Detection and Mitigation of Corrupted Information in Distributed Model Predictive Control Based on Resource Allocation

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- Energy Distribution System
- Traffic management
- Heat distribution
- Water distribution





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Geographically distributed

- Coupled by constraints (energy)
- Optimization objectives
 - Energy
 - User satisfaction
 - . . .
- Solution → Model Predictive Control





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- Solution \rightarrow Model Predictive Control



optimize
$$u(k:k+N_p-1|k)$$
subject to
$$\begin{aligned} & J(\boldsymbol{x}(k),\boldsymbol{u}(k)) \\ & \boldsymbol{x}(\xi+1) = f(\boldsymbol{x}(\xi),\boldsymbol{u}(\xi)) \\ & g_i(\boldsymbol{x}(\xi),\boldsymbol{u}(\xi)) \leq 0 \\ & h_j(\boldsymbol{x}(\xi),\boldsymbol{u}(\xi)) = 0 \end{aligned} \end{aligned} \begin{cases} \forall \xi \in \{1,\ldots,N_p\} \\ \forall i \in \{1,\ldots,m\} \\ \forall j \in \{1,\ldots,p\} \end{cases}$$



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$$J(\boldsymbol{x}(k), \boldsymbol{u}(k))$$

$$\forall \xi \in \{1, \dots, N_p\}$$

$$\forall i \in \{1, \dots, m\}$$

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optimize
$$u(k:k+N_p-1|k)$$

$$\mathbf{x}(\xi+1) = f(\mathbf{x}(\xi), \mathbf{u}(\xi))$$
subject to
$$\begin{aligned} \mathbf{x}(\xi+1) &= f(\mathbf{x}(\xi), \mathbf{u}(\xi)) \\ g_i(\mathbf{x}(\xi), \mathbf{u}(\xi)) &\leq 0 \\ h_j(\mathbf{x}(\xi), \mathbf{u}(\xi)) &= 0 \end{aligned} \} \begin{cases} \forall \xi \in \{1, \dots, N_p\} \\ \forall i \in \{1, \dots, m\} \\ \forall j \in \{1, \dots, p\} \end{cases}$$



minimize
$$\mathbf{u}_{(k:k+N_p-1|k)} = \sum_{j=1}^{N_p} \|\mathbf{v}(k+j|k)\|_Q^2 + \|\mathbf{u}(k+j-1|k)\|_R^2$$
subject to
$$\mathbf{x}(\xi+1) = f(\mathbf{x}(\xi), \mathbf{u}(\xi))$$

$$g_i(\mathbf{x}(\xi), \mathbf{u}(\xi)) \le 0$$

$$h_j(\mathbf{x}(\xi), \mathbf{u}(\xi)) = 0$$

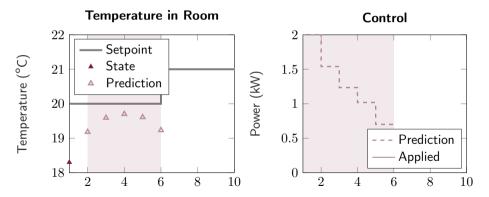
$$\forall \xi \in \{1, \dots, N_p\}$$

$$\forall i \in \{1, \dots, m\}$$

$$\forall j \in \{1, \dots, p\}$$

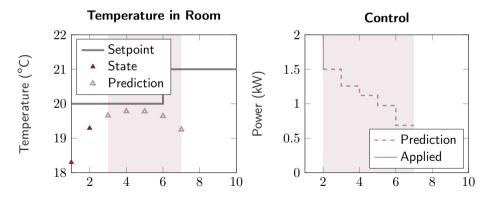


Find optimal control sequence



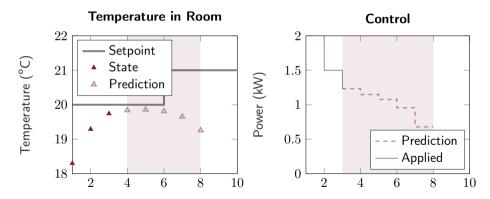


Find optimal control sequence, apply first element



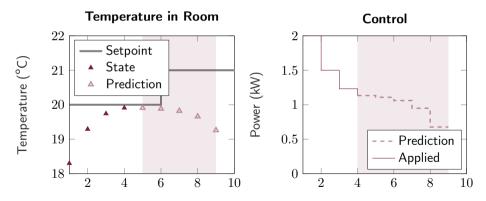


Find optimal control sequence, apply first element, rinse repeat





Find optimal control sequence, apply first element, rinse repeat ightarrow Receding Horizon



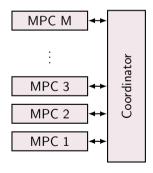


- ullet Problem: Complexity depends on N_p, m, p and sizes of $oldsymbol{x}$ and $oldsymbol{u}$
- Solution: Divide and Conquer

MPC



- Problem: Complexity depends on N_p, m, p and sizes of x and u
- Solution: Divide and Conquer





$$\underbrace{\frac{J_G(k)}{J_i(k)}}_{\substack{u_i(k:k+N_p-1|k) \\ \text{subject to}}} \sum_{i=1}^{M} \sum_{j=1}^{N_p} \|\boldsymbol{v}_i(k+j|k)\|_{Q_i}^2 + \|\boldsymbol{u}_i(k+j-1|k)\|_{R_i}^2 \\ \boldsymbol{x}_i(k+1) = A_i \boldsymbol{x}_i(k) + B_i \boldsymbol{u}_i(k) \} \ \forall i \in \{1, \dots, M\} \\ \sum_{i=1}^{M} \Gamma_i \boldsymbol{u}_i(k) = \boldsymbol{u}_{\max} \} \ \forall j \in \{1, \dots, N_p\}$$



$$J_{i}^{\star}(\boldsymbol{\theta}_{i}(k)) = \underset{\boldsymbol{u}_{i}(k:k+N_{p}-1|k)}{\operatorname{minimize}} J_{i}(k)$$
s.t. $\boldsymbol{x}_{i}(k+1) = A_{i}\boldsymbol{x}_{i}(k) + B_{i}\boldsymbol{u}_{i}(k)$

$$\Gamma_{i}\boldsymbol{u}_{i}(k) = \boldsymbol{\theta}_{i}(k) : \boldsymbol{\lambda}_{i}(k)$$

$$\forall i \in \{1, \dots, M\}$$

$$J^{\star} = \underset{\boldsymbol{\theta}(k:k+N_p-1|k)}{\text{minimize}} \sum_{i=1}^{M} J_i^{\star}(\boldsymbol{\theta}_i(k))$$
s.t.
$$\sum_{i=1}^{M} \boldsymbol{\theta}_i(k) = \boldsymbol{u}_{\text{max}}$$



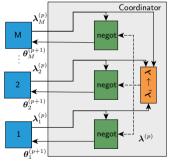
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$$\forall i \in \{1, \dots, M\}$$

$$\boldsymbol{\theta}_i^{(p+1)} = \boldsymbol{\theta}_i^{(p)} + \rho \left(\boldsymbol{\lambda}_i^{\star}(\boldsymbol{\theta}_i^{(p)}) - \frac{1}{M} \sum_{j=1}^{M} \boldsymbol{\lambda}_j^{\star}(\boldsymbol{\theta}_j^{(p)}) \right)$$





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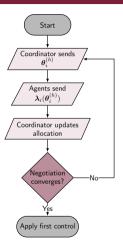


Figure 1: Quantity decomposition based DMPC



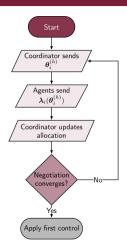


Figure 1: Quantity decomposition based DMPC



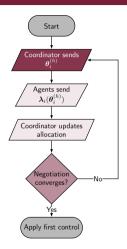


Figure 1: Quantity decomposition based DMPC



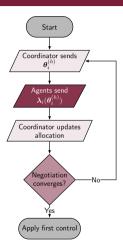


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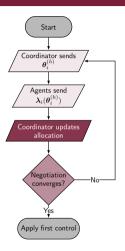


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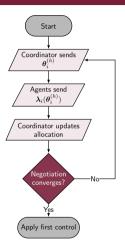


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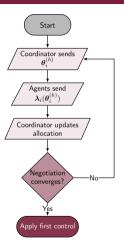


Figure 1: Quantity decomposition based DMPC



What if agents send a non-agreed λ_i ?



Outline

- Vulnerabilities in distributed MPC based on Resource Allocation Attacks
 Consequences
- 2 Securing the DMPC Analysis of Subproblems Detection Mechanism Mitigation Mechanism Complete Mechanism
- 3 Results



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Securing the DMPC
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How can a non-cooperative agent attack?

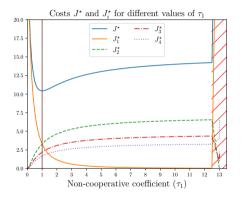
- λ_i is the only interface with coordination
- Non-cooperative agent sends $\tilde{\lambda}_i = \gamma_i(\lambda_i)$



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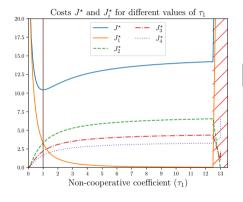
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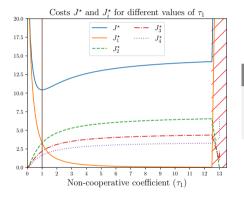
- Agent 1 is non-cooperative
- It uses $\tilde{\lambda}_1 = \gamma_1(\lambda_1) = \tau_1 I \lambda_1$





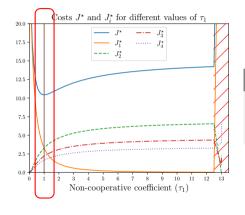
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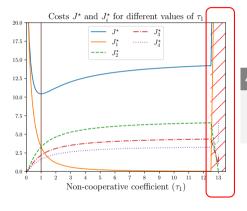
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$$\underbrace{ \begin{aligned} & J_i(k) \\ & \underset{\boldsymbol{u}_i(k:k+N_p-1|k)}{\text{minimize}} & \sum_{j=1}^{N_p} \|\boldsymbol{v}_i(k+j|k)\|_{Q_i}^2 + \|\boldsymbol{u}_i(k+j-1|k)\|_{R_i}^2 \\ & \text{s.t.} & \underbrace{ \begin{aligned} & \boldsymbol{x}_i(\xi+1) &= A_i \boldsymbol{x}_i(\xi) + B_i \boldsymbol{u}_i(\xi) \\ & \Gamma_i \boldsymbol{u}_i(\xi) &= \boldsymbol{\theta}_i(\xi) : \boldsymbol{\lambda}_i(\xi) \end{aligned} } \forall \xi \in \{1, \dots, N_p\}$$



minimize
$$\frac{J_i(\boldsymbol{\theta}_i)}{\frac{1}{2}\boldsymbol{U}_i(k)^T H_i \boldsymbol{U}_i(k) + \boldsymbol{f}_i(k)^T \boldsymbol{U}_i(k)}$$
s.t. $\Theta_i \boldsymbol{U}_i(k) = \boldsymbol{\theta}_i : \boldsymbol{\lambda}_i$



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$$\begin{aligned} & \underset{\boldsymbol{U}_{i}(k)}{\text{minimize}} & & \overbrace{\frac{1}{2}\boldsymbol{U}_{i}(k)^{T}\boldsymbol{H}_{i}\boldsymbol{U}_{i}(k) + \boldsymbol{f}_{i}(k)^{T}\boldsymbol{U}_{i}(k)}^{J_{i}(k)} \\ & \text{s.t.} & & \boldsymbol{\Theta}_{i}\boldsymbol{U}_{i}(k) = \boldsymbol{\theta}_{i}: \boldsymbol{\lambda}_{i} \\ & & \boldsymbol{\lambda}_{i} = -P_{i}\boldsymbol{\theta}_{i} - \boldsymbol{s}_{i}(k) \\ & & \text{where } P_{i} = \left(\boldsymbol{\Theta}_{i}\boldsymbol{H}_{i}^{-1}\boldsymbol{\Theta}_{i}^{T}\right)^{-1} \text{ and } \boldsymbol{s}_{i}(k) = P_{i}\boldsymbol{\Theta}_{i}\boldsymbol{H}_{i}^{-1}\boldsymbol{f}_{i}(k) \end{aligned}$$



$$\begin{array}{ll} J_i(\boldsymbol{\theta}_i) \\ \underset{\boldsymbol{U}_i(k)}{\text{minimize}} & \overline{\frac{1}{2}\boldsymbol{U}_i(k)^TH_i\boldsymbol{U}_i(k) + \boldsymbol{f}_i(k)^T\boldsymbol{U}_i(k)} \\ \text{s.t.} & \Theta_i\boldsymbol{U}_i(k) = \boldsymbol{\theta}_i: \boldsymbol{\lambda}_i \\ \\ \boldsymbol{\lambda}_i = -\underline{\boldsymbol{P}_i}\boldsymbol{\theta}_i - \boldsymbol{s}_i(k) \\ \\ \text{where } \underline{\boldsymbol{P}_i} = \left(\Theta_iH_i^{-1}\Theta_i^{\mathrm{T}}\right)^{-1} \text{ and } \boldsymbol{s}_i(k) = P_i\Theta_iH_i^{-1}\boldsymbol{f}_i(k) \end{array}$$



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Assumption

We know nominal \bar{P}_i

Assumption

Attacker chooses
$$\tilde{\lambda}_i = \gamma_i(\lambda_i) = T_i(k)\lambda_i \rightarrow -T_i(k)P_i\theta_i - T_i(k)s_i(k)$$

• We can estimate \hat{P}_i and $\hat{\tilde{s}}_i(k)$ such as:

$$\widetilde{\boldsymbol{\lambda}}_i = \gamma_i(\boldsymbol{\lambda}_i(\boldsymbol{\theta}_i)) = -\widehat{\widetilde{P}}_i(k)\boldsymbol{\theta}_i - \widehat{\widetilde{\boldsymbol{s}}}_i(k)$$

• If
$$\widehat{\tilde{P}}_i(k) \neq \bar{P}_i \to \mathsf{Attack}$$

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¹Using Recursive Least Squares

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 Attack

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¹Using Recursive Least Squares

- ullet We estimate \hat{P}_i and $\widehat{ ilde{s}}_i(k)$ simultaneously using Recursive Least Squares
- Problem: Estimation during negotiation fails
 - ullet Consecutive $oldsymbol{\lambda}_i^p$ and $oldsymbol{ heta}_i^p$ are linearly dependent o low input excitation
- Solution: Send sequence of random values of θ_i until estimates converge



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- Error $E_i(k) = \|\widehat{\tilde{P}}_i(k) \bar{P}_i\|_F$
- ullet Create threshold ϵ_P
- Indicator $d_i \in \{0,1\}$ detects the attack in agent i.
- $d_i = 1$ if $E_i(k) > \epsilon_P$, 0 otherwise



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• Main idea: Reconstruct λ_i and use in negotiation

Assumption

We suppose $\tilde{\lambda}_i = 0$ only if $\lambda_i = 0$, which implies $T_i(k)$ invertible

• Estimate the inverse of $T_i(k)$

$$\widehat{T_i(k)^{-1}} = \bar{P}_i \widehat{\tilde{P}}_i(k)^{-1}$$

$$\lambda_{i \text{rec}} = \widehat{T_i(k)^{-1}} \tilde{\hat{s}}_i = -\bar{P}_i \theta_i - \widehat{T_i(k)^{-1}} \hat{\tilde{s}}_i(k)$$



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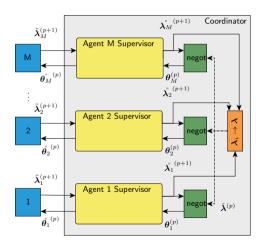
• Estimate the inverse of $T_i(k)$

$$\widehat{T_i(k)^{-1}} = \bar{P}_i \widehat{\tilde{P}}_i(k)^{-1}$$

$$\boldsymbol{\lambda}_{i\mathrm{rec}} = \widehat{T_i(k)^{-1}} \tilde{\boldsymbol{\lambda}}_i = -\bar{P}_i \boldsymbol{\theta}_i - \widehat{T_i(k)^{-1}} \hat{\tilde{\boldsymbol{s}}}_i(k)$$



Complete Mechanism



Two phases:

- Detect which agents are non-cooperative
- **2** Reconstruct λ_i and use in negotiation



Secure DMPC

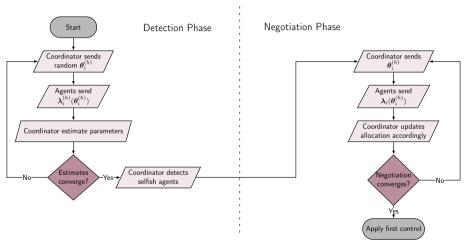


Figure 2: Secure DMPC



Secure DMPC

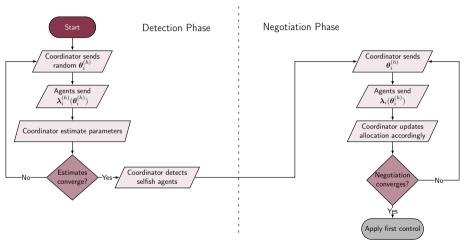


Figure 2: Secure DMPC



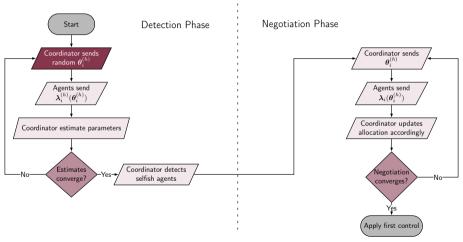


Figure 2: Secure DMPC



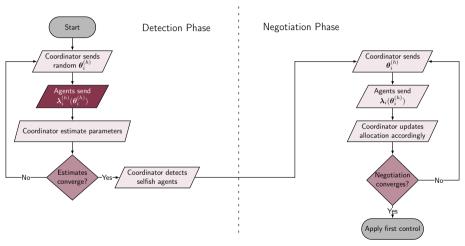


Figure 2: Secure DMPC



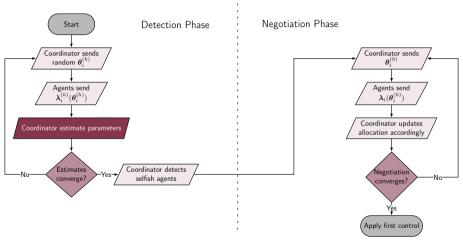


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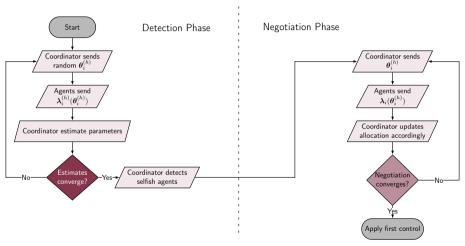


Figure 2: Secure DMPC



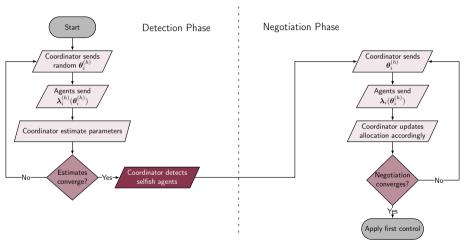


Figure 2: Secure DMPC



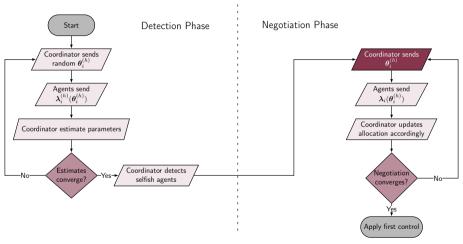


Figure 2: Secure DMPC



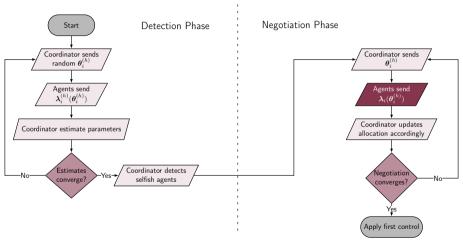


Figure 2: Secure DMPC



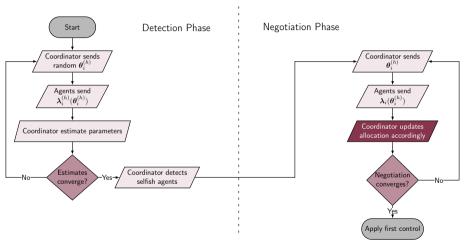


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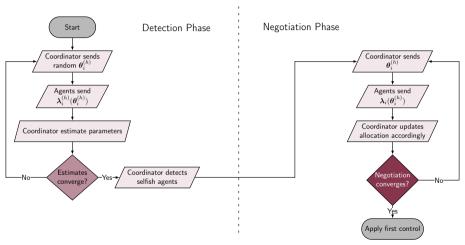


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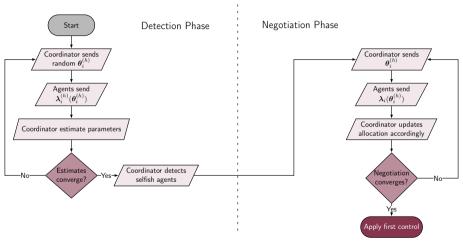


Figure 2: Secure DMPC



Outline

- Vulnerabilities in distributed MPC based on Resource Allocation Attacks Consequences
- Securing the DMPC
 Analysis of Subproblems
 Detection Mechanism
 Mitigation Mechanism
 Complete Mechanism
- 3 Results



Example

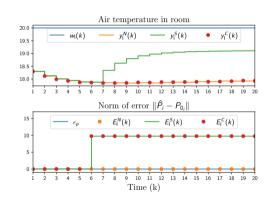
Temperature Control of 4 Distinct Rooms Under Power Scarcity

- 4 distinct rooms modeled using 3R-2C
- Initial temperature under 20°C
- ullet Not enough power to achieve setpoint $\left(\sum_{i=1}^4 oldsymbol{u}_i(k) \leq 4 \mathrm{kW} \right)$
- Simulated for a period of 5h
- ZOH $T_s = 0.25h$



Results

Temporal



- N Nominal
- S Selflish behavior
- C selfish behavior with Correction



Results

Table 1: Comparison of costs J_i^N and J_G^N

Agent	Nominal	Selfish	Selfish + correction
1	103	64	104
Ш	73	91	73
Ш	100	123	101
IV	132	154	131
Global	408	442	409



- Resource allocation based
 DMPC is vulnerable to attacks.
- Sub-problems' structure has time invariant parameters.
- Attacks can be detected using these parameters.
- 4 Effects can be mitigated.

- Outlook
 - Inequality Constraints yield Hybrid behavior
 - Non-linear attack model



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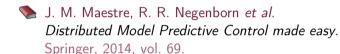


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For Further Reading I



P. Velarde, J. M. Maestre, H. Ishii, and R. R. Negenborn, "Scenario-based defense mechanism for distributed model predictive control," 2017 IEEE 56th Annual Conference on Decision and Control (CDC). IEEE, Dec 2017, pp. 6171–6176.



Thank you!

Repository https://github.com/Accacio/SysTol-21

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