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## OPERATING SYSTEM (GATE 2023) - REPORTS

[OVERALL ANALYSIS](#) [COMPARISON REPORT](#) **SOLUTION REPORT**

ALL(33) [CORRECT\(23\)](#) [INCORRECT\(9\)](#) [SKIPPED\(1\)](#)

**Q. 11**

 [Have any Doubt ?](#)



Consider a system has four free non-adjacent memory partitions (holes) of size 200 KB, 400 KB, 600 KB and 300 KB. System allocate variable size partitions with contiguous memory allocation, and uses worst-fit strategy for memory allocation. Assume, system has memory requests from five processes  $P_0, P_1, P_2, P_3$  and  $P_4$  of size 245 KB, 50 KB, 175 KB, 190 KB and 210 KB respectively in given order ( $P_0$  requests first and  $P_4$  requests last). The size of free memory partition (hole) created on allocation of memory to process  $P_2$  is \_\_\_\_\_ KB.

180

Your answer is **Correct** 180

**Solution :**  
180

Initial	200 KB	400 KB	600 KB	300 KB
$P_0$			-245 KB	
$P_1$		400 KB -50 KB	355 KB	
$P_2$		350 KB	355 KB -175 KB	
$P_3$		350 KB -190 KB	180 KB	
$P_4$		160 KB		300 KB -210 KB
				90 KB

 [QUESTION ANALYTICS](#)



**Q. 12**

 [Have any Doubt ?](#)



Which of the following is/are correct?

**A** Two or more processes can concurrently execute the same program.

Your option is **Correct**

**B** A process and its child process share the code/program at physical memory because it is static.

**C** A process in user mode cannot execute privileged instructions.

Your option is **Correct**

**D** Threads of a process share the code/program at physical memory because it is static.

Your option is **Correct**

**YOUR ANSWER - a,c,d**

**CORRECT ANSWER - a,c,d**

**STATUS - ✓**

**Solution :**

(a, c, d)

- (a) Parent and child can execute the same program.
- (b) Child creation make copy of physical memory. Also consider exec() call.

 [QUESTION ANALYTICS](#)



**Q. 13**

 [Have any Doubt ?](#)



Consider a system has four processes  $P_0, P_1, P_2$  and  $P_3$ . System want to synchronize these processes, such that

- (i) First  $P_0$  start its body/code and finish it before all other processes.
- (ii)  $P_3$  start at last after all other processes finishes their execution.
- (iii)  $P_1, P_2$  execute code concurrently.

System uses three counting semaphores  $S_1, S_2$  and  $S_3$  with initial value of 0 for all ( $S_1 = 0, S_2 = 0, S_3 = 0$ ). Which of the following synchronize implementation/pseudocode is satisfying above condition?

<b>(a) <math>P_0</math> :</b>	<b><math>P_1</math> :</b>	<b><math>P_2</math> :</b>	<b><math>P_3</math> :</b>
{	{	{	{
code;	$P(S_1);$	$P(S_2);$	$P(S_3);$
$V(S_1);$	code;	code;	code;
}	$V(S_2);$	$V(S_3);$	}
	}	}	

<b>(b) <math>P_0</math> :</b>	<b><math>P_1</math> :</b>	<b><math>P_2</math> :</b>	<b><math>P_3</math> :</b>
{	{	{	{
code;	$P(S_1);$	$P(S_2);$	$P(S_3);$
....			

```

    V(S1);
    V(S2);
    V(S3);
}
(c) P0 : { code; } code; } code;
{ code;
V(S1);
V(S1);
}
P1 : { P(S1); code; } code;
V(S2);
}
P2 : { P(S1); code; } code;
V(S3);
}
P3 : { P(S2); code; }
}
(d) All of the above

```

(d) All of the above

A a

B b

C c

Your option is Correct

D d

YOUR ANSWER - c

CORRECT ANSWER - c

STATUS -

Solution :

- (a) P<sub>1</sub> and P<sub>2</sub> can execute only sequentially (P<sub>1</sub> then P<sub>2</sub>), but not concurrently.
- (b) Here P<sub>2</sub> can finish before P<sub>1</sub> and P<sub>2</sub>.
- (c) satisfy all conditions.

QUESTION ANALYTICS

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Q. 14

Have any Doubt ?

Q

Consider a Unix-like file system in which an index node (inode) of a file contains 10 direct, 2 single indirect, and 1 double indirect pointers. Disk block size is 8 KBytes and block address is 64 bit long. Assume a file F stored in file system and inode of F has 10 direct pointers and 2 single indirect pointers, but 1 double indirect pointer is NULL. And blocks pointed by 2 single indirect pointers are completely filled (no pointer is NULL). The size of file without considering inode block is \_\_\_\_\_ MBytes. (Upto 2 decimal places)

16.078 [16.07 - 16.08]

Your answer is **Correct** 16.08

Solution :

$$\begin{aligned}
 & 16.078 [16.07 - 16.08] \\
 & \text{Block size (BS)} = 8 \text{ KB} \\
 & \text{Block address (BA)} = 64 \text{ bits} = 8 \text{ B} \\
 & \text{Addresses per block} = \frac{8 \text{ KB}}{8 \text{ B}} = 1 \text{ K} = 2^{10} \\
 & \text{Size of file F} = [(\text{Number of direct pointer}) + (\text{Number of single indirect pointer}) \\
 & \quad \times (\text{Address per block})] \times \text{BS} \\
 & = (10 + 2 \times 2^{10}) \times 2^{13} \text{ B} \\
 & = \left( \frac{10}{128} + 2^4 \right) 2^{20} \text{ B} = 16.078 \text{ MB}
 \end{aligned}$$

QUESTION ANALYTICS

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Q. 15

See your Answers

Q

Consider following snapshot of a system with P<sub>0</sub>, P<sub>1</sub>, P<sub>2</sub> and P<sub>3</sub> processes and R<sub>0</sub>, R<sub>1</sub> and R<sub>2</sub> resources types. In following table, 'Allocation' column shows number of resource types allocated to processes, and 'Max' column shows maximum requirement of a resource type by a process.

Process	Allocation			Max		
	R <sub>0</sub>	R <sub>1</sub>	R <sub>2</sub>	R <sub>0</sub>	R <sub>1</sub>	R <sub>2</sub>
P <sub>0</sub>	1	0	2	3	4	5
P <sub>1</sub>	0	1	0	3	4	5
P <sub>2</sub>	0	3	1	2	5	2
P <sub>3</sub>	1	1	2	2	6	6

A Above snapshot of system is in safe state.

Your option is Correct

B Only one order of processes possible to complete the execution of all processes.

Your answer is IN-CORRECT

C At this snapshot, P<sub>1</sub> requests for (R<sub>0</sub>, R<sub>1</sub>, R<sub>2</sub>) = (0, 0, 1) instances and system allocate requested resources to P<sub>1</sub>. Now, system still in safe state.

Your answer is IN-CORRECT

D At this snapshot, P<sub>0</sub> requests for (R<sub>0</sub>, R<sub>1</sub>, R<sub>2</sub>) = (0, 0, 1) instances and system allocate requested resources to P<sub>0</sub>. Now, system still in safe state.

Correct Option

YOUR ANSWER - a,b,c

CORRECT ANSWER - a,d

STATUS - ✘

**Solution :**

(a,d)

Use Banker's algorithm:

Total allocated resources  $(R_0, R_1, R_2) = (2, 5, 5)$ Available resources at snapshot  $(R_0, R_1, R_2) = (2, 2, 2)$ 

Required/need table:

Process	Need		
	$R_0$	$R_1$	$R_2$
$P_0$	2	4	3
$P_1$	3	3	5
$P_2$	2	2	1
$P_3$	1	5	4

Possible order for complete execution

1.  $P_2, P_0, P_1, P_3$ 2.  $P_2, P_0, P_3, P_1$ (c)  $(0, 0, 1)$  allocated to  $P_1$ , now available resources  $= (2, 2, 1)$ . Only  $P_2$  can run. After  $P_2$ , available resources  $= (2, 5, 2)$ . Now, it is deadlock.(d)  $(0, 0, 1)$  allocated to  $P_0$ , now system in safe state, with two possible orders as shown above.

## QUESTION ANALYTICS

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Q. 16

Have any Doubt ?

↗

Which of the following is correct?

(A) In multi-level feedback queue scheduling processes are not allowed to move from one queue to another.

(B) Priority inversion is a condition that occurs in real-time system where lower priority process is starved because higher priority processes have gained hold of CPU.

Your answer is IN-CORRECT

(C) User threads have higher thread scheduling overhead than Kernel threads.

(D) When fork() system call creates a child process, then child process inherits all the open file descriptors of parent process.

Correct Option

**Solution :**  
(d)

## QUESTION ANALYTICS

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Q. 17

Have any Doubt ?

↗

Consider a system uses TLB, two-level page table, and stores page table in memory. TLB and memory access time are 20 nanoseconds (ns) and 100 ns respectively. TLB have a cache hit ratio of 95%. What is effective access time of memory in nanoseconds?

130

Your answer is Correct 130

**Solution :**  
130

$$H_T = \text{Hit ratio of TLB}$$

$$M_T = \text{Miss ratio of TLB}$$

$$T_T = \text{Access time of TLB}$$

$$T_M = \text{Access time of memory}$$

$$L = \text{Levels in page table}$$

Average access time of memory

$$\begin{aligned} &= H_T \times (T_T + T_M) + M_T (T_T + L \times T_M + T_M) \\ &= 0.95 \times (20 + 100) + 0.05 (20 + 2 \times 100 + 100) \\ &= 0.95 \times 120 + 0.05 (320) = 114 + 16 = 130 \text{ ms} \end{aligned}$$

## QUESTION ANALYTICS

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Q. 18

Have any Doubt ?

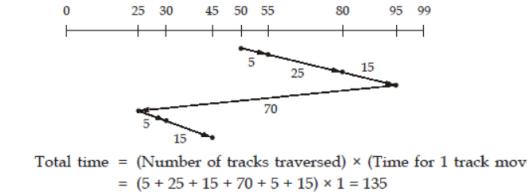
↗

Consider a disk in which disk scheduler uses C-LOOK scheduling algorithm. At current instance disk has following requests (in track numbers).

55, 45, 95, 25, 30, 80

Disk has 100 tracks (0 to 99), currently disk head positioned at track 50 and pointing toward track 99. Assume head takes 1 ms to move from a track to its adjacent track and each request service time is negligible. The time taken by disk to service all disk requests shown above is \_\_\_\_\_ ms.

**Solution :**  
135



$$\text{Total time} = (\text{Number of tracks traversed}) \times (\text{Time for 1 track move}) \\ = (5 + 25 + 15 + 70 + 5 + 15) \times 1 = 135$$

QUESTION ANALYTICS

Q. 19

[Have any Doubt ?](#)

Consider a system with 48 bits virtual address space and 4 GB byte addressable physical memory. The operating system uses hierarchical paging where inner most level page table contains frame number and page tables at level other than inner most level contain address of page tables at its next-level, towards inner most table. Page table entry size at each level is 4 Bytes and page size is 4 KB. Page table at outer most level uses only one page, the amount of this page which is not used for page table is \_\_\_\_\_ bytes.

**Solution :**  
3840

$$\text{Virtual address length} = 48 \text{ bit}$$

$$\text{Virtual memory size} = 2^{48} \text{ bytes}$$

$$\text{Page size (S)} = 4 \text{ KB}$$

$$\text{Total number of pages in virtual memory} = \frac{2^{48} \text{ B}}{4 \text{ KB}} = 2^{36} \text{ pages}$$

$$\text{Number of page table entries (PTE)} = \text{Number of virtual pages} = 2^{36}$$

$$\text{Number of pages at level-1 (inner most level)}$$

$$= \frac{(\text{Number of PTE}) \times (\text{PTE size})}{\text{Page size}} = \frac{2^{36} \times 4 \text{ B}}{4 \text{ KB}} = 2^{26}$$

$$\text{Number of pages at level-2} = \frac{(\text{Number of pages at level-1}) \times (\text{PTE size})}{\text{Page size}} = \frac{2^{26} \times 4 \text{ B}}{4 \text{ KB}} = 2^{16}$$

$$\text{Number of pages at level-3} = \frac{(\text{Number of pages at level-2}) \times (\text{PTE size})}{\text{Page size}} = \frac{2^{16} \times 4 \text{ B}}{4 \text{ KB}} = 2^6$$

$$\text{Number of pages at level-4} = \left\lceil \frac{(\text{Number of pages at level-3}) \times (\text{PTE size})}{\text{Page size}} \right\rceil = \left\lceil \frac{2^6 \times 4 \text{ B}}{4 \text{ KB}} \right\rceil = 1 \text{ page}$$

$$\text{Space used at level-4 page} = (\text{Number of pages at level-3}) \times (\text{PTE size}) \\ = 2^6 \times 4 \text{ B}$$

$$\text{Unused space in page at level-4}$$

$$= (4 \times 2^{10} - 4 \times 2^6) \text{ Bytes}$$

$$= 4 \times 2^6 (2^4 - 1) = 3840 \text{ bytes}$$

QUESTION ANALYTICS

Q. 20

[Have any Doubt ?](#)

Which of the following is false?

**A** User-level thread switching requires user-mode to Kernel-mode conversion of their process.

Correct Option

**Solution :**

(a)

User-level threads managed in user space using user-thread libraries. There is no need of Kernel in user-thread libraries.

**B** Threads of a process created within the address space of this process.

**C** In user-level threads of a process, if one thread blocked then entire process will block.

**D** Threads of a process share the protected access to other processes which are connected to process of these threads as inter-process communication.

Your answer is IN-CORRECT

QUESTION ANALYTICS