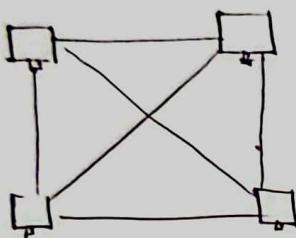


ARPA NET

—agen advanced research



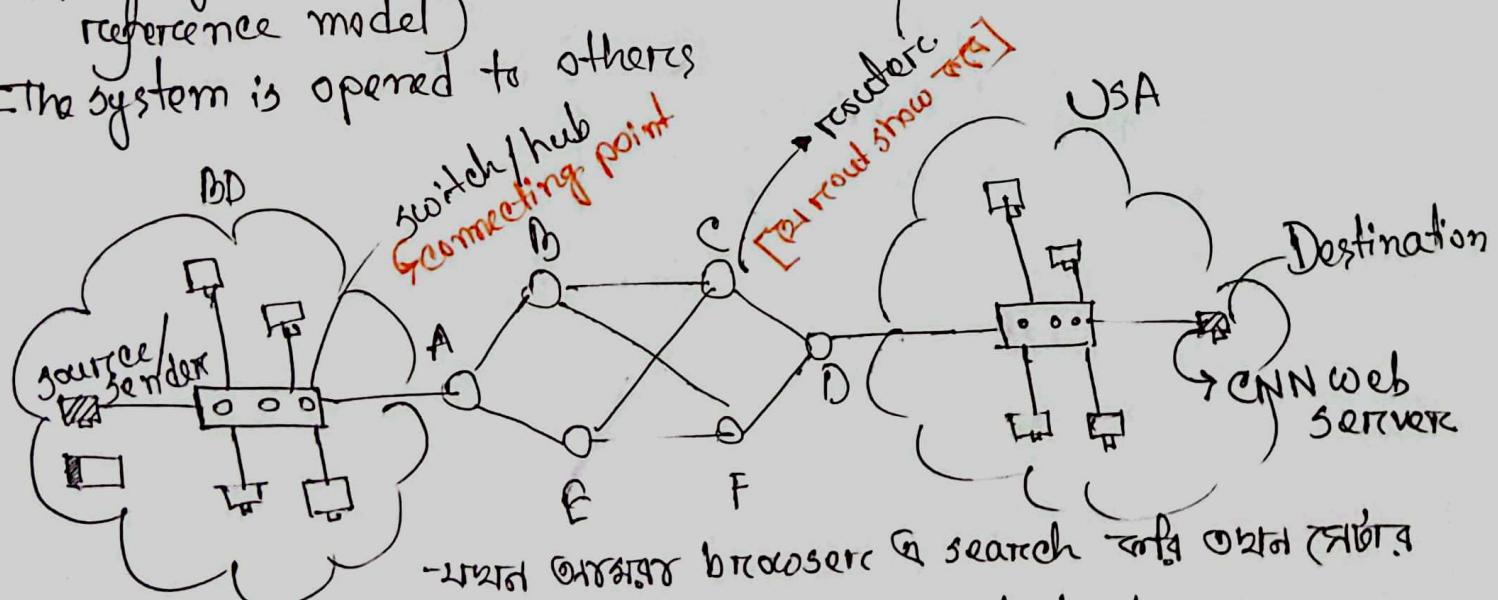
standard → O₃[↑] ref. Model

ଆପ୍ଲିକେସନ୍ ମ୍ୟୁ ସାଥେ
communicate କରିବାର ପାଇଁ
(open system interconnection
reference model).

The system is opened to others
hub point

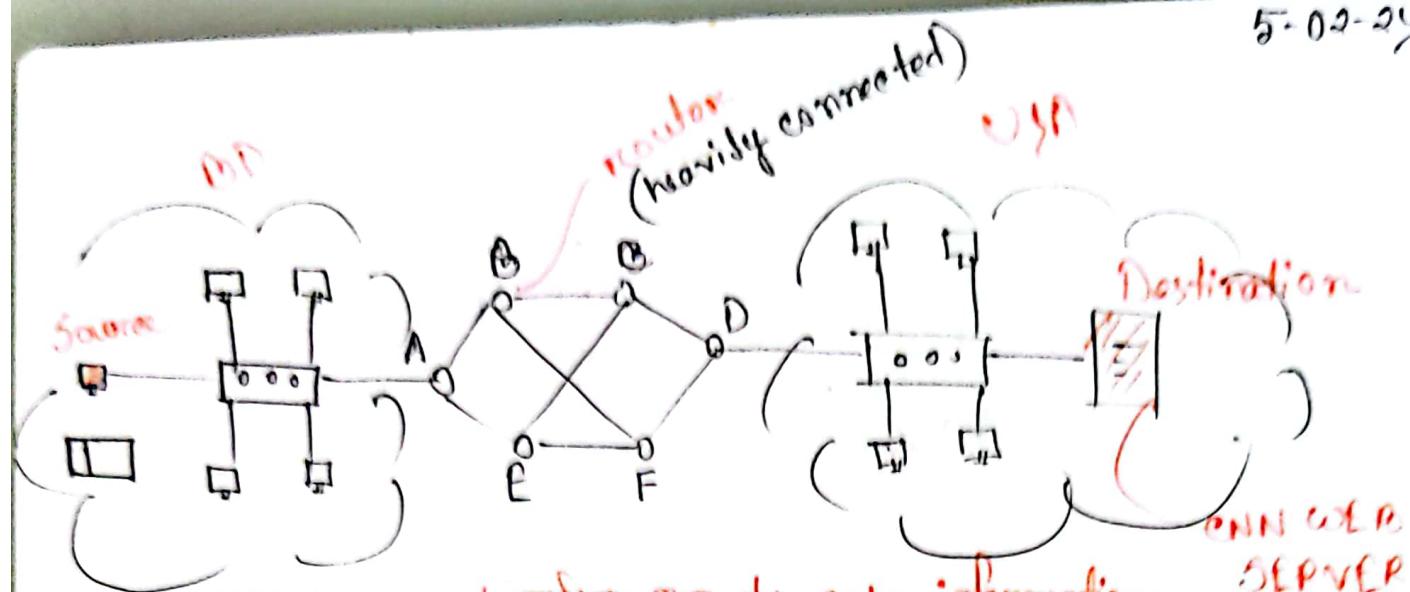
এখন সবাই same algorithm use করত-
ওজনে communicate করতে পারে।
অসম টেলিভিশন ইন্টেরনেট

7 routers ଥିଲେ heavy



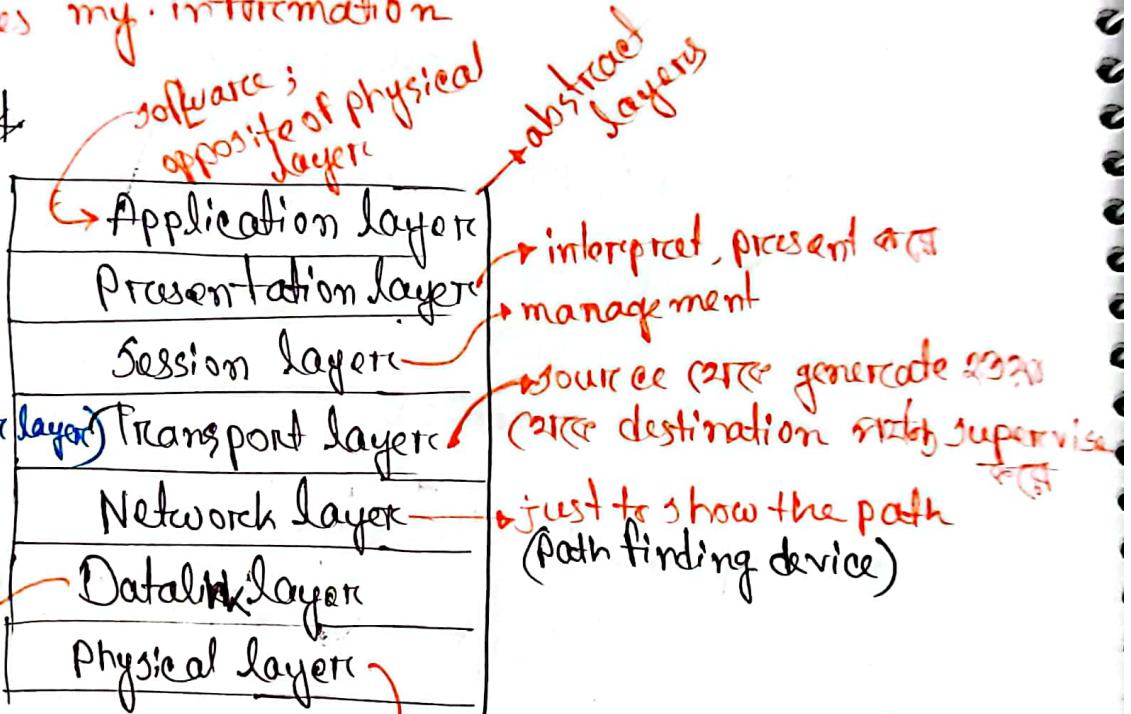
- যখন আপনার browser-এ search করলে তখন সোচ্চার webpage এর address সহ website show করে।

आनान्दग्रन्थाना connection के एक point a connect करते हैं तो switch ; it works like a multiplexer



- ~~DATA request~~ ~~DATA~~ goes to gain information;
- ~~DATA request~~ to generate ~~DATA~~ ~~DATA~~ send ~~DATA~~ ~~DATA~~
- Router carries my. information

Standard
ISO Ref. model
(Layers & layers)

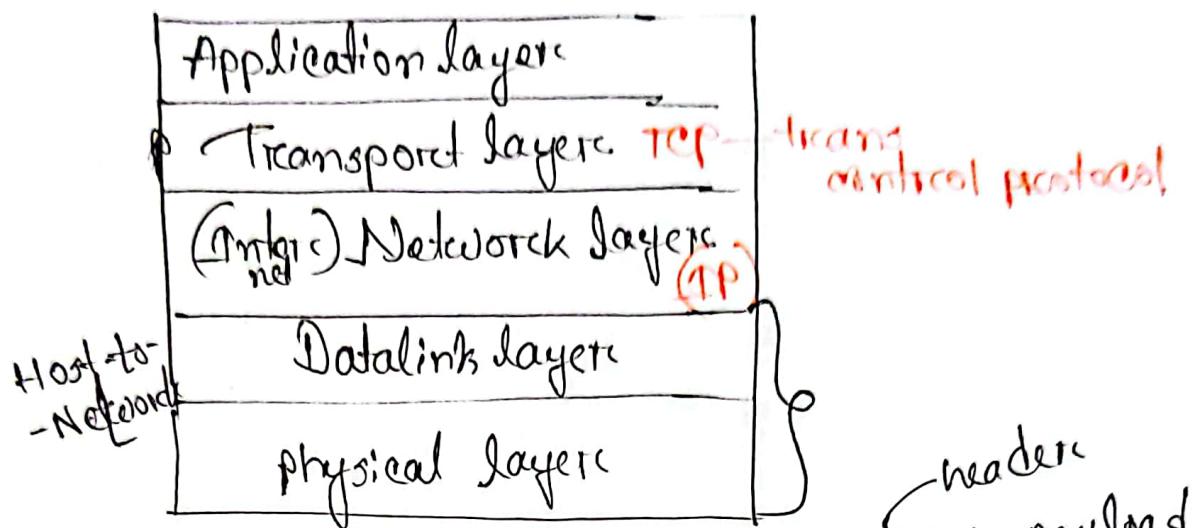


* adjacent devices
* error-free data delivery

* data link ensures
error-free networks to
adjacent devices communication

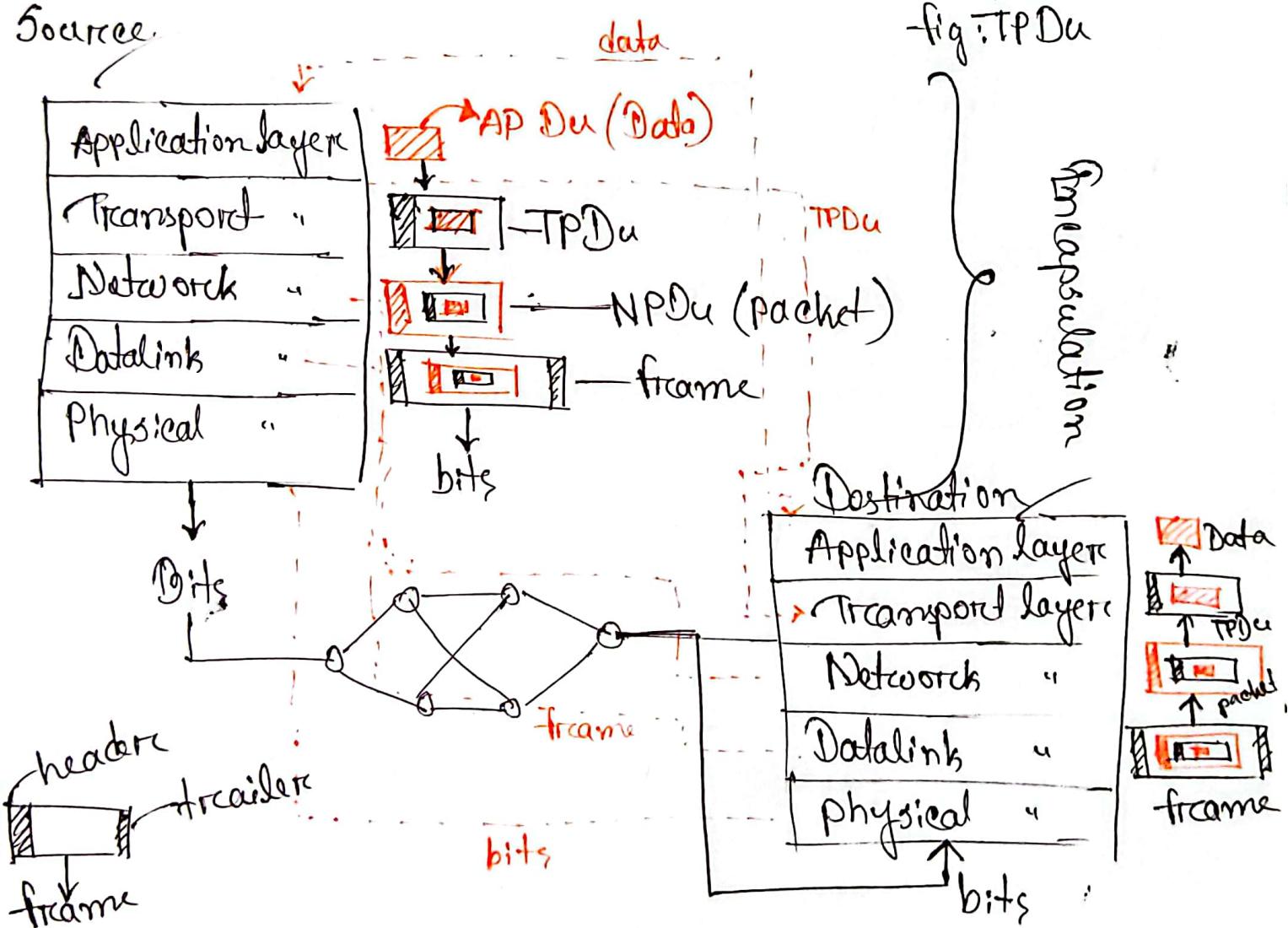
* one device multiple layers as per requirement

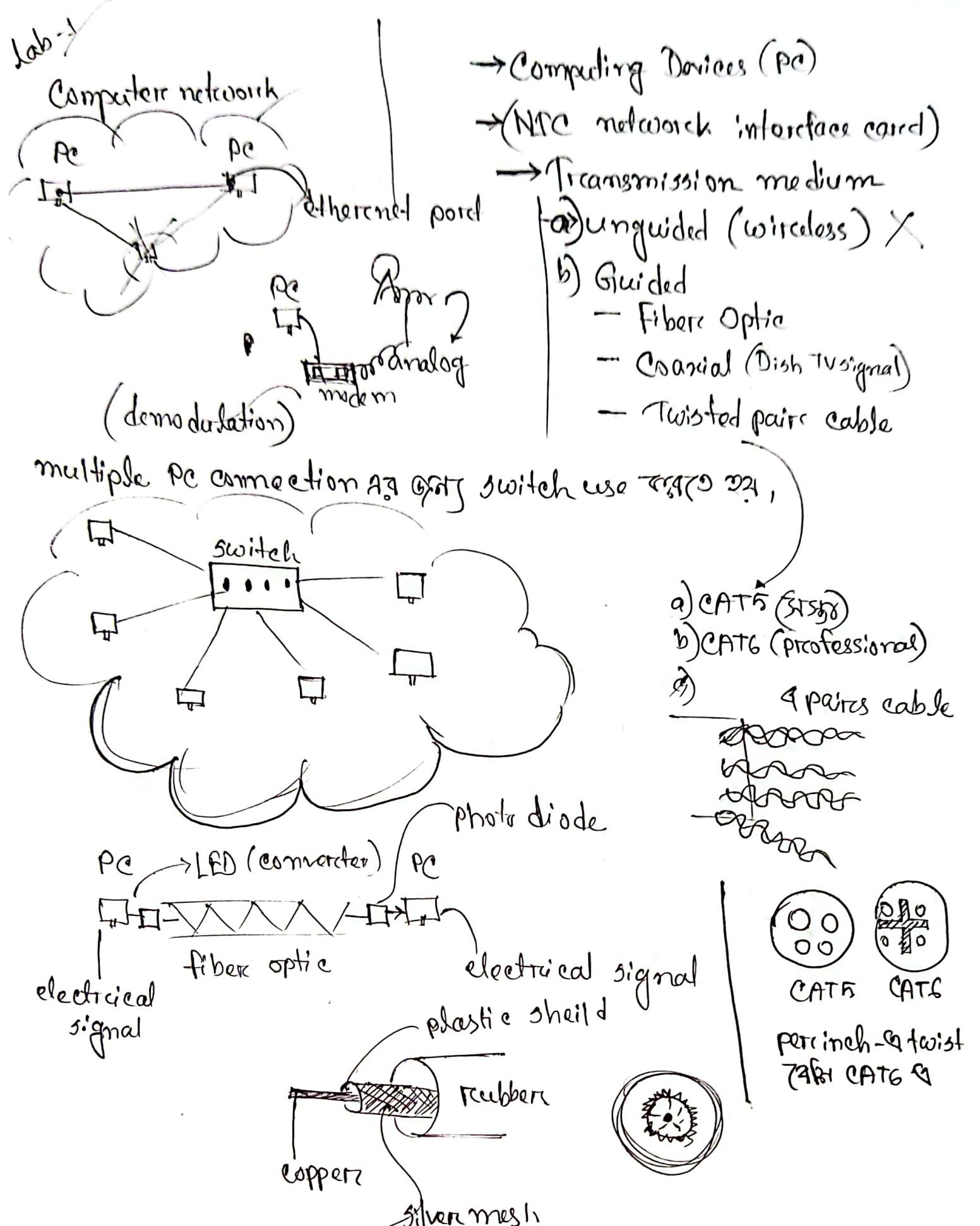
TCP/IP Ref. model *smaller/derivative form of OSI*



Deep Encapsulation

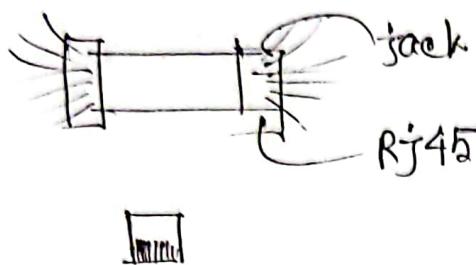
Source:





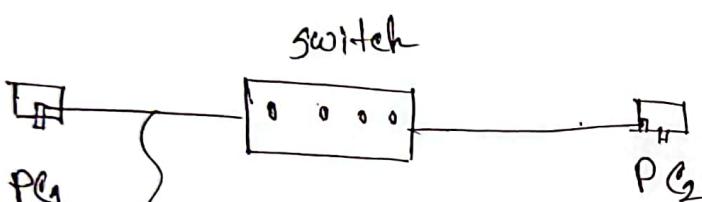
Twisted

- 4 solid colors } Green, blue, orange, brown }
- 4 Mixed with white } white - Brown
white - Green
white - Blue
white - Orange }



Next lab
CAT5 - 2m
RJ45 - 5/6 pos

Crimping tool



straight
through cable

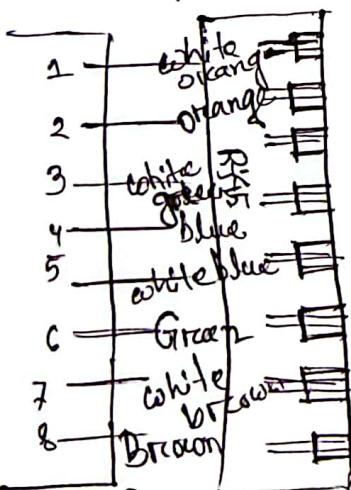
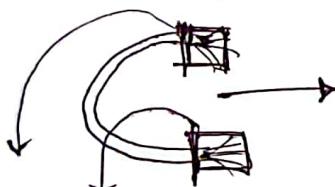
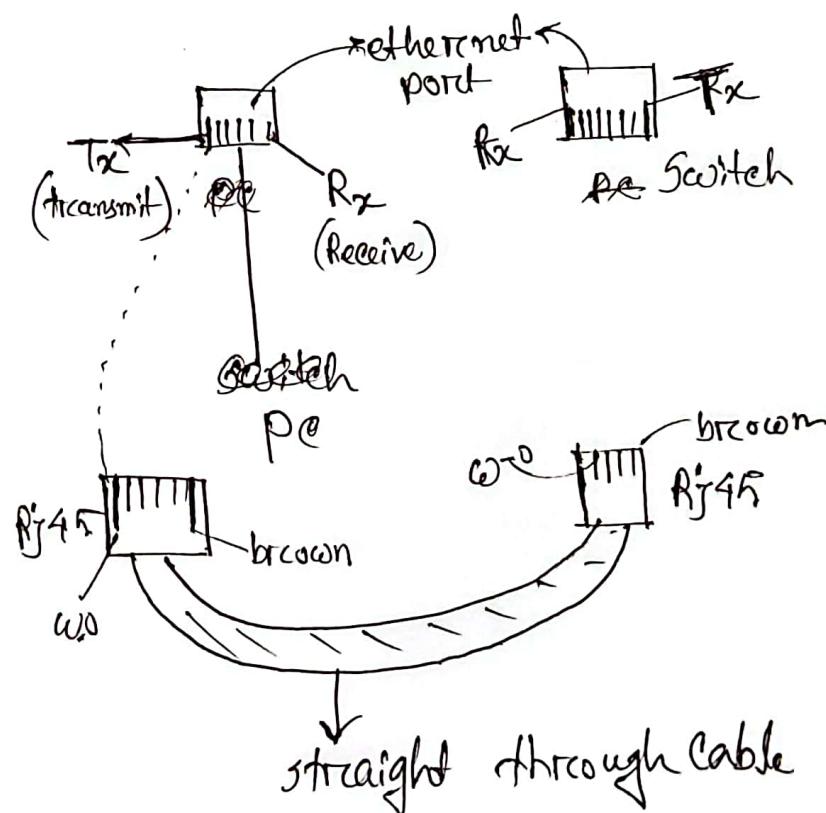
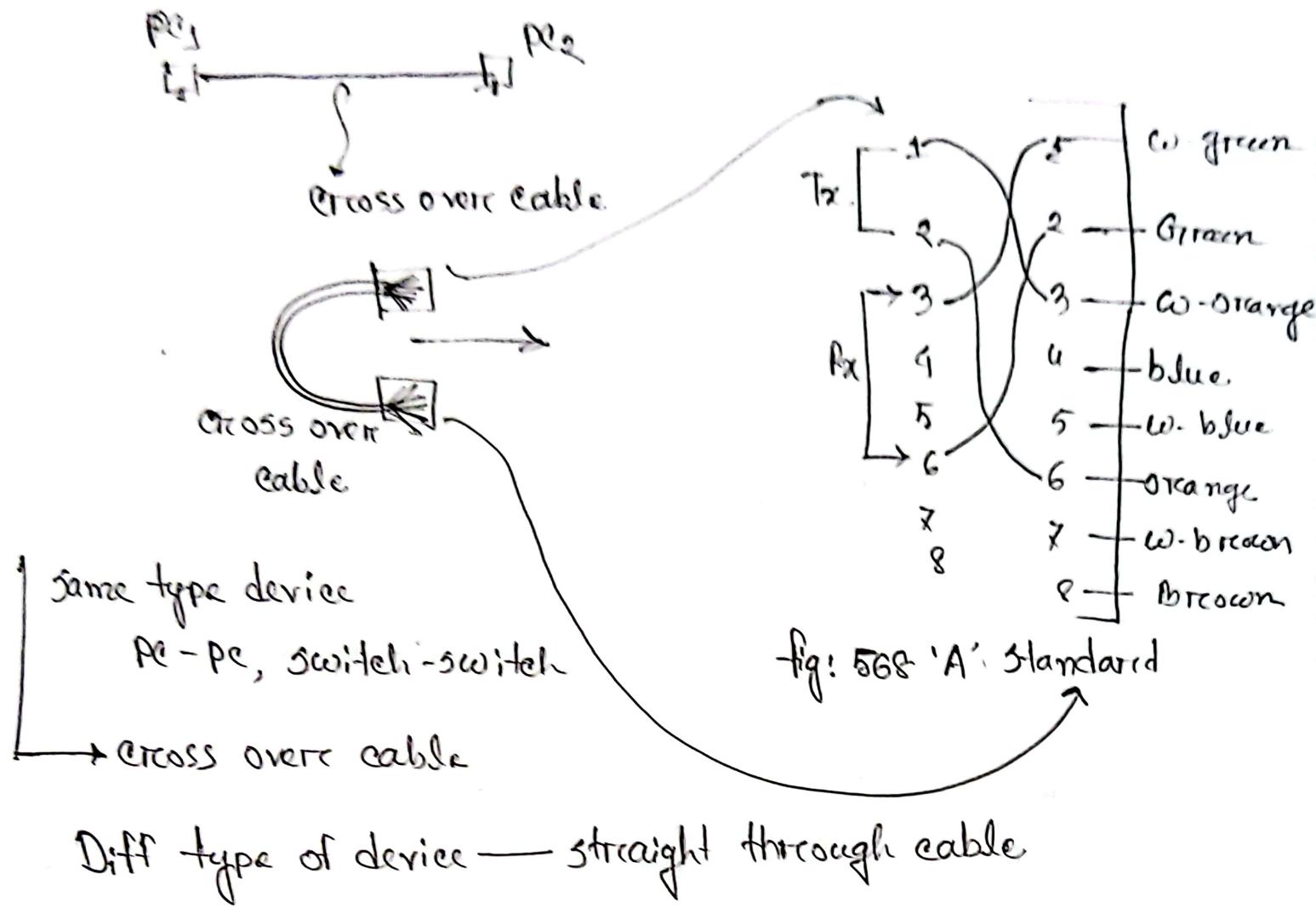
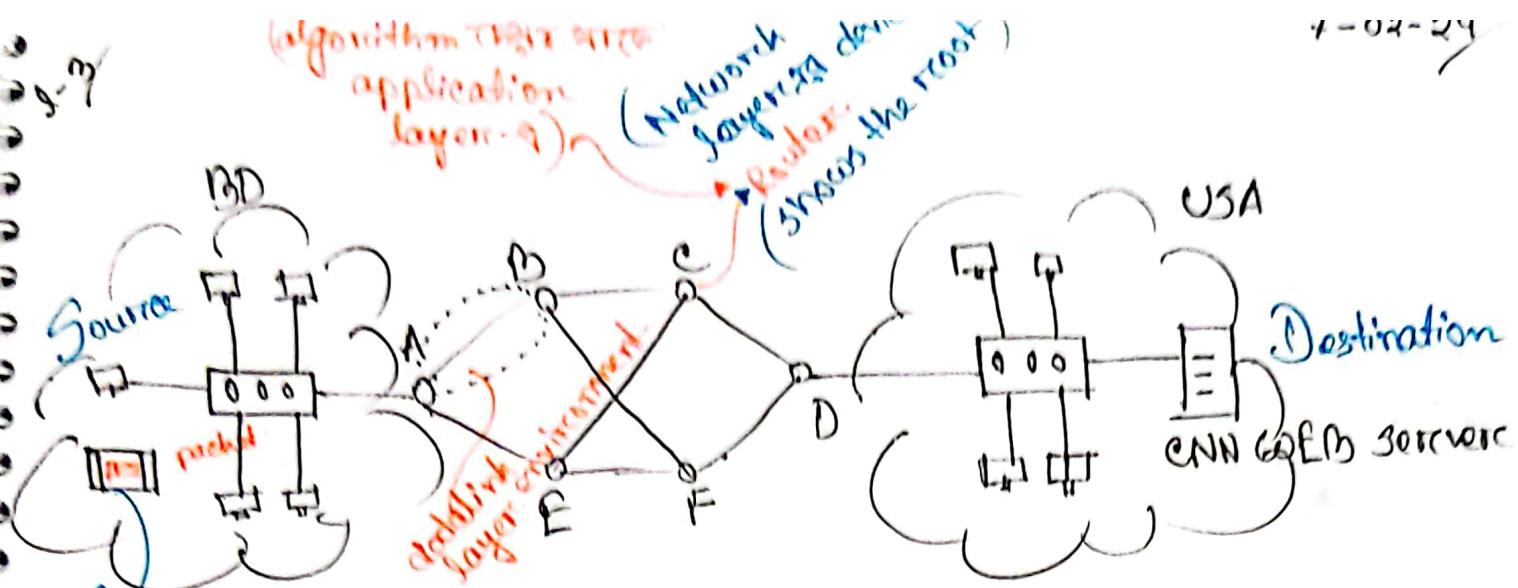


fig: 568'B' standard







এই frame এর মধ্যে bit
পালনে data গাঠন করে next device (switch → router ... → destination)

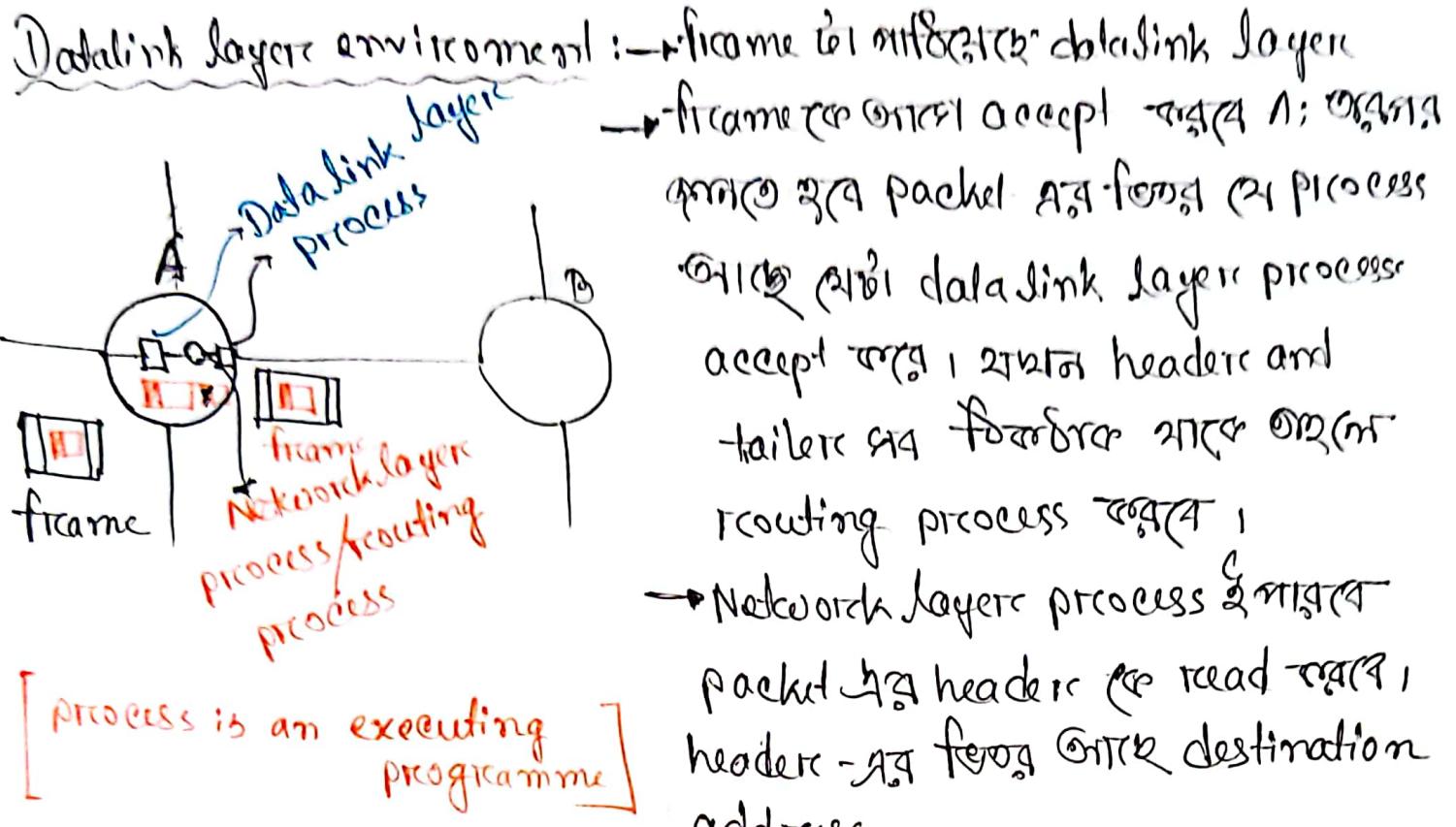
Router define করুন ইন্টের্নেট → physical layer এ

কার্ড কলেজ ইন্টের্নেট → Network layer এর মত

algorithm করুন তাহে → application layer এ

Packet - এই packet এর header-তের মধ্যে আছে **destination as address**
Destination address (IP address)

frame এখন A এবং কলেজ থাবে, A এর দ্রুততা ওয়ে frame এর
address-ের কি? address কোম্পিউট এ ক্লেক্সিমাপ্প অথবা algorithm
করে route করাবে; Packet এর header করুণ পারে
Network layer। এ packet device to device travel করার পর
destination-এ কোথায়ে, অথবা কোন route/ link-এ ফের ফের
আবে ও depend করুণ packet এর destination এর address এর পরে
A থেকে B টো যে environment create হবে. (for data link layer)
environment.



* frame A has to change its address before reaching host B.

Framing techniques

- i) Character count
- ii) Byte stuffing
- iii) Bit stuffing
- iv) Physical layer coding violation

Character count

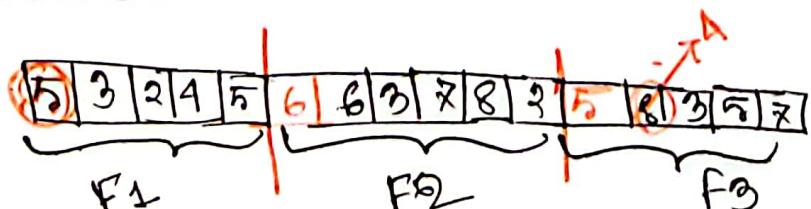


total count + character

$$= 13 + 1 \\ = 14$$

→ all technique follow same rule
 total 14 for character read
 next router - 1 byte, ...

error-free solve ??



Bhared → creates LAN

→ frame: 5460 Byte (max)

ATM Network (asynchronous ...)

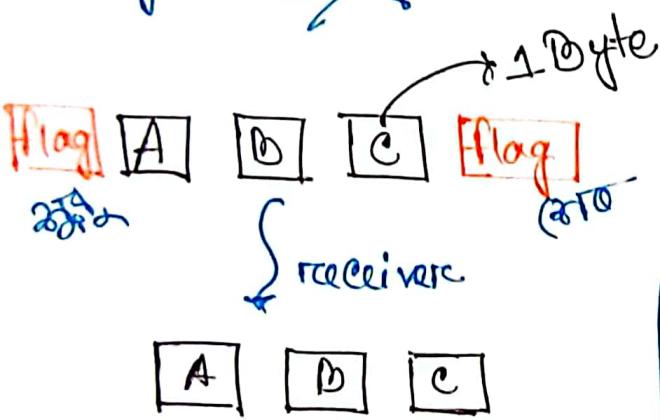
backbone use ATM

frame (cell): 53 Byte (max)

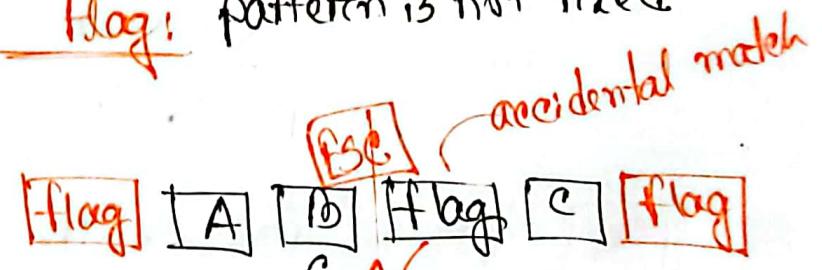
Sender ৰাখিবলৈ receiver read কৰিবে, receiver ৰেad

কৰিবে (চন্দ্ৰ কম্পুটা) [Router] এখন net ৰেad কৰিবে। ৰেad ৰেad কৰিবে। (১) read কৰিবে, ফিলু একোৱ দুই character ৰেad কৰিবে তাৰে flag ও দুইটা পাঠ না।

④ Byte stuffing আপে লিহৈ flag insert



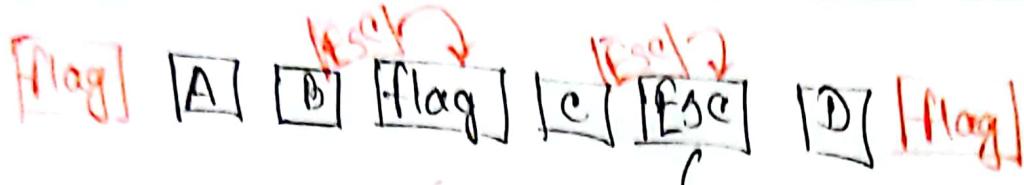
Flag: pattern is not fixed



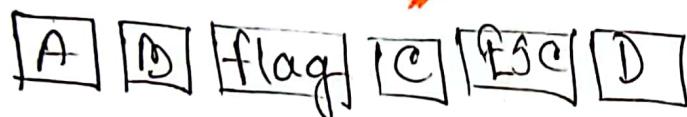
ESC = you should skip it

অথবা flag নহ' আপি (ESC) character ব্যৱ কৰিবলৈ দুইটা গুৰুত্ব মানেৱ flag (ৰেad count কৰিব),

starting [A] [B] flag [C] ending flag



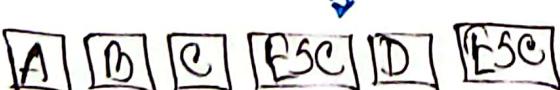
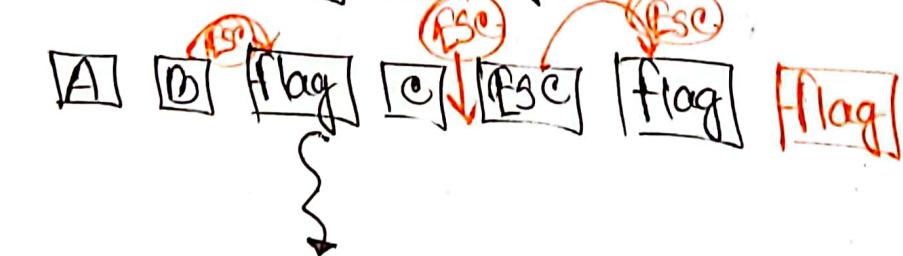
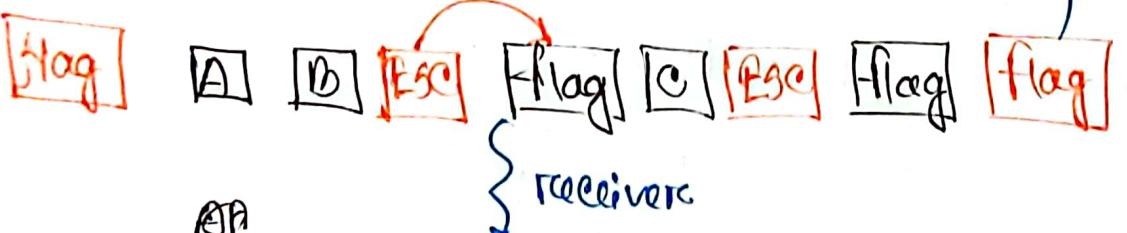
receiversকে যাতে ক্ষয় না থান্তে ESC মার্কিং data. এর প্রয়োগের
জন্মে [ESC] (skip character) হিঁতে হবে.



প্রয়োগের [ESC] হিঁতে হবে.

not fixed

{-flag : 10101010
ESC : 00001111}



Bit Stuffings

[Flag] 1010111010011 [Flag]
 (Data)

- * flag fixed
- * pattern fixed
- flag (fixed): 0111110

0111110 1010111010011 0111110

 (Receiver)

1010111010011  0111110
 (Data) 

0111110 1010110101111010 0111110
 flag  flag লাঠিকার আপে দেশকু
 address কর্তৃত হয়

1010110101111010
 instruction part
 Sender

after every 5 consecutive 1's,
 insert a '0'

আপে - শিখ flag insert করার ↑

receiver instruction

flag read কর

- 5 consecutive 1 এর পরে 0 করা পারি কর্তৃত হবে

0111110 1010110101111010 0111110

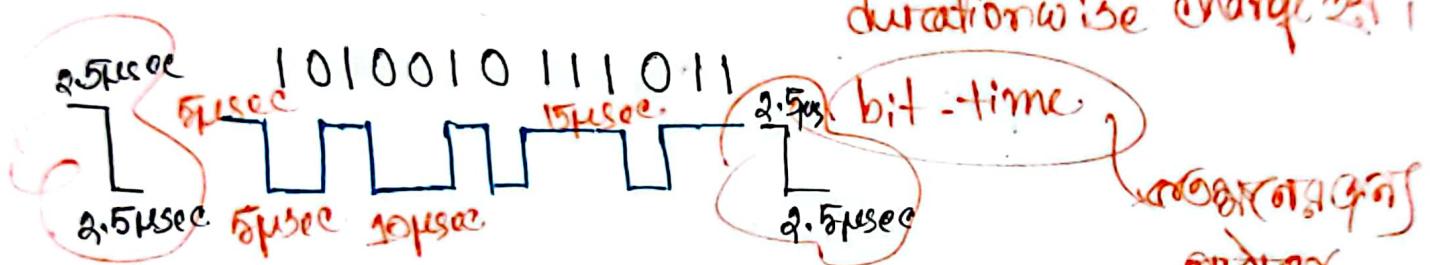
 receiver

1010110101111010 0111110

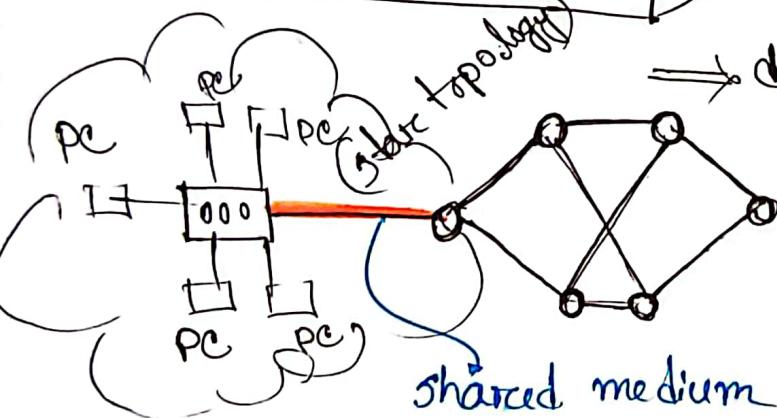
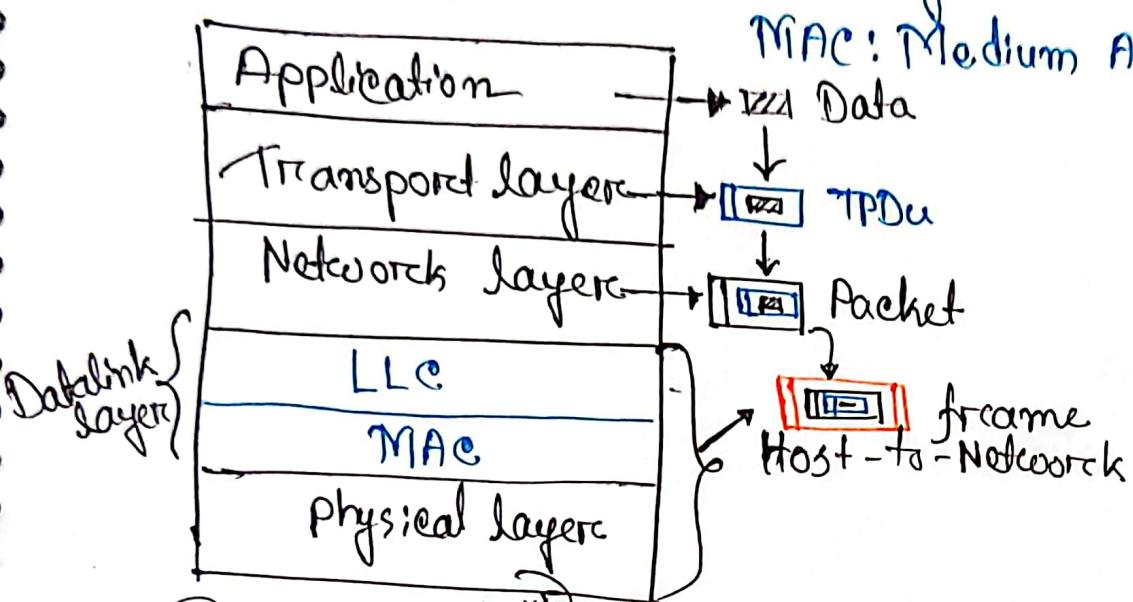
দেখানো Data কর্তৃত 11111 থাকবেন।

01111110 10101101011111001010 | 0111110
(flag) { (flag)

Physical layer Coding violation



lec-5/

TCP/IP Ref. Model

LLC: Logical Link Control sublayer
MAC: Medium Access Control sublayer

Data

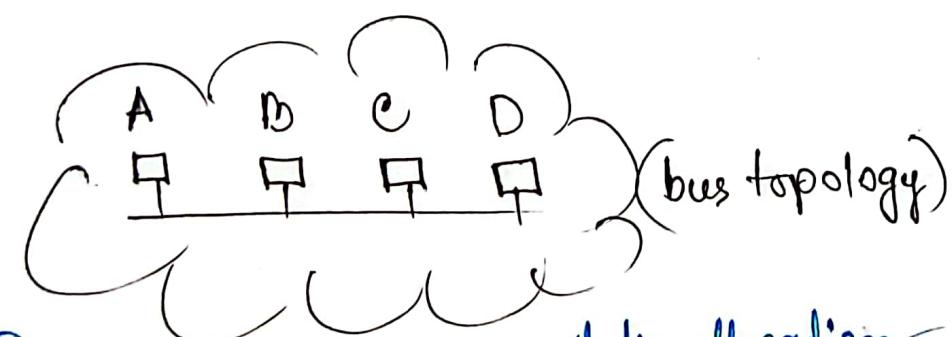
TPDU

Packet

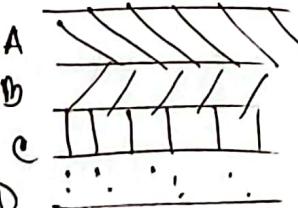
frame

Host-to-Network

datalink layer গুরুত্বপূর্ণ device থেকে
জন্য device → data frame শারী
কর্তৃত এবং প্রয়োগ নির্দেশনা
জন্য control access থাকে যেটা
জন্য MAC .

Dynamic allocation

A → অধিকারী
B → সম্মত
C → উন্নত
D → প্রাপ্তি
সাধাৰণ

Static allocation

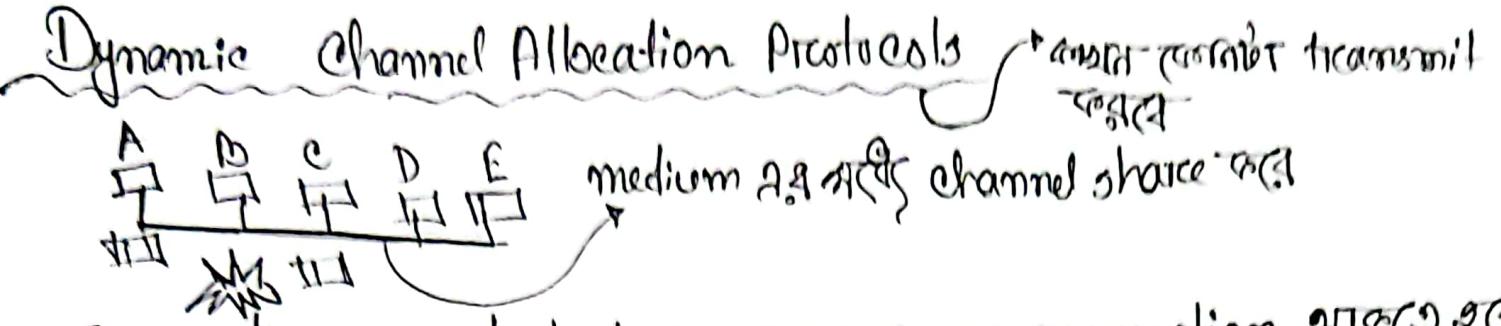
এভোবো জনক data পাঠাব

পাইলে, specific ক্ষমতা

মাঝে ; ওই এটা static

allocation (জন্য ক্ষয়ে দেখা)

It is not very useful



Assumptions: practical design কর্তৃপক্ষ উদ্দেশ্যে assumption আছে।

- i) Station Model (There are ~~n~~ sum of independent stations/nodes (integers))
 - ii) Single channel
 - iii) Collision Time [continuous data এবং you can transmit]
 - iv) Collision [Slotted]
 - continuous data এবং you can transmit
 - Slotted
 - v) Sense [frame এবং turn এবং unit এবং]
- (Carrier sensing capability)

- ALOHA → মানব ধারণা মানবে ;

+



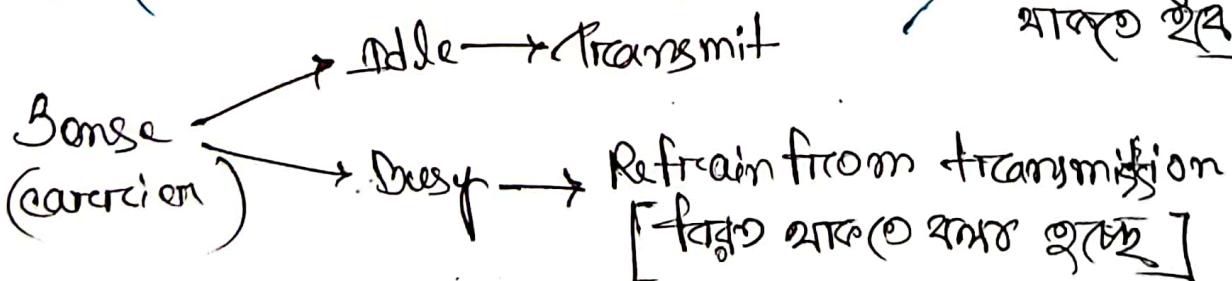
অন্য সে data মাঠাব। অন্য data থেকে :

মানব অন্য ক্লাউড এবং ক্লাউড মাঠে collision
শুধুই হবে। ক্লাউড collision হলে ক্লাউড
collision হওয়ার possibility এখন থাকে।

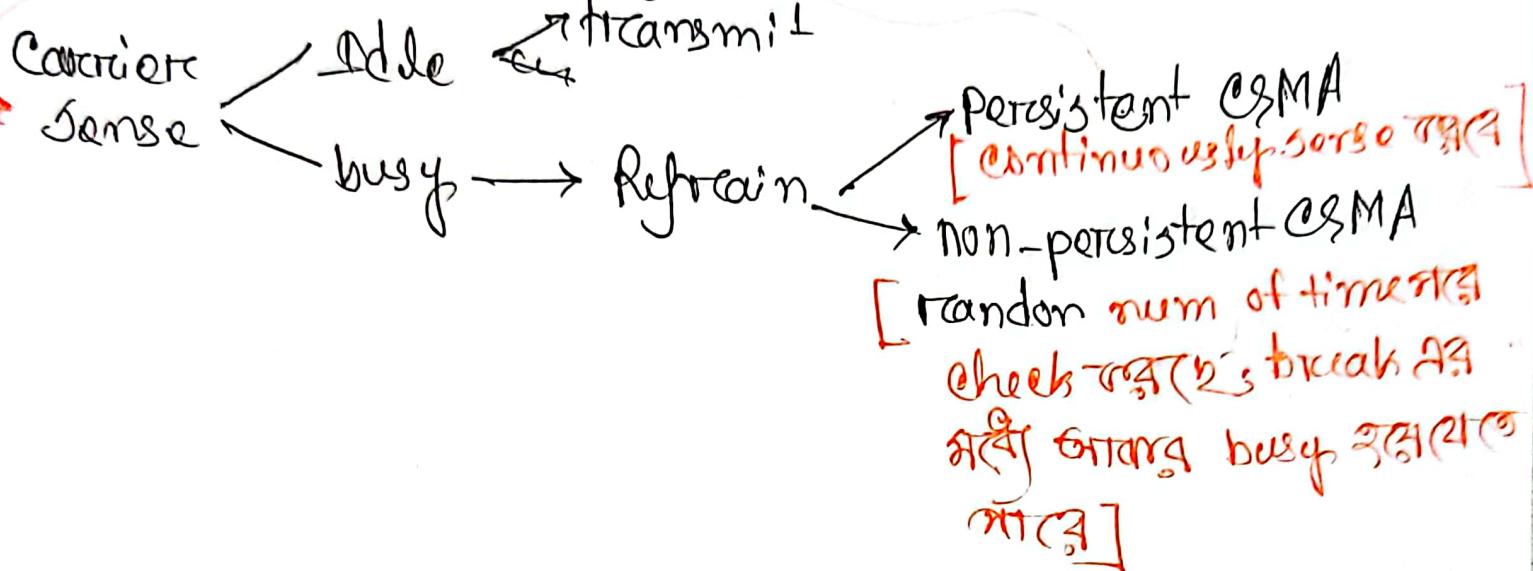
অন্য free slot মানুষের মানুষের হওয়া না।

→ Slotted ALOHA number of nodes এখন মানবে এটি preferable

- CSMA (carrier sense multiple access) [Sensing capability মানব ইন্ডিকেশন]



Idle হলে collision ঘটার chance কমাওয়ার ব্যবস্থা transmit
করতে পারবে, busy হলে retransmit করতে পারবে।



20-04-24

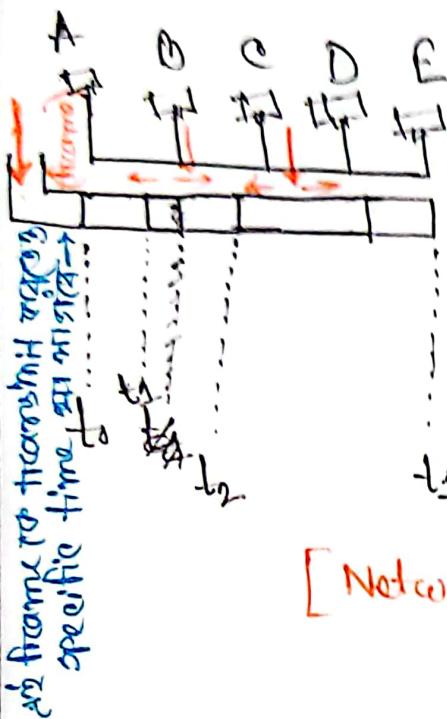
lec 6 CSMA (Carrier Sense Multiple Access) Protocol

[Collision chance ০%]

Sense → Idle → Transmit

Sense → Busy → Refrain from transmission.

→ এই ক্ষেত্রে ২৩ টাইম নিয়ে data ও
সেন্স ফলো ফিল্ড কো



[Collision হলি]

কলিজন ঘটেছে (১)
frame আকার
কলিজন মূল্য
চালনা দরকার
কলিজন হলি

→ At 1st time → bit B (১)
কলিজন হলি B (১) idle
পার্শ্ব পার্শ্ব

→ Sense ক্ষেত্রে দুর্দেশন
signal আবেক্ষণ কো, if there
is going any current is
going on.

at time = t₂ → C,D,E idle; transmit কো

ফলো কো busy নিয়ে - that's why

B can't send any data.



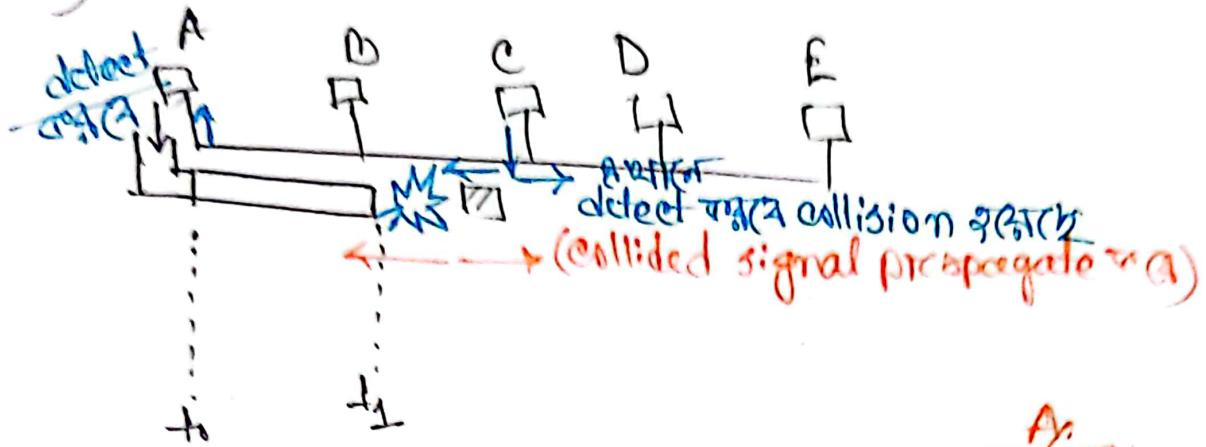
→ CSMA কে collision হও গালি, (tiny window অবৈ)

by detecting current
or voltage in the carrier

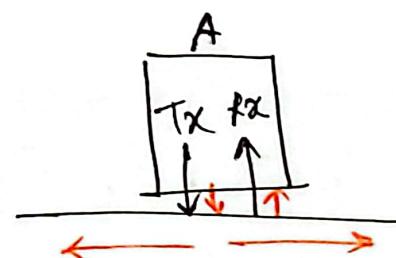
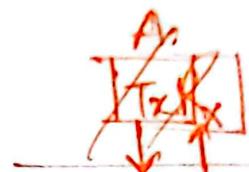
at time t₁ কে পিলে (পৌছান কোরে আছে)

collision হও না ; একজন দরকার থাকে "data নির্দেশ করে হও নেতৃত্বে"

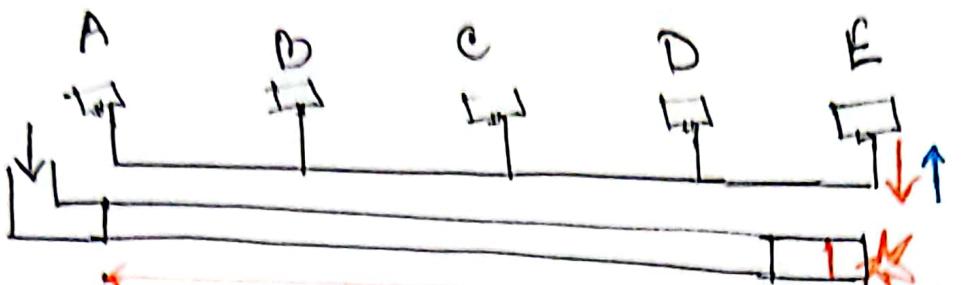
CSMA/CD (Carrier sense multiple access with collision detection) ?
 (better than CSMA)



Transmitted signal and received signal compare কোনো তথ্য
 same হলে no collision
 different হলে collision



মনি transmitted signal = received signal
 \therefore collision হলে
 মনি equal এবং হলে
 collision হবে



collide as signal A (A)

ପାଇଁ ଅନ୍ତର୍ଗତ ଏ କ୍ଷେତ୍ର

collision ହେଲିବା

ϵ (sliding value)

$$(t_f - \epsilon) = (t_f - \frac{\epsilon}{n_E})$$

$t_2 - t_1$ (ପାଇଁ କେବେଳିକିମିଟ
ଦୂରବଳା)

- $t_f - \epsilon$ is the last possible time of be collision

ଯେତେ ମନ୍ତ୍ର + ଆମତେ ମନ୍ତ୍ର \Rightarrow ଫର୍ତ୍ତେ collision ହେଲିବା ହେଲା

$$(t_f - \epsilon) + (t_f - \epsilon)$$

A (ଏ କେବେଳିକିମିଟ ମନ୍ତ୍ର)

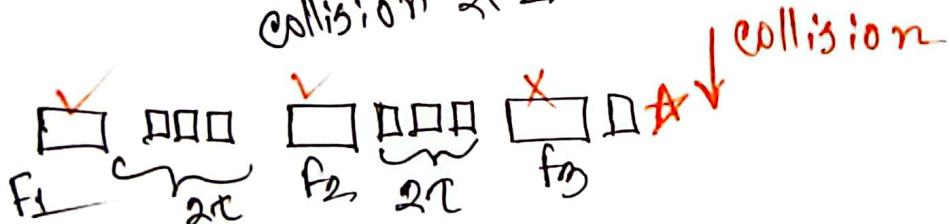
$$= 2t_f - \epsilon = 2t_f = 2T_{A-E} \quad | \begin{matrix} A-E = \text{farthest node} \\ = \text{Contention Period} \end{matrix}$$

Propagation delay = T_{A-E} (ପ୍ରତିକାଳିକା କାଲୀମାଣ)

$$\text{Contention period} = 2T_{A-E}$$

[ଏହି ମନ୍ତ୍ରର ପାଇଁ collision
ହେଲିବା କାହିଁବା ହେଲା
ଦୂରବଳା]

frame c $\square \square \square \square$
ଘରେ ପାଇଁ କାଲୀମାଣ ହେଲା
ଘରେ ପାଇଁ କାଲୀମାଣ ହେଲା
ଘରେ ପାଇଁ କାଲୀମାଣ ହେଲା



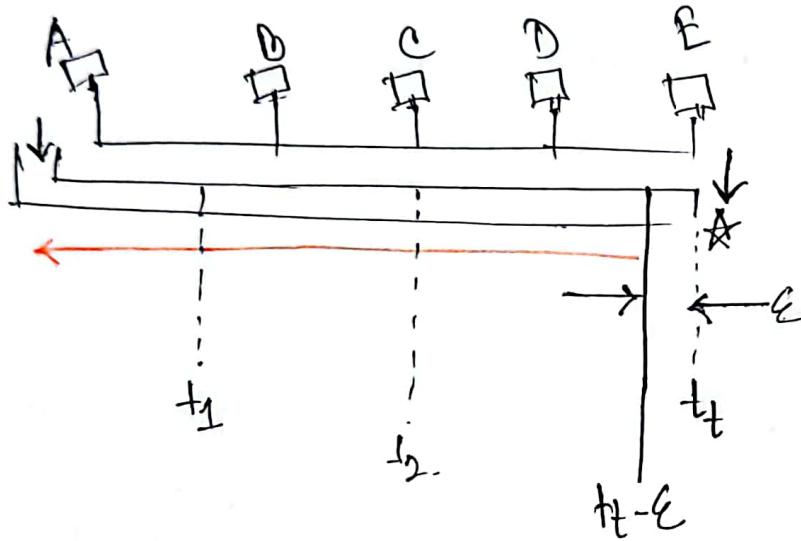
~~Sec 7~~ / CSMA/CD is always designed as half duplex



1: propagation Delay

* How it detects collision?

* Contention period = $2T$



Destination C দূরত্ব

সত্ত্বেও signal A (মধ্যে E
পর্যন্ত থাই, (here every
elements are covered)

ঠিক মানের (contention
period) এলাজে প্রদৰ
propagation delay

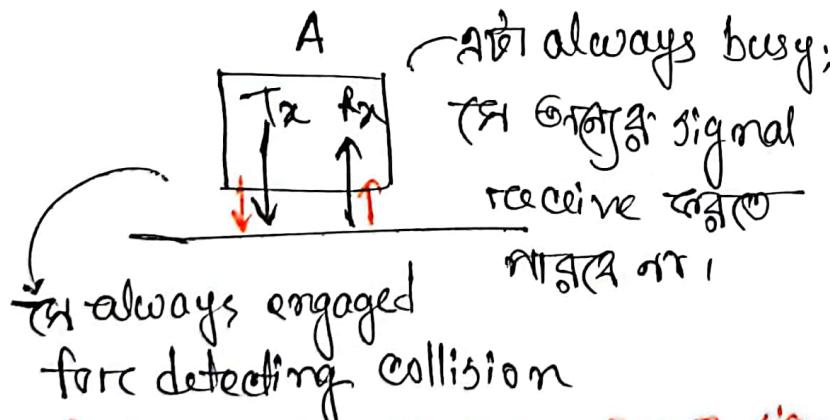
যোগাযোগ

$$\therefore \text{contention period} = 2T_{A-E}$$

$$= 2T_{A-C}$$

between 2 farthest nodes

যাইতু আনন্দের exact
position গুণিত প্রদৰ
medium এর propagation
time।



2nd full duplex কোথা থাই
কোথা থাই
কোথা থাই
কোথা থাই



collided signal যেকোন medium-এ propagate.
ক্ষয়ে, collision হবে।
তারা ক্ষেত্র-ক্ষেত্রে signal
পাঠাবে ক্ষয় ক্ষয়ের আপেক্ষ
সব frame এর উপর দাও।

যখন কোন signal পাঠাবে।

collision হল ক্ষেত্র, তখন BEB apply
যখন কোন ক্ষেত্র। (multiple collision অব্যুক্ত)
BEB (Binary Exponential Backoff) Algorithm :

collision হলেই মানে frame এর ফিল্ড গুরুত্বের সময়ের ক্ষেত্রে ক্ষেত্রে ক্ষেত্রে।

যখন এই algo run ক্ষেত্রে অথবা গুরুত্ব ক্ষেত্রে।

Set = $\{0, \dots, 2^n - 1\}$ n: # of collision (counter)

(pick an element from the set randomly) \times (512 bit time)
= waiting time

For A,

Collision # 1

$$\text{Set} = \{0, \dots, 2^1 - 1\} = \{0, \dots, 1\} = \{0, 1\}$$

0×512 Bit time = 0 time (waiting time doesn't need to wait
ক্ষেত্রে সার্বত্র transmit করো)

1×512 Bit time = 512 Bit time (waiting time)

Forc D, collision #1

$$\text{Set} = \{0, 1\}$$

0×512 Bit time = 0 time (waiting time)

1×512 " " = 512 Bit time ("")

50% chance for next collision — worst case.

~~Forc A~~

Collision #2 ($n=2$)

$$\begin{aligned}\text{Set} &= \{0, \dots, 2^m - 1\} \\ &= \{0, 1, 2, 3\}\end{aligned}$$

0×512 Bit time = 0 time (waiting time)

$$1 \times 512 = ?$$

$$2 \times 512 = ?$$

$$3 \times 512 = ?$$

forc D

Collision #12 ($n=2$)

$$\text{Set} = \{0, 1, 2, 3\}$$

$0/1/2/3 \times 512$ Bit time = ? (waiting time)

Others same num picks \rightarrow 50% chance 50%.

~~Forc A~~

Collision #13 ($n=3$)

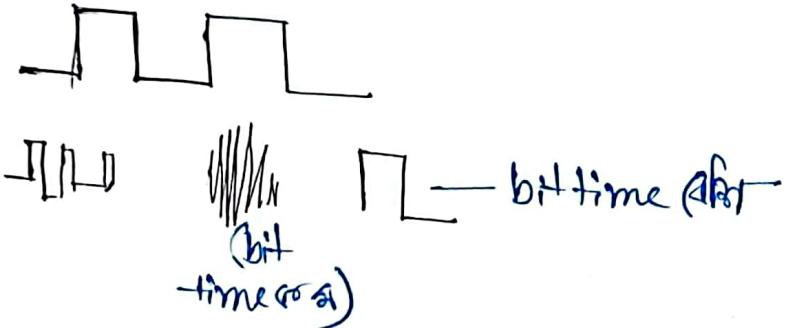
$$\begin{aligned}\text{Set} &= \{0, \dots, 2^3 - 1\} \\ &= \{0, 1, 2, 3, 4, 5, 6, 7\}\end{aligned}$$

$0/1/2/3/4/5/6/7 \times 512$ Bit time = ? (waiting time)

Others same num picks \rightarrow 50% chance 12.5%.

10th collision * consecutive 10 ସବୁ collision ହେବୁ ପ୍ରକାର
ଏହି ଚାଲୁକୀ ଡର୍ଜ କରିବାକୁ ପାଇଁ ,
* ୨ଟିମେତ୍ର ଯେ ଏକାଏକ ଅତି ଶାଖା

4-collision କ୍ଷେତ୍ରର time = $t_1 - \epsilon$



Lab-03 Simulation

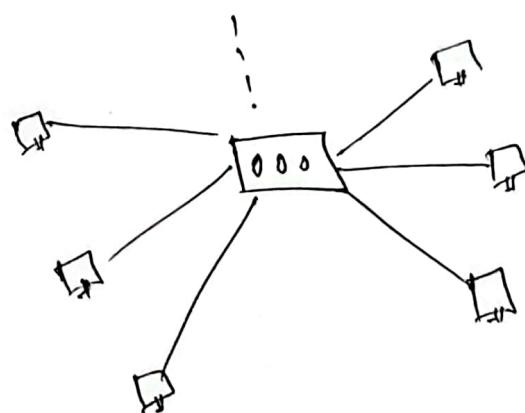
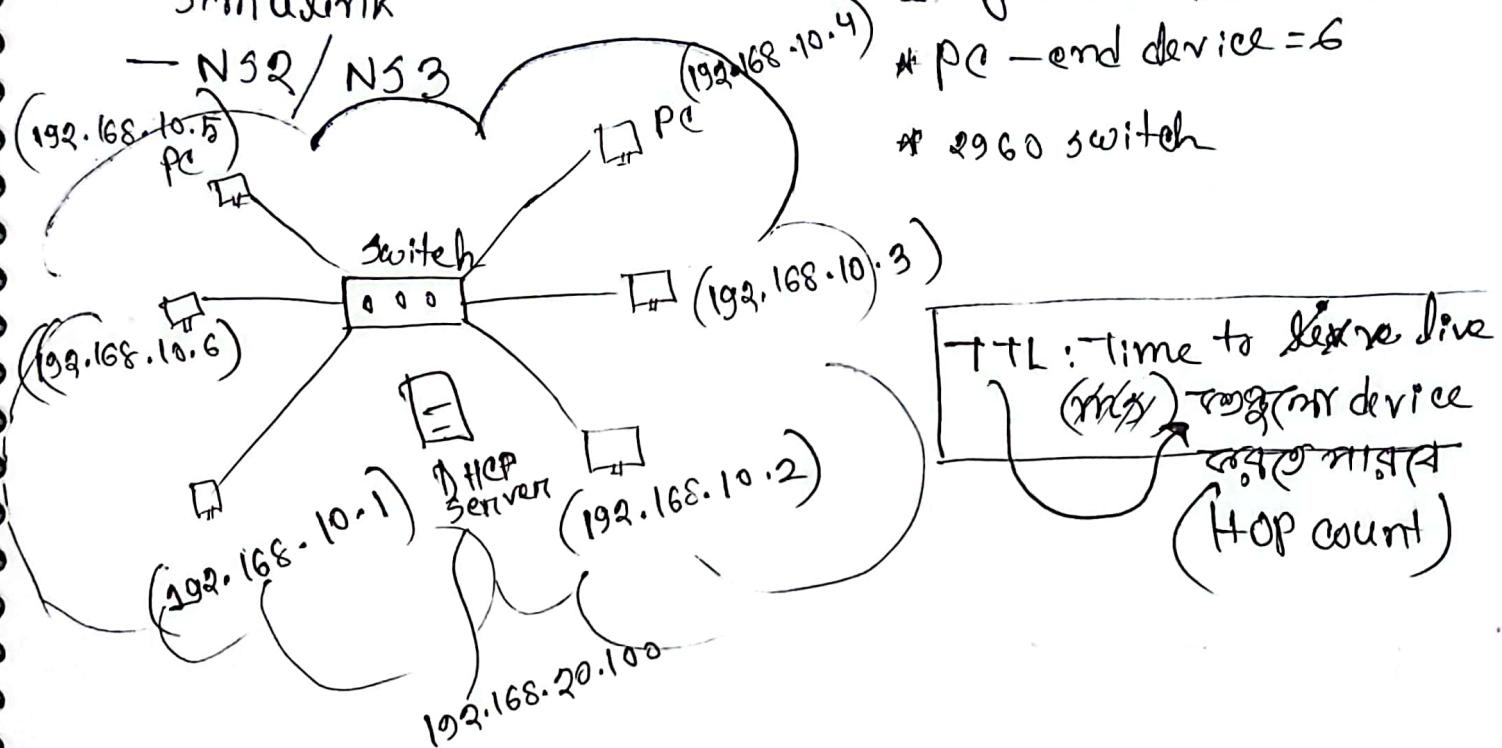
CISCO packet Tracer (pre-defined OTC3)

- OPNET

- OMNET

- Simulink

- NS2 / NS3



IP Address
Internet Protocol (IP)

Date: / /

specific address for the device
total 2^{32}

0000 000(0)

110 00000 000001010 ...

1111 1111(255)

↓

192.10.15.5

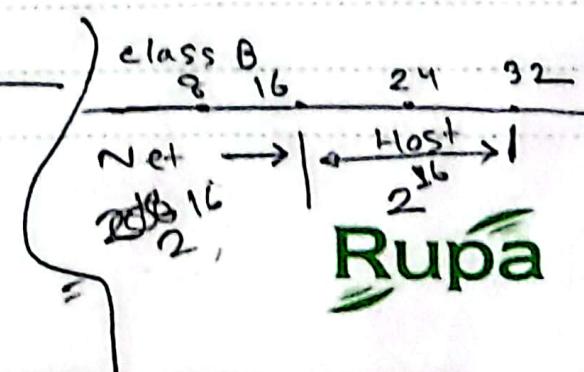
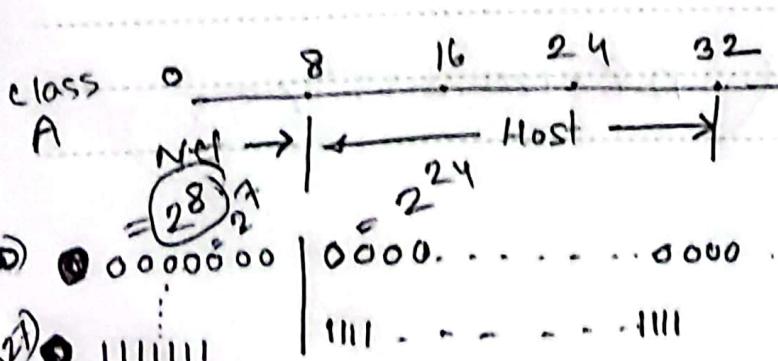
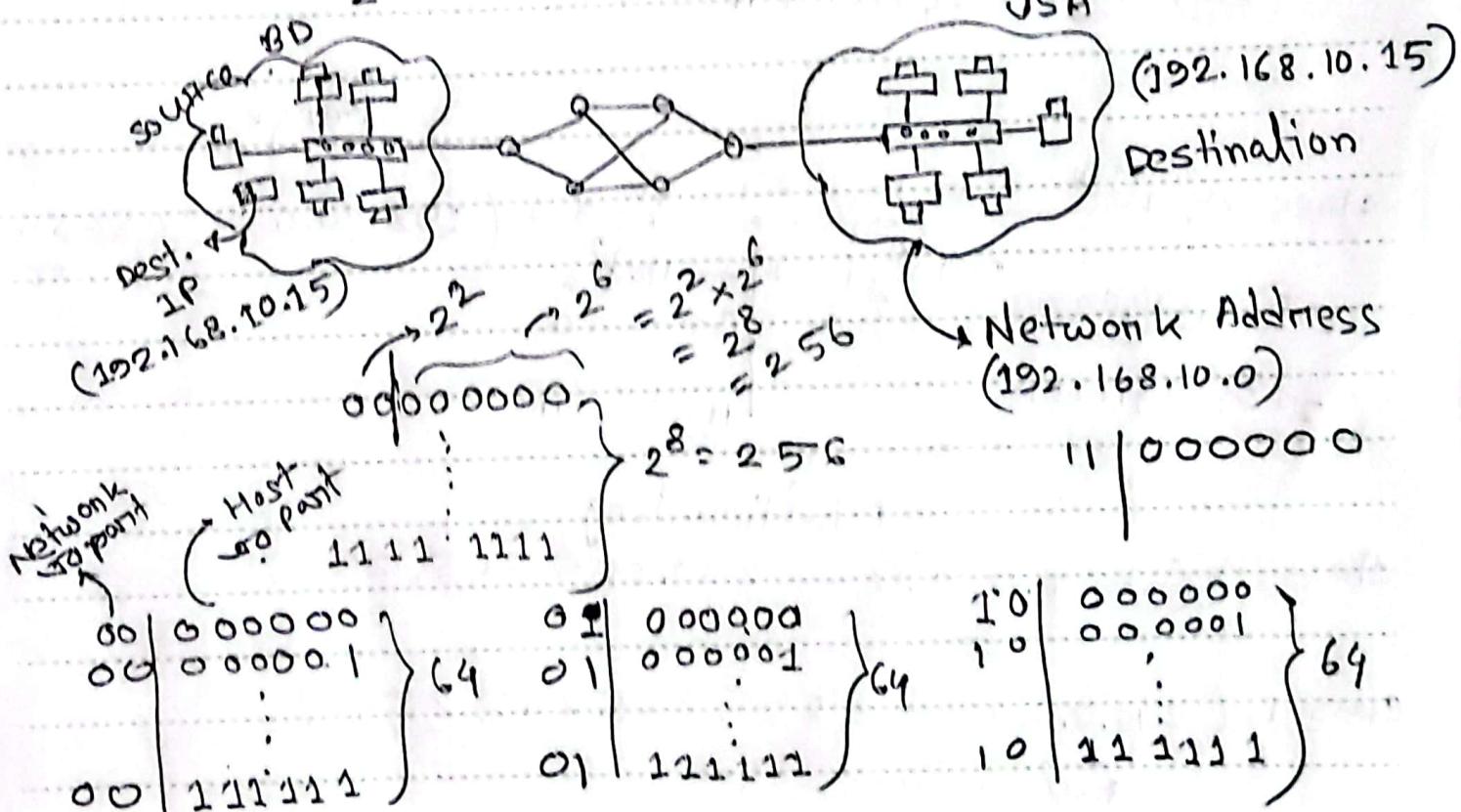
(decimal notation)

IPv6 (128 Bits long)

2^{128}

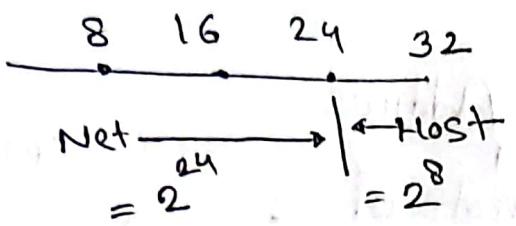
Quiz #1

13 March '29



Rupa

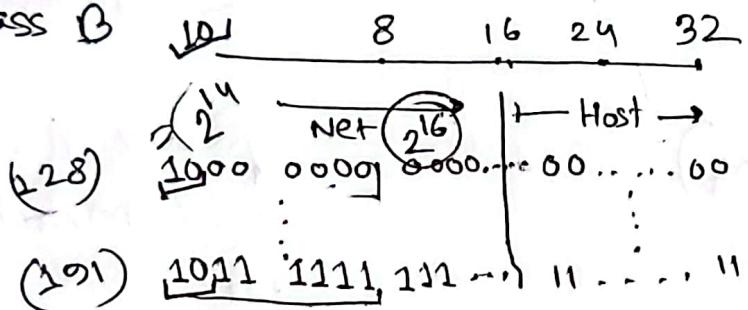
Class C



$$2^{15} - 2^0 \\ 32 \ 16 \ 8 \ 4 \ 2^0$$

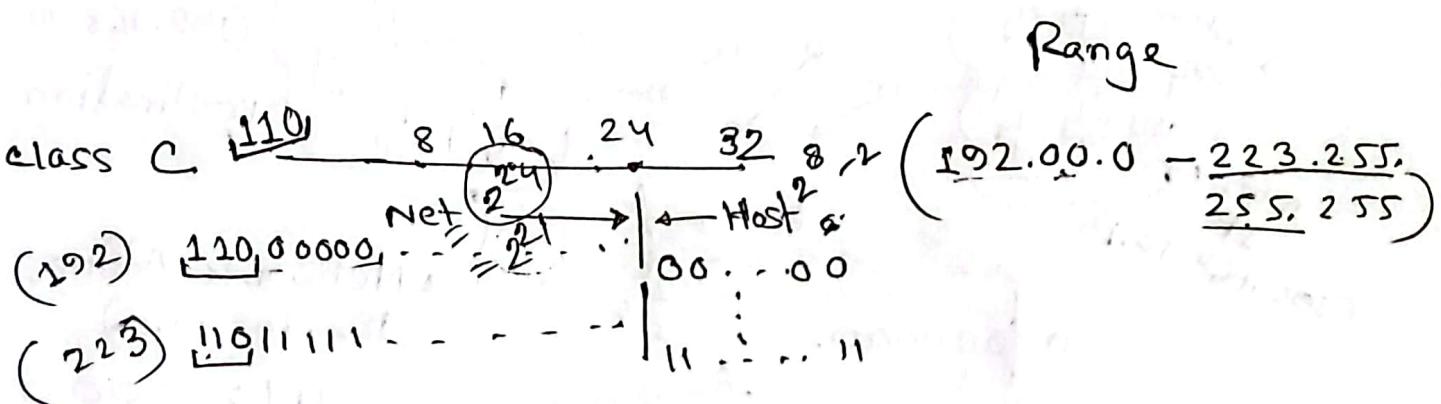
Class A Range (0.0.0.0 — 127.255.255.255)

Class B



Range

$$128.0.0.0 - 191.255.255.255$$



Range

Class D (224.0.0.0 — 239.255.255.255)

Class E (240.0.0.0 — 255.255.255.255)

Future use

$$(128.120.255.255)$$

00...00

Destination IP : 192.168.10.15
Net Host → 11111111 (255)

Date → 0900000000 (°)

class C

Network IP Address : 192.168.10.0

Broadcast IP Address : 192.168.10.255

Subnet Mask IP Address :

Network Address

11...11 11...11 11...11 00...00
255.255.255.0

X (192.168.10.0)

192.168.10.1

Total
254
(Host)

Dest IP : 000.00.00.00

192.168.10.254

X 192.168.10.255

156.140.35.10
Net Host

11...11 11...11

class : B

Network IP Address : 156.140.0.0

Broadcast IP Address : 156.140.255.255

Subnet Mask IP : 255.255.0.0

Rupa

Lijenta

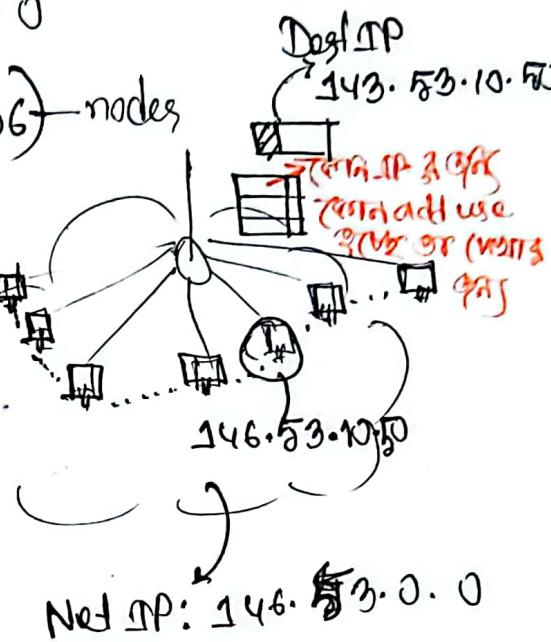
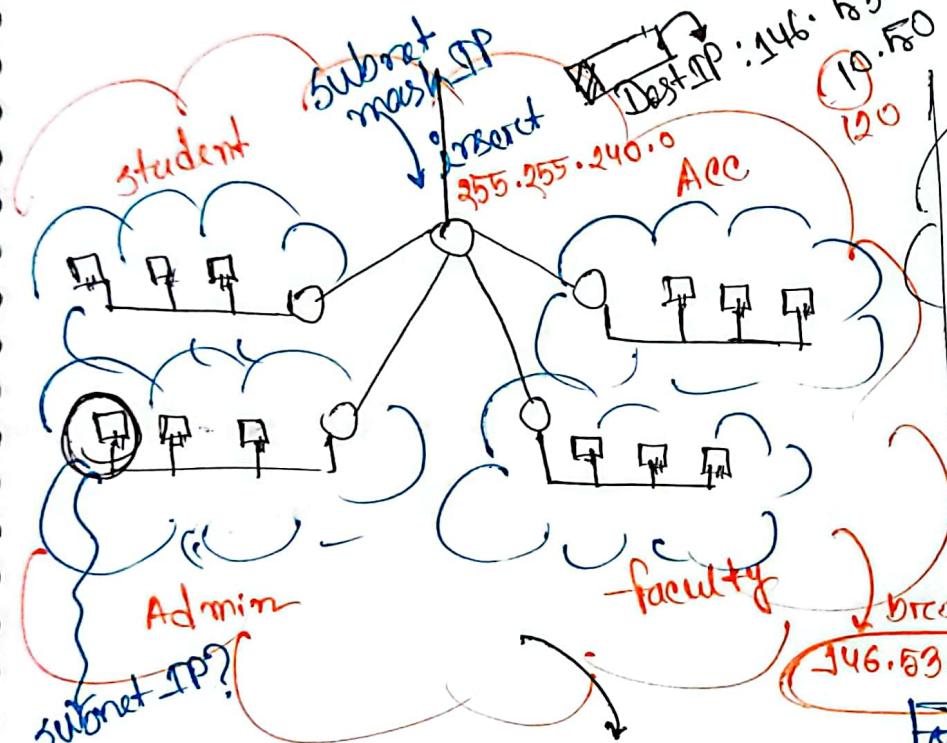
Dec. 9 /

Net.
Destination IP: 146.183.10.50
(class:B) Host

* address in database layer
as subject in main address
first

Network IP address : ~~192.168.1~~ 146.83.0.0

Total # of host : 2^{16} (65536) nodes



Net DP: 146. ~~3~~ 3. 0. 0

ଏମାନ୍ search କରିବାରେ

ଶିଳ୍ପ ମୁଦ୍ରଣ - disrupted

एक networks द्वारा बनाये जाएँ
फिलहाल Sub networks बनाते
होते, Subnetwork द्वारा (Sub
network)

এখানে এটা এটা select করে প্রত্যাক্ষণ subnet করে
down করে লাগু করার সাথে। এটা use করা
বিশেষ।

Network part unchanged

દ્વારા રચિત

15+ sub :

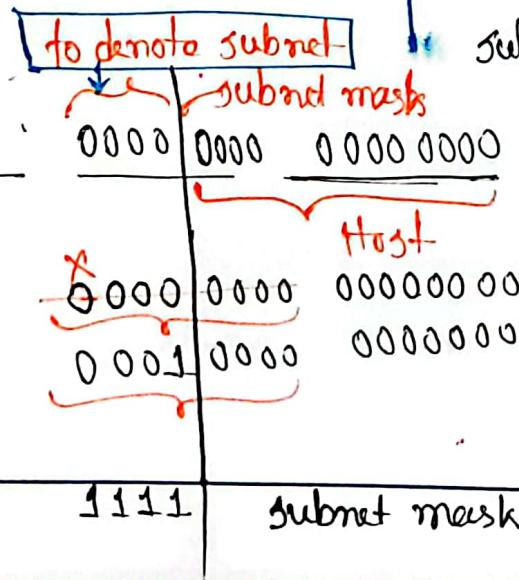
~~1st~~ 2nd sub

100-3

1

Last sub : 1111

subnet mask = separation



$$2^4 \times 2^{12} = 2^{16}$$

-Host

~~1st subnet IP : 146.53.0.0~~

subnet পরি ওপরে
লোগো change কৰ
শক্তি না.

~~2nd subnet IP : 146.53.16.0~~

প্রথম 1st subnet 0111

Net IP same; CIDR 24.

ব্রডকাস্ট

| | | | |
|------------|-----------|---------------------|-------------------------|
| <u>146</u> | <u>53</u> | 0 0 0 0 0 0 0 0 | 0 0 0 0 0 0 0 0 |
| ← Net | → | | |
| 1st sub | | : 0 0 0 1 0 0 0 0 | 0 0 0 0 0 0 0 → 1st sub |
| 2nd sub | | : 0 0 1 0 0 0 0 0 | 0 0 0 0 0 0 0 → 2nd sub |

1st subnet IP : 146.53.16.0

[প্রথম num network এর
সাথে শেষ num broadcast
করা হয়ে match হওয়া]

2nd subnet IP : 146.53.32.0

3rd subnet IP : 146.53.48.0

| | | | |
|------------|-----------|---------------------|-----------------------------|
| <u>146</u> | <u>53</u> | 0 0 0 0 0 0 0 0 | 0 0 0 0 0 0 0 0 |
| | | | |
| 1st subnet | | : 0 0 0 1 0 0 0 0 | 0 0 0 0 0 0 0 → 1st sub |
| | | : 0 0 0 1 0 0 0 0 | 0 0 0 0 0 0 0 → 1st host |
| | | | 0 0 0 0 0 0 1 → 1st host |
| | | | 0 0 0 0 0 1 0 → 2nd host |
| | | | 1 1 1 1 1 1 1 1 → last host |

1st subnet 1st host IP : 146.53.16.0

146.53.16.1 → last host

1st subnet 2nd host IP : 146.53.16.2

:

Broadcast

1st subnet last host IP : 146.53.31.255

:

1st subnet last host : 146.53.31.254

→ 1st last number

→ in broadcast addr

→ 1st last host of

→ first 3 bits of number

→ 1100 - 1000

subnet 21 goes on 6. last which 59 :

| | | | |
|-----|----|------------------|-----------------------------|
| 146 | 53 | <u>0000 0000</u> | <u>00000000</u> |
| | | <u>0011 0000</u> | <u>00000001</u> → 1st host |
| | | <u>1011 1111</u> | <u>11111110</u> → last host |

3rd subnet 1st host IP : 146.53.48.1

3rd subnet last host IP : 146.53.63.254

| | | | |
|-----|----|------------------|----------------------------|
| 146 | 53 | <u>0000 0000</u> | <u>00000000</u> |
| | | <u>1111 0000</u> | <u>00000000</u> |
| | | <u>X111 1111</u> | <u>11111111</u> |
| | | <u>1110 0000</u> | <u>00000000</u> → last sub |

last subnet IP : 146.53.840.0 X

last subnet Broadcast IP : 146.53.255.255

because this broadcast IP is same as the mother

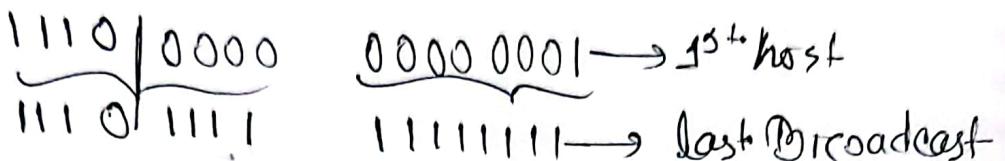
broadcast IP

∴ last subnet IP : 146.53.53.224.0

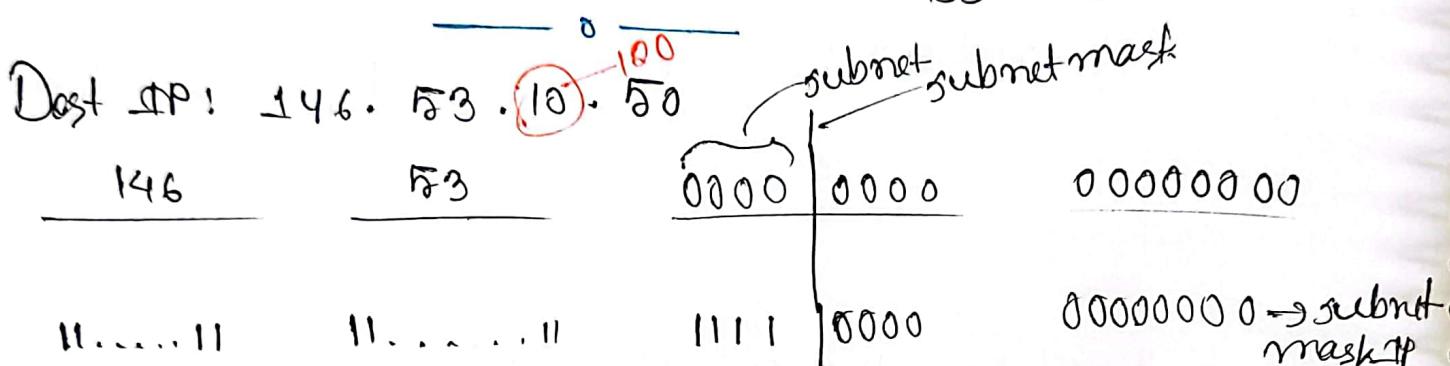
$$\text{Subnet} = 2^4 - 2 = 16 - 2 = 14 \text{ Subnet}$$

→ 1 bit subnet হতে পারবেনা ; atleast 2 bits নিতে হবে $2^2 - 2 = 2$

Last subnet 1st host : 146. 53. 224. 1



Last subnet Broadcast IP : 146. 53. 239. 255



Subnet mask IP (decimal notation) : network + subnet portion = 1
+ host part = 0

→ 146. 255. 255. 240. 0

Subnet mask IP (CDR) : 146. 53. 10. 50 / 20 → fact separation
- 010

(class less interc)

Domain

Subnet mask IP : 255. 255. 255. 240. 0 AND

Dest : 146. 53. 10. 50

AND

10010010 00110101

146 53

0111000 00110010

00001010 00110010

01110000 00110010

01110000 00110010

01110000 00110010

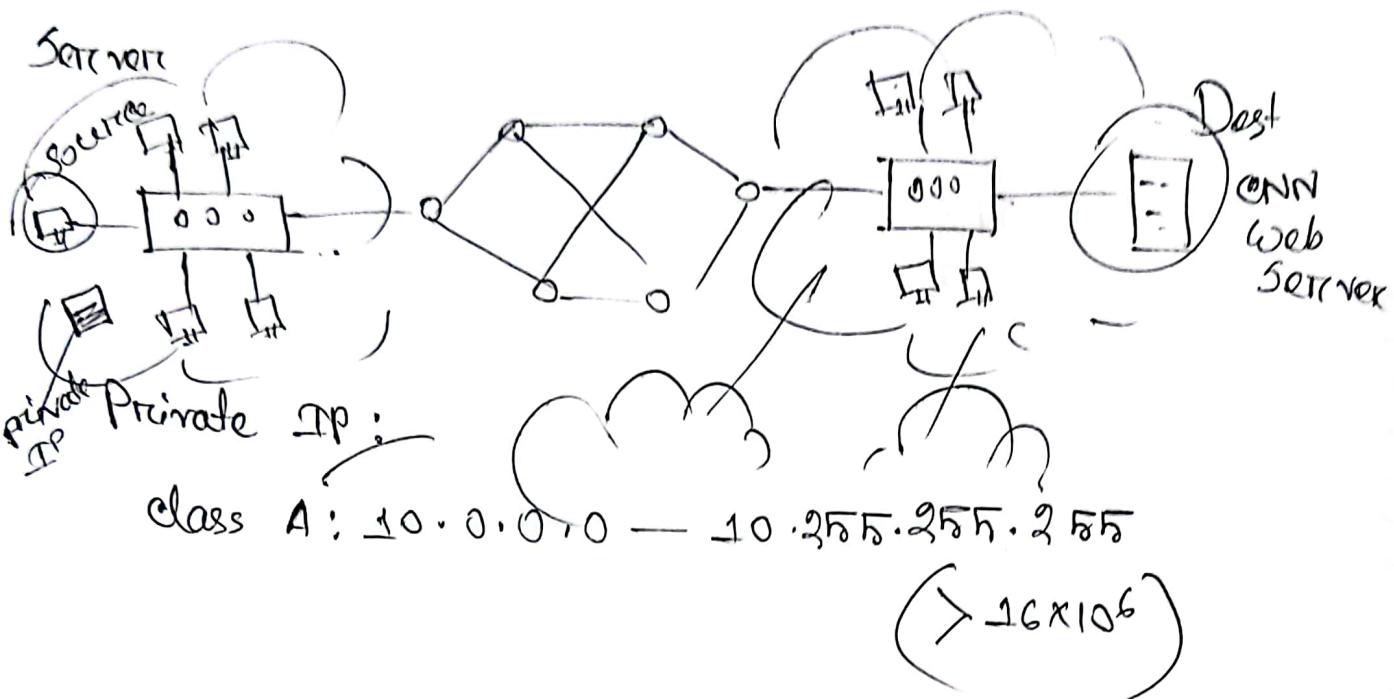
01110000 00110010

Subnet IP : 146. 53. 112. 0

11111111 11111111 11110000 00000000
10010010 00110101 01111000 00110010

~~Lab - 9~~
~~lec - 10~~

IP v4 (32 Bits) Quick fix? → IP v6 (128 Bits)
total = 2^{128}



Class B: 172.16.0.0 — 172.31.255.255 ($> 1 \times 10^6$)

Class C: 192.168.0.0 — 192.168.255.255

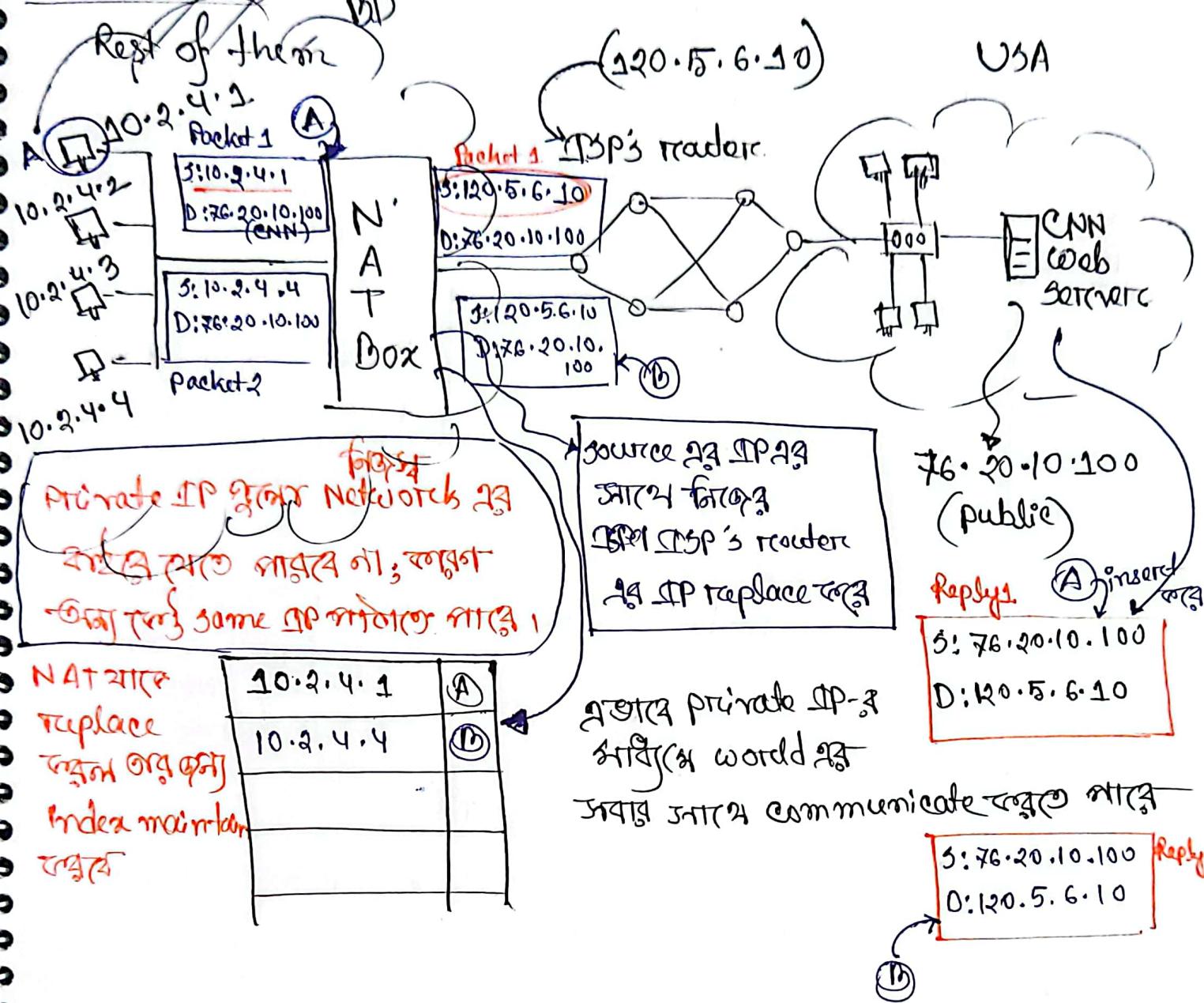
($> 65 \times 10^3$)

Real (Public)
Rest of them

NAT (Network Address Translation)

(feature not help to fix)

Real (Public)



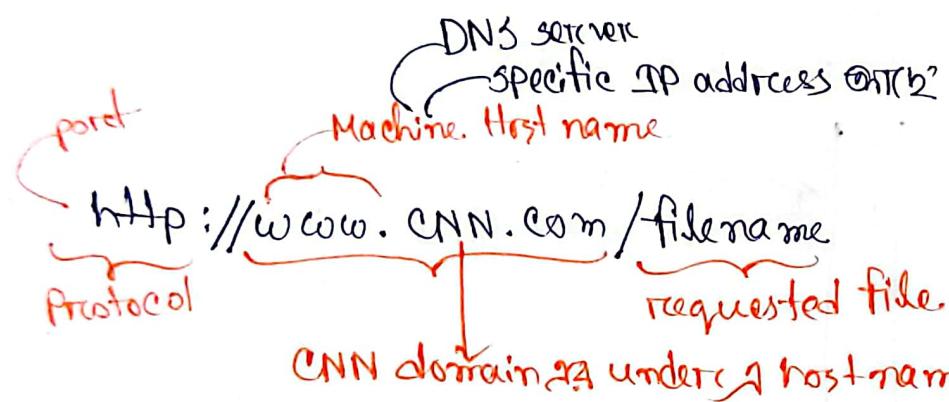
Port address
(16 bits long)
header

- IPDU-ରୁ କୌଣସି ଥାଏ ପୋର୍ଟ ଅଠେସ୍
- frame ରୁ କୌଣସି ଥାଏ ନେଟ୍‌ଵେବ୍ ଅଠେସ୍

memory location
specify ...
location

| Source port | Dest Port |
|-------------|-----------|
| | |
| | |

16



2¹⁶ ports

= 65536 ports

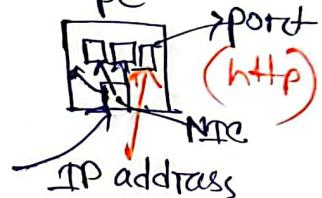
ଆହେ ଉପରେ device

ଏହି ରେଣ୍ଡର କରାଯାଇଲା

port ଏବେଳାବିଲ୍ ଅଛି

ତୁମ୍ହଙ୍କ କାହାରେ

- SIP addିଶ୍ୟ ପୂର୍ବପାଦେ
- ଫିର୍ତ୍ତୁ service ନିବାଳି
- ଅନ୍ତର୍ଭାବରେ port ପୂର୍ବପାଦେ



DNS - Domain name system server

କେବଳିକୁ କାମିଦିନୀ କରିବାରେ କରିବାରେ କରିବାରେ

IP ନିଯମ ଅନୁଯାୟୀ କରିବାରେ କରିବାରେ

• Reserved port (0 - 1023) → Well known services ଏହି କରିବାରେ

ex: http: (port 80)

Telnet: 23

Ftp: 21

Smtp: 110

• Rest → free (not for specific use) (idle)

