

Detection and Mitigation of Corrupted Information in Distributed Model Predictive Control Based on Resource Allocation

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<https://git.io/JEFGW>





- Energy Distribution System
- Traffic management
- Heat distribution
- Water distribution
- ...



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- Geographically distributed
- Coupled by constraints (energy)
- Optimization objectives
 - Energy
 - User satisfaction
 - ...
- Solution → Model Predictive Control



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Objective: Find control input sequence that optimizes an objective function

$$\begin{array}{ll} \underset{\mathbf{u}(k:k+N_p-1|k)}{\text{optimize}} & J(\mathbf{x}(k), \mathbf{u}(k)) \\ \text{subject to} & \left. \begin{array}{l} \mathbf{x}(\xi + 1) = f(\mathbf{x}(\xi), \mathbf{u}(\xi)) \\ g_i(\mathbf{x}(\xi), \mathbf{u}(\xi)) \leq 0 \\ h_j(\mathbf{x}(\xi), \mathbf{u}(\xi)) = 0 \end{array} \right\} \begin{array}{l} \forall \xi \in \{1, \dots, N_p\} \\ \forall i \in \{1, \dots, m\} \\ \forall j \in \{1, \dots, p\} \end{array} \end{array}$$

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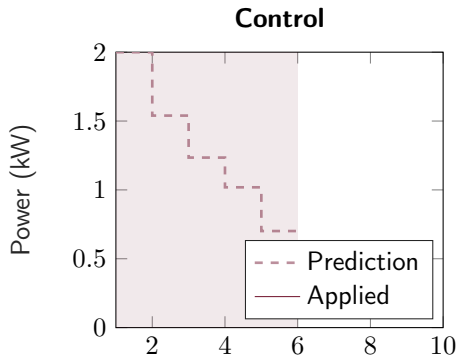
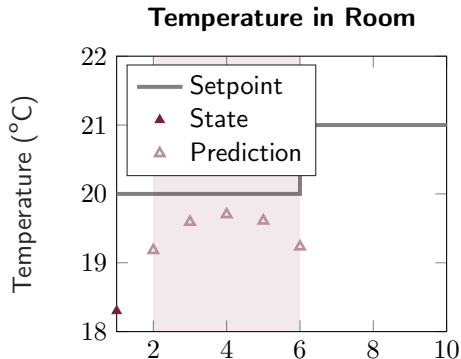
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$$\begin{array}{ll} \underset{\mathbf{u}(k:k+N_p-1|k)}{\text{minimize}} & \sum_{j=1}^{N_p} \|\mathbf{v}(k+j|k)\|_Q^2 + \|\mathbf{u}(k+j-1|k)\|_R^2 \\ \text{subject to} & \left. \begin{array}{l} \mathbf{x}(\xi+1) = f(\mathbf{x}(\xi), \mathbf{u}(\xi)) \\ g_i(\mathbf{x}(\xi), \mathbf{u}(\xi)) \leq 0 \\ h_j(\mathbf{x}(\xi), \mathbf{u}(\xi)) = 0 \end{array} \right\} \begin{array}{l} \forall \xi \in \{1, \dots, N_p\} \\ \forall i \in \{1, \dots, m\} \\ \forall j \in \{1, \dots, p\} \end{array} \end{array}$$

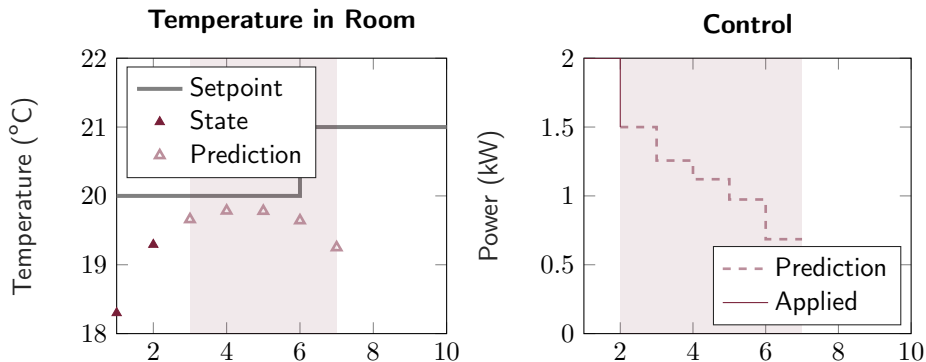
Model Predictive Control

Find optimal control sequence



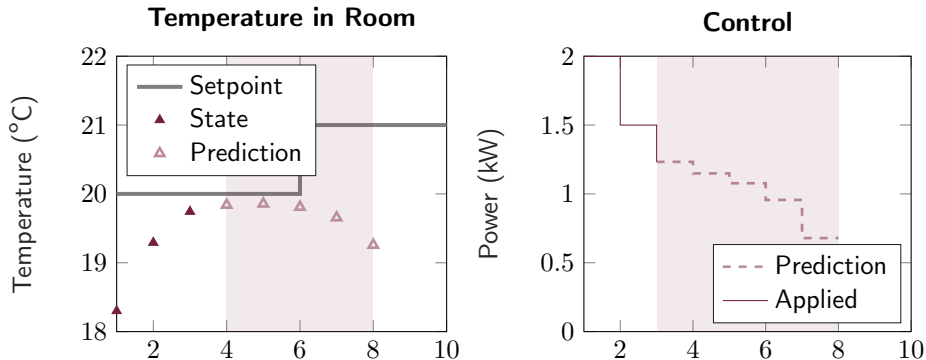
Model Predictive Control

Find optimal control sequence, apply first element



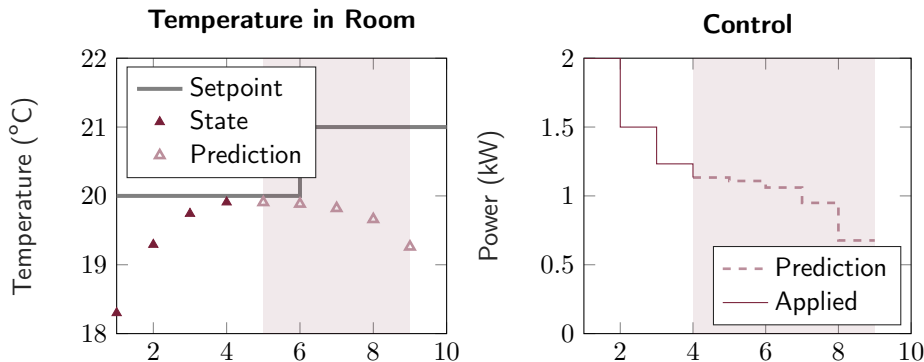
Model Predictive Control

Find optimal control sequence, apply first element, rinse repeat



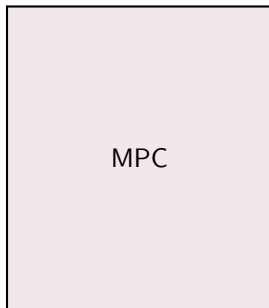
Model Predictive Control

Find optimal control sequence, apply first element, rinse repeat → Receding Horizon



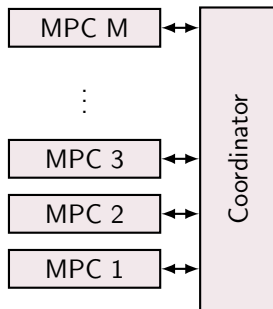
Distributed Model Predictive Control

- Problem: Complexity depends on N_p, m, p and sizes of x and u
- Solution: Divide and Conquer



Distributed Model Predictive Control

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Distributed Model Predictive Control

Quantity Decomposition | Resource Allocation

$$\begin{aligned} & \text{minimize}_{\mathbf{u}_i(k:k+N_p-1|k)} \quad \overbrace{\sum_{i=1}^M \sum_{j=1}^{N_p} \|\mathbf{v}_i(k+j|k)\|_{Q_i}^2 + \|\mathbf{u}_i(k+j-1|k)\|_{R_i}^2}^{J_G(k)} \\ & \text{subject to} \quad \left. \begin{aligned} \mathbf{x}_i(k+1) &= A_i \mathbf{x}_i(k) + B_i \mathbf{u}_i(k) \\ \sum_{i=1}^M \Gamma_i \mathbf{u}_i(k) &= \mathbf{u}_{\max} \end{aligned} \right\} \begin{aligned} & \forall i \in \{1, \dots, M\} \\ & \forall j \in \{1, \dots, N_p\} \end{aligned} \end{aligned}$$



Distributed Model Predictive Control

Quantity Decomposition | Resource Allocation

$$\left. \begin{aligned} J_i^*(\theta_i(k)) &= \underset{\mathbf{u}_i(k:k+N_p-1|k)}{\text{minimize}} && J_i(k) \\ \text{s.t.} \quad \mathbf{x}_i(k+1) &= A_i \mathbf{x}_i(k) + B_i \mathbf{u}_i(k) \\ \Gamma_i \mathbf{u}_i(k) &= \theta_i(k) : \lambda_i(k) \end{aligned} \right\} \begin{aligned} &\forall i \in \{1, \dots, M\} \\ &\forall j \in \{1, \dots, N_p\} \end{aligned}$$

$$\begin{aligned} J^* &= \underset{\theta(k:k+N_p-1|k)}{\text{minimize}} \sum_{i=1}^M J_i^*(\theta_i(k)) \\ \text{s.t.} \quad \sum_{i=1}^M \theta_i(k) &= \mathbf{u}_{\max} \end{aligned}$$



Distributed Model Predictive Control

Quantity Decomposition | Resource Allocation

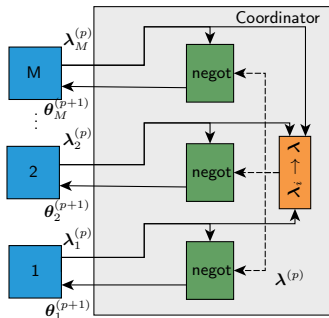
$$\left. \begin{aligned} J_i^*(\boldsymbol{\theta}_i(k)) &= \underset{\mathbf{u}_i(k:k+N_p-1|k)}{\text{minimize}} \quad J_i(k) \\ \text{s.t.} \quad \mathbf{x}_i(k+1) &= A_i \mathbf{x}_i(k) + B_i \mathbf{u}_i(k) \\ \Gamma_i \mathbf{u}_i(k) &= \boldsymbol{\theta}_i(k) : \boldsymbol{\lambda}_i(k) \end{aligned} \right\} \begin{aligned} &\forall i \in \{1, \dots, M\} \\ &\forall j \in \{1, \dots, N_p\} \end{aligned}$$

$$\boldsymbol{\theta}_i^{(p+1)} = \boldsymbol{\theta}_i^{(p)} + \rho \left(\boldsymbol{\lambda}_i^*(\boldsymbol{\theta}_i^{(p)}) - \frac{1}{M} \sum_{j=1}^M \boldsymbol{\lambda}_j^*(\boldsymbol{\theta}_j^{(p)}) \right)$$



Distributed Model Predictive Control

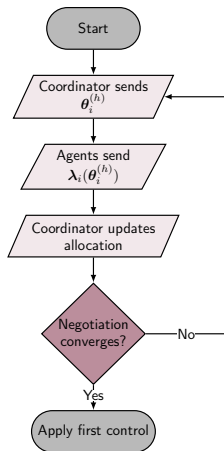
Quantity Decomposition | Resource Allocation



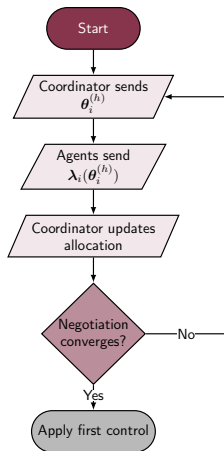
$$J_i^*(\theta_i(k)) = \begin{cases} \text{minimize}_{\mathbf{u}_i(k:k+N_p-1|k)} J_i(k) \\ \text{s.t. } \mathbf{x}_i(k+1) = A_i \mathbf{x}_i(k) + B_i \mathbf{u}_i(k) \\ \Gamma_i \mathbf{u}_i(k) = \theta_i(k) : \lambda_i(k) \end{cases} \quad \left. \begin{matrix} \forall i \in \{1, \dots, M\} \\ \forall j \in \{1, \dots, N_p\} \end{matrix} \right\}$$

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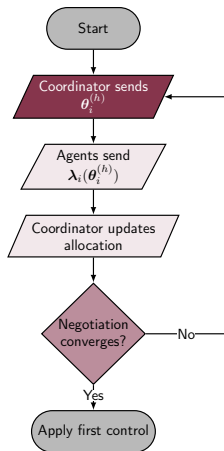
Quantity Decomposition | Resource Allocation



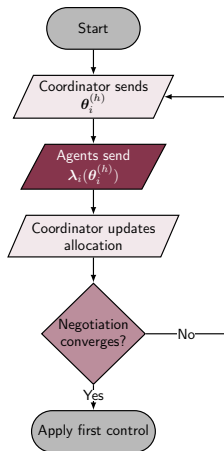
Quantity Decomposition | Resource Allocation



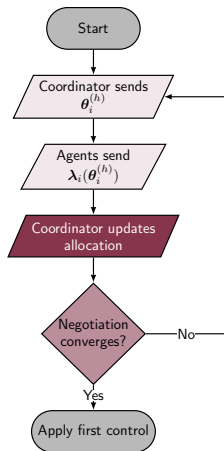
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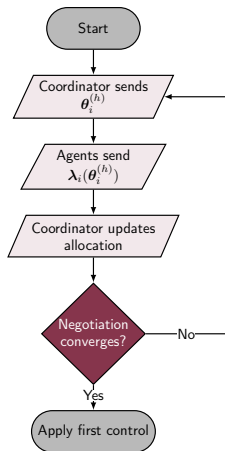
Quantity Decomposition | Resource Allocation



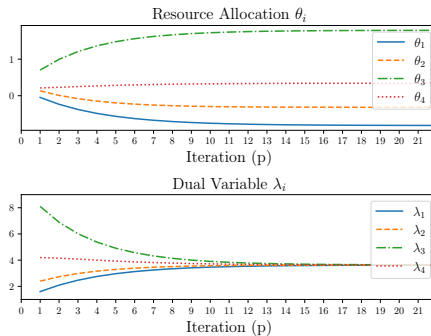
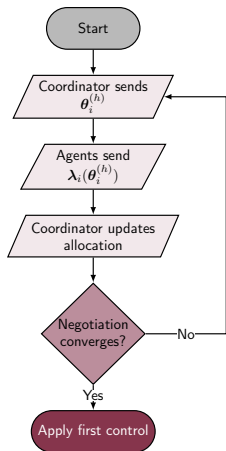
Quantity Decomposition | Resource Allocation



Quantity Decomposition | Resource Allocation



Quantity Decomposition | Resource Allocation



What if agents send a non-agreed λ_i ?



- 1 Vulnerabilities in distributed MPC based on Resource Allocation
Attacks
Consequences
- 2 Securing the DMPC
Analysis of Subproblems
Detection Mechanism
Mitigation Mechanism
Complete Mechanism
- 3 Results

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How can a non-cooperative agent attack?

- λ_i is the only interface with coordination
- Non-cooperative agent sends $\tilde{\lambda}_i = \gamma_i(\lambda_i)$

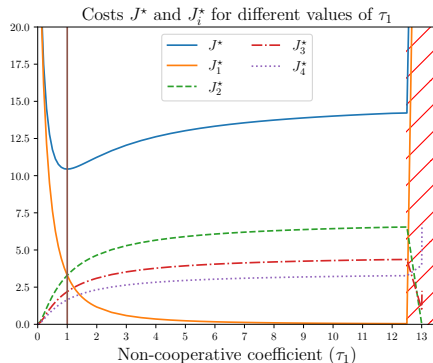


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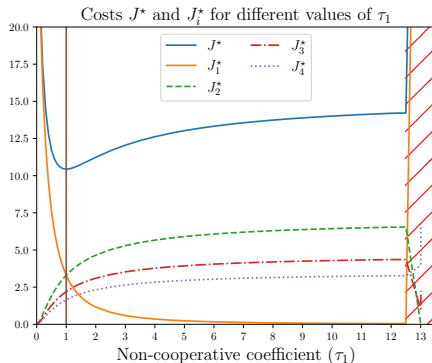
Example



4 distinct agents

- Agent 1 is non-cooperative
- It uses $\tilde{\lambda}_1 = \gamma_1(\lambda_1) = \tau_1 I \lambda_1$

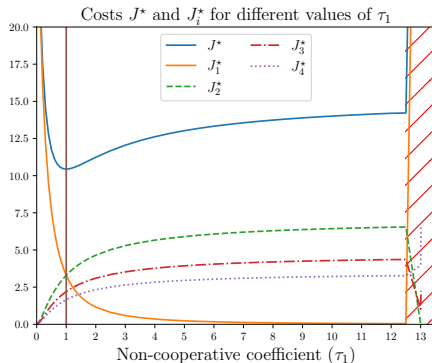
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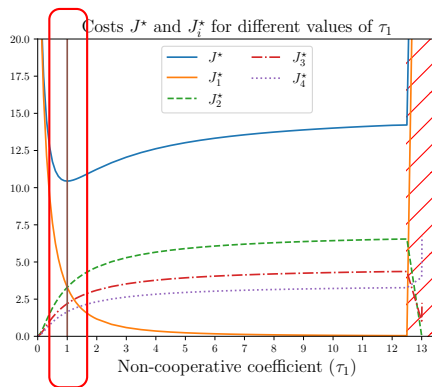
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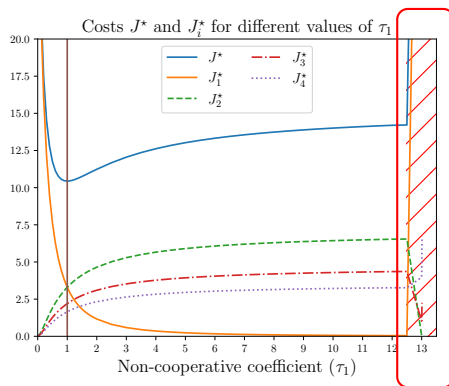
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Quadratic Case

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 & \text{s.t.} \quad \left. \begin{aligned} \mathbf{x}_i(\xi+1) &= A_i \mathbf{x}_i(\xi) + B_i \mathbf{u}_i(\xi) \\ \Gamma_i \mathbf{u}_i(\xi) &= \boldsymbol{\theta}_i(\xi) : \boldsymbol{\lambda}_i(\xi) \end{aligned} \right\} \forall \xi \in \{1, \dots, N_p\}
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$$\begin{array}{ll} \underset{\mathbf{U}_i(k)}{\text{minimize}} & \overbrace{\frac{1}{2} \mathbf{U}_i(k)^T \mathbf{H}_i \mathbf{U}_i(k) + \mathbf{f}_i(k)^T \mathbf{U}_i(k)}^{J_i(\boldsymbol{\theta}_i)} \\ \text{s.t.} & \boldsymbol{\Theta}_i \mathbf{U}_i(k) = \boldsymbol{\theta}_i : \boldsymbol{\lambda}_i \end{array}$$

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 \end{aligned}$$

$$\boldsymbol{\lambda}_i = -\mathbf{P}_i \boldsymbol{\theta}_i - \mathbf{s}_i(k)$$

where $\mathbf{P}_i = (\boldsymbol{\Theta}_i \mathbf{H}_i^{-1} \boldsymbol{\Theta}_i^T)^{-1}$ and $\mathbf{s}_i(k) = \mathbf{P}_i \boldsymbol{\Theta}_i \mathbf{H}_i^{-1} \mathbf{f}_i(k)$

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Detection

Assumption

We know nominal \bar{P}_i

Assumption

*Attacker chooses $\tilde{\lambda}_i = \gamma_i(\lambda_i) = T_i(k)\lambda_i$
 $-T_i(k)P_i\theta_i - T_i(k)s_i(k) \rightarrow -\tilde{P}_i\theta_i - \tilde{s}_i(k)$*

- We can estimate¹ \hat{P}_i and $\hat{\tilde{s}}_i(k)$ such as:

$$\tilde{\lambda}_i = \gamma_i(\lambda_i(\theta_i)) = -\hat{\tilde{P}}_i(k)\theta_i - \hat{\tilde{s}}_i(k)$$

- If $\hat{\tilde{P}}_i(k) \neq \bar{P}_i \rightarrow \text{Attack}$

¹Using Recursive Least Squares

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¹Using Recursive Least Squares



About Estimation

- We estimate \hat{P}_i and $\hat{s}_i(k)$ simultaneously using Recursive Least Squares
- Problem: Estimation during negotiation fails
 - Consecutive λ_i^p and θ_i^p are linearly dependent \rightarrow low input excitation
- Solution: Send sequence of random values of θ_i until estimates converge



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Detection

In detail

- Error $E_i(k) = \|\hat{\tilde{P}}_i(k) - \bar{P}_i\|_F$
- Create threshold ϵ_P
- Indicator $d_i \in \{0, 1\}$ detects the attack in agent i .
- $d_i = 1$ if $E_i(k) > \epsilon_P$, 0 otherwise



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Mitigation

- Main idea: Reconstruct λ_i and use in negotiation

Assumption

We suppose $\tilde{\lambda}_i = 0$ only if $\lambda_i = 0$, which implies $T_i(k)$ invertible.

- Estimate the inverse of $T_i(k)$

$$\widehat{T_i(k)^{-1}} = \bar{P}_i \hat{\tilde{P}}_i(k)^{-1}$$

- Reconstruct λ_i

$$\lambda_{i\text{rec}} = \widehat{T_i(k)^{-1}} \tilde{\lambda}_i = -\bar{P}_i \theta_i - \widehat{T_i(k)^{-1}} \hat{\tilde{s}}_i(k)$$



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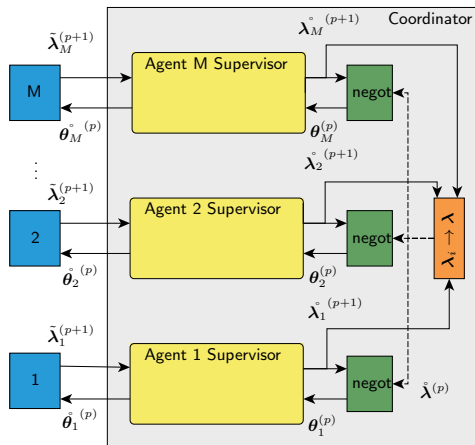
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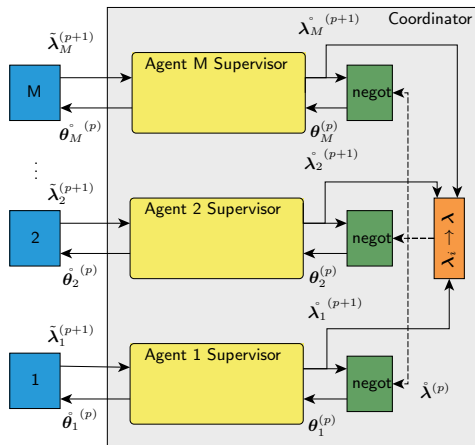
Complete Mechanism



Two phases:

- 1 Detect which agents are non-cooperative
- 2 Reconstruct λ_i and use in negotiation

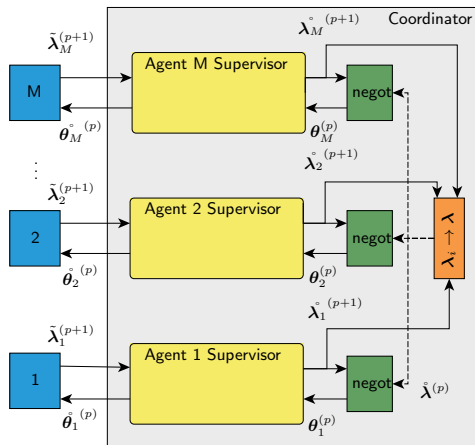
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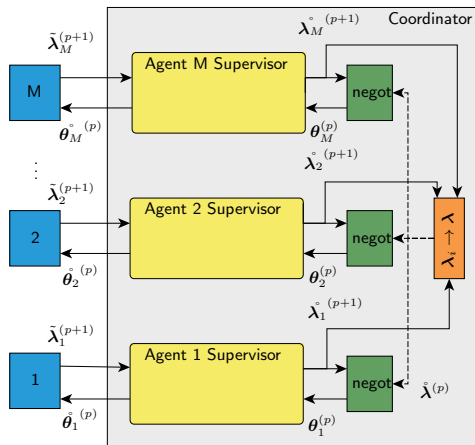
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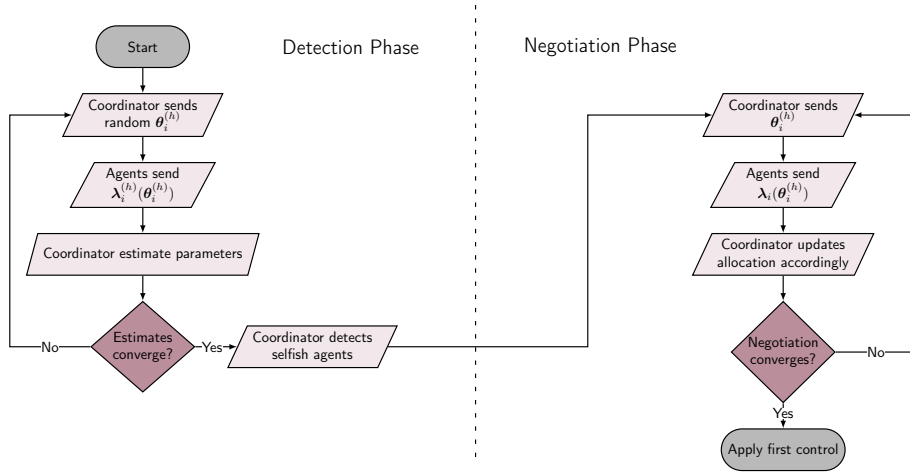
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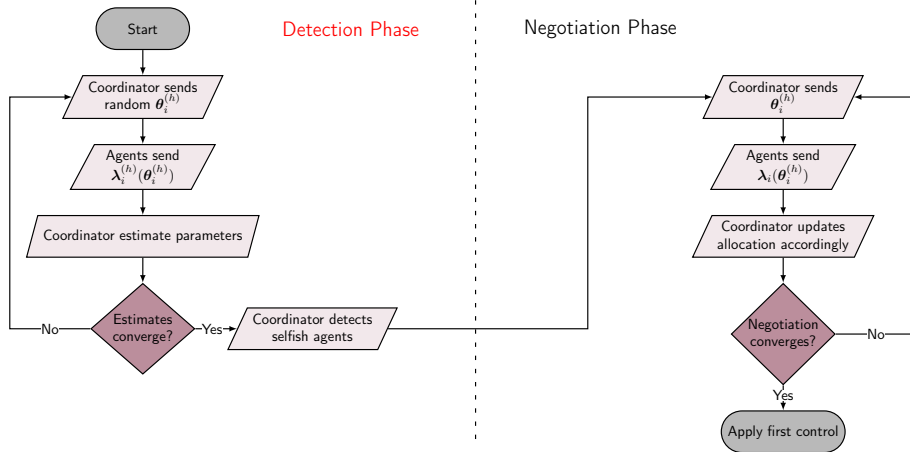
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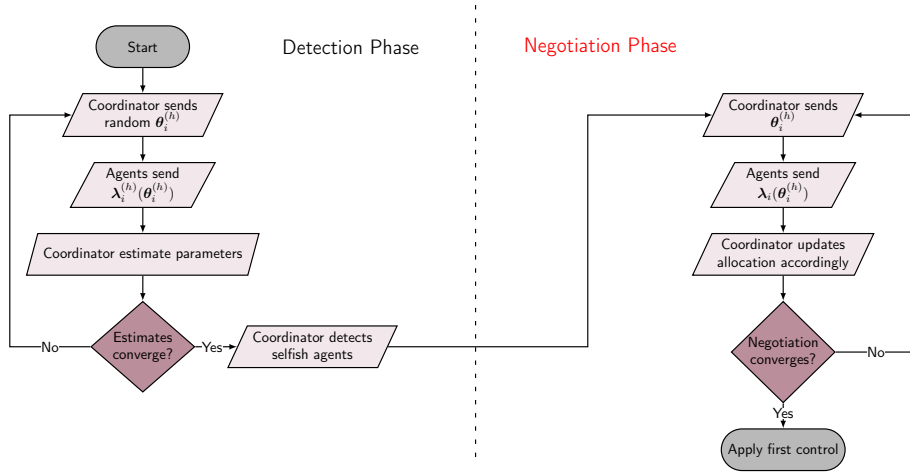
Secure DMPC



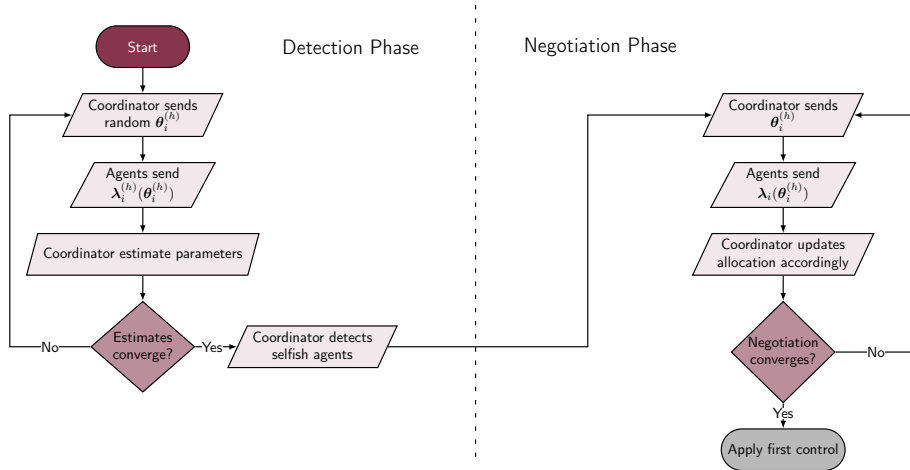
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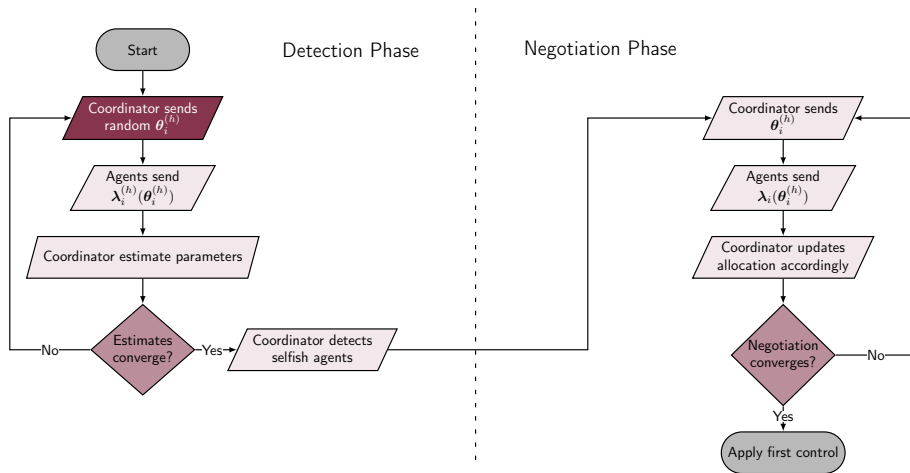
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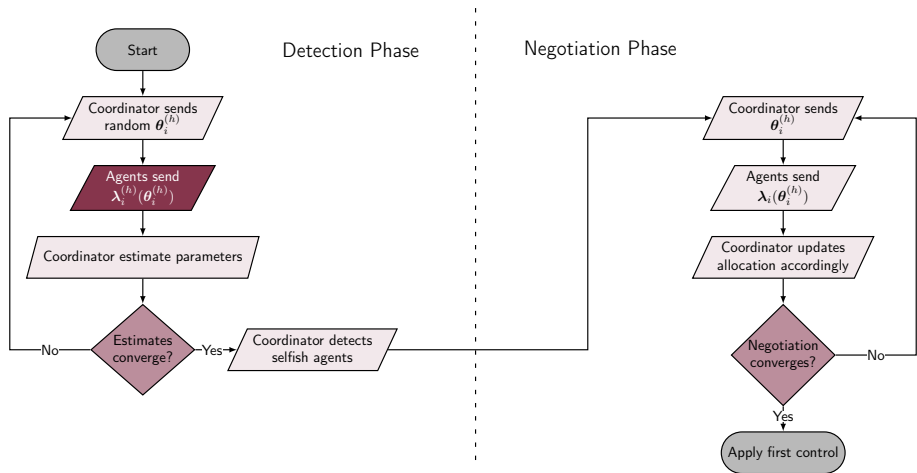
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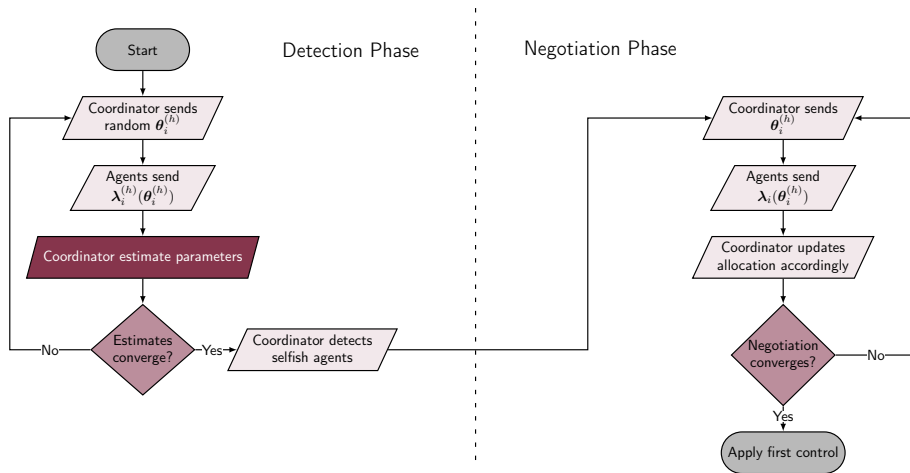
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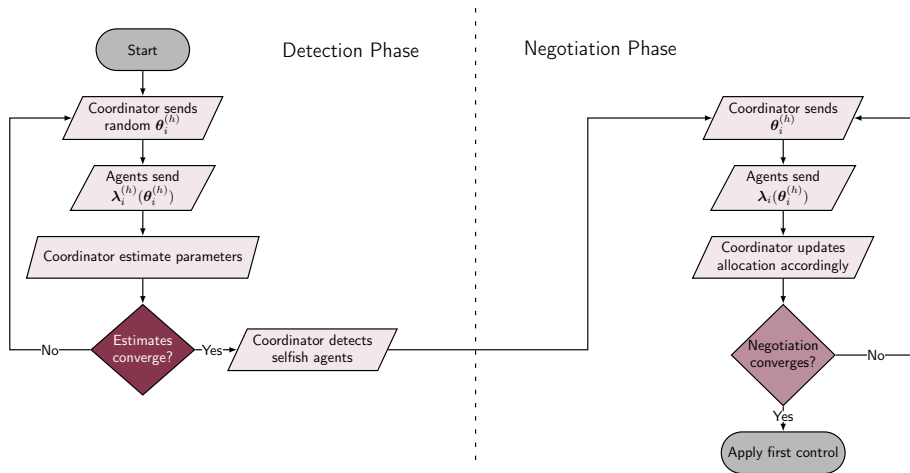
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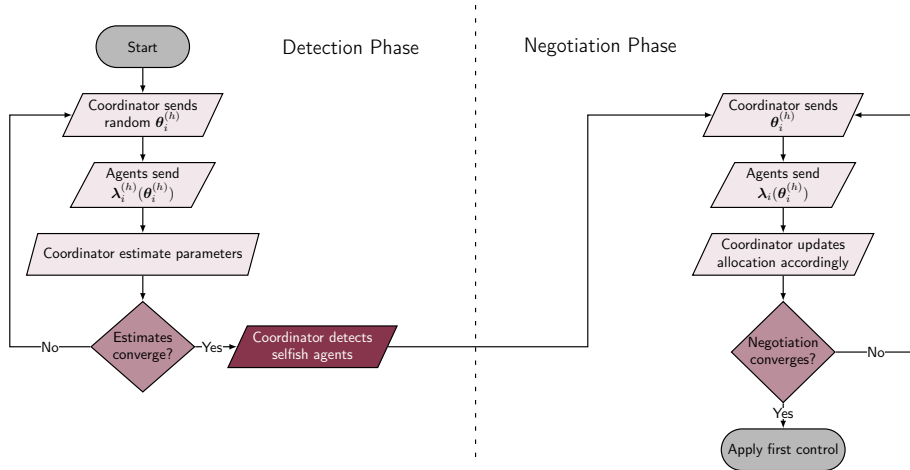
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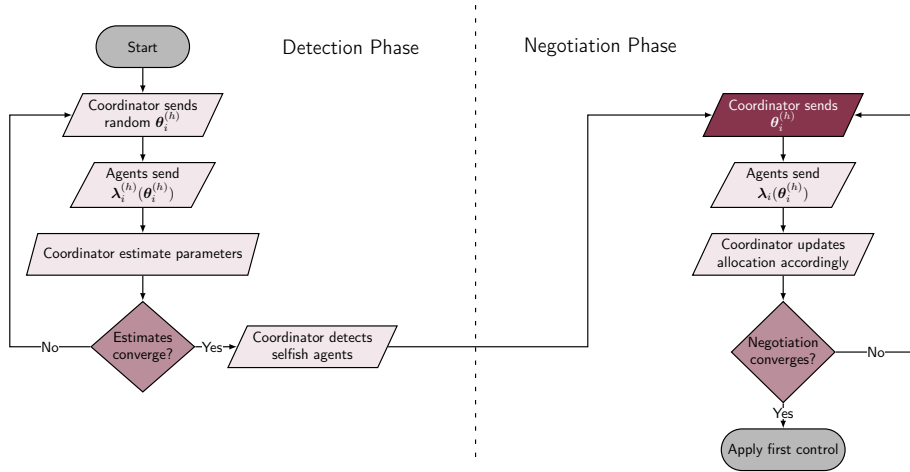
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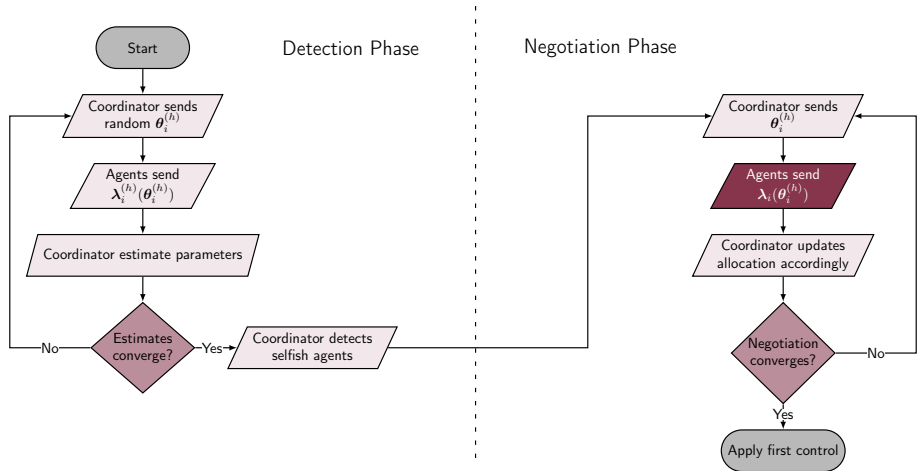
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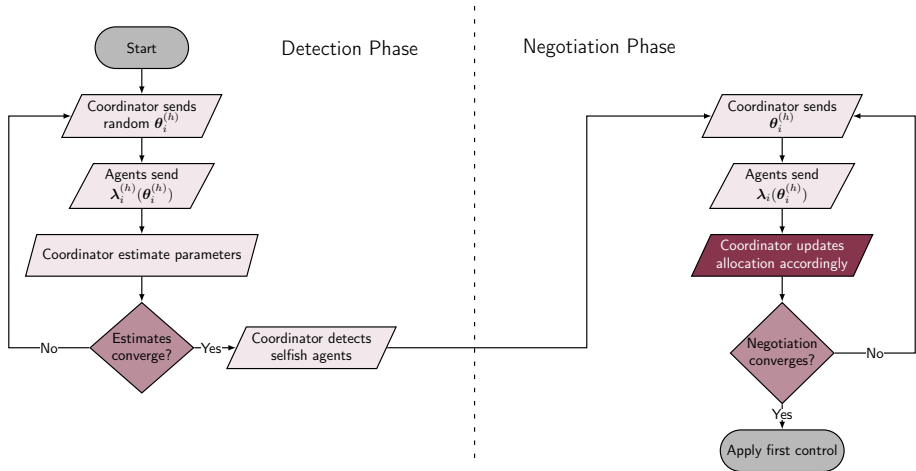
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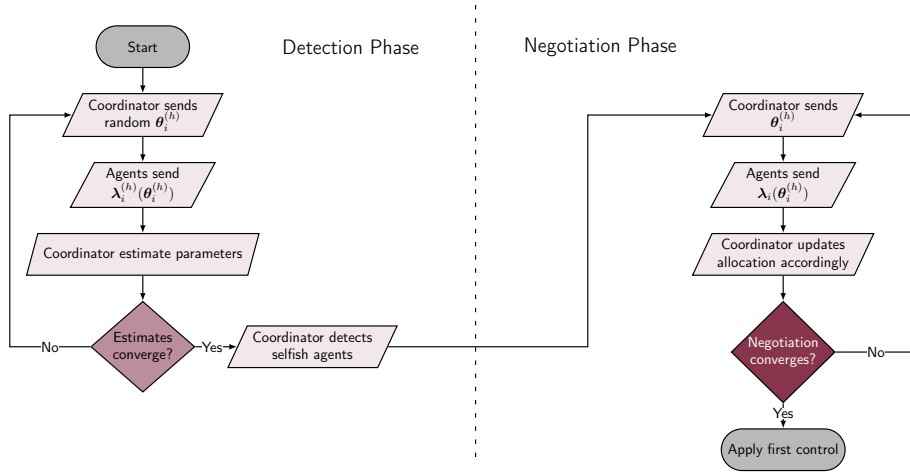
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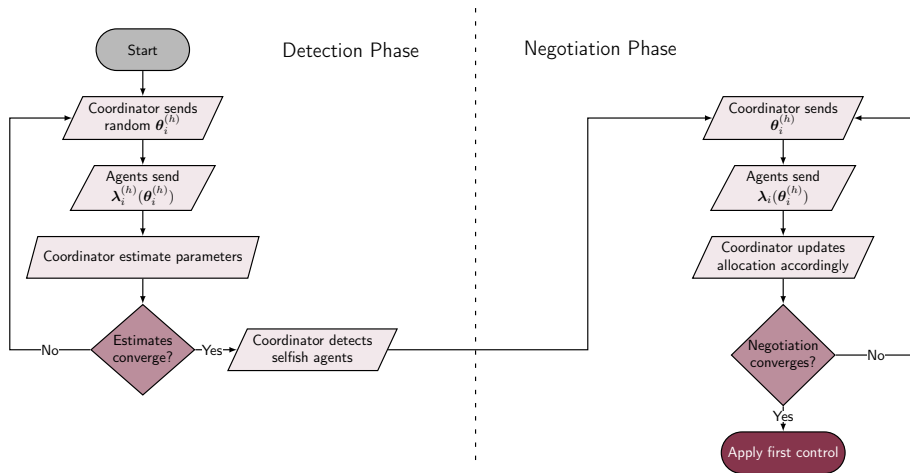
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Outline

- ① Vulnerabilities in distributed MPC based on Resource Allocation
 - Attacks
 - Consequences
- ② Securing the DMPC
 - Analysis of Subproblems
 - Detection Mechanism
 - Mitigation Mechanism
 - Complete Mechanism
- ③ Results



Example

Temperature Control of 4 Distinct Rooms Under Power Scarcity

- 4 distinct rooms modeled using 3R-2C
- Initial temperature under 20°C
- Not enough power to achieve setpoint $\left(\sum_{i=1}^4 u_i(k) \leq 4\text{kW}\right)$
- Simulated for a period of 5h
- ZOH $T_s = 0.25\text{h}$
- 3 scenarios
 - ① Nominal
 - ② Agent I non cooperative from $k>6$ with $T=4^{\circ}\text{I}$
 - ③ Similar but with secure algorithm



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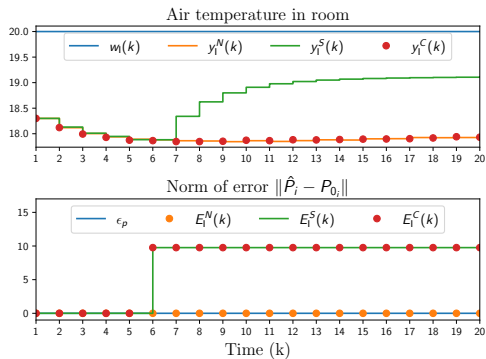
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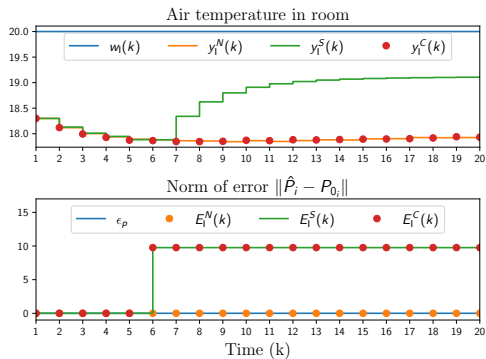
N Nominal

S Selfish behavior

C selfish behavior with Correction

Results

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Table 1: Comparison of costs J_i^N and J_G^N

Agent	Nominal	Selfish	Selfish + correction
I	103	64	104
II	73	91	73
III	100	123	101
IV	132	154	131
Global	408	442	409

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Summary

- ① Resource allocation based DMPC is vulnerable to attacks.
 - ② Sub-problems' structure has time invariant parameters.
 - ③ Attacks can be detected using these parameters.
 - ④ Effects can be mitigated.
- Outlook
 - Inequality Constraints yield Hybrid behavior
 - Non-linear attack model



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For Further Reading I

-  J. M. Maestre, R. R. Negenborn *et al.*
Distributed Model Predictive Control made easy.
Springer, 2014, vol. 69.
-  P. Velarde, J. M. Maestre, H. Ishii, and R. R. Negenborn,
“Scenario-based defense mechanism for distributed model predictive control,”
2017 IEEE 56th Annual Conference on Decision and Control (CDC). IEEE,
Dec 2017, pp. 6171–6176.

Thank you!

Repository

<https://github.com/Accacio/SysTol-21>



Contact

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