

## Lagrangian

**standard form problem** (not necessarily convex)

$$\begin{array}{ll}\text{minimize} & f_0(x) \\ \text{subject to} & f_i(x) \leq 0, \quad i = 1, \dots, m \\ & h_i(x) = 0, \quad i = 1, \dots, p\end{array}$$

variable  $x \in \mathbf{R}^n$ , domain  $\mathcal{D}$ , optimal value  $p^\star$

**Lagrangian:**  $L : \mathbf{R}^n \times \mathbf{R}^m \times \mathbf{R}^p \rightarrow \mathbf{R}$ , with  $\text{dom } L = \mathcal{D} \times \mathbf{R}^m \times \mathbf{R}^p$ ,

$$L(x, \lambda, \nu) = f_0(x) + \sum_{i=1}^m \lambda_i f_i(x) + \sum_{i=1}^p \nu_i h_i(x)$$

- weighted sum of objective and constraint functions
- $\lambda_i$  is Lagrange multiplier associated with  $f_i(x) \leq 0$
- $\nu_i$  is Lagrange multiplier associated with  $h_i(x) = 0$

# Lagrange dual function

**Lagrange dual function:**  $g : \mathbf{R}^m \times \mathbf{R}^p \rightarrow \mathbf{R}$ ,

$$\begin{aligned} g(\lambda, \nu) &= \inf_{x \in \mathcal{D}} L(x, \lambda, \nu) \\ &= \inf_{x \in \mathcal{D}} \left( f_0(x) + \sum_{i=1}^m \lambda_i f_i(x) + \sum_{i=1}^p \nu_i h_i(x) \right) \end{aligned}$$

$g$  is concave, can be  $-\infty$  for some  $\lambda, \nu$

**lower bound property:** if  $\lambda \succeq 0$ , then  $g(\lambda, \nu) \leq p^*$

proof: if  $\tilde{x}$  is feasible and  $\lambda \succeq 0$ , then

$$f_0(\tilde{x}) \geq L(\tilde{x}, \lambda, \nu) \geq \inf_{x \in \mathcal{D}} L(x, \lambda, \nu) = g(\lambda, \nu)$$

minimizing over all feasible  $\tilde{x}$  gives  $p^* \geq g(\lambda, \nu)$

# The dual problem

## Lagrange dual problem

$$\begin{array}{ll}\text{maximize} & g(\lambda, \nu) \\ \text{subject to} & \lambda \succeq 0\end{array}$$

- finds best lower bound on  $p^*$ , obtained from Lagrange dual function
- a convex optimization problem; optimal value denoted  $d^*$
- $\lambda, \nu$  are dual feasible if  $\lambda \succeq 0, (\lambda, \nu) \in \mathbf{dom} g$
- often simplified by making implicit constraint  $(\lambda, \nu) \in \mathbf{dom} g$  explicit

**example:** standard form LP and its dual (page 5–5)

$$\begin{array}{ll}\text{minimize} & c^T x \\ \text{subject to} & Ax = b \\ & x \succeq 0\end{array}$$

$$\begin{array}{ll}\text{maximize} & -b^T \nu \\ \text{subject to} & A^T \nu + c \succeq 0\end{array}$$

## Weak and strong duality

**weak duality:**  $d^* \leq p^*$

- always holds (for convex and nonconvex problems)
- can be used to find nontrivial lower bounds for difficult problems  
for example, solving the SDP

$$\begin{array}{ll} \text{maximize} & -\mathbf{1}^T \nu \\ \text{subject to} & W + \mathbf{diag}(\nu) \succeq 0 \end{array}$$

gives a lower bound for the two-way partitioning problem on page 5–7

**strong duality:**  $d^* = p^*$

- does not hold in general
- (usually) holds for convex problems
- conditions that guarantee strong duality in convex problems are called **constraint qualifications**

# Karush-Kuhn-Tucker (KKT) conditions

the following four conditions are called KKT conditions (for a problem with differentiable  $f_i, h_i$ ):

1. primal constraints:  $f_i(x) \leq 0, i = 1, \dots, m, h_i(x) = 0, i = 1, \dots, p$
2. dual constraints:  $\lambda \succeq 0$
3. complementary slackness:  $\lambda_i f_i(x) = 0, i = 1, \dots, m$
4. gradient of Lagrangian with respect to  $x$  vanishes:

$$\nabla f_0(x) + \sum_{i=1}^m \lambda_i \nabla f_i(x) + \sum_{i=1}^p \nu_i \nabla h_i(x) = 0$$

from page 5–17: if strong duality holds and  $x, \lambda, \nu$  are optimal, then they must satisfy the KKT conditions

## Sufficient conditions for KKT at optima:

- **Slater's condition:** for **convex** problems, strong duality holds and KKT holds at optimum if there is a *strictly feasible* point  $x$  where all  $f_i(x) < 0$  (or  $= 0$  if  $f_i$  is affine).
- **LICQ** (linearly independent constraint qualification): even for **nonconvex** problems, KKT must hold at any local optima if the  $\{\nabla f_i\}$  are linearly independent for *active* constraints
  - **Active** constraints:  $f_i(x) = 0$