



西北工业大学
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Lab report

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Experiment 7

Experiment No:7

Transaction and Concurrency Control

Goal:

1. To grasp the concept transaction, and how to create the transaction.
2. Understand the data inconsistency problems in concurrency operation, and can use lock and ioslation mechanisms.

Content:

According to the requirements, try to complete the following experiments based on the database above:

1. Write a transaction to achieve the following operations: a student(student number is 20200032) transfers 200 yuan from the bank card to the campus card, and if there is a failure during the transfer process, it will be rolled back. (10 points)

The screenshot shows a MySQL query editor window titled "Query 1". The SQL code is as follows:

```
1 • create table icbc_card(  
2   icbcid int,  
3   sno varchar(8),  
4   balance int  
5 )ENGINE = INNODB;  
6 • create table campus_card(  
7   sno varchar(8),  
8   balance int  
9 )ENGINE = INNODB;  
10 • insert into campus_card values ('20200032', 1);  
11 • insert into campus_card values ('20200033', 100);  
12 • insert into icbc_card values (1, '20200032', 300);  
13 • insert into icbc_card values (2, '20200033', 400);
```

Below the query editor is the "Output" section, which displays the "Action Output" table. The table has columns for #, Time, Action, Message, and Duration / Fetch. The output shows the execution of the SQL statements and the number of rows affected.

#	Time	Action	Message	Duration / Fetch
✓ 1	17:07:45	create table icbc_card(icbcid int, sno varchar(8), b...	0 row(s) affected	0.031 sec
✓ 2	17:07:45	create table campus_card(sno varchar(8), balance ...	0 row(s) affected	0.032 sec
✓ 3	17:07:45	insert into campus_card values ('20200032', 1)	1 row(s) affected	0.000 sec
✓ 4	17:07:45	insert into campus_card values ('20200033', 100)	1 row(s) affected	0.000 sec
✓ 5	17:07:45	insert into icbc_card values (1, '20200032', 300)	1 row(s) affected	0.015 sec
✓ 6	17:08:01	insert into icbc_card values (2, '20200033', 400)	1 row(s) affected	0.016 sec

1 • `SELECT * FROM students.campus_card;`

Result Grid | Filter Rows: | Export: | Wrap Cell Content: [IA](#)

	sno	balance
▶	20200032	1
	20200033	100

Output

Action Output

#	Time	Action	Message
✓ 1	17:09:55	SELECT * FROM students.campus_card LIMIT 0, ...	2 row(s) returned

Fig: Student Campus Card

Query 1 | icbc_card x

Limit to 1000 rows

1 • `SELECT * FROM students.icbc_card;`

Result Grid | Filter Rows: | Export: | Wrap Cell Content: [IA](#)

	icbcid	sno	balance
▶	1	20200032	300
	2	20200033	400

Output

Action Output

#	Time	Action	Message	Duration / Fetch
✓ 1	17:09:55	SELECT * FROM students.campus_card LIMIT 0, ...	2 row(s) returned	0.000 sec / 0.000 sec
✓ 2	17:10:36	SELECT * FROM students.icbc_card LIMIT 0, 1000	2 row(s) returned	0.015 sec / 0.000 sec

Fig: Student Icbc Card

Query 1 icbc_card campus_card x

Limit to 1000 rows

1 • `SELECT * FROM students.campus_card;`

Result Grid

sno	balance
20200032	201
20200033	100

campus_card 1 x

Output

Action Output

#	Time	Action	Message	Duration / Fetch
1	17:16:55	<code>SELECT * FROM students.icbc_card LIMIT 0, 1000</code>	2 row(s) returned	0.000 sec / 0.000 sec
2	17:19:41	<code>SELECT * FROM students.campus_card LIMIT 0, ...</code>	2 row(s) returned	0.000 sec / 0.000 sec

Fig: After transaction the updated value in campus card

Query 1 x

Limit to 1000 rows

1 • `use students;`
 2 • `SET autocommit = 1;`
 3 • `START TRANSACTION;`
 4 • `UPDATE icbc_card`
 5 • `SET balance = balance - 200`
 6 • `WHERE sno = '20200032';`
 7 •
 8 • `update campus_card`
 9 • `set balance = balance + 200`
 10 • `WHERE sno = '20200032';`
 11 •
 12 • `ROLLBACK;`

Output

Action Output

#	Time	Action	Message	Duration / Fetch
1	17:23:58	<code>ROLLBACK</code>	0 row(s) affected	0.000 sec

Fig: Doing Rollback

Limit to 1000 rows

```

2 • SET autocommit = 1;
3 • START TRANSACTION;
4 • UPDATE icbc_card
5     SET balance = balance - 200
6     WHERE sno = '20200032';
7
8 • update campus_card
9     set balance = balance + 200
10    WHERE sno = '20200032';
11
12 • ROLLBACK;
13 • SELECT * FROM students.campus_card;

```

Result Grid

sno	balance
20200032	1
20200033	100

campus_card 1 x

Output

Action Output

#	Time	Action	Message
1	17:26:37	SELECT * FROM students.campus_card LIMIT 0, 1000	2 row(s) returned

Fig: Value rolled back to the previous value

```

2 • SET autocommit = 1;
3 • START TRANSACTION;
4 • UPDATE icbc_card
5     SET balance = balance - 200
6     WHERE sno = '20200032';
7
8 • update campus_card
9     set balance = balance + 200
10    WHERE sno = '20200032';
11
12 • ROLLBACK;
13 • SELECT * FROM students.icbc_card;

```

Result Grid

icbcid	sno	balance
1	20200032	300
2	20200033	400

icbc_card 2 x

Output

Action Output

#	Time	Action	Message
1	17:26:37	SELECT * FROM students.campus_card LIMIT 0, 1000	2 row(s) returned
2	17:29:29	SELECT * FROM students.icbc_card LIMIT 0, 1000	2 row(s) returned

Fig: Value rolled back to the previous value

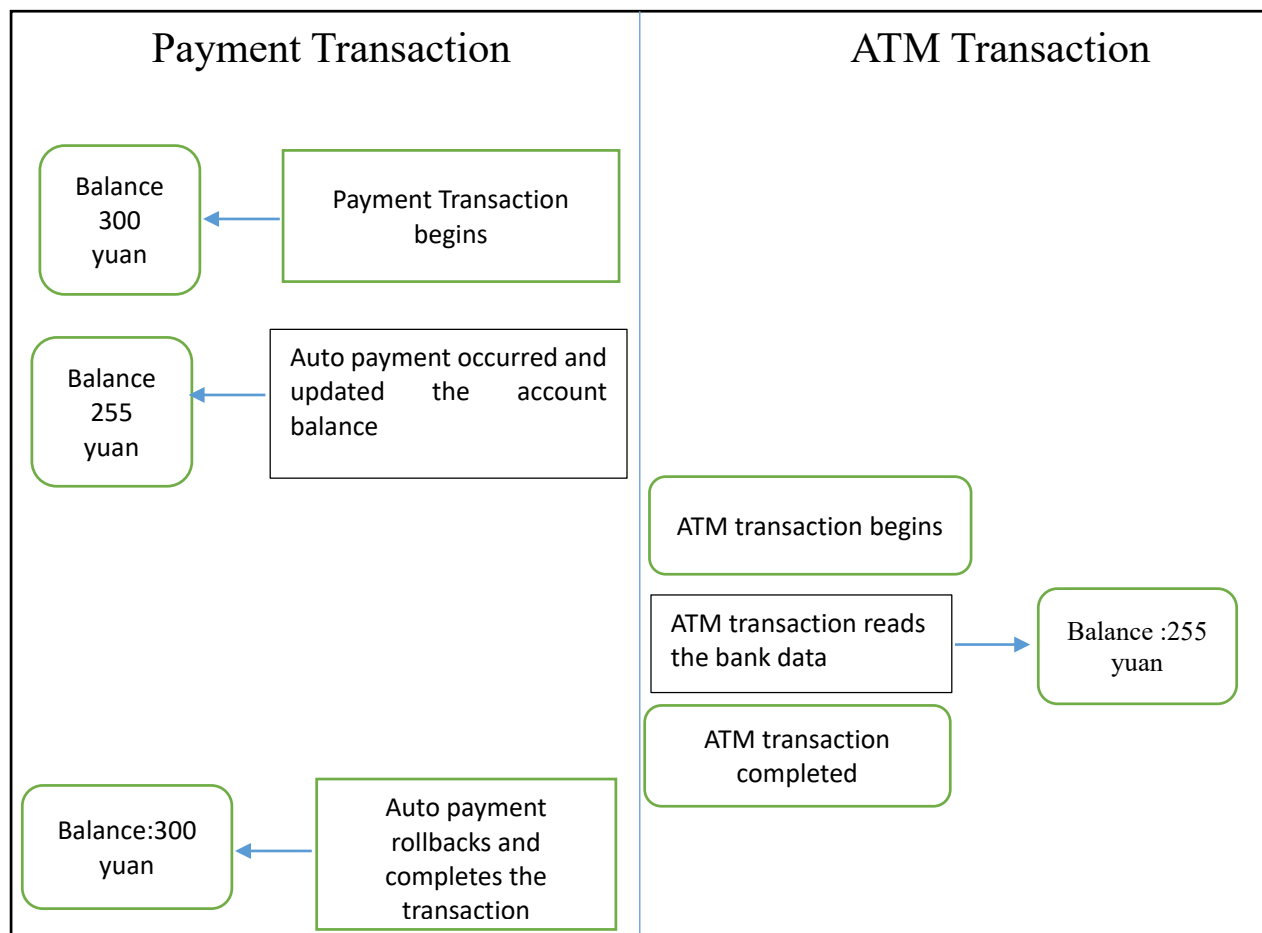
2. According to the database and tables, use specific examples to show several data inconsistency problems: such as missing and modifying, reading dirty data, non repeatable reading and phantom reading (deletion and insertion). If there is any situation that cannot be displayed, please explain the reasons. (20 points, 10 points for each data inconsistency)

Solution:

Dirty read

The simplest explanation of the dirty read is the state of reading uncommitted data. In this circumstance, we are not sure about the consistency of the data that is read because we don't know the result of the open transaction(s). After reading the uncommitted data, the open transaction can be completed with rollback. On the other hand, the open transaction can complete its actions successfully. The data that is read in this ambiguous way is defined as dirty data.

In this scenario, **student number = 20200032** has 300 yuan in the bank account, and the automatic payment system withdraws 45 yuan from Betty's account for the electric bill for dorm. At that time, **20200032** wants to check the bank account on the ATM, and she notices 255 yuan in her bank account. However, if the electric bill payment transaction is rollbacked for any reason, the bank account balance will be turned to 300 yuan again, so the data read by **20200032** is dirty data. In this case, **20200032** will be confused. The following diagram illustrates this dirty read scenario in a clearer manner.



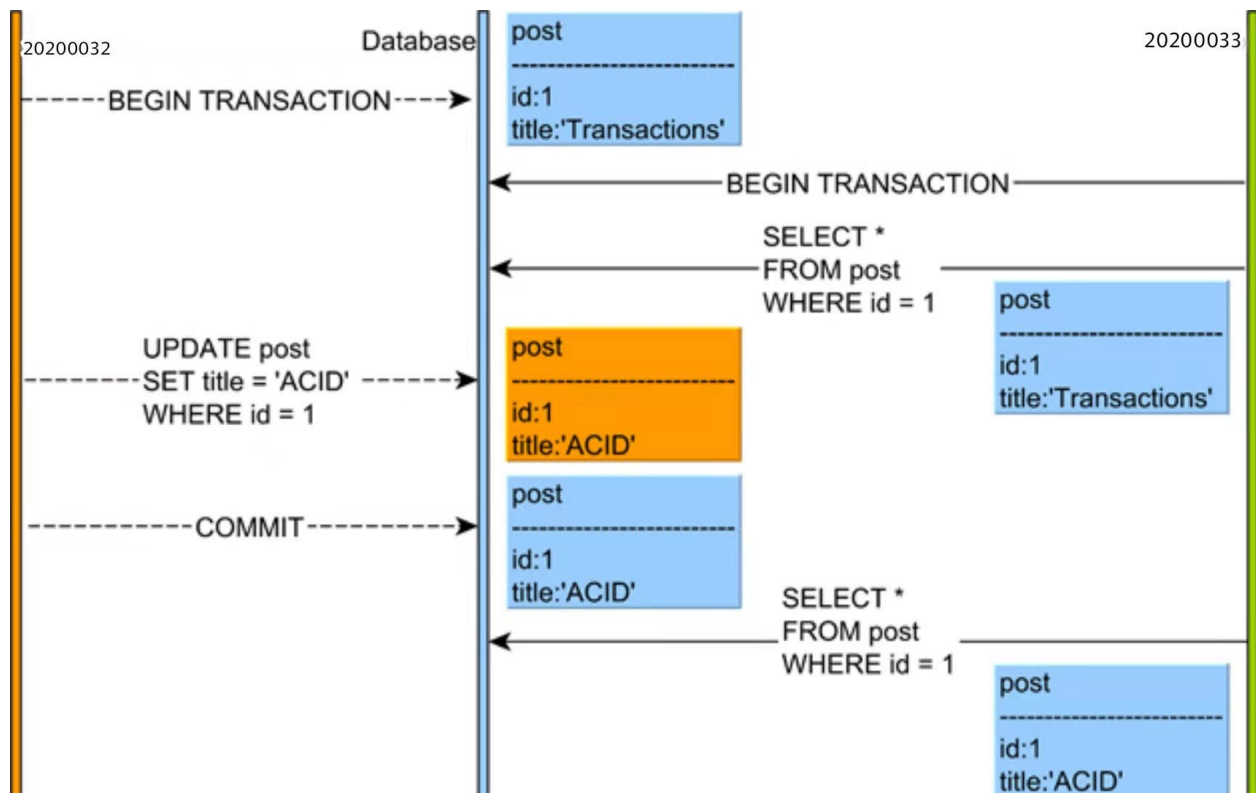
We can see that, the values are inconsistent. Therefore, we can see inconsistency.

Non-repeatable Reads:

A non-repeatable read occurs when a transaction reads the same row twice but gets different data each time. For example, suppose transaction 1 reads a row. Transaction 2 updates or deletes that row and commits the update or delete. If transaction 1 rereads the row, it retrieves different row values or discovers that the row has been deleted.

The **Non-Repeatable Read** anomaly looks as follows:

1. 20200032 and 20200033 start two database transactions.
2. 20200033's reads the post record and title column value is Transactions.
3. 20200032 modifies the title of a given post record to the value of ACID.
4. 20200032 commits her database transaction.
5. If 20200033's re-reads the post record, he will observe a different version of this table row

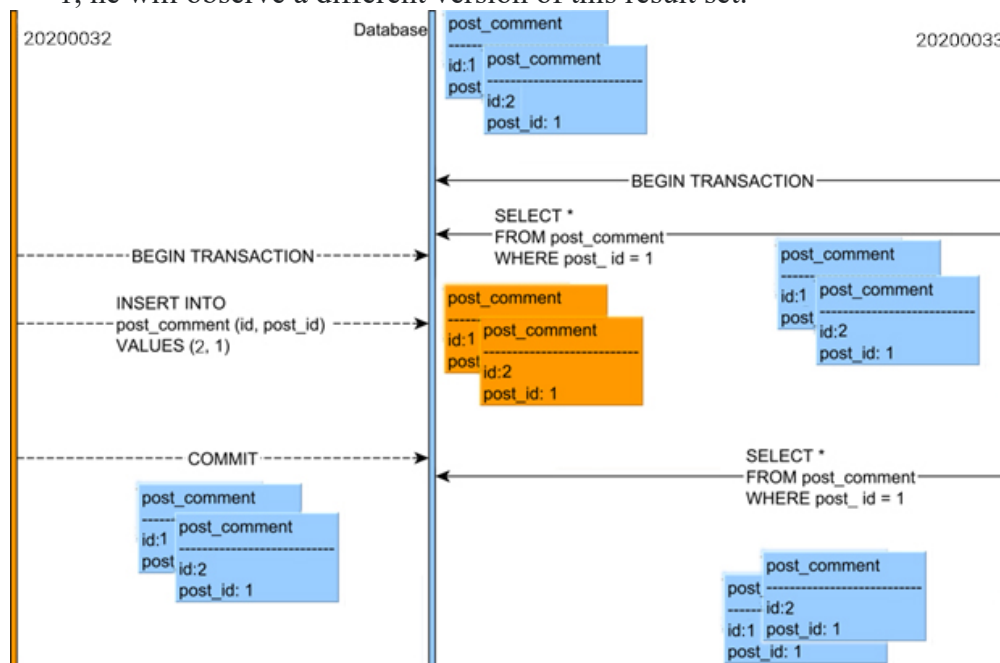


1. 20200032 and 20200033 start two database transactions.
2. 20200033's reads all the post_comment records associated with the post row with the identifier value of 1.
3. 20200032 adds a new post_comment record which is associated with the post row having the identifier value of 1.
4. 20200032 commits her database transaction.

5. If 20200033's re-reads the post_comment records having the post_id column value equal to 1, he will observe a different version of this result set.

The **Phantom Read** anomaly can happen as follows:

1. 20200032 and 20200033 start two database transactions.
2. 20200033's reads all the post_comment records associated with the post row with the identifier value of 1.
3. 20200032 adds a new post_comment record which is associated with the post row having the identifier value of 1.
4. 20200032 commits her database transaction.
5. If 20200033's re-reads the post_comment records having the post_id column value equal to 1, he will observe a different version of this result set.



So, while the **Non-Repeatable Read** applies to a single row, the **Phantom Read** is about a range of records which satisfy a given query filtering criteria.

3. By using the isolation levels or lock mechanism of the database, design solutions to solve the data inconsistency problems you have set in question 2. (20 points, 5 points for each data inconsistency)

Solution:

The SQL standard, which has been adopted by both ANSI and ISO/IEC, defines four levels of **transaction isolation**. These levels have differing degrees of impact on transaction processing throughput.

These isolation levels are defined in terms of phenomena that must be prevented between concurrently executing transactions. The preventable phenomena are:

- Dirty reads

A transaction reads data that has been written by another transaction that has not been committed yet.

- Nonrepeatable (fuzzy) reads

A transaction rereads data it has previously read and finds that another committed transaction has modified or deleted the data. For example, a user queries a row and then later queries the same row, only to discover that the data has changed.

- Phantom reads

A transaction reruns a query returning a set of rows that satisfies a search condition and finds that another committed transaction has inserted additional rows that satisfy the condition.

For example, a transaction queries the number of employees. Five minutes later it performs the same query, but now the number has increased by one because another user inserted a record for a new hire. More data satisfies the query criteria than before, but unlike in a fuzzy read the previously read data is unchanged.

Preventable Read Phenomena by Isolation Level

Isolation Level	Dirty Read	Nonrepeatable Read	Phantom Read
Read uncommitted	Possible	Possible	Possible
Read committed	Not possible	Possible	Possible
Repeatable read	Not possible	Not possible	Possible
Serializable	Not possible	Not possible	Not possible

Conflicting Writes in Read Committed Transactions

In a read committed transaction, a **conflicting write** occurs when the transaction attempts to change a row updated by an uncommitted concurrent transaction, sometimes called a **blocking transaction**. The read committed transaction waits for the blocking transaction to end and release its row lock. The options are as follows:

- If the blocking transaction rolls back, then the waiting transaction proceeds to change the previously locked row as if the other transaction never existed.
- If the blocking transaction commits and releases its locks, then the waiting transaction proceeds with its intended update to the newly changed row.

Read-Only Isolation Level

The **read-only isolation** level is similar to the serializable isolation level, but read-only transactions do not permit data to be modified in the transaction unless the user is SYS. Thus, read-only transactions are not susceptible to the ORA-08177 error. Read-only transactions are useful for generating reports in which the contents must be consistent with respect to the time when the transaction began.

Oracle Database achieves read consistency by reconstructing data as needed from the undo segments. Because undo segments are used in a circular fashion, the database can overwrite undo data. Long-running reports run the risk that undo data required for read consistency may have been reused by a different transaction, raising a snapshot too old error. Setting an **undo retention period**, which is the minimum amount of time that the database attempts to retain old undo data before overwriting it, appropriately avoids this problem.

REPEATABLE READ isolation level

SET TRANSACTION ISOLATION LEVEL READ COMMITTED;

to

SET TRANSACTION ISOLATION LEVEL REPEATABLE READ;

and execute it. The code for second session can be used as is.

With this simple change, we have a lock that is held by first session. This leads the second session to wait for the first one to complete before actually modify row

4. Construct two transactions and update one tuple in the database at the same time. Try to use the following SQL commands to view and understand the feed back information of transaction and lock status in the current system. (10 points)

```
mysql> show engine innodb status\G
***** 1. row *****
  Type: InnoDB
  Name:
  Status:
=====
2021-11-07 17:37:28 0x115c80 INNODB MONITOR OUTPUT
=====
Per second averages calculated from the last 12 seconds
-----
BACKGROUND THREAD
-----
srv_master_thread loops: 15 srv_active, 0 srv_shutdown, 15963 srv_idle
srv_master_thread log flush and writes: 0
-----
SEMAPHORES
-----
OS WAIT ARRAY INFO: reservation count 36
OS WAIT ARRAY INFO: signal count 36
RW-shared spins 0, rounds 0, OS waits 0
RW-excl spins 0, rounds 0, OS waits 0
RW-sx spins 0, rounds 0, OS waits 0
Spin rounds per wait: 0.00 RW-shared, 0.00 RW-excl, 0.00 RW-sx
-----
TRANSACTIONS
-----
Trx id counter 93590
Purge done for trx's n:o < 93589 undo n:o < 0 state: running but idle
History list length 0
LIST OF TRANSACTIONS FOR EACH SESSION:
---TRANSACTION 284282236696736, not started
0 lock struct(s), heap size 1128, 0 row lock(s)
---TRANSACTION 284282236695960, not started
0 lock struct(s), heap size 1128, 0 row lock(s)
---TRANSACTION 284282236695184, not started
0 lock struct(s), heap size 1128, 0 row lock(s)
---TRANSACTION 284282236693632, not started
0 lock struct(s), heap size 1128, 0 row lock(s)
---TRANSACTION 284282236692856, not started
0 lock struct(s), heap size 1128, 0 row lock(s)
---TRANSACTION 284282236692080, not started
0 lock struct(s), heap size 1128, 0 row lock(s)
---TRANSACTION 93589, ACTIVE 1312 sec
4 lock struct(s), heap size 1128, 6 row lock(s), undo log entries 2
MySQL thread id 9, OS thread handle 840340, query id 631 localhost 127.0.0.1 root
Trx read view will not see trx with id >= 93590, sees < 93590
-----
FILE I/O
-----
I/O thread 0 state: wait Windows aio (insert buffer thread)
```

5. Construct a deadlock situation.

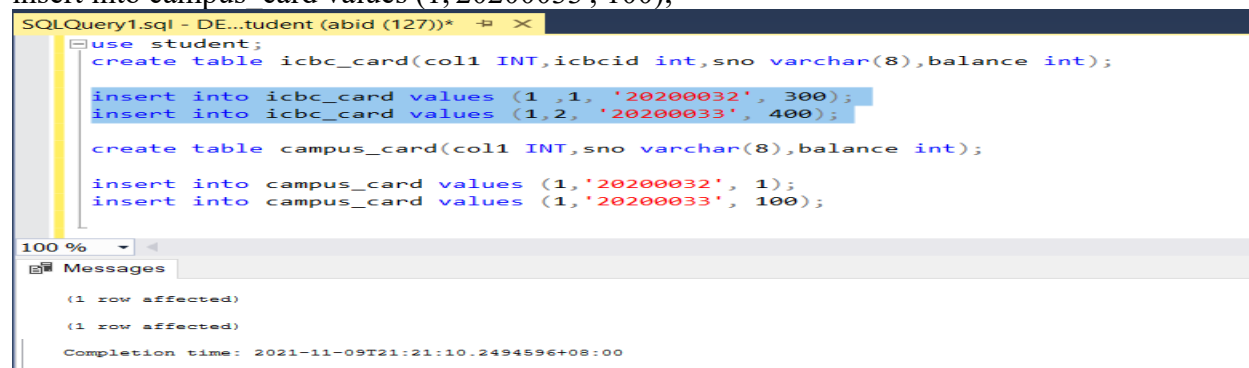
Object creation:

```
use student;  
create table icbc_card(col1 INT,icbcid int,sno varchar(8),balance int);
```

```
insert into icbc_card values (1,1, '20200032', 300);  
insert into icbc_card values (1,2, '20200033', 400);
```

```
create table campus_card(col1 INT,sno varchar(8),balance int);
```

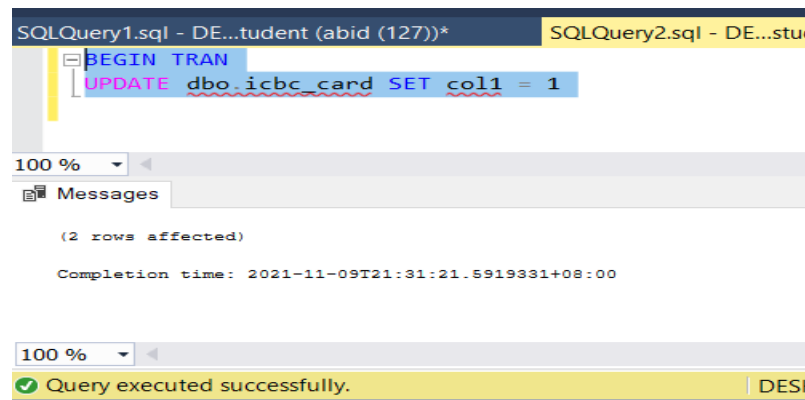
```
insert into campus_card values (1,'20200032', 1);  
insert into campus_card values (1,'20200033', 100);
```



```
SQLQuery1.sql - DE...tudent (abid (127))*  
use student;  
create table icbc_card(col1 INT,icbcid int,sno varchar(8),balance int);  
  
insert into icbc_card values (1,1, '20200032', 300);  
insert into icbc_card values (1,2, '20200033', 400);  
  
create table campus_card(col1 INT,sno varchar(8),balance int);  
  
insert into campus_card values (1,'20200032', 1);  
insert into campus_card values (1,'20200033', 100);  
  
100 %  
Messages  
  
(1 row affected)  
(1 row affected)  
  
Completion time: 2021-11-09T21:21:10.2494596+08:00
```

Then,we open up a new query window and paste this code and execute it

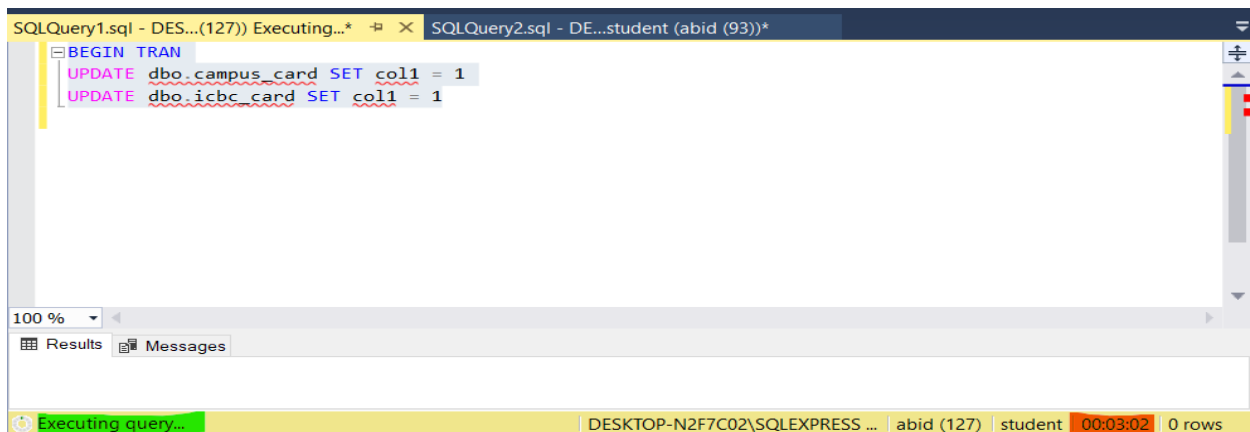
```
BEGIN TRAN  
UPDATE dbo.icbc_card SET col1 = 1
```



```
SQLQuery1.sql - DE...tudent (abid (127))* SQLQuery2.sql - DE...stuc  
BEGIN TRAN  
UPDATE dbo.icbc_card SET col1 = 1  
  
100 %  
Messages  
  
(2 rows affected)  
  
Completion time: 2021-11-09T21:31:21.5919331+08:00  
  
100 %  
Query executed successfully. DES
```

Again, we open up a new query window and paste this code and execute it

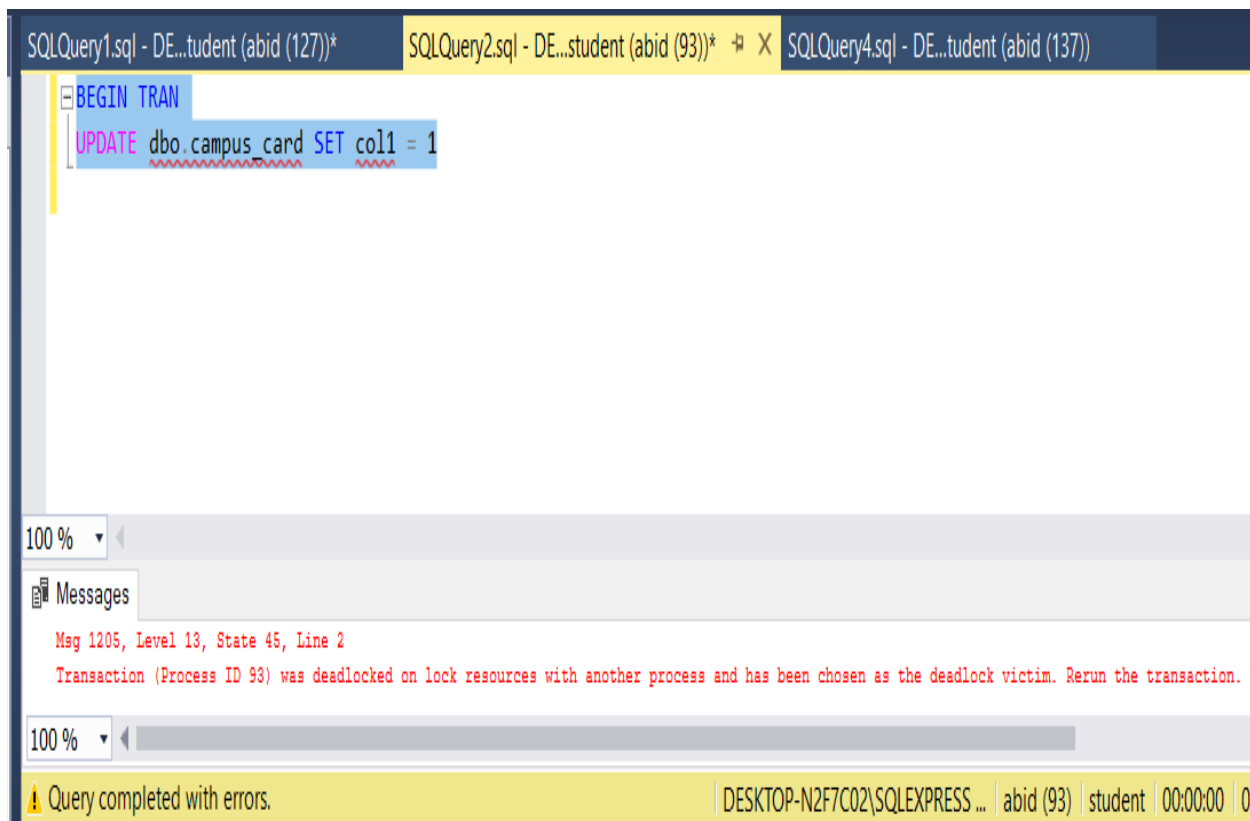
```
BEGIN TRAN  
UPDATE dbo.campus_card SET col1 = 1  
UPDATE dbo.icbc_card SET col1 = 1
```



After waiting certain moment of time,still the deadlock exist.

00:08:03

Then, we go back to your first query window (with the first BEGIN TRAN statement) and execute this code:



If we read the warning ,we can see that,we have successfully caused deadlock.

6. Construct the transaction containing some 'savepoint' and roll back to a savepoint at a certain time.

Solution:

Code:

```
GO
BEGIN TRANSACTION

SELECT*FROM dbo.icbc_card
insert into icbc_card values (3, '20200034', 400)
SAVE TRANSACTION SP1;

SELECT*FROM dbo.icbc_card;
insert into icbc_card values (4, '20200035', 500);
SAVE TRANSACTION SP2;

SELECT*FROM dbo.icbc_card;
insert into icbc_card values (5, '20200036', 600);
SAVE TRANSACTION SP3;

SELECT*FROM dbo.icbc_card;
insert into icbc_card values (6, '20200037', 700);
SAVE TRANSACTION SP4;

SELECT*FROM dbo.icbc_card;
insert into icbc_card values (7, '20200038', 800);
SAVE TRANSACTION S5;
```


GO

BEGIN TRANSACTION

```
SELECT*FROM dbo.icbc_card  
insert into icbc_card values (3, '20200034', 400)  
SAVE TRANSACTION SP1;
```

```
SELECT*FROM dbo.icbc_card;  
insert into icbc_card values (4, '20200035', 500);  
SAVE TRANSACTION SP2;
```

```
SELECT*FROM dbo.icbc_card;  
insert into icbc_card values (5, '20200036', 600);  
SAVE TRANSACTION SP3;
```

```
SELECT*FROM dbo.icbc_card;  
insert into icbc_card values (6, '20200037', 700);  
SAVE TRANSACTION SP4;
```

```
SELECT*FROM dbo.icbc_card;  
insert into icbc_card values (7, '20200038', 800);  
SAVE TRANSACTION S5;
```

GO

BEGIN TRANSACTION

```
SELECT*FROM dbo.icbc_card  
insert into icbc_card values (3, '20200034', 400)  
SAVE TRANSACTION SP1;
```

```
SELECT*FROM dbo.icbc_card;  
insert into icbc_card values (4, '20200035', 500);  
SAVE TRANSACTION SP2;
```

```
SELECT*FROM dbo.icbc_card;  
insert into icbc_card values (5, '20200036', 600);  
SAVE TRANSACTION SP3;
```

```
SELECT*FROM dbo.icbc_card;  
insert into icbc_card values (6, '20200037', 700);  
SAVE TRANSACTION SP4;
```

```
SELECT*FROM dbo.icbc_card;  
insert into icbc_card values (7, '20200038', 800);  
SAVE TRANSACTION S5;
```

```
SELECT * FROM dbo.icbc_card;  
insert into icbc_card values (7, '20200038', 800);  
SAVE TRANSACTION S5;
```

100 %

Results Messages

	icbcid	sno	balance
1	1	20200032	300
2	2	20200033	400
3	3	20200034	400
4	4	20200035	500
5	5	20200036	600
6	6	20200037	700
7	7	20200038	800

> ✓ Query executed successfully.

Code

ROLLBACK TO SP3;

SQLQuery1.sql - DE...tudent (abid (135))* X SQLQuery4.sql - DE...student (abid (9

```
ROLLBACK TRANSACTION SP3  
SELECT * FROM dbo.icbc_card;
```

100 %

Messages

Commands completed successfully.

Completion time: 2021-11-10T10:46:01.9303152+08:00

SQLQuery1.sql - DE...tudent (abid (135))* X SQLQuery4.sql - D

```

SAVE TRANSACTION S5;

ROLLBACK TRANSACTION SP3
SELECT * FROM dbo.icbc_card;

```

100 %

Results Messages

	icbcid	sno	balance
1	1	20200032	300
2	2	20200033	400
3	3	20200034	400
4	4	20200035	500
5	5	20200036	600

We can see that, only save point till SP3 is showing because there was 7 bank id. Now, after going to SP3 savepoint we can see 5 id.

Therefore, we have successfully created savepoints.

- Through experiments try to check all kinds of logs in MySQL: query log, error log and slow query log.

query log/general query log

Query 1 icbc_card x

```

1 • SET global general_log = 1;
2 • SET global log_output = 'table';
3 • CREATE TABLE `general_log` (
4   `event_time` timestamp NOT NULL DEFAULT CURRENT_TIMESTAMP
5     ON UPDATE CURRENT_TIMESTAMP,
6   `user_host` mediumtext NOT NULL,
7   `thread_id` bigint(21) unsigned NOT NULL,
8   `server_id` int(10) unsigned NOT NULL,
9   `command_type` varchar(64) NOT NULL,
10  `argument` mediumtext NOT NULL
11 ) ENGINE=CSV DEFAULT CHARSET=utf8 COMMENT='General log'

```

Limit to 1000 rows

Output

#	Time	Action	Message	Duration / Fetch
1	10:20:09	ROLLBACK TO SP3	Error Code: 1305. SAVEPOINT SP3 does not exist	0.000 sec
2	10:56:46	CREATE TABLE `general_log` (0 row(s) affected, 3 warning(s): 1681 integer display	0.015 sec
3	10:57:41	SET global general_log = 1	0 row(s) affected	0.000 sec
4	10:57:41	SET global log_output = 'table'	0 row(s) affected	0.000 sec
5	10:57:49	CREATE TABLE `general_log` (Error Code: 1050. Table 'general_log' already exists	0.016 sec

Query 1 icbc_card x

Limit to 1000 rows

```

11 ) ENGINE=CSV DEFAULT CHARSET=utf8 COMMENT='General log'
12
13 SELECT
14 *
15 FROM
16 mysql.general_log;

```

Result Grid

	event_time	user_host	thread_id	server_id	command_type	argument
▶	2021-11-10 10:57:49.013103	root[root] @ localhost [127.0.0.1]	9	1	Query	BLOB
	2021-11-10 10:59:41.262618	root[root] @ localhost [127.0.0.1]	9	1	Query	BLOB

general_log 2 x Read Only

Output

Action Output

#	Time	Action	Message	Duration / Fetch
✓ 1	10:59:41	SELECT * FROM mysql.general_log LIMIT 0,...	2 row(s) returned	0.000 sec / 0.000 sec

slow query log

```

mysql> show global variables like '%slow%';
+-----+-----+
| Variable_name | Value |
+-----+-----+
| log_slow_admin_statements | OFF |
| log_slow_extra | OFF |
| log_slow_replica_statements | OFF |
| log_slow_slave_statements | OFF |
| slow_launch_time | 2 |
| slow_query_log | ON |
| slow_query_log_file | DESKTOP-N2F7C02-slow.log |
+-----+-----+
7 rows in set (0.02 sec)

```

```

mysql> set global slow_query_log = 1;

```

```
mysql> select @@global.long_query_time;
+-----+
| @@global.long_query_time |
+-----+
| 10.000000 |
+-----+
1 row in set (0.00 sec)
```

```
mysql> set global long_query_time = 0;
Query OK, 0 rows affected (0.00 sec)

mysql> set global slow_query_log = 0;
Query OK, 0 rows affected (0.01 sec)
```

```
mysql> show create table mysql.slow_log \G
***** 1. row *****
      Table: slow_log
Create Table: CREATE TABLE `slow_log` (
  `start_time` timestamp(6) NOT NULL DEFAULT CURRENT_TIMESTAMP(6) ON UPDATE CURRENT_TIMESTAMP(6)
1.  `user_host` mediumtext NOT NULL,
..  `query_time` time(6) NOT NULL,
..  `lock_time` time(6) NOT NULL,
:lc  `rows_sent` int NOT NULL,
NU  `rows_examined` int NOT NULL,
NU  `db` varchar(512) NOT NULL,
NU  `last_insert_id` int NOT NULL,
..  `insert_id` int NOT NULL,
..  `server_id` int unsigned NOT NULL,
..  `sql_text` mediumblob NOT NULL,
..  `thread_id` bigint unsigned NOT NULL
.. ) ENGINE=CSV DEFAULT CHARSET=utf8mb3 COMMENT='Slow log'
1 row in set (0.01 sec)
```

View

I Disk (C:) > ProgramData > MySQL > MySQL Server 8.0 > Data >

Search Data

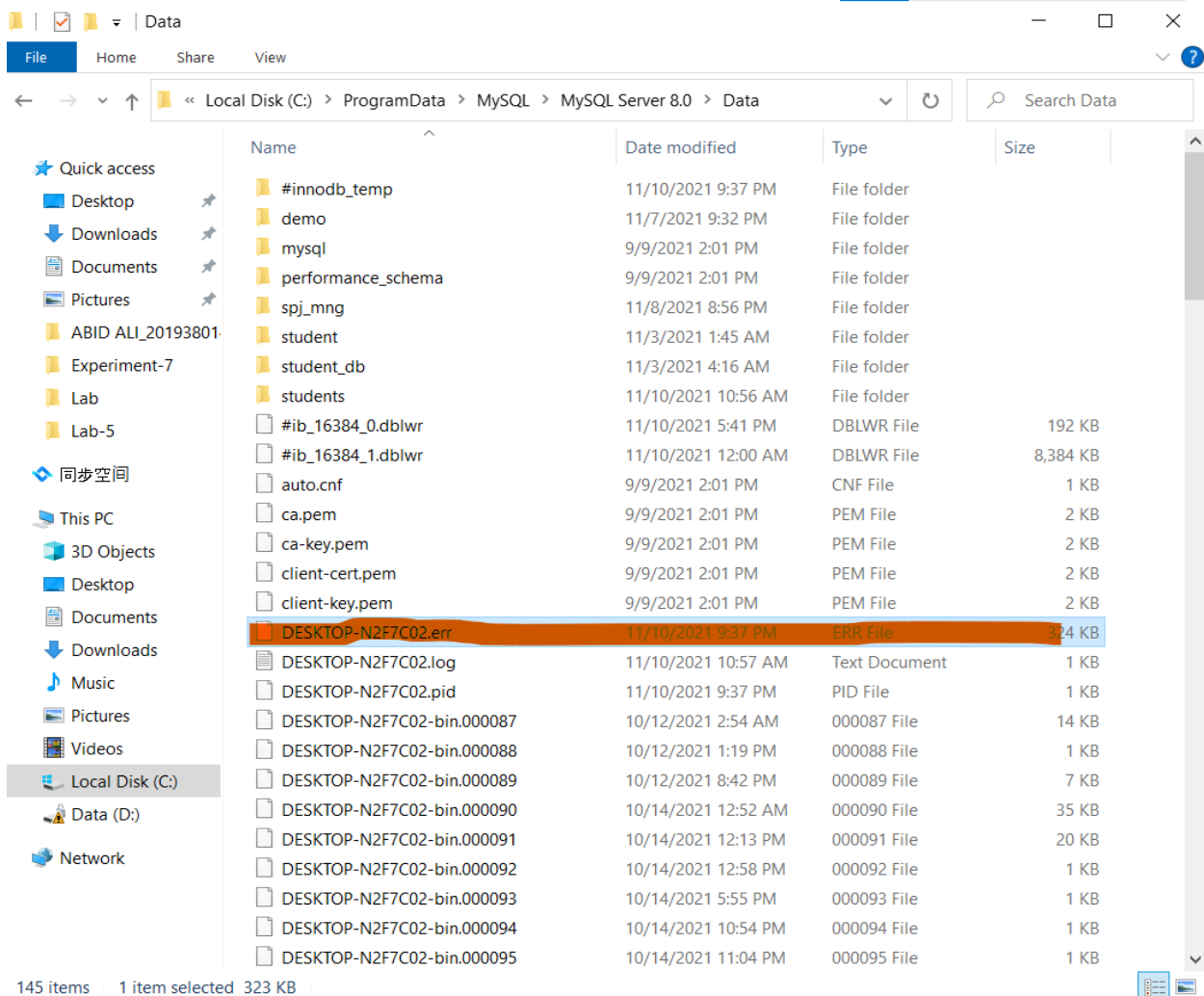
Name	Date modified	Type	Size
DESKTOP-N2F7C02-bin.000184	11/8/2021 3:38 PM	000184 File	1 KB
DESKTOP-N2F7C02-bin.000185	11/8/2021 10:12 PM	000185 File	4 KB
DESKTOP-N2F7C02-bin.000186	11/9/2021 3:55 AM	000186 File	1 KB
DESKTOP-N2F7C02-bin.000187	11/9/2021 4:29 PM	000187 File	1 KB
DESKTOP-N2F7C02-bin.000188	11/9/2021 5:10 PM	000188 File	1 KB
DESKTOP-N2F7C02-bin.000189	11/9/2021 5:10 PM	000189 File	1 KB
DESKTOP-N2F7C02-bin.000190	11/9/2021 10:47 PM	000190 File	1 KB
DESKTOP-N2F7C02-bin.000191	11/10/2021 12:00 AM	000191 File	1 KB
DESKTOP-N2F7C02-bin.000192	11/10/2021 1:46 AM	000192 File	1 KB
DESKTOP-N2F7C02-bin.000193	11/10/2021 11:47 AM	000193 File	5 KB
DESKTOP-N2F7C02-bin.000194	11/10/2021 3:47 PM	000194 File	1 KB
DESKTOP-N2F7C02-bin.000195	11/10/2021 5:26 PM	000195 File	1 KB
DESKTOP-N2F7C02-bin.000196	11/10/2021 5:39 PM	000196 File	1 KB
DESKTOP-N2F7C02-bin.000197	11/10/2021 6:10 PM	000197 File	1 KB
DESKTOP-N2F7C02-bin.000198	11/10/2021 9:37 PM	000198 File	1 KB
DESKTOP-N2F7C02-bin.000199	11/10/2021 9:37 PM	000199 File	1 KB
DESKTOP-N2F7C02-bin.index	11/10/2021 9:37 PM	INDEX File	4 KB
DESKTOP-N2F7C02-slow.log	11/10/2021 9:37 PM	Text Document	39 KB

Error log

```
mysql> show variables like '%error%';
```

Variable_name	Value
binlog_error_action	ABORT_SERVER
error_count	0
log_error	.\DESKTOP-N2F7C02.err
log_error_services	log_filter_internal; log_sink_internal
log_error_suppression_list	
log_error_verbosity	2
max_connect_errors	100
max_error_count	1024
performance_schema_error_size	4946
replica_skip_errors	OFF
slave_skip_errors	OFF

11 rows in set (0.01 sec)



8. Use mysqlbinlog to view the transaction log of the database, and try to recover the data according to the following scenarios.

```
mysql> use db1;
Database changed
mysql> create table t1(id int);
Query OK, 0 rows affected (0.05 sec)

mysql> create table t2(id int);
Query OK, 0 rows affected (0.03 sec)

mysql> insert into t1 values(11);
Query OK, 1 row affected (0.01 sec)

mysql> insert into t1 values(12);
Query OK, 1 row affected (0.00 sec)

mysql> insert into t1 values(13);
Query OK, 1 row affected (0.00 sec)

mysql> insert into t2 values(21);
Query OK, 1 row affected (0.00 sec)

mysql> insert into t2 values(22);
Query OK, 1 row affected (0.01 sec)

mysql> insert into t2 values(23);
Query OK, 1 row affected (0.00 sec)

mysql> drop table t1;
Query OK, 0 rows affected (0.02 sec)

mysql> insert into t2 values(24);
Query OK, 1 row affected (0.00 sec)

mysql>
```

We can see the binlog transaction file

```
mysql> show master status;
+-----+-----+-----+-----+-----+
| File                               | Position | Binlog_Do_DB | Binlog_Ignore_DB | Executed_Gtid_Set |
+-----+-----+-----+-----+-----+
| DESKTOP-N2F7C02-bin.000201        | 156      |              |                  |                  |
+-----+-----+-----+-----+-----+
1 row in set (0.00 sec)
```

```
mysql> SHOW BINLOG EVENTS IN 'DESKTOP-N2F7C02-bin.000201';
```

Log_name	Pos	Event_type	Server_id	End_log_pos	Info
DESKTOP-N2F7C02-bin.000201	4	Format_desc	1	125	Server ver: 8.0.26, Binlog ver: 4
DESKTOP-N2F7C02-bin.000201	125	Previous_gtid	1	156	
DESKTOP-N2F7C02-bin.000201	156	Anonymous_Gtid	1	233	SET @@SESSION.GTID_NEXT= 'ANONYMOUS'
DESKTOP-N2F7C02-bin.000201	233	Query	1	342	use `db1`; create table t1(id int) /* xid=49 */
DESKTOP-N2F7C02-bin.000201	342	Anonymous_Gtid	1	419	SET @@SESSION.GTID_NEXT= 'ANONYMOUS'
DESKTOP-N2F7C02-bin.000201	419	Query	1	528	use `db1`; create table t2(id int) /* xid=50 */
DESKTOP-N2F7C02-bin.000201	528	Anonymous_Gtid	1	607	SET @@SESSION.GTID_NEXT= 'ANONYMOUS'
DESKTOP-N2F7C02-bin.000201	607	Query	1	681	BEGIN
DESKTOP-N2F7C02-bin.000201	681	Table_map	1	728	table_id: 96 (db1.t1)
DESKTOP-N2F7C02-bin.000201	728	Write_rows	1	768	table_id: 96 flags: STMT_END_F
DESKTOP-N2F7C02-bin.000201	768	Xid	1	799	COMMIT /* xid=51 */
DESKTOP-N2F7C02-bin.000201	799	Anonymous_Gtid	1	878	SET @@SESSION.GTID_NEXT= 'ANONYMOUS'
DESKTOP-N2F7C02-bin.000201	878	Query	1	952	BEGIN
DESKTOP-N2F7C02-bin.000201	952	Table_map	1	999	table_id: 96 (db1.t1)
DESKTOP-N2F7C02-bin.000201	999	Write_rows	1	1039	table_id: 96 flags: STMT_END_F
DESKTOP-N2F7C02-bin.000201	1039	Xid	1	1070	COMMIT /* xid=52 */
DESKTOP-N2F7C02-bin.000201	1070	Anonymous_Gtid	1	1149	SET @@SESSION.GTID_NEXT= 'ANONYMOUS'
DESKTOP-N2F7C02-bin.000201	1149	Query	1	1223	BEGIN
DESKTOP-N2F7C02-bin.000201	1223	Table_map	1	1270	table_id: 96 (db1.t1)
DESKTOP-N2F7C02-bin.000201	1270	Write_rows	1	1310	table_id: 96 flags: STMT_END_F
DESKTOP-N2F7C02-bin.000201	1310	Xid	1	1341	COMMIT /* xid=53 */
DESKTOP-N2F7C02-bin.000201	1341	Anonymous_Gtid	1	1420	SET @@SESSION.GTID_NEXT= 'ANONYMOUS'
DESKTOP-N2F7C02-bin.000201	1420	Query	1	1494	BEGIN
DESKTOP-N2F7C02-bin.000201	1494	Table_map	1	1541	table_id: 97 (db1.t2)
DESKTOP-N2F7C02-bin.000201	1541	Write_rows	1	1581	table_id: 97 flags: STMT_END_F
DESKTOP-N2F7C02-bin.000201	1581	Xid	1	1612	COMMIT /* xid=54 */
DESKTOP-N2F7C02-bin.000201	1612	Anonymous_Gtid	1	1691	SET @@SESSION.GTID_NEXT= 'ANONYMOUS'
DESKTOP-N2F7C02-bin.000201	1691	Query	1	1765	BEGIN
DESKTOP-N2F7C02-bin.000201	1765	Table_map	1	1812	table_id: 97 (db1.t2)
DESKTOP-N2F7C02-bin.000201	1812	Write_rows	1	1852	table_id: 97 flags: STMT_END_F
DESKTOP-N2F7C02-bin.000201	1852	Xid	1	1883	COMMIT /* xid=55 */
DESKTOP-N2F7C02-bin.000201	1883	Anonymous_Gtid	1	1962	SET @@SESSION.GTID_NEXT= 'ANONYMOUS'
DESKTOP-N2F7C02-bin.000201	1962	Query	1	2036	BEGIN
DESKTOP-N2F7C02-bin.000201	2036	Table_map	1	2083	table_id: 97 (db1.t2)
DESKTOP-N2F7C02-bin.000201	2083	Write_rows	1	2123	table_id: 97 flags: STMT_END_F
DESKTOP-N2F7C02-bin.000201	2123	Xid	1	2154	COMMIT /* xid=56 */
DESKTOP-N2F7C02-bin.000201	2154	Anonymous_Gtid	1	2231	SET @@SESSION.GTID_NEXT= 'ANONYMOUS'
DESKTOP-N2F7C02-bin.000201	2231	Query	1	2356	use `db1`; DROP TABLE `t1` /* generated by server */ /* xid=57 */
DESKTOP-N2F7C02-bin.000201	2356	Anonymous_Gtid	1	2435	SET @@SESSION.GTID_NEXT= 'ANONYMOUS'
DESKTOP-N2F7C02-bin.000201	2435	Query	1	2509	BEGIN
DESKTOP-N2F7C02-bin.000201	2509	Table_map	1	2556	table_id: 97 (db1.t2)
DESKTOP-N2F7C02-bin.000201	2556	Write_rows	1	2596	table_id: 97 flags: STMT_END_F
DESKTOP-N2F7C02-bin.000201	2596	Xid	1	2627	COMMIT /* xid=58 */

Delete

We can see the operation that was performed.

```
mysql> mysqlbinlog.exe DESKTOP-N2F7C02-bin.000201 > test_000201;
```

We restored the information in text file

```
66 /*!0098 set@@session.immediate_server_version =6090*//*!*/;
67 table_id: 97 (db1.t2)
69 table_id: 97 flags: STMT_END_F
70 SET @@SESSION.GTID_NEXT= 'ANONYMOUS'
70 use `db1`; DROP TABLE `t1` /* generated by server */ /* xid=57 */
71 /*!*/;
72 # at 2356
73 SET @@SESSION.GTID_NEXT= 'ANONYMOUS'
74 #original_commit_timestamp 22:22:43 server id 1 end_log_pos 4433 CRJ54 0xed43
76 #immediatel_commit_timestamp 22:22:43 server id 1 end_log_pos 4433 CRJ54 0xed43
77 table_id: 97 (db1.t2)
78 #original_commit_timestamp 22:22:43 server id 1 end_log_pos 4499 CRJ54 0xed43
79 #immediate_commit_timestamp 22:22:43 server id 1 end_log_pos 4499 CRJ54 0xed43
80 table_id: 97 flags: STMT_END_F
81 COMMIT /* xid=58 */
```

We can see the delete operation that we used in the mysql query and restore our data from there.

Problems:

At first ,I was having problem with the concept transaction, and how to create the transaction. Didn't know how to implement data inconsistency problems in concurrency operation, and can use lock and ioslation mechanisms. These topics were mixed of theoretical and practical concept was struggling at the beginning.

Solutions:

To solve these problems which I faced during doing this practical, I took help from internet especially YouTube, StackOverflow and W3school to get information about these errors for the solution. I also asked the teacher to help me understand them. And provided instructions helped to solve some of my errors during the experiment.

Attachments:

1) DB7_2019380141_ABID ALI.pdf

References:

- 1) <https://www.w3schools.com/>
- 2) <https://stackoverflow.com/>
- 3) <https://youtube.com/>
- 4) <https://www.thegeekstuff.com/2017/08/mysqlbinlog-examples/>
- 5) <https://www.youtube.com/watch?v=iCizaSoJd5w>
- 6) <https://www.youtube.com/watch?v=xYysvuDAX70>

