

## 4. Root finding

## Last time

- Approximating functions by polynomials
- Taylor series
- Taylor polynomials + Lagrange remainder

## Goal for today

- Methods for finding **roots** of functions
- Convergence of those methods

# Solving equations

- Suppose we need to solve the equation  $g(x) = h(x)$
- Write  $f(x) := g(x) - h(x)$
- Equivalent to solve  $f(x) = 0$  instead
- Solutions  $x^*$  such that  $f(x^*) = 0$ : **roots** or **zeros** of  $f$
- We may want to find a single root or all roots

# Roots of polynomials

- Suppose  $p(x) = a_0 + a_1x + \dots + a_nx^n$  is **polynomial**
- How many roots does  $p$  have?
- **Fundamental theorem of algebra:**  
*a degree- $n$  polynomial has exactly  $n$  roots in  $\mathbb{C}$*
- So can factor  $p$  as

$$p(x) = (x - x_1)^{m_1}(x - x_2)^{m_2} \dots (x - x_k)^{m_k}$$

where  $x_i$  are the roots and  $\sum_{i=1}^k m_i = n$

- $m_i$  is the **multiplicity** of root  $x_i$

## Roots of polynomials II

- Some of the roots may be **multiple roots**
- E.g.  $p = (x - 1)^2(x^2 - 2)$  has roots 1, 1,  $\sqrt{2}$  and  $-\sqrt{2}$
- Roots may be complex even if coefficients  $a_i$  are real
- E.g.  $x^2 + 1$  has roots  $x_{1,2} = \pm i$  and *no* real roots

# Can we find roots of polynomial analytically?

- Suppose  $p$  is a degree- $n$  polynomial with coefficients  $a_i$
- $n = 2$  (quadratic): exact formula for roots in terms of  $a_i$
- $n = 3, 4$ : **more complicated formulae** exist
- **Abel–Ruffini theorem:**  
for  $n \geq 5$ , in general ***no such formula exists!***

## Numerical methods for root finding

- Need **numerical methods** to **approximate** roots
- Some methods for *all* roots of a polynomial simultaneously, e.g. **Aberth method** (PS 2)
- Here we focus on methods to find single root



# Fixed-point iteration

- Numerical method for roots must be **iterative**:

$$x_{n+1} = g(x_n)$$

- algorithm  $g$  and initial guess  $x_0$
- Produces a sequence  $x_0, x_1, x_2, \dots$
- If iteration converges,  $x_n \rightarrow x^*$  as  $n \rightarrow \infty$ , then

$$g(x^*) = x^*$$

provided  $g$  is **continuous**

- $x^*$  is a **fixed point** of  $g$ ; solves  $f(x) := g(x) - x = 0$

## Existence of fixed point

- If  $g$  is continuous and maps  $[a, b]$  into itself,  $\exists$  a fixed point.
- If  $|g'(x)| \leq k \forall x \in [a, b]$  with  $k < 1$ , then unique.
- Called a **contraction mapping**

## Dynamics of iterations

- Such an iteration is a **discrete-time dynamical system**
- Assume there is a fixed point  $x^*$
- What condition do we need for  $(x_n)$  to converge?
- Look at distance  $\delta_n := x_n - x^*$

# Dynamics II

- We have

$$\delta_{n+1} = x_{n+1} - x^*$$

$$= g(x_n) - x^* = g(x^* + \delta_n) - x^* \simeq \delta_n g'(x^*)$$

- So (asymptotically)  $\delta_{n+1} = \alpha \delta_n$  with  $\alpha = g'(x^*)$
- $\delta_{n+1} = \alpha \delta_n \Rightarrow \delta_n = \alpha^n$
- Decays (**stable fixed point**) if  $|\alpha| < 1$
- Grows (**unstable fixed point**) if  $\alpha > 1$

# Rate of convergence

- For  $\sqrt{\phantom{x}}$ , Babylonian converges **faster** than bisection
- How can we make this precise?
- If  $x_n \rightarrow x^*$  then  $x_n - x^* \rightarrow 0$
- Estimate **distance** or **error**  $|x_n - x^*|$  as function of  $n$
- In bisection, error reduces by constant factor 2 at each step

## Rate of convergence II

- If there are constants  $\lambda$  and  $\alpha$  such that

$$\lim_{n \rightarrow \infty} \frac{|x_{n+1} - x^*|}{|x_n - x^*|^\alpha} = \lambda$$

then the iteration is of **order**  $\alpha$

- $\alpha = 1$ : **linearly convergent**
- $\alpha = 2$ : **quadratic convergence**
- Note that “linear” actually means that the sequence converges exponentially fast!

# Newton method

- Want to solve  $f(x) = 0$ ; have guess  $x_0$
- Let  $x = x_0 + \delta$ ; solve  $f(x_0 + \delta) = 0$
- Taylor theorem:  $f(x_0) + \delta f'(x_0) + \frac{1}{2}\delta^2 f''(\xi) = 0$
- So  $\delta \simeq \frac{-f(x_0)}{f'(x_0)}$
- **Newton method:**  $x_{n+1} = x_n - \frac{f(x_n)}{f'(x_n)}$
- **Quadratically convergent!**
- Babylon is special case

## Newton method viewed as a fixed-point iteration

- How **design** quadratically-convergent algorithm for  $f(x) = 0$ ?
- Look for fixed-point algorithm with  $g(x) := x - \phi(x)f(x)$
- Impose  $g'(x^*) = 0$  at root  $f(x^*) = 0$
- $g'(x) = 1 - \phi'(x)f(x) - \phi(x)f'(x)$
- So  $g'(x^*) = 1 - \phi(x^*)f'(x^*) = 0$
- So need  $\phi(x^*) = \frac{1}{f'(x^*)}$
- Take  $\phi(x) := \frac{1}{f'(x)}$  – gives Newton method



## Convergence of Newton method

- Newton is very powerful *when it works*
- Generalises
- When “sufficiently close” to a root, Newton converges fast
- Otherwise, *not guaranteed* to converge; may fail
- There are known conditions for Newton to converge:
  - Smale  $\alpha$  theory
  - Radii polynomials
- Interval Newton method: Use interval arithmetic to guarantee convergence

# Summary

- There are many methods for finding **roots** of functions
- Rewrite as **fixed-point iteration**
- Do not always converge
- Methods differ in **convergence rate**