# Algorithmics Correction Midterm #4 (C4)

Undergraduate  $2^{nd}$  year (S4) — Epita  $6 \; March \; 2018 - 14:45$ 

## Solution 1 (Cut points, cut edges - 3 points)

1. Spanning forest for the DFS of the graph  $G_1$ :

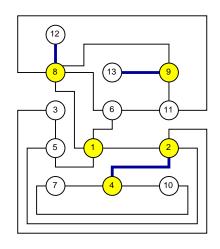


Figure 1: Graphe  $G_1$ 

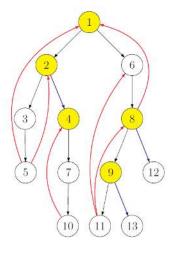


Figure 2: Forêt couvrante

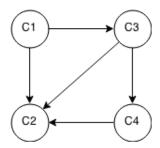
- 2. Cut points of  $G_1$ : 1, 2, 4, 8, 9
- 3. Cut edges of  $G_1$ : (2,4), (8,12), (9, 13).

## Solution 2 (SCC and reduced digraphs - 5 points)

- 1. Digraph  $\rightarrow$  condensation
  - (a) Strongly connected components of the digraph  $G_2$ :

 $C_1: \{0, 1\} \\ C_2: \{2, 3, 4, 6, 7, 8\} \\ C_3: \{5\} \\ C_4: \{9\}$ 

(b) Condensation of  $G_2$ :



1

- (c) The addition of a single edge from component  $C_2$  to the component  $C_1$  builds a cycle that traverses all components. Thus it makes the digraph strongly connected.
- 2. Condensation  $\rightarrow$  digraph
  - (a) Vertices in the component  $C_2$  (from 4 to 7) are unreachable from the vertex 0.
  - (b) Among the following paths, which ones can not exist in  $G_3$ ?
    - 3 → 7
    - $4 \rightsquigarrow 21$  existe
    - 18 → 2
    - 11 → 15
  - (c) Adding two edges to  $G_3$  is sufficient to make it strongly connected. For instance  $x_1 \to y_1$  with  $x_1 \in C_6$ ,  $y_1 \in C_2$  and  $x_2 \to y_2$  with  $x_2 \in C_4$ ,  $y_2 \in C_6$ .

# Solution 3 ("Global Connectivity Indicators" - 6 points)

Global connectivity indicators measure the subdivision degree of a graph into connected components separated from each other.

1. The weighted connectivity index expresses the probability that two random vertices can be connected by a path (i.e. belong to the same connected component).

#### 2. Specifications:

The functions indexes (G) computes both "connectivity indexes" simple  $(IC_1)$  and weighted  $(IC_2)$  of the graph G.

```
def __nbVertexDFS(G, s, M):
                   M[s] = True
                   nb = 1
                   for adj in G.adjlists[s]:
                       if not M[adj]:
                           nb += __nbVertexDFS(G, adj, M)
6
                   return nb
               def connectivity(G):
                   M = [False]*G.order
                   k = 0
                   IC2 = 0
                   for s in range(G.order):
                       if not M[s]:
14
                           k += 1
                           nb = __nbVertexDFS(G, s, M)
                           IC2 += nb*nb
                   IC1 = (G.order - k) / (G.order - 1)
18
                   IC2 = IC2 / (G.order * G.order)
19
                   return (IC1, IC2)
```

## Solution 4 (Strongly Connected? - 7 points)

1. Property(ies) of the first component root met:

The digraph is strongly connected if the first component root met during the traversal is the root of the unique spanning tree.

### 2. Specifications:

The function  $is\_strong(G)$  tests whether the digraph G is strongly connected.

```
def __isStronglyConnected(G, x, pref, cpt):
            cpt += 1
            pref[x] = cpt
            return_x = pref[x]
            for y in G.adjlists[x]:
                if pref[y] == 0:
                    (ret_y, cpt) = __isStronglyConnected(G, y, pref, cpt)
                    if ret_y == -1:
                       return (-1, cpt)
9
                    return_x = min(return_x, ret_y)
                else:
11
                    return_x = min(return_x, pref[y])
12
            if return_x == pref[x]:
14
                return (-1, cpt)
16
17
             return (return_x, cpt)
18
```

```
def isStronglyConnected(G):
    pref = [0]*G.order
    cpt = 0
    (r, cpt) = __isStronglyConnected(G, 0, pref, cpt)
    return (r != -1) and (cpt == G.order) # all vertices have been encountered
```