

# **Chapter 8**Switching

Figure 8.1 Switched network

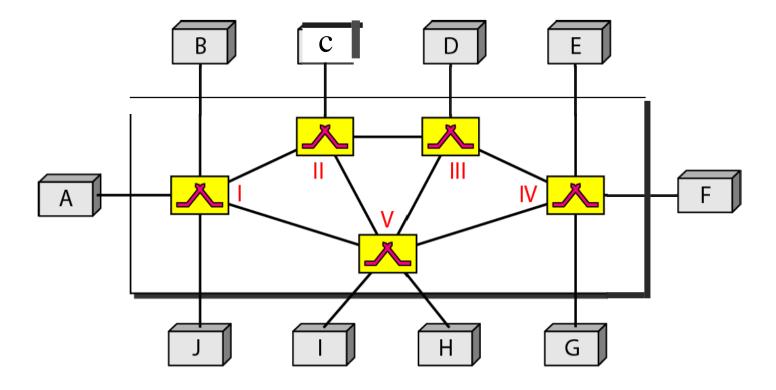
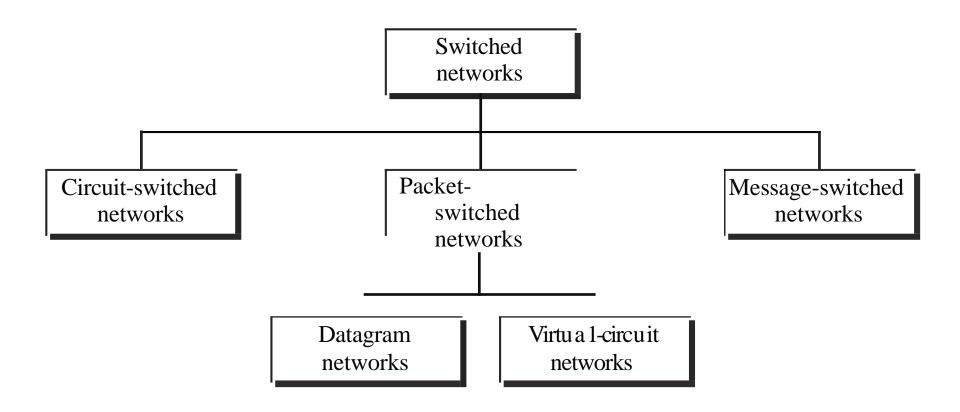


Figure 8.2 Taxonomy of switched networks



#### 8-1 CIRCUIT-SWITCHED NETWORKS

A circuit-switched network consists of a set of switches connected by physical links. A connection between two stations is a dedicated path made of one or more links. However, each connection uses only one dedicated channel on each link. Each link is normally divided into n channels by using FDM or TDM.

#### Topics discussed in this section:

**Three Phases** 

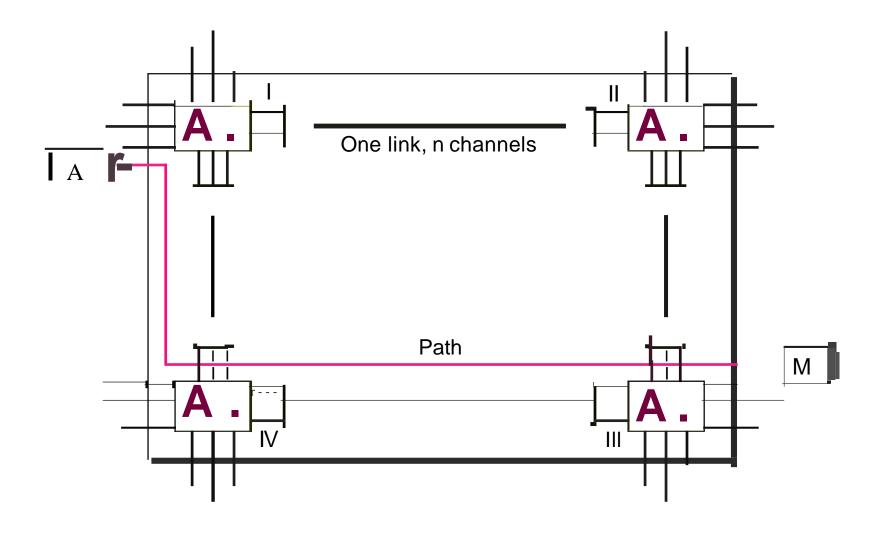
**Efficiency** 

**Delay** 

**Circuit-Switched Technology in Telephone Networks** 

A circuit-switched network is made of a set of switches connected by physical links, in which each link is divided into *n* channels.

Figure 8.3 A trivial circuit-switched network

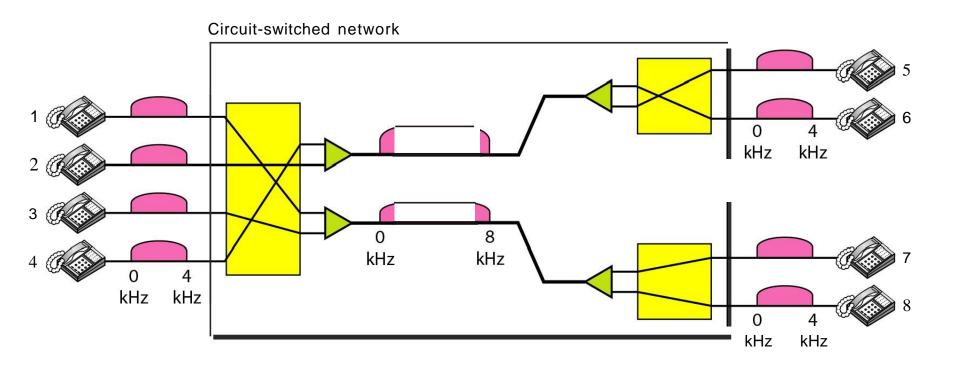


In circuit switching, the resources need to be reserved during the setup phase; the resources remain dedicated for the entire duration of data transfer until the teardown phase.

### Example 8.1

As a trivial example, let us use a circuit-switched network connect eight telephones in a small area. Communication is through 4-kHz voice channels. We assume that each link uses FDM to connect a maximum of two voice channels. The bandwidth of each link is then 8 kHz. Figure 8.4 shows the situation. Telephone 1 is connected to telephone 7; 2 to 5; 3 to 8; and 4 to 6. Of course the situation may change when new connections are made. The switch controls the connections.

Figure 8.4 Circuit-switched network used in Example 8.1



### Example 8.2

As another example, consider a circuit-switched network that connects computers in two remote offices of a private company. The offices are connected using a T-1 line leased from a communication service provider. There are two  $4 \times 8$  (4 inputs and 8 outputs) switches in this network. For each switch, four output ports are folded into the input ports to allow communication between computers in the same office. Four other output ports allow communication between the two offices. Figure 8.5 shows the situation.

#### Figure 8.5 Circuit-switched network used in Example 8.2

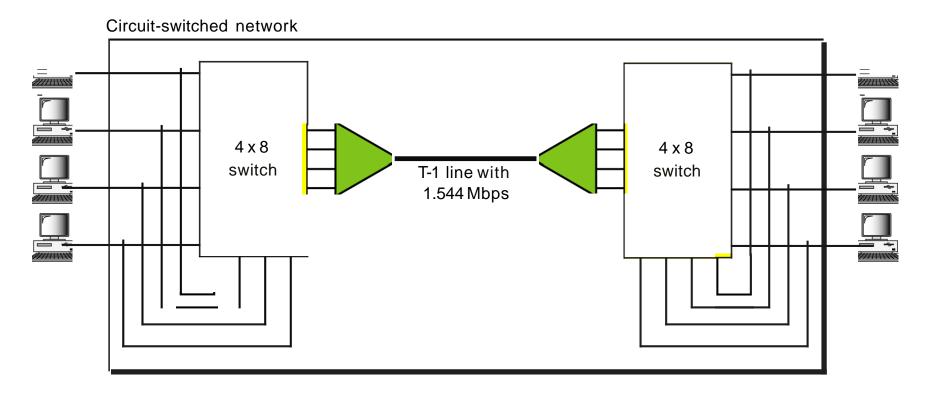
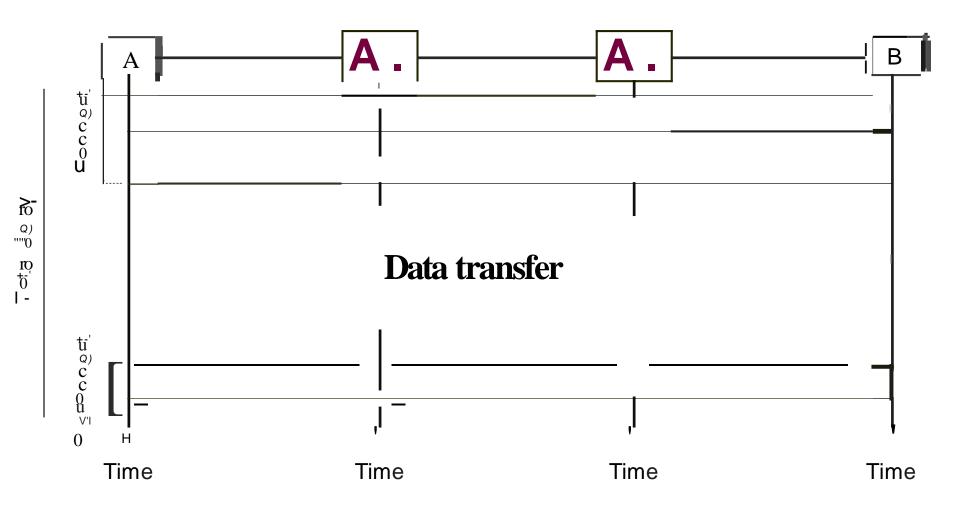


Figure 8.6 Delay in a circuit-switched network



Switching at the physical layer in the traditional telephone network uses the circuit-switching approach.

#### 8-2 DATAGRAM NETWORKS

In data communications, we need to send messages from one end system to another. If the message is going to pass through a packet-switched network, it needs to be divided into packets of fixedor variable size. The size of the packet is determined by the network and the governing protocol.

#### Topics discussed in this section:

Routing Table Efficiency Delay

**8D4**tagram Networks in the Internet



## In a packet-switched network, there is no resource reservation; resources are allocated on demand.

Figure 8.7 A datagram network with four switches (routers)

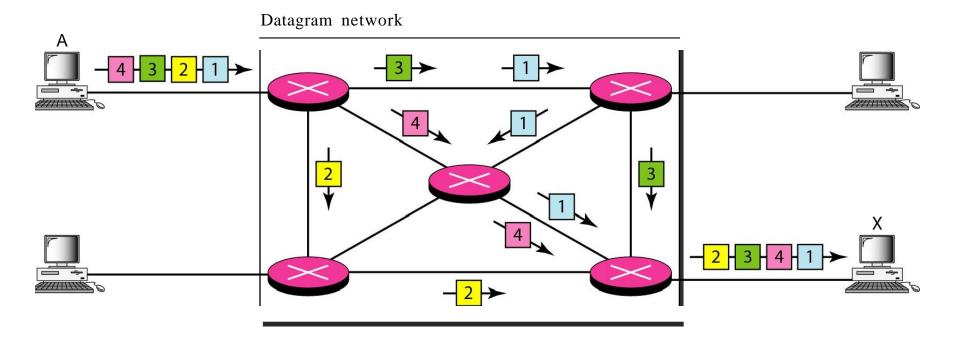


Figure 8.8 Routing table in a datagram network

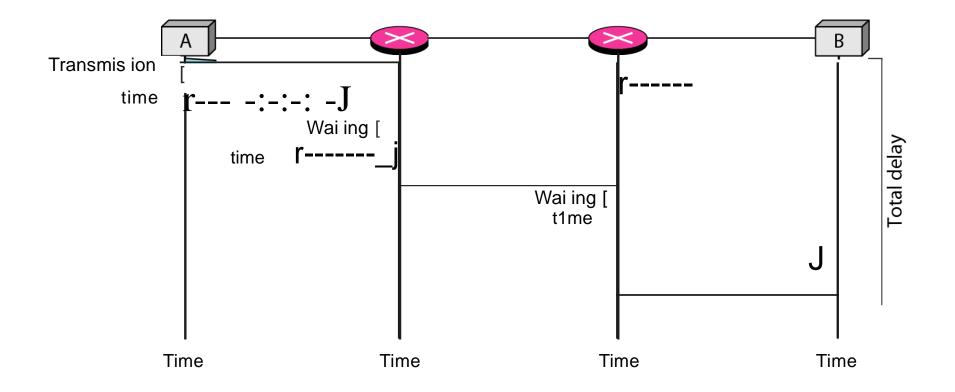
Destination		Output
address		port
1232		1
4150		2
9130		3
,		
1		4
	2	3



A switch in a datagram network uses a routing table that is based on the destination address.

The destination address in the header of a packet in a datagram network remains the same during the entire journey of the packet.

Figure 8.9 Delay in a datagram network



Switching in the Internet is done by using the datagram approach to packet switching at the network layer.

#### 8-3 VIRTUAL-CIRCUIT NETWORKS

A virtual-circuit network is a cross between a circuitswitched network and a datagram network. It has some characteristics of both.

#### Topics discussed in this section:

Addressing
Three Phases
Efficiency

**Delay** 

**Circuit-Switched Technology in WANs** 

Figure 8.10 Virtual-circuit network

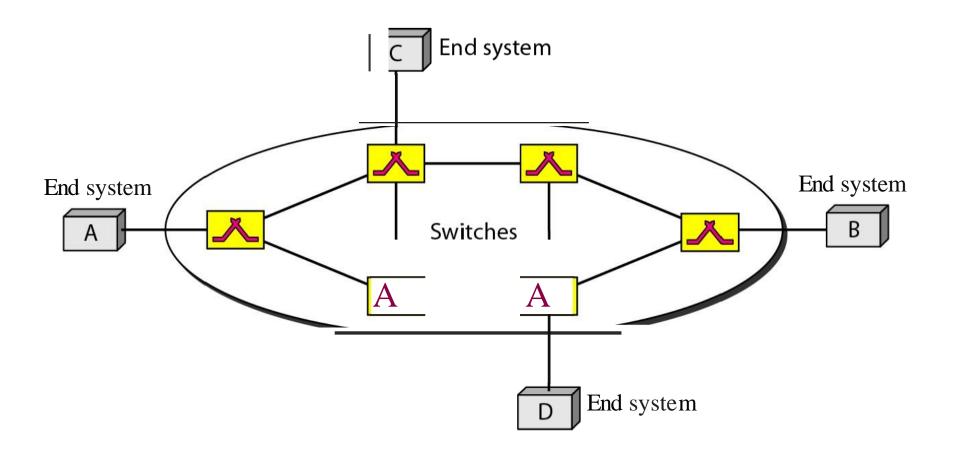


Figure 8.11 Virlual-circuit identifier

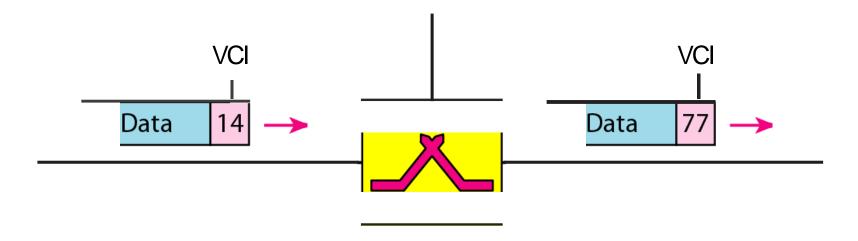


Figure 8.12 Switch and tables in a virtual-circuit network

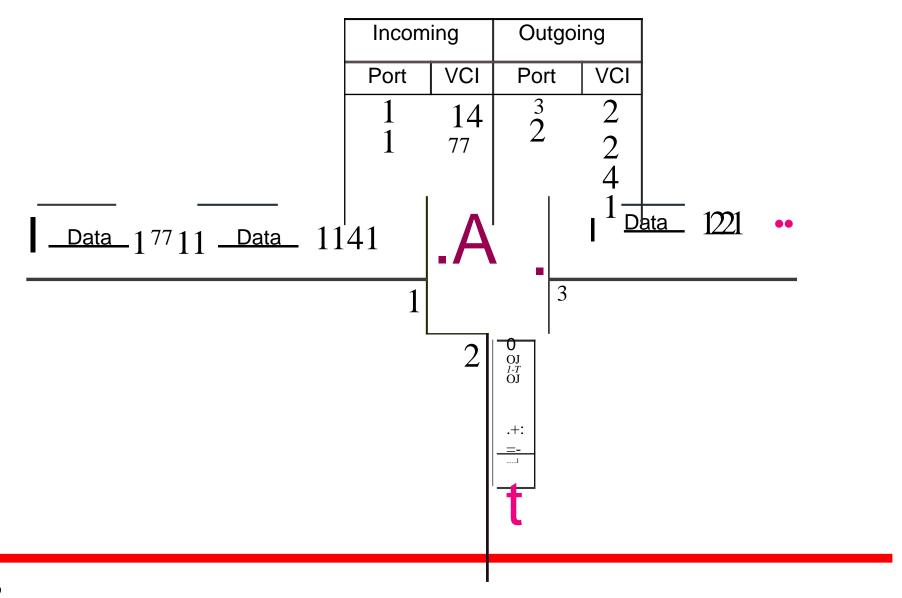


Figure 8.13 Source-to-destination data transfer in a virtual-circuit network

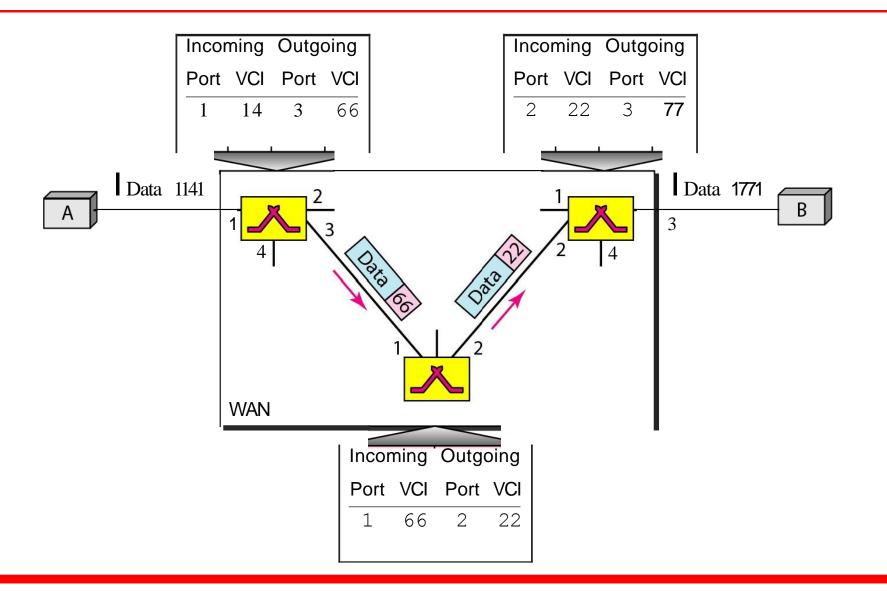


Figure 8.14 Setup request in a virtual-circuit network

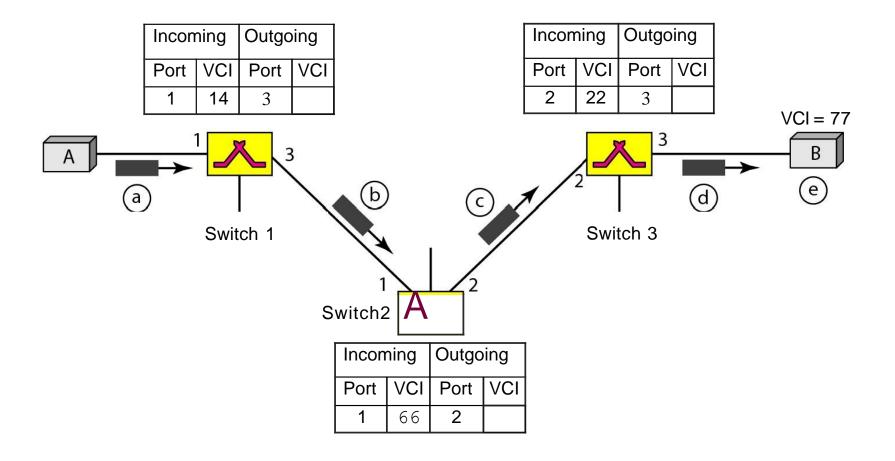
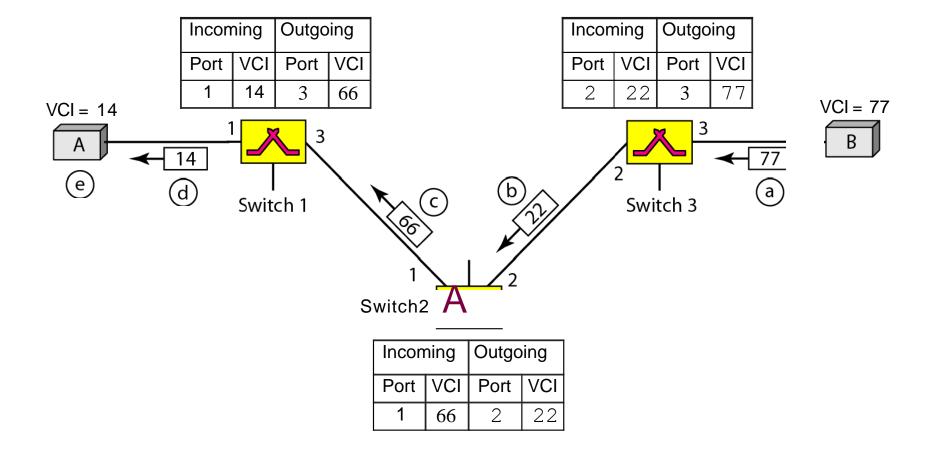
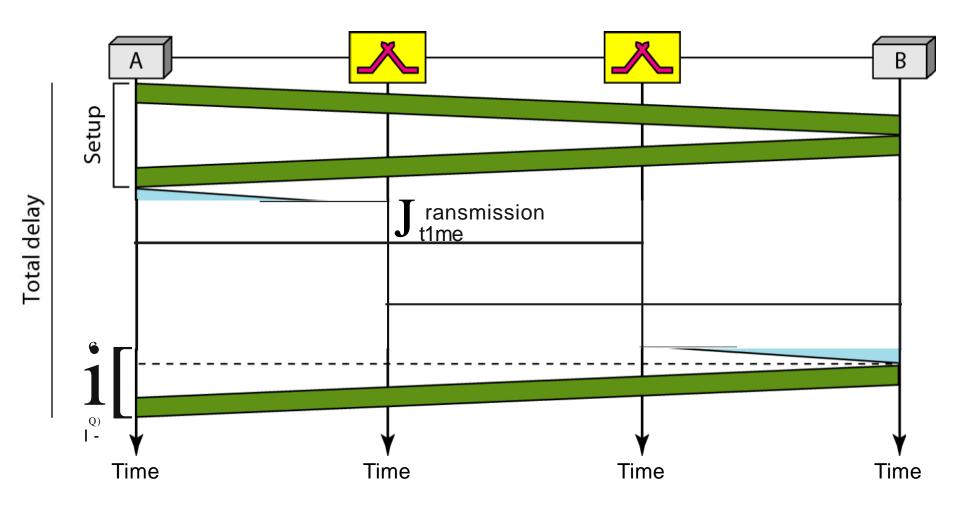


Figure 8.15 Setup acknowledgment in a virtual-circuit network



In virtual-circuit switching, all packets belonging to the same source and destination travel the same path; but the packets may arrive at the destination with different delays if resource allocation is on demand.

Figure 8.16 Delay in a virtual-circuit network





Switching at the data link layer in a switched WAN is normally implemented by using virtual-circuit techniques.

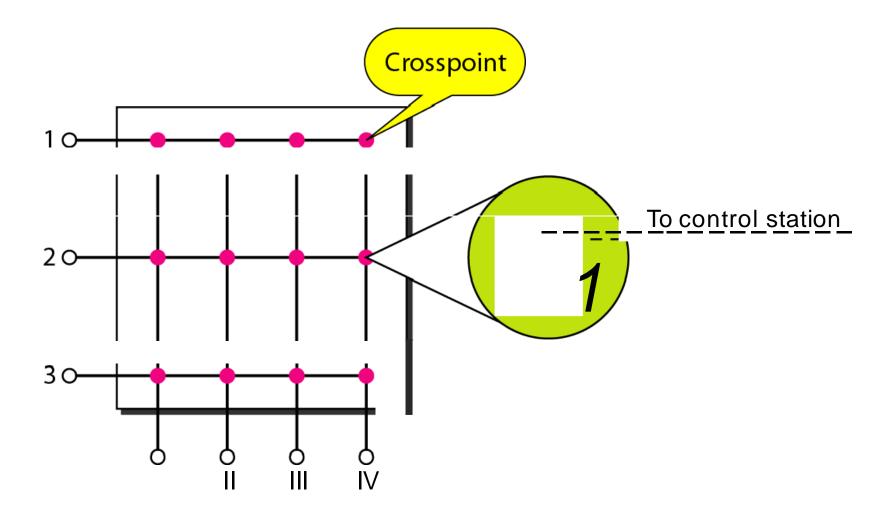
#### 8-4 STRUCTURE OF A SWITCH

We use switches in circuit-switched and packetswitched networks. In this section, we discuss the structures of the switches used in each type of network.

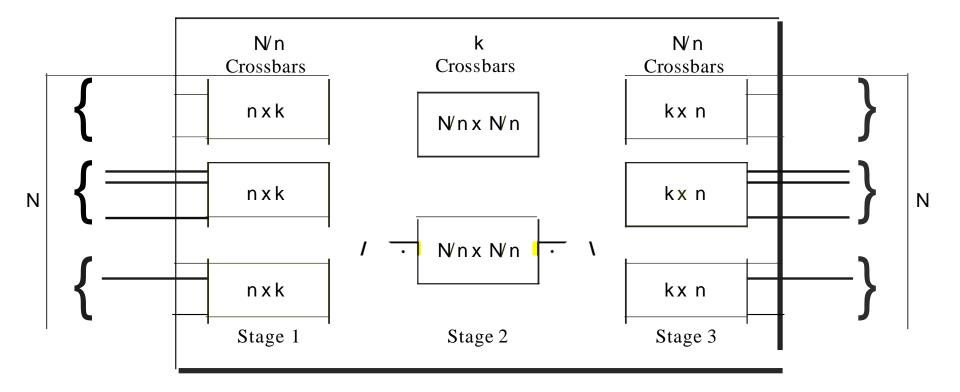
#### Topics discussed in this section:

**Structure of Circuit Switches Structure of Packet Switches** 

Figure 8.17 Crossbar switch with three inputs and four outputs



#### Figure 8.18 Multistage switch





## In a three-stage switch, the total number of crosspoints is $\frac{2kN + k(N/n)^2}{}$

which is much smaller than the number of crosspoints in a single-stage switch  $(N^2)$ .

#### Example 8.3

8.36age switch  $(200 \times 200 = 40,000)$ .

• Design a three-stage,  $200 \times 200$  switch (N = 200) with k = 4 and n = 20.

## • Soluti on

• In the first stage we have N/n or 10 crossbars, each of size  $20 \times 4$ . In the second stage, we have 4 crossbars, each of size  $10 \times 10$ . In the third stage, we have 10 crossbars, each of size  $4 \times 20$ . The total number of crosspoints is  $2kN + k(N/n)^2$ , or 2000 crosspoints. This is 5 percent of the number of crosspoints in a single-

- According to the Clos criterion:
- $\bullet \ n = (N/2)^{1/2}$
- k > 2n-1
- Crosspoints  $\geq 4N [(2N)^{1/2}-1]$

## Example 8.4

Redesign the previous three-stage, 200 × 200 switch, using the Clos criteria with a minimum number of crosspoints.

#### Solution

We let  $n = (200/2)^{1/2}$ , or n = 10. We calculate k = 2n - 1 = 19. In the first stage, we have 200/10, or 20, crossbars, each with  $10 \times 19$  crosspoints. In the second stage, we have 19 crossbars, each with  $10 \times 10$  crosspoints. In the third stage, we have 20 crossbars each with  $19 \times 10$  crosspoints. The total number of crosspoints is  $20(10 \times 19) + 19(10 \times 10) + 20(19 \times 10) = 9500$ .

#### Figure 8.19 Time-slot interchange

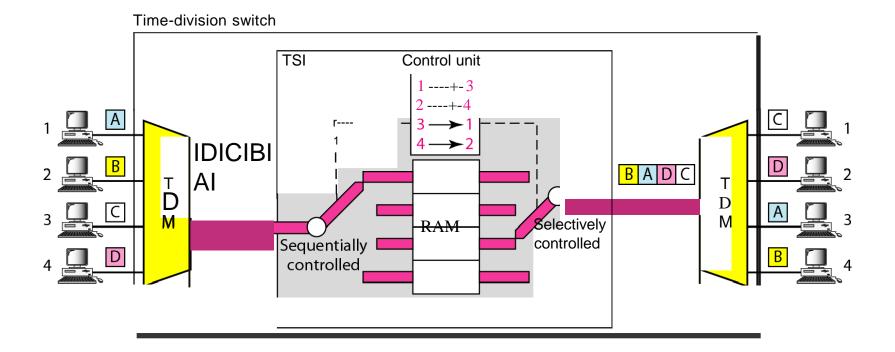


Figure 8.20 Time-space-time switch

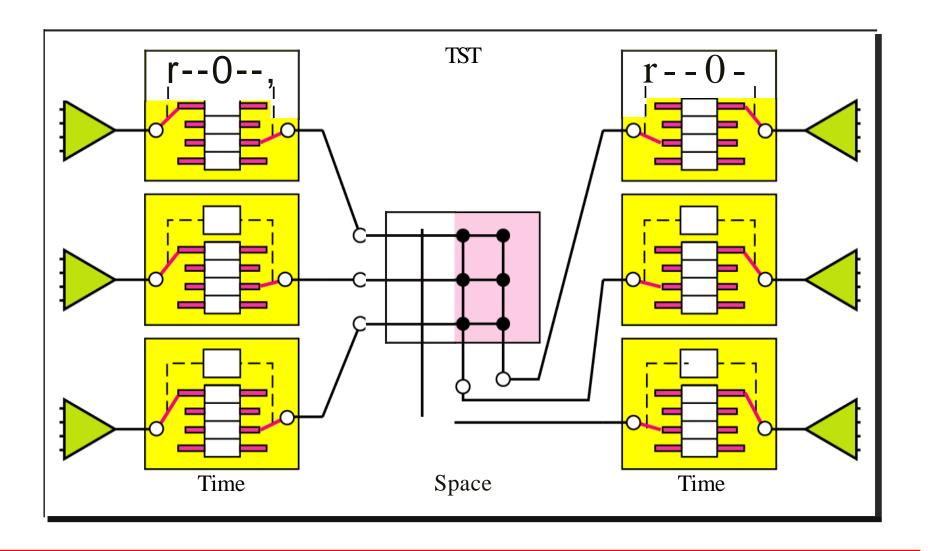
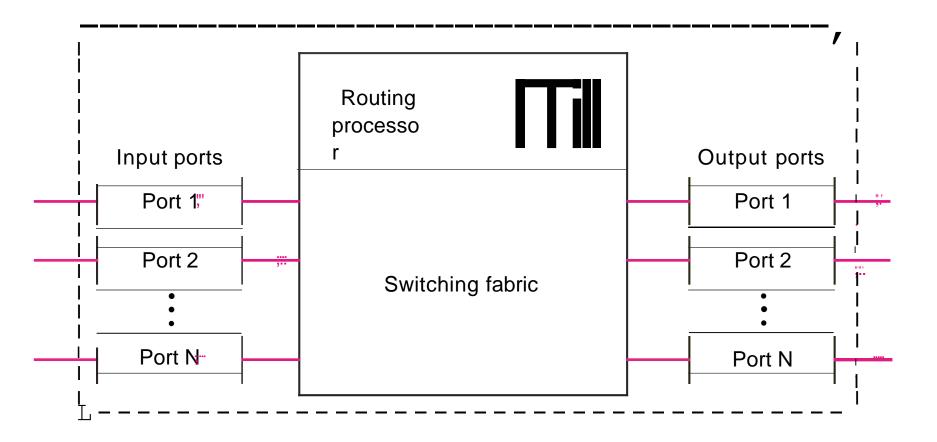
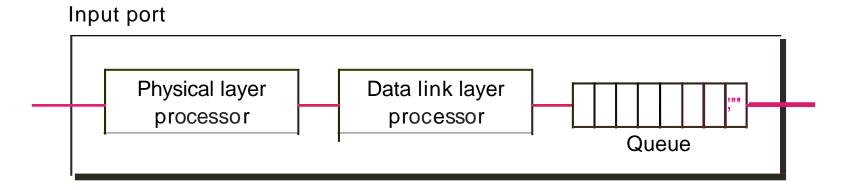


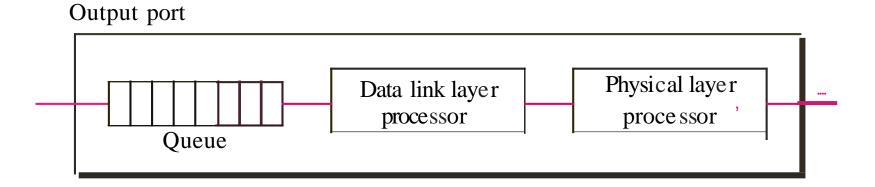
Figure 8.21 Packet switch components



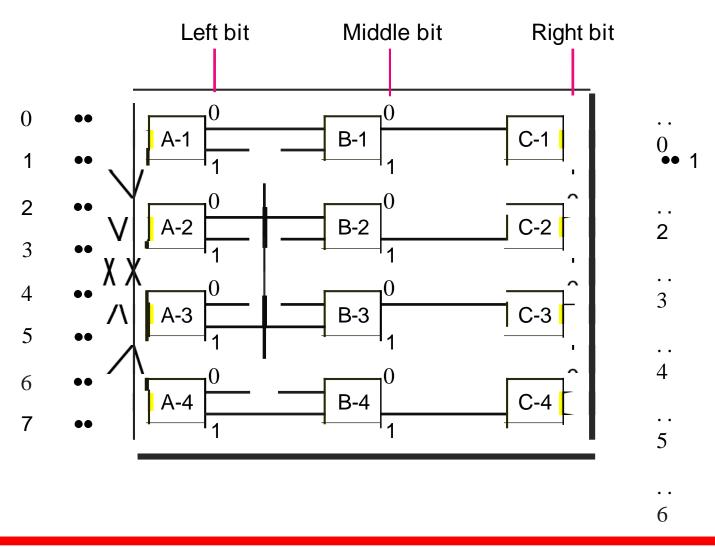
#### Figure 8.22 Input port



#### Figure 8.23 Outputport

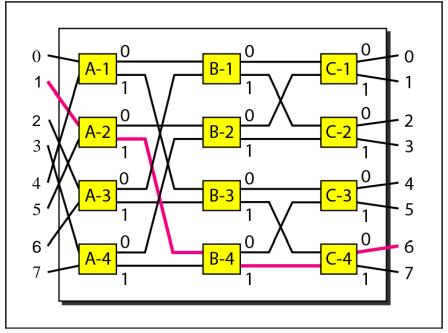


#### Figure 8.24 A banyan switch

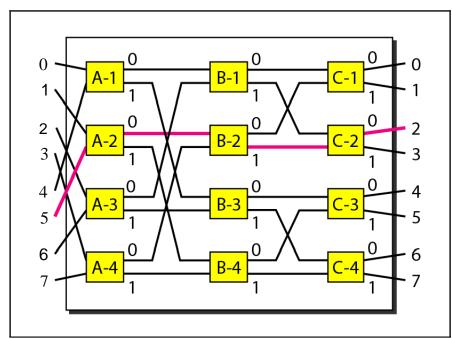


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#### Figure 8.25 Examples of routing in a banyan switch



a. Input 1 sending a cell to output 6 (110)



b. Input 5 sending a cell to output 2 (010)

#### Figure 8.26 Butcher-banyan switch

