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## THEORY OF COMPUTATION-2: (GATE 2022) - REPORTS

OVERALL ANALYSIS COMPARISON REPORT **SOLUTION REPORT**

ALL(17) CORRECT(11) INCORRECT(5) SKIPPED(1)

Q. 11

FAQ

Solution Video

Have any Doubt ?



Let  $G = (V, T, P, S)$  be any context free grammar without any  $\epsilon$  productions or unit productions. Let  $m$  be the maximum number of symbols on the right of any production in  $P$ . The maximum number of production rules for any equivalent grammar in Chomsky Normal Form is given by ( $|\cdot|$  denotes the cardinality of the set).

A  $m|P| + |T|$

B  $(m-1)|P| + |T|$

Your answer is Correct

Solution :

(b)

This is the standard theorem in Peter Linz and answer is  $(m-1)|P| + |T|$ .

C  $m|P| + |T| - 1$

D  $(m-1)|P| + |T| - 1$

QUESTION ANALYTICS



Q. 12

Solution Video

Have any Doubt ?



Which of the following grammar can generate the language  $L = \{a^i b^j c^k \mid i > k, 0 \leq j < 3, k \geq 0\}$ ?

A  $X \rightarrow aXc \mid X \rightarrow aX \mid c$

B  $S \rightarrow aS \mid aSc \mid a \mid ab \mid abb \mid abbb$

C  $S \rightarrow aS \mid aSc \mid a \mid ab \mid abb$

Your answer is Correct

Solution :

(c)

$$L = \{a^i b^j c^k \mid i > k, 0 \leq j < 3, k \geq 0\}$$

The context-free grammar for the given language can be

$S \rightarrow aS$

$S \rightarrow aSc$

$S \rightarrow a$

$S \rightarrow ab$

$S \rightarrow abb$

D None of these

QUESTION ANALYTICS



Q. 13

FAQ

Solution Video

Have any Doubt ?



Consider  $L = L_1 \cap L_2$  where

$$L_1 = \{0^m 1^n 2^m \mid m, n \geq 0\}$$

$$L_2 = \{0^m 1^n 2^k \mid m, n, k \geq 0\}$$

Then, the language  $L$  is

A Recursively enumerable but not context free

Your answer is IN-CORRECT

B Regular

C Context free but not regular

Correct Option

Solution :

(c)

(u)

$L = L_1 \cap L_2$   
 $L = \{0^m 1^m \mid m \geq 0\}$  which is CFL but not regular because there is a infinite comparison.

**D** Not recursive

QUESTION ANALYTICS

Q. 14

FAQ

Solution Video

Have any Doubt ?



Consider the statements:

- I. The problem "Is  $L_1 \wedge L_2 = \phi$ " is undecidable for context sensitive languages  $L_1$  and  $L_2$ .  
 II. The problem "Is  $w \in L$ ?" is decidable for context sensitive language  $L$ . (where  $w$  is a string).

The number of correct statements is/are \_\_\_\_\_.

**2**

Your answer is Correct2

Solution :

2

Both statement are true.

- Checking emptiness or intersection between 2 language in CSL is Undecidable.
- Checking membership problem for CSL is decidable.

QUESTION ANALYTICS

Q. 15

Solution Video

Have any Doubt ?

Given language  $L = \{x^m y^m \mid m > 0\}$ . Now consider the given statements:

- I.  $E \rightarrow xEy \mid xy$   
 II.  $xy \mid x^+ xy y^+$   
 III.  $x^+ y^+$

The number of the correct statements which generates  $L$  is/are \_\_\_\_\_.

**1**

Correct Option

Solution :

1

Given language  $L$  is CFL, hence it can not be generated either through regular grammar and or regular expression. Statement I is CFG and generates  $L$ .

**1**

Your Answer is 2

QUESTION ANALYTICS

Q. 16

FAQ

Solution Video

Have any Doubt ?



Which of the following are true?

**A** The complement of an RE language can be not-RE.

Your option is Correct

**B** The complement of an REC language can be RE.

Correct Option

**C** The complement of a non-DCFL can be a DCFL.

**D** The complement of a language accepted by a NFA can be acceptable by a DFA.

Your option is Correct

YOUR ANSWER - a,d

CORRECT ANSWER - a,b,d

STATUS - ✖

Solution :

(a, b, d)

QUESTION ANALYTICS

Q. 17

FAQ

Solution Video

Have any Doubt ?



Consider the following languages:

$$L_1 = \{a^i b^j c^k\} \cup \{a^i b^j c^k\} \mid i, j \geq 0.$$

$$L_2 = \{w w^R \mid w \in \{a, b\}^*\} \text{ where } R \text{ represents reversible operation.}$$

$$L_3 = \{a^m b^n c^k \mid m, n, k \geq 0\}$$

Which one of the following is(are) inherently ambiguous language(s)?

**A**  $L_1$

Your option is Correct

**B**  $L_2$

**C**  $L_3$

**D** None of these

YOUR ANSWER - a

CORRECT ANSWER - a

STATUS - ✓

**Solution :**

(a)

- If there is no context free grammar for  $L$  which is unambiguous. Hence  $L$  is inherently ambiguous or in other words "A language for which every grammar is ambiguous is called inherently ambiguous language".

Note: If a language is regular or DCFL it will surely can be written in unambiguous grammar.



- So, surely  $L_3$  can't be inherently ambiguous language as it is regular language.
- $L_2$  is not inherently ambiguous language as we can write grammar for  $L_2$  which is unambiguous.  
 $S \rightarrow aSa \mid bSb \mid \epsilon$
- $L_1$  is CFL and it is inherently ambiguous language because  $L_1$  is language with union of two DCFL. So the grammar will always have OR operation which will make them ambiguous. Let's see how we can write grammar.

$$\begin{aligned} S &\rightarrow X \mid Y \\ X &\rightarrow aXc \mid P \\ P &\rightarrow bP \mid \epsilon \\ Y &\rightarrow aYb \mid Q \\ Q &\rightarrow cQ \mid \epsilon \end{aligned}$$

If we generate string 'abc' two parse tree possible either through  $X$  or  $Y$ .  
Hence it is ambiguous and leads to inherently ambiguous language.

QUESTION ANALYTICS

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