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THEORY OF COMPUTATION-1: (GATE 2022) - REPORTS

OVERALL ANALYSIS COMPARISON REPORT **SOLUTION REPORT**

ALL(17) CORRECT(10) INCORRECT(5) SKIPPED(2)

Q. 11

FAQ

Solution Video

Have any Doubt ?



Let us consider the pushdown automaton $\langle Q, \Sigma, \Gamma, q_0, Z_0, \{q_2\}, \delta \rangle$, where $Q = \{q_0, q_1, q_2\}$, $\Sigma = \{a, b, c\}$, $\Gamma = \{a, b, Z_0\}$ and let δ be as given in the following table.
[Note: σ represent alphabet a or b].

State	Input	Top of stack	Move
q_0	a	Z_0	(q_0, aZ_0)
q_0	b	Z_0	(q_0, bZ_0)
q_0	a	σ	$(q_0, a\sigma)$
q_0	b	σ	$(q_0, b\sigma)$
q_0	c	σ	(q_1, σ)
q_1	a	a	(q_1, ϵ)
q_1	b	b	(q_1, ϵ)
q_1	c	Z_0	(q_2, Z_0)

Which of the following language is accepted by the given PDA?

A $\{w c w^R \mid w \in (a, b)^* \mid |w| \geq 0\}$

B $\{w w^R \mid w \in (a, b)^*\}$

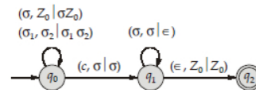
C $\{w c w \mid w \in (a, b)^* \mid |w| \geq 0\}$

D $\{w c w^R \mid w \in (a, b)^* \text{ and } |w| \geq 1\}$

Correct Option

Solution :
(d)

The transition diagram of the PDA is as shown below. In the figure σ , σ_1 and σ_2 represent a or b.



PDA accepting $\{w c w^R \mid w \in (a, b)^* \text{ and } |w| \geq 1\}$.

QUESTION ANALYTICS



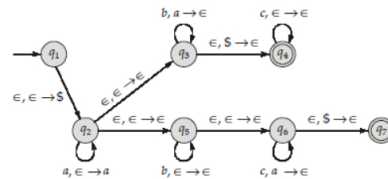
Q. 12

Solution Video

Have any Doubt ?



Consider the following automata that accepts the language L:



Which of the following is correct?

A $L = \{a^i b^j c^k \mid i, j, k \geq 1 \text{ and } i = j \text{ or } i = k\}$

B $L = \{a^i b^j c^k \mid i, j, k \geq 1 \text{ and } i = j \text{ and } i = k\}$

C $L = \{a^i b^j c^k \mid i, j, k \geq 0 \text{ and } i = j \text{ or } i = k\}$

Correct Option

Solution :
(c)

Given PDA is NPDA, hence two comparison with or is possible. If you observe the automata, at q_2 for every input 'a' there is two transition on epsilon which proves it is NPDA. The upper branch comparing a with b and lower branch comparing a with 'c' and both leads to the final states. L accepts epsilon as well. Hence $L = \{a^i b^j c^k \mid i, j, k \geq 0 \text{ and } i = j \text{ or } i = k\}$.

D $L = \{a^i b^j c^k \mid i, j, k \geq 0 \text{ and } i = j\}$

Q. 13

FAQ

Solution Video

Have any Doubt ?



Which of the following language generated by given grammar?

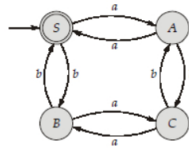
 $S \rightarrow aA \mid bB \mid \epsilon$ $A \rightarrow bC \mid aS$ $B \rightarrow aC \mid bS$ $C \rightarrow aB \mid bA$
☒ A $L = \{w : n_a(w) \text{ and } n_b(w) \text{ both are even}\}$

Your answer is Correct

Solution :

(a)

Since grammar is right linear regular grammar, convert it to machine.

Language is (a) $n_a(w)$ and $n_b(w)$ both are even.
☐ B $L = \{w : n_a(w) \text{ and } n_b(w) \text{ both are odd}\}$
☐ C $L = \{w : n_a(w) \text{ or } n_b(w) \text{ are even}\}$
☐ D $L = \{w : n_a(w) \text{ or } n_b(w) \text{ are odd}\}$

Q. 14

FAQ

Solution Video

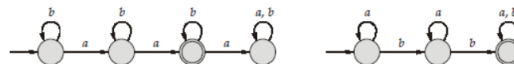
Have any Doubt ?

The minimum number of DFA states required for the languages $\{w \in (a, b)^* \mid \text{has exactly two } a\text{'s and at least two } b\text{'s}\}$ is _____.
☒ 10

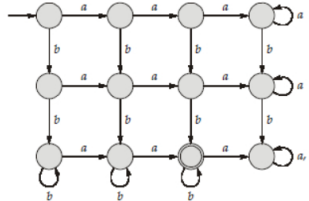
Correct Option

Solution :

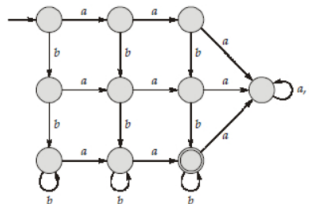
10

The following are DFAs for the two language $\{w \mid w \text{ has exactly two } a\text{'s}\}$ and $w \mid w \text{ has at least two } b\text{'s}\}$:

Combining them using the intersection construction gives the DFA:



Certain steps can be minimized



Hence total 10 states required.

Your Answer is 6

Q. 15

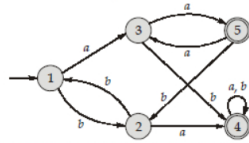
FAQ

Solution Video

Have any Doubt ?



Consider the following DFA:



The minimum number of states in the minimized DFA is _____.

5

Correct Option

Solution :

5

Make final non-final partition

[1, 2, 3] [4, 5]

Now, break because of $1 \xrightarrow{a} 3$ and 2, 3 goes to different states. So partition becomes

[1] [2, 3] [4, 5]

Break 4, 5 due to different transition on a . So partition becomes

[1] [2, 3] [4] [5]

Break [2] [3] because $2 \xrightarrow{a} 4$
 $3 \xrightarrow{a} 5$

So [1] [2] [3] [4] [5]

Minimum 5 states will be required.



Your Answer is 4



QUESTION ANALYTICS



Q. 16

FAQ

Have any Doubt ?



Consider the following grammar G:

 $S \rightarrow PQ$ $P \rightarrow aPb \mid ab$ $Q \rightarrow bQc \mid \epsilon$

Which of the following is true?

A $L(G) = \{a^m b^{m+n} c^n \mid m \geq 0, n \geq 0\}$ B $L(G) = \{a^m b^{m+n} c^n \mid m \geq 1, n \geq 0\}$

Your option is Correct

C $L(G) = \{a^m b^n c^{m+n} \mid m \geq 1, n \geq 0\}$ D The string "ab" belongs to $L(G)$

Your option is Correct

YOUR ANSWER - b,d

CORRECT ANSWER - b,d

STATUS - ✓

Solution :

(b, d)

 $P \rightarrow aPb \mid ab$ mean $P \rightarrow \{a^m b^m \mid m \geq 1\}$ $Q \rightarrow bQc \mid \epsilon$ mean $Q \rightarrow \{b^n c^n \mid n \geq 0\}$

So

 $S \rightarrow PQ \rightarrow \{a^m b^m b^n c^n \mid m \geq 1, n \geq 0\}$
 $\rightarrow \{a^m b^{m+n} c^n \mid m \geq 1, n \geq 0\}$

So option (b) is true.

Putting $m = 1$ and $n = 0$ "ab" belongs to $L(G)$ is true.

QUESTION ANALYTICS



Q. 17

FAQ

Solution Video

Have any Doubt ?



Which of the following statement(s) is/are true?

A If L is regular language and F is a finite language (i.e. a language with finite number of words) the $L \cup F$ must be a regular language.

Your option is Correct

B Regular expression that do not contain the star operator can represent infinite languages.


C Regular expression that do not contain the star operation can represent only finite language.

Your option is **Correct**

D If L is a regular language on F is a finite language then $L \cup F$ may or may not be regular language.

YOUR ANSWER - a,c

CORRECT ANSWER - a,c

STATUS - 

Solution :

(a, c)



QUESTION ANALYTICS



Item 11-17 of 17

 previous

1

2

next 