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## THEORY OF COMPUTATION-2: (GATE 2022) - REPORTS

OVERALL ANALYSIS    COMPARISON REPORT    **SOLUTION REPORT**

ALL(17)    CORRECT(11)    INCORRECT(5)    SKIPPED(1)

Q. 1

FAQ

Solution Video

Have any Doubt ?



The collection of decidable languages is not closed under the operation of

A Union

B Complementation

C Intersection

**D None of the above**

Your answer is Correct

**Solution :**

(d)  
Decidable languages are closed under union, complementation, concatenation, intersection, star operation.

QUESTION ANALYTICS



Q. 2

FAQ

Solution Video

Have any Doubt ?



Which of the following grammars is a (right) regular grammar with the same language as the regular expression  $a^* + b^* + ab$ ?

A  $S \rightarrow AB$   
 $A \rightarrow aA \mid \epsilon$   
 $B \rightarrow bB \mid \epsilon$

B  $S \rightarrow ab \mid A \mid B$   
 $A \rightarrow Aa \mid \epsilon$   
 $B \rightarrow Bb \mid \epsilon$

C  $S \rightarrow A \mid B$   
 $A \rightarrow aA \mid b$   
 $B \rightarrow Eb \mid a$

**D  $S \rightarrow ab \mid A \mid B$**   
 **$A \rightarrow aA \mid \epsilon$**   
 **$B \rightarrow bB \mid \epsilon$**

Your answer is Correct

**Solution :**

(d)

QUESTION ANALYTICS



Q. 3

FAQ

Solution Video

Have any Doubt ?



Let  $L_1$  and  $L_2$  be languages over  $\Sigma = \{a, b\}$  represented by the regular expressions  $(a^* + b)^*$  and  $(a + b)^*$  respectively. Which of the following is true with respect to the two languages?

A  $L_1 \subset L_2$

B  $L_2 \subset L_1$

**C  $L_1 = L_2$**

Your answer is Correct

**Solution :**

(c)  
Both  $L_1$  and  $L_2$  generate all the strings over 'a' and 'b'. Hence both are same.

D  $L_1 \cap L_2 = \phi$

Q. 4

FAQ

Solution Video

Have any Doubt ?



Let  $L_1$  be a decidable language and  $L_2$  be a Turing recognizable but not decidable language. Which of the following statements are true?

**A**  $L_2 \setminus L_1$  is a Turing recognizable language

Correct Option

**Solution :**  
(a)

**B**  $L_2 \cap L_1$  is a Turing decidable language

**C**  $L_2' \cap L_1'$  is a Turing recognizable language

**D**  $L_1 \setminus L_2$  is a Turing decidable language

Q. 5

Solution Video

Have any Doubt ?



Which of the following statements is true?

**A** The union of two context free languages is context free.

Your answer is Correct

**Solution :**  
(a)

- Context free languages are not closed under complementation and intersection. Hence option (b) and (c) is false.
- DPDA is less powerful than PDA. Hence there is CFL language which can not be accepted by DPDA. Hence option (d) is false.

**B** The intersection of two context free languages is context free

**C** The complement of a context free language is context free.

**D** If a language is context free, it can always be accepted by a deterministic pushdown automaton.

Q. 6

Solution Video

Have any Doubt ?



Consider the following statements:

$S_1$  : Language  $L$  is Turing-recognizable if a non-deterministic Turing Machine recognizes it.

$S_2$  : Every CFL is Turing-decidable.

$S_3$  :  $\text{REG}_{\text{TM}} = \{(M) \mid M \text{ is a TM and } L(M) \text{ is regular}\}$  is decidable.

The number of the correct statements is/are \_\_\_\_\_.

**B** 2

Correct Option

**Solution :**

- 2  
 $S_1$  : True  
 $S_2$  : True  
 $S_3$  : False, checking regularity of TM is undecidable.

Your Answer is 3

Q. 7

FAQ

Solution Video

Have any Doubt ?



Consider the following languages:

I.  $\{a^m b^n c^p d^q \mid m + p = n + q, \text{ where } m, n, p, q \geq 0\}$

- II.  $\{a^m b^n c^p d^q \mid mn = p + q, \text{ where } m, n, p, q \geq 0\}$   
 III.  $\{a^m b^n c^p d^q \mid m = n \text{ and } p = q, \text{ where } m, n, p, q \geq 0\}$   
 IV.  $\{a^m b^n c^p d^q \mid m = n = p \text{ and } p = q, \text{ where } m, n, p, q \geq 0\}$

How many of the above languages are context free?

2

Your answer is Correct

Solution :

2

- I.  $\{a^m b^n c^p d^q \mid m + p = n + q\}$  is CFL because we can rearrange the equation as  $m - n + p - q = 0$  which can be done by push, pop, push and pop respectively and check if stack is empty or not.  
 II.  $\{a^m b^n c^p d^q \mid mn = p + q\}$  is not a CFL, since  $mn$  involves multiplying number of a's and number of b's which cannot be done by a PDA.  
 III.  $\{a^m b^n c^p d^q \mid m = n \text{ and } p = q\}$  is CFL since one comparison at a time can be done by PDA.  
 IV.  $\{a^m b^n c^p d^q \mid m = n = p \text{ and } p = q\}$  is not CFL since  $m = n = p$  is a double comparison which can not be done by PDA.

QUESTION ANALYTICS

Q. 8

FAQ

Solution Video

Have any Doubt ?



Which of the following is/are correct for Turing recognizable?

- A. TM halts in an accepting configuration if  $w$  is in the language. Your option is Correct  
 B. TM may halt in a rejecting configuration or go on indefinitely if  $w$  is not in the language. Your option is Correct  
 C. TM halts in a rejecting configuration if  $w$  is not in the language. Your answer is IN-CORRECT  
 D. TM whether halts or not depends on the given input. Your option is Correct

YOUR ANSWER - a,b,c,d

CORRECT ANSWER - a,b,d

STATUS - ✖

Solution :

(a, b, d)

Turing-recognizable languages:

- TM halts in an accepting configuration if  $w$  is in the language.
- TM may halt in a rejecting configuration or go on indefinitely if  $w$  is not in the language.

Turing-decidable languages:

- TM halts in an accepting configuration if  $w$  is in the language.
- TM halts in a rejecting configuration if  $w$  is not in the language.

QUESTION ANALYTICS

Q. 9

FAQ

Solution Video

Have any Doubt ?



Consider  $L_1$  is regular language,  $L_2$  is context free language,  $L_3$  is recursive language and  $L_4$  is recursively enumerable language.

- A.  $\overline{L_3} \cup L_4$  is recursively enumerable language. Your option is Correct  
 B.  $L_1 \cdot L_2$  is regular language.  
 C.  $L_1^* \cap L_2$  is context free language. Your option is Correct  
 D.  $\overline{L_2} \cup L_3$  is context free language.

YOUR ANSWER - a,c

CORRECT ANSWER - a,c

STATUS - ✔

Solution :

(a, c)

- (a)  $\overline{L_3} \cup L_4 = \overline{REC} \cup RE = RE$   
 (b)  $L_1 \cdot L_2 = Reg \cdot CFL = CFL$   
 (c)  $L_1^* \cap L_2 = (Reg)^* \cap CFL = Reg \cap CFL = CFL$   
 (d)  $\overline{L_2} \cup L_3 = \overline{CFL} \cup REC = CSL \cup REC = REC$   
 So, only (a) and (c) are correct.

QUESTION ANALYTICS

Q. 10

[? FAQ](#)[▶ Solution Video](#)[🔔 Have any Doubt ?](#)

Let  $G_1$  and  $G_2$  be arbitrary context free languages and  $R$  be an arbitrary regular language. Consider the following problems:

I. Is  $L(G_1) = L(G_2)$ ?

II. Is  $L(G_2) \subseteq L(G_1)$ ?

III. Is  $L(G_1) = R$ ?

Which of the above problems are undecidable?

☐ A I only

☐ B II only

☐ C I and II only

☒ D I, II and III

Your answer is Correct

**Solution :**

(d)

Checking CFL is equivalence, equality and subset problems for CFL are undecidable.



QUESTION ANALYTICS



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