

Quantum Error Correction

Sender
Alice

Noisy channel

Receiver
Bob

$$\begin{array}{ccccccc} 1 & 0 & 0 & 1 & \underline{1} & 0 & 1 & 0 \\ 1 & 0 & 0 & 1 & 0 & 0 & 1 & 0 \\ 1 & 0 & 1 & 1 & 0 & 0 & 1 & 0 \\ = & = & & & & & & \\ 1 & 0 & 0 & 1 & 0 & 1 & 1 & 0 \end{array}$$

consecutive bits
are corrupted

Bits is sent ✓

single bit error ✓

multiple bit error ✓

Burst error

$$\boxed{1011001}0 \quad \text{Parity bit}$$

$$\text{no. of 1} = 4$$

$$\text{10100010} \rightarrow \text{not acceptable}$$

Logical Bits

$$\begin{cases} 0_L = \underline{\underline{000}} \\ 1_L = \underline{\underline{111}} \end{cases}$$

Suppose you get : 101 $\xrightarrow{\text{then correct it to}}$ 111

Errors in quantum communication

Decoherence

$$\frac{|0\rangle + |1\rangle}{\sqrt{2}} \longrightarrow \frac{|0\rangle - |1\rangle}{\sqrt{2}}$$

Bit flips

continuous evolution of quantum states
also introduce phase errors

$$|\psi\rangle = \frac{|0\rangle + |1\rangle}{\sqrt{2}}; \quad |\psi'\rangle = e^{i\alpha} \left(\frac{|0\rangle + |1\rangle}{\sqrt{2}} \right)$$

Say, a phase perturbation converts it to

$$|\psi'\rangle = \frac{|0\rangle + e^{i\alpha} |1\rangle}{\sqrt{2}}$$

A msmt of σ_z will still give you
|0> or |1> with probability $\frac{1}{2}$

But if you make a msmt of σ_x
it will project ± 1 with probability

$$|\psi\rangle = \frac{|0\rangle + |1\rangle}{\sqrt{2}} = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 \\ 1 \end{pmatrix}$$

$$\begin{aligned} |\psi'\rangle &= \frac{|0\rangle + e^{i\alpha} |1\rangle}{\sqrt{2}} = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 \\ e^{i\alpha} \end{pmatrix} \\ &= \frac{1}{\sqrt{2}} e^{i\alpha/2} \begin{pmatrix} e^{-i\alpha/2} \\ e^{i\alpha/2} \end{pmatrix} \end{aligned}$$

$$\sigma_x = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}, \quad \lambda = \pm 1$$

$$\lambda = +1, \quad |\underline{\lambda=1}\rangle = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 \\ 1 \end{pmatrix}$$

$$\begin{aligned} P_{+1} &= | +1 \rangle \langle +1 | = \frac{1}{2} \begin{pmatrix} 1 & 1 \end{pmatrix} \begin{pmatrix} 1 \\ 1 \end{pmatrix} \\ &= \frac{1}{2} \begin{pmatrix} 1 & 1 \\ 1 & 1 \end{pmatrix} \end{aligned}$$

$$P_{+1} |\psi'\rangle = \frac{1}{\sqrt{2}} \begin{pmatrix} \cos \alpha/2 \\ \cos \alpha/2 \end{pmatrix} e^{i\alpha/2} = \underline{\underline{\cos \alpha/2 e^{i\alpha/2}}} \boxed{\frac{1}{\sqrt{2}} \begin{pmatrix} 1 \\ 1 \end{pmatrix}}$$

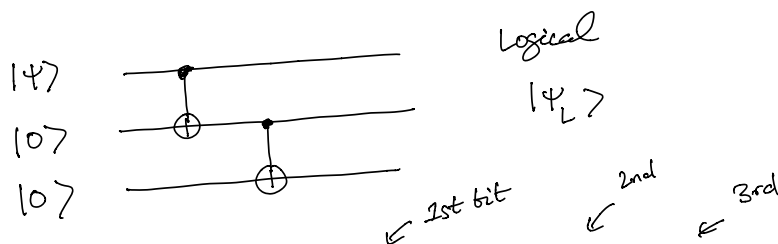
$$\cos^2 \frac{\alpha}{2}$$

$$a_L = 000$$

Sender

Alice

$$|\psi\rangle = \alpha|0\rangle + \beta|1\rangle$$



$$|\psi\rangle_L = (\alpha|0\rangle + \beta|1\rangle) |0\rangle |0\rangle$$

$$= (\alpha|00\rangle + \beta|11\rangle) |0\rangle$$

$$|\psi_L\rangle = \alpha|000\rangle + \beta|111\rangle$$

Alice is sending ψ_L

Say, noise in the channel flips a bit with probability p and leaving it unchanged is $1-p$

The state received by Bob via the noisy channel is $|\psi_2\rangle_L$

We will have the following possibilities

Bob receives

$$\alpha|000\rangle + \beta|111\rangle$$

$$\alpha|100\rangle + \beta|011\rangle$$

$$\alpha|010\rangle + \beta|101\rangle$$

$$\dots \dots \dots \alpha|110\rangle$$

Probability of occurrence

$$(1-p)^3$$

$$p(1-p)^2$$

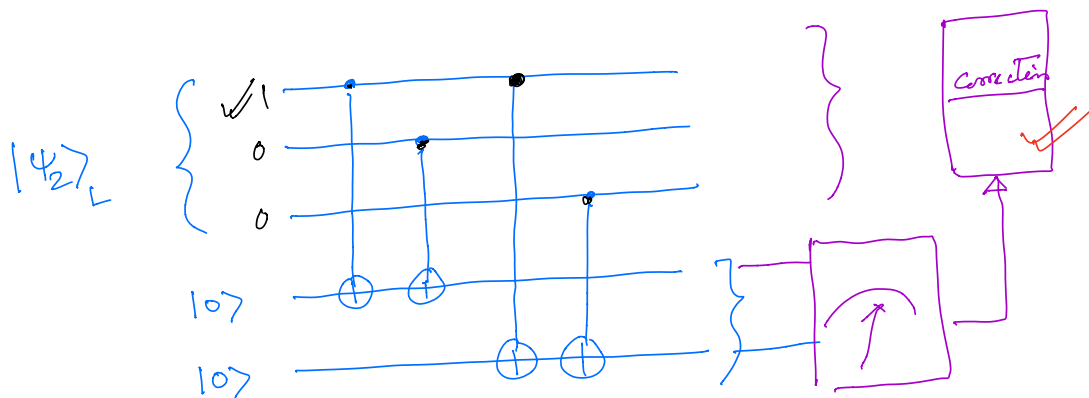
$$p(1-p)^2$$

$$p(1-p)^2$$

$$\propto |001\rangle + p|111\rangle \quad \text{Total} = 3p(1-p)^2$$

$$\begin{aligned} &\propto |110\rangle + \beta|001\rangle \\ &\propto |101\rangle + \beta|010\rangle \\ &\propto |011\rangle + \beta|100\rangle \\ &\propto |111\rangle + \beta|000\rangle \end{aligned} \quad \left. \begin{array}{l} \rightarrow p^2(1-p) \\ \text{Total} = \frac{3p^2(1-p)}{p^3} \end{array} \right\}$$

$$|\psi\rangle_L \quad \left\{ \begin{array}{l} \text{---} \\ \text{---} \\ \text{---} \\ 10\rangle \\ 10\rangle \end{array} \right.$$



$$(\alpha|000\rangle + \beta|111\rangle)|00\rangle = (\alpha|000\rangle + \beta|111\rangle)|00\rangle$$

\Rightarrow ancilla bits are not affected at all.

$$|\psi_2\rangle_L = \underline{(\alpha|100\rangle + \beta|011\rangle)} \underline{|00\rangle}$$

=

Bob's action results in $|4_3\rangle_L$

State received	After Bob's gate operation	Probability
$ 4_2\rangle_L$	$ 4_3\rangle_L$	
$\alpha 000\rangle + \beta 111\rangle$	$(\alpha 000\rangle + \beta 111\rangle) \underline{ 00\rangle}$	$(1-p)^3$
$\alpha 100\rangle + \beta 011\rangle$	$(\alpha 100\rangle + \beta 011\rangle) 11\rangle$	$p(1-p)^2$
$\alpha 010\rangle + \beta 101\rangle$	$(\alpha 010\rangle + \beta 101\rangle) 10\rangle$	$p(1-p)^2$
$\alpha 001\rangle + \beta 110\rangle$	$(\alpha 001\rangle + \beta 110\rangle) \underline{ 01\rangle}$	$p(1-p)^2$
$\alpha(1110\rangle + \beta 001\rangle)$	$(\alpha 1110\rangle + \beta 001\rangle) \underline{ 01\rangle}$	$p^2(1-p)$
$\alpha 101\rangle + \beta 010\rangle$	_____ $ 10\rangle$	$p^2(1-p)$
$\alpha 011\rangle + \beta 100\rangle$	_____ $ 11\rangle$	$p^2(1-p)$
$\alpha 111\rangle + \beta 000\rangle$	_____ $\underline{ 00\rangle}$	p^3

msmt of ancilla may yield
 $|00\rangle, |01\rangle, |10\rangle, |11\rangle$

Say Bob get $|00\rangle$
 $\alpha|000\rangle + \beta|111\rangle$ with prob. $(1-p)^3$ (without errors)
 $\alpha|111\rangle + \beta|000\rangle$ with prob. p^3 (errors)

✓✓ Bob takes no action

Now, say Bob gets $|01\rangle$
 $\alpha|001\rangle + \beta|110\rangle$ with prob. $p(1-p)^2$ (with one error)
 $\alpha|110\rangle + \beta|001\rangle$ with prob. $p^2(1-p)$ (with two errors)

h.o.

Then

Bob applies a σ_x on the
third qubit

$$\begin{aligned}\sigma_x \xrightarrow{\text{3rd}} (\alpha |001\rangle + \beta |110\rangle) &\longrightarrow \underline{\underline{\alpha |000\rangle + \beta |111\rangle}} \\ \sigma_x \xrightarrow{\text{3rd}} (\alpha |110\rangle + \beta |001\rangle) &\longrightarrow \underline{\underline{\alpha |111\rangle + \beta |000\rangle}}\end{aligned}$$

One is getting corrected fully
while the other one is fully getting wrong!