Incorporating Trilinos Data Classes into an Application

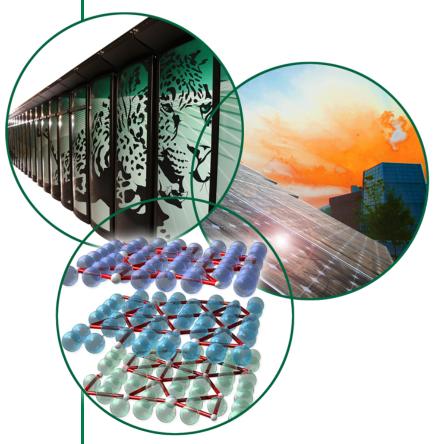
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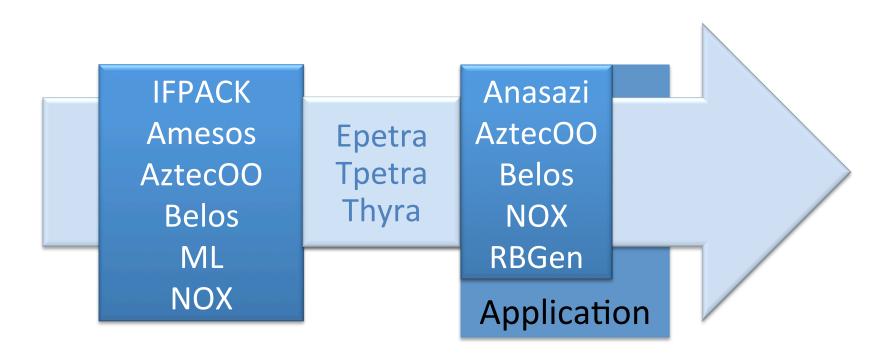






Trilinos Overview

- POV: Trilinos is all about operators.
- Operator abstraction is the heart of the solvers, and the main interface between applications and Trilinos.





What is an Operator?

- Operator is a Petra Object Model concept
- Operator encapsulates the effect of some action (typically linear) from one vector into another.

```
Epetra
virtual Epetra Operator::Apply(const Epetra MultiVector &X,
                                     Epetra MultiVector &Y) const;
                                                                Tpetra
virtual Tpetra::Operator<S>::apply(
        const Tpetra::MultiVector<S> &X,
              Tpetra::MultiVector<S> &Y,
        ...) const;
virtual Thyra::LinearOpBase<S>::apply(...,
         const Thyra::MultiVectorBase<S> &X,
               Thyra::MultiVectorBase<S> *Y,
         ...) const;
```

Example Operators

- Direct operator implementations:
 - sparse matrix (Epetra/Tpetra)
 - dense matrix (Epetra/Tpetra)
 - direct solve (Amesos, IFPACK, Epetra/Tpetra)
- Indirect operator implementations:
 - iterator linear solve (AztecOO, Belos)
 - composition of operators (Epetra, Thyra)
- Application-specific:
 - **???**



Operator Interface

Operator defines three main methods:

```
void apply(const Vector x, Vector y);
Map getDomainMap();
Map getRangeMap();
```

- apply () obviously realizes the effect of the operator.
- The others describe the properties of the input/output vectors, namely, the domain and range of the operator.

```
x = createVector<S>( op->getDomainMap() );
y = createVector<S>( op->getRangeMap() );
op->apply( *x, *y );
```



Whither the Map

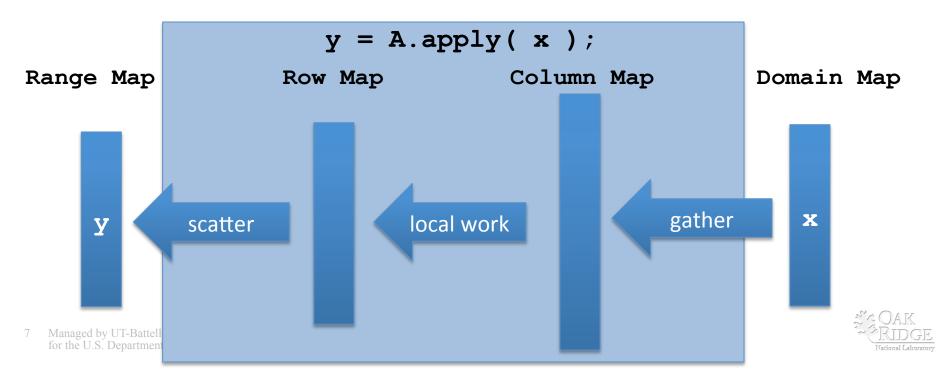
- The Petra Object Model describes the construction and manipulation of distributed memory objects.
- The Map concept addresses two main concerns:
 - distribution of a global object across local memories
 - translation between global indices and local indices
- A related concept is the Distributed Object, an object that is distributed in this manner.
- A Distributed Object is distributed according to a Row Map:
 - examples: Vector, MultiVector, CrsMatrix, CrsGraph
- This describes a row distribution of the data.



Other Common Maps

In addition to Domain Map and Range Map, Operator objects often utilize two other Map objects:

- Row Map indicates which rows of the output are computed, a row distribution of the work.
- Column Map indicates the rows of the input vector necessary to perform that work, a function of the Operator and the Row Map



Supporting cast

- The gather/scatter in a communicating Operator is facilitated by Import and Export objects.
- These contain information needed to transfer data between Distributed Objects described by different Maps.
- These are constructed in advance, analyzing source and destination Maps for significant quantities:
 - remote indices, owners
 - local/permuted indices
 - total transfer size (for allocating buffers)



Impact on App Developers

- What does this mean for app developers?
- For current Epetra users moving to Tpetra, the concepts are the same.
 - The semantics and implementations are mostly the same, but translation of syntax is needed in many places.
- For developers new to Trilinos, you may need some understanding of these.
 - This is necessary, at the least, when instantiating vectors and defining operators.



Side by side: Creating Maps

- Initialize MPI, create a communicator
- Create a uniformly distributed map with contiguously allocated elements, query the list of elements

```
MPI_Init(&argc,&argv);
Epetra_MpiComm comm( MPI_COMM_WORLD );
Epetra_Map map(NumGlobalElements, 0, comm);
const int NumMyElements = map.NumMyElements();
std::vector<int> MyGlobalElements(NumMyElements);
map.MyGlobalElements( &MyGlobalElements[0] );
```

```
Teuchos::GlobalMPISession mpiSession(&argc,&argv,...);
Platform &plat = Tpetra::DefaultPlatform::getDefaultPlatform();
RCP<const Teuchos::Comm<int> > comm = platform.getComm();
RCP<const Tpetra::Map<int> > map;
map = Tpetra::createUniformContigMap<int>(numGlobalElements, comm);
const size_t numMyElements = map->getNodeNumElements();
Teuchos::ArrayView<const int> myGlobalElements;
myGlobalElements = map->getNodeElementList();
```

Side by side: Creating Maps

- Create a map with specified number of elements
- Create a map with a specified list of elements

```
Epetra_Map mapContig(numGlobal, numLocal, 0, comm);

std::vector<int> elemList(numLocal);
// ... fill elemList ...
Epetra_Map mapFromList(numGlobal, numLocal, &elemList[0], 0, comm);
```

```
RCP<const Tpetra::Map<int> > mapContig, mapFromList;
mapContig = Tpetra::createContigMap<int>(numGlobal, numLocal, comm);

Teuchos::Array<int> elemList(numLocal);
// ... fill elemList ...
mapFromList = Tpetra::createNonContigMap<int>(elemList(), comm);
```



Side by side: Vectors and MultiVectors

- Create Vectors or MultiVectors from maps
- Extract Vectors from MultiVectors
- Extract data pointers from Vectors

```
Epetra_Vector vec(map);
Epetra_MultiVector mvec(map, numVecs);
for (i=0; i<numVecs; ++i) {
   Epetra_Vector * vecptr = mvec(i);
   double * vecdata = mvec[i];
}</pre>
```

```
RCP< Tpetra::Vector<double> > vec;
vec = Tpetra::createVector<double>(map);
RCP< Tpetra::MultiVector<double> > mvec;
mvec = Tpetra::createMultiVector<double>(map, numVecs);
for (i=0; i<numVecs; ++i) {
   RCP< Tpetra::Vector<double> > veci = mvec->getVector(i);
   ArrayRCP<double> vecdata = veci->get1dViewNonConst();
}
```

Side by side: Writing an Operator

A simple, non-communicating diagonal Operator

```
class DiagOp : public Epetra_Operator {
  int Apply(const Epetra_MultiVector &X, Epetra_MultiVector &Y) {
    for (int v=0; v < X.NumVectors(); ++v)
        for (int i=0; i < X.MyLength(); ++i)
        Y[v][i] = diag_[i] * X[v][i];
    return 0;
}
};</pre>
```

Side by side: Using Anasazi/Belos

 Switching Anasazi/Belos from Epetra to Tpetra requires only modifying the data initialization.

Side by side: Communicating Operators

Communicating Operators with Import/Export

```
Epetra
MyOp::MyOp (const Epetra Map &rowMap,
           const Epetra Map &rangeMap,
           const Epetra Map &domainMap)
  rowMap (rowMap);
  // ... build colMap from rowMap and operator specifics ...
  Importer = new Epetra Import(colMap , domainMap);
 Exporter = new Epetra Export(rowMap , rangeMap);
  ImportVector = new Epetra MultiVector(colMap , blockSize);
 ExportVector = new Epetra MultiVector(rowMap , blockSize);
int MyOp::Apply(const Epetra MultiVector &X, Epetra MultiVector &Y)
  ImportVector ->Import(x, *Importer , Insert);
  xp = (const double*) ImportVector ->Values();
  yp = ( double*) ExportVector ->Values();
  // ... apply operator from xp to yp[r] for each r in rowMap ...
  Y.Export(*ExportVector , *Exporter , Add));
  return 0;
```

Side by side: Communicating Operators

Tpetra is conceptually the same... with caveats.

- POM concepts direct the code to a similar outline.
- Support for multi-core/GPUs begins to highlight weaknesses in the grab-the-pointer-and-run model.



Addressing OpenMP Nodes: Epetra

- New releases of Epetra support OpenMP for sparse matrix multiply and for vector operations; GPUs soon.
- Therefore, custom Epetra_Operator objects in an OpenMP-enabled build should also support OpenMP.

```
int MyOp::Apply(const Epetra MultiVector &X, Epetra MultiVector &Y)
  ImportVector ->Import(x, *Importer , Insert);
 xp = (const double*) ImportVector ->Values();
 vp = ( double*) ExportVector ->Values();
#ifdef EPETRA HAVE OMP
#pragma omp parallel for ...
  // ... apply operator from xp to yp[r] for each r in rowMap ...
#else
 // ... apply operator from xp to yp[r] for each r in rowMap ...
#endif
 Y.Export(*ExportVector , *Exporter , Add));
  return 0;
```

Addressing Generic Nodes: Tpetra

• Previous slide showed all template parameters in Tpetra::Operator

Tpetra::Operator <t, go,="" lo,="" node=""></t,>					
T	LO	GO	Node		
Scalar field over which the operator applies	Ordinal type for local indices	Ordinal type for global indices	Kokkos node for shared-memory parallel node		

- Scalar type is generic, supporting broad capability.
- Ordinals templated separately to simultaneously maximize efficiency and capability.
- Node type is a template parameter in order to utilize Kokkos template meta-programming shared-memory API.



Impact of Templated Classes on Apps

- Templated classes force apps into an immediate decision:
 - whether or not to use a generic programming approach?
- Option #1: utilize typedefs to hide templates
 - app avoids templates, but is hard-coded to this choice

• Option #2: embrace templates, suffer the consequences

```
template <class Scalar>
void everyMethodInMyWholeApplication(Operator<Scalar> &tedium) {
   // ... insert syntactic sugar here ...
}
```



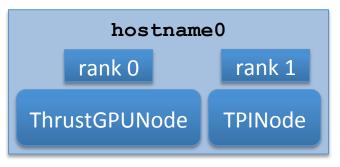
Tpetra::HybridPlatform

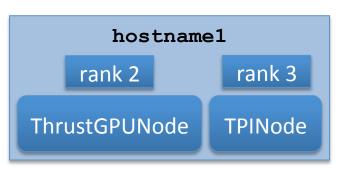
Encapsulate "main" in a templated class method:

- HybridPlatform maps the communicator rank to the Node type, instantiates a node and the user routine.
 - main() becomes boilerplate code.

HybridPlatform Machine File

round-robin assignment	interval assignment	explicit assignment	default
% M=N	[M,N]	=N	default





```
<ParameterList>
  <ParameterList name="%2=0">
    <Parameter name="NodeType"</pre>
                                       type="string"
                                                       value="Kokkos::ThrustGPUNode"/>
    <Parameter name="Verbose"</pre>
                                                       value="1"/>
                                       type="int"
    <Parameter name="Device Number"</pre>
                                       type="int"
                                                       value="0"/>
                                                       value="4"/>
                                       type="int"
    <Parameter name="Node Weight"</pre>
  </ParameterList>
  <ParameterList name="%2=1">
    <Parameter name="NodeType"</pre>
                                       type="string"
                                                       value="Kokkos::TPINode"/>
    <Parameter name="Verbose"
                                                       value="1"/>
                                       type="int"
    <Parameter name="Num Threads"</pre>
                                       type="int"
                                                       value="15"/>
                                                       value="15"/>
    <Parameter name="Node Weight"</pre>
                                       type="int"
  </ParameterList>
```

</ParameterList>

What about Parallel Tpetra Operators?

- These can be written in the same way as Tpetra does:
 - Kokkos Shared-Memory Parallel Node API

```
template <class T, class LO, class GO, class Node>
void MyOp<T,LO,GO,Node>::apply(
          const Tpetra::MultiVector<T,LO,GO,Kokkos::SerialNode> &X,
                Tpetra::MultiVector<T,LO,GO,Kokkos::SerialNode> &Y,
          ...) {
  RCP<Node> node = X.getRowMap()->getNode();
  importMV ->doImport(X, *importer, INSERT);
  for (int v=0; v < X.getNumVectors(); ++v) {</pre>
    ArrayRCP<const T> xp = importMV ->getLocalMV()->getValues(v);
   ArrayRCP<T> yp;
    yp = exportMV ->getLocalMVNonConst()->getValuesNonConst(v);
   MyKernel<T,LO> kern(xp,yp,...);
    node->parallel for(0, X.getLocalLength(), kern);
  Y.doExport(*exportMV , *exporter, ADD);
```

All but the kernel-specific part is boilerplate code...



Kokkos Parallel Constructs

- Parallel for: execute loop iterations in parallel
- User-defined struct (work-data pair) contains:
 - the necessary data and execute (int iter)
- Parallel reduce: reduce implicit set of elements in parallel via user-specified associative binary operation
 - typedef ReductionType
 - ReductionType identity()
 - ReductionType generate(int i)
 - ReductionType reduce(ReductionType a, ReductionType b)
- Template meta-programming fuses generic loop skeleton with user data and kernel specifications.

```
Node::parallel_for <WDP>(int beg, int end, WDP args);
Node::parallel_reduce<WDP>(int beg, int end, WDP args);
```



Kokkos parallel_for example

• Consider simple vector axpy: $y = \alpha * x + y$

```
template <class Scalar>
struct AxpyOp {
  Scalar alpha;
  const Scalar *x;
  Scalar *y;
  inline void execute(int i) {
    y[i] += alpha * x[i];
};
AxpyOp<double> daxpy( ... );
Node::parallel for(0,N,daxpy);
AxpyOp<complex<float> > caxpy( ... );
Node::parallel for(0,N,caxpy);
```



Kokkos parallel_reduce example

• Consider real-valued vector inner product: $\alpha = x^T y$

```
template <class Scalar>
struct DotOp {
  const Scalar *x, *y;
  typedef Scalar ReductionType;
  Scalar identity() { return 0; }
  Scalar generate(int i) {
    return x[i]*y[i];
  Scalar reduce(Scalar a, Scalar b) {
    return a+b;
};
DotOp<float> fdot( ... );
float f = Node::parallel reduce(0,N,fdot);
DotOp<qd real> qddot( ... );
qd real q = Node::parallel reduce(0,N,qddot);
```

Tpetra RTI Operator Methods

 Tpetra Reduction/Transformation Interface provides convenience methods/macros for applying user Kokkos kernels to Tpetra Vectors/MultiVectors.

```
RCP< Tpetra::Map<LO,GO,Node> > domMap, rngMap, rowMap, colMap;
RCP< Tpetra::Import<LO,GO,Node> > importer = ...;
RCP< Tpetra::Export<LO,GO,Node> > exporter = ...;
MyKernel<T,LO> kern(...);
RCP< Tpetra::Operator<T,LO,GO,Node> > op;
op = Tpetra::RTI::kernelOp<T>(kern,domMap,rngMap,importer,exporter);
op->apply(x, y);
```

- Also wrappers for applying general functors.
 - e.g.: simple diagonal operator using a C++11 lambda function

```
RCP< Tpetra::Map<LO,GO,Node> > map;
RCP< Tpetra::Operator<T,LO,GO,Node> > op;
op = Tpetra::RTI::binaryOp<T>( [](T, T x) {return 2.0 * x;} , map );
op->apply(x, y);
```



Tpetra RTI Vector Methods

- Set of stand-alone non-member methods, e.g.:
 - unary_transform<UOP>(Vector &v, UOP op)
 - binary_transform<BOP>(Vector &v1, const Vector &v2, BOP op)
 - reduce<G>(const Vector &v1, const Vector &v2, G op_glob)
- This level provides maximal expressiveness, but coarser levels brings convenience.



Easy parallel algorithm development

- Inline templated hybrid-parallel conjugate gradient.
 - Fun game: Find the MPI or threading!

```
for (k=0; k<numIters; ++k) {</pre>
                                             //Ap = A*p
 A->apply(*p,*Ap);
  S pAp = TPETRA REDUCE2(
            p, Ap,
             p*Ap, ZeroOp<S>, plus<S>() ); // p'*Ap
                                             // alpha = r'*r/p'*Ap
 const S alpha = rr / pAp;
  TPETRA BINARY TRANSFORM
             x, p,
                                             // x = x + alpha*p
             x + alpha*p);
  S rrold = rr;
  rr = TPETRA BINARY PRETRANSFORM REDUCE (
             r, Ap,
                                             // fused kernels
                                             // r - alpha*Ap
             r - alpha*Ap,
             r*r, ZeroOp<S>, plus<S>() ); // sum r'*r
  const S beta = rr / rrold;
                                             // beta = r'*r/old(r'*r)
  TPETRA BINARY TRANSFORM
            p, r,
                                             // p = z + beta*p
             r + beta*p);
```

Conclusion

- Tpetra and Epetra use cases are very similar.
- Capabilities have diverged somewhat, although many of the same issues exist regarding programming model.
- Understanding these classes is critical to leveraging most Trilinos solver libraries.
- Incorporation of these classes has application value independent of downstream package use.
- Still experimenting with programming models for multi-core platforms.

