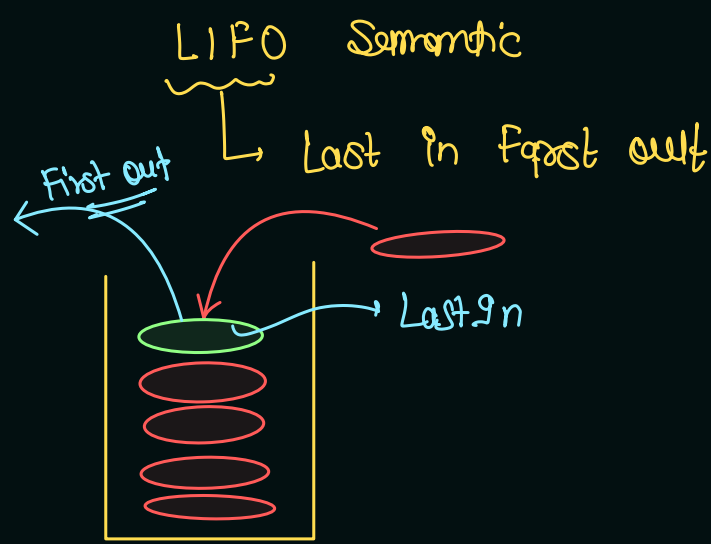


Stack

DS } Org. of Data
 { get
 { set
 { retrieve



Integer data

get → peek / top
add → push } add on top of stack
remove → pop();

Functions of Stacks →

Usage of Stack dependent
- Queue uses Stack
- Stack implementation generic
- Queue implementation generic
- Adapter

①

get] → st. peek() → $O(1)$

②

set] → add → st. push(val); → $O(1)$

remove → st. pop(); → $O(1)$

③

size → st. size(); → $O(1)$

④

isEmpty → st. isEmpty(); → $O(1)$
 { True
 { False

remove top data and return that value

Stack → horizontal [] ← add → remove

```
Stack<Integer> st = new Stack<>();
System.out.println(st.isEmpty());
st.push(10);
System.out.println(st);
st.push(20);
System.out.println(st);
st.push(30);
System.out.println(st);
st.push(40);
System.out.println(st);
st.push(50);
System.out.println("Size : " + st.size());
System.out.println(st);
System.out.println(st.pop());
System.out.println(st);
System.out.println(st.pop());
System.out.println(st);
System.out.println(st.isEmpty());
```

[] ✓

True ✓

✓ ✓

[10]; ✓

✓ ✓

[10, 20]; ✓

✓ ✓

[10, 20, 30]; ✓

✓ ✓

[10, 20, 30, 40]; ✓

✓ ✓

Size : 5 ✓

[10, 20, 30, 40, 50] ✓

pop → 50 ✓

[10, 20, 30, 40]; ✓

pop → 40 ✓

[10, 20, 30]; ✓

False ✓

[10, 20, 30]

Stack → st

Output is correct

1	true
2	[10]
3	[10, 20]
4	[10, 20, 30]
5	[10, 20, 30, 40]
6	Size : 5
7	[10, 20, 30, 40, 50]
8	50
9	[10, 20, 30, 40]
10	40
11	[10, 20, 30]
12	false

Duplicate brackets:

→ Expression is balanced

No. of opening brackets are equal to closing brackets & they are perfectly arranged.

① $(a+b) + (c+d) \rightarrow \underline{\text{True}}$

② $((a+b) + c) + d \rightarrow \text{False}$

③ $\Rightarrow ((a+b) + (c+d)) \rightarrow \text{True}$

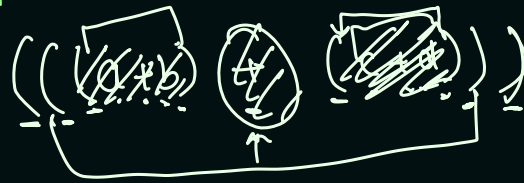
④ $((a+b) + c) + (d+e) + (f+g) \rightarrow \text{False}$

Example -

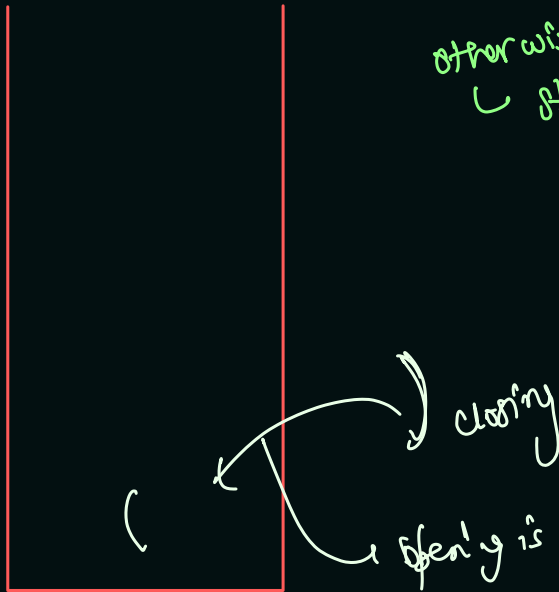
$(a+b) + (c+d)$ → Duplicate,

Action
└ closing bracket
otherwise
└ stack push

$((a+b) + c) + (d+e) + (f+g)$ → No duplicacy,



opening is at top] → duplicacy is present
other way



```

Stack<Character> st = new Stack<>();

for(int i = 0; i < str.length(); i++) {
    char ch = str.charAt(i);
    if(ch == ' ') continue;
    if(ch != ')') {
        st.push(ch);
    } else {
        if(st.peek() == '(') {
            return true;
        } else {
            while(st.peek() != '(') {
                st.pop();
            }
            st.pop();
        }
    }
}

return false;

```

~~((a+b)*c)~~ + ~~(d+e)~~ + ~~(f+g)~~

+

Balanced bracket

hint

L opening bracket
 L push

L closing
 L

some push
 opening
 peek() pop

Balanced brackets =

e.g.

$[(a + b) + \{(c + d) * (e / f)\}] \rightarrow \text{true}$

$[(a + b) + \{(c + d) * (e / f)\}] \rightarrow \text{false}$

$[(a + b) + \{(c + d) * (e / f)\}] \rightarrow \text{false}$

$[(a + b) + \{(c + d) * (e / f)\}] \rightarrow \text{false}$

opening bracket \rightarrow push

closing bracket \rightarrow check if counter
opening is available
or not

op: (\rightarrow) close

$\{ \rightarrow \}$

$[\rightarrow]$

NOTE!

if at last $\text{str.size()} > 0$
that means there
are more opening
bracket then closing
bracket

$(\underline{[\underline{\{ a + b \}}] }) \rightarrow \text{OR}$

$(\underline{[\underline{a + b}] + c }) \underline{] } \}$

closing \rightarrow if $\text{str.size()} = 0$

More closing bracket

$\frac{E}{L}$

next greater on Right:

$\xrightarrow{\text{Right max}}$ $\xrightarrow{14}$ $\text{next greater} = 12$

0	1	2	3	4	5	6	7	8
10	6	12	5	3	11	14	8	9

Right
max

next
greater

next
greater
on Right

value
= next
greater

0	1	2	3	4	5	6	7	8
2. 12	2. 12	6. 14	5. 11	5. 11	6. 14	-1	8. 9	-1

8-9	8-9
7-8	7-8
6-14	6-14
5-11	5-11
4-3	4-3
3-5	3-5
2-12	2-12
1-6	1-6
0-10	0-10

Index - value

Stack
Maintain \rightarrow decreasing
order of
value.

no pair \rightarrow Index carry \rightarrow we can access value as well

array →

0	1	2	3	4	5	6	7	8
2	5	9	3	1	12	6	8	7

ngr res →

0	1	2	3	4	5	6	7	8
5	9	12	12	12	-1	8	-1	-1

n = 9

```

int n = arr.length;
int[] res = new int[n];
Stack<Integer> st = new Stack<>();
st.push(0); // stack hold index of values
for(int i = 1; i < n; i++) {
    // pop smaller value from stack and mark their next greater
    while(st.size() > 0 && arr[st.peek()] < arr[i]) {
        int indx = st.pop();
        res[indx] = arr[i]; // place value in res
    }
    st.push(i);
}
while(st.isEmpty() != true) {
    res[st.pop()] = -1;
}
return res;

```

complexity →

$O(2n)$ max
have 841
 $O(n)$

8-7	8-7
7-8	7-8
6-6	6-6
5-12	5-12
4-1	4-1
3-3	3-3
2-9	2-9
1-5	1-5
0-2	0-2

Stack - st

3-3
2-9
1-5
0-2

next greater / smaller } left
Right

Arrays : [10, 6, 12, 5, 3, 11, 14, 8, 9]

ngr -> [12, 12, 14, 11, 11, 14, -1, 9, -1]

nsr -> [6, 5, 5, 3, -1, 8, 8, -1, -1]

ngl -> [-1, 10, -1, 12, 5, 12, -1, 14, 14]

ns1 -> [-1, -1, 6, -1, -1, 3, 11, 3, 8]

Stock Span:

Span \rightarrow No. of days before today, when stock price was higher than today.

for the array [2 5 9 3 1 12 6 8 7]

✓ span for 2 is 1

✓ span for 5 is 2

✓ span for 9 is 3

✓ span for 3 is 1

✓ span for 1 is 1

✓ span for 12 is 6

✓ span for 6 is 1

✓ span for 8 is 2

✓ span for 7 is 1

next greater on left

price

nagl. index

no. of days \rightarrow $\text{indx} - \text{naglindex}$

0	1	2	3	4	5	6	7	8
2	5	9	3	1	12	6	8	7
-1	-1	-1	2	3	-1	5	5	7
1	2	3	1	1	6	1	2	1

nagl not exist \rightarrow -1

for final Res \rightarrow $\text{nagl}[i] = i - \text{nagl}[i]$

Daily temperature:

0 1 2 3 4 5 6 7
[73, 74, 75, 71, 69, 72, 76, 73]

ngr ind → 1 2 6 5 5 6 6 7

res → 1 1 4 2 1 1 0 0

$\text{ngr}[i] - i$
diff

next greater index

next greater on right index

res $\left[\text{ngr}[i] - \text{ngr}[i] - i \right]$

if ngr is not available res[i] = i