

Consider a binary source 'S' with symbols  $s_1$  &  $s_2$   
2nd Extension: with probabilities  $p_1$  &  $p_2$  where  $p_1 + p_2 = 1$ .

Second extension of a binary source will have  $2^2 = 4$  symbols given by.

$s_1 s_1$  with probabilities;  $p_1 p_1 = p_1^2$

$s_1 s_2$  " "  $p_1 p_2$

$s_2 s_1$  " "  $p_2 p_1$  and

$s_2 s_2$  " "  $p_2 p_2 = p_2^2$

Sum of all probabilities of the 2nd extended source.

$$p_1^2 + p_1 p_2 + p_2 p_1 + p_2^2 = p_1^2 + 2p_1 p_2 + p_2^2 = (p_1 + p_2)^2 = 1$$

Entropy of basic binary source;

$$H(S) = \sum_{k=1}^2 p_k \log \frac{1}{p_k}$$

$$= p_1 \log \frac{1}{p_1} + p_2 \log \frac{1}{p_2} \quad \text{--- ①}$$

Entropy of the 2nd extended source;

$$H(S^2) = \sum_{k=1}^4 p_k \log \frac{1}{p_k}$$

$$= p_1^2 \log \frac{1}{p_1^2} + p_1 p_2 \log \frac{1}{p_1 p_2} + p_1 p_2 \log \frac{1}{p_1 p_2} + p_2^2 \log \frac{1}{p_2^2}$$

$$= 2p_1^2 \log \frac{1}{p_1^2} + 2p_1 p_2 \log \frac{1}{p_1 p_2} + 2p_2^2 \log \frac{1}{p_2^2}$$

$$= 2p_1^2 \log \frac{1}{p_1} + 2p_1 p_2 \log \frac{1}{p_1} + 2p_1 p_2 \log \frac{1}{p_2} + 2p_2^2 \log \frac{1}{p_2}$$

$$= 2p_1 \left( \frac{p_1 + p_2}{1} \right) \log \frac{1}{p_1} + 2p_2 \left( \frac{p_1 + p_2}{1} \right) \log \frac{1}{p_2}$$

$$= 2 \left[ p_1 \log \frac{1}{p_1} + p_2 \log \frac{1}{p_2} \right]$$

$$\boxed{H(S^2) = 2 H(S)} \quad \text{from eq ①}$$

Similarly the entropy of 3rd extended source can be

$$\boxed{H(S^3) = 3 \cdot H(S)}$$

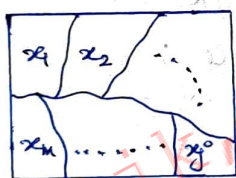
Generalizing, the  $n$ th extension of the basic binary source will have  $2^n$  symbols and the entropy of the  $n$ th extended source is given by

$$H(S^n) = n \cdot H(S)$$

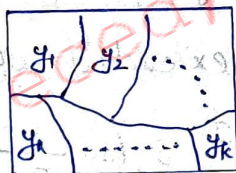
## JOINT ENTROPY AND CONDITIONAL ENTROPY

So far we considered only a single probability scheme i.e. one dimensional probability scheme; probability belong either to the transmitter or to the receiver. That means we are able to study the behaviour of either transmitter or receiver.

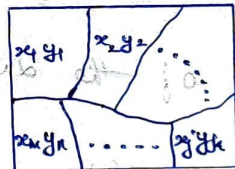
But to study the behaviour of a communication system, we must simultaneously study the behaviour of the transmitter and receiver. This will give the concept of two dimensional probability scheme.



Sample space,  $S_1$



$S_2$



$S = S_1 S_2$

There are two finite discrete sample spaces  $S_1$  and  $S_2$  and their product space will be  $S = S_1 S_2$ .

$$\text{Let } [X] = [x_1, x_2, \dots, x_m]$$

$$\text{and } [Y] = [y_1, y_2, \dots, y_k]$$

be the set of events in  $S_1$  and  $S_2$  respectively.

Each event  $x_j$  of  $S_1$  occur in connection (relation) with any event  $y_k$  of  $S_2$ .

The complete set of events in  $S = S_1 S_2$  is  $[XY]$ .

$$[xy] = \begin{bmatrix} x_1 y_1 & x_1 y_2 & \dots & x_1 y_k & \dots & x_1 y_n \\ x_2 y_1 & x_2 y_2 & \dots & x_2 y_k & \dots & x_2 y_n \\ \vdots & \vdots & & \vdots & & \vdots \\ x_j y_1 & x_j y_2 & \dots & x_j y_k & \dots & x_j y_n \\ \vdots & \vdots & & \vdots & & \vdots \\ x_m y_1 & x_m y_2 & \dots & x_m y_k & \dots & x_m y_n \end{bmatrix}$$

The three sets of complete probability schemes:

$$p(x) = [p(x_j)] = [p(x_1) \ p(x_2) \ \dots \ p(x_j) \ \dots \ p(x_m)]$$

$$p(y) = [p(y_k)] = [p(y_1) \ p(y_2) \ \dots \ p(y_k) \ \dots \ p(y_n)]$$

$$p(xy) = [p(x_j, y_k)] =$$

$$\begin{matrix} & y_1 & y_2 & \dots & y_k & \dots & y_n \\ \begin{matrix} x_1 \\ x_2 \\ \vdots \\ x_j \\ \vdots \\ x_m \end{matrix} & \begin{bmatrix} p(x_1, y_1) & p(x_1, y_2) & \dots & p(x_1, y_k) & \dots & p(x_1, y_n) \\ p(x_2, y_1) & p(x_2, y_2) & \dots & p(x_2, y_k) & \dots & p(x_2, y_n) \\ \vdots & \vdots & & \vdots & & \vdots \\ p(x_j, y_1) & p(x_j, y_2) & \dots & p(x_j, y_k) & \dots & p(x_j, y_n) \\ \vdots & \vdots & & \vdots & & \vdots \\ p(x_m, y_1) & p(x_m, y_2) & \dots & p(x_m, y_k) & \dots & p(x_m, y_n) \end{bmatrix} \end{matrix}$$

$p(xy)$  is called "Joint Probability Matrix". Properties of

JPM are :

① By adding the elements of JPM columnwise, we can obtain the ~~probability~~ probability of output symbols.

$$\sum_{j=1}^m p(x_j, y_k) = p(y_k).$$

② By adding the elements of JPM rowwise, we can obtain the probability of input symbols.

$$\sum_{k=1}^n p(x_j, y_k) = p(x_j).$$

③ The sum of all the elements of JPM is equal to unity

$$\sum_{j=1}^m \sum_{k=1}^n p(x_j, y_k) = 1$$