



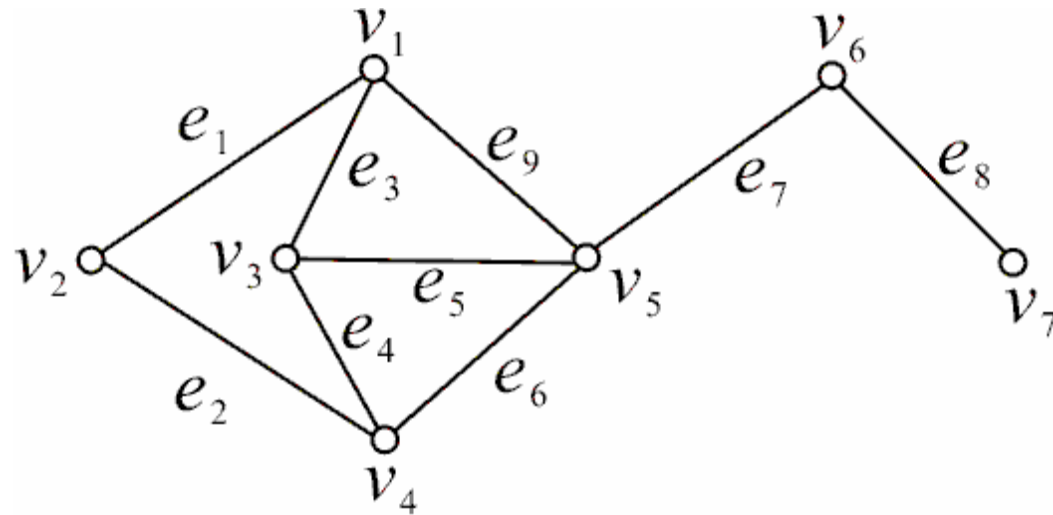
Paths

- **Definition**
- Let n be a nonnegative integer and G an undirected graph. A path of length n from u to v in G is a sequence of n edges e_1, \dots, e_n of G for which there exists a sequence $x_0 = u, x_1, \dots, x_n = v$ of vertices such that e_i has the endpoints x_{i-1} and x_i , when the graph is simple, we denote this path by its vertex sequence.
- **circuit** ---- begins and ends at the same vertex
- **Simple path** ---- does not contain the same edge more than once.
- **length of a path** ---- the number of edges on the path



Example

Discrete
Mathematics



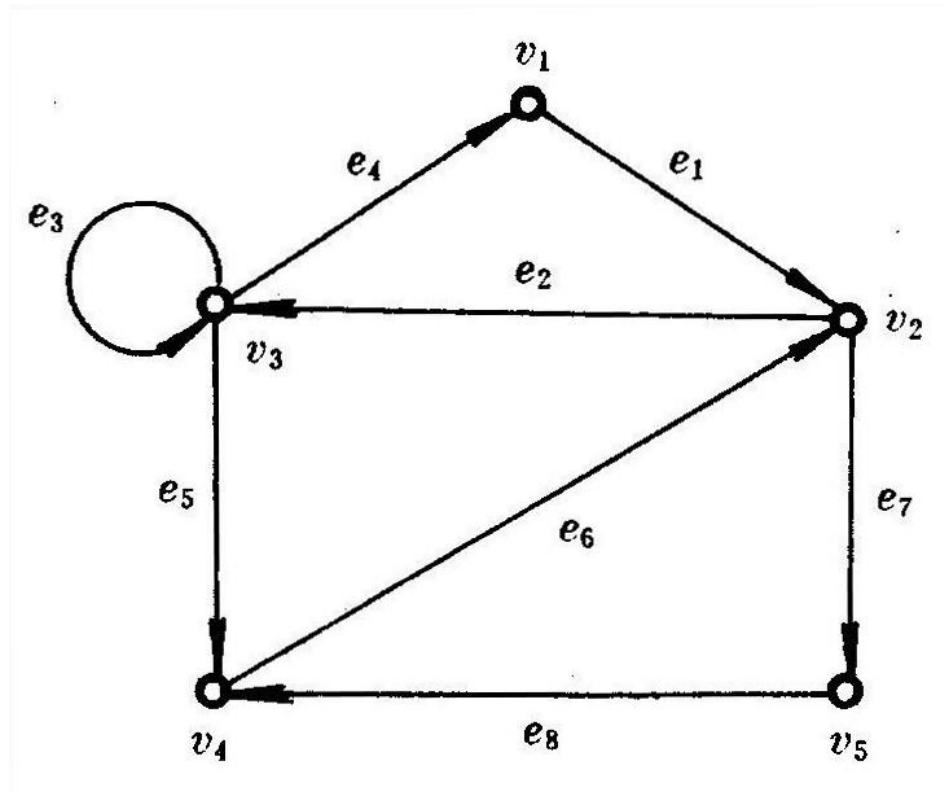


- **Definition**
- Let n be a nonnegative integer and G a directed graph. A path from a to b in G is a sequence of edges $(x_0, x_1), (x_1, x_2), \dots, (x_{n-1}, x_n)$, where $x_0 = a, x_n = b$. that is, a sequence of edges where the terminal vertex of an edge is the same as the initial vertex in the next edge in the path.
- Using $e_1, e_2 \dots e_n$ to denote path.
- A path can pass through a vertex more than once, can traverse an edge more than once.



Example

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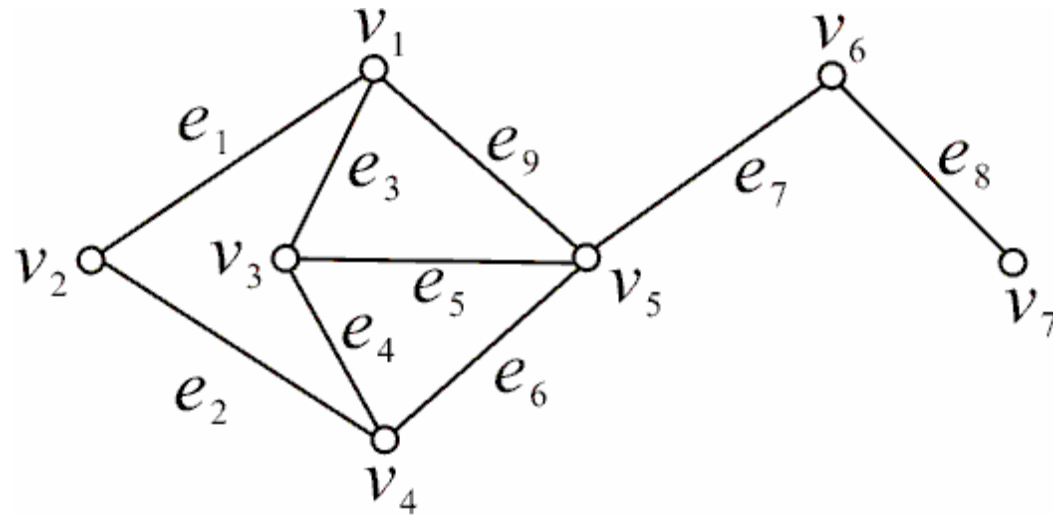


Connectedness

- **Definition:**
- An undirected graph is called connected if there is a path between every pair of distinct vertices of the graph. An undirected graph that is not connected is called disconnected. We say that we disconnect a graph when we remove vertices or edges, to produce a disconnected subgraph.



Example



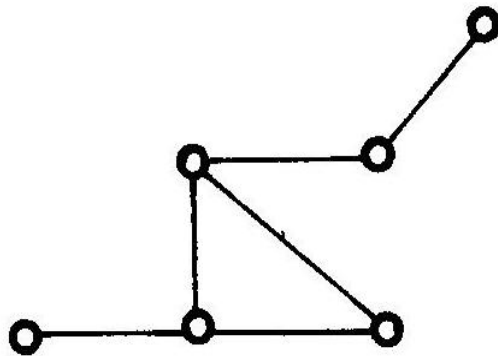


- Theorem
- There is a simple path between every pair of distinct vertices of a connected undirected graph.

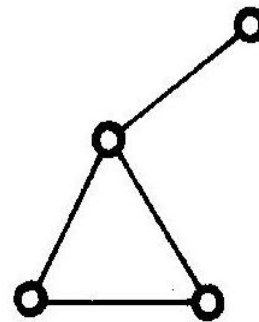


Connected components

- A **connected component** of a graph G is a connected graph of G that is not a proper subgraph of another subgraph of G . That is, a connected component of a graph G is a maximal connected graph of G .
- A graph that is not connected has two or more connected component that are disjoint and have G as their union.



(a)

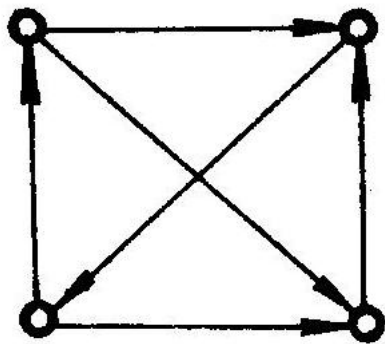


(b)

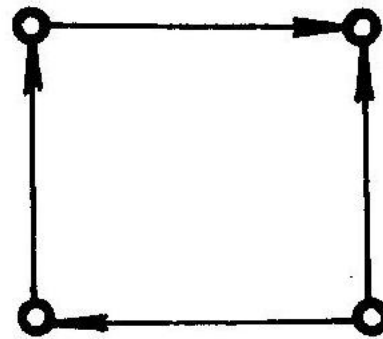


Connectedness

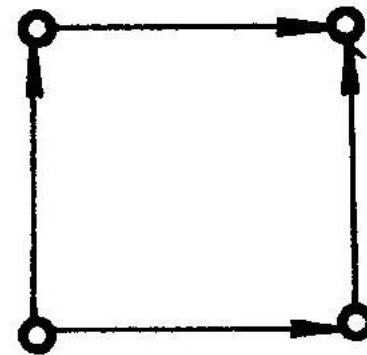
- **Definition:**
- A directed graph is strongly connected if there is a path from a to b and from b to a whenever a and b are vertices in the graph.
- A directed graph is weakly connected if there is a path between every two vertices in the underlying undirected graph.
- **G is strongly connected if and only if there is a circuit passing through every vertex.**



(a)



(b)



(c)



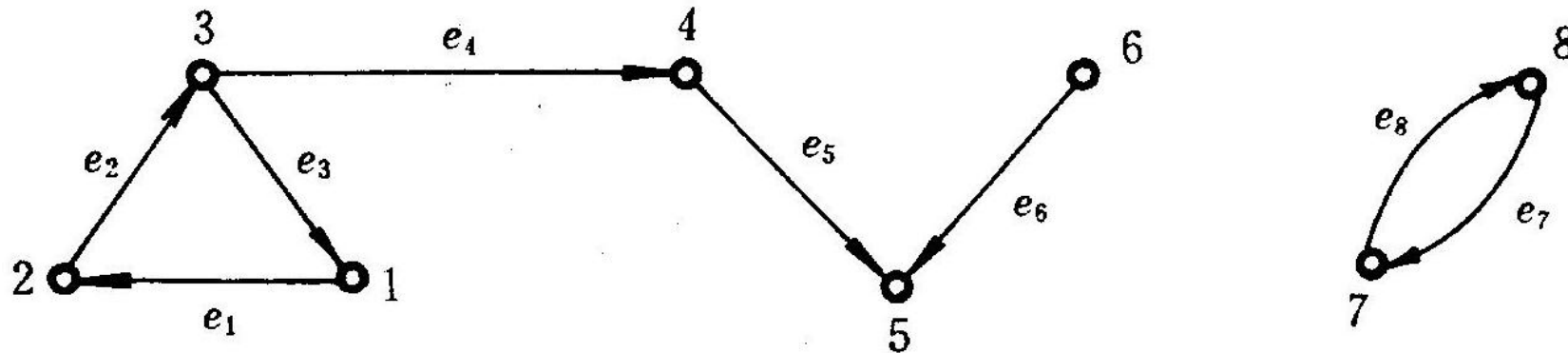
Strong components

- The subgraph of a directed graph G that are strongly connected but not contained in larger strongly connected subgraphs, that is, the maximal strongly connected subgraphs, are called the strongly connected components or strong components of G .



Strong components

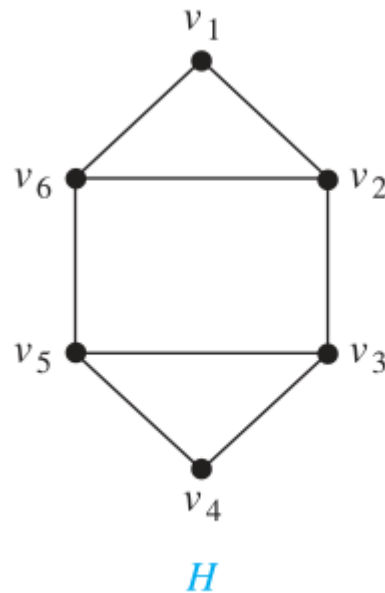
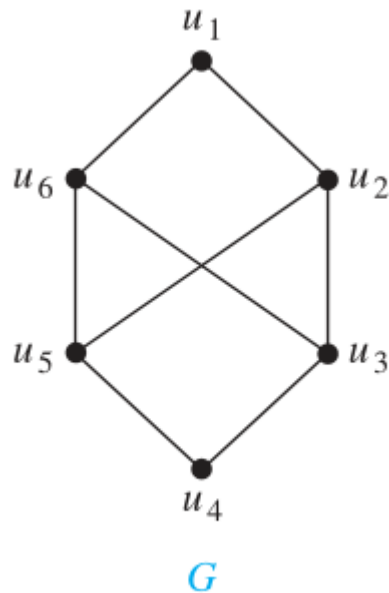
- $\{ \langle \{1,2,3\}, \{e_1, e_2, e_3\} \rangle , \langle \{4\}, \varphi \rangle , \langle \{5\}, \varphi \rangle , \langle \{6\}, \varphi \rangle , \langle \{7,8\}, \{e_7, e_8\} \rangle \}$





Path and Isomorphism

- Path can be used to determine whether two graphs are isomorphic or not
- the existence of a simple circuit of length k



H has a simple circuit of length three, namely, v_1, v_2, v_6, v_1 , whereas G has no simple circuit of length three



Counting paths

Theorem Let G be a graph with adjacency matrix A with respect to the ordering v_1, v_2, \dots, v_n of the vertices of the graph (with directed or undirected edges, with multiple edges and loops allowed). The number of different paths of length r from v_i to v_j , where r is a positive integer, equals the (i, j) th entry of A^r .

$$b_{i1}a_{1j} + b_{i2}a_{2j} + \dots + b_{in}a_{nj},$$



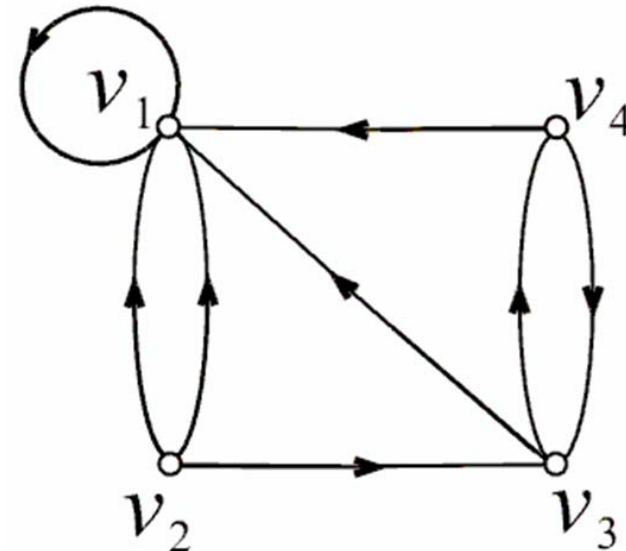
Example

$$A = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 2 & 0 & 1 & 0 \\ 1 & 0 & 0 & 1 \\ 1 & 0 & 1 & 0 \end{bmatrix}$$

$$A^3 = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 4 & 0 & 1 & 0 \\ 3 & 0 & 0 & 1 \\ 3 & 0 & 1 & 0 \end{bmatrix}$$

$$A^2 = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 3 & 0 & 0 & 1 \\ 2 & 0 & 1 & 0 \\ 2 & 0 & 0 & 1 \end{bmatrix}$$

$$A^4 = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 5 & 0 & 0 & 1 \\ 4 & 0 & 1 & 0 \\ 4 & 0 & 0 & 1 \end{bmatrix}$$



The number of different paths of length 1 is 8, the circuit is 1.

The number of different paths of length 2 is 11, the circuit is 3.

The number of different paths of length 3 is 14, the circuit is 1.

The number of different paths of length 4 is 17, the circuit is 3.



Shortest-path problems

Discrete
Mathematics

- In many applications, each edge of a graph has an associated numerical value, called a weight. The weight of an edge is often referred to as the "cost" of the edge.
- In applications, the weight may be a measure of distance, time, cost etc.
- Usually, the edge weights are nonnegative integers.



Shortest Paths

- shortest-path: Given two vertices A and B, there are more than one paths from A to B. The path with minimum cost is called shortest path.
- There are several different algorithms to find a shortest path.



Dijkstra's Algorithm

- We assume that there is a path from the source vertex v_0 to every other vertex in the graph.
- Let S be the set of vertices whose minimum distance from the source vertex has been found. Initially S contains only the source vertex.
- The algorithm is iterative, adding one vertex to S at each step.
- We maintain an array D such that for each vertex v , $D[v]$ is the minimum distance from the source vertex to v via vertices that are already in S .
- Every subpath is the shortest path in this whole shortest path.



- STEPS:

1. let $S=\{v_0\}$, compute $D[i]$ for each vertex v_i as following:

$$D[i] = \begin{cases} 0 & \text{if } i=0 \\ w_{0i} & \text{if } i \neq 0, \text{ and } \langle v_0, v_i \rangle \text{ is an edge, } w_{si} \text{ is the weight} \\ \infty & \text{if } i \neq 0, \text{ and } \langle v_0, v_i \rangle \text{ is not an edge} \end{cases}$$

2. choose the vertex v_j such that:

$$D[j] = \min \{ D[k] \mid v_k \notin S \}$$

then the v_j is starting of the next shortest path and $D[j]$ is its cost.



STEPS:

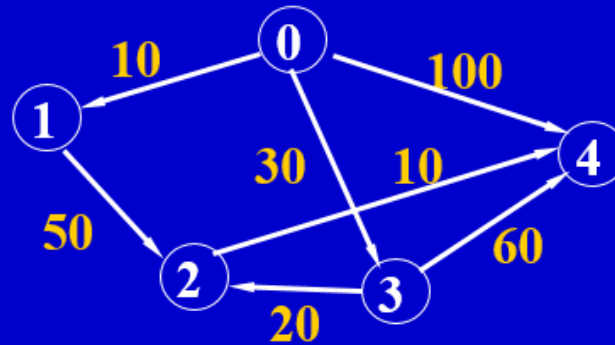
3. Place v_j in S . That is

$$S = S \cup \{v_j\}$$

4. For each $v_k \notin S$, modify the $D[k]$:

$$D[k] = \min \{ D[k], D[j] + \text{weight}(\langle v_j, v_k \rangle) \}$$

5. Repeat 2---4 until all vertices have been added in S .



Steps	S	D[0]	D[1]	D[2]	D[3]	D[4]
begin	{ 0 }	0	10	∞	30	100
1	{ 0, 1 }	0	10	60	30	100
2	{ 0, 1, 3 }	0	10	50	30	90
3	{ 0, 1, 3, 2 }	0	10	50	30	60
4	{ 0, 1, 3, 2, 4 }	0	10	50	30	60



homework

Discrete
Mathematics

- 10.4 P724
3, 4, 5, 11, 14(a)(b), 19(b)(c), 20,
- 10.6 p751
2, 17