



Wounded Triage, Transport and Coordination of Emergency Operation in Urban Emergencies



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Abstract: For the wounded triage, transport and cooperative scheduling problem of emergency surgery in urban emergency rescue, this study uses the idea of supply chain collaborative scheduling, considers factors such as the number of the wounded, rescue vehicle capacity and hospital operation time to achieve the optimization goals of the shortest rescue response time and the most economical transportation capacity, establishes a mixed integer programming model, and designs a two-stage scheduling algorithm to solve the model. It uses the relative gap between the maximum time span of the entire rescue process and the optimal performance under ideal conditions to measure the performance of the algorithm. The simulation experiments show that the two-stage scheduling algorithm has better problem-solving ability for scenarios with larger number of the wounded and stronger carrying capacity, and has better performance than MFF algorithm and MBF algorithm.

Keywords: Urban emergency; Collaborative scheduling; Mixed integer programming model; Two-stage scheduling algorithm

1. Introduction

In the emergency rescue, the wounded triage, transshipment and emergency surgery constitute the complete process of emergency rescue. Due to the shortage of resources involved in the rescue, it is necessary to strengthen the close cooperation among various rescue jobs and coordinate the implementation in order to complete it efficiently. rescue work. With the development of network information technology, all kinds of data can be shared well, and real-time information of each rescue link can be obtained conveniently. Therefore, it is necessary to organically combine the wounded triage, transshipment and emergency surgery to form a complete rescue chain for collaborative scheduling.

The rescue chain in emergency rescue is similar to the supply chain in industrial production. It belongs to the collaborative scheduling problem of rescue transportation under the condition of limited vehicles. It can use a similar supply chain processing method to achieve collaborative optimization. At present, there is no literature on combining all links in emergency rescue to achieve collaborative scheduling, let alone using supply chain methods to study the optimization of emergency rescue processes, so this study herein is theoretically innovative.

Many scholars have studied from different perspectives on the problem of production and transportation collaborative scheduling in the supply chain, and have obtained many good results. Liu [1] studied the production and transportation collaborative scheduling problem and divided the transportation into two stages; with the research goal of minimizing the sum of the manufacturing span time and the total installation cost, he proposed two genetic algorithms, and illustrated the improved genetic algorithm through simulation experiments. The algorithm can solve the problem very well and has better performance than the basic genetic algorithm. Garcia and Lozano [2] studied the collaborative scheduling under multiple constraints such as limited productivity, vehicle quantity and order time window, gave an integer programming model, and used the tabu search algorithm to solve it. Niu and Han [3] proposed a two-stage scheduling algorithm based on the unit manufacturing problem, which solves the problem through process decomposition and algorithm optimization. The scheduling algorithm uses a combination of "accurate" calculation and "approximate" solution. This method not only improves the

computational efficiency but also takes into account the global optimization goal. Hu et al. [4] studied the single-machine ordering problem of collaborative scheduling of processing and transportation. The optimization goal is to minimize the arrival time of the last workpiece sent to the customer, and an optimal algorithm for this problem is given, and is proved that the worst case bound is $3/2$. Ng et al. [5] studied the single-machine, serial batch, and total completion time scheduling problems with priority constraints, same release date and processing time, and proposed a dynamic programming algorithm with time complexity as $O(n^5)$. Liu et al. [6] studied the batch scheduling problem under the condition of unlimited capacity of the batch machine, with the goal of minimizing the total completion time, and proved that the problem is NP.

This study uses the idea of supply chain of industrial manufacturing to analyze the collaborative scheduling of each rescue link in emergency rescue, fully considers the number of wounded and the capacity of rescue vehicles, realizes the scheduling process with the shortest rescue response time and the most economical capacity, establishes a mathematical model, designs an optimization algorithm, and analyzes the performance of the algorithm, and finally tests the actual effect of the algorithm through simulation.

2. Problem Description

After an emergency occurs, the rescue force goes deep into the disaster area, and first triages the wounded rescued. Therefore, it is necessary to open up a field as a place for triage. Due to the shortage of medical rescue forces, the wounded must enter the triage field for queuing and the triage adopts the combination of fuzzy qualitative method and quantitative scoring method. The wounded who are triaged according to their injuries are transported to the hospital in batches by rescue vehicles for further surgical treatment. This process forms a complete rescue chain, and only the collaborative scheduling of each rescue link can ensure efficient and fast rescue. The rescue chain process is shown in Figure 1.

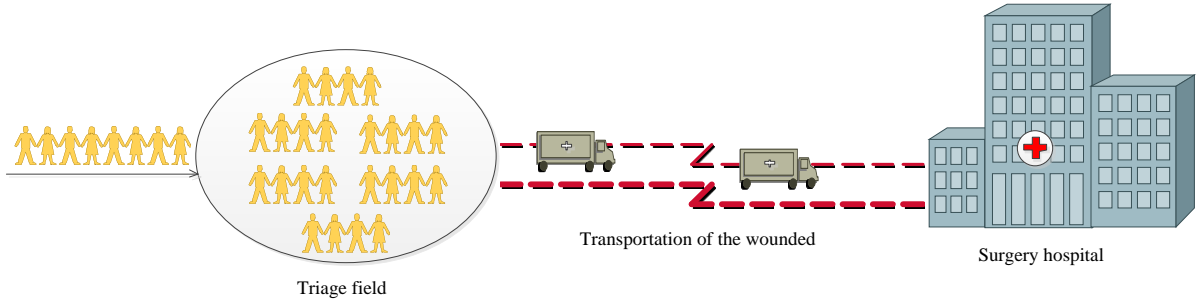


Figure 1. Emergency rescue chain process

Assuming that n wounded persons $\{W_1, W_2, \dots, W_n\}$ arrives at the triage field dynamically, these wounded need to undergo triage in the triage field. Due to the urgency of the rescue, the triage operation is a continuous batch operation, the wounded are sent batch by batch and receive triage one by one, and the end time of the triage is equal to the completion time of the last one in this batch of wounded. The ordering count value of the wounded W_i and the time spent in the triage process are recorded as s_i and p_i , the time to arrive at the triage field is recorded as r_i . Denote the triage time of b_k batch of wounded who received triage on the triage field as P^k , then

$$P^k = \sum_{W_i \in b_k} p_i, i = 1, 2, \dots, n \quad (1)$$

Assume that the carrying capacity of the vehicle is c . Due to the limited number of medical personnel, each batch of wounded must queue up for examination, and the waiting time for the wounded is set to be s . Obviously, the time of the last wounded entering the triage field in a batch of wounded is the waiting time of this batch of wounded. When a batch of wounded has been examined, the batch of wounded are immediately transferred by rescue vehicles. Assume that the one-way travel time of the rescue vehicle from the triage field to the operation hospital is T . When a batch of wounded arrive at the hospital, the batch of wounded will be operated on in the operating room of the hospital. Due to the existence of multiple operating rooms, the emergency rescue operation can be operated on multiple wounded in the same batch at the same time. This mode is a parallel processing machine [7] mode. The operation time of each batch of wounded is equal to a constant P . According to the actual situation of the operation of the wounded, it is assumed $p_i + s \geq P, i = 1, 2, \dots, n$.

3. Problem Modeling

First, the parameters of the given mathematical model are as follows:

γ : total number of wounded; i is the number of the wounded, $i=1, 2, \dots, n$; s_i is the counted number of the wounded i ; r_i is the time when the wounded i arrives at the triage field; p_i is the examination time of the wounded i in a certain triage group on the triage field; P is the operation time of each batch of wounded in the hospital operating room; c is the maximum number of wounded in a batch in the triage field, or the carrying capacity of a transfer vehicle; L is the total number of batches, satisfying $\lceil s_i/c \rceil \leq L \leq n$; k, f is the batch number, $k=1, 2, \dots, L$; s is the waiting time of wounded in line in the triage field; T is the one-way travel time of the rescue vehicle from the triage field to the operation hospital; M is a large positive integer. The decision variables defined are as follows:

x_{ik} : If the wounded i is assigned to the k batch for triage and transfer, then $x_{ik}=0$; otherwise $x_{ik}=1$; y_{kf} : If the wounded of the k batch in the triage field are triaged before the f batch, then $y_{kf}=0$; otherwise $y_{kf}=1$; z_{kf} : If the wounded of the k batch in the hospital operating room are operated before the f batch, then $z_{kf}=0$; otherwise $z_{kf}=1$; P^k : The sum of the triage time of the wounded in the k batch of the triage field; S_{1k} : Triage start time of the wounded in the k batch of the triage field; C_{1k} : Triage end time of the wounded in the k batch of triage field; S_{2k} : Start time of operation of the wounded in the k batch of operating room of the hospital; C_{2k} : End k time of operation of the wounded in the k batch of operating room of the hospital; C_{max} : The maximum time span of whole rescue process or the duration during which the severest wound in the rescue chain is treated.

The mixed integer programming model of the problem is as follows:

$$\text{Min } C_{\max} \quad (2)$$

s.t.

$$\sum_{k=1}^L x_{ik} = 1, \quad i = 1, 2, \dots, n \quad (3)$$

$$\sum_{i=1}^n s_i \cdot x_{ik} \leq c, \quad k = 1, 2, \dots, L \quad (4)$$

$$S_{ik} \geq r_i \cdot x_i, \quad i = 1, 2, \dots, n; k = 1, 2, \dots, L \quad (5)$$

$$C_{ik} = S_{1k} + s + \sum_{i=1}^n x_{ik} \cdot p_i, \quad k = 1, 2, \dots, L \quad (6)$$

$$C_{2k} = S_{2k} + P, \quad k = 1, 2, \dots, L \quad (7)$$

$$C_{2k} = S_{2k} + P, \quad k = 1, 2, \dots, L \quad (8)$$

$$C_{1k} - C_{1f} + P^f + s - (1 - y_{kf})M \leq 0, \quad k = 1, 2, \dots, L; f = 1, 2, \dots, L; k \neq f \quad (9)$$

$$z_{kf} + z_{fk} = 1, \quad k = 1, 2, \dots, L; f = 1, 2, \dots, L; k \neq f \quad (10)$$

$$C_{\max} \geq C_{2k}, \quad k = 1, 2, \dots, L \quad (11)$$

$$x_{ik} \in \{0, 1\}, \quad \forall i, k \quad (12)$$

$$y_{kf} \in \{0, 1\}, \quad \forall k, f \quad (13)$$

$$z_{kf} \in \{0, 1\}, \quad \forall k, f \quad (14)$$

In the model, Formula (2) is the objective function, which represents the span time to minimize the whole rescue. Formulas (3) ~ (14) are constraint conditions, in which, Formula (3) indicates that any wounded can only be assigned to one rescue process batch, and Formula (4) indicates that the total number of all wounded in any batch

cannot exceed the transportation capacity of vehicles c , Formula (5) indicates that the triage of a batch of wounded must begin on the premise that all the wounded in the batch can arrived, Formula (6) indicates the completion time of a batch of wounded in the triage field, Formulas (7) and (8) represent the start time and end time of a batch of wounded in the operating room of the hospital respectively. Formula (9) shows that there is no overlap in the triage process of the wounded in the triage field. Formula (10) indicates that due to the limitation of hospital surgical resources, any two rescue batches at a certain moment cannot be operated simultaneously in the surgical hospital. Formula (11) indicates the nature of the entire rescue process time, and Formulas (12) ~ (14) indicate the range of values of the decision variable.

For the above mathematical model, according to the actual situation of emergency rescue, the following theory is considered:

Lemma 1 If the order of the wounded in the same batch that participates in the triage is interchanged, the scheduling result will not be affected.

Corollary 1 For any optimal scheduling scheme, if the arrival time of any existing batch of wounded is sorted in increasing order and then undergoes triage, this operation will not affect the optimal nature of the scheduling scheme.

Lemma 2 In the optimal scheduling scheme, there is an optimal solution, which contains the batch allocation strategy for the wounded. If all the batches are sorted according to the increasing order of preparation time, this operation does not affect the optimal nature of the scheduling scheme, where the batch preparation time is the maximum arrival time of the wounded in the batch.

Prove Assume that $r_{max} = \max_{1 \leq i \leq n} \{r_i\}$ represents the maximum time for the wounded to arrive at the triage field; in the scheduling scheme, there is a certain batch b_l receiving triage at the triage field; at the same time, there are two adjacent wounded batches b_k and b_f behind the batch b_l . Considering the universality, assuming that $r_k > r_f$ represents the time to arrive at the triage field, it is only necessary to prove that it is still the optimal solution when the triage start time is $S_{1k} > S_{1f}$. If the end time of the triage of the wounded in the l batch of the triage field is C_{1l} , there are two situations:

(1) $C_{1l} \geq r_{max}$

If $S_{1k} < S_{1f}$, the end time of triage of the batch b_k is $C'_{1k} = S_{1k} + 2s + P^f + P^k = C_{1f}$, obviously if the batches b_k and b_f are interchanged, the solution of the scheduling scheme is still the optimal solution, which does not affect the final result at all, so it is still set $S_{1k} > S_{1f}$.

(2) $C_{1l} < r_{max}$

When $r_f < r_k \leq C_{1l}$, the same as (1), it can be concluded that the scheduling scheme is still optimal when $S_{1k} > S_{1f}$. When $r_f \leq C_{1l} < r_k$, if $S_{1k} < S_{1f}$, then $C_{1f} = r_k + 2s + P^f + P^k$; when the batches b_k and b_f are interchanged, the end time of the batches b_k and b_f is respectively $C'_{1f} = C_{1l} + s + P^f$ and $C'_{1k} = \max\{r_k + s + P^k, C_{1l} + 2s + P^f + P^k\}$, it can be concluded that $C'_{1k} < C_{1f}$. Obviously, this conclusion does not conform to the nature of the optimal solution, and contradicts the assumption, so the optimal scheduling scheme should satisfy $S_{1k} > S_{1f}$. When $C_{1l} < r_f < r_k$, if $S_{1k} < S_{1f}$, then $C_{1f} = r_k + 2s + P^f + P^k$; when batches b_k and b_f are interchanged, the end time of triage of batches b_f and b_k is respectively $C'_{1f} = r_f + s + P^f$ and $C'_{1k} = \max\{r_k + s + P^k, r_f + 2s + P^f + P^k\}$, it can be concluded that $C'_{1k} < C_{1f}$. Obviously, this conclusion does not conform to the nature of the optimal solution, and contradicts the assumption, so the optimal scheduling scheme should satisfy $S_{1k} > S_{1f}$.

In summary, when $S_{1k} > S_{1f}$, it is still the optimal solution, and the lemma is proved.

Lemma 3 Suppose there is a scheduling scheme G , $G\{\dots, b_k, b_{k+1}, \dots\}$, if $b_{k+1} = \{W_j\}$, $j=1, 2, \dots, n$, satisfies $S_{1k} < r_j < S_{1k} + s$ and $\sum_{W_i \in b_k} S_i + W_j \leq c$, then when the wounded W_j join the batch b_k , the scheduling scheme is better than the original scheme.

Prove From $r_j < S_{1k} + s$ and $C_{1k} = S_{1k} + s + P^k$, $r_j < C_{1k}$ can be obtained; therefore, $C_{1(k+1)} = S_{1(k+1)} + s + P^{k+1} = S_{1k} + 2s + P^k + p_j$. Since $\sum_{W_i \in b_k} S_i + W_j \leq c$, there is enough queue space in the b_k batch for the wounded W_j to join it. If the wounded W_j is added to the batch b_k , the start and end time of triage of the b_k batch in the triage field are updated to $S'_{1k} = r_j$ and $C'_{1k} = S'_{1k} + s + P^k + p_j = r_j + s + P^k + p_j$ respectively. Therefore, $C'_{1k} - C_{1(k+1)} = r_j + s + P^k + p_j - (S_{1k} + 2s + P^k + p_j) = r_j - S_{1k} - s$. Since $r_j < S_{1k} + s$, $C'_{1k} < C_{1(k+1)}$. Lemma 3 is proved.

Lemma 4 Assume that the wounded of two adjacent batches b_k and b_f are triaged at the triage field, this scheme π^* is the optimal scheduling scheme. If $S_{1k} > S_{1f}$, then $S_{2k} > S_{2f}$.

Prove Since the operation mode of the hospital's operating room belongs to the parallel batch machine mode, it is assumed that the earliest time of operation in the hospital's operating room is E , and the earliest operation time that the batch b_k and b_f can be operated is marked as T_k and T_f respectively, $T_f = C_{1f} + T$ and $T_k = C_{1k} + T$ can be obtained, so $T_k > T_f$. If $S_{2k} < S_{2f}$, then $C_{2f} = \max\{T_k, E\} + 2P$, where P is a constant, representing the operation time of a batch of wounded. If b_k and b_f are interchanged, the end time of the operation of the b_k batch is $C'_{2k} = \max\{\max\{T_f, E\} + P, T_k\} + P$, then $C'_{2k} \leq C_{2f}$; this conclusion does not conform to the nature of the optimal solution, and contradicts the hypothesis. Therefore, $S_{2k} > S_{2f}$, the lemma is proved.

In the triage, some special situations may occur, for example, all wounded have the same arrival time r , and this special situation is called "the same release moment".

Lemma 5 In the case of "the same release time", there are L^* batches of wounded in an optimal dispatching

scheme π^* , then $C_{\max}(\pi^*) = r + L^* \cdot s + \sum_{i=1}^n p_i + T + P$.

Prove Considering this situation, all wounded at time of r can be triaged at the triage field. Then, for b_1 , $C_{11} = S_{11} + s + P^1 = r + s + P^1$ and $C_{21} = S_{21} + P = r + s + P^1 + T + P$. In the same way, for b_2 , $C_{12} = S_{12} + s + P^2 = r + 2s + P^1 + P^2$. If it is recorded C_{21} as the triage start time of any new batch, then based on $p_i + s \geq P$, $C_{12} + T \geq C_{21}$, get $C_{22} = S_{22} + P = r + 2s + P^1 + P^2 + T + P$; similarly, for the l batch, b_l , $l=3, 4, \dots, L^*$, $C_{1l} = r + l \cdot s + \sum_{k=1}^l P^k$ and $C_{2l} = S_{2l} + P = r + l \cdot s + \sum_{k=1}^l P^k + T + P$. Obviously, $C_{\max} = C_{2L^*} = r + L^* \cdot s + \sum_{k=1}^{L^*} P^k + T + P$, because $\sum_{k=1}^{L^*} P^k = \sum_{i=1}^n p_i$, get $C_{\max} = r + L^* \cdot s + \sum_{i=1}^n p_i + T + P$.

From Lemma 5, there are:

Corollary 2 In the case of "the same release moment", if there are L batches in a scheduling scheme π , $C_{\max}(\pi) \geq r + L^* \cdot s + \sum_{i=1}^n p_i + T + P$.

Corollary 3 In the case of "the same release moment", $C_{\max} \geq r + \lceil s_i/c \rceil \cdot s + \sum_{i=1}^n p_i + T + P$.

Lemma 6 In the case of "the same release moment", if there is a scheduling scheme π with no extra idle time in the triage and wounded transfer process, then the fewer batches designed in the scheme is, the better the scheduling solution is.

Prove Assume that in two scheduling schemes π and π' , there is no extra idle time in the process of triage and wounded transfer. The number of batches contained in π is L , the number of batches contained in π' is L' , and by Lemma 3, $C_{\max}(\pi) = r + L \cdot s + \sum_{i=1}^n p_i + T + P$ and $C_{\max}(\pi') = r + L' \cdot s + \sum_{i=1}^n p_i + T + P$, then $C_{\max}(\pi) - C_{\max}(\pi') = r + L \cdot s + \sum_{i=1}^n p_i + T + P - (r + L' \cdot s + \sum_{i=1}^n p_i + T + P)$. If $L \geq L'$, then $C_{\max}(\pi) \geq C_{\max}(\pi')$; otherwise, $C_{\max}(\pi) < C_{\max}(\pi')$.

From Lemma 6, there are:

Corollary 4 In the case of "the same release time", the optimal scheduling scheme has the least number of batches among all scheduling schemes.

Lemma 7 In the case of "the same release time", there is an optimal scheduling scheme. If the order of any two wounded batches in this scheme is interchanged, the scheme is still the optimal scheduling scheme.

Prove Considering the universality, it is assumed that there is an optimal scheduling scheme, which has batches b_k and b_f and the triage order of the b_k batch is close to b_f . If all the wounded can receive triage on the triage field at r , then $C_{1f} = S_{1f} + s + P^f$ and $C_{1k} = S_{1k} + s + P^k = S_{1f} + 2s + P^f + P^k$. When the batches b_k and b_f are interchanged on the triage field, the start and end time of triage of the batch b_k is updated to $S'_{1k} = S_{1f}$ and $C'_{1k} = S'_{1k} + s + P^k$ respectively. The end time of triage of the b_f batch is updated to $C'_{1f} = C'_{1k} + s + P^f = S_{1f} + 2s + P^k + P^f$, then $C_{1k} = C'_{1f}$. According to Lemma 5, the end time of triage of the b_f batch is $C_{2f} = C'_{2k}$. Therefore, this scheme remains optimal.

4. Algorithm Design and Performance Analysis

Based on the problem nature of the whole rescue chain of triage, wounded transfer and emergency surgery, this study designs a two-stage scheduling algorithm (TS-A). The first stage of the algorithm is used to solve the scheduling problem when the end time of triage of the k batch of wounded in the triage is $C_{1k} < r_{\max}$, where $r_{\max} = \max_{i=1,2,\dots,n} \{r_i\}$. The second stage is used to solve the scheduling problem when the start time of triage of the k batch of wounded in the triage is $S_{1k} \geq r_{\max}$. In the scheduling scheme, the situation of the wounded joining the batch and the order of the batches in the scheme has been determined, the problem of the second stage is equivalent to the case of "the same release time", then all the unscheduled wounded have reached the triage field at the time r_{\max} .

The first stage of TS-A algorithm is based on Lemma 3, and the second stage is based on Lemma 5, 6.6, 6.7 and FFD algorithm [8]. It's essential to define the ordering rules first. Rule 1: Sort the wounded in the non-decreasing order of the time they arrived at the triage field. If multiple wounded arrive at the same time, then sort them in the non-increasing order of the count value of the wounded. Rule 2: Sort the wounded in the non-increasing order of the count value of wounded arriving at the triage field. The parameters of the TS-A algorithm are defined as: $W = \{W_1, W_2, \dots, W_n\}$ the set of wounded; $A(t)$ is the set of effectively dispatched wounded at the time t ; $U(t)$ represents the set of wounded who do not participate in the scheduling scheme at time t ; $K(t)$ represents the set of wounded who do not participate in the scheduling scheme and not belong to $A(t)$ at time t , then $K(t) = U(t) \setminus A(t)$; B represents a set of batches in a scheduling; S^k represents the total count value of the wounded in the b_k batch; Q represents the batch number of the wounded who participate in the triage last in the first stage.

The steps of the first stage of the algorithm are as follows:

Step 1: Set $k=1$, $b=\emptyset$, $S^1=0$, $S_{11}=r_1$, $P^1=0$, $t=r_1$, $U(t)=W$, $A(t) = \{W_i \mid r_i = t\}$, and $K(t)=U(t) \setminus A(t)$. The wounded in the sets $U(t)$, $A(t)$ and $K(t)$ are sorted according to rule 1.

Step 2: If $t \geq r_{\max}$, jump to the second stage; otherwise, jump to Step 3.

Step 3: According to the ordering rules, determine whether each wounded in the set $A(t)$ can be added to the b_k batch. If $S^k + s_i \leq c$, then $b_k \cup b_k \{W_i\}$, $S^k = S^k + s_i$ and $P^k = P^k + p_i$. Update the sets $U(t)$, $A(t)$ and $K(t)$.

Step 4: If $S^k = c$, jump to Step 7; otherwise, jump to Step 5.

Step 5: According to the ordering rules, determine whether each wounded in the set $K(t)$ can be added to the b_k batch. If $S^k + s_i \leq c$ and $r_i \leq \alpha s + t$, where the parameter $\alpha (0 \leq \alpha \leq 1)$ has a certain impact on the performance of the algorithm, it is usually taken $\alpha=0.3$ and s is the queuing waiting time; if the condition is met, then jump to Step 6;

otherwise, jump to Step 7.

Step 6: Set $b_k \cup b_k \{W_i\}$, $S^k = S^k + s_i$, $P^k = P^k + p_i$ and $S_{1k} = t$, update the set $U(t)$, jump to Step 8.

Step 7: For the b_k batch for triage, set $B = B \cup b_k$, let $t = r + s + P^k$ and $C_{1k} = t$. If $U(t) = \emptyset$, the algorithm stops; otherwise, let $k = k + 1$, $b_k = \emptyset$, $S^k = 0$ and $P^k = 0$, jump to Step 8.

Step 8: Set $t = \max\{\min\{r_i \mid W_i \in U(t)\}, t\}$ and $S_{1k} = t$, update the sets $A(t)$ and $K(t)$, jump to the second stage.

The steps of the second stage of the algorithm are as follows:

Step 1: Sort the wounded in the set $U(t)$ according to ordering rule 2. Set $Q = k$ and $k = k + 1$, add the wounded with the set $U(t)$ with the number of 1 into the b_k batch, and then update the set $U(t)$.

Step 2: If $U(t) = \emptyset$, stop; otherwise, jump to Step 3.

Step 3: Set $i = i + 1$, add the i wounded to the batch whose batch number is the smallest and greater than Q , and the count value of the wounded should not be greater than $c - s_i$, update the set $U(t)$, and jump to Step 2.

According to Lemma 2, 4 and 7, when all the wounded are arranged in the first and second stages, these batches start the scheduling process according to the order in which they are generated.

In order to evaluate the performance of the TS-A algorithm, two lower bounds of the problem under ideal conditions are considered. The first lower bound means that if the problem is specialized that all the wounded have arrived at the triage field at the time r_{\min} , then $r_{\min} = \min_{i=1,2,\dots,n} \{r_i\}$, denoted as LB_1 , its lower bound is less than or equal to the lower bound of the original problem. According to Corollary 3, it can be obtained:

$$LB_1 = r_{\min} + \lceil s_i / c \rceil \cdot s + \sum_{i=1}^n p_i + T + P$$

The second lower bound assumes that there is no additional idle time during the operation in the hospital operating room, calculate the minimum arrival time of the wounded, the end time of triage of the first wound batch, the one-way transportation time between the triage field and the hospital, and the sum of the operation time of all batches in the hospital, i.e., $LB_2 = r_{\min} + P^1 + T + L \cdot P$. As $P^1 \geq \min_{i=1,2,\dots,n} \{p_i\}$, $L \geq \lceil s_i / c \rceil$, then:

$$LB_2 = r_{\min} + s + \min_{i=1,2,\dots,n} \{p_i\} + T + \lceil s_i / c \rceil \cdot P$$

Therefore, for the whole process of the rescue chain, the optimal performance under ideal conditions is the maximum value of the above two lower bounds, i.e.,

$$LB = \max \{LB_1, LB_2\} = \max \left\{ r_{\min} + \lceil s_i / c \rceil \cdot s + \sum_{i=1}^n p_i + T + P, r_{\min} + s + \min_{i=1,2,\dots,n} \{p_i\} + T + \lceil s_i / c \rceil \cdot P \right\}$$

5. Simulation Demonstration and Analysis

In order to test the performance of the two-stage algorithm TS-A, the analysis is carried out through simulation experiments. In the simulation experiment, the number of wounded and the time to arrive at the triage field are represented by discrete values and interval values respectively, such as the number of wounded, $n \in \{300, 500, 700, 1000, 1500, 2000\}$ and the time $r_i \in [1, 200]$ when the wounded arrive at the triage field, $i = 1, 2, \dots, n$. In the same way, the triage capacity of the triage team or the transport capacity of the transport vehicle $c \in \{150, 280\}$ in the triage field, the count value of the wounded $s_i \in [1, n]$, $i = 1, 2, \dots, n$, the triage time of the wounded $p_i \in [1, 20]$, the waiting time of the wounded in the triage field is $s \in [0, 20]$, the transfer time of the wounded between the triage field and the hospital is $T \in [30, 80]$, and the operation time of the wounded in the operating room of the hospital is $P \in [45, 130]$.

In order to evaluate the performance of the algorithm, the TS-A algorithm is compared with the MFF algorithm and the MBF algorithm [9]. The relative gap between the maximum time span of the whole rescue process obtained by the two-stage algorithm TS-A and the optimal performance under ideal conditions (the maximum of the two lower bounds) can be used to measure the performance of the algorithm. The relative gap percentage is as follows:

$$RG = \frac{C_{\max}^H - LB}{LB} \cdot 100\%$$

Considering the factors that usually vary in simulation experiments, the number of wounded and the number of wounded triaged in each batch or the capacity of rescue vehicles can be used as experimental factors, and the performance of the algorithm is tested with variables n and c . n represents the number of wounded, c represents the capacity of each batch for triage or the carrying capacity of the vehicle, repeat the experiments for ten times

with the values of n and c .

Because $c \in \{150, 280\}$ and $n \in \{300, 500, 700, 1000, 1500, 2000\}$, the experiment is divided into two cases: $c=150$ and $c=280$.

(1) When $c=150$, after 10 repeated experiments, the average value of RG obtained is shown in Figure 2.

It can be seen from the figure that when the number of wounded increases, the average value of RG calculated by the TS-A algorithm decreases. When the number of wounded ranges from 300 to 2000, the average value of RG ranges from 0.75% to 3.23%. It shows that when the capacity of each batch of triage or the carrying capacity of the transport vehicle is 150, as the number of wounded increases, the performance of the algorithm is closer to the optimal performance under ideal conditions. The experimental results also show that the TS-A algorithm has better performance than the MFF algorithm and MBF algorithm.

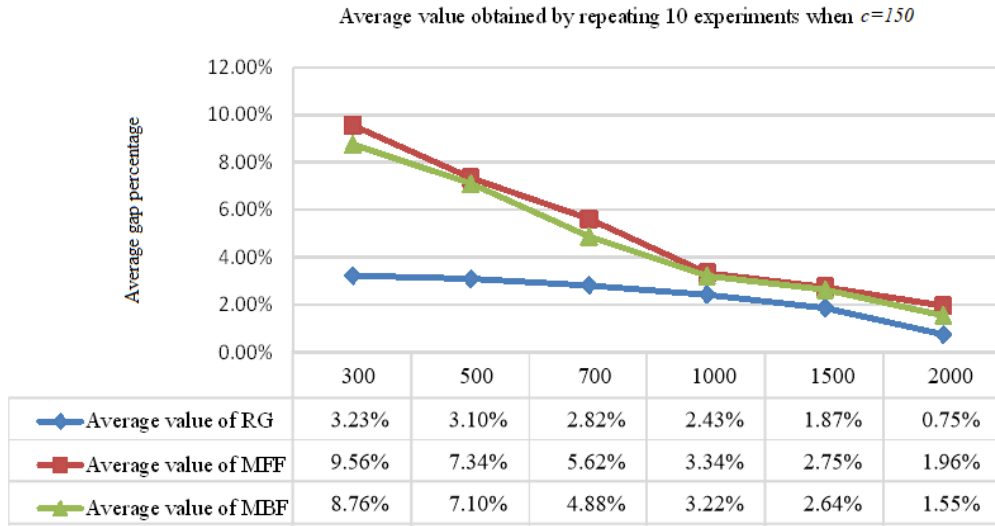


Figure 2. The experimental results of the average gap percentage when $c=150$

(2) When $c=280$, after 10 repeated experiments, the average value of RG obtained is shown in Figure 3.

It can be seen from the figure that when the number of wounded increases, the average value of RG calculated by the TS-A algorithm decreases. When the number of wounded ranges from 300 to 2,000, the average value of RG ranges from 0.23% to 2.94%. It shows that when the capacity of each batch of triage or the carrying capacity of the transport vehicle is 280, as the number of wounded increases, the performance of the algorithm is closer to the optimal performance under ideal conditions. The experimental results also show that the TS-A algorithm has better performance than the MFF algorithm and MBF algorithm.

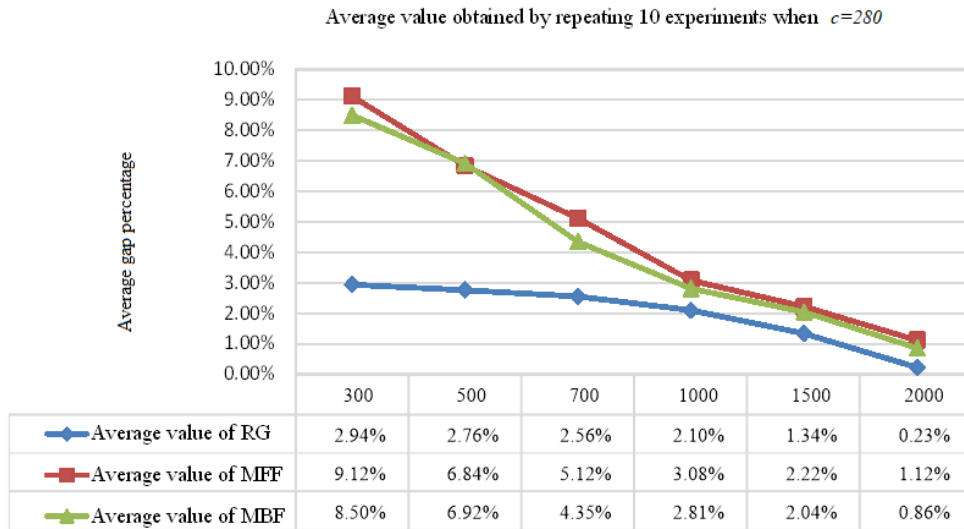


Figure 3. The experimental results of the average gap percentage when $c=280$

Similarly, when the capacity of each batch of triage or the carrying capacity of the transport vehicle increases, the average gap percentage also decreases, indicating that the increase in the capacity of each batch of triage or the carrying capacity of the transport vehicle will improve the performance of the algorithm. The above two graphs show that the TS-A algorithm proposed herein has a better solution and better performance for the situation where the number of wounded is larger and the carrying capacity is stronger.

6. Conclusions

For the wounded triage, transport and cooperative scheduling problem of emergency surgery in urban emergency rescue, this study adopts a supply chain-like processing method to realize the collaborative scheduling of each rescue link. It mainly considers the collaborative scheduling problem of triage, transport and surgery under the constraint factors such as the different time when the wounded arrive at the triage field and the different transport capacity of the transport vehicles. The research goal is to minimize the whole rescue time span of different batches of wounded in the rescue process; it establishes a mathematical model to arrange batches of wounded according to the capacity of the triage field or the transport capacity of the transport vehicle, analyzes the impact of batches on the whole process time, and gives the nature of the problem in general and special cases.

In order to evaluate the performance of the TS-A algorithm, this study finally gives a two-stage algorithm TS-A and constructs the lower bound of the problem domain under the ideal state. The simulation experiment shows that the TS-A algorithm proposed herein is superior to the other two algorithms mentioned in references. At the same time, for the case of a large number of wounded and a large rescue vehicle capacity, the experiment can obtain better results, which shows that the algorithm can handle large-scale collaborative scheduling problems.

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Data Availability

The data used to support the findings of this study are available from the corresponding author upon request.

Conflicts of Interest

The authors declare that they have no conflicts of interest.

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