(4)

of code B are not distinct and rence it become singular

3. Uniquely Decodable codes;

A non-singular code is said to be "uniquely de codable" is every code would de codable " or " uniquely de cipherable" is every code would present in a long received sequence can be uniquely identified.

eg: Received a sequence at the receiver on the absence of noise; R = 001100.

O code A is used:

Decoded ou : S, Sq S,

-> It is the only possible way.

@ Code B is wed :

Decoded on 8 82 84 82 or

8, 8, 84 8, 8, or

81 81 84 82 07

S. S. S.

-> notricinequely decodable.

4. Instentaneous Codes:

A uniquely decodable well is said to be instentaneous if it is possible to recognize the end of any cocle word in any received sequence, with out reference to the succeeding symbols. There is no time delay in the process of decoding of decoding is einstentaneous.

eg: Consider three codes

Source symbols	Code C	Code D	Code F
31	00	0	0
S2	, O 1:10	10 -	01/
W W MO W	~34 o	llo o	ا دال جورا
SA	11	1110	OIII

Received sequence; R = 001100
using différent coding schemes:

Codec : SI SA SI : - Laston

Code D : SI SI S3 SI

Code E : S, S, S, S,

Receiver, it can cimmediately decoded as SI, because no other code woords are started with a o' A second zero is decoded as SI again. Here when a o' arriver, it ends each code would, it and changinary "comma" may be curerted at the receiver after the receiver of every zero. Then it is decoded. So it is celled a "COHMA CODE"

5. Optimal Codes : 000

An contentaneous code is said to be optimal code " if it has "ninimum average length L" for a some with a given purbability assignment for the some symbol.

code -> Non block

-> Block -> dengular

-> Non singular -> Non-uniquely

-> decodable
-> curiquely -> Non enst-entaneous

-> decodable
-> instentaneous -> Non optimal
-> optimal

Code peopoily as tree diagram:

Prefix of a code:

Let $X = x_4 x_2 - x_m$ be a code coold of some code, then the sequence $x_4 x_2 - x_4 y$ for all $y \ge x_m$ one calced 'prefixes' of code X.

Let
$$x = 0$$
 | | $m = 4$

when
$$j=1$$
; ≥ 0 ≥ 1 ≥ 1

Test por contantaneous Property (prefix property):

A recessary and sufficient condition for a cuniquely decodable code to be constantaneous is that "No complète word of a code would be a prefix of any other code word"

- 18 prefixes au not present: code is instantaneous.

Even if one prezix is present, the code will not be cirstantanéous.

Applying lest to code Er for Sy = 0111, prefixes are 0,01,011, which one the code words of si,siffs. So code E is not instantaneous.

Apply to code D: $S_q = 1110$; prefixes one 1,11,111, and which are not code woulds; do code D is instantaneous.

t properly in our antifical :