Lexer

COP-3402 Systems Software Paul Gazzillo

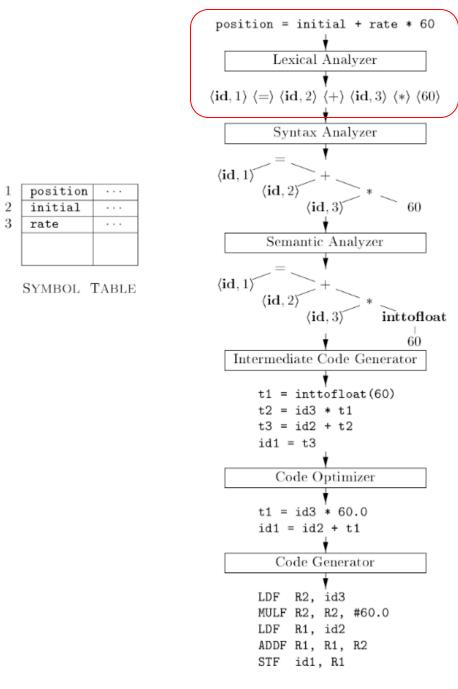


Figure 1.7: Translation of an assignment statement

Core Idea: Convert Characters to Tokens

- Regular expressions specify tokens
 - As sets of strings
- Finite automata *recognize* tokens
 - Via state transitions
- A lexer *produces* a stream of tokens
 - From a stream of individual characters

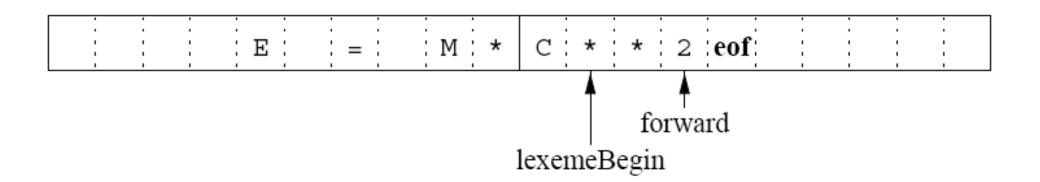


Figure 3.3: Using a pair of input buffers

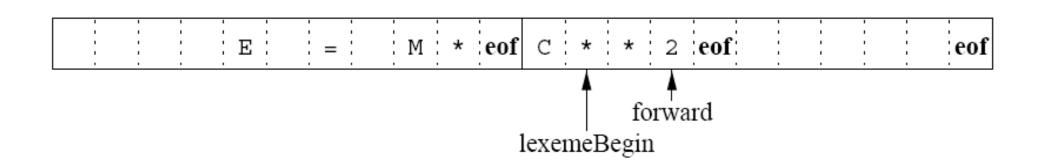


Figure 3.4: Sentinels at the end of each buffer

Token	Informal Description	SAMPLE LEXEMES
if	characters i, f	if
${f else}$	characters e, 1, s, e	else
comparison	< or $>$ or $<=$ or $>=$ or $!=$	<=, !=
\mathbf{id}	letter followed by letters and digits	pi, score, D2
\mathbf{number}	any numeric constant	3.14159, 0, 6.02e23
literal	anything but ", surrounded by "'s	"core dumped"

Figure 3.2: Examples of tokens

LEXEMES	Token Name	ATTRIBUTE VALUE
Any ws		_
if	if	
then	${f then}$	_
else	else	_
Any id	\mathbf{id}	Pointer to table entry
Any number	${f number}$	Pointer to table entry
<	\mathbf{relop}	LT
<=	${f relop}$	LE
=	${f relop}$	EQ
<>	relop	NE
>	$\overline{\text{relop}}$	GT
>=	${f relop}$	GE

Figure 3.12: Tokens, their patterns, and attribute values

Implementing a Lexer

- Regex compiler
 - Regex -> NFA -> DFA -> state table
- Hand-coded lexer
 - Lexer reads one character at a time, e.g., fgetc
 - Concatenation is sequence of statements
 - Union is a conditional
 - Closure is a while loop

Concatenation is sequence of statements, e.g., ab

```
char c;
c = fgetc(lexerin);
assert 'a' == c
c = fgetc(lexerin);
assert 'b' == c
```

Union is a conditional, e.g., a|b

```
char c;
c = fgetc(lexerin);
if ('a' == c) {
    // ...
} else if ('b' == c) {
    // ...
} else {
    error(...)
}
```

Closure is a while loop, e.g., a*

```
char c;
c = fgetc(lexerin);
while('a' == c) {
   c = fgetc(lexerin);
}
// c is now the next character
```

Demo

Lexing PL/0

EBNF Standardizes Language Specifications

- EBNF = Extended Backus-Naur Format
- Different notation for regular expressions
 - a* -> { a }
 - $a | \varepsilon \rightarrow [a]$
 - ab -> a b
 - a | b -> a | b

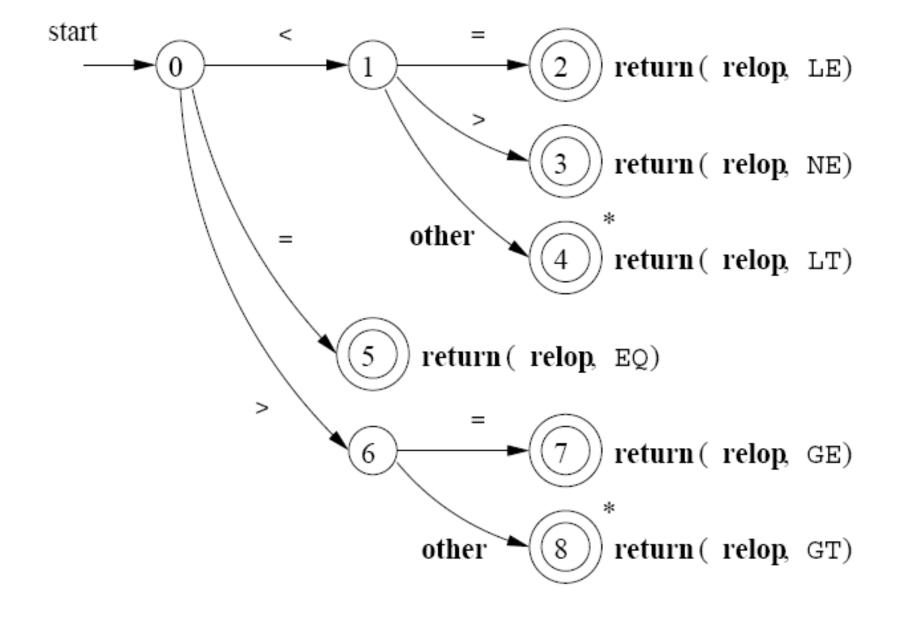


Figure 3.13: Transition diagram for **relop**