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Homework -7

Operating System

1 Consider the following snapshot of a system:

	<u>Allocation</u>	<u>Max</u>	<u>Available</u>
	<u>A B C D</u>	<u>A B C D</u>	<u>A B C D</u>
P_0	0 0 1 2	0 0 1 2	1 5 2 0
P_1	1 0 0 0	1 7 5 0	
P_2	1 3 5 4	2 3 5 6	
P_3	0 6 3 2	0 6 5 2	
P_4	0 0 1 4	0 6 5 6	

Answer the following questions using the banker's algorithm:

- What is the content of the matrix Need?
- Is the system in a safe state?
- If a request from process P_1 arrives for (0,4,2,0), can the request be granted immediately?

Answer No:1

a. Need = Max.- Allocation

①

	Allocation	Max	Available	Need	
	A B C D	A B C D	A B C D	A B C D	
P_0	0 0 1 2	0 0 1 2	1 5 2 0	0 0 0 0	①
P_1	1 0 0 0	1 7 5 0	1 5 3 2	0 7 5 0	⑤
P_2	1 3 5 4	2 3 5 6	2 8 8 6	1 0 0 2	②
P_3	0 6 3 2	0 6 5 2	2 1 4 11 8	0 0 2 0	③
P_4	0 0 1 4	0 6 5 6	2 1 4 12 12 3 1 4 12 12	0 6 4 2	④
Sum =	2 9 10 12				

Total
A B C D
3 14 12 12

Safe Sequence: $P_0 P_2 P_3 P_4 P_1$

Yes, the system is in a safe state. Because, Total and final result of available is same. We can say that, system is in a safe state.

The values of Need for processes P_0 through P_4 respectively are (0, 0, 0, 0), (0, 7, 5, 0), (1, 0, 0, 2), (0, 0, 2, 0), and (0, 6, 4, 2).

b. The system is in a safe state, Yes. With **Available** being equal to (1, 5, 2, 0), either process P_0 or P_3 could run. Once process P_3 runs, it releases its resources, which allow all other existing processes to run.

①

	Allocation	Max	Available	Need	
	ABCD	ABCD	ABCD	ABCD	
✓ P_0	0012	0012	1520	0000	①
✓ P_1	1000	1750	1532	0750	⑤
✓ P_2	1354	2356	2886	1002	②
✓ P_3	0632	0652	214118	0020	③
✓ P_4	0014	0656	2141212	0642	④
Sum =	291012		3141212		

Ans No: 1(a)

Total
ABCD
3 14 12 12

Safe Sequence: $P_0 P_2 P_3 P_4 P_1$

Yes, the system is in a safe state. Because, Total and final result of available is same. We can say that, system is in a safe state.

c.

②

$$P_1 \rightarrow (0, 4, 2, 0)$$

$$\text{Need}(P_1) \rightarrow \text{Max-Allocation}$$

$$= 1750 - 1000$$

$$= 0750$$

$$\textcircled{1} \text{ Request} \leq \text{need} \quad \checkmark$$

$$0420 \leq 0750$$

$$\textcircled{2} \text{ Request} \leq \text{available} \quad \checkmark$$

$$0420 \leq 1520$$

③

$$* \text{available} = \text{available} - \text{request}$$

$$= 1520 - 0420$$

$$= 1100$$

$$* \text{allocation} = \text{allocation} + \text{request}$$

$$= 1000 + 0420$$

$$= 1420$$

$$* \text{Need} = \text{Need} - \text{Request}$$

$$= 0750 - 0420$$

$$= 0330$$

③						①
	Allocation	Max	Available	Need		
	ABCD	ABCD	ABCD	ABCD		
✓ P ₀	0012	0012	1100	0000	①	
P ₁	1420	1750	1112 (P ₀ +A)	0330	⑤	
✓ P ₂	1354	2356	2466 (P ₂ +A)	1002	②	
✓ P ₃	0632	0652	21098 (P ₃ +A)	0020	③	
✓ P ₄	0014	0656	2101012 (P ₄ +A)	0642	④	
	2131212		3141212			
						Total
						ABCD
						3141212

Safe Sequence: P₀ P₂ P₃ P₄ P₁

The request can be granted immediately. This results in the value of **Available** being (1, 1, 0, 0). One ordering of processes that can finish is P₀, P₂, P₃, P₁, and P₄.

2 Consider the following snapshot of a system:

	<u>Allocation</u>	<u>Max</u>	<u>Available</u>
	<u>A B C D</u>	<u>A B C D</u>	<u>A B C D</u>
P_0	2 0 0 1	4 2 1 2	3 3 2 1
P_1	3 1 2 1	5 2 5 2	
P_2	2 1 0 3	2 3 1 6	
P_3	1 3 1 2	1 4 2 4	
P_4	1 4 3 2	3 6 6 5	

Answer the following questions using the banker's algorithm:

- Illustrate that the system is in a safe state by demonstrating an order in which the processes may complete.
- If a request from process P_1 arrives for (1, 1, 0, 0), can the request be granted immediately?
- If a request from process P_4 arrives for (0, 0, 2, 0), can the request be granted immediately?

Answer No:2

(a)

$Max = Allocation + Need$

	Allocation	Max	Available	Need
	A B C D	A B C D	A B C D	A B C D
✓ P ₀	2 0 0 1	4 2 1 2	3 3 2 1	2 2 1 1
✓ P ₁	3 1 2 1	5 2 5 2	5 3 2 2	2 1 3 1
P ₂	2 1 0 3	2 3 1 6	6 6 3 4	0 2 1 3
✓ P ₃	1 3 1 2	1 4 2 4	7 10 6 6	0 1 1 2
✓ P ₄	1 4 3 2	3 6 6 5	10 11 8 7	2 2 3 3
Sum =	9 9 6 9		12 12 8 10	

Safe Sequence: P₀ P₃ P₄ P₁ P₂

⇒ Allocation + Available

$Need = Max - Allocation$

$Total = Sum of Allocation + Available$

Total
A B C D
12 12 8 10

Yes, System is in Safe State

(b)

② P₀ ✓

P₁ → (1, 1, 0, 0)

Need (P₁) → Max - Allocation

= 5 2 5 2 - 3 1 2 1

= 2 1 3 1

① Request ≤ need
1 1 0 0 ≤ 2 1 3 1 ✓

② Request ≤ available
1 1 0 0 ≤ 3 3 2 1 ✓

③ * available = available - request
= 3 3 2 1 - 1 1 0 0
= 2 2 2 1

* allocation = allocation + request
= 3 1 2 1 + 1 1 0 0
= 4 2 2 1

* Need = Need - Request
= 2 1 3 1 - 1 1 0 0
= 1 0 3 1

③

	Allocation	Max	Available	Need
	A B C D	A B C D	A B C D	A B C D
✓ P ₀	2 0 0 1	4 2 1 2	2 2 2 1	2 2 1 1 ①
✓ P ₁	4 2 2 1	5 2 5 2	4 2 2 2	1 0 3 1 ④
✓ P ₂	2 1 0 3	2 3 1 6	5 5 3 4	0 2 1 3 ⑤
✓ P ₃	1 3 1 2	1 4 2 4	6 9 6 6	0 1 1 2 ②
✓ P ₄	1 4 3 2	3 6 6 5	10 11 8 7	2 2 3 3 ③
	<u>10 10 6 9</u>		12 12 8 10	

Safe Sequence : P₀ P₃ P₄ P₁ P₂

Total
 A B C D
12 12 8 10

The request can be granted immediately. This results in the value of **Available** being (2, 2, 2, 1). One ordering of processes that can finish is P₀, P₃, P₄, P₁, and P₂.

(c)

④

$$P_4 \rightarrow (0, 0, 2, 0)$$

2

$$\begin{aligned} \text{Need}(P_4) &\rightarrow \text{Max. Allocation} \\ &= 3665 - 1432 \\ &= 2233 \end{aligned}$$

① Request \leq need ✓

$$0020 \leq 2233$$

② Request \leq available ✓

$$0020 \leq 3321$$

③ * available = available - request

$$= 3321 - 0020$$

$$\begin{aligned} &= 3301 \\ * \text{allocation} &= \text{allocation} + \text{request} \\ &= 1432 + 0020 \end{aligned}$$

$$\begin{aligned} &= 1452 \\ * \text{Need} &= \text{Need} - \text{Request} \\ &= 2233 - 0020 \\ &= 2213 \end{aligned}$$

	Allocation	Max	Available	Need
	A B C D	A B C D	A B C D	
P ₀	2 0 0 1	4 2 1 2	2 2 1 3	2 2 1 1 ①
P ₁	3 1 2 1	5 2 5 2	4 2 1 4	2 1 3 1 ⑤
P ₂	2 1 0 3	2 3 1 6	6 3 1 7	0 2 1 3 ②
P ₃	1 3 1 2	1 4 2 4	7 6 2 9	0 1 1 2 ③
P ₄	1 4 5 2	3 6 6 5	8 10 7 11	2 2 1 3 ④
	9 9 8 9		11 11 9 12	
				Total
				A B C D
				11 11 9 12

P₀ P₂ P₃ P₄ P₁

The request can be granted immediately. This results in the value of **Available** being (2, 2, 1, 3). One ordering of processes that can finish is P₀, P₂, P₃, P₄, and P₁.