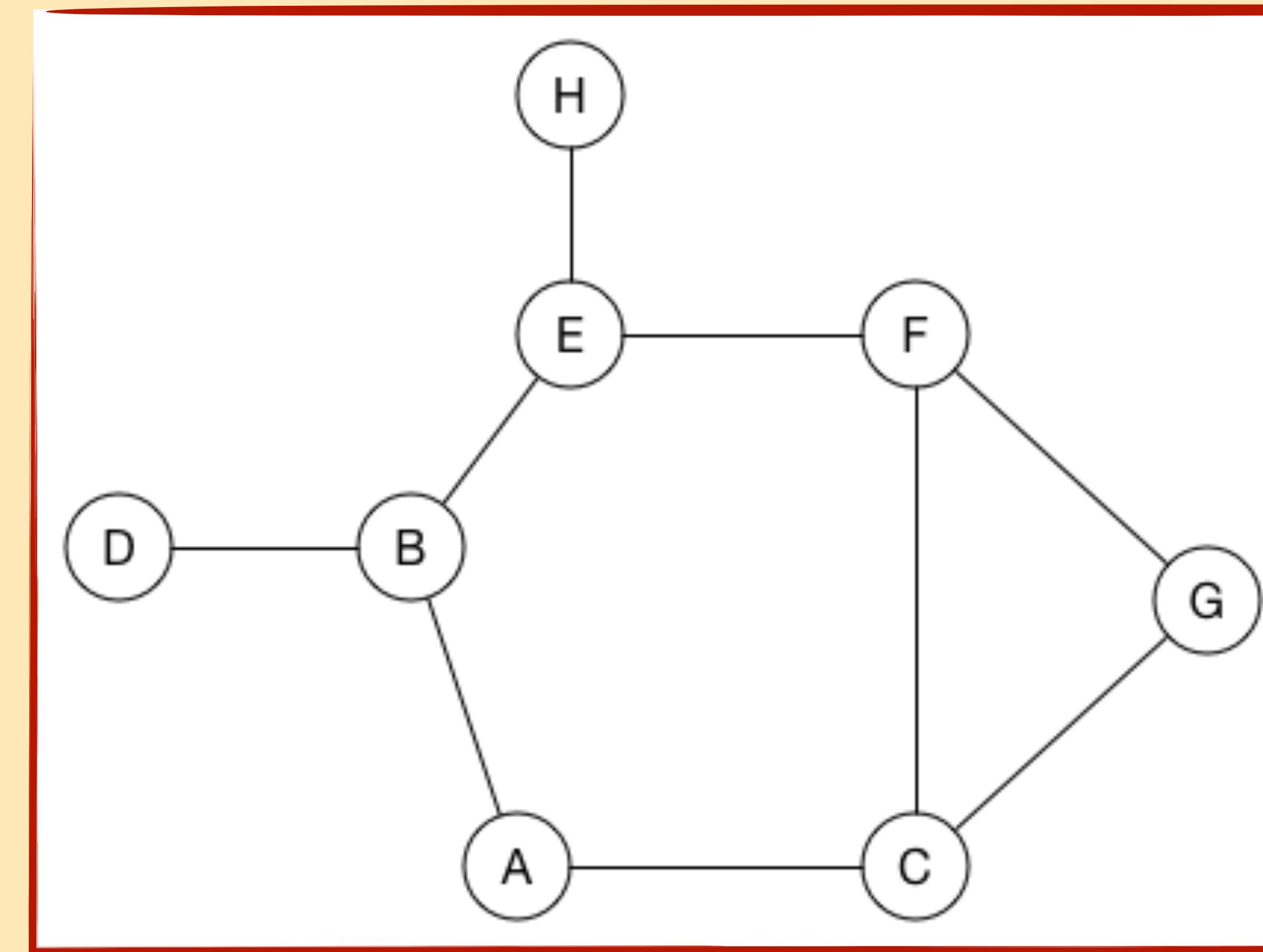


# Graph Theory II

## Seminar 13.

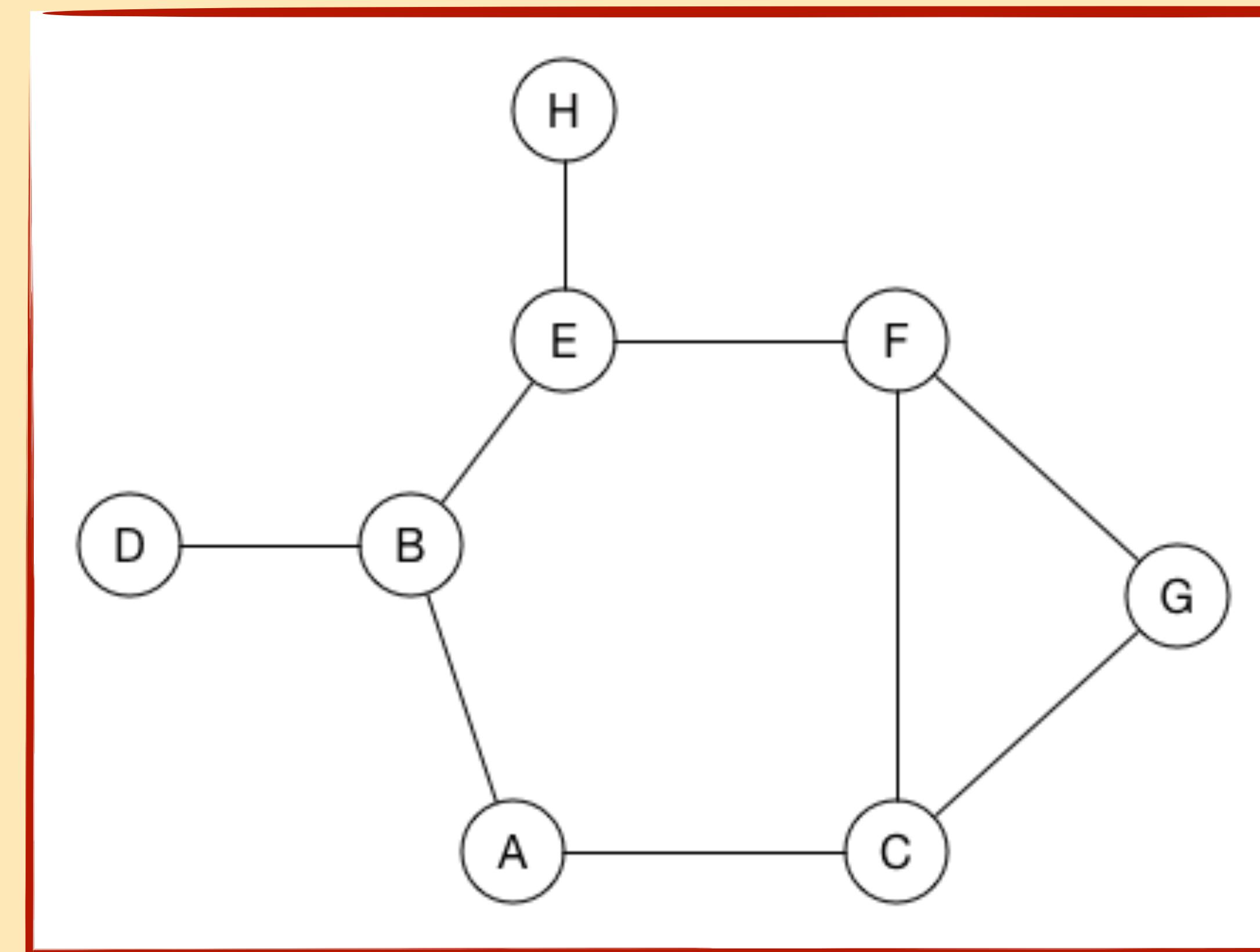
# Graph Search algorithms

## Breadth First Search



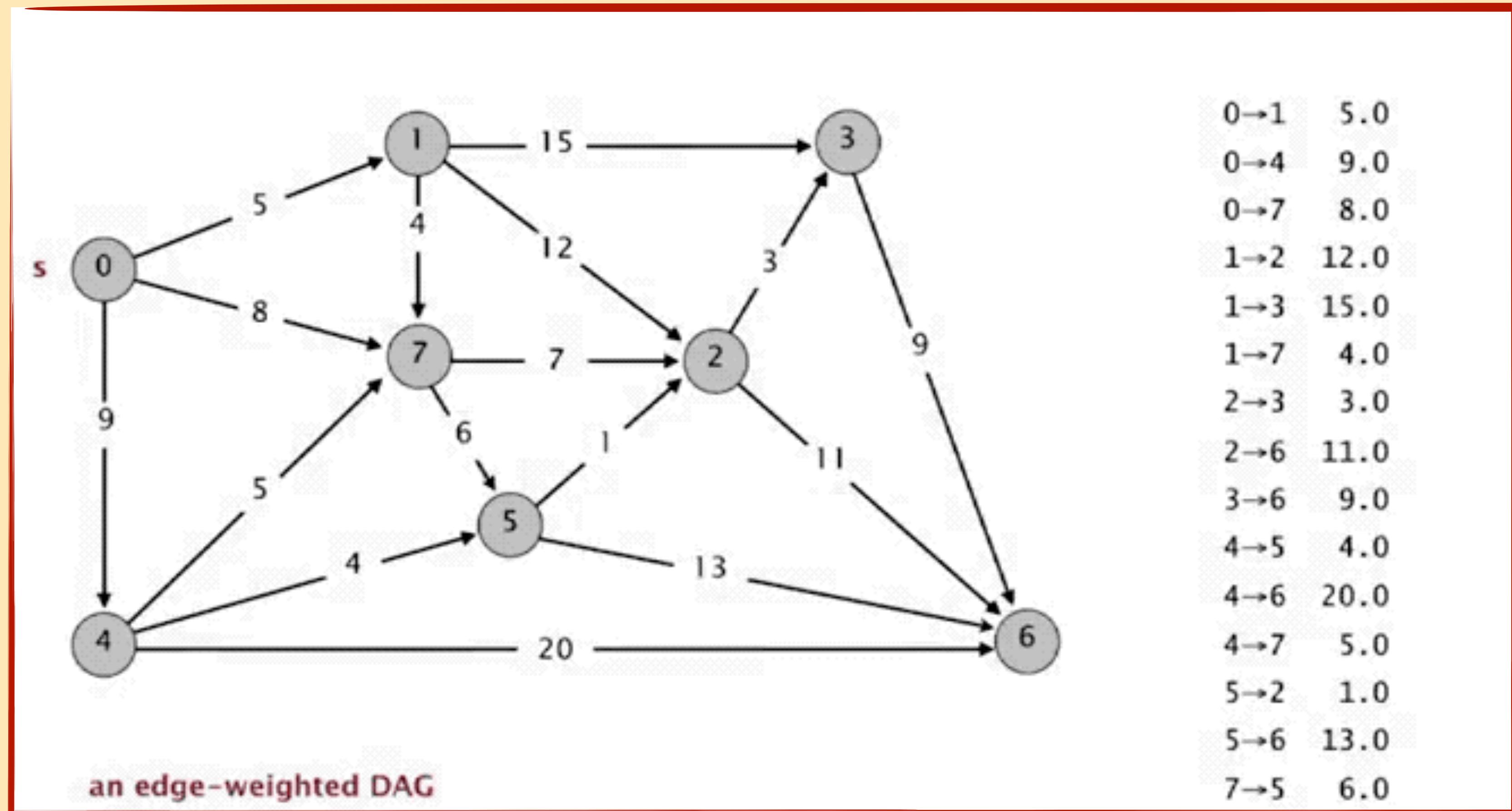
# Graph Search algorithms

## Depth First Search



# Minimum distances

## Dijkstra algorithm



# Minimum distances

## Matrix multiplication

- *Task I:* № of paths of fixed length
- *Task II:* All pairs shortest path

$\mathbf{g}$  – adjacency matrix

# Minimum distances

## Matrix multiplication

- *Task I:* № of paths of fixed length   $d_{k+1}[i][j] = \sum_p d_k[i][p] \cdot g[p][j]$
- *Task II:* All pairs shortest path

**g** – adjacency matrix

# Minimum distances

## Matrix multiplication

- *Task II:* All pairs shortest path  ameliorate?

### The algorithm

Consider a matrix  $D$  with  $d_{ij} = c_{ij}$  if there exists an arc from  $i$  to  $j$ ,  $\infty$  otherwise.

Let  $d_{ii} = 0$  for all  $i$ . Then we define the “dot” operation as

$$D_{ij}^{(2)} = (D \odot D^T)_{ij} = \min\{d_{i1} + d_{1j}, \dots, d_{in} + d_{nj}\}.$$

$D_{ij}^{(2)}$  represents the shortest path with 2 or fewer edges from  $i$  to  $j$ .

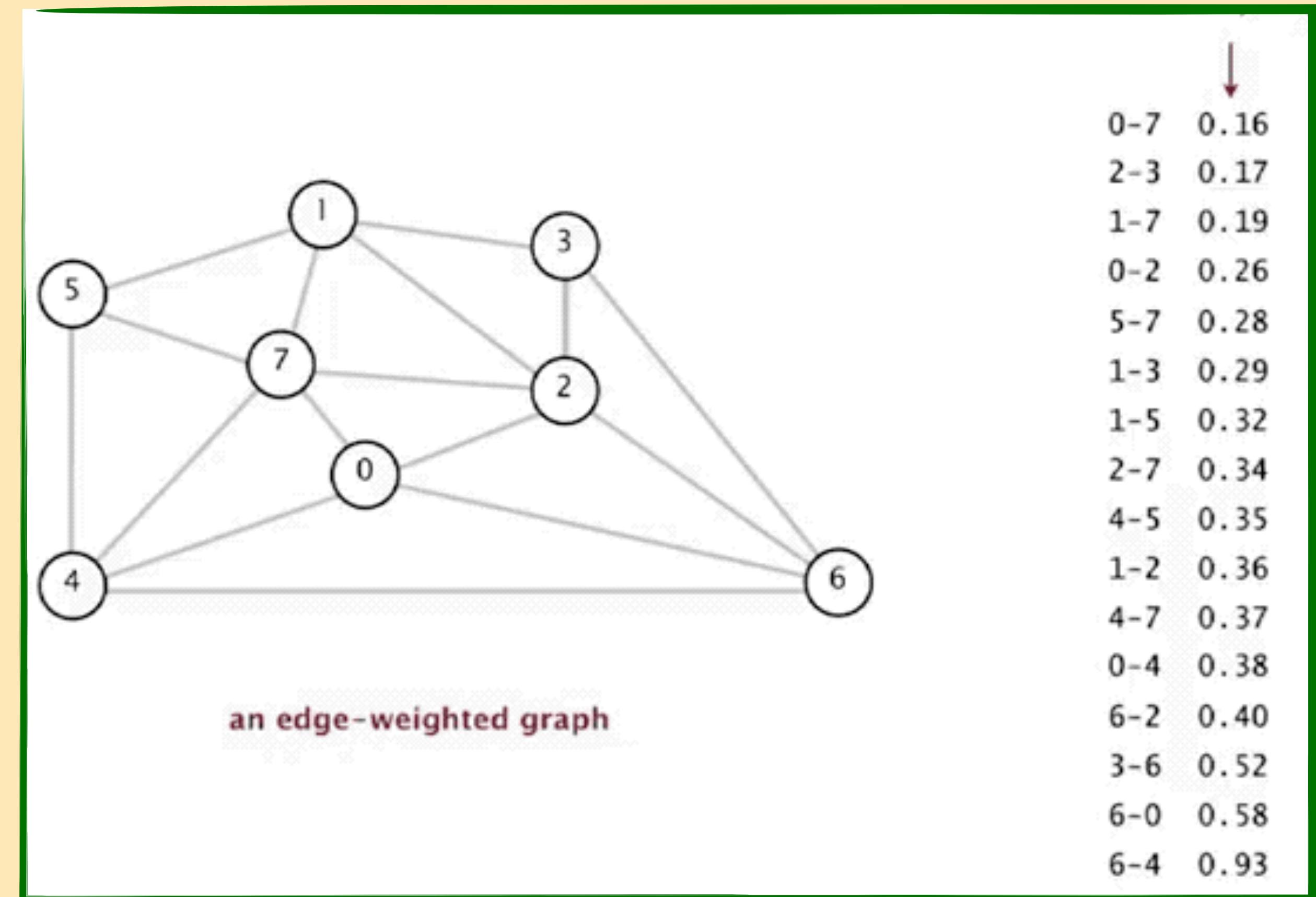
$D_{ij}^{(3)} = (D \odot D^{(2)T})_{ij} =$  shortest path with number of edges  $\leq 3$  from  $i$  to  $j$

⋮

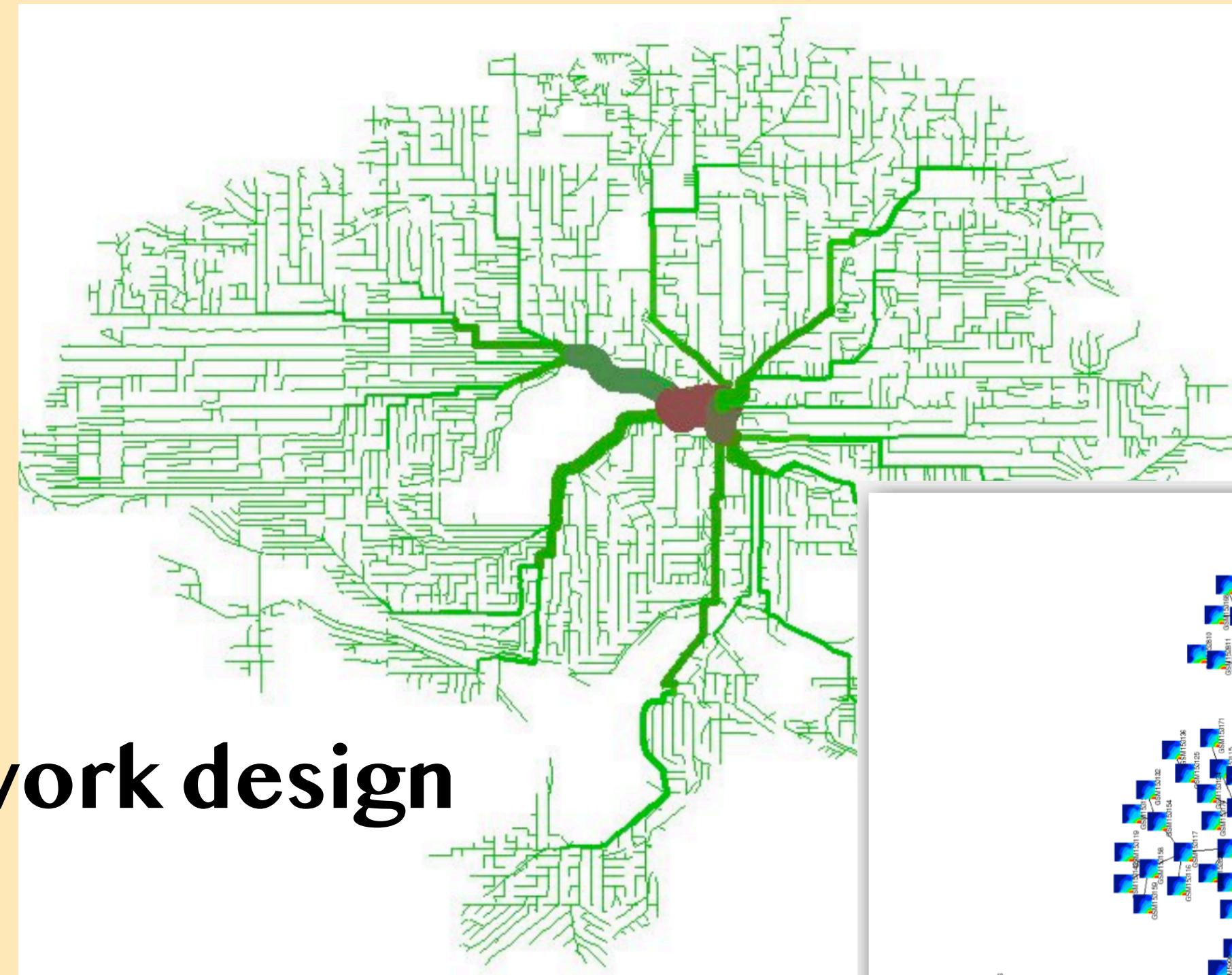
$D_{ij}^{(n)} = (D \odot D^{(n-1)T})_{ij} =$  shortest path with number of edges  $\leq n$  from  $i$  to  $j$

# Minimum Spanning Tree

- *Tree*: acyclic graph with  $n$  nodes and  $n-1$  edges with only one path existing between nodes
- *Minimum spanning*: tree with minimum overall weight

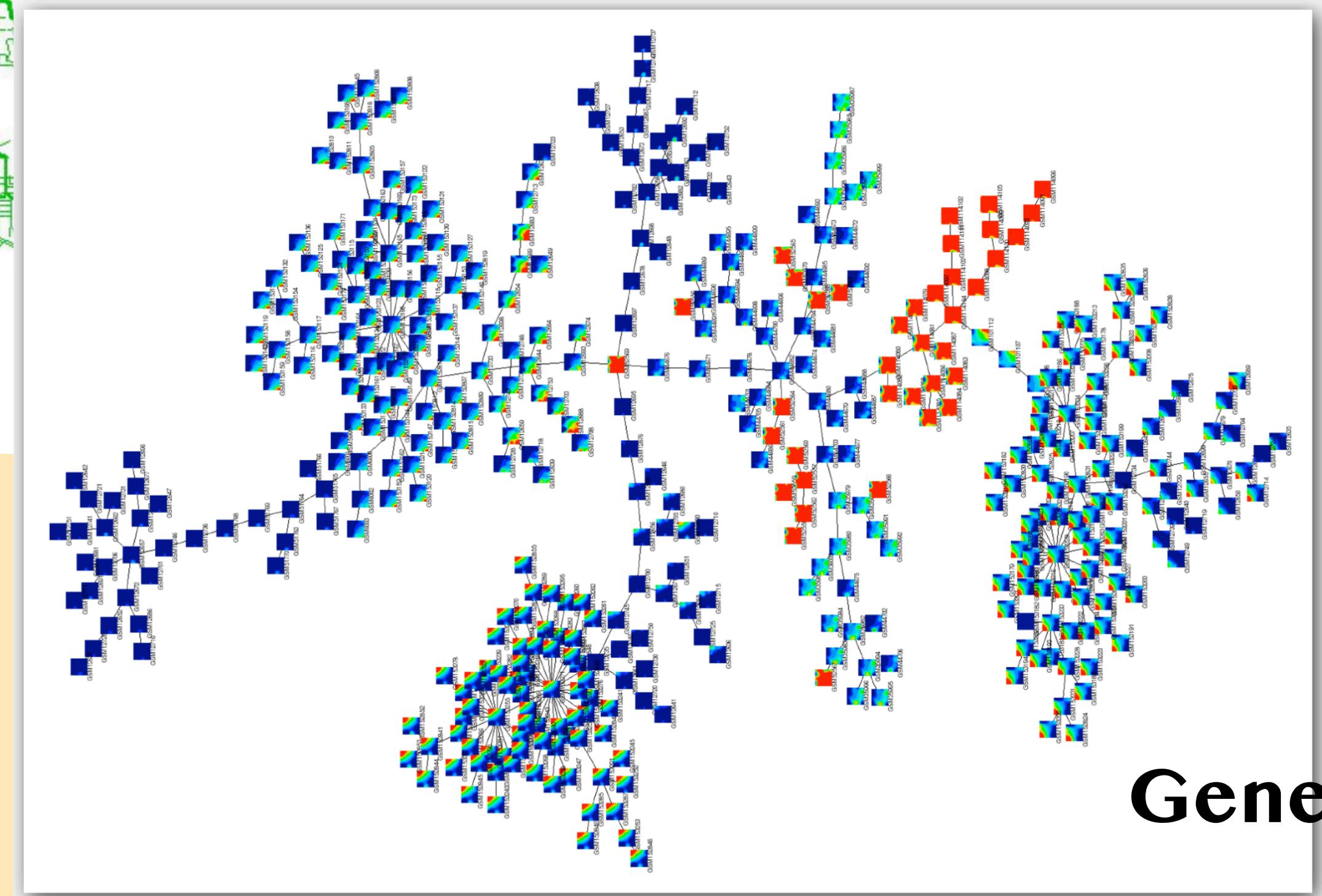
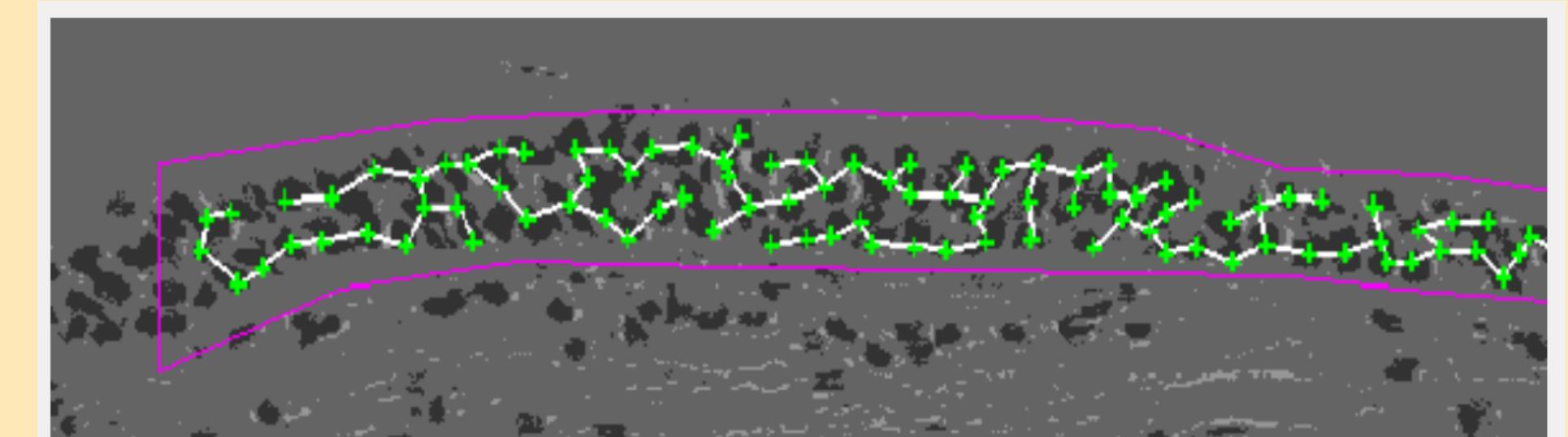


# Minimum Spanning Tree Applications



Network design

Medical images



Genetics

# Minimum Spanning Tree

## Prim's algorithm

- Choose source vertex
- Launch identical to Dijkstra
- Add the minimum edge to MST

**minKey()** – find minimum distance value unvisited neighbour

**prim()** – choose neighbour - update the following's distances and ancestors

# Minimum Spanning Tree

## Kruskal's algorithm

- Iterate through edges in non-decreasing order
- Check for safe edge => Add to the MST subgraph if TRUE

**find()** – identify the root index of particular node

**union()** – join two subsets

**kruskal()** – take edges - check for cycles - unite

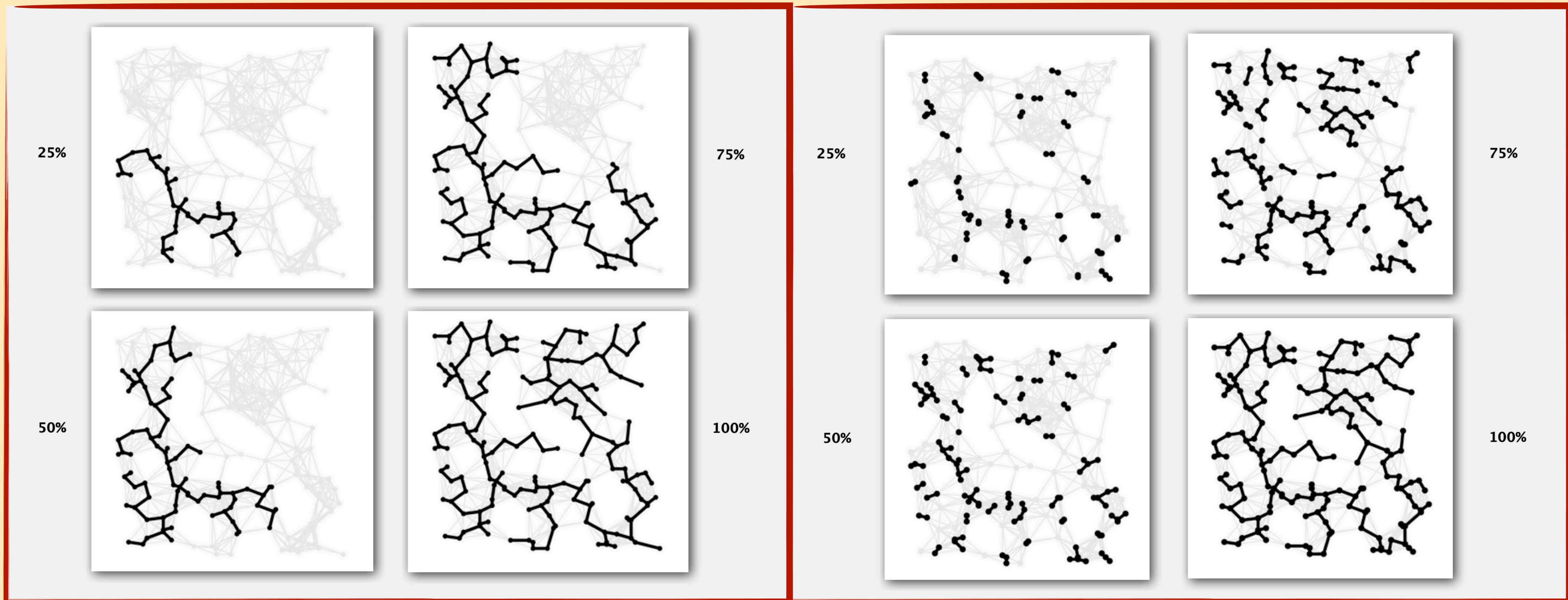
# Minimum Spanning Tree

## Algorithms' comparison

Prim's Algorithm	Kruskal's Algorithm
It starts to build the Minimum Spanning Tree from any vertex in the graph.	It starts to build the Minimum Spanning Tree from the vertex carrying minimum weight in the graph.
It traverses one node more than one time to get the minimum distance.	It traverses one node only once.
Prim's algorithm has a time complexity of $O(V^2)$ , $V$ being the number of vertices and can be improved up to $O(E \log V)$ using Fibonacci heaps.	Kruskal's algorithm's time complexity is $O(E \log V)$ , $V$ being the number of vertices.
Prim's algorithm gives connected component as well as it works only on connected graph.	Kruskal's algorithm can generate forest(disconnected components) at any instant as well as it can work on disconnected components

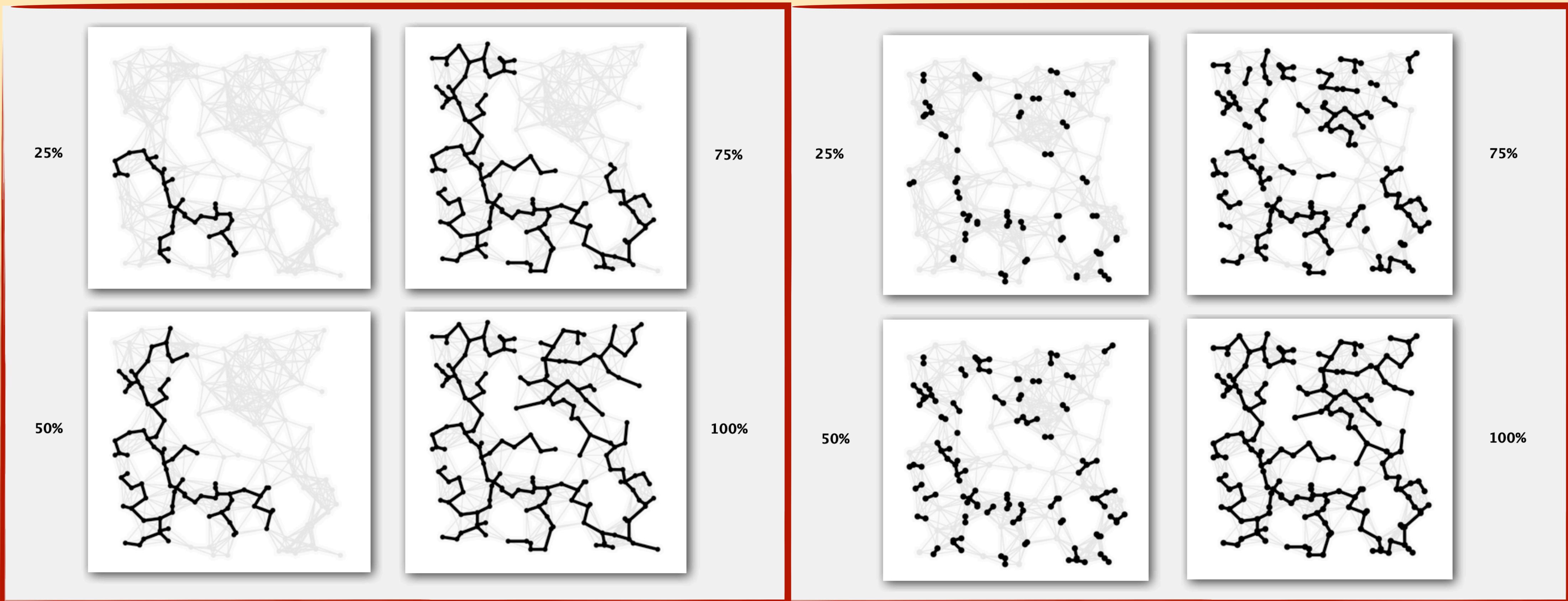
# Minimum Spanning Tree

## Algorithms' comparison



# Minimum Spanning Tree

## Algorithms' comparison



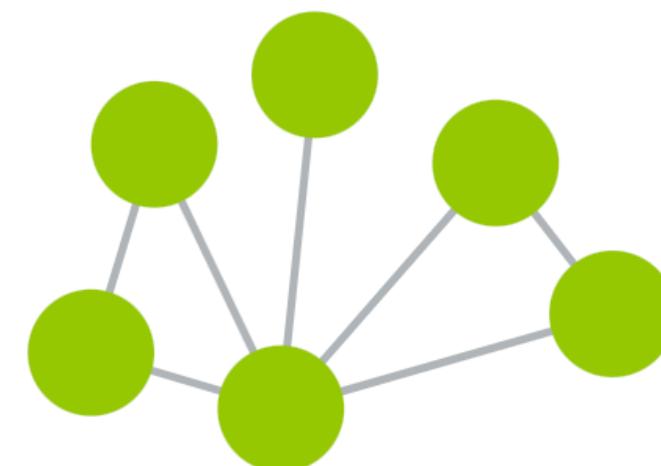
Prim

Kruskal

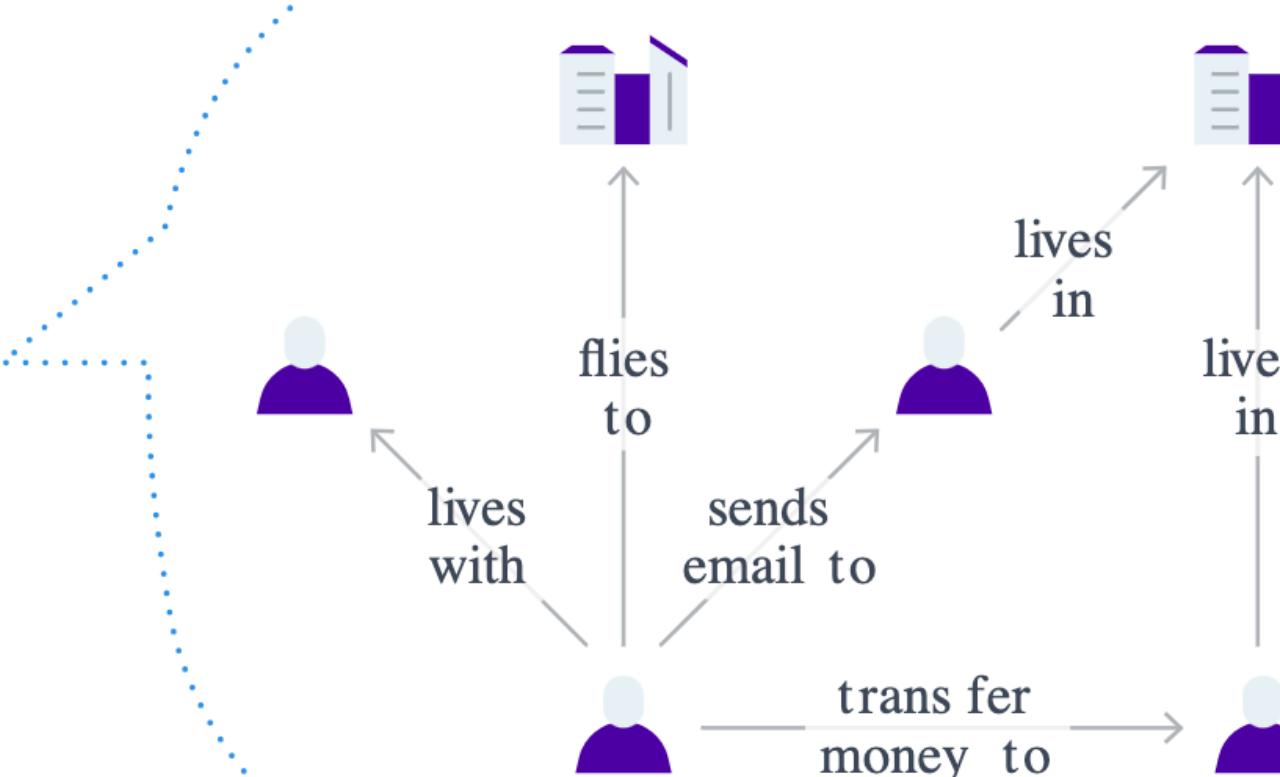
# Graphs & ML

## Types of graphs

Homogeneous

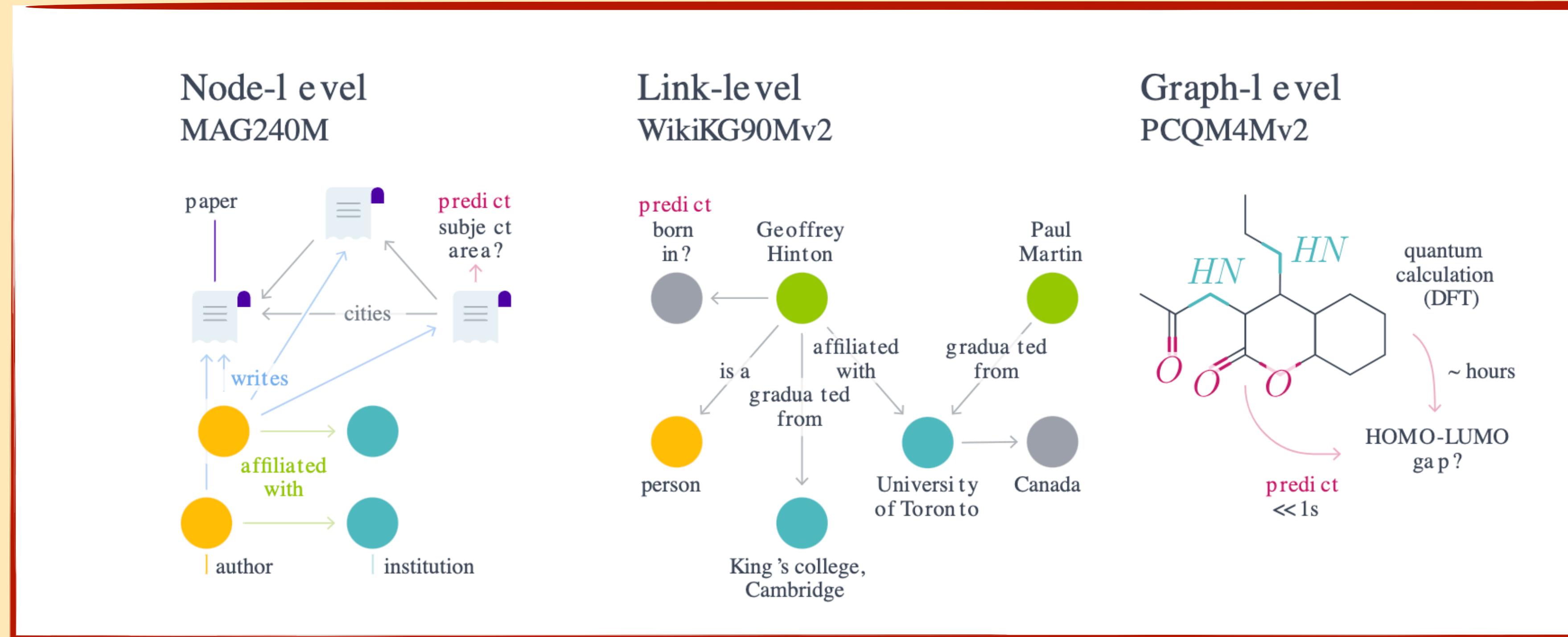


Heterogeneous



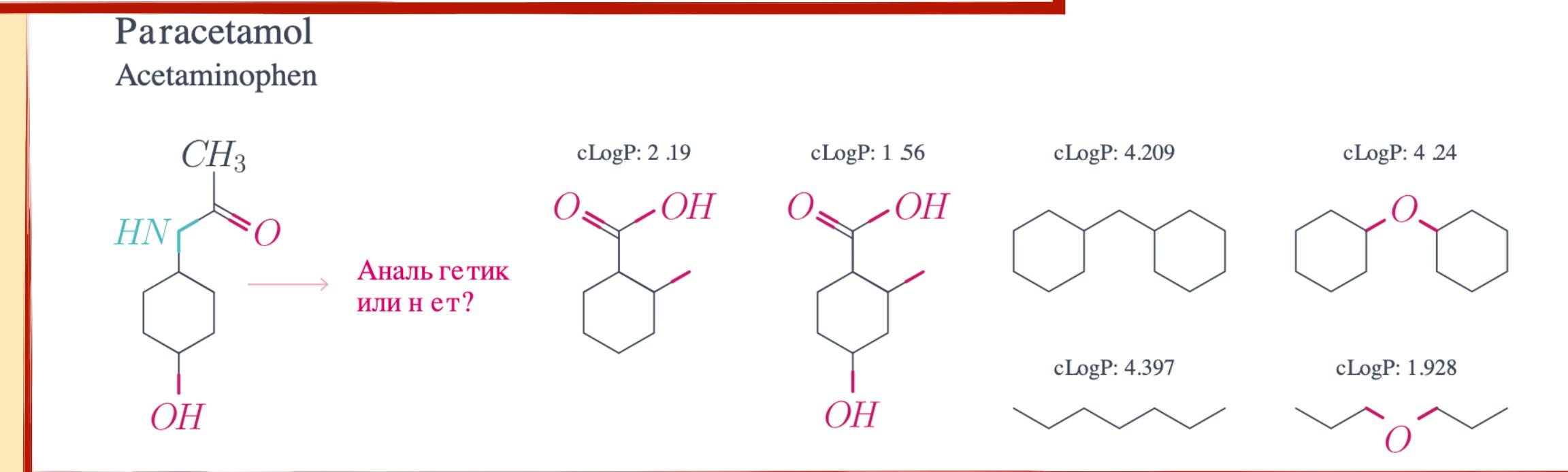
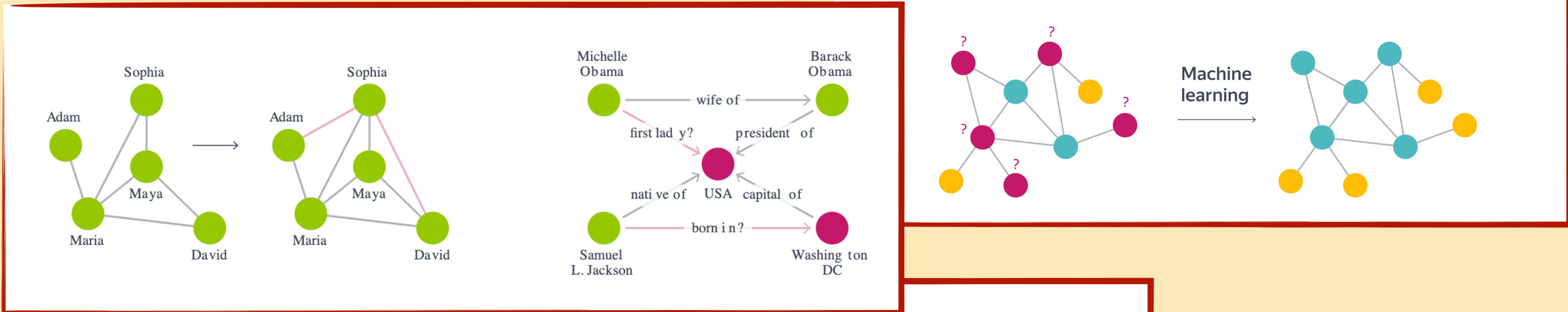
# Graphs & ML

## Types of tasks



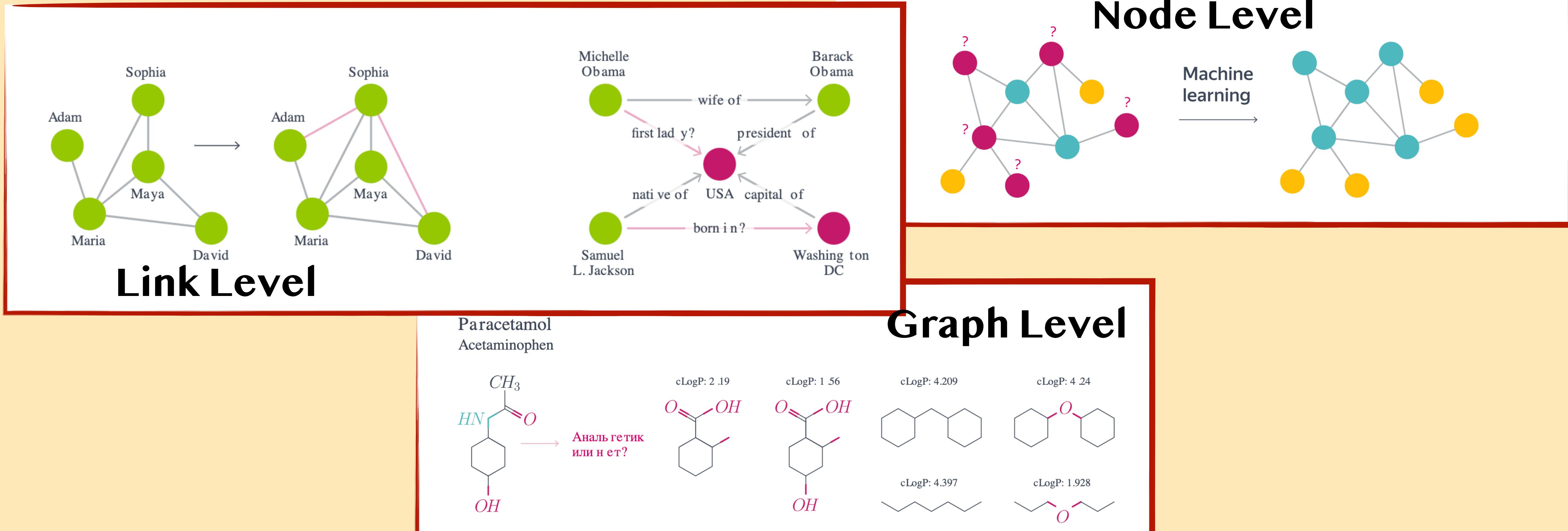
# Graphs & ML

## Types of tasks



# Graphs & ML

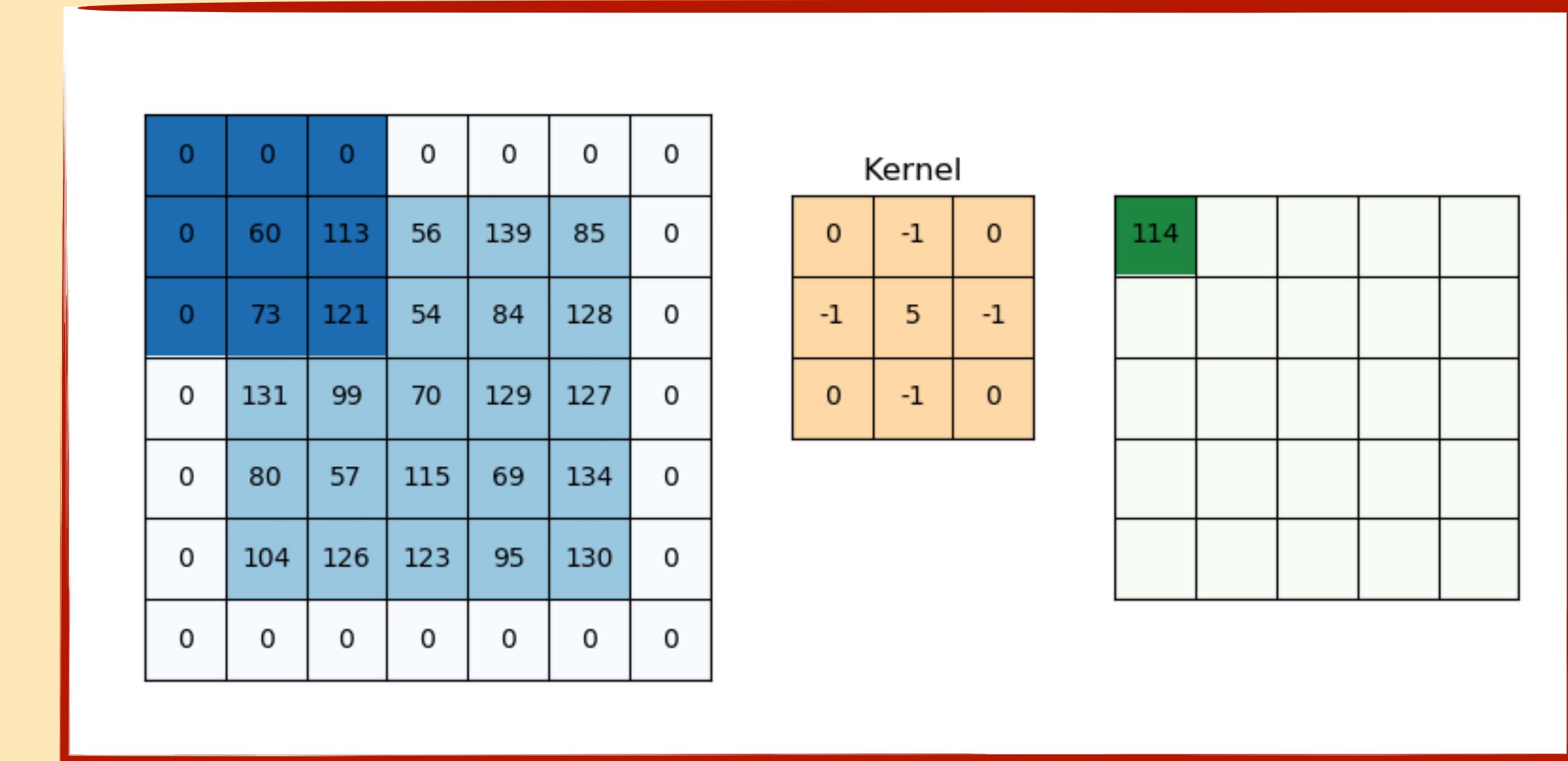
## Types of tasks



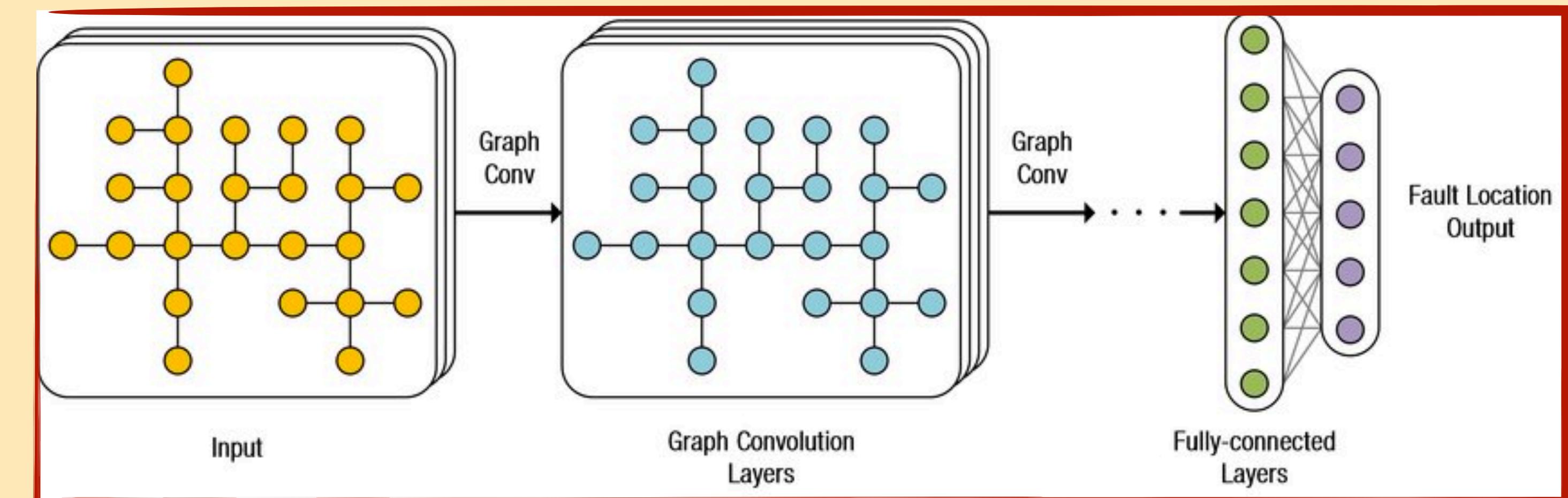
# Graphs & ML

## Neural Network: Arch

For images



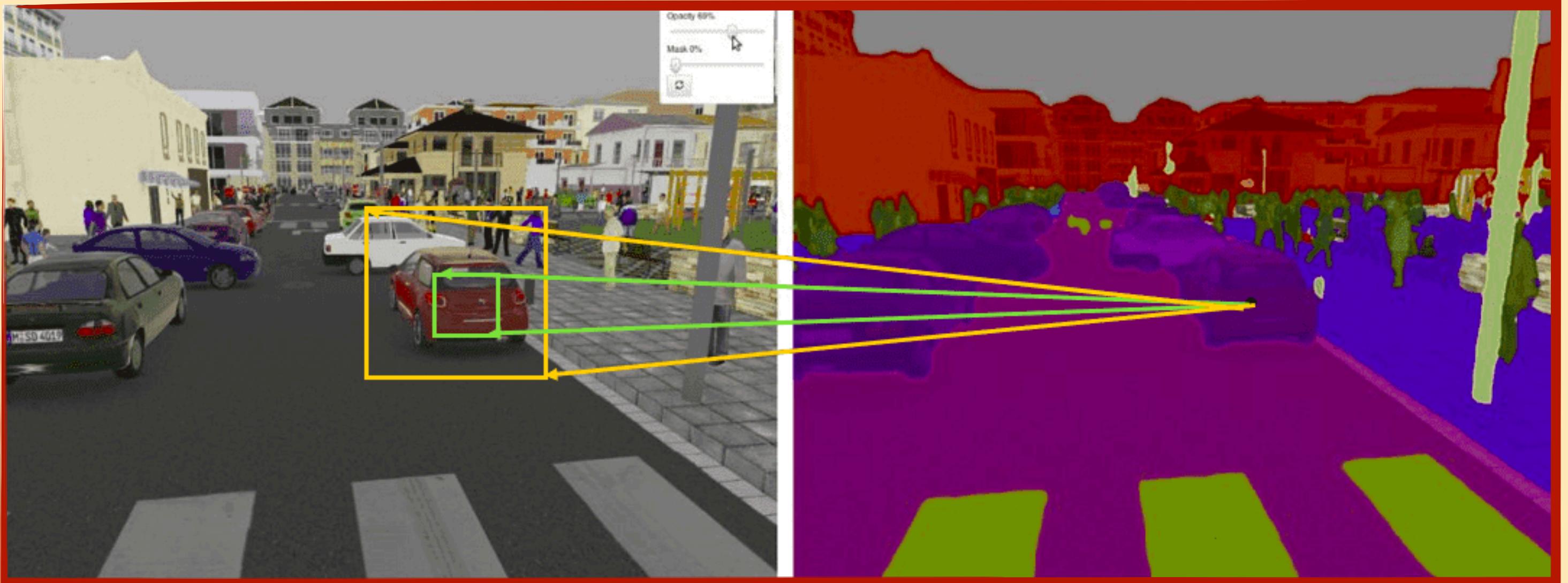
For graphs



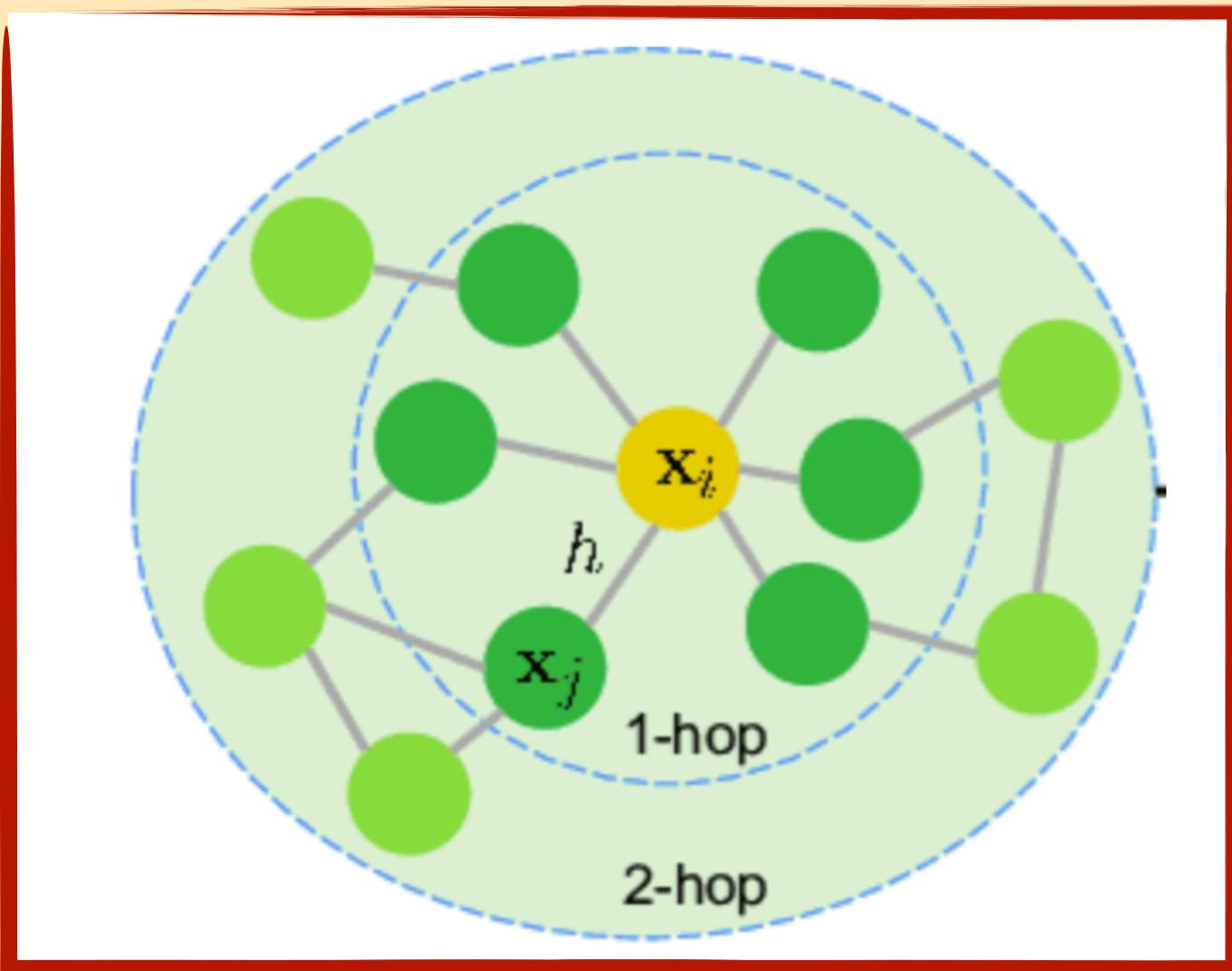
# Graphs & ML

## Neural Network: receptive field

For images



For graphs



# Graphs & ML

## Knowledge graphs & RL

