

11-09 Doubt Session

Sunday, 11 September 2022 10:12 AM

$$T(n) = n * T(n-1) \quad (1)$$

$$T(n-1) = (n-1) * T(n-2) \quad (2)$$

2 into 1

$$T(n) = n * (n-1) * T(n-2)$$

Putting $n=2$ in 1 $T(n-2) = (n-2) * T(n-3) \quad (3)$

$$T(n) = n * (n-1) * (n-2) * T(n-3)$$

$$= n * (n-1) * (n-2) * (n-3) * \dots * T(n-n)$$

$$= n * (n-1) * (n-2) * (n-3) * \dots * T(1)$$

$$= n * (n-1) * (n-2) * (n-3) * \dots * 1$$

$$= n * n * \left(1 - \frac{1}{n}\right) * n * \left(1 - \frac{2}{n}\right) * n * \left(1 - \frac{3}{n}\right) * \dots * n * \left(\frac{2}{n}\right) * n * \left(\frac{1}{n}\right)$$

$$= n * n * n * n * n * n * \dots$$

$$= n^n$$



$$2 + 2 + 3$$

$$2 + 2$$

1. O notation:-

- This notation is used to express upper bound of an algorithm's running time. tight
- It represents the worst case of algo's time complexity. i.e. the longest amount of time an algo possibly can take.

$$f(n) = O(g(n)) ; \text{ if there exist a value } n_0 \text{ and } c > 0 \text{ such that } f(n) \leq c g(n) \text{ for all } n > n_0$$

$$f(n) \leq c g(n) \quad \forall n > n_0$$

$$f(n) = 2n^2 + 5n + 1 \quad g(n) = n^2$$

$$g(n) = n^2$$

$$f(n) \leq c g(n)$$

$$2n^2 + 5n + 1 \leq 3n^2 \quad \text{for } n=1, 2+5+1=8 \leq 3 \times 1=3 \quad \text{false}$$

$$n=1 \quad 8 \leq 3 \quad \text{false}$$

$$2n^2 + 5n + 1 \leq 8n^2$$

$$n=1$$

$$2+5+1=8$$

$$8 \leq 8 \quad \text{true}$$

$$C=8$$

$$g(n) = n^2$$

$$2n^2 + 5n + 1 \leq 4n^2$$

$$n=1$$

$$8 \leq 4 \quad \text{false}$$

$$n=6$$

$$2n^2 + 5n + 1 \leq 3n^2$$

$$n=1 \quad 2+5+1=8 \leq 3 \quad \text{false}$$

$$2n^2 + 5n + 1 \leq 3n^2$$

$$\boxed{n \geq 6}$$

$$\frac{73}{30}$$

$$2(6) + 5(6) + 1 \leq 3(6)$$

$$72 + 30 + 1 \leq 108$$

$$103 \leq 108 \Rightarrow$$

$$n=5$$

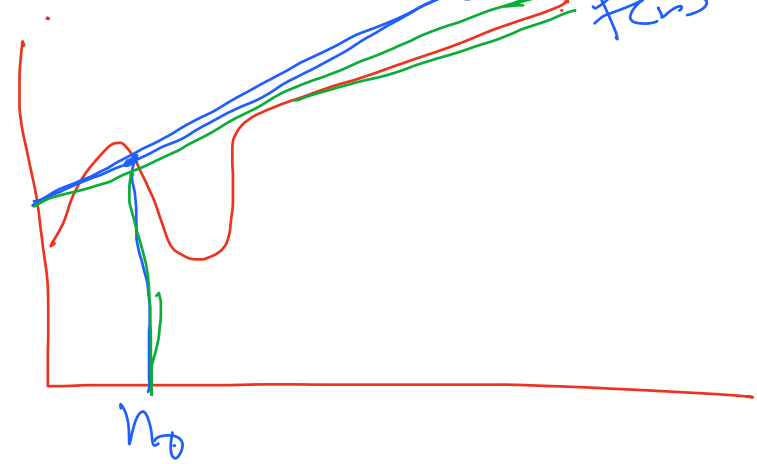
$$2(25) + 25 + 1 \leq 6(25)$$

$$50 + 25 + 1 \leq 150 \Rightarrow C=3$$

$$n=4$$

$$2(16) + 20 + 1 \leq 6(16)$$

$$32 + 20 + 1 \leq 96$$



$$n=1 \quad 8 \leq 4$$

$$2n^2 + 5n + 1 \leq 8n^2$$

$$n=1$$

$$2 + 5 + 1 \leq 8$$

$$8 \leq 8$$

$$C=8$$

$$g(n) = n^3$$

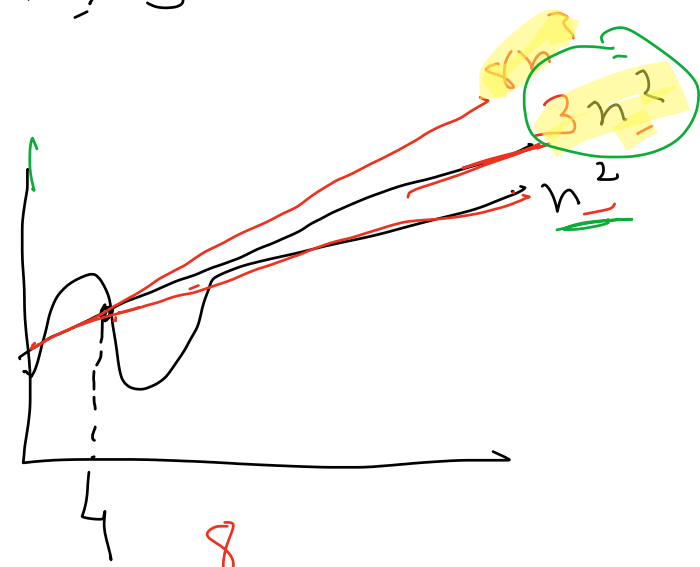
$$n \geq 1$$

$$2n^2 + 5n + 1 \leq 3n^3$$

$$C=3$$

$$g(n) = n^2$$

$$n \geq 5$$



$$(a) \quad C = g(n)$$

$$(b) \quad C = g(n)$$

$$(c) \quad C = g(n)$$

$$(d) \quad C = g(n)$$

$$Q \quad f(n) = 2n^2 + 5n + 1$$

$$g(n) = n^3$$

$$C =$$

$$2n^2 + 5n + 1 \leq C n^3$$

$$2n^2 + 5n + 1 \leq C n^3$$

$$C=8$$

$$C=4$$

$$2n^2 + 5n + 1 \leq 4n^3$$

$$n=1$$

$$2 + 5 + 1 \leq 4$$

$$8 \leq 4 \quad \text{false}$$

$$n=1$$

$$2 + 5 + 1 \leq 4 \quad \text{false}$$

$$4n^3$$

$$n=1$$

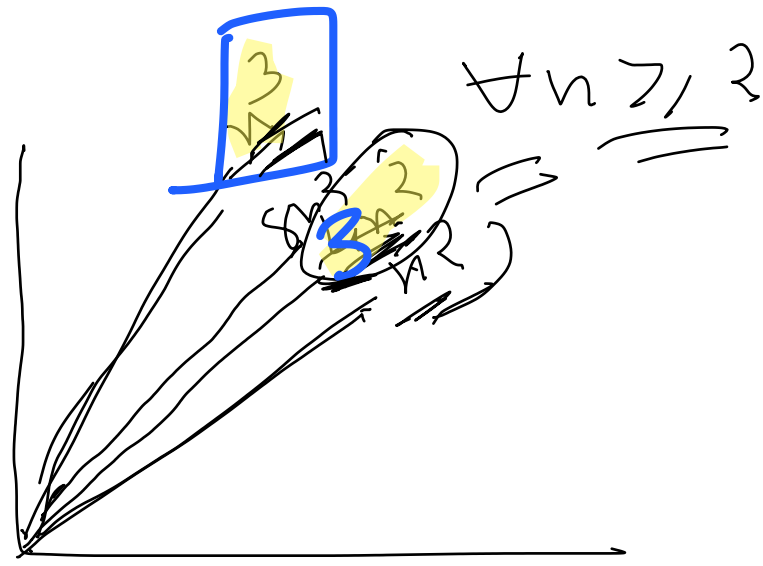
$$n \geq 1$$

$$n=2$$

$$2 \times 4 + 5 \times 2 + 1 \leq 4 \times 8$$

$$8 + 10 + 1 \leq 32 \quad \text{true}$$

yw



$$C = 8, n = 1 \quad 8n^3 \geq 4n^3$$

(1) Loop:-

```
for(i=1; i<=n; i++)
{
    x = y + 2 // C
}
```

$O(Cn)$
 $O(\underline{\underline{Cn}})$

(2) nested loop:

```
n times {
    for(i=1; i<=n; i++)
    {
        for(j=1; j<=n; j++) { // n times
            x = y + 2 // constant time
        }
    }
}
```

$O(Cn \times n)$ $O(Cn^2)$
 $O(\underline{\underline{Cn^2}})$

(3) Sequential

```
a = a + b; // C1
for(i=1; i<=n; i++) // n
{
    x = y + z; // C2
}
for(j=1; j<=n; j++) // n
{
    c = d + e; // C3
}
```

$C_1 + nC_2 + nC_3$ $\underline{\underline{n}}$ $O(n)$

(4) if-else

```
if (condition)
{
    // ... // O(n)
}
else
{
    // ... // O(n^2)
}
```

$O(n^2)$ $O(\underline{\underline{Cn^2}})$

Q A child is running up staircase with n steps and can hop either 1 step, 2 steps or 3 steps at a time. Implement a method to count

how many possible ways a child can run up the stairs.

$n=1$ } possible no. of :- 1 (step of 1)
case

$n=2$ } 1+1 (2 steps of 1 each)
2 } 2 (1 step of 2)

$n=3$ } 1+1+1 (3 steps of 1 each)
4 } { 2+1 (first \rightarrow 2, second 1)
1+2
3 (one step of 3) \rightarrow 1 } } = 4

$n=4$ } 1+1+1+1, 1+1+2, 1+2+1, 2+1+1,
7 } 3+1, 1+3, 2+2

JS

Adv. JS

Arrays

Alg. ana

function findStep(n)

$n=3$

5

{ 101 010 000 }
100 001 }

function