Detection and Mitigation of Corrupted Information in Distributed Model Predictive Control Based on Resource Allocation

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https://git.io/JEFGW







Smart City

- Energy Distribution System
- Traffic management
- Heat distribution
- Water distribution





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- Geographically distributed
- Coupled by constraints (energy)
- Optimization objectives
 - Energy
 - User satisfaction
 - . . .
- Solution → Model Predictive Control





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optimize
$$u(k:k+N_p-1|k)$$
subject to
$$\begin{aligned} & J(\boldsymbol{x}(k),\boldsymbol{u}(k)) \\ & \boldsymbol{x}(\xi+1) = f(\boldsymbol{x}(\xi),\boldsymbol{u}(\xi)) \\ & g_i(\boldsymbol{x}(\xi),\boldsymbol{u}(\xi)) \leq 0 \\ & h_j(\boldsymbol{x}(\xi),\boldsymbol{u}(\xi)) = 0 \end{aligned} \end{aligned} \begin{cases} \forall \xi \in \{1,\ldots,N_p\} \\ \forall i \in \{1,\ldots,m\} \\ \forall j \in \{1,\ldots,p\} \end{cases}$$



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optimize
$$u(k:k+N_p-1|k)$$

$$subject to$$

$$x(\xi+1) = f(x(\xi), u(\xi))$$

$$g_i(x(\xi), u(\xi)) \le 0$$

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$$\forall \xi \in \{1, \dots, N_p\}$$

$$\forall i \in \{1, \dots, m\}$$

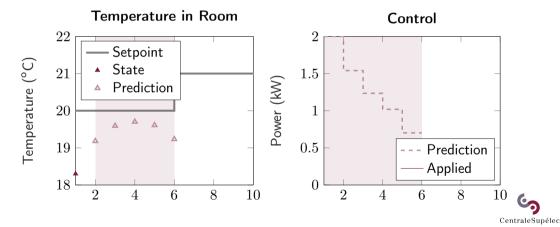
$$\forall j \in \{1, \dots, p\}$$



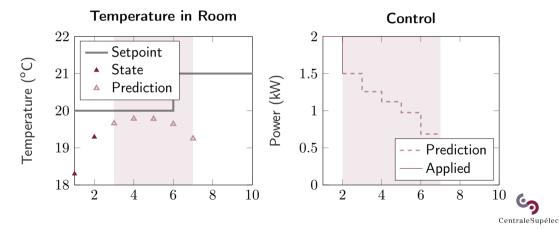
$$\begin{array}{ll} \underset{\boldsymbol{u}(k:k+N_p-1|k)}{\text{minimize}} & \sum_{j=1}^{N_p} \|\boldsymbol{v}(k+j|k)\|_Q^2 + \|\boldsymbol{u}(k+j-1|k)\|_R^2 \\ \text{subject to} & \boldsymbol{x}(\xi+1) = f(\boldsymbol{x}(\xi),\boldsymbol{u}(\xi)) \\ & g_i(\boldsymbol{x}(\xi),\boldsymbol{u}(\xi)) \leq 0 \\ & h_j(\boldsymbol{x}(\xi),\boldsymbol{u}(\xi)) = 0 \end{array} \right\} \stackrel{\forall \xi \in \{1,\ldots,N_p\}}{\forall i \in \{1,\ldots,p\}}$$



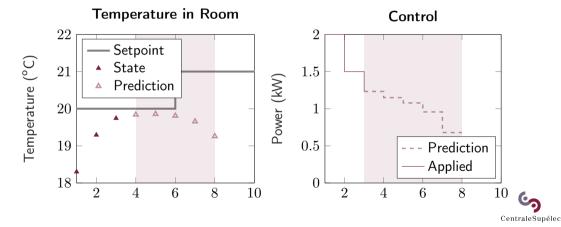
Find optimal control sequence



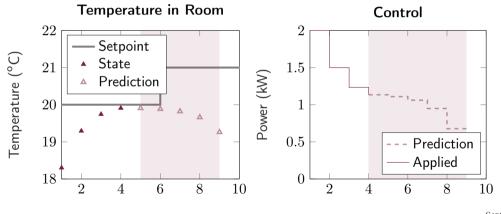
Find optimal control sequence, apply first element



Find optimal control sequence, apply first element, rinse repeat



Find optimal control sequence, apply first element, rinse repeat ightarrow Receding Horizon

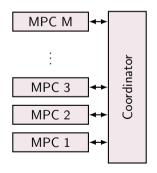


- ullet Problem: Complexity depends on N_p, m, p and sizes of $oldsymbol{x}$ and $oldsymbol{u}$
- Solution: Divide and Conquer

MPC



- Problem: Complexity depends on N_p, m, p and sizes of x and u
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$$\underbrace{\frac{J_G(k)}{J_i(k)}}_{U_i(k:k+N_p-1|k)} = \underbrace{\sum_{i=1}^{M} \sum_{j=1}^{N_p} \|\boldsymbol{v}_i(k+j|k)\|_{Q_i}^2 + \|\boldsymbol{u}_i(k+j-1|k)\|_{R_i}^2}_{Subject to} \times \underbrace{\boldsymbol{x}_i(k+1) = A_i\boldsymbol{x}_i(k) + B_i\boldsymbol{u}_i(k)}_{\sum_{i=1}^{M} \Gamma_i\boldsymbol{u}_i(k) = \boldsymbol{u}_{\max}} \forall i \in \{1,\dots,M\}$$



$$J_{i}^{\star}(\boldsymbol{\theta}_{i}(k)) = \underset{\boldsymbol{u}_{i}(k:k+N_{p}-1|k)}{\operatorname{minimize}} J_{i}(k)$$
s.t. $\boldsymbol{x}_{i}(k+1) = A_{i}\boldsymbol{x}_{i}(k) + B_{i}\boldsymbol{u}_{i}(k)$

$$\Gamma_{i}\boldsymbol{u}_{i}(k) = \boldsymbol{\theta}_{i}(k) : \boldsymbol{\lambda}_{i}(k)$$

$$\forall i \in \{1, \dots, M\}$$

$$\forall j \in \{1, \dots, N_{p}\}$$

$$J^* = \underset{\boldsymbol{\theta}(k:k+N_p-1|k)}{\text{minimize}} \sum_{i=1}^{M} J_i^*(\boldsymbol{\theta}_i(k))$$

s.t. $\sum_{i=1}^{M} \boldsymbol{\theta}_i(k) = \boldsymbol{u}_{\text{max}}$



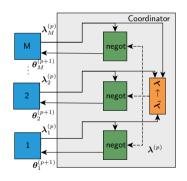
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$$\forall i \in \{1, \dots, M\}$$

$$\forall j \in \{1, \dots, N_p\}$$

$$\boldsymbol{\theta}_i^{(p+1)} = \boldsymbol{\theta}_i^{(p)} + \rho \left(\boldsymbol{\lambda}_i^{\star}(\boldsymbol{\theta}_i^{(p)}) - \frac{1}{M} \sum_{j=1}^{M} \boldsymbol{\lambda}_j^{\star}(\boldsymbol{\theta}_j^{(p)}) \right)$$





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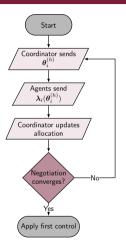


Figure 1: Quantity decomposition based DMPC



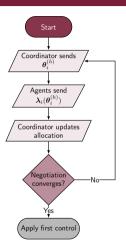


Figure 1: Quantity decomposition based DMPC



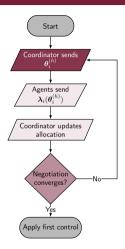


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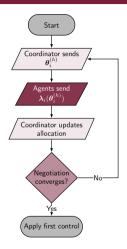


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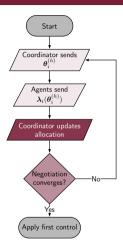


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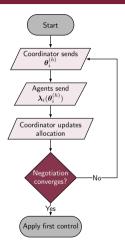


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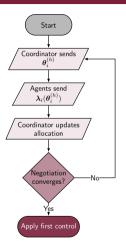


Figure 1: Quantity decomposition based DMPC



What if agents send a non-agreed λ_i ?



Outline

- Vulnerabilities in distributed MPC based on Resource Allocation Attacks Consequences
- Securing the DMPC
 Analysis of Subproblems
 Detection Mechanism
 Mitigation Mechanism
 Complete Mechanism
- 3 Results



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How can a non-cooperative agent attack?

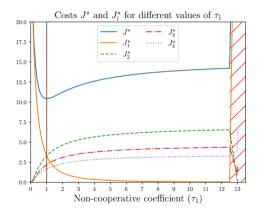
- ullet $oldsymbol{\lambda}_i$ is the only interface with coordination
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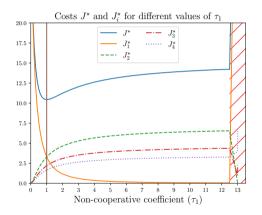
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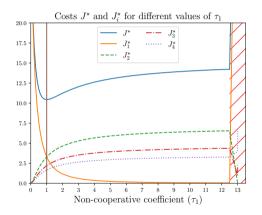
- Agent 1 is non-cooperative
- It uses $\tilde{\lambda}_1 = \gamma_1(\lambda_1) = \tau_1 I \lambda_1$





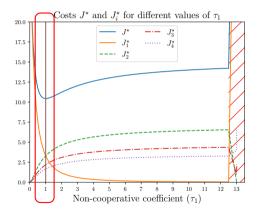
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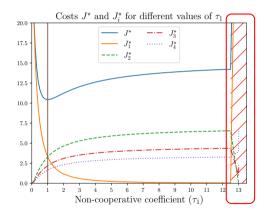
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$$\underbrace{\frac{J_i(k)}{\mathbf{w}_i(k:k+N_p-1|k)}}_{\mathbf{u}_i(k:k+N_p-1|k)} \underbrace{\sum_{j=1}^{N_p} \|\mathbf{v}_i(k+j|k)\|_{Q_i}^2 + \|\mathbf{u}_i(k+j-1|k)\|_{R_i}^2}_{\text{s.t.}}$$

$$\text{s.t.} \underbrace{\frac{\mathbf{x}_i(\xi+1) = A_i\mathbf{x}_i(\xi) + B_i\mathbf{u}_i(\xi)}{\Gamma_i\mathbf{u}_i(\xi) = \boldsymbol{\theta}_i(\xi) : \boldsymbol{\lambda}_i(\xi)}}_{\boldsymbol{I}_i(\xi)} \forall \xi \in \{1, \dots, N_p\}$$



minimize
$$\frac{J_i(\boldsymbol{\theta}_i)}{\frac{1}{2}\boldsymbol{U}_i(k)^T H_i \boldsymbol{U}_i(k) + \boldsymbol{f}_i(k)^T \boldsymbol{U}_i(k)}$$
s.t. $\Theta_i \boldsymbol{U}_i(k) = \boldsymbol{\theta}_i : \boldsymbol{\lambda}_i$



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$$\begin{aligned} & \underbrace{\frac{J_i(\boldsymbol{\theta}_i)}{2}\boldsymbol{U}_i(k)}^{} & \underbrace{\frac{1}{2}\boldsymbol{U}_i(k)^TH_i\boldsymbol{U}_i(k) + \boldsymbol{f}_i(k)^T\boldsymbol{U}_i(k)}_{\text{S.t.}} \\ & \text{S.t.} & \boldsymbol{\Theta}_i\boldsymbol{U}_i(k) = \boldsymbol{\theta}_i: \boldsymbol{\lambda}_i \\ & \boldsymbol{\lambda}_i = -P_i\boldsymbol{\theta}_i - \boldsymbol{s}_i(k) \end{aligned}$$
 where $P_i = \left(\boldsymbol{\Theta}_iH_i^{-1}\boldsymbol{\Theta}_i^{\mathrm{T}}\right)^{-1}$ and $\boldsymbol{s}_i(k) = P_i\boldsymbol{\Theta}_iH_i^{-1}\boldsymbol{f}_i(k)$



$$\begin{aligned} & \underset{\boldsymbol{U}_{i}(k)}{\text{minimize}} & & \overbrace{\frac{1}{2}\boldsymbol{U}_{i}(k)^{T}H_{i}\boldsymbol{U}_{i}(k) + \boldsymbol{f}_{i}(k)^{T}\boldsymbol{U}_{i}(k)} \\ & \text{s.t.} & & \boldsymbol{\Theta}_{i}\boldsymbol{U}_{i}(k) = \boldsymbol{\theta}_{i}:\boldsymbol{\lambda}_{i} \\ & & \boldsymbol{\lambda}_{i} = -\boldsymbol{P}_{i}\boldsymbol{\theta}_{i} - \boldsymbol{s}_{i}(k) \\ \end{aligned}$$
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Assumption

We know nominal \bar{P}_i

Assumption

Attacker chooses
$$ilde{m{\lambda}}_i = \gamma_i(m{\lambda}_i) = T_i(k)m{\lambda}_i o -T_i(k)P_im{ heta}_i - T_i(k)m{s}_i(k)$$

• We can estimate \hat{P}_i and $\hat{\tilde{s}}_i(k)$ such as:

$$\widetilde{\boldsymbol{\lambda}}_i = \gamma_i(\boldsymbol{\lambda}_i(\boldsymbol{\theta}_i)) = -\widehat{\tilde{P}}_i(k)\boldsymbol{\theta}_i - \widehat{\tilde{\boldsymbol{s}}}_i(k)$$

• If
$$\widehat{\tilde{P}}_i(k) \neq \bar{P}_i \to \mathsf{Attack}$$

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¹ Ising Recursive Least Squares

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 - Consecutive λ_i^p and θ_i^p are linearly dependent \to low input excitation
- Solution: Send sequence of random values of θ_i until estimates converge



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- Error $E_i(k) = \|\widehat{\tilde{P}}_i(k) \bar{P}_i\|_F$
- ullet Create threshold ϵ_P
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Mitigation

ullet Main idea: Reconstruct $oldsymbol{\lambda}_i$ and use in negotiation

Assumption

We suppose $\tilde{\lambda}_i = \mathbf{0}$ only if $\lambda_i = \mathbf{0}$, which implies $T_i(k)$ invertible

• Estimate the inverse of $T_i(k)$

$$\widehat{T_i(k)^{-1}} = \bar{P}_i \widehat{\tilde{P}}_i(k)^{-1}$$

Reconstruct λ_i

$$\lambda_{i \text{rec}} = \widehat{T_i(k)^{-1}} \tilde{\lambda}_i = -\bar{P}_i \theta_i - \widehat{T_i(k)^{-1}} \hat{\tilde{s}}_i(k)$$



Mitigation Mechanism

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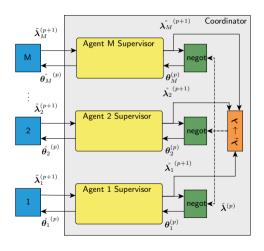
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Complete Mechanism



Two phases:

- Detect which agents are non-cooperative
- **2** Reconstruct λ_i and use in negotiation



Secure DMPC

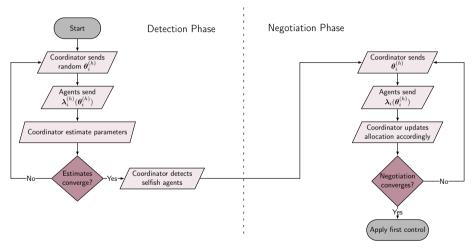




Figure 2: Secure DMPC

Secure DMPC

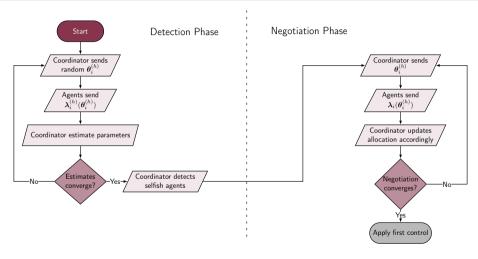




Figure 2: Secure DMPC

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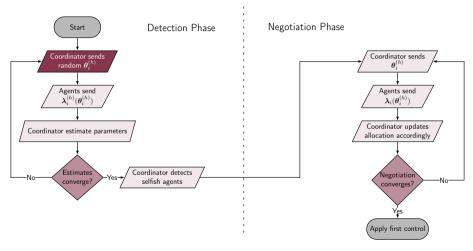




Figure 2: Secure DMPC

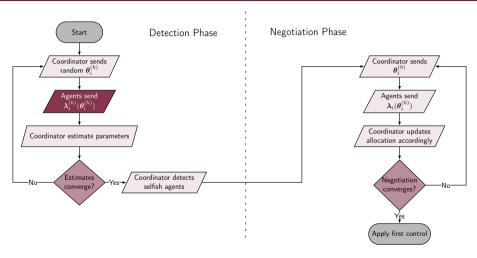




Figure 2: Secure DMPC

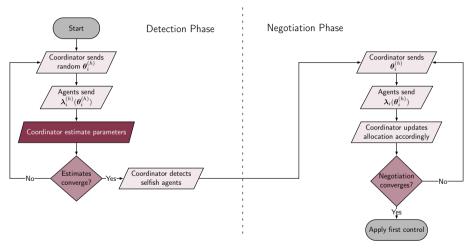




Figure 2: Secure DMPC

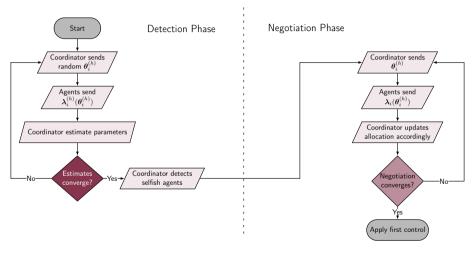




Figure 2: Secure DMPC

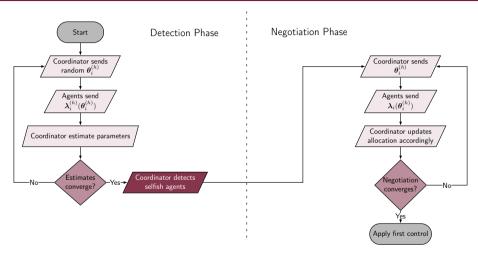




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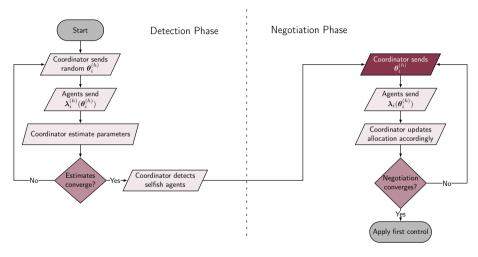




Figure 2: Secure DMPC

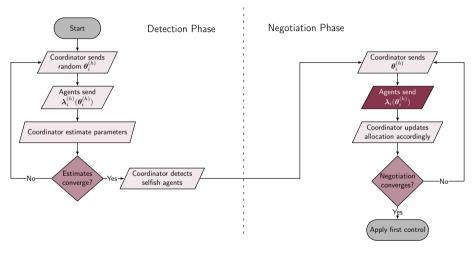




Figure 2: Secure DMPC

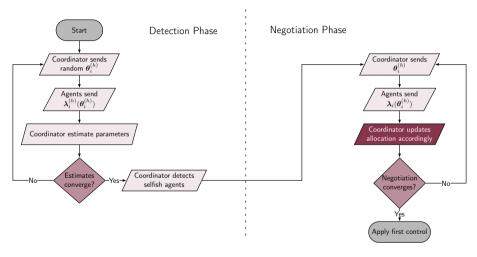




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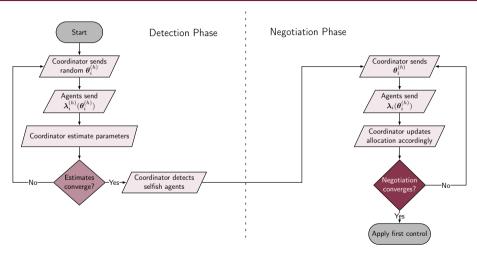




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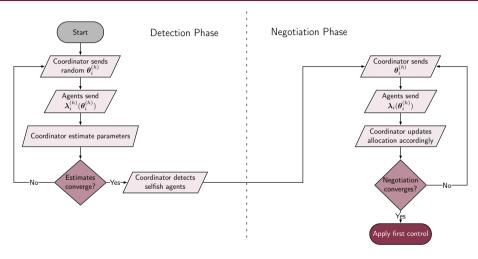




Figure 2: Secure DMPC

Outline

- Vulnerabilities in distributed MPC based on Resource Allocation Attacks
 Consequences
- Securing the DMPC
 Analysis of Subproblems
 Detection Mechanism
 Mitigation Mechanism
 Complete Mechanism
- 3 Results



Example

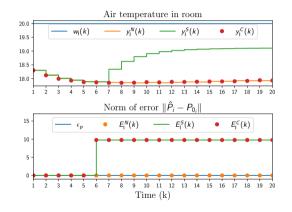
Temperature Control of 4 Distinct Rooms Under Power Scarcity

- 4 distinct rooms modeled using 3R-2C
- Initial temperature under 20°C
- ullet Not enough power to achieve setpoint $\left(\sum_{i=1}^4 oldsymbol{u}_i(k) \leq 4 \mathrm{kW}
 ight)$
- Simulated for a period of 5h
- ZOH $T_s = 0.25h$



Results

Temporal



- Nominal
- S Selflish behavior
- C selfish behavior with Correction



Results

Table 1: Comparison of costs J_i^N and J_G^N

Agent	Nominal	Selfish	Selfish + correction
1	103	64	104
П	73	91	73
Ш	100	123	101
IV	132	154	131
Global	408	442	409



- Resource allocation based DMPC is vulnerable to attacks.
- Sub-problems' structure has time invariant parameters.
- Attacks can be detected using these parameters.
- 4 Effects can be mitigated.

- Outlook
 - Inequality Constraints yield Hybrid behavior
 - Non-linear attack model



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For Further Reading I



J. M. Maestre, R. R. Negenborn et al. Distributed Model Predictive Control made easy. Springer, 2014, vol. 69.



P. Velarde, J. M. Maestre, H. Ishii, and R. R. Negenborn. "Scenario-based defense mechanism for distributed model predictive control." 2017 IEEE 56th Annual Conference on Decision and Control (CDC). IEEE, Dec 2017, pp. 6171–6176.



Questions?

 ${\it Repository } \\ {\it https://github.com/Accacio/SysTol-21}$



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