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1  |----- MODULE Skeen -----|
   | The specification of Skeen's protocol for atomic multicast; see Section III of the DSN2019 paper |
   | "White-Box Atomic Multicast" by Alexey Gotsman, Anatole Lefort, and Gregory Chockler. |
7  | EXTENDS Naturals, Sequences, FiniteSets, TLC |
8  |-----|
9  |  $Injective(f) \triangleq \forall a, b \in \text{DOMAIN } f : (a \neq b) \Rightarrow (f[a] \neq f[b])$  |
11 |  $Max(a, b) \triangleq \text{IF } a > b \text{ THEN } a \text{ ELSE } b$  |
12 |-----|
13 | CONSTANTS |
14 |   Msg,      the set of messages, ranged over by m |
15 |   Proc,      the set of processes, ranged over by p |
16 |   Dest       $Dest[m] \subseteq Proc$ : the set of destination processes of  $m \in Msg$  |
18 | ASSUME |
19 |    $\wedge Dest \in [Msg \rightarrow \text{SUBSET } Proc]$  |
21 |  $Priority \triangleq \text{CHOOSE } f \in [Proc \rightarrow 1 \dots Cardinality(Proc)] : Injective(f)$  |
22 |-----|
23 | VARIABLES |
24 |   clock,       $clock[p]$ : the clock at process  $p \in Proc$  |
25 |   phase,       $phase[p][m]$ : the phase of the message  $m \in Msg$  at process  $p \in Proc$  |
26 |   localTS,     $localTS[p][m]$ : the local ts of the message  $m \in Msg$  at process  $p \in Proc$  |
27 |   globalTS,    $globalTS[p][m]$ : the global ts of the message  $m \in Msg$  at process  $p \in Proc$  |
28 |   delivered,   $delivered[p][m]$ : has  $m \in Msg$  been delivered at process  $p \in Proc$  |
29 |   incoming,   $incoming[p] \subseteq Message$  (defined below): the incoming channel of process  $p \in Proc$  |
30 |   sent        $sent \subseteq Msg$ : the set of messages that have been multicast; only for TLC |
32 |  $pvars \triangleq \langle clock, phase, localTS, globalTS, delivered \rangle$  |
33 |  $vars \triangleq \langle clock, phase, localTS, globalTS, delivered, incoming, sent \rangle$  |
34 |-----|
35 |  $MaxCounter \triangleq Cardinality(Msg) * Cardinality(Proc)$  |
37 |  $TS \triangleq [c : 0 \dots MaxCounter, p : Proc] \text{ c for counter}$  |
39 |  $GT(u, v) \triangleq \text{Is } u > v?$  |
40 |    $\vee u.c > v.c$  |
41 |    $\vee \wedge u.c = v.c$  |
42 |    $\wedge Priority[u.p] > Priority[v.p]$  |
44 |  $MaxV(vs) \triangleq \text{CHOOSE } u \in vs : \forall v \in vs : u \neq v \Rightarrow GT(u, v)$  |
45 |-----|
46 |  $Message \triangleq [type : \{ "MULTICAST" \}, m : Msg]$  |
47 |    $\cup [type : \{ "PROPOSE" \}, m : Msg, p : Proc, lts : TS]$  |
49 |  $Send(msg) \triangleq \text{Send } msg \in Message \text{ to its destination processes}$  |
50 |    $incoming' = [p \in Proc \mapsto$  |
51 |     IF  $p \in Dest[msg.m]$  THEN  $incoming[p] \cup \{msg\}$  |

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ELSE  $incoming[p]$ 

Send  $msg \in Message$  to its destination processes and remove  $rmsg \in Message$  from  $incoming[sender]$

Precondition:  $sender \in Dest[msg.m]$

59  $SendAndRemove(msg, sender, rmsg) \triangleq$ 60  $incoming' = [p \in Proc \mapsto$ 61 IF  $p = sender$  THEN  $(incoming[sender] \cup \{msg\}) \setminus \{rmsg\}$ 62 ELSE IF  $p \in Dest[msg.m]$  THEN  $incoming[p] \cup \{msg\}$ 63 ELSE  $incoming[p]$ 64  $]$ 65  $TypeOK \triangleq$ 66  $\wedge clock \in [Proc \rightarrow 0 \dots MaxCounter]$ 67  $\wedge phase \in [Proc \rightarrow [Msg \rightarrow \{\text{"START"}, \text{"PROPOSED"}, \text{"COMMITTED"}\}]]$ 68  $\wedge localTS \in [Proc \rightarrow [Msg \rightarrow TS]]$ 69  $\wedge globalTS \in [Proc \rightarrow [Msg \rightarrow TS]]$ 70  $\wedge delivered \in [Proc \rightarrow [Msg \rightarrow BOOLEAN]]$ 71  $\wedge incoming \in [Proc \rightarrow SUBSET Message]$ 72  $\wedge sent \subseteq Msg$ 73  $]$ 74  $Init \triangleq$ 75  $\wedge clock = [p \in Proc \mapsto 0]$ 76  $\wedge phase = [p \in Proc \mapsto [m \in Msg \mapsto \text{"START"}]]$ 77  $\wedge localTS = [p \in Proc \mapsto [m \in Msg \mapsto [c \mapsto 0, p \mapsto p]]]$ 78  $\wedge globalTS = [p \in Proc \mapsto [m \in Msg \mapsto [c \mapsto 0, p \mapsto p]]]$ 79  $\wedge delivered = [p \in Proc \mapsto [m \in Msg \mapsto FALSE]]$ 80  $\wedge incoming = [p \in Proc \mapsto \{\}]$ 81  $\wedge sent = \{\}$ 82  $]$ 83  $Multicast(m) \triangleq$  Multicast  $m \in Msg$ 84  $\wedge m \in Msg \setminus sent$ 85  $\wedge sent' = sent \cup \{m\}$ 86  $\wedge Send([type \mapsto \text{"MULTICAST"}, m \mapsto m])$ 87  $\wedge UNCHANGED pvars$ 89  $Propose(p) \triangleq$  When  $p \in Proc$  receives a *MULTICAST* for some  $m \in Msg$ 90  $\exists msg \in incoming[p] :$ 91  $\wedge msg.type = \text{"MULTICAST"}$ 92  $\wedge LET m \triangleq msg.m$ 93 IN  $\wedge Assert(p \in Dest[m], \text{"p should be one of the destination process of m"})$ 94  $\wedge clock' = [clock \text{ EXCEPT } ![p] = @ + 1]$ 95  $\wedge localTS' = [localTS \text{ EXCEPT } ![p][m] = [c \mapsto clock'[p], p \mapsto p]]$ 96  $\wedge phase' = [phase \text{ EXCEPT } ![p][m] = \text{"PROPOSED"}]$ 97  $\wedge SendAndRemove([type \mapsto \text{"PROPOSE"}, m \mapsto m, p \mapsto p,$ 98  $lts \mapsto localTS'[p][m]], p, msg)$

99  $\wedge \text{UNCHANGED } \langle \text{globalTS}, \text{delivered}, \text{sent} \rangle$

101  $\text{Deliver}(p) \triangleq$  When  $p \in \text{Proc}$  receives all *PROPOSE* for some  $m \in \text{Msg}$

102  $\quad \exists m \in \text{Msg} :$

103  $\quad \text{LET } \text{msgofm} \triangleq \{ \text{msg} \in \text{incoming}[p] : \text{msg.type} = \text{"PROPOSE"} \wedge \text{msg.m} = m \}$

104  $\quad \text{destofm} \triangleq \{ \text{msg.p} : \text{msg} \in \text{msgofm} \}$

105  $\quad \text{ltsofm} \triangleq \{ \text{msg.lts} : \text{msg} \in \text{msgofm} \}$

106  $\quad \text{IN } \wedge \text{destofm} = \text{Dest}[m]$

107  $\quad \wedge \text{globalTS}' = [\text{globalTS} \text{ EXCEPT } ![p][m] = \text{MaxV}(\text{ltsofm})]$

108  $\quad \wedge \text{clock}' = [\text{clock} \text{ EXCEPT } ![p] = \text{Max}(\text{clock}[p], \text{globalTS}'[p][m].c)]$

109  $\quad \wedge \text{phase}' = [\text{phase} \text{ EXCEPT } ![p][m] = \text{"COMMITTED"}]$

110  $\quad \wedge \text{LET } \text{readym} \triangleq \{ rm \in \text{Msg} :$

111  $\quad \quad \wedge \text{phase}'[p][rm] = \text{"COMMITTED"}$

112  $\quad \quad \wedge \text{delivered}[p][rm] = \text{FALSE}$

113  $\quad \quad \wedge \forall pm \in \text{Msg} :$

114  $\quad \quad \text{phase}'[p][pm] = \text{"PROPOSED"}$

115  $\quad \quad \Rightarrow \text{GT}(\text{localTS}[p][pm], \text{globalTS}'[p][rm]) \}$

116  $\quad \text{IN } \text{delivered}' = [\text{delivered} \text{ EXCEPT } ![p] = [pm \in \text{Msg} \mapsto$

117  $\quad \quad \text{IF } pm \in \text{readym} \text{ THEN TRUE ELSE } @[pm]]]$

118  $\quad \quad \text{TODO: deliver in } \text{globalTS}[p][m] \text{ order (using deliver sequence)}$

119  $\quad \wedge \text{UNCHANGED } \langle \text{localTS}, \text{sent}, \text{incoming} \rangle$

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120  $\text{Next} \triangleq$

121  $\quad \vee \exists m \in \text{Msg} : \text{Multicast}(m)$

122  $\quad \vee \exists p \in \text{Proc} :$

123  $\quad \quad \vee \text{Propose}(p)$

124  $\quad \quad \vee \text{Deliver}(p)$

125

127  $\text{Spec} \triangleq \text{Init} \wedge \Box[\text{Next}]_{\text{vars}}$

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128 Invariant: Global timestamps are unique for each  $m \in \text{Msg}$ ; see Section III.

132  $\text{UniqueGTS} \triangleq$

133  $\quad \forall p \in \text{Proc}, m1, m2 \in \text{Msg} :$

134  $\quad (\text{m1} \neq \text{m2} \wedge \text{phase}[p][m1] = \text{"COMMITTED"} \wedge \text{phase}[p][m2] = \text{"COMMITTED"})$

135  $\quad \Rightarrow \text{globalTS}[p][m1] \neq \text{globalTS}[p][m2]$

Invariant: Each  $m \in \text{Msg}$  is assigned a single global timestamp.

140  $\text{SameGTS} \triangleq$

141  $\quad \forall p1, p2 \in \text{Proc}, m \in \text{Msg} :$

142  $\quad (\text{phase}[p1][m] = \text{"COMMITTED"} \wedge \text{phase}[p2][m] = \text{"COMMITTED"})$

143  $\quad \Rightarrow \text{globalTS}[p1][m] = \text{globalTS}[p2][m]$

TODO: Invariant: atomic multicast

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149  $\text{THEOREM } \text{TypeTheorem} \triangleq \text{Spec} \Rightarrow \Box \text{TypeOK}$

151 THEOREM  $UniqueGTSTheorem \triangleq Spec \Rightarrow \Box UniqueGTS$

153 THEOREM  $SameGTSTheorem \triangleq Spec \Rightarrow \Box SameGTS$

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