On Xoodoo

Gilles VAN ASSCHE1

based on joint work with Joan DAEMEN², Seth HOFFERT and Ronny VAN KEER¹

¹STMicroelectronics ²Radboud University

Advances in Permutation-Based Cryptography Milano, Italy, October 2018

Outline

- **1** X00D00
- 2 Trail bounds
- 3 XOOFFF

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What is Xoopoo?



Xoodoo · [noun, mythical] · /zu: du:/ · Alpine mammal that lives in compact herds, can survive avalanches and is appreciated for the wide trails it creates in the landscape. Despite its fluffy appearance it is very robust and does not get distracted by side channels.

Xoodoo



XOODOO cookbook: https://eprint.iacr.org/2018/767 [KECCAK team with Seth Hoffert]

- 384-bit permutation KECCAK philosophy ported to Gimli shape
- Main purpose: usage in Farfalle: XOOFFF
 - Achouffe configuration
 - Efficient on wide range of platforms

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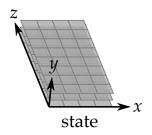
Xoodoo



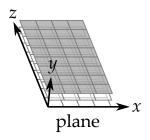
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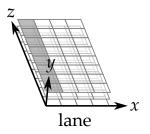
Xoodoo state



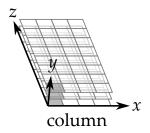
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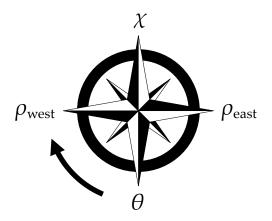
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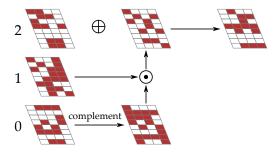
XOODOO round function



Iterated: n_r rounds that differ only by round constant

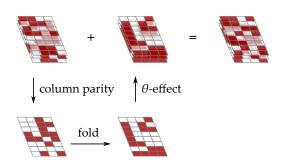
Nonlinear mapping χ

Effect on one plane:



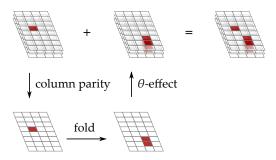
- \blacksquare χ as in Keccak-p, operating on 3-bit columns
- Involution and same propagation differentially and linearly

Mixing layer θ



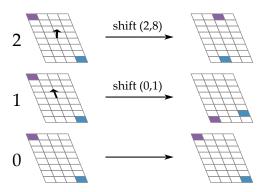
- Column parity mixer: compute parity, fold and add to state
- Good average diffusion, identity for states in kernel

Mixing layer θ



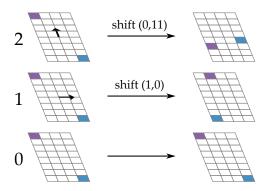
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Plane shift ρ_{east}



- After χ and before θ
- Shifts planes y = 1 and y = 2 over different directions

Plane shift ρ_{west}



- \blacksquare After θ and before χ
- Shifts planes y = 1 and y = 2 over different directions

Xoodoo pseudocode

 n_r rounds from $i = 1 - n_r$ to 0, with a 5-step round function:

$$\begin{array}{l} \theta: \\ P \leftarrow A_0 + A_1 + A_2 \\ E \leftarrow P \lll (1,5) + P \lll (1,14) \\ A_y \leftarrow A_y + E \text{ for } y \in \{0,1,2\} \\ \\ \rho_{\text{west}}: \\ A_1 \leftarrow A_1 \lll (1,0) \\ A_2 \leftarrow A_2 \lll (0,11) \\ \iota: \\ A_{0,0} \leftarrow A_{0,0} + C_i \\ \chi: \\ B_0 \leftarrow \overline{A_1} \cdot A_2 \\ B_1 \leftarrow \overline{A_2} \cdot A_0 \\ B_2 \leftarrow \overline{A_0} \cdot A_1 \\ A_y \leftarrow A_y + B_y \text{ for } y \in \{0,1,2\} \\ \\ \rho_{\text{east}}: \\ A_1 \leftarrow A_1 \lll (0,1) \\ A_2 \leftarrow A_2 \lll (2,8) \end{array}$$

X00D00 software performance

	width	cycles/byte per round			
		ARM Inte			
	bytes	Cortex M3	Skylake		
$KECCAK-p[1600, n_r]$	200	2.44	0.080		
ChaCha	64	0.69	0.059		
Gimli	48	0.91	0.074*		
XOODOO	48	1.10	0.083		

* on Intel Haswell

- XOODOO has slower rounds than Gimli but ...
- ... requires less rounds for equal security objectives!

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- 1 XOODOO
- 2 Trail bounds
- 3 XOOFFF

# rounds:	1	2	3	4	5	6
differential:	2	8	36	≥ 54	≥ 56	<u>≥ 104</u>
linear:	2	8	36	\geq 54	\geq 56	\geq 104

$$w(Q) = w_{rev}(a_1) + w(b_1) + w(b_2)$$

$$\xrightarrow{a_1} b_1 \xrightarrow{a_2} b_2$$

rounds: 1 2 3 4 5 6 differential: 2 8 36
$$\geq$$
 54 \geq 56 \geq 104 linear: 2 8 36 \geq 54 \geq 56 \geq 104

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■ Generating (a_1, b_1)

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- Generating (a_1, b_1)
- Extending forward by one round till weight 50

# rounds:	1	2	3	4	5	6
differential:	2	8	36	≥ 54	≥ 56	<u>≥ 104</u>
linear:	2	8	36	\geq 54	\geq 56	\geq 104

$$w(Q) = w_{rev}(a_2) + w(b_2)$$

$$\xrightarrow{a_2} b_2$$

■ Generating (a_2, b_2)

# rounds:	1	2	3	4	5	6
differential:	2	8	36	≥ 54	≥ 56	<u>≥ 104</u>
linear:	2	8	36	\geq 54	\geq 56	\geq 104



- Generating (a_2, b_2)
- Extending backward by one round till weight 50

rounds: 1 2 3 4 5 6 differential: 2 8 36
$$\geq$$
 54 \geq 56 \geq 104 linear: 2 8 36 \geq 54 \geq 56 \geq 104

$$w(Q) = w_{rev}(a_1) + w(b_1) + w(b_2)$$

$$\xrightarrow{a_1 \quad b_1 \quad a_2 \quad b_2} \lambda \xrightarrow{b_2} \lambda$$

Extending all 3-round trail cores to 6 rounds till weight 102

Using the tree-search approach

Set U of units with a total order relation \prec

Tree

■ Node: subset of U, represented as a unit list

$$a = (u_i)_{i=1,\ldots,n}$$
 $u_1 \prec u_2 \prec \cdots \prec u_n$

Children of a node a:

$$a \cup \{u_{n+1}\} \quad \forall \ u_{n+1} : u_n \prec u_{n+1}$$

■ Root: the empty set $a = \emptyset$

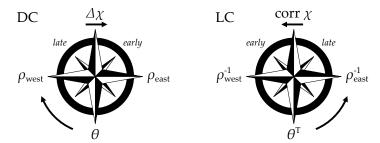
[Mella, Daemen, Van Assche, FSE 2017]

Definition of units

Units represent one bit at a time:

- \blacksquare Active bit in odd column (x, y, z)
- Bit in affected column (x, y, z, value 0/1)
- Active bit of an orbital (x, y, z)
- ⇒ allows for finer-grained bounding

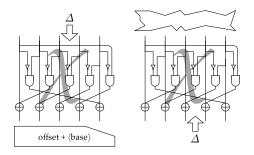
Properties of the trail search



Difference and mask propagation in χ follow the same rule \Rightarrow differential and linear trail search are almost identical

Properties of the trail search

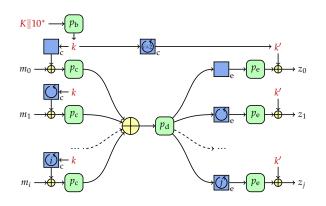
Compared to trail search in Keccak-p:



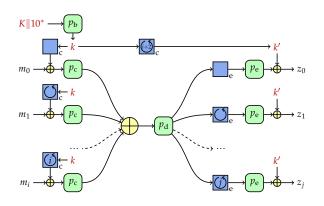
In Xoodoo, both χ and χ^{-1} have algebraic degree 2 \Rightarrow affine-space extension in both directions

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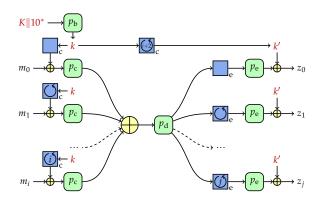
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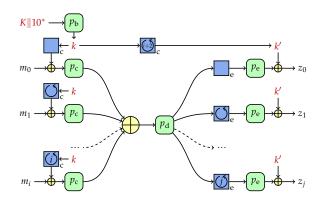
- $p_b = p_c = p_d = p_e = X00000[6]$
- Input mask rolling with LFSR, state rolling with NLFSR
- Target security: 128 bits, incl. multi-target and quantum adv.



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XOOFFF applications and implementations

The Xoopoo Cookbook also specifies:

- XOOFFF-SANE: session AE relying on user nonce
- XOOFFF-SANSE: session AE using SIV technique
- XOOFFF-WBC: tweakable wide block cipher

KECCAK Code Package

↓

eXtended KECCAK Code Package

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KECCAK Code Package ↓ eXtended KECCAK Code Package

Any questions?

Thanks for your attention!

■ More information https://eprint.iacr.org/2018/767

■ Some implementations

```
https://github.com/XoodooTeam/Xoodoo/ (ref. code in C++ and Python)
https://github.com/XKCP/XKCP (C, Assembler)
https://tinycrypt.wordpress.com/2018/02/06/... (C, Assembler)
```