

Chapter 5:

Logical Database Design and the Relational Model

Modern Database Management
8th Edition

Jeffrey A. Hoffer, Mary B. Prescott,
Fred R. McFadden

Objectives

- Definition of terms
- List five properties of relations
- State two properties of candidate keys
- Define first, second, and third normal form
- Describe problems from merging relations
- Transform E-R and EER diagrams to relations
- Create tables with entity and relational integrity constraints
- Use normalization to convert anomalous tables to well-structured relations

is different from relationship
and it's equal to table

Relation

- **Definition:** A **relation** is a named, two-dimensional **table** of data
- **Table** consists of **rows** (records) and **columns** (attribute or field)
- **Requirements for a table to qualify as a relation:**
 - It must have a **unique name**
 - Every **attribute value** must be **atomic** (not multivalued, not composite)
 - Every **row** must be **unique** (can't have two rows with exactly the same values for all their fields)
 - Attributes (**columns**) in tables must have **unique names**
 - The **order** of the **columns** must be **irrelevant**
 - The **order** of the **rows** must be **irrelevant**

NOTE: all *relations* are in ***1st Normal form***

Example of Relation

Figure 5-1 EMPLOYEE1 relation with sample data

EMPLOYEE1

Emp_ID	Name	Dept_Name	Salary
100	Margaret Simpson	Marketing	48,000
140	Allen Beeton	Accounting	52,000
110	Chris Lucero	Info Systems	43,000
190	Lorenzo Davis	Finance	55,000
150	Susan Martin	Marketing	42,000

we can express the structure of the employee relation as:

Employee1 (Emp_ID, Name, Dept_Name, Salary)

foreign key is dash under lined also

Correspondence with E-R Model

- Relations (**tables**) correspond with **entity types** and with many-to-many relationship types
 - Rows correspond with entity **instances** and with many-to-many relationship instances
 - Columns correspond with **attributes**
-
- **NOTE:** The word ***relation*** (in relational database) is **NOT** the same as the word ***relationship*** (in E-R model)

Key Fields



- **Keys** are special fields that serve two main **purposes**:
 - **Primary keys** are unique identifiers of the relation in question.
- **Keys** can be **simple** (a single field) or **composite** (more than one field)
 - *This is how we can guarantee that all rows are unique*
 - **Examples:**

Simple Attribute:

Employee(Emp_ID, Name, Dept_Name, Salary)

Composite Key:

Order_Line(Order_No, Product_No, Quantity)

Key Fields



- **Foreign keys** are identifiers that enable a dependent relation (on the many side of a relationship) to refer to its parent relation (on the one side of the relationship)

- Example:

Relationship between Employee and Department Tables

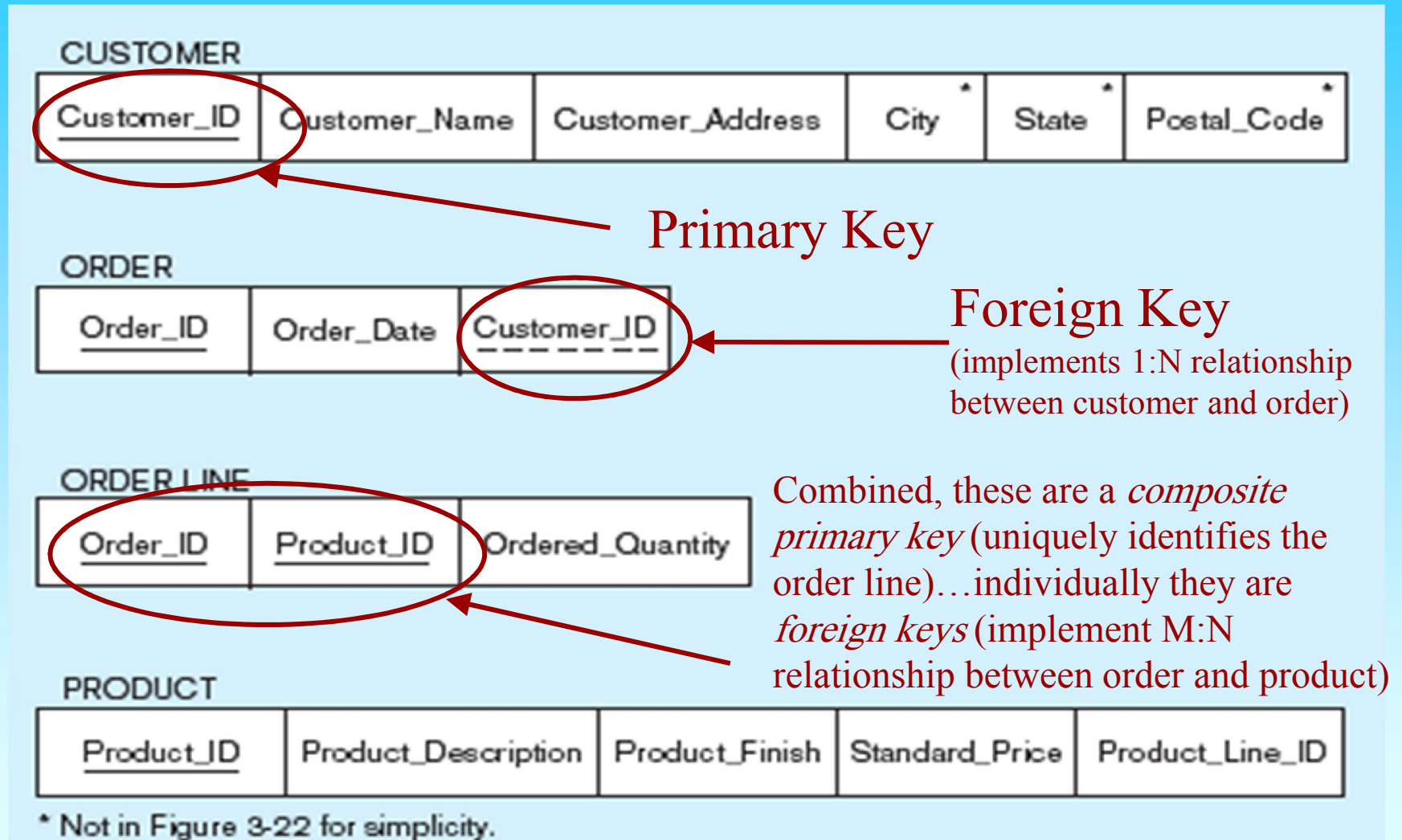
Employee(Emp_ID, Name, Dept_No, Salary)

Department(Dept_No, Location, Fax)

- Keys usually are used as indexes to speed up the response to user queries (More on this in Ch. 6)

important shape / tables / schema of relation

Figure 5-3 Schema for four relations (Pine Valley Furniture Company)



alternate key: when two attributes can be used as primary key and we can use any of them one of them is primary key and the other one is the alternate key the two are called candidate ID

Integrity Constraints

- Relational model contains several types of constraints (**Business Rules**) to facilitate maintaining the accuracy and Integrity of data , such as:
- **Domain Constraints**
 - A **domain** is the set of values assigned to an attribute
 - **Domain definition** consists of:
 - Domain **name**
 - Domain **meaning**
 - data **type**
 - Data **size**
 - Allowable **values**

Table 5-1 Domain Definitions for INVOICE Attributes

<i>Attribute</i>	<i>Domain Name</i>	<i>Description</i>	<i>Domain</i>
Customer_ID	Customer_IDs	Set of all possible customer IDs	character: size 5
Customer_Name	Customer_Names	Set of all possible customer names	character: size 25
Customer_Address	Customer_Addresses	Set of all possible customer addresses	character: size 30
City	Cities	Set of all possible cities	character: size 20
State	States	Set of all possible states	character: size 2
Postal_Code	Postal_Codes	Set of all possible postal zip codes	character: size 10
Order_ID	Order_IDs	Set of all possible order IDs	character: size 5
Order_Date	Order_Dates	Set of all possible order dates	date format mm/dd/yy
Product_ID	Product_IDs	Set of all possible product IDs	character: size 5
Product_Description	Product_Descriptions	Set of all possible product descriptions	character size 25
Product_Finish	Product_Finishes	Set of all possible product finishes	character: size 15
Standard_Price	Unit_Prices	Set of all possible unit prices	monetary: 6 digits
Product_Line_ID	Product_Line_IDs	Set of all possible product line IDs	integer: 3 digits
Ordered_Quantity	Quantities	Set of all possible ordered quantities	integer: 3 digits

Domain definitions enforce domain integrity constraints

Integrity Constraints

- **Entity Integrity** important
 - Each **relation** must have a **primary key**
 - No primary key attribute may be **null**.
 - All **primary key** fields **MUST** have **data**
 - Null value may be assigned to **non-key** attribute when no other **value applies** or when the **applicable value** is unknown
- **Action Assertions**
 - **Business rules**. Recall from Ch. 4

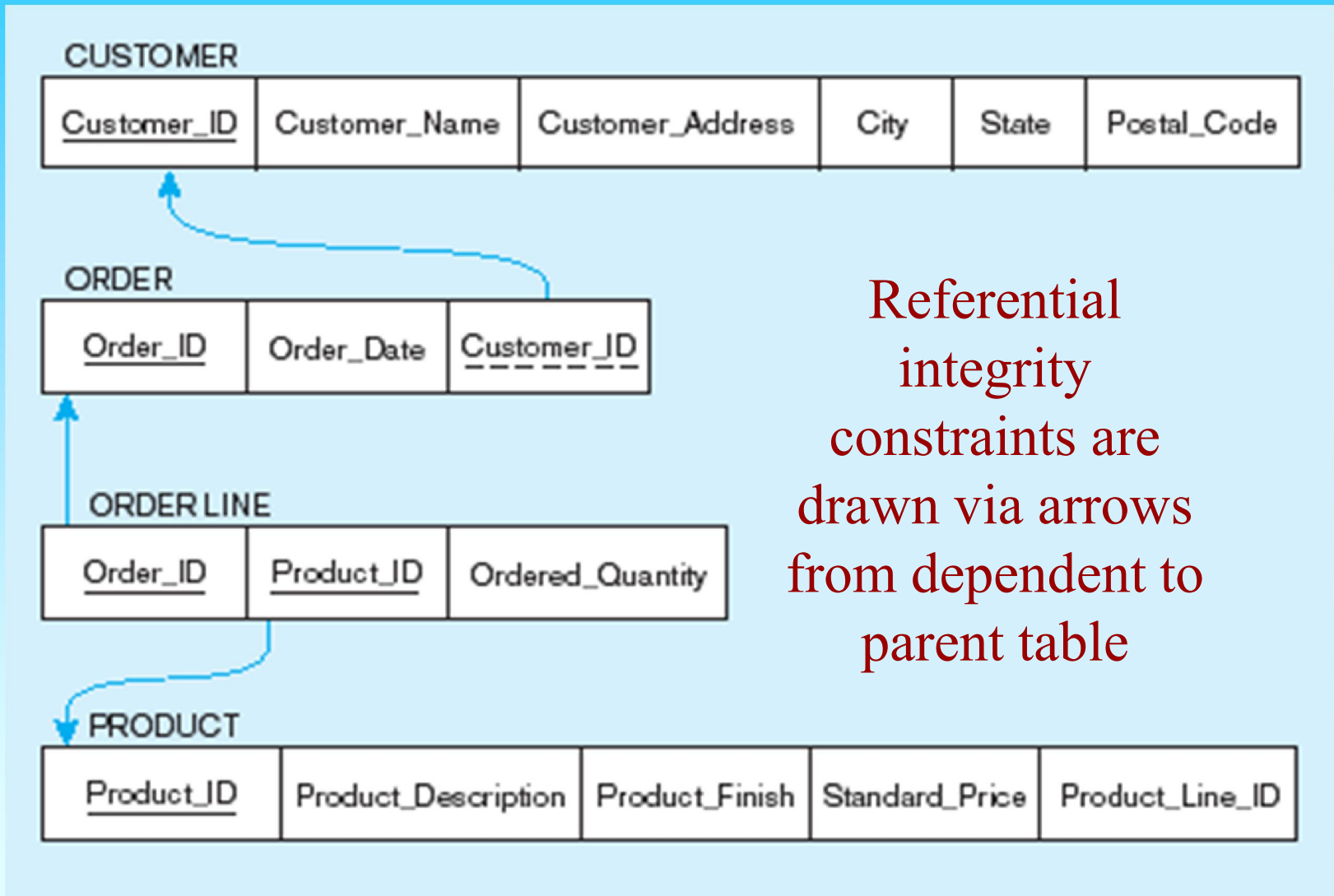
Integrity Constraints

■ Referential Integrity

- A rule states that any foreign key value (on the relation of the many side) **MUST match** a primary key value in the relation of the one side. (Or the foreign key can be null)
- For example: **Delete Rules**
 - **Restrict** – don't allow delete of "parent" side if related rows exist in "dependent" side
 - **Cascade** – automatically delete "dependent" side rows that correspond with the "parent" side row to be deleted
 - **Set-to-Null** – set the foreign key in the dependent side to null if deleting from the parent side → **not allowed for weak entities**

Figure 5-5

Referential integrity constraints (Pine Valley Furniture)



Creating Relational Tables

SOL table definitions

```
CREATE TABLE CUSTOMER
(CUSTOMER_ID          VARCHAR(5)          NOT NULL,
CUSTOMER_NAME         VARCHAR(25)         NOT NULL,
CUSTOMER ADDRESS      VARCHAR(30)         NOT NULL,
CITY                  VARCHAR(20)         NOT NULL,
STATE                 CHAR(2)              NOT NULL,
POSTAL_CODE           CHAR(10)            NOT NULL,
PRIMARY KEY (CUSTOMER_ID);

CREATE TABLE ORDER
(ORDER_ID             CHAR(5)              NOT NULL,
ORDER DATE            DATE                 NOT NULL,
CUSTOMER_ID           VARCHAR(5)          NOT NULL,
PRIMARY KEY (ORDER_ID),
FOREIGN KEY (CUSTOMER_ID) REFERENCES CUSTOMER (CUSTOMER_ID);

CREATE TABLE ORDER_LINE
(ORDER_ID             CHAR(5)              NOT NULL,
PRODUCT_ID            CHAR(5)              NOT NULL,
ORDERED_QUANTITY      INT                  NOT NULL,
PRIMARY KEY (ORDER_ID, PRODUCT_ID),
FOREIGN KEY (ORDER_ID) REFERENCES ORDER (ORDER_ID),
FOREIGN KEY (PRODUCT_ID) REFERENCES PRODUCT (PRODUCT_ID);

CREATE TABLE PRODUCT
(PRODUCT_ID           CHAR(5)              NOT NULL,
PRODUCT_DESCRIPTION    VARCHAR(25),
PRODUCT_FINISH        VARCHAR(12),
STANDARD_PRICE        DECIMAL(8,2)        NOT NULL,
PRODUCT_LINE_ID       INT                  NOT NULL,
PRIMARY KEY (PRODUCT_ID);
```

Referential
integrity
constraints are
implemented
with foreign key
to primary key
references

Transforming EER Diagrams into Relations until here

Transforming (**mapping**) ER-Diagrams to Relations is straightforward with a well-defined set of rules.

Types of Entities

1. **Regular** Entities: have an **independent** existence such as persons, products,
2. **Weak** Entities: **dependent** and cannot exist without identifying relationship with an owner (Regular Entity)
3. **Associative** Entities: formed from **many-to-many** relationships between other entities

Transforming EER Diagrams into Relations

Mapping Regular Entities to Relations

1. **Simple** attributes: E-R attributes map directly onto the relation
2. **Composite** attributes: Use only their simple, component attributes
3. **Multivalued** Attribute—Becomes a separate relation with a foreign key taken from the superior entity

Figure 5-8 Mapping a regular entity

**(a) CUSTOMER
entity type with
simple
attributes**



(b) CUSTOMER relation

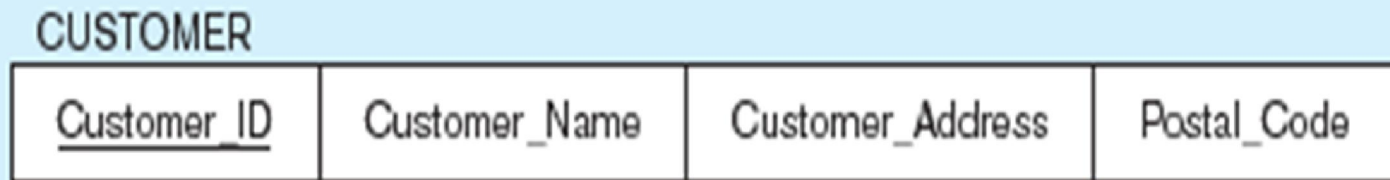


Figure 5-9 Mapping a composite attribute

(a) CUSTOMER
entity type with
composite
attribute



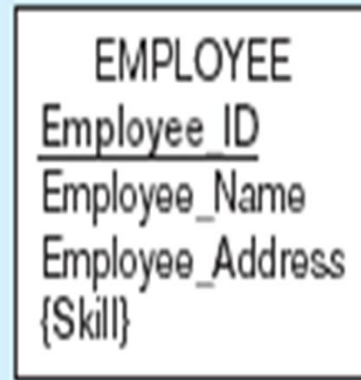
CUSTOMER

(b) CUSTOMER relation with address detail

<u>Customer_ID</u>	Customer_Name	Street	City	State	Postal_Code
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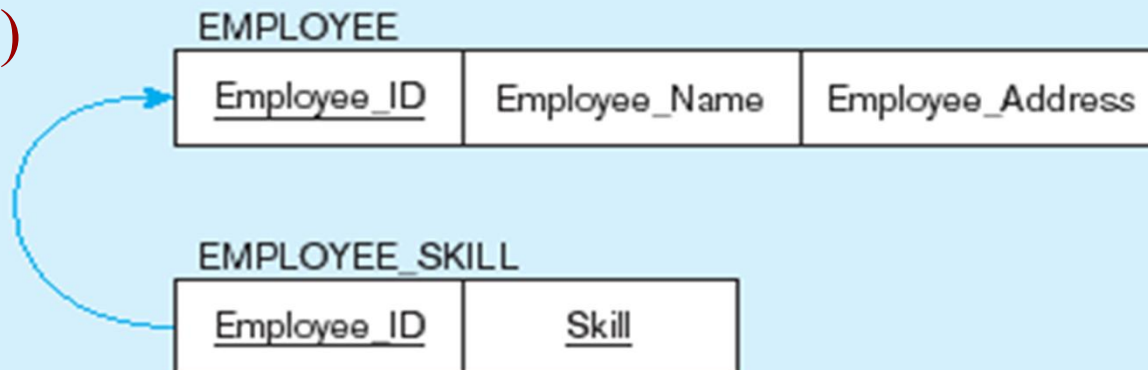
Figure 5-10 Mapping an entity with a multivalued attribute

(a)



Multivalued attribute becomes a separate relation with foreign key

(b)



One-to-many relationship between original entity and new relation

Transforming EER Diagrams into Relations (cont.)

Mapping Weak Entities

- Weak entity cannot exist independently
- It does not have a complete Identifier
- Becomes a separate relation with a foreign key taken from the superior entity
- Primary key composed of:
 - Partial identifier of weak entity
 - Primary key of identifying relation (strong entity)

Figure 5-11 Example of mapping a weak entity

a) **Weak entity** DEPENDENT

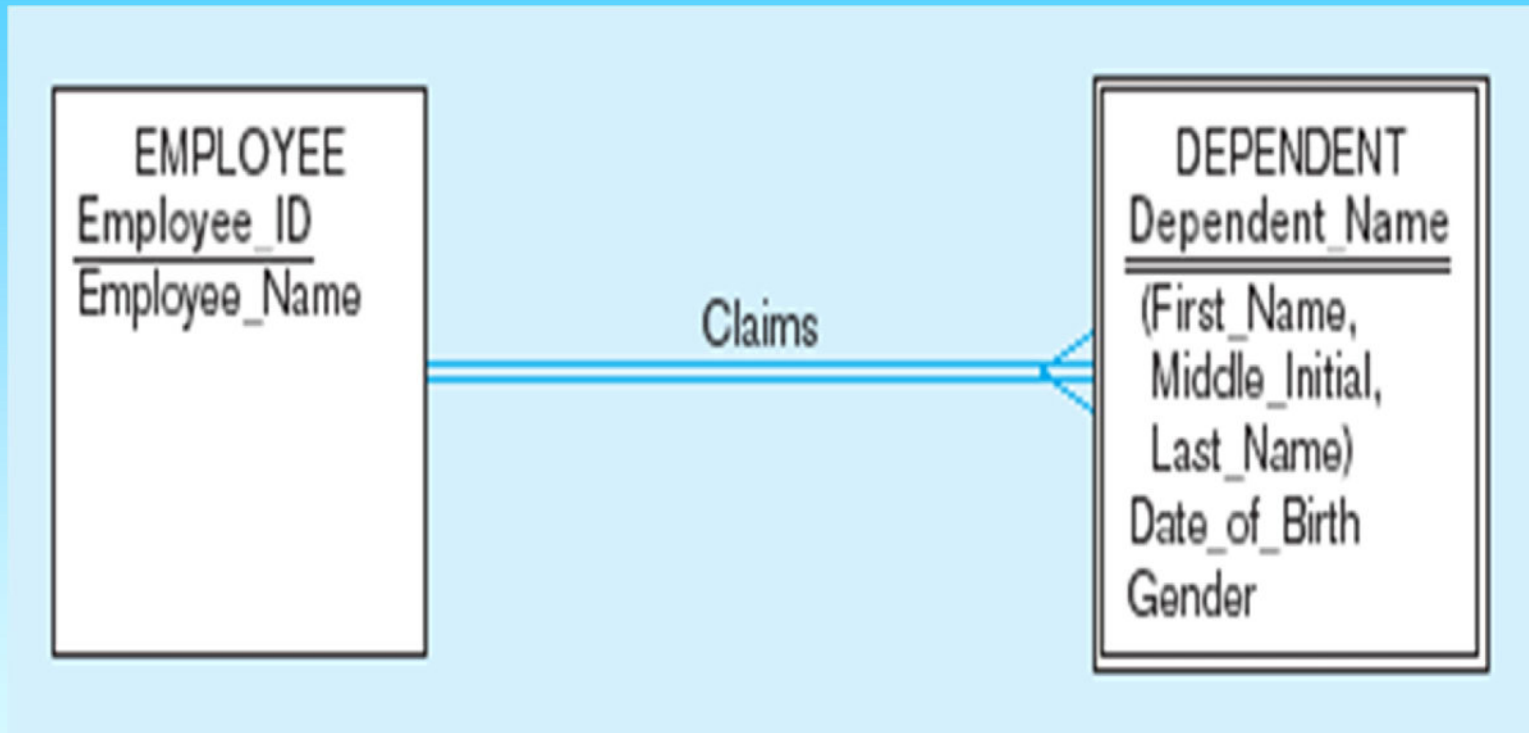
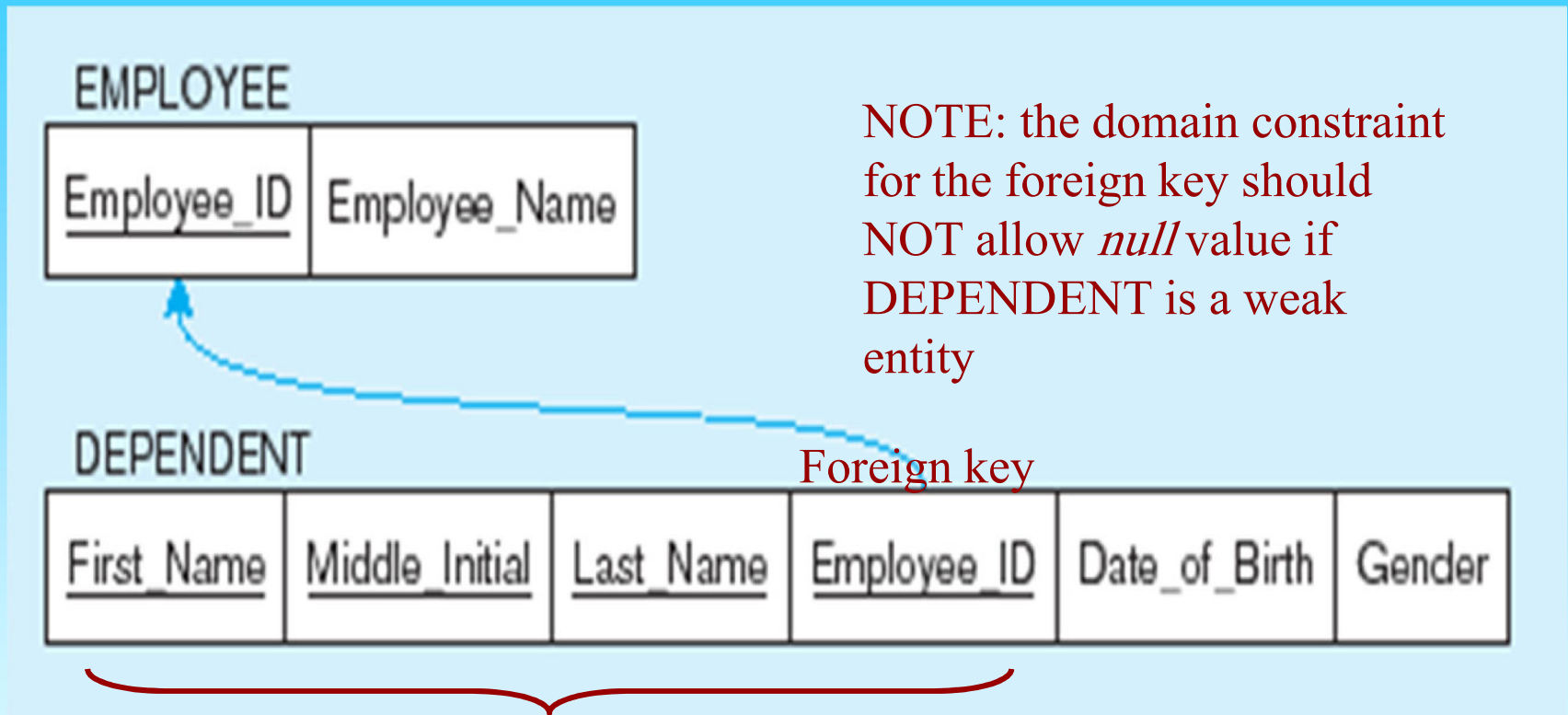


Figure 5-11 Example of mapping a weak entity (cont.)

b) Relations resulting from weak entity



Composite primary key

Alternate approach:

Dependent(Employee_ID, Dependent#, First_Name, Middle_Name, Last_Name, Date_of_Birth, Gender)

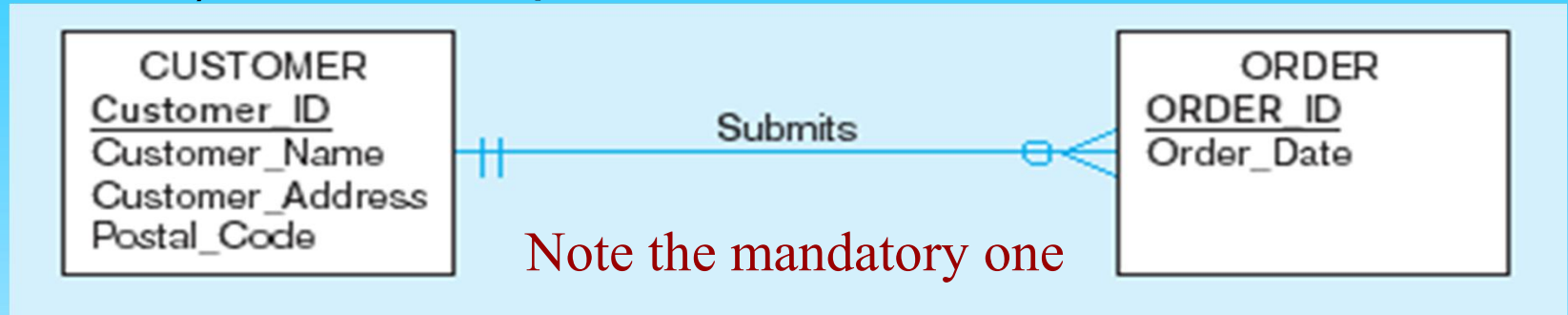
Transforming EER Diagrams into Relations (cont.)

Mapping **Binary** Relationships

- **One-to-Many** – **Primary** key on the **one** side becomes a **foreign** key on the **many** side
- **Many-to-Many** – Create a ***new relation*** with the primary keys of the two entities as its primary key
- **One-to-One** – **Primary** key on the **mandatory** side becomes a **foreign** key on the **optional** side

Figure 5-12 Example of mapping a 1:M relationship

a) Relationship between customers and orders



b) Mapping the relationship

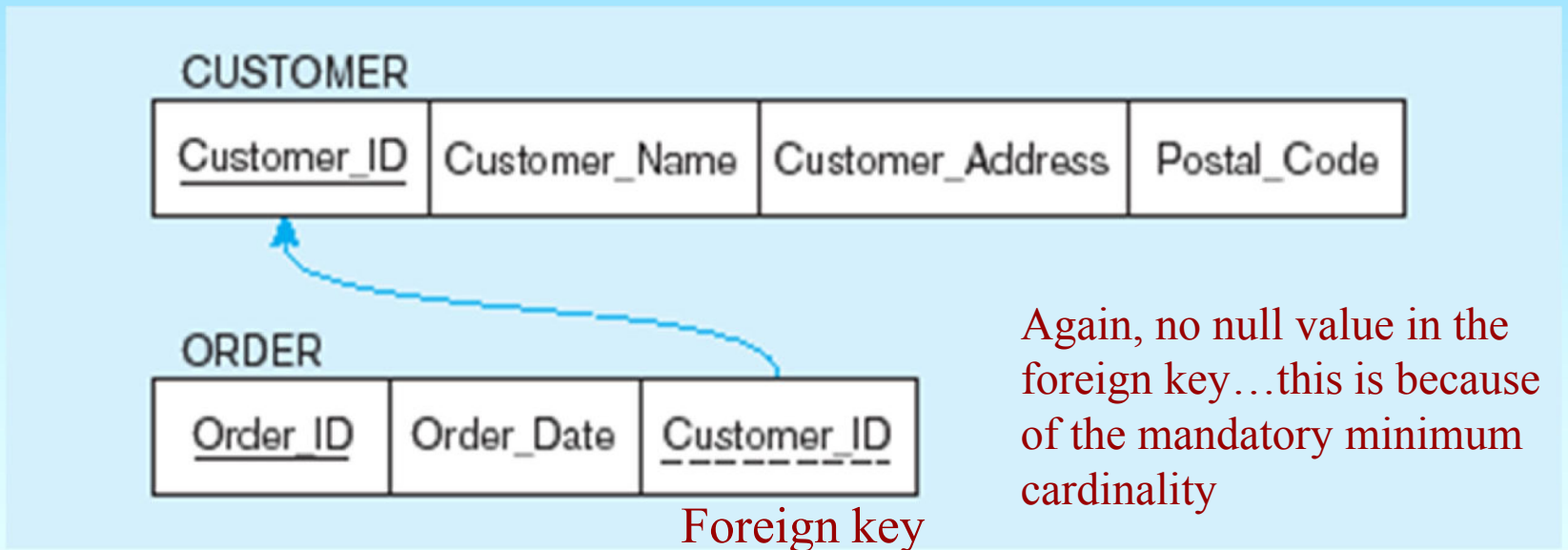
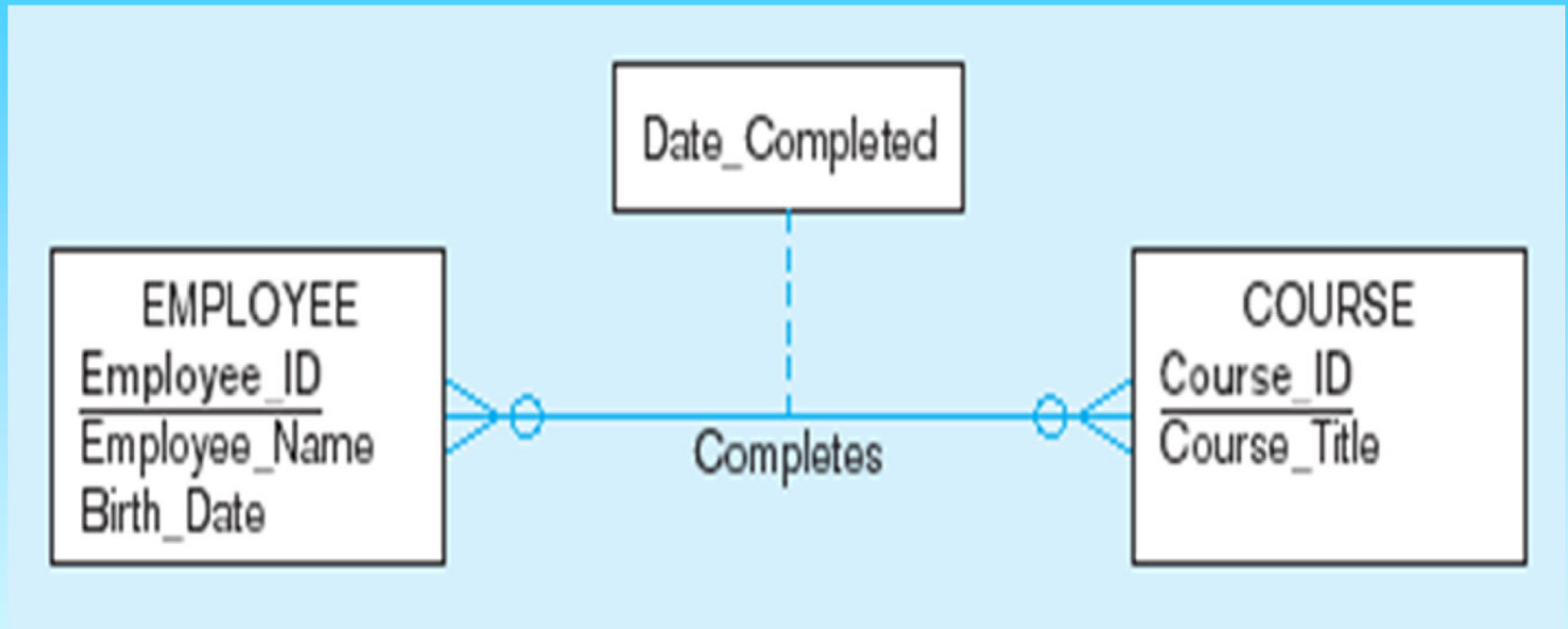


Figure 5-13 Example of mapping an **M:N** relationship

a) Completes relationship (M:N)



The *Completes* relationship will need to become a separate relation

Figure 5-13 Example of mapping an M:N relationship (cont.)

b) **Three** resulting relations

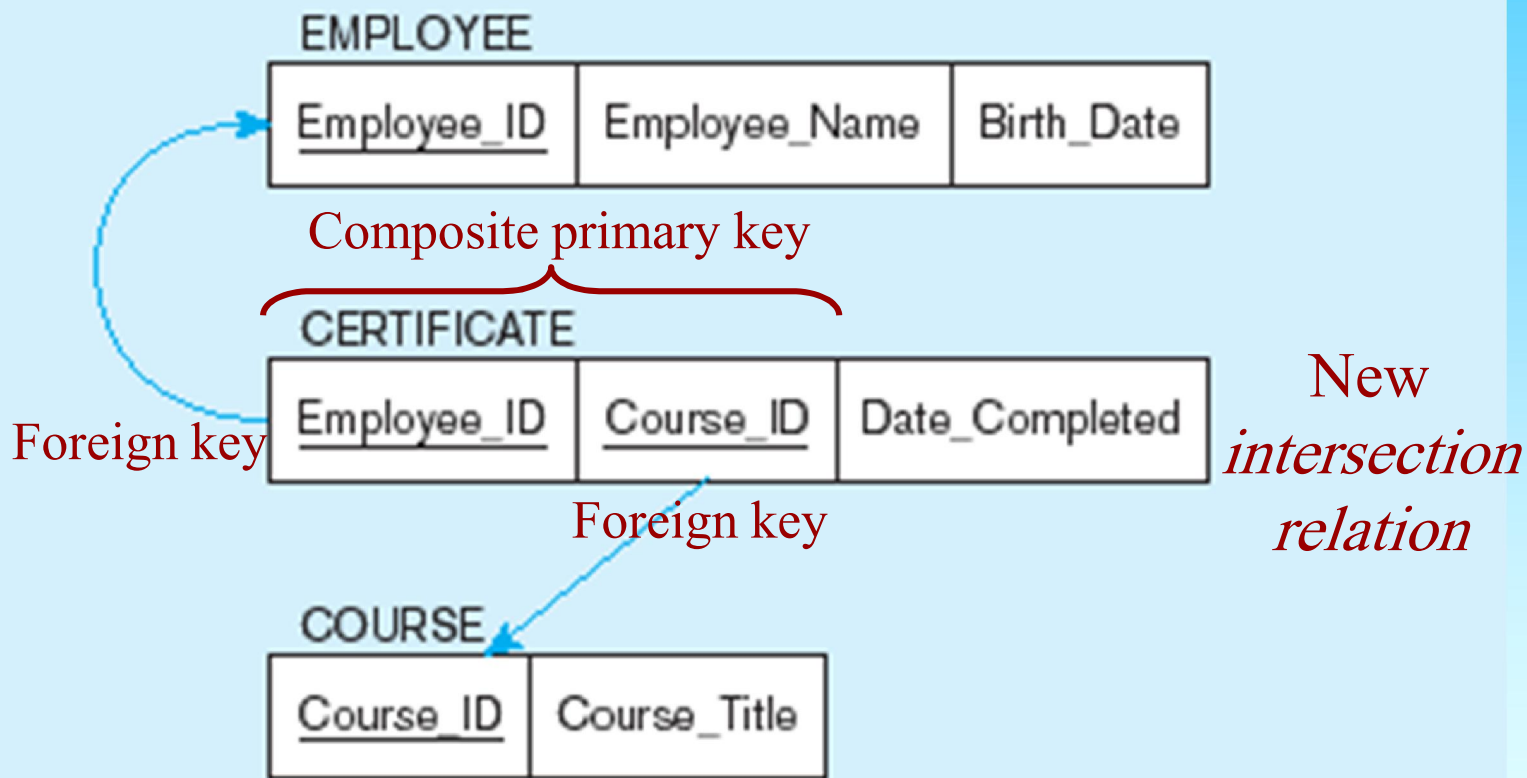
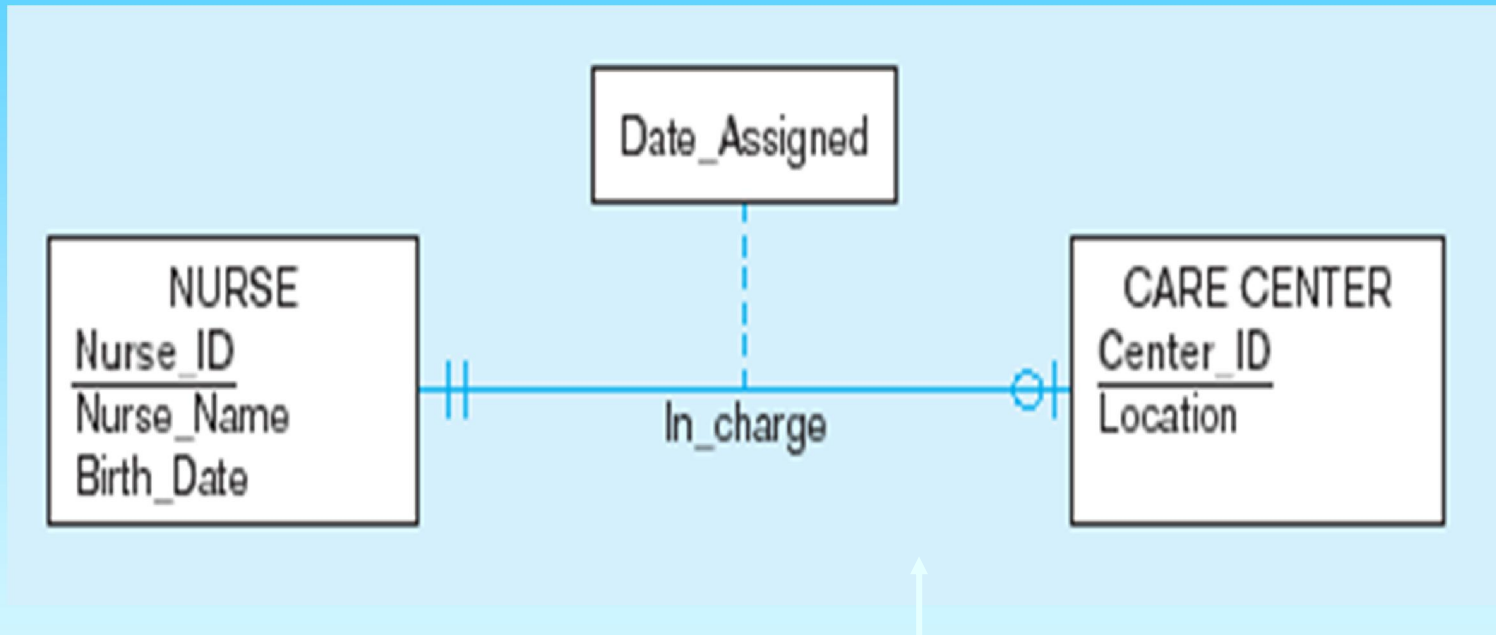


Figure 5-14 Example of mapping a binary 1:1 relationship

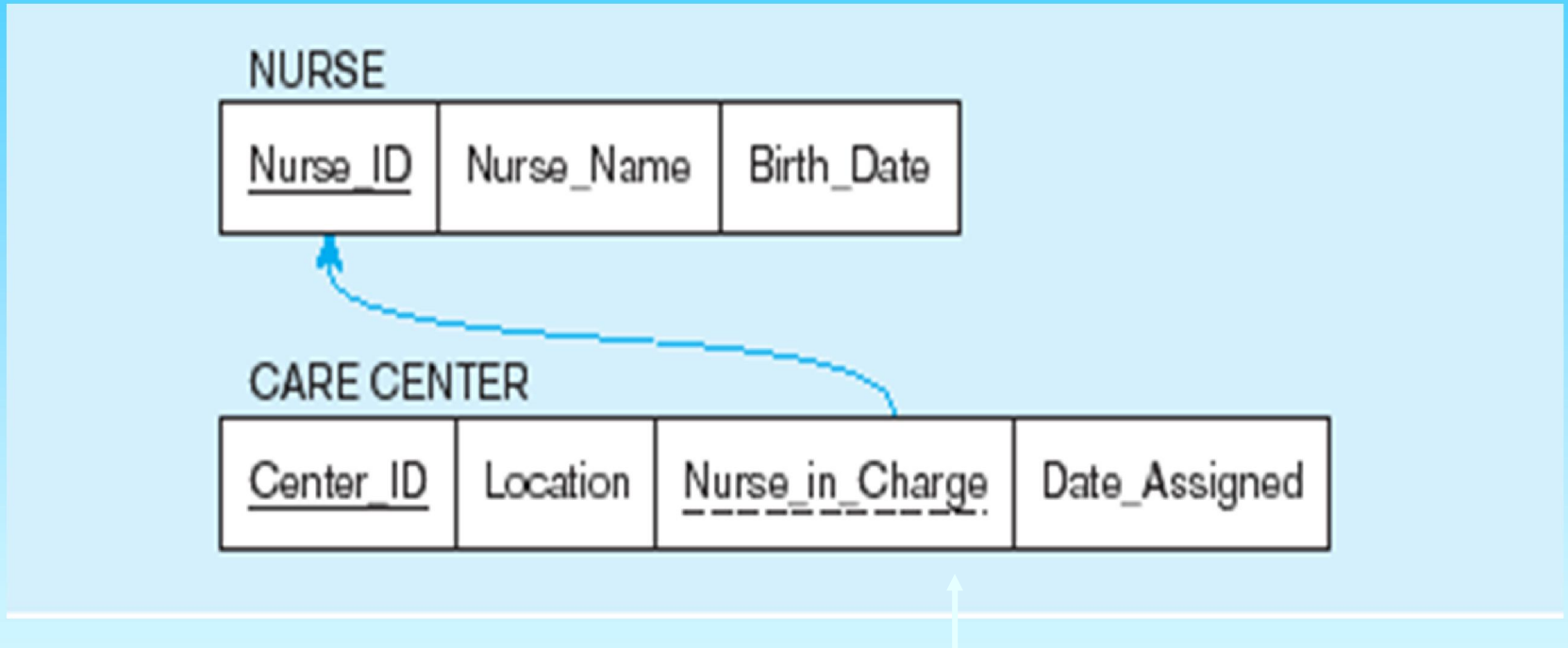
a) In_charge relationship (1:1)



Often in 1:1 relationships, one direction is optional.

Figure 5-14 Example of mapping a binary 1:1 relationship (cont.)

b) Resulting relations



Foreign key goes in the relation on the optional side,
Matching the primary key on the mandatory side

Transforming EER Diagrams into Relations (cont.)

Mapping **Associative** Entities

- Identifier **Not** Assigned
 - **Default** primary key for the association relation is **composed** of the primary keys of the **two** entities (as in M:N relationship)
- Identifier **Assigned**
 - It is **natural** and familiar to end-users
 - Default identifier **may not** be **unique**

Figure 5-15 Example of mapping an **associative** entity

a) An associative entity

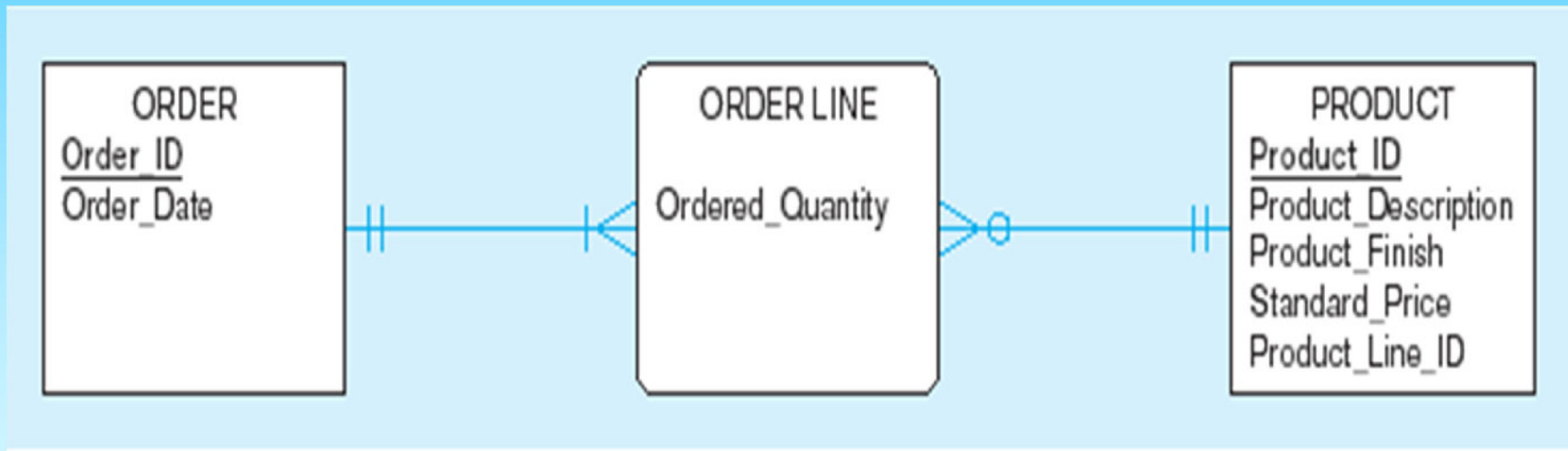


Figure 5-15 Example of mapping an associative entity (cont.)

b) **Three** resulting relations

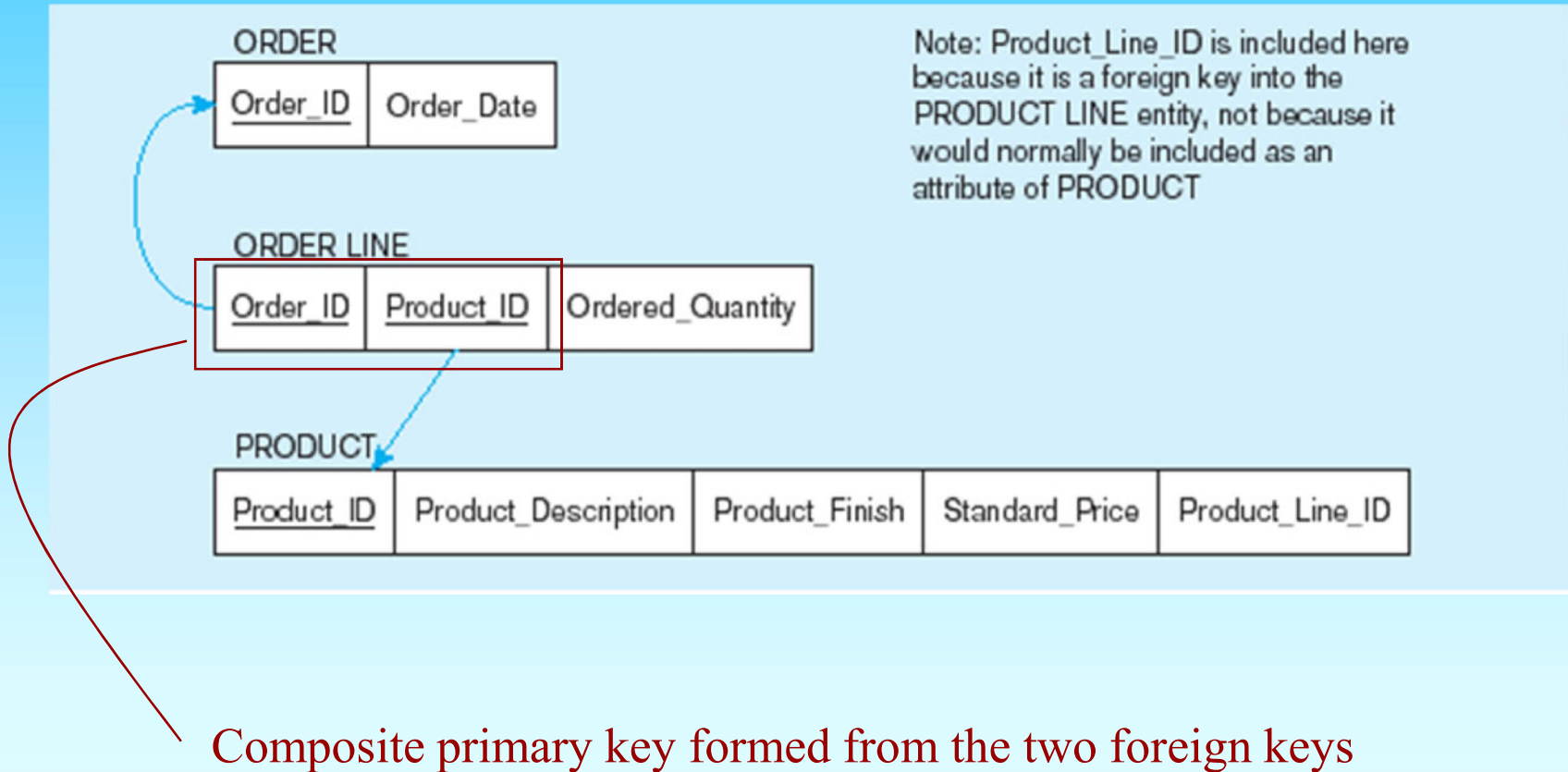


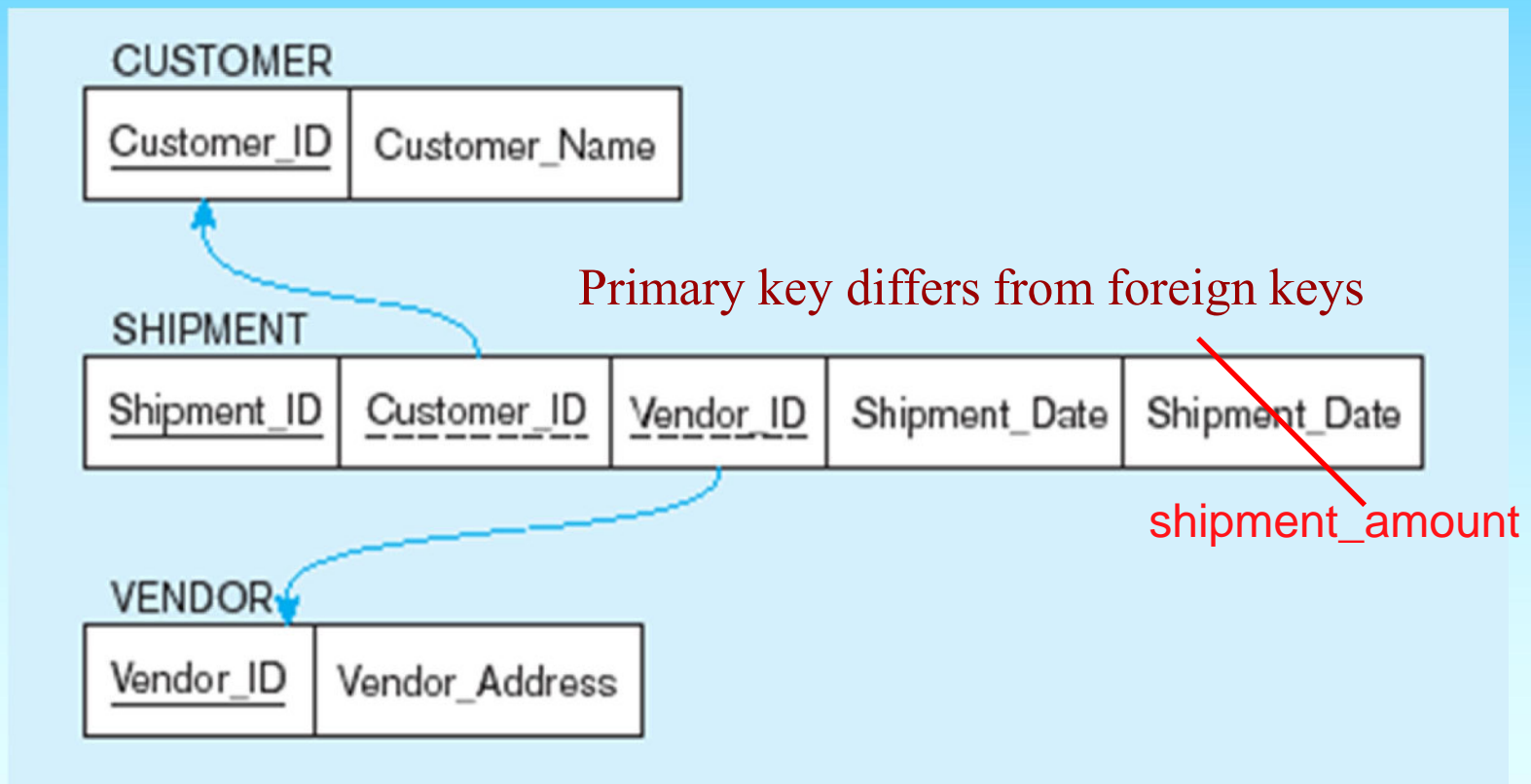
Figure 5-16 Example of mapping an **associative** entity with an **identifier**

a) SHIPMENT associative entity



Figure 5-16 Example of mapping an associative entity with an identifier (cont.)

b) **Three** resulting relations

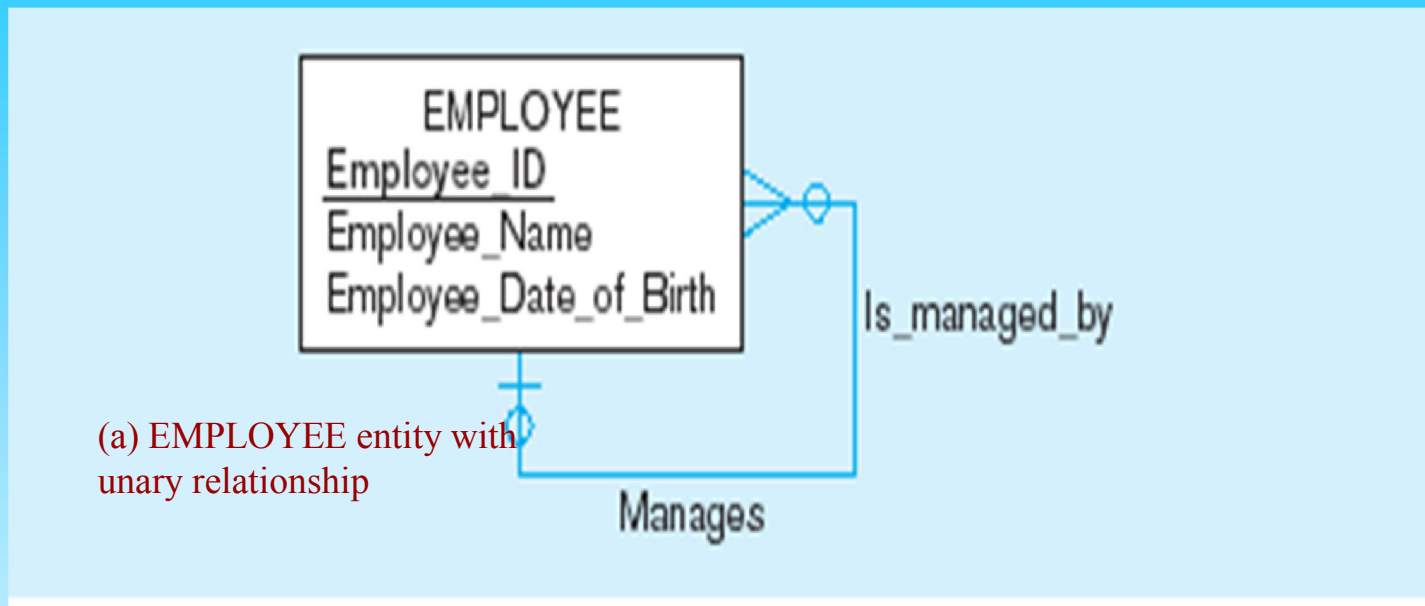


Transforming EER Diagrams into Relations (cont.)

Mapping **Unary** Relationships

- One-to-Many – **Recursive** foreign key in the same relation
- Many-to-Many – **Two** relations:
 - **One** for the **entity** type
 - **One** for an **associative** relation in which the **primary** key has **two** attributes, **both** taken from the **primary** key of the **entity**

Figure 5-17 Mapping a **unary 1:N** relationship



(b) EMPLOYEE relation with recursive foreign key

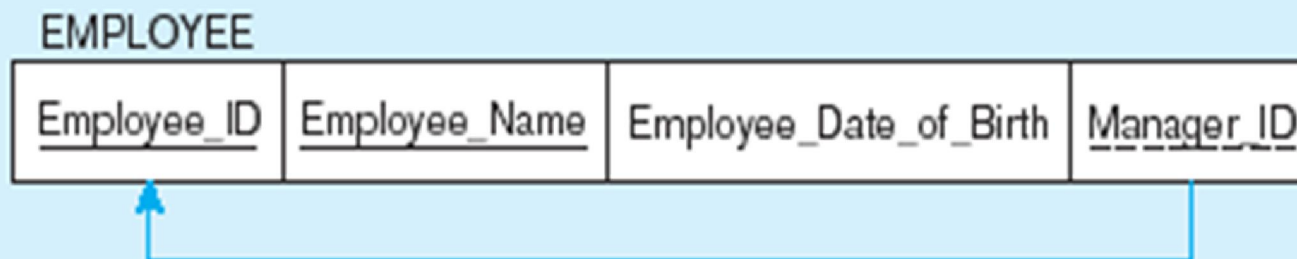
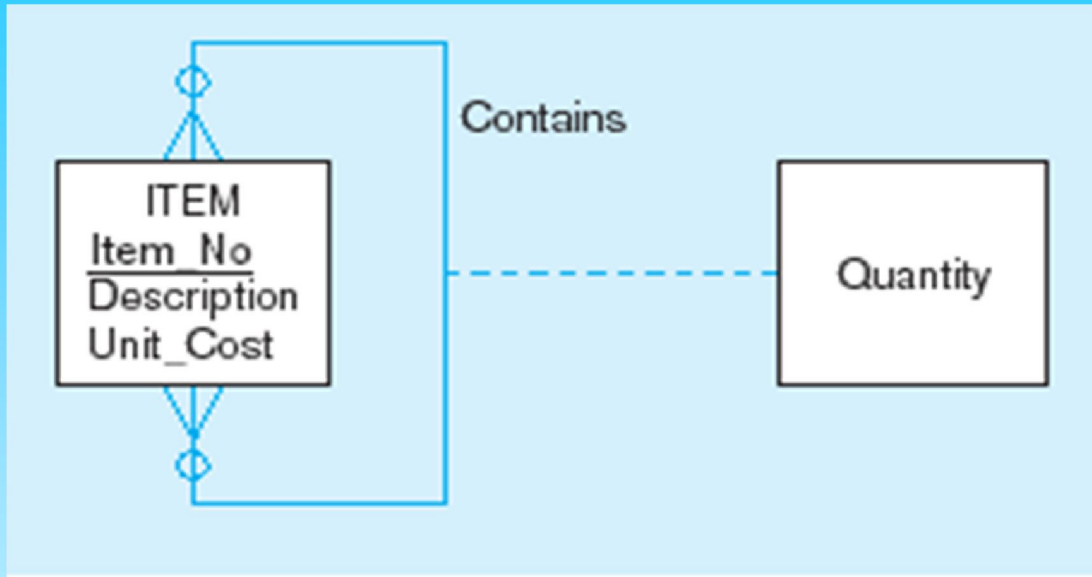


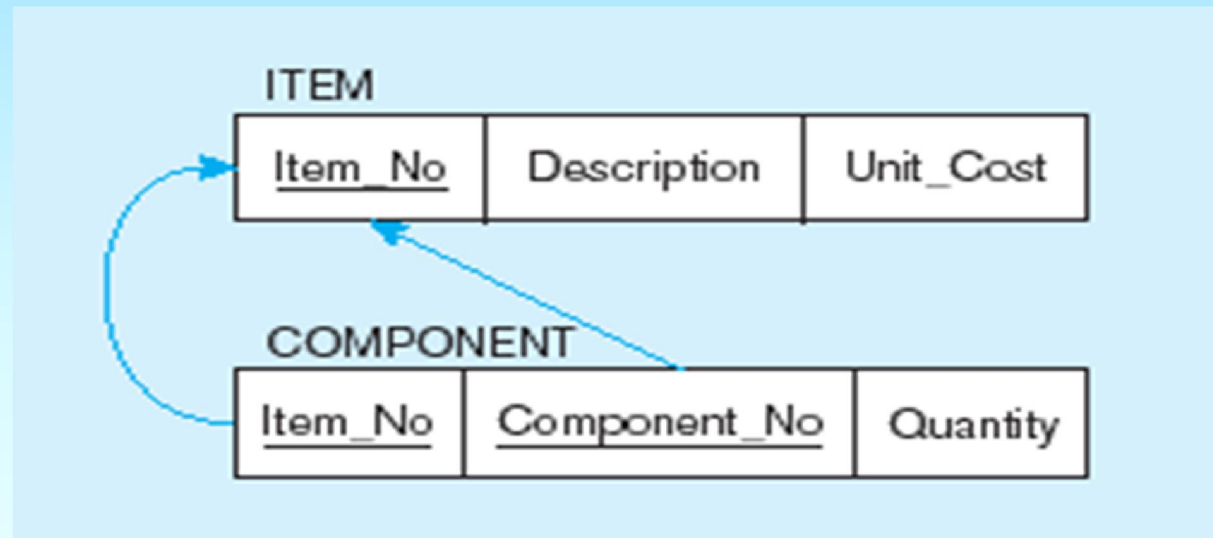
Figure 5-18 Mapping a **unary M:N** relationship



(a) Bill-of-materials relationships (M:N)

any item can be an item or component of items like in furniture

(b) ITEM and COMPONENT relations



Transforming EER Diagrams into Relations (cont.)

Mapping Ternary (and n-ary) Relationships

- One relation for each entity and one for the associative entity
- Associative entity has foreign keys to each entity in the relationship

Figure 5-19 Mapping a **ternary** relationship

a) PATIENT TREATMENT Ternary relationship with **associative** entity

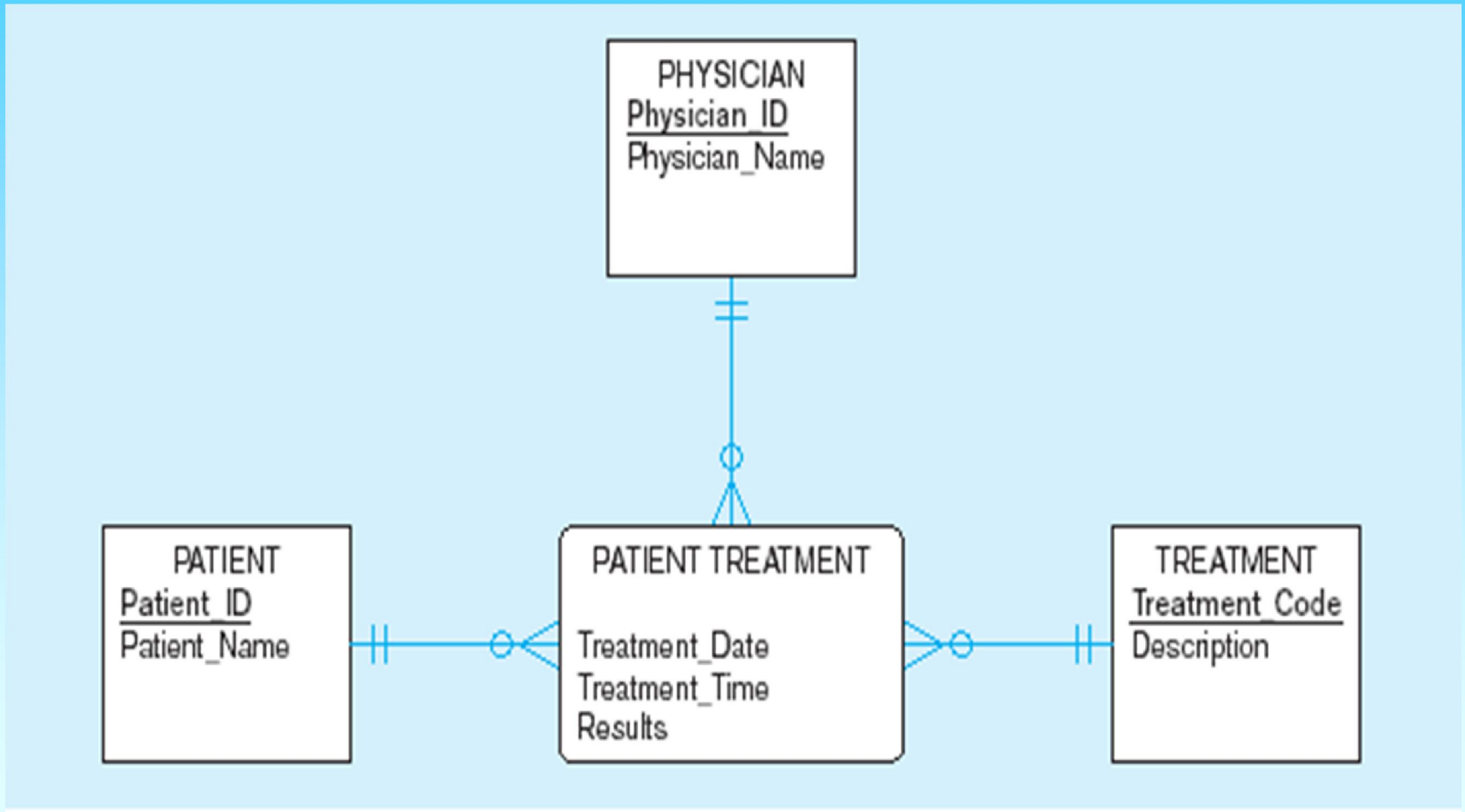
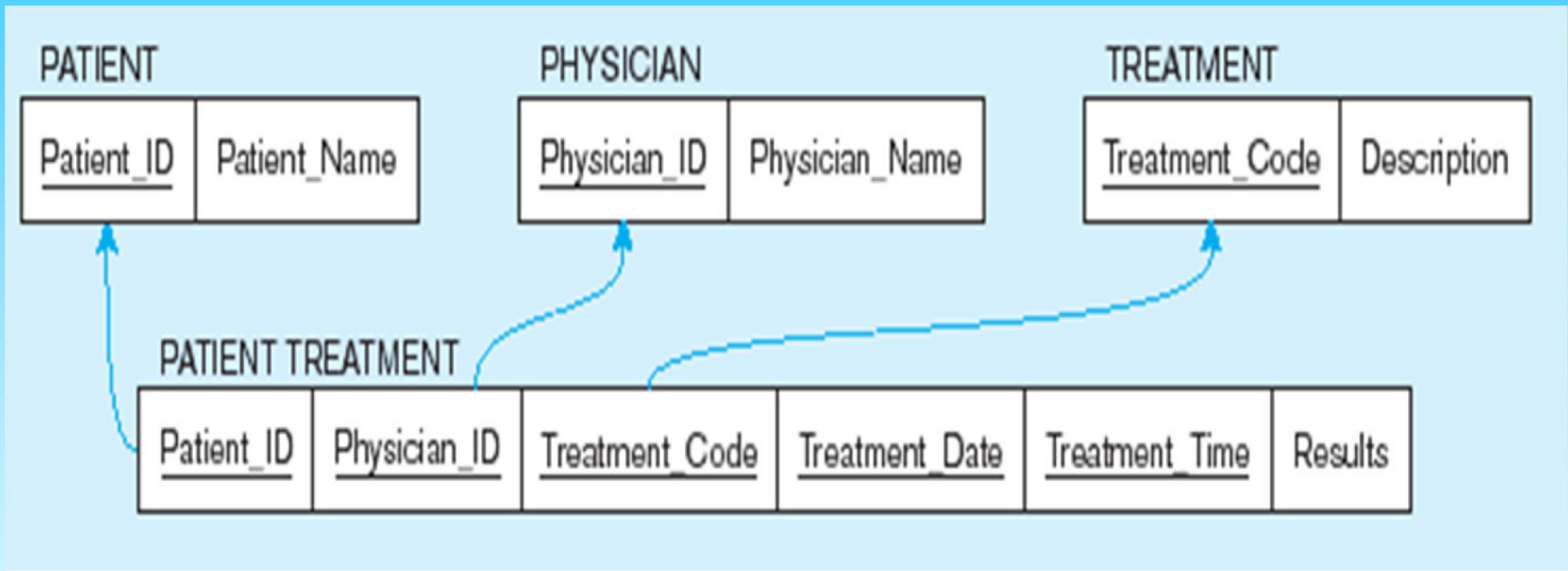


Figure 5-19 Mapping a **ternary** relationship (cont.)

b) Mapping the ternary relationship PATIENT TREATMENT



Remember
that the
primary key
MUST be
unique

This is why
treatment date
and time are
included in the
composite
primary key

But this makes a
very
cumbersome
key...

It would be
better to create a
Alternate key
like **Treatment#**

Transforming EER Diagrams into Relations (cont.)

Mapping Supertype/Subtype Relationships

- One relation for supertype and for each subtype
- Supertype attributes (including identifier and subtype discriminator) go into supertype relation
- Subtype attributes go into each subtype; primary key of supertype relation also becomes primary key of subtype relation
- 1:1 relationship established between supertype and each subtype, with supertype as primary table

Figure 5-20 **Supertype/subtype** relationships

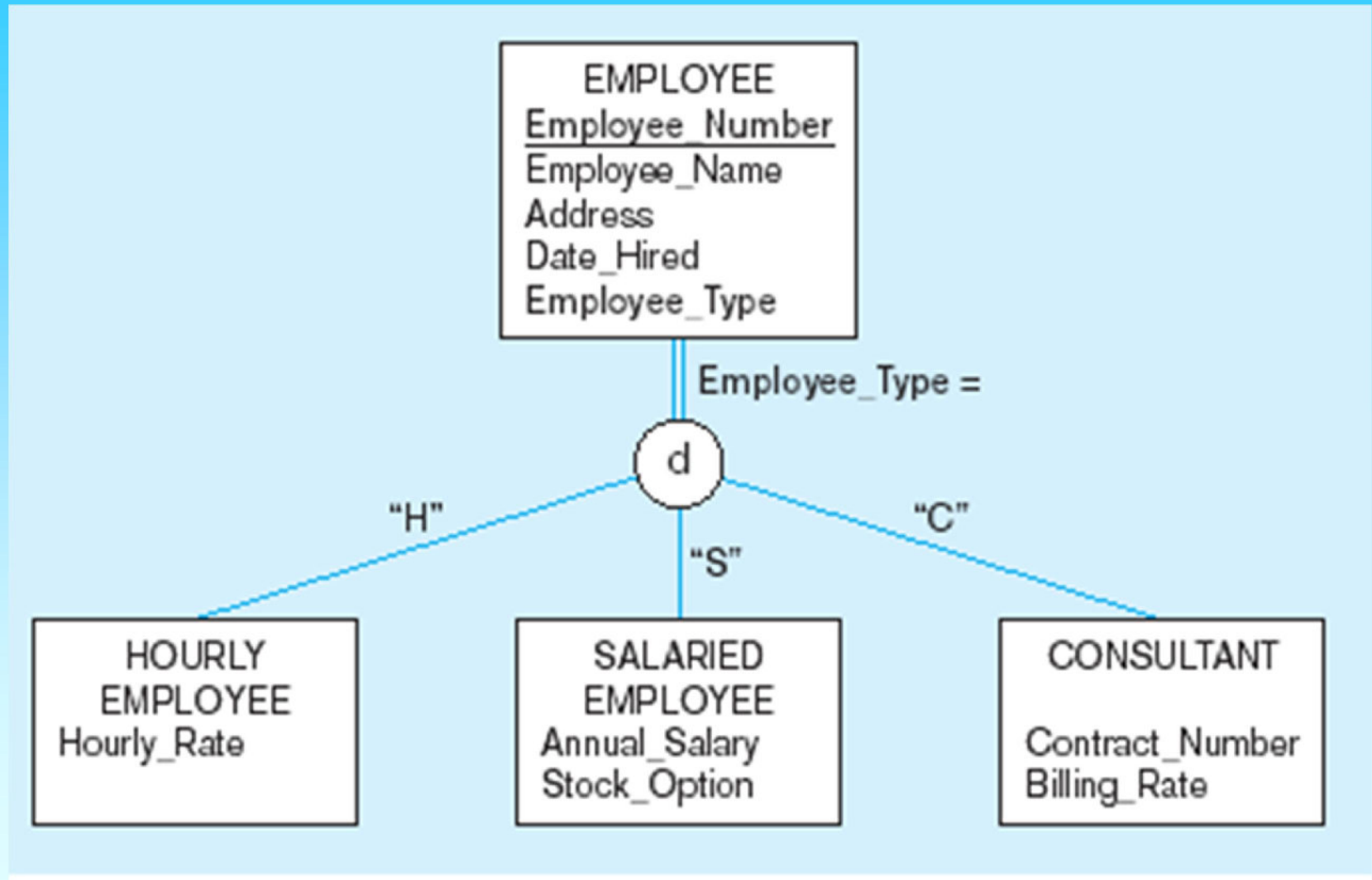
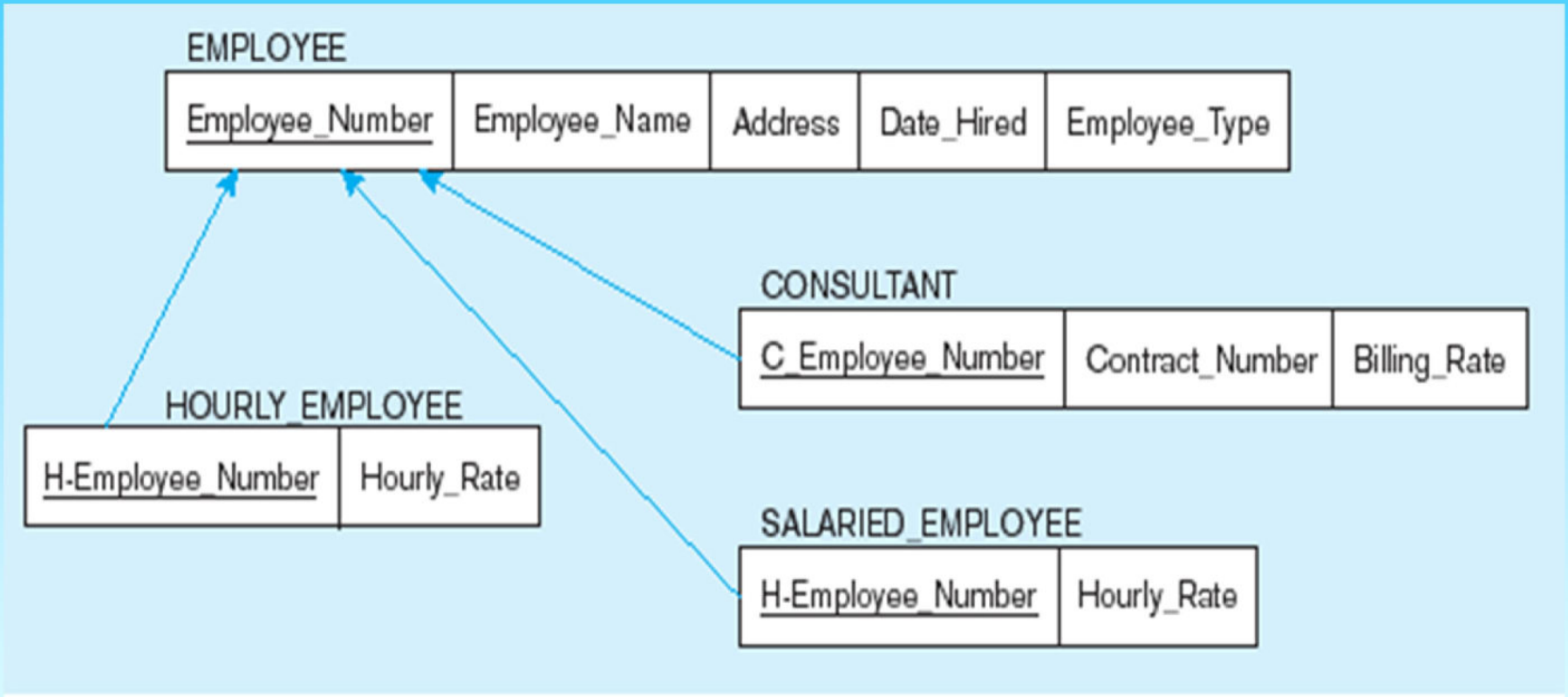


Figure 5-21

Mapping **Supertype/subtype** relationships to relations



These are implemented as **one-to-one** relationships