

*A Mini Project report on*

## **RABIN CRYPTOSYSTEM**

*submitted in partial fulfillment of the course*

CSE-1007: Introduction to Cryptography

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## **Introduction**

The Rabin cryptosystem is an asymmetric cryptographic technique, whose security, like that of RSA, is related to the difficulty of factorization. However, the Rabin cryptosystem has the advantage that the problem on which it relies has been proved to be as hard as integer factorization, which is not currently known to be true of the RSA problem. It has the disadvantage that each output of the Rabin function can be generated by any of four possible inputs; if each output is a cipher text, extra complexity is required on decryption to identify which of the four possible inputs was the true plaintext. The process was published in January 1979 by Michael O. Rabin. The Rabin cryptosystem was the first asymmetric cryptosystem where recovering the entire plaintext from the cipher text could be proven to be as hard as factoring.

## Theoretical study

Like all asymmetric cryptosystems, the Rabin system uses a key pair: a public key for encryption and a private key for decryption. The public key is published for anyone to use, while the private key remains known only to the recipient of the message.

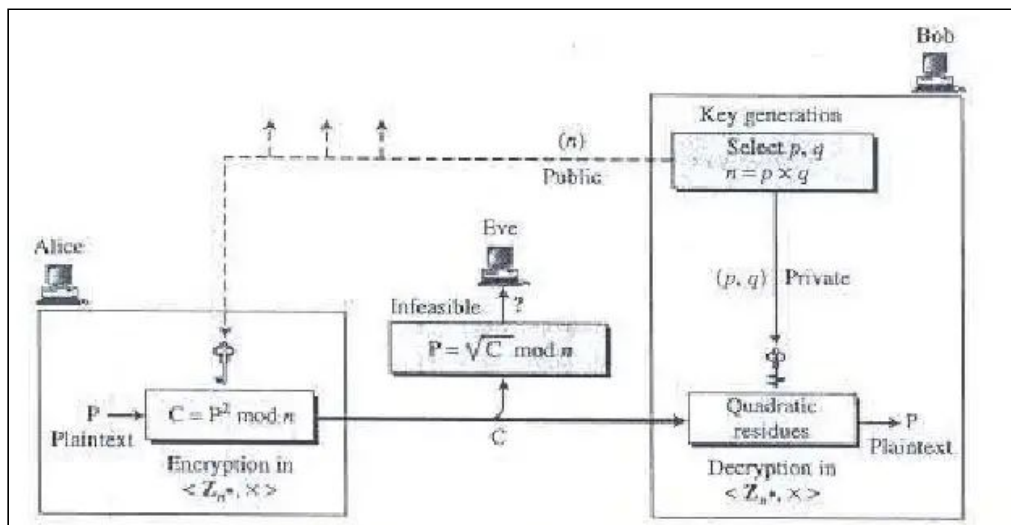


Figure 1: Rabin Encryption and Decryption

### Key generation

The keys for the Rabin cryptosystem are generated as follows:

1. Choose two large distinct prime numbers  $p$  and  $q$  such that  $p \equiv 3 \pmod{4}$  and  $q \equiv 3 \pmod{4}$
2. Compute  $n = p \times q$ .

Then  $n$  is the public key and the pair  $(p, q)$  is the private key.

### Encryption

A message  $M$  can be encrypted by first converting it to a number  $m < n$  using a reversible mapping, then computing  $c = m^2 \bmod n$ . The ciphertext is  $c$ .

### Decryption

The message  $m$  can be recovered from the ciphertext  $c$  by taking its square root modulo  $n$  as follows.

Compute the square root  $c$  of modulo  $p$  and  $q$  using these formulas:

$$m_p = c^{1/4(p+1)} \bmod p$$

$$m_q = c^{1/4(q+1)} \bmod q$$

Use the extended Euclidean algorithm to find  $y_p$  and  $y_q$  such that  $y_p \cdot p + y_q \cdot q = 1$ .

Use the Chinese remainder theorem to find the four square roots of  $c$  modulo  $n$ :

$$r_1 = (y_p \cdot p \cdot m_q + y_q \cdot q \cdot m_p) \bmod n$$

$$r_2 = n - r_1$$

$$r_3 = (y_p \cdot p \cdot m_q - y_q \cdot q \cdot m_p) \bmod n$$

$$r_4 = n - r_3$$

One of these four values is the original plaintext  $m$ , although which of the four is the correct one cannot be determined without additional information.

Computing square roots

We can show that the formulas in step 1 above actually produce the square roots of  $c$  as follows. For the first formula, we want to prove that  $m_p^2 \equiv c \bmod p$ . Since  $p \equiv 3 \bmod 4$  the exponent  $1/4(p+1)$  is an integer. The proof is trivial if  $c \equiv 0 \bmod p$  so we may assume that  $c$  does not divide  $p$ . Note that  $c \equiv m^2 \bmod pq$  implies that  $c \equiv m^2 \bmod q$ , so  $c$  is a quadratic residue modulo  $p$ . Then

$$m_p^2 \equiv c^{1/2(p+1)} \equiv c \cdot c^{1/2(p-1)} \equiv c \cdot 1 \bmod p$$

The last step is justified by Euler's criterion.

### Example

As an example, take

$P=7$  and  $q=11$ , then  $n=77$ . Take  $m=20$  as our plaintext.

The ciphertext is thus  $c = m^2 \bmod n = 400 \bmod 77 = 15$ .

Decryption proceeds as follows:

Compute

$$m_p = c^{1/4(p+1)} \bmod p = 15^2 \bmod 7 = 1$$

And

$$m_q = c^{1/4(q+1)} \bmod q = 15^3 \bmod 11 = 9$$

Use the extended Euclidean algorithm to compute  $y_p = -3$  and  $y_q = 2$ . We can confirm that  $y_p \cdot p + y_q \cdot q = (-3 \cdot 7) + (2 \cdot 11) = 1$ .

Compute the four plaintext candidates:

$$r_1 = (-3 \cdot 7 \cdot 9 + 2 \cdot 11 \cdot 1) \bmod 77 = 64$$

$$r_2 = 77 - 64 = 13$$

$$r_3 = (-3 \cdot 7 \cdot 69 - 2 \cdot 11 \cdot 1) \bmod 77 = 20$$

$$r_4 = 77 - 20 = 50$$

and we see that  $r_3$  is the desired plaintext. Note that all four candidates are square roots of 15 mod 77. That  $r_i^2 \bmod 77 = 15$  is, for each candidate, so each  $r_i$  encrypts to the same value, 15.



## ASCII Table

Decimal	Hexadecimal	Binary	Octal	Char	Decimal	Hexadecimal	Binary	Octal	Char	Decimal	Hexadecimal	Binary	Octal	Char
0	0	0	0	[NULL]	48	30	110000	60	0	96	60	1100000	140	`
1	1	1	1	[START OF HEADING]	49	31	110001	61	1	97	61	1100001	141	a
2	2	10	2	[START OF TEXT]	50	32	110010	62	2	98	62	1100010	142	b
3	3	11	3	[END OF TEXT]	51	33	110011	63	3	99	63	1100011	143	c
4	4	100	4	[END OF TRANSMISSION]	52	34	110100	64	4	100	64	1100100	144	d
5	5	101	5	[ENQUIRY]	53	35	110101	65	5	101	65	1100101	145	e
6	6	110	6	[ACKNOWLEDGE]	54	36	110110	66	6	102	66	1100110	146	f
7	7	111	7	[BELL]	55	37	110111	67	7	103	67	1100111	147	g
8	8	1000	10	[BACKSPACE]	56	38	111000	70	8	104	68	1101000	150	h
9	9	1001	11	[HORIZONTAL TAB]	57	39	111001	71	9	105	69	1101001	151	i
10	A	1010	12	[LINE FEED]	58	3A	111010	72	:	106	6A	1101010	152	j
11	B	1011	13	[VERTICAL TAB]	59	3B	111011	73	;	107	6B	1101011	153	k
12	C	1100	14	[FORM FEED]	60	3C	111100	74	<	108	6C	1101100	154	l
13	D	1101	15	[CARRIAGE RETURN]	61	3D	111101	75	=	109	6D	1101101	155	m
14	E	1110	16	[SHIFT OUT]	62	3E	111110	76	>	110	6E	1101110	156	n
15	F	1111	17	[SHIFT IN]	63	3F	111111	77	?	111	6F	1101111	157	o
16	10	10000	20	[DATA LINK ESCAPE]	64	40	1000000	100	@	112	70	1110000	160	p
17	11	10001	21	[DEVICE CONTROL 1]	65	41	1000001	101	A	113	71	1110001	161	q
18	12	10010	22	[DEVICE CONTROL 2]	66	42	1000010	102	B	114	72	1110010	162	r
19	13	10011	23	[DEVICE CONTROL 3]	67	43	1000011	103	C	115	73	1110011	163	s
20	14	10100	24	[DEVICE CONTROL 4]	68	44	1000100	104	D	116	74	1110100	164	t
21	15	10101	25	[NEGATIVE ACKNOWLEDGE]	69	45	1000101	105	E	117	75	1110101	165	u
22	16	10110	26	[SYNCHRONOUS IDLE]	70	46	1000110	106	F	118	76	1110110	166	v
23	17	10111	27	[ENG OF TRANS. BLOCK]	71	47	1000111	107	G	119	77	1110111	167	w
24	18	11000	30	[CANCEL]	72	48	1001000	110	H	120	78	1111000	170	x
25	19	11001	31	[END OF MEDIUM]	73	49	1001001	111	I	121	79	1111001	171	y
26	1A	11010	32	[SUBSTITUTE]	74	4A	1001010	112	J	122	7A	1111010	172	z
27	1B	11011	33	[ESCAPE]	75	4B	1001011	113	K	123	7B	1111011	173	{
28	1C	11100	34	[FILE SEPARATOR]	76	4C	1001100	114	L	124	7C	1111100	174	
29	1D	11101	35	[GROUP SEPARATOR]	77	4D	1001101	115	M	125	7D	1111101	175	}
30	1E	11110	36	[RECORD SEPARATOR]	78	4E	1001110	116	N	126	7E	1111110	176	~
31	1F	11111	37	[UNIT SEPARATOR]	79	4F	1001111	117	O	127	7F	1111111	177	[DEL]
32	20	100000	40	[SPACE]	80	50	1010000	120	P					
33	21	100001	41	!	81	51	1010001	121	Q					
34	22	100010	42	"	82	52	1010010	122	R					
35	23	100011	43	#	83	53	1010011	123	S					
36	24	100100	44	\$	84	54	1010100	124	T					
37	25	100101	45	%	85	55	1010101	125	U					
38	26	100110	46	&	86	56	1010110	126	V					
39	27	100111	47	'	87	57	1010111	127	W					
40	28	101000	50	(	88	58	1011000	130	X					
41	29	101001	51	)	89	59	1011001	131	Y					
42	2A	101010	52	*	90	5A	1011010	132	Z					
43	2B	101011	53	+	91	5B	1011011	133	[					
44	2C	101100	54	,	92	5C	1011100	134	\					
45	2D	101101	55	-	93	5D	1011101	135	]					
46	2E	101110	56	.	94	5E	1011110	136	^					
47	2F	101111	57	/	95	5F	1011111	137	_					

Table 1: ASCII Table

## **Pseudo code**

```
p= input (p)

q= input(q)

n=p*q

filename= input(Enter the file name in which plaintext is present)

call FileReading function and pass filename

plaintext = FileReading(filename)

print plaintext

FOR (i=0 to plaintext length)

    store each character in ch

    typecast ch in inetger and store in PlainText

    call Encrytion function and pass PlainText

    store return value in variable Cipher

    store Cipher in an integer array cip

END LOOP

print cip

store cip in string variable in ciphertext

filename=input(Enter the file name where the ciphertext is to be stored)

call FileWriting function and pass the filename and ciphertext

for(i=0 to plaintext length)

    call Decryption function and pass each value of cip array
```

```

END LOOP

print dec1,dec2,dec3,dec4

initialize decrypted_str

FOR (i=0 to plainttext length)

    IF(dec1 has character between A to Z or white space or between 0 to 9)

        THEN add that character to decrypted_str

    do same for dec2,dec3,dec3

END LOOP

filename=input("Enter the file name where decrypted text to be stored")

call FileWriting function and pass filename, decrypted_str

print decryted_str

FUNCTION FileReading(filename)

    scan the file

    read the plain text from file

    while (till last line )

        read each line

        store in a variable text

    END LOOP

    close file

    return text

END FUNCTION

```

FUNCTION FileWriting(filename,text)

    create a text file

    read the content from text

    print in the file

END FUNCTION

FUNCTION Encryption(PlainText)

    call FindT function and pass Plaintext,e,n

    store the return value in cipher

    return cipher

END FUNCTION

FUNCTION FindT(a,m,n)

    calculating  $a=(a^m) \% n$

    return a

END FUNCTION

FUNCTION Decryption(Cipher)

    initialize P1,P2,P3,P4,a1,a2,a3,a4

    calculate a1 and a3 using FindT function

    calculate  $a2 = p - a1$

    calculate  $b2 = q - b1$

    calculate P1,P2,P3,P4 using ChineseRemainderTheorem function

store all the value in p1,p2,p3,p4 array

convert p1,p2,p3,p4 to string and store in dec1,dec2,dec3,dec4  
respectively

END FUNCTION

FUNCTION ChineseRemainderTheorem(a,b,m1,m2)

intialize M, M1, M2, M1\_inv, M2\_inv

intialize result

$M = m1 * m2$

$M1 = M / m1$

$M2 = M / m2$

$M1\_inv = \text{inverse}(m1, M1)$

$M2\_inv = \text{inverse}(m2, M2)$

$\text{result} = (a * M1 * M1\_inv + b * M2 * M2\_inv) \% M$

return result

END FUNCTION

FUNCTION inverse(a,b)

int inv=0

intialize q, r, r1 = a, r2 = b, t, t1 = 0, t2 = 1

while (r2 greater than 0)

$q = r1 / r2$

$r = r1 - q * r2$

```

        r1 = r2

        r2 = r

        t = t1 - q * t2;

        t1 = t2

        t2 = t

    END LOOP

    if (r1 == 1)

        inv = t1

        if (inv < 0)

            inv = inv + a

        return inv

    END FUNCTION

```

## Implementation

```
package rabinusingascii;

import java.io.*;

import java.util.*;

import java.lang.*;

public class AsciiRabin {

    static int e = 2, n;

    static int p, q;

    static int[] cip=new int[100];

    static char[] ch=new char[100];

    static int[] newch=new int[100];

    static int[] p1=new int[100];

    static int[] p2=new int[100];

    static int[] p3=new int[100];

    static int[] p4=new int[100];

    static char ch1[]=new char[100];

    static char ch2[]=new char[100];

    static char ch3[]=new char[100];

    static char ch4[]=new char[100];

    static String dec1="";

    static String dec2="";
```

```

static String dec3="";

static String dec4="";

static int FindT(int a, int m, int n)
{
    int r;

    int[] y =new int[1] ;

    int[] x =new int[1] ;

    x[0]=a;

    y[0]=1;

    while (m > 0) {

        r = m % 2;

        FastExponention(r, n, y, x);

        m = m / 2;

    }

    return y[0];

}

static void FastExponention(int bit, int n, int[] y, int[] x)
{

    if (bit == 1)

        y[0] = (y[0] * x[0]) % n;

```



```

        x[0] = (x[0] * x[0]) % n;
    }

static int Encryption(int PlainText)
{
    int cipher = FindT(PlainText, e, n);

    return cipher;
}

static int inverse(int a, int b)
{
    int inv=0;

    int q, r, r1 = a, r2 = b, t, t1 = 0, t2 = 1;

    while (r2 > 0) {
        q = r1 / r2;

        r = r1 - q * r2;

        r1 = r2;

        r2 = r;

        t = t1 - q * t2;

        t1 = t2;

        t2 = t;
    }

    if (r1 == 1)

```

```

        inv = t1;

    if (inv < 0)

        inv = inv + a;

    return inv;

}

static void Decryption(int Cipher)

{

    int i=0;

    int P1, P2, P3, P4;

    int a1, a2, b1, b2;

    a1 = FindT(Cipher, (p + 1) / 4, p);

    a2 = p - a1;

    b1 = FindT(Cipher, (q + 1) / 4, q);

    b2 = q - b1;

    P1 = ChineseRemainderTheorem(a1, b1, p, q);

    P2 = ChineseRemainderTheorem(a1, b2, p, q);

    P3 = ChineseRemainderTheorem(a2, b1, p, q);

    P4 = ChineseRemainderTheorem(a2, b2, p, q);


    p1[i]=P1;

    p2[i]=P2;

```

```

    p3[i]=P3;

    p4[i]=P4;

    dec1+=(char)p1[i];

    dec2+=(char)p2[i];

    dec3+=(char)p3[i];

    dec4+=(char)p4[i];

    i++;
}

static int ChineseRemainderTheorem(int a, int b, int m1, int m2)
{
    int M, M1, M2, M1_inv, M2_inv;

    int result;

    M = m1 * m2;

    M1 = M / m1;

    M2 = M / m2;

    M1_inv = inverse(m1, M1);

    M2_inv = inverse(m2, M2);

    result = (a * M1 * M1_inv + b * M2 * M2_inv) % M;

    return result;
}

```

```

public static void FileWriting(String filename, String text) {

    PrintWriter outputstream=null;

    try {

        outputstream=new PrintWriter(new FileOutputStream(filename,false));

    }

    catch(FileNotFoundException e) {

        System.out.println("Error in opening file"+filename);

        System.exit(0);

    }

    outputstream.println(text);

    outputstream.close();

}

```

```

public static String FileReading(String filename) {

    Scanner inputstream=null;

    try {

        inputstream=new Scanner(new File(filename));

    }

    catch(FileNotFoundException e) {

        System.out.println("The file \""+filename+"\" could not be opened");

        System.exit(0);

    }

}

```

```

    }

    String text="";

    int c=0;

    System.out.println("\n");

    while(inputstream.hasNextLine()) {

        text=inputstream.nextLine();

    }

    inputstream.close();

    return text;

}

public static void main(String args[])

{

    Scanner sc=new Scanner(System.in);

    System.out.println("Enter the private keys p and q");

    p=sc.nextInt();

    q=sc.nextInt();

    n=p*q;

    System.out.println("Enter the filename that cointains the plaintext");

    sc=new Scanner(System.in);

    String filename=sc.nextLine();

    String plaintext=FileReading(filename);

```

```

System.out.println("Plaintext is = "+plaintext);

for(int i=0;i<plaintext.length();i++) {

char ch=plaintext.charAt(i);

int PlainText=(int)ch;

    int Cipher = Encryption(PlainText);

    cip[i]=Cipher;

}

System.out.print("Ciphertext is = ");

String ciphertext="";

for(int i=0;i<plaintext.length();i++) {

System.out.print(cip[i]+" ");

ciphertext+=cip[i]+" ";

}

System.out.println();

    System.out.println("Enter the filename where the ciphertext is to be
stored: ");

    sc=new Scanner(System.in);

    filename=sc.nextLine();

    FileWriting(filename,ciphertext);

for(int i=0;i<plaintext.length();i++) {

    Decryption(cip[i]);

```

```

    }

    System.out.println("dec1 = "+dec1);

    System.out.println("dec2 = "+dec2);

    System.out.println("dec3 = "+dec3);

    System.out.println("dec4 = "+dec4);


    String decrypted_str="";

    for(int i=0;i<plaintext.length();i++) {

        if((dec1.charAt(i)>='A'&&dec1.charAt(i)<='Z')||(dec1.charAt(i)==' '))
        {
            decrypted_str+=dec1.charAt(i);
        }

        if((dec2.charAt(i)>='A'&&dec2.charAt(i)<='Z')||(dec2.charAt(i)==' '))
        {
            decrypted_str+=dec2.charAt(i);
        }

        if((dec3.charAt(i)>='A'&&dec3.charAt(i)<='Z')||(dec3.charAt(i)==' '))
        {
            decrypted_str+=dec3.charAt(i);
        }
    }

```

```

        if((dec4.charAt(i)>='A'&&dec4.charAt(i)<='Z')||(dec4.charAt(i)=='
')||(dec4.charAt(i)>='0'&&dec4.charAt(i)<='9')) {

            decrypted_str+=dec4.charAt(i);

        }

    }

    System.out.println("Enter the filename where the decryptrd String is to be
stored: ");

    sc=new Scanner(System.in);

    filename=sc.nextLine();

    FileWriting(filename,decrypted_str);

    System.out.println("\nDecrypted String = "+decrypted_str);

}

}

```



## Results

```
Output - Project (run)

RUN:
Enter the private keys p and q
383
59
Enter the filename that contains the plaintext
plaintext.txt

Plaintext is = THIS IS RABIN CRYPTOSYSTEM 007
Ciphertext is = 7056 5184 5329 6889 1024 5329 6889 1024 6724 4225 4356 5329 6084 1024 4489 6724 7921 6400 7056 6241 6889 7921 6889 7056 4761 5529 1024 2304 2304 3025
Enter the filename where the ciphertext is to be stored:
ciphertext.txt
dec1 = T且向等口向等口端口，向性口川培培器T器端坡等T欲增D00德
dec2 = 庚HL I, XA星I机 CX等堆页航序页E管 007
dec3 = 口增城友整海友整口明E增W整空口ED0友友口限口整美离境
dec4 = 福Om斯as增B航+eL增>移ASYS移口所理理口
Enter the filename where the decrypted String is to be stored:
decrypted.txt

Decrypted String = THIS IS RABIN CRYPTOSYSTEM 007
BUILD SUCCESSFUL (total time: 28 seconds)
```

Figure 2: Output in NetBeans IDE




plaintext.txt - Notepad

File Edit Format View Help

THIS IS RABIN CRYPTOSYSTEM 007

Figure 3: Reading from text file plaintext.txt



ciphertext.txt - Notepad

File Edit Format View Help

7056 5184 5329 6889 1024 5329 6889 1024 6724 4225 4356 5329 6084 1024 4489 6724  
7921 6400 7056 6241 6889 7921 6889 7056 4761 5929 1024 2304 2304 3025

Figure 4: Writing text file ciphertext.txt



Figure 5: Writing text file in decrypted.txt

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