50 60 62 70 ~~37~~ 87

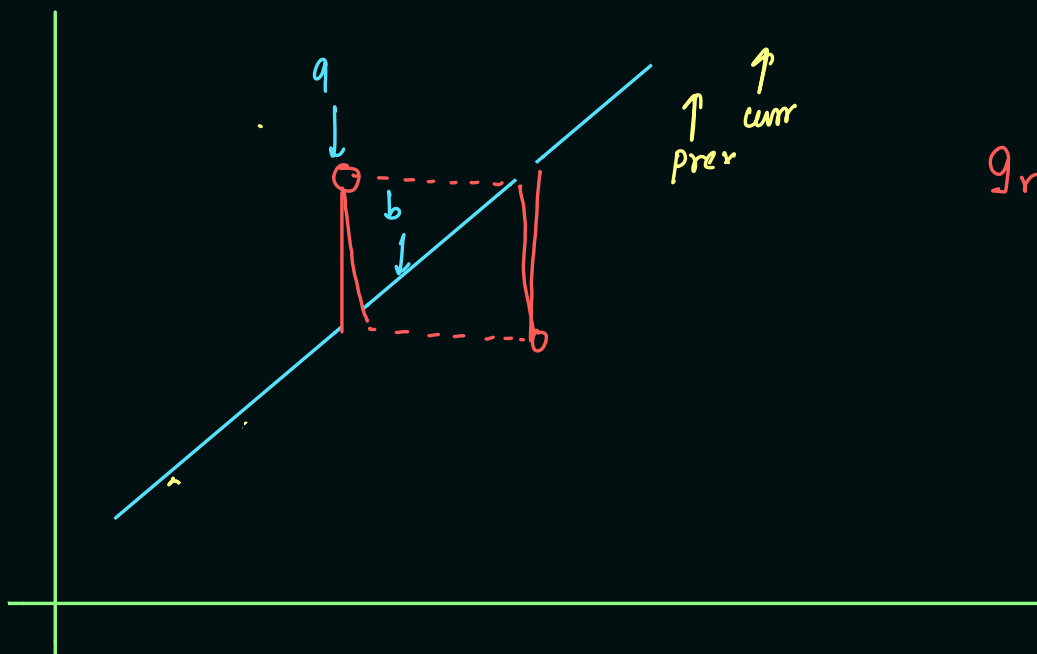
Fix faulty node in BST.

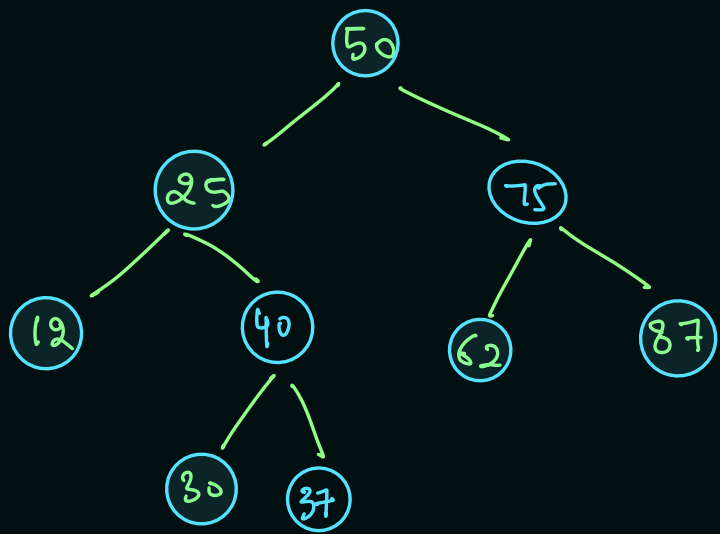
Initially tree was BST but one pair was somehow swapped. Now we have to fix this tree.

→ If Inorder is not sorted then BST is faulty.



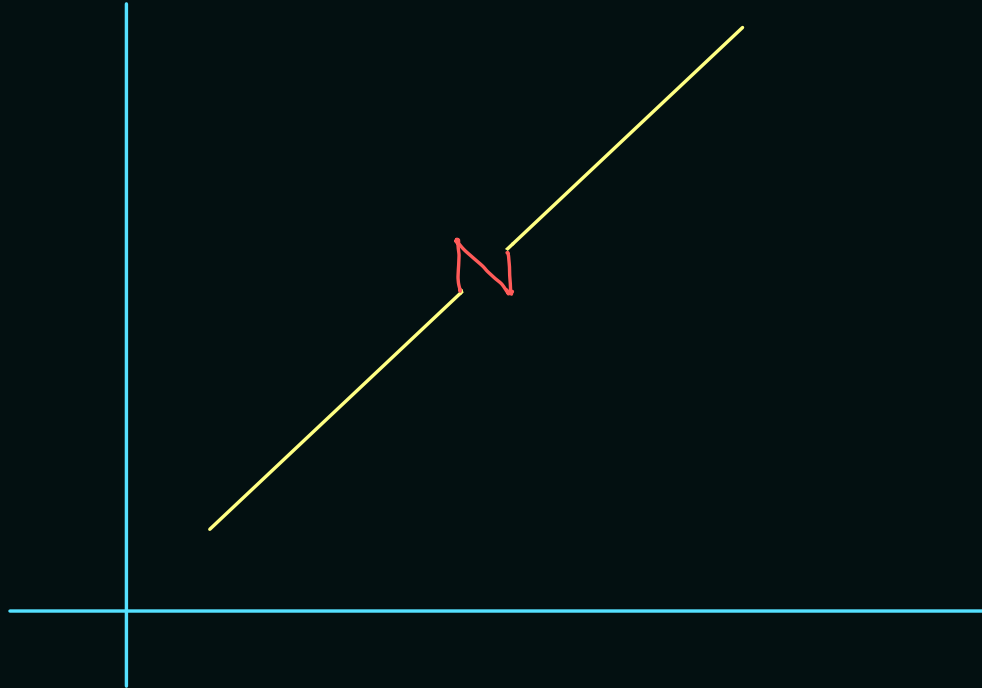
Inorder of BST





Inorder → 12 25 30 40 37 50 62 75 87

Annotations:
- Above 40: 9, 40, 5
- Above 37: 5
- Red arrows: from 40 to 37 (left), from 37 to 40 (right)
- Red bracket under 37 and 40
- Red arrows labeled 'prev' and 'cur' pointing up at the end of the sequence



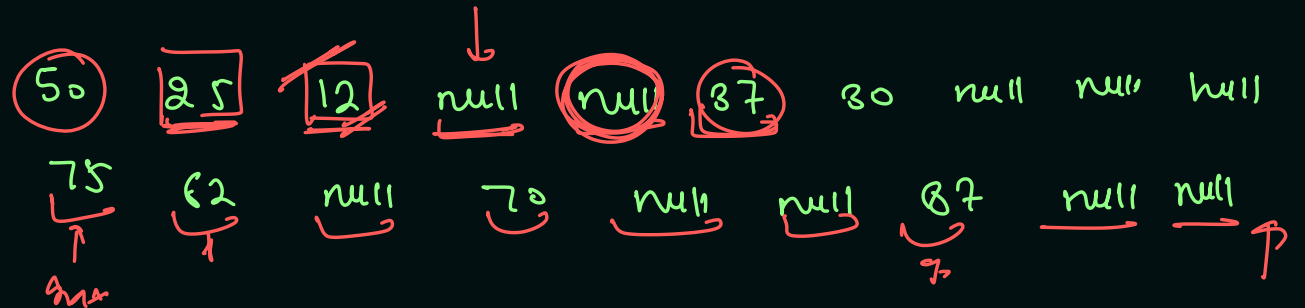
Serialize and deserialize :



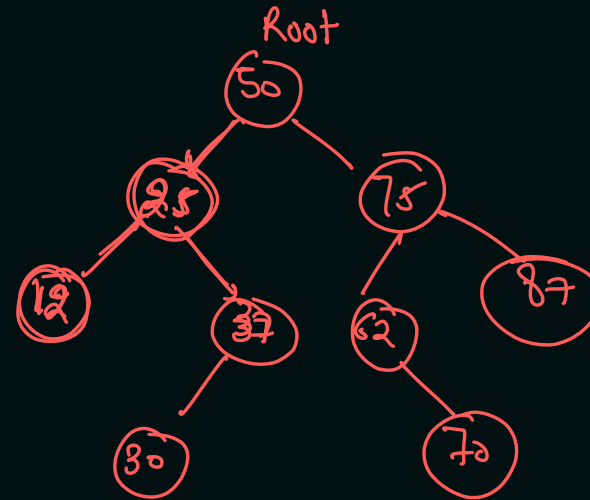
50 25 12 -1 -1 37 30 -1 -1 -1 75 62 -1
70 -1 -1 87 -1 -1

String → serialize -

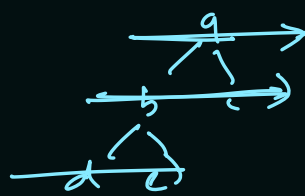
50 # 25 # 12 # null # null # 37 # 30 # null # null # null #
75 # 62 # null # 70 # null # null # 87 # null # null #



37, 0
25, 12
50, 1



left view and Right view :-



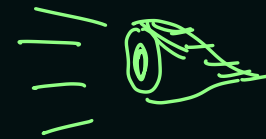
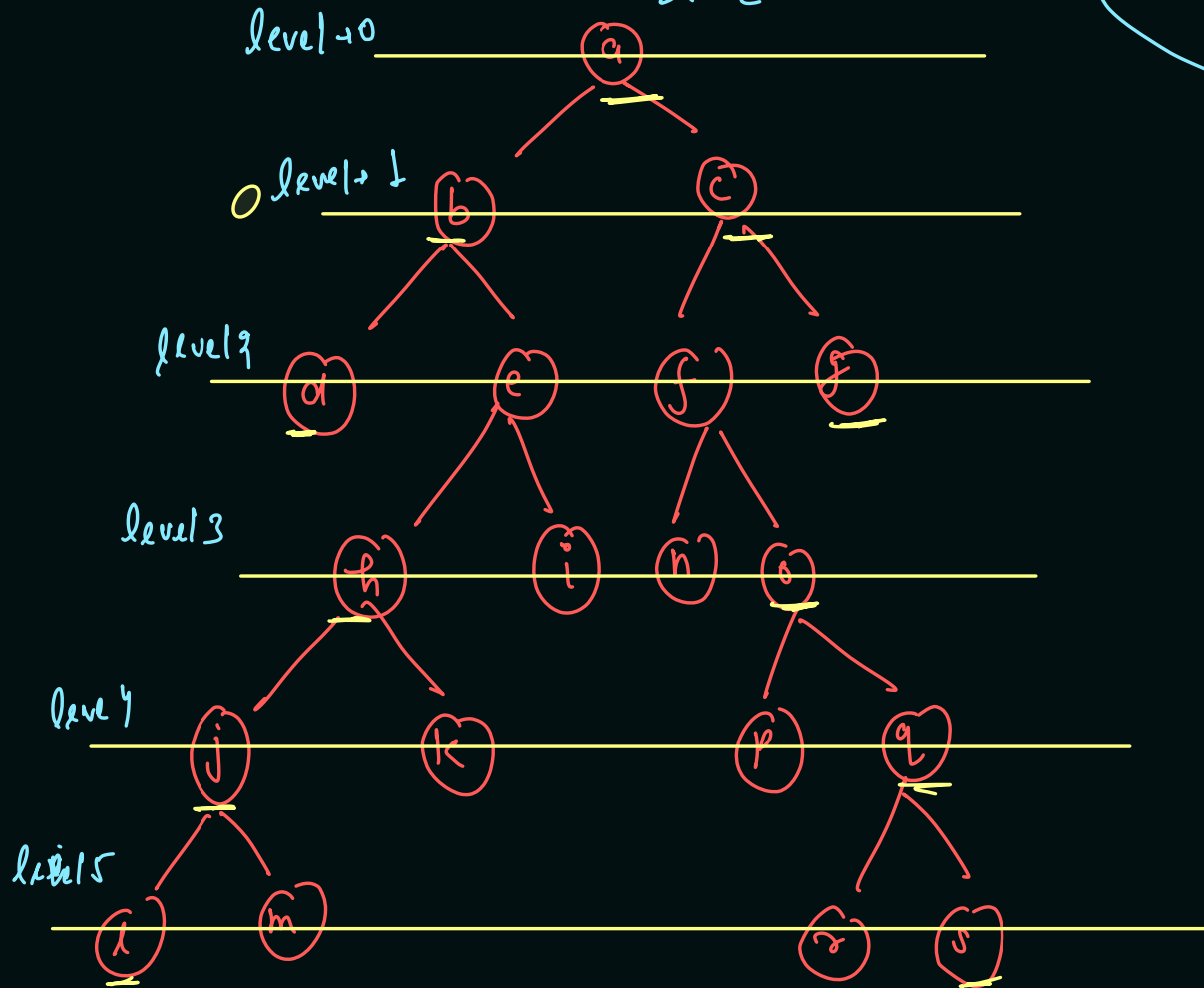
level order (1) →
 level order (2) → $\left[\begin{array}{c} \text{---} \\ \text{---} \\ \text{---} \end{array} \right]$
 using null delimiter
 why two queue
 ↓
 size based while

left view



a
b
d
h
j
l

left view



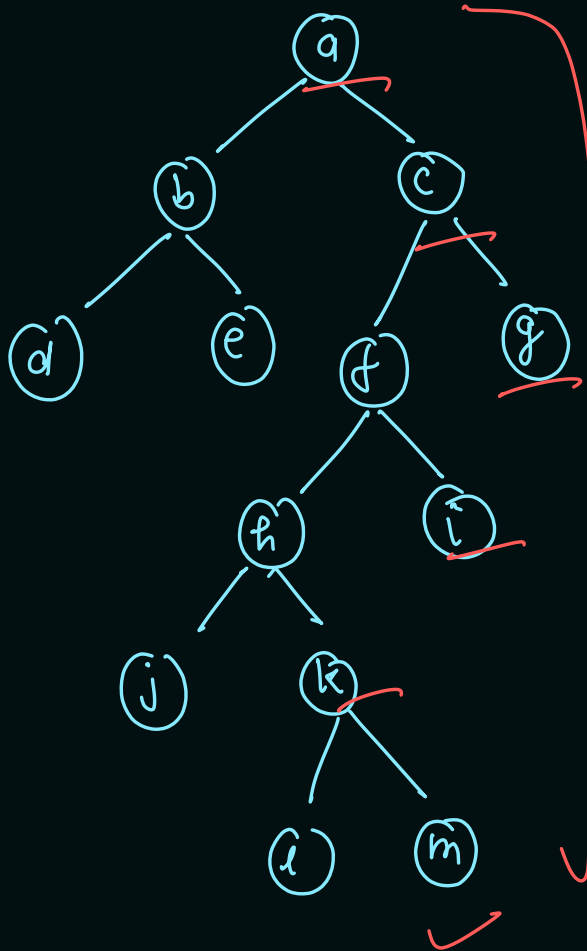
a
c
g
o
q
s

Right view

l0 → a
 l1 → b c
 l2 → d e f g
 l3 → h i n o ...

level order

left view → first node of level
 right view → last node of level



Right
view

~~l~~ / ~~m~~

Size = ~~1~~ ~~2~~ ~~4~~ ~~2~~ ~~2~~

a

b c

d e f g

h i

j k

l m

a

c

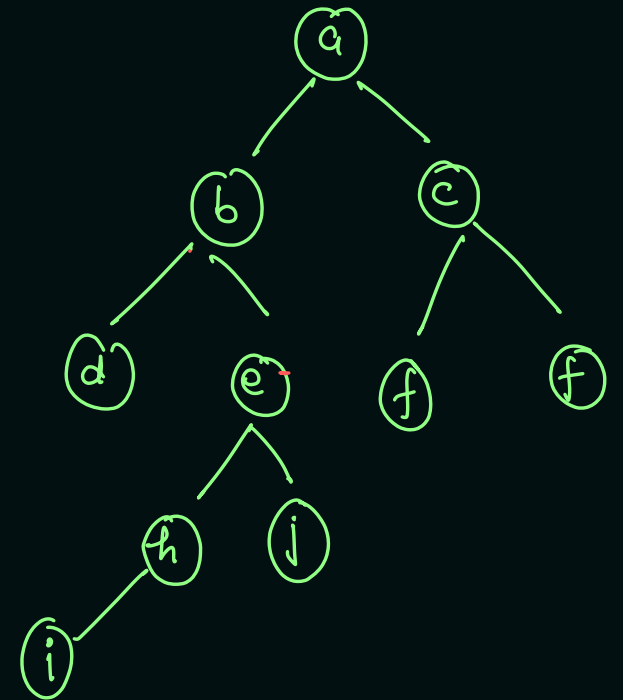
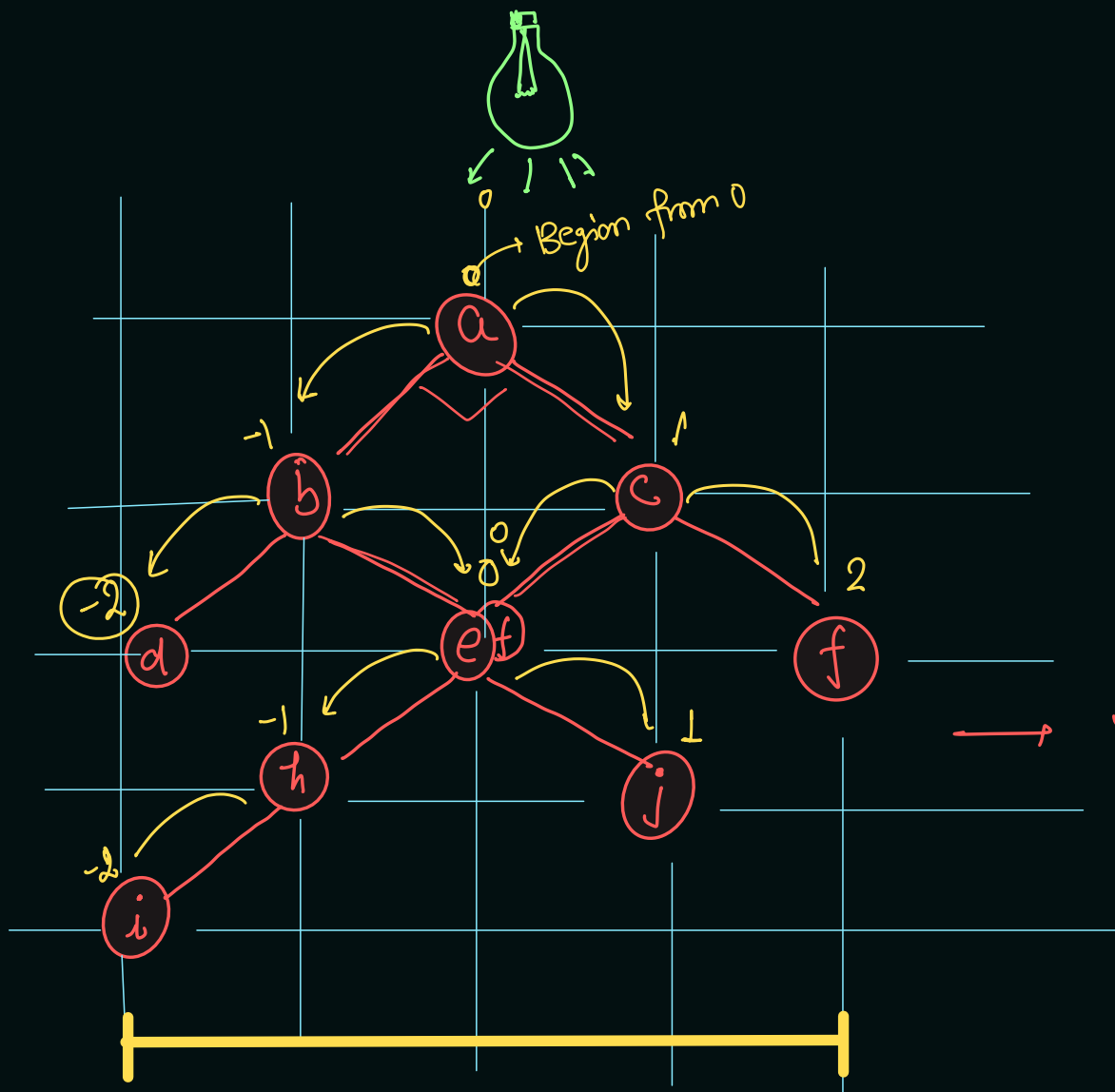
g

i

k

n

Width of Shadow of Binary Tree:-



Shadow length = 4 on basis of edges.

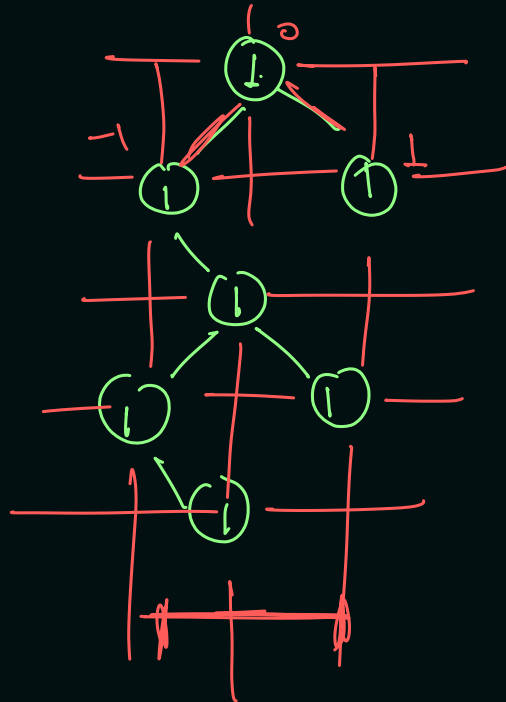
left = -2
right = +2
right - left

max left = |2|

max Right = 2

on the basis of node \rightarrow right - left + 1

15
1
1
-1
1
1
-1
1
-1
-1
1
-1
-1
-1
-1



on the basis of node
Result = Right - left + 1

~~Vertical order~~

Top view

Bottom view

✓ Mom's traversal

Pre

Platz

Any Node to any Node

Ind

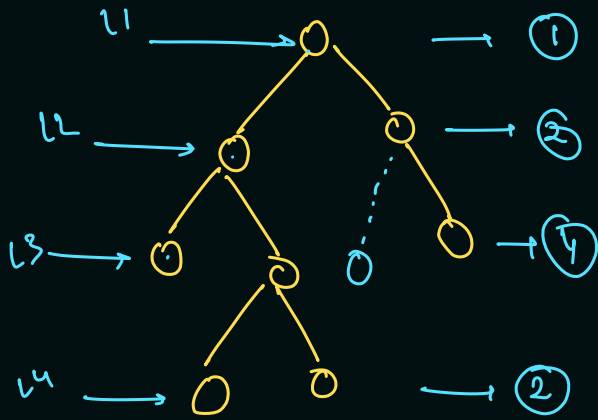
Max
8 in

[BST Iterator]

Max. width of Binary Tree: → * On a particular level, leftmost node & rightmost node should be not null

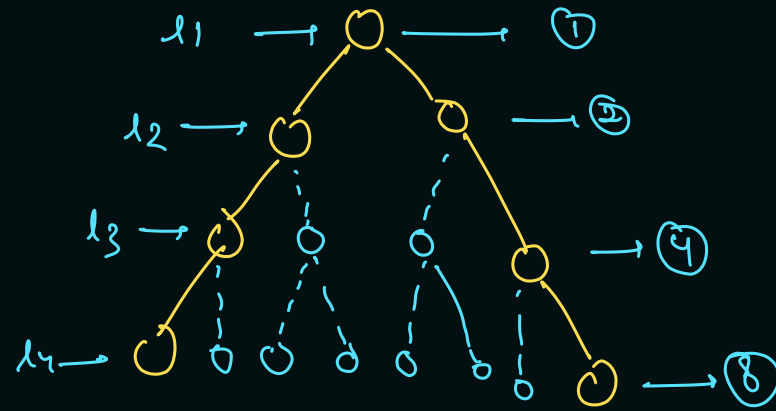
* Space of node on a particular level is width of that level

* Max. width??



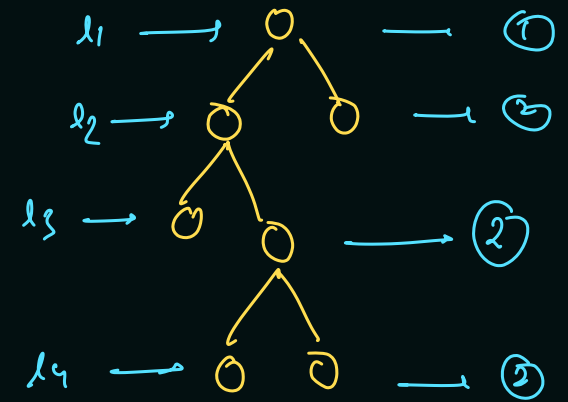
Max. width = 4

Shadow = 5



max width = 8

Shadow = 7



Max = 2

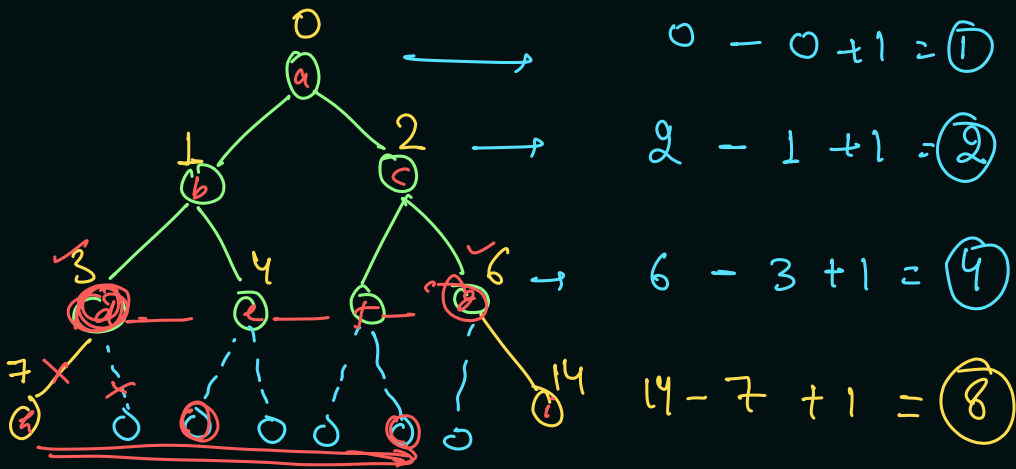
Shadow = 4

Task ⇒ Provide index to a complete Binary tree??

$rmi = \text{right most index}$

$lmi = \text{left most index}$

$$w = \overset{\text{width}}{rmi - lmi + 1}$$



~~h, 7~~ ~~i, 11~~

Size = ~~1~~ ~~2~~ ~~4~~ ~~8~~

$lmi = \min = 0$ ~~1~~ ~~3~~ ~~7~~

$rmi = \max = 14$ ~~1~~ ~~3~~ ~~4~~ ~~6~~ ~~7~~ ~~11~~

width = ~~1~~ ~~2~~ ~~4~~ ~~8~~

Indexing -

parent index = i

left child index = $2*i + 1$

Right child index = $2*i + 2$

Task

① level wise solve-

② parent index i

$lci = 2*i + 1$

$rmi = 2*i + 2$

$\min = lmi$
 $\max = rmi$

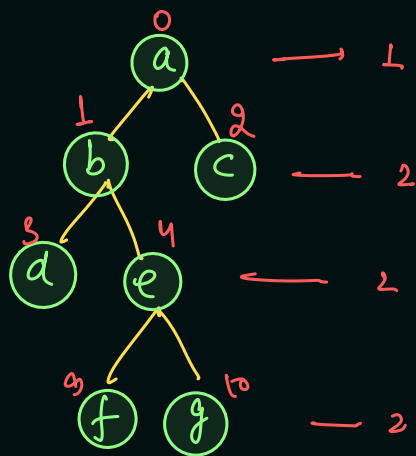
during solution of a particular

③ width at particular level

$rmi - lmi + 1$

④ maximise width.

Node
Index



Size = ~~1~~ 2 ~~2~~ 2

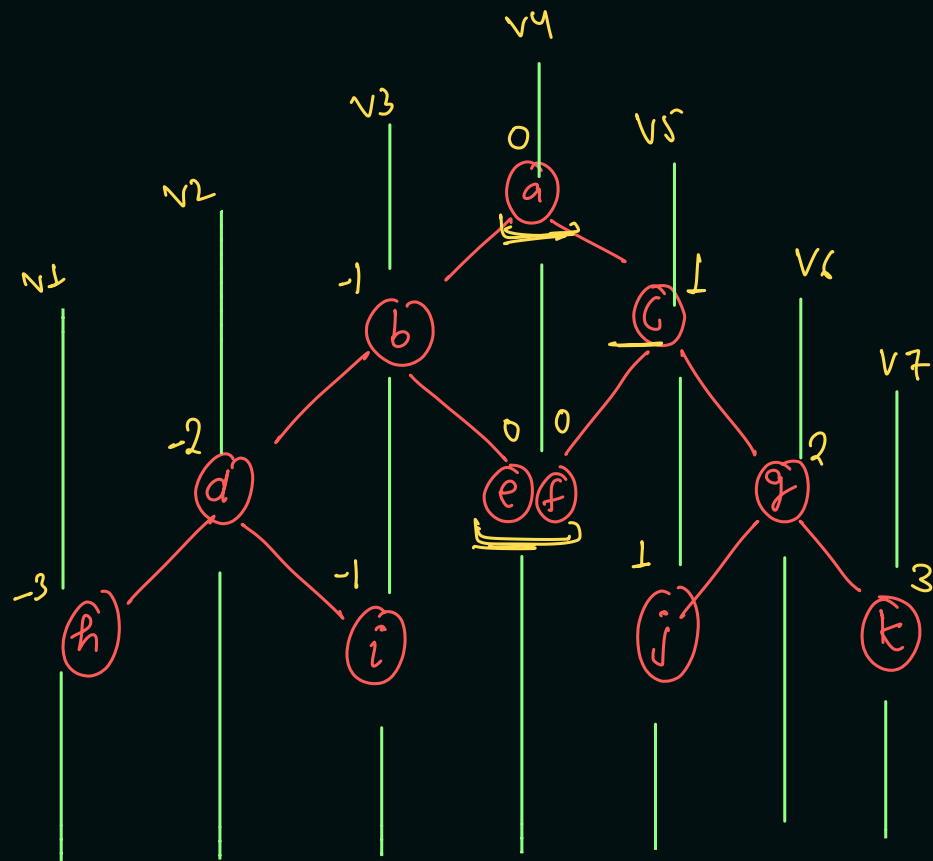
Max width 2 2

min index
lmi = ~~0~~ 1 3 9
rmi = ~~0~~ 1 2 3 4 9 10
max index

width = rmi - lmi + 1 = ~~1~~ 2 ~~2~~ 2

Overall width = ~~0~~ 1 2

Vertical order traversal : \rightarrow



$v_1 \rightarrow -3 \rightarrow h$

$v_2 \rightarrow -2 \rightarrow d$

$v_3 \rightarrow -1 \rightarrow b, i$

$v_4 \rightarrow 0 \rightarrow a, e, f$

$v_5 \rightarrow 1 \rightarrow c, j$

$v_6 \rightarrow 2 \rightarrow g$

$v_7 \rightarrow 3 \rightarrow k$

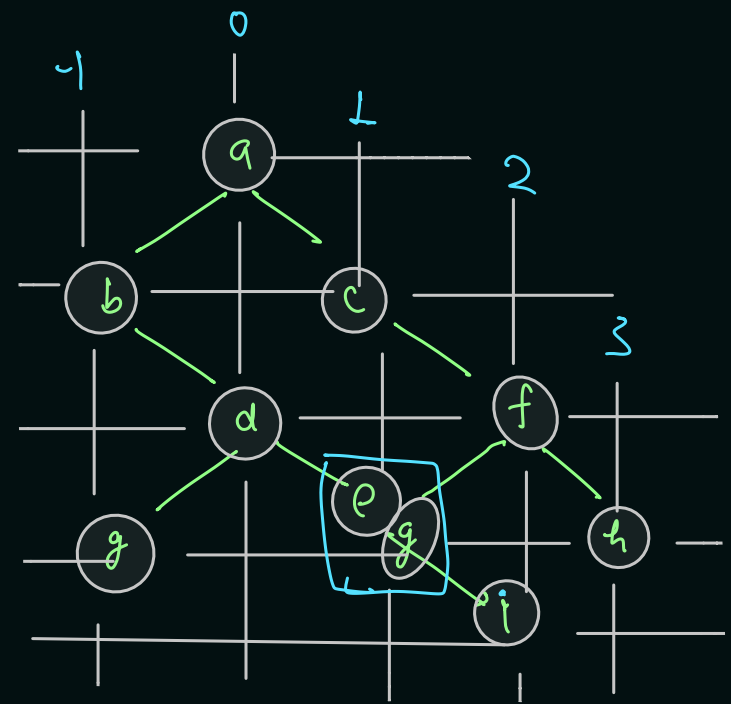
$lh = -3$

$rh = 3$

-3 to 3

HashMap < horizontal index, list(char) >

DFS



Size = ~~1~~ ~~2~~ ~~2~~ ~~4~~ 1

HashMap

-1 → b, g
0 → a, d
1 → c, e, f
2 → f, i
3 → h

vertical order (I)

lh = -1
rh = 3

-1 to 3
for every key -
array list get

-1 → b, g
0 → a, d
1 → e, c, f
2 → i, f
3 → h

problem

DFS is not good

we have to fill vertical
order level wise.

$v_1 \rightarrow 10$

$v_2 \rightarrow 3$

$v_3 \rightarrow 7, (13), (15)$

$v_4 \rightarrow 2, 6, 8$

$v_5 \rightarrow 5, (1), (3)$

$v_6 \rightarrow 9$

$v_7 \rightarrow 11$

queue \rightarrow priority queue

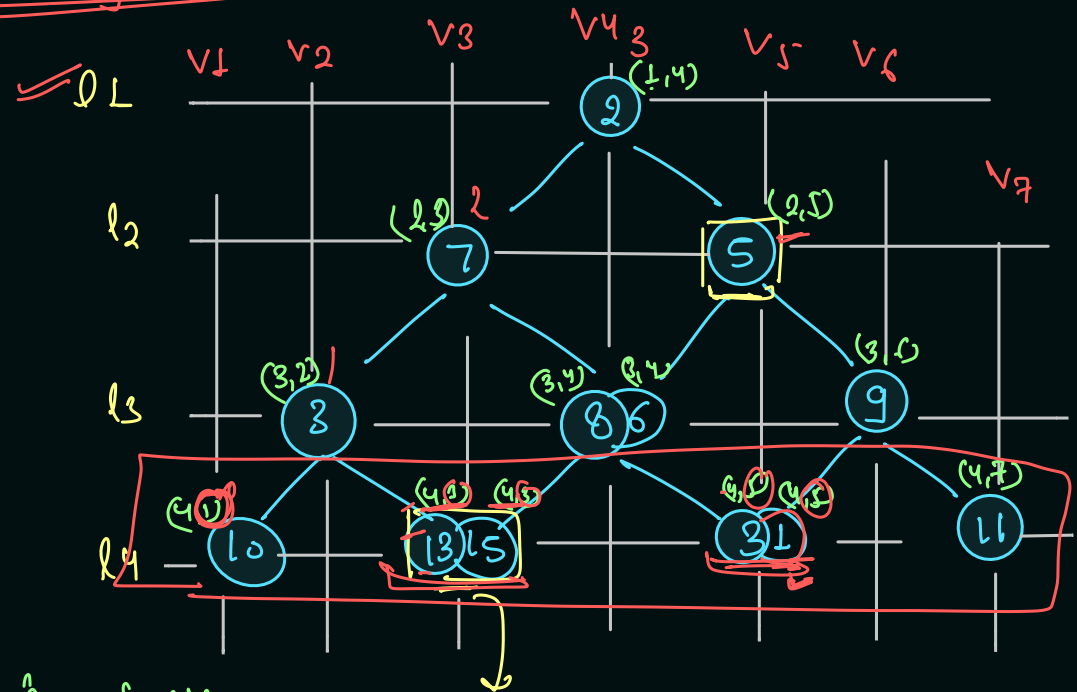
$lh = 3$

By default \rightarrow due to method of

Iteration \rightarrow level is prioritized

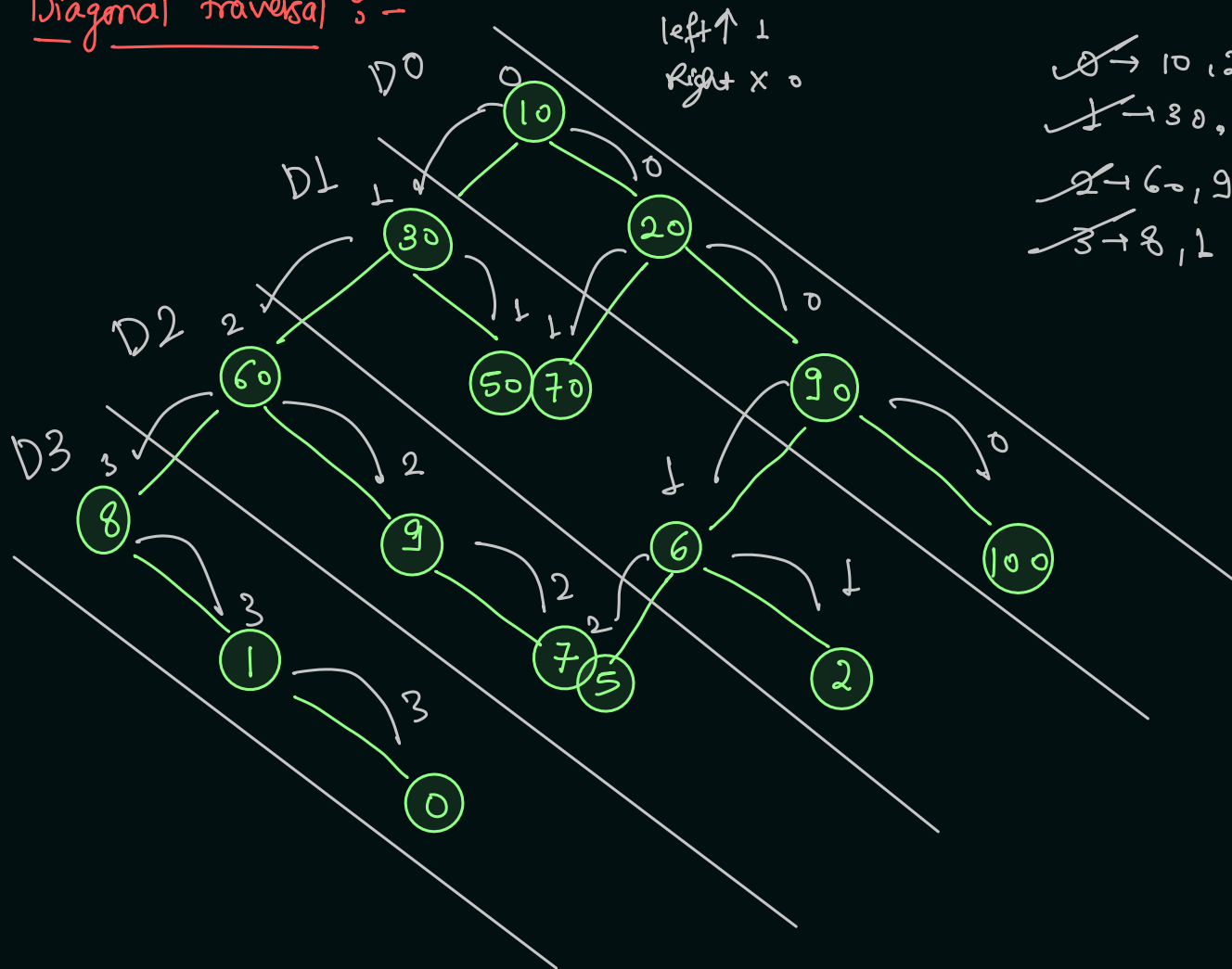
For a particular level \rightarrow first priority \rightarrow vertical index

If v.index is same then value is on priority.

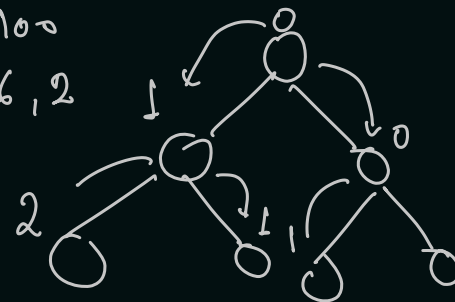


on the basis of
priority -
priority \rightarrow Min.

Diagonal traversal :-



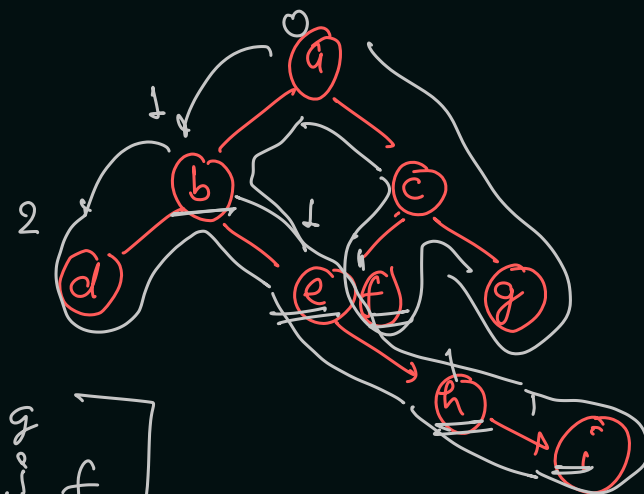
0 → 10, 20, 90, 100
1 → 30, 50, 70, 6, 2
2 → 60, 9, 7, 5
3 → 8, 1, 0



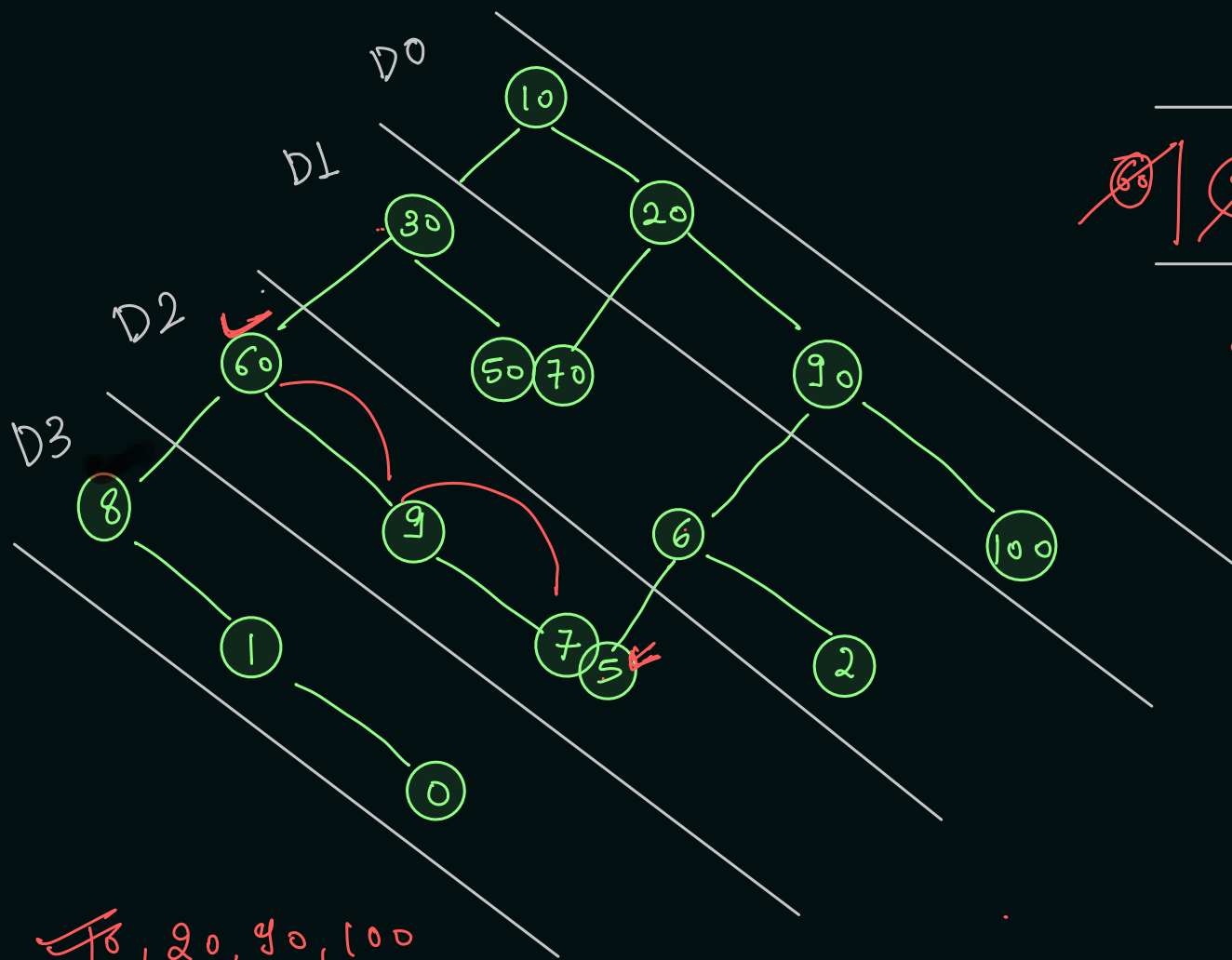
left (1) - increment
right no increment

BFS + Hash Map → Done

0 → a, c, g
1 → b, e, h, i, f
2 → d



DFS doesn't work



~~10~~ | ~~50~~ | 9

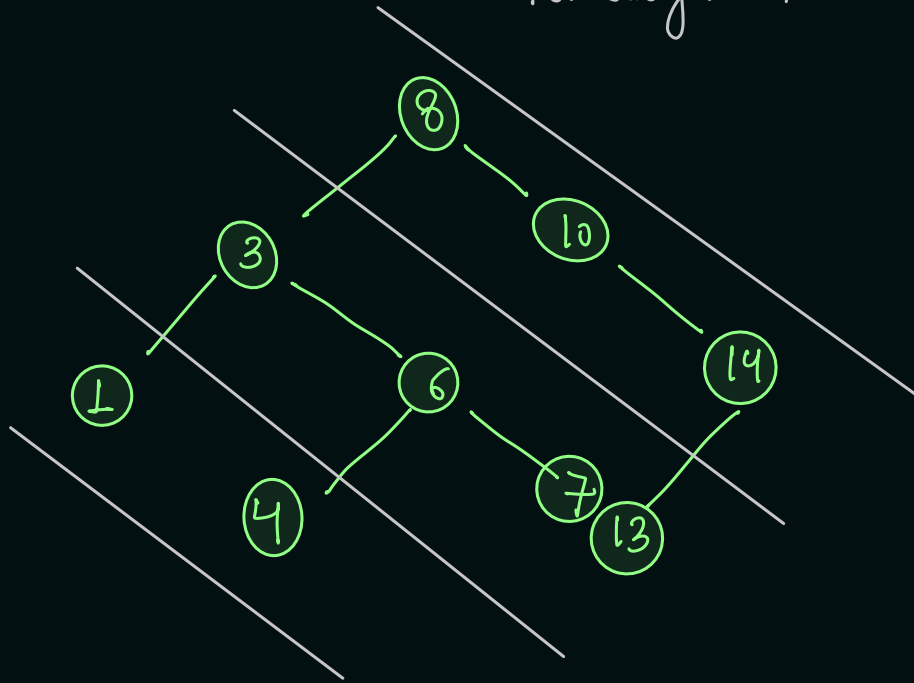
size 2

~~10, 20, 90, 100~~

~~30, 50, 70, 6, 2~~

~~60, 9, 7, 5~~

For diagonal, move toward Right & add left in queue.



```
public ArrayList<Integer> diagonal(Node root) {
    Queue<Node> que = new LinkedList<>();
    que.add(root);
    ArrayList<Integer> res = new ArrayList<>();
    while(que.size() > 0) {
        int sz = que.size();
        while(sz-- > 0) {
            Node rem = que.remove();
            while(rem != null) {
                res.add(rem.data);
                if(rem.left != null) que.add(rem.left);
                rem = rem.right;
            }
        }
    }
    return res;
}
```

~~1~~ ~~4~~

sz = 2 1 2 1

Rem = ~~8~~ 3

res = [8, 10, 14, 3, 6, 7, 13, 1, 4]

Anti diagonal → How?? → move toward left
→ add Right Node in Queue.

Morning → Tree Complete

Evening →] → AVL

weekday → Thursday] (AVL)

left Right Height]
L.H. — R.H.]

AVL

10

20

30

AVL

10

20

30

40

optimize - BST

40

10

20

30

40

50

week days

10

20

30

40

50

① Burning Tree ① 2 ②

② Max. path sum

③ Morris traversal → in & pre

④ BST iterator

creation

1 class
content

Thursday

AVL → Self Balancing
Binary Search
Tree

LL Rotation

RR Rotation

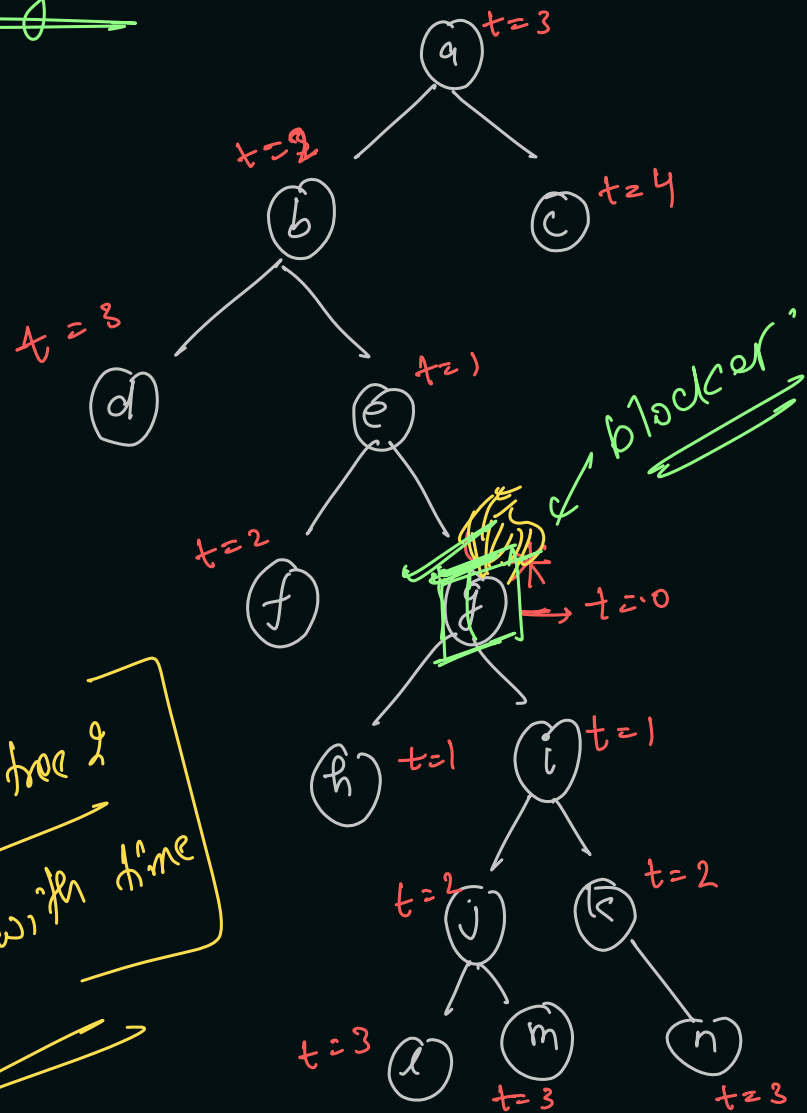
LR "

RL "

2 hrs

add → remove

Burning Tree



spreading Rate \rightarrow Adjacent node per second.

Max. time Required to complete
burning of tree.

Max time = 4 sec

Burning tree &
Node with time

Rotten Oranges

fire in the city

Infection spreading

graph

but Not
in tree

Node to Root path. \rightarrow

max depth \rightarrow
begin from root

g
0
(3)

e
1
(2)

b
2
(3)

a
3
(4)

Max = (4)

soln