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[BOOKMARKS](#)
[MY PROFILE](#)
[REPORTS](#)
[BUY PACKAGE](#)
[NEWS](#)
[TEST SCHEDULE](#)

OPERATING SYSTEM-1 (GATE 2023) - REPORTS

[OVERALL ANALYSIS](#) [COMPARISON REPORT](#) **SOLUTION REPORT**

ALL(17) CORRECT(10) INCORRECT(3) SKIPPED(4)

Q. 1
[Have any Doubt ?](#)


Consider a disk in a system, P is the probability of occurrence of page fault on a request to disk. Memory access time is 10 μs and effective access time is 109 μs for this system. Page fault time is 1 ms for this system. The value of P is

 A 0.12

 B 0.1

Your answer is Correct

Solution :

(b)

$$\begin{aligned} \text{Effective access time} &= (1 - P) \times (\text{Memory access time}) + P \times (\text{Page fault time}) \\ \Rightarrow 109 \mu s &= (1 - P) \times 10 \mu s + P \times 1000 \mu s \\ \Rightarrow 109 &= 10 - 10P + 1000P \\ \Rightarrow 990P &= 99 \\ \Rightarrow P &= \frac{99}{990} = 0.1 \end{aligned}$$

 C 0.05

 D 0.01

QUESTION ANALYTICS


Q. 2
[Have any Doubt ?](#)


Which of the following is incorrect?

 A External fragmentation not possible in paging.

 B Internal fragmentation possible in paging.

 C Segmentation does not require memory compaction.

Your answer is Correct

Solution :

(c)

External fragmentation possible in segmentation, so it need compaction also.

 D On page fault, if page request is valid then data in page table changes on page fault handling.

QUESTION ANALYTICS


Q. 3
[Have any Doubt ?](#)


Consider a system using grouping method to implement free-space-list to keep track of free/ unallocated block of disk. Disk block size is 2 KB and length for address of a location in disk is 8 bytes. Assume, at an instance, 3 disk blocks are used for free-space-list and last block from these 3 block contains only 15 addresses of free blocks. How many blocks are currently free in the system?

 A 527

 B 524

 C 530

 D 525

Correct Option

Solution :

(d)

Block size (BS) = 2 KB

Address length (AL) = 8 bytes

$$\text{Maximum addresses can stored in a block} = \frac{\text{BS}}{\text{AL}} = \frac{2 \text{ KB}}{8 \text{ bytes}} = 2^8$$

First block in free-space-list contains $2^8 - 1$ address of free blocks and last address for next block in free-space-list.

Same way second block contains $2^8 - 1$ bytes.

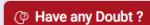
Now third which is last block contains only 15 address.

$$\begin{aligned} \text{So, Total free blocks} &= (2^8 - 1) + (2^8 - 1) + 15 \\ &= 255 + 255 + 15 = 525 \end{aligned}$$

 QUESTION ANALYTICS

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Q. 4



↗

In which of the following situation deadlock will definitely occur?

- A Each resource in system has single instance and current resource-allocation graph contains cycles.

Your answer is Correct

Solution :

- (a)
- (b) If multiple instance for resources, then cycle is necessary but not sufficient condition for deadlock.
- (c) Here, mutual exclusive property not hold, so no deadlock.
- (d) Here, no-preemption does not hold, so no deadlock.

- B Each resource in system has multiple instances and current resource-allocation graph contains cycles.

- C All resources in system are sharable.

- D All processes in system have different priorities and high priority process can preempt a lowpriority process.

 QUESTION ANALYTICS

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Q. 5



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Consider a disk queue in a system contains requests for following tracks:

60, 40, 145, 150, 85, 98

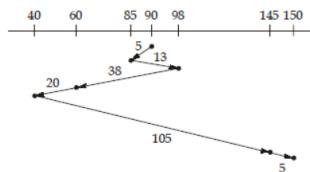
Disk has 200 tracks numbered from 0 to 199. Initially head of disk located at track number 90. Moving of head from one track to its adjacent track takes 1 ms. Disk uses shortest-seek-time-first scheduling algorithm. How much time will require to complete the requests in given queue?

- A 186 ms

Correct Option

Solution :

- (a)
- For SSTF, following will be order of request



Total tracks reached = $5 + 13 + 38 + 20 + 105 + 5 = 186$

Time for 186 tracks = $1 \times 186 \text{ ms}$

- B 160 ms

- C 180 ms

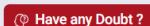
- D 176 ms

Your answer is IN-CORRECT

 QUESTION ANALYTICS

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Q. 6



↗

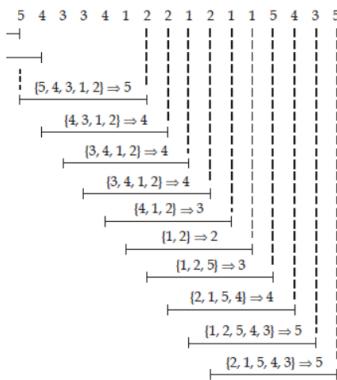
Consider a system uses working-set model for frame allocations. Let working-set window size is 7. Let following is sequence of virtual page references:
5, 4, 3, 3, 4, 1, 2, 2, 1, 2, 1, 1, 5, 4, 3, 5

Let x and y are minimum and maximum number of elements respectively in a working set on execution of above sequence. The value of $x + y$ is _____.

Solution :

7

Working set window (Δ) = 7



Here minimum cardinality working set is $\{1, 2\}$, so $x = 2$.

Maximum cardinality working sets have cardinality 5 and they are more than one.

So,

$$y = 5$$

$$x + y = 2 + 5 = 7$$

QUESTION ANALYTICS

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Q. 7

Have any Doubt ?

?

Consider a system uses First-In-First-Out (FIFO) page replacement policy. Let following is sequence of virtual page references:

1, 2, 3, 1, 4, 5, 2, 5, 3, 4, 1

If system has 3 frames available then x number of page faults will occur for above sequence. If system has 4 frames available then y number of page faults will occur for above sequence. Assume, initially all frames are empty. The value of $x - y$ is _____.

3

Your answer is Correct

Solution :

3

Number of frames = 3

1	2	3	1	4	5	2	5	3	4	1
			3	3	4	5	2	2	3	4
			2	2	2	3	4	5	5	2
1	1	1	X	2	X	4	X	X	X	3

F F F F F F F F F = 9 faults = x

Number of frames = 4

1	2	3	1	4	5	2	5	3	4	1
				4	5	5	5	5	5	1
				3	3	3	4	4	4	5
1	1	1	1	1	2	2	2	2	2	3

F F F F F F F F F = 6 faults = y

$x - y = 9 - 6 = 3$

QUESTION ANALYTICS

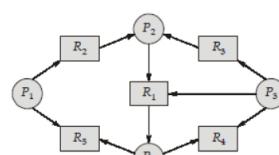
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Q. 8

Have any Doubt ?

?

Consider a system has four processes P_1, P_2, P_3, P_4 and five resources R_1, R_2, R_3, R_4, R_5 . At an instance, following is resource allocation graph of these processes and resources:



Let $G(V, E)$ is wait-for graph of above given resource-allocation graph. Which of the following is/are correct?

A G has nodes for P_1, P_2, P_3 , and P_4 processes

Correct Option

B G has nodes for R_1, R_2, R_3 resources

- D** (P₁, P₄) ∈ E and (P₂, P₄) ∈ E

YOUR ANSWER - NA

CORRECT ANSWER - a,c

STATUS - SKIPPED

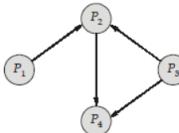
Solution :

(a,c)

In wait-for graph, there is no resource node. So (b) is incorrect.

 $P_i \rightarrow P_j$ is edge in wait-for graph iff there exist a resource R_k for which $P_i \rightarrow R_k$ and $R_k \rightarrow P_j$ are edges in resource-allocation graph.

Wait-for graph of above graph:



$$\{(P_1, P_2), (P_2, P_4), (P_3, P_2), (P_2, P_4)\} = E$$

$$\{P_1, P_2, P_3, P_4\} = V$$

QUESTION ANALYTICS

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Q. 9

Have any Doubt ?

□

Which of the following statements is/are correct for frame table in an operating system which uses paging system?

- A** Frame table maintains allocated frames.

Your option is Correct

- B** Frame table maintains free frames.

Correct Option

- C** If a frame is allocated, then frame table has an entry for process id to which it allocated.

Your option is Correct

- D** Frame table accessed and modified only by operating system.

Your option is Correct

YOUR ANSWER - a,c,d

CORRECT ANSWER - a,b,c,d

STATUS - ✘

Solution :

(a, b, c, d)

QUESTION ANALYTICS

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Q. 10

Have any Doubt ?

□

Consider the following statements regarding contiguous memory allocation methods. Which of the following statement is incorrect?

- A** There can be the situation where worst-fit algorithm allocates memory to all processes, but not possible for best-fit.

- B** There can be the situation where next-fit algorithm allocates memory to all processes, but not possible for best-fit.

- C** There can be the situation where first-fit algorithm allocates memory to all processes, but not possible for best-fit.

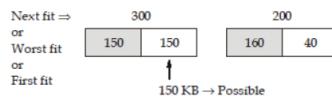
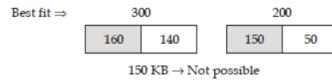
- D** None of the above is incorrect

Your answer is Correct

Solution :

(d)

Processes of size 150 KB, 160 KB, 150 KB in order, with 300 KB and 200 KB memory



QUESTION ANALYTICS

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