## Team Notebook

### December 17, 2018

Contents			13 Geometry	10	$26 \ { m Random}_G enerator$	19
1	2-SAT	2	14 HLD	11	<b>27</b> Sparse <sub>M</sub> $atrix_(RMQ)$	19
2	$2_{C} losest_{P} oints_{i} n_{2} D_{P} lane_{(} N_{l} og_{N}^{2})$	2	15 $\mathbf{Hashing}_{(Strings)}$	11	${\bf 28 \ Strongly}_{C} onnected_{C} omponent$	19
3	$\mathbf{Bridge}_T ree_i n_G raph$	3	16 Intervals $_H$ and $ling$	12	${\bf 29} \ {\bf Tree}_{C}onstruction_with_Specific_Vertices$	20
4	${\bf Centroid}_{D} ecomposition$	3	17 KMP	12	30 $\mathbf{Z}_{A} lgorithm$	20
5	$\mathbf{Convex}_{H}ull_{(}Dynamic)$	3	18 LineSweep-AreaofRectangles	13	31 blockcuttree	20
6	$\mathbf{Convex}_{H}ull_{(}Graham's_{S}can)$	5	19 LineSweepClosestPair	14	32 boruvkamst	21
7	$\mathbf{DFS}_{-C}ycle_{D}etection_{i}n_{D}irected_{G}raph$	5	<b>20</b> Matching $(Hopcroft-Karp)_i n_B ipartite_G rap$	ph	34 dpoptimization	21
8	$\mathbf{DSU}_O n_T rees$	6	21 Matrix <sub>S</sub> truct	15	34 euler	22
9	Dijkstra	6	22 $\mathbf{MaxFlow}_{-P}ush_{R}elabel[V^{3}]$	16	35 implicitsegtree	23
10	$\mathbf{Extended}_{E}uclidean_{A}lgorithm_{(}Extensive)$	7	$23 \ \mathrm{Min}_{C}ost_{M}ax_{F}low_{-D}ijkstra$	17	36 manacher	24
11	$\mathbf{FFT}_(Iterative)$	8	$24 \text{ Mo's}_{A} lgorithm$	18	37 maths	24
<b>12</b>	$\mathbf{Gaussian}_{E} limination$	10	<b>25</b> Ordered <sub>S</sub> $tatistic_Tree_(PBDS)$	18	38 persistent	25

39 suffixarray

26

### 1 2-SAT

```
struct TwoSAT
static const int MAXV=1e5+5;
int n, cnt;
vector<int> g[MAXV], rg[MAXV]; //g=forward, rg=
    backward
bool vis[MAXV];
int order[MAXV], comp[MAXV];
void init(int curn)
 n=curn:
 for(int i=0;i<n;i++)</pre>
  g[i].clear();
  rg[i].clear();
void add(int u, int v)
 g[u].push_back(v);
 rg[v].push_back(u);
 void dfs1(int u)
 vis[u] = true;
 for(auto it:g[u])
  if(!vis[it])
   dfs1(it);
 order[cnt++] = u;
void dfs2(int u, int c)
 comp[u] = c;
 for(auto it:rg[u])
  if(comp[it]==-1)
   dfs2(it, c);
}
int solve(vector<int> &ans)
```

```
{
cnt=0:
memset(vis, 0, sizeof(vis));
for(int i=0;i<n;i++)</pre>
 if(!vis[i])
  dfs1(i);
memset(comp, -1, sizeof(comp));
int grp=0;
for(int i=n-1;i>=0;i--)
 int u=order[i];
 if(comp[u] == -1)
  dfs2(u, grp++);
for(int i=0:i<n:i+=2)</pre>
 if(comp[i] == comp[i^1])
  return 0;
ans.clear();
for(int i=0;i<n;i+=2)</pre>
 int choose = (comp[i] > comp[i^1]) ? i : (i^1)
  ans.push_back(choose);
return 1;
```

### $egin{array}{ccc} oldsymbol{2} & oldsymbol{2}_{C} losest_{P}oints_{i}n_{2}D_{P}lane_{(}N_{l}og_{N}^{2}) \end{array}$

```
struct Point
{
  int x, y;
  Point operator -(Point p)
  {
   return {x-p.x, y-p.y};
  }
```

```
int dist()
 return x*x + y*y;
bool by_x(Point &a, Point &b)
return a.x < b.x;</pre>
bool by_y(Point &a, Point &b)
return a.y < b.y;</pre>
int n, ans=1e18;
int a[N], pref[N];
Point pt[N];
int solve(int L, int R)
{
if(L==R)
 return 1e18:
int M=(L+R)/2:
sort(pt+L, pt+R+1, by_x);
 int d=min(solve(L, M), solve(M+1, R));
 int midx=pt[L+(R-L+1)/2].x;
vector<Point> v;
for(int i=L;i<=R;i++)</pre>
 if(Point{pt[i].x-midx, 0}.dist()<d)</pre>
  v.push_back(pt[i]);
 sort(v.begin(), v.end(), by_y);
for(int i=0;i<v.size();i++)</pre>
 for(int j=i+1;j<v.size();j++)</pre>
  if(Point{0, v[i].y-v[j].y}.dist()>d)
   break;
  d=min(d, (v[i]-v[j]).dist());
```

```
}
return d;
}
```

### 3 $\mathbf{Bridge}_T ree_i n_G raph$

```
int tim=0, grp=1;
int u[N], v[N], comp[N];
bool vis[N], vis2[N], isBridge[M];
int tin[N], tout[N], minAncestor[N];
queue<int> Q[N];
vector<pair<int, int> > g[N];
vector<int> tree[N], vertices[N]; //Tree stores
   Bridge Tree, vertices stores the nodes in
   each component
void dfs(int k, int par)//bridges
vis[k]=1;
tin[k]=++tim;
minAncestor[k]=tin[k];
for(auto it:g[k])
 if(it.first==par)
  continue:
 if(vis[it.first])
  minAncestor[k]=min(minAncestor[k], tin[it.
      firstl):
  continue:
 dfs(it.first, k);
 minAncestor[k]=min(minAncestor[k], minAncestor[
     it.first]);
 if(minAncestor[it.first]>tin[k])
  isBridge[it.second]=1;
 tout[k]=tim;
```

```
void dfs2(int k)
int comp=grp;
Q[comp].push(k);
 vis2[k]=1;
while(!Q[comp].empty())
 int u=Q[comp].front();
 Q[comp].pop();
 vertices[comp].push_back(u);
 for(auto it:g[u])
  int v=it.first;
  int edgeidx=it.second;
  if(vis2[v])
   continue:
  if(isBridge[edgeidx])
   grp++;
   tree[comp].push_back(grp);
   tree[grp].push_back(comp);
   dfs2(v);
  else
   Q[comp].push(v);
   vis2[v]=1;
```

### $oxed{4}$ Centroid Decomposition

```
int subtree[N], parentcentroid[N];
set<int> g[N];
void dfs(int k, int par)
```

```
nodes++:
subtree[k]=1:
for(auto it:g[k])
 if(it==par)
  continue;
 dfs(it, k);
 subtree[k]+=subtree[it];
int centroid(int k, int par)
for(auto it:g[k])
 if(it==par)
  continue:
 if(subtree[it]>(nodes>>1))
  return centroid(it, k);
return k:
void decompose(int k, int par)
nodes=0;
dfs(k, k);
 int node=centroid(k, k);
 parentcentroid[node]=par;
for(auto it:g[node])
 g[it].erase(node);
 decompose(it, node);
```

### 5 $\mathbf{Convex}_Hull_(Dynamic)$

struct ConvexHullDynamic

```
{
static const int INF=1e18:
struct Line
 int a, b; //y = ax + b
 double xLeft; //Stores the intersection wiith
     previous line in the convex hull. First
     line has -INF
 enum Type {line, maxQuery, minQuery} type;
 int val;
 explicit Line(int aa=0, int bb=0): a(aa), b(bb)
     , xLeft(-INF), type(Type::line), val(0) {}
 int valueAt(int x) const
 {
  return a*x + b;
 friend bool isParallel(const Line &11. const
     Line &12)
  return 11.a == 12.a;
 friend double intersectX(const Line &11, const
     Line &12)
  return isParallel(11, 12)?INF:1.0*(12.b-11.b)
      /(l1.a-l2.a);
 bool operator<(const Line& 12) const
  if(12.type == line)
   return this->a < 12.a:
  if(12.type == maxQuery)
   return this->xLeft < 12.val;</pre>
  if(12.type == minQuery)
   return this->xLeft > 12.val;
 }
}:
```

```
bool isMax:
set<Line> hull:
bool hasPrev(set<Line>::iterator it)
return it!=hull.begin();
bool hasNext(set<Line>::iterator it)
return it!=hull.end() && next(it)!=hull.end();
bool irrelevant(const Line &11, const Line &12,
    const Line &13)
{
return intersectX(11, 13) <= intersectX(11, 12)</pre>
}
bool irrelevant(set<Line>::iterator it)
return hasPrev(it) && hasNext(it) && (
  (isMax && irrelevant(*prev(it), *it, *next(it)
     ))
  || (!isMax && irrelevant(*next(it), *it, *prev
      (it))));
//Updates xValue of line pointed by it
set<Line>::iterator updateLeftBorder(set<Line>::
    iterator it)
if(isMax && !hasPrev(it) || !isMax && !hasNext(
    it))
 return it;
double val=intersectX(*it, isMax?(*prev(it)):(*
    next(it)));
Line temp(*it);
 it=hull.erase(it):
temp.xLeft=val;
it=hull.insert(it, temp);
return it;
}
```

```
explicit ConvexHullDynamic(bool isMax): isMax(
   isMax) {}
void addLine(int a, int b) //Add ax + b in logN
Line 13=Line(a, b);
auto it=hull.lower_bound(13);
//If parallel liune is already in set, one of
    the lines becomes irrelevant
if(it!=hull.end() && isParallel(*it, 13))
 if(isMax && it->b<b || !isMax && it->b>b)
  it=hull.erase(it);
 else
  return:
it=hull.insert(it, 13);
if(irrelevant(it))
 hull.erase(it):
 return:
 //Remove lines which became irrelevant after
    inserting
 while(hasPrev(it) && irrelevant(prev(it)))
 hull.erase(prev(it));
 while(hasNext(it) && irrelevant(next(it)))
 hull.erase(next(it));
//Update xLine
it=updateLeftBorder(it);
if(hasPrev(it))
 updateLeftBorder(prev(it));
if(hasNext(it))
 updateLeftBorder(next(it));
int getBest(int x)
Line q;
```

```
q.val=x;
q.type = isMax?Line::Type::maxQuery : Line::
        Type::minQuery;
auto bestLine=hull.lower_bound(q);
if(isMax)
    --bestLine;
return bestLine->valueAt(x);
};
```

### 6 $Convex_Hull_(Graham's_Scan)$

```
struct point //Replace double with int if not
   required
double x, y;
point () {}
point(int x, int y) : x(x), y(y) {}
void operator =(const point &p)
 x=p.x, y=p.y;
 bool operator <(const point&p)</pre>
 if(x==p.x)
  return y<p.y;</pre>
 return x<p.x;</pre>
point operator +(const point&p) const
 point pt(x + p.x, y + p.y);
 return pt;
point operator -(const point&p) const
 point pt(x - p.x, y - p.y);
 return pt;
```

```
double crossProduct(const point &p) const
 return x * p.y - y * p.x;
 int dotProduct(const point &p) const
 return x * p.x + y * p.y;
double dist()
 return x*x + y*y;
bool comp(point &p1, point &p2)
if(p1.x!=p2.x)
 return p1.x<p2.x;</pre>
return p1.y<p2.y;</pre>
bool cw(point &a, point &b, point &c)
int area=a.x*(b.y-c.y) + b.x*(c.y-a.y) + c.x*(a.
    v-b.v);
return area<0;</pre>
bool ccw(point &a, point &b, point &c)
int area=a.x*(b.y-c.y) + b.x*(c.y-a.y) + c.x*(a.
    y-b.y);
return area>0:
vector<point> convex_hull(vector<point> &v)
if(v.size()==1)
 return v;
```

```
sort(v.begin(), v.end(), comp);
point p1=v[0], p2=v.back();
vector<point> up, down;
up.push_back(p1);
down.push_back(p1);
for(int i=1;i<v.size();i++)</pre>
 if(i==v.size()-1 || cw(p1, v[i], p2))
  while(up.size()>=2 && !cw(up[up.size()-2], up[
      up.size()-1], v[i]))
   up.pop_back();
 up.push_back(v[i]);
 if(i==v.size()-1 || ccw(p1, v[i], p2))
  while(down.size()>=2 && !ccw(down[down.size()
      -2], down[down.size()-1], v[i]))
   down.pop_back();
  down.push_back(v[i]);
for(int i=down.size()-2;i>0;i--)
 up.push_back(down[i]);
return up;
```

### $egin{array}{ccc} \mathbf{7} & \mathbf{DFS}_{-C}ycle_Detection_in_Directed_Ground 1. \end{array}$

```
bool findLoop(int v)
{
  if(vis[v]==1)
   return 1;
  if(vis[v]==2)
   return 0;
  vis[v]=1;
  for(auto &it:g[v])
  {
```

```
if(findLoop(it))
  return 1;
}
vis[v]=2;
return 0;
}
bool checkLoop()
{
  fill(vis+1, vis+n+1, 0);
  for(int i=1;i<=n;i++)
  {
    if(!vis[i] && findLoop(i))
     return 1;
}
return 0;
}</pre>
```

### 8 $\mathbf{DSU}_O n_T rees$

```
f[cnt[col[v]]]-=col[v]:
cnt[col[v]]+=x:
active.insert(cnt[col[v]]);
f[cnt[col[v]]]+=col[v];
for(auto u:g[v])
 if(u!=p && !big[u])
  add(u, v, x);
void computeans(int v)
int maxf=*(--active.end());
ans[v]=f[maxf];
void dfs(int v, int p, int keep)
int mx = -1, bigChild = -1;
for(auto u:g[v])
 if(u!=p && subtree[u]>mx)
  mx=subtree[u], bigChild=u;
for(auto u:g[v])
 if(u!=p && u!=bigChild)
  dfs(u, v, 0); //Run DFS on small children and
      clear them
if(bigChild!=-1)
 dfs(bigChild, v, 1);
 big[bigChild]=1;
 }
 add(v, p, 1);
//Now we have the information of subtree of v
 computeans(v);
if(bigChild!=-1)
 big[bigChild]=0;
if(keep==0)
```

```
add(v, p, -1);
}
```

### 9 Dijkstra

```
Dijkstra with Path:
int arrival[N], departure[N], vis[N], parent[N];
vector<pair<int, int> > g[N];
void dijkstra(int source, int destination)
for(int i=1;i<=n;i++)</pre>
 arrival[i]=1e18:
 departure[i]=1e18;
 vis[i]=0:
 arrival[source]=0;
set<pair<int, int> > s;
 s.insert({0, source});
 while(!s.empty())
 auto x = *(s.begin());
 s.erase(x);
 vis[x.second]=1;
 departure[x.second] = arrival[x.second];
 for(auto it:g[x.second])
  if(arrival[it.first] > departure[x.second] +
      it.second)
   s.erase({arrival[it.first], it.first});
   arrival[it.first] = departure[x.second] + it.
   s.insert({arrival[it.first], it.first});
    parent[it.first] = x.second;
```

```
}
}
if(!vis[destination])
{
   cout<<"-1";
   return;
}
int v=destination;
vector<int> ans;
while(parent[v])
{
   ans.push_back(v);
   v=parent[v];
}
ans.push_back(source);
reverse(ans.begin(), ans.end());
for(auto it:ans)
   cout<<it<<" ";
}</pre>
```

# 10 Extended<sub>E</sub> $uclidean_Algorithm_{\ell}(Extensive)$

```
int modular inverse(int a, int m)
int x, y;
int g=xgcd(a, m, x, y);
if(g!=1)
 return -1;
 else
 x=(x\%m + m)\%m:
 return x;
void shift_solution(int &x, int &y, int a, int b,
    int cnt)
x+=cnt*b:
y-=cnt*a;
bool find_any_solution(int a, int b, int c, int &
 int g=xgcd(abs(a), abs(b), x0, y0);
if(c%g!=0)
 return false;
x0 *= c/g;
 v0 *= c/g;
 if(a<0)
 x0*=-1:
 if(b<0)
 y0*=-1;
return true;
int find_all_solutions(int a, int b, int c, int
   minx, int maxx, int miny, int maxy) //Returns
    number of solutions with x[minx, maxx], y[
   miny, maxy]
```

```
int x, y, g;
if(!find_any_solution(a, b, c, x, y, g))
return 0:
a /= g;
b /= g;
int sign_a = a>0 ? +1 : -1;
int sign_b = b>0 ? +1 : -1;
shift_solution(x, y, a, b, (minx - x) / b);
if (x < minx) shift_solution(x, y, a, b, sign_b)</pre>
if (x > maxx) return 0;
int lx1 = x;
shift_solution(x, y, a, b, (maxx - x) / b);
if (x > maxx) shift_solution(x, y, a, b, -sign_b
    );
int rx1 = x:
shift_solution(x, y, a, b, - (miny - y) / a);
if (y < miny) shift_solution(x, y, a, b, -sign_a</pre>
if (y > maxy) return 0;
int 1x2 = x:
shift_solution(x, y, a, b, - (maxy - y) / a);
if (y > maxy) shift_solution(x, y, a, b, sign_a)
int rx2 = x;
if (1x2 > rx2)
 swap (1x2, rx2);
int lx = max (lx1, lx2);
int rx = min(rx1, rx2):
return (rx - lx) / abs(b) + 1;
```

### 11 FFT<sub>(</sub>Iterative)

```
4 Call FFT with MOD:
typedef complex<double> base;
const double PI = acos(-1.01);
const int N = 8e5+5;
const int Maxb = 19:
const int Maxp = 450;
const int MOD=13313;
vector<int> rev;
vector<base> omega;
void calc_rev(int n, int log_n) //Call this
    before FFT
 omega.assign(n, 0);
rev.assign(n, 0);
for(int i=0;i<n;i++)</pre>
 rev[i]=0;
 for(int j=0; j<log_n; j++)</pre>
  if((i>>j)&1)
   rev[i] |= 1<<(log_n-j-1);
}
}
void fft(vector<base> &A, int n, bool invert)
for(int i=0;i<n;i++)</pre>
 if(i<rev[i])</pre>
  swap(A[i], A[rev[i]]);
```

```
for(int len=2:len<=n:len<<=1)</pre>
 double ang=2*PI/len * (invert?-1:+1);
 int half=(len>>1);
 base curomega(cos(ang), sin(ang));
 omega[0]=base(1, 0);
 for(int i=1;i<half;i++)</pre>
  omega[i]=omega[i-1]*curomega;
 for(int i=0;i<n;i+=len)</pre>
  base t:
  int pu = i,
   pv = i+half,
   pu_end = i+half,
   pw = 0;
  for(; pu!=pu_end; pu++, pv++, pw++)
   t=A[pv] * omega[pw];
   A[pv] = A[pu] - t;
   A[pu] += t;
if(invert)
 for(int i=0;i<n;i++)</pre>
  A[i]/=n;
void multiply(int n, vector<base> &A, vector<base</pre>
    > &B, vector<int> &C)
fft(A, n, false);
fft(B, n, false);
for(int i=0;i<n;i++)</pre>
 A[i] *= B[i];
```

```
fft(A, n, true);
for(int i=0:i<n:i++)</pre>
 C[i] = (int)(A[i].real() + 0.5);
 C[i] %= MOD;
void Solve(int n, vector<int> &coeffA, vector<int</pre>
   > &coeffB, vector<int> &result, bool big1,
   bool big2) //Call 4 times: 00, 01, 10, 11
vector<base> A(n), B(n);
for(int i=0;i<n;i++)</pre>
 A[i]=big1?coeffA[i]/Maxp : coeffA[i]%Maxp;
 B[i]=0:
for(int i=0;i<n;i++)</pre>
 B[i]=big2?coeffB[i]/Maxp : coeffB[i]%Maxp;
vector<int> C(n):
multiply(n, A, B, C);
for(int i=0;i<n;i++)</pre>
 int add=C[i];
 if(big1)
  add*=Maxp;
 if(big2)
  add*=Maxp;
 add%=MOD;
 result[i]+=add;
 result[i]%=MOD:
}
void do_FFT(vector<int> &A, vector<int> &B,
   vector<int> &result)
```

```
rev[i]=0:
                                                                                                               }
int n=1, bits=0;
while(n<2*A.size() || n<2*B.size())</pre>
                                                        for(int j=0; j<log_n; j++)</pre>
 n<<=1, bits++;
result.assign(n, 0);
                                                         if((i>>j)&1)
                                                                                                              if(invert)
 calc_rev(n, bits);
                                                          rev[i] |= 1<<(log_n-j-1);
                                                                                                               for(int i=0;i<n;i++)</pre>
vector<int> tempA(A.begin(), A.end());
                                                                                                                A[i]/=n;
vector<int> tempB(B.begin(), B.end());
tempA.resize(n);
tempB.resize(n);
                                                                                                             void multiply(int n, vector<base> &A, vector<base</pre>
 for(int i=0;i<2;i++)</pre>
                                                      void fft(vector<base> &A, int n, bool invert)
                                                                                                                 > &B, vector<int> &C)
                                                       for(int i=0;i<n;i++)</pre>
                                                                                                              fft(A, n, false);
 for(int j=0;j<2;j++)</pre>
                                                                                                              fft(B, n, false);
                                                        if(i<rev[i])</pre>
                                                                                                              for(int i=0;i<n;i++)</pre>
  Solve(n, tempA, tempB, result, i, j);
                                                         swap(A[i], A[rev[i]]);
                                                                                                              A[i] *= B[i]:
}
                                                                                                              fft(A, n, true);
                                                       for(int len=2;len<=n;len<<=1)</pre>
                                                                                                              for(int i=0:i<n:i++)</pre>
                                                        <del>double-ang=2*PI/len-*-(invert?-1:+1)</del>-;------
                                                                                                             --G[i]-=-(int)(A[i]-real()-+-0.5);
                                                        int half=(len>>1);
Single Call without MOD:
                                                        base curomega(cos(ang), sin(ang));
typedef complex<double> base;
                                                        omega[0]=base(1, 0);
                                                                                                             void Solve(int n, vector<int> &coeffA, vector<int</pre>
const double PI = acos(-1.01);
                                                        for(int i=1;i<half;i++)</pre>
                                                                                                                 > &coeffB, vector<int> &result)
const int N = 8e5+5;
                                                         omega[i]=omega[i-1]*curomega;
const int Maxb = 19;
                                                                                                              vector<base> A(n), B(n);
const int Maxp = 450;
                                                        for(int i=0;i<n;i+=len)</pre>
                                                                                                              for(int i=0;i<n;i++)</pre>
const int MOD=13313;
                                                        ₹
                                                                                                               A[i]=coeffA[i];
                                                         base t;
                                                                                                               B[i]=0;
vector<int> rev;
                                                         int pu = i,
vector<base> omega;
                                                          pv = i+half,
                                                          pu_end = i+half,
                                                                                                              for(int i=0;i<n;i++)</pre>
void calc_rev(int n, int log_n) //Call this
                                                          : 0 = wq
    before FFT
                                                         for(; pu!=pu_end; pu++, pv++, pw++)
                                                                                                               B[i]=coeffB[i]:
omega.assign(n, 0);
                                                          t=A[pv] * omega[pw];
                                                                                                              vector<int> C(n);
rev.assign(n, 0);
                                                          A[pv] = A[pu] - t;
                                                                                                              multiply(n, A, B, C);
for(int i=0;i<n;i++)</pre>
                                                          A[pu] += t;
                                                                                                              for(int i=0;i<n;i++)</pre>
```

### 12 $Gaussian_E limination$

```
//Logic: https://math.stackexchange.com/questions
    /48682/maximization-with-xor-operator
struct Gaussian
{
    int no_of_bits = 20;
    vector<int> v;
    int set, origsize=0, redsize=0;

void push(int val)
    {
        origsize++;
        if(val)
        v.push_back(val);
    }

void clear()
```

```
v.clear():
set=0, redsize=0;
void eliminate()
set = redsize = 0;
for(int bit=0;bit<=no_of_bits;bit++)</pre>
 bool check=false;
  for(int i=redsize;i<v.size();i++)</pre>
  if((v[i]>>bit)&1)
   swap(v[i], v[redsize]);
   check=true:
   break;
  if (check)
  for(int i=redsize+1;i<v.size();i++)</pre>
   if((v[i]>>bit)&1)
    v[i]^=v[redsize];
  redsize++;
v.resize(redsize);
for(auto it:v)
  set |=it:
Gaussian& operator =(Gaussian &orig)
v = orig.v;
set = orig.set;
redsize = orig.redsize;
```

```
origsize = orig.origsize;
return *this;
}
};
```

### 13 Geometry

```
struct point
int x, y, idx;
};
//Finds squared euclidean distance between two
   points
int dist(point &a, point &b)
return (a.x-b.x)*(a.x-b.x) + (a.y-b.y)*(a.y-b.y)
//Checks if angle ABC is a right angle
int isOrthogonal(point &a, point &b, point &c)
return (b.x-a.x) * (b.x-c.x) + (b.y-a.y) * (b.y-a.y)
    c.v) == 0;
//Checks if ABCD form a rectangle (in that order)
int isRectangle(point &a, point &b, point &c,
   point &d)
return isOrthogonal(a, b, c) && isOrthogonal(b,
    c, d) && isOrthogonal(c, d, a);
//Checks if ABCD form a rectangle, in any
   orientation
```

```
int isRectangleAnyOrder(point &a, point &b, point
    &c, point &d)
return isRectangle(a, b, c, d) || isRectangle(b,
     c, a, d) | isRectangle(c, a, b, d);
//Checks if ABCD form a square (in that order)
int isSquare(point &a, point &b, point &c, point
   &d)
return isRectangle(a, b, c, d) && dist(a, b) ==
    dist(b, c);
}
//Checks if ABCD form a square, in any
   orientation
int isSquareAnyOrder(point &a, point &b, point &c
    , point &d)
return isSquare(a, b, c, d) || isSquare(b, c, a,
     d) | isSquare(c, a, b, d);
```

### 14 HLD

```
void init() {
    chainNo = 1;
    pos = 0;
    for (int i = 1; i <= n; i++) {
        chainHead[i] = -1;
        g[i].clear();
    }
}
void HLD(int k, int cost, int par) {
    if (chainHead[chainNo] == -1)
        chainHead[chainNo] = k;
    pos++;</pre>
```

```
chainInd[k] = chainNo:
   basePos[k] = pos;
   a[pos] = cost:
   int special = -1, maxval = -1;
   for (auto it: g[k]) {
       if (it.ff == par)
           continue;
       if (subtree[it.ff] > maxval) {
           maxval = it.ss:
           special = it.ff;
   if (special != -1)
       HLD(special, maxval, k);
   for (auto it: g[k]) {
       if (it.ff == par || it.ff == special)
           ICPC Notebook
       continue:
       chainNo++;
       HLD(it.ff, it.ss, k);
int queryUp(int u, int v) {
   int left = chainInd[u];
   int right = chainInd[v];
   int ans = 0;
   while (1) {
       if (u == v)
           break:
       if (left == right) {
           ans = max(ans, query(1, 1, pos,
              basePos[v] + 1, basePos[u]));
           break;
       ans = max(ans, query(1, 1, pos, basePos[u
           ], basePos[chainHead[left]]));
       u = chainHead[left];
       u = parent[0][u];
       left = chainInd[u];
```

```
return ans;
}
int querypath(int u, int v) {
   int lca = LCA(u, v);
   int ans1 = queryUp(u, lca);
   int ans2 = queryUp(v, lca);
   return max(ans1, ans2);
}
void rearrange()
for (int i = 1; i <= n - 1; i++)
   if (level[edge[i].ss] < level[edge[i].ff])
      swap(edge[i].ff, edge[i].ss);
void change(int pos1, int val) {
   int node = edge[pos1].ss;
   update(1, 1, pos, basePos[node], val);
}</pre>
```

### $|15 \quad \mathbf{Hashing}(Strings)|$

```
struct Hashs
vector<int> hashs;
vector<int> pows;
int P;
int MOD;
Hashs() {}
Hashs(string &s, int P, int MOD) : P(P), MOD(MOD
 {
 int n = s.size();
 pows.resize(n+1, 0):
 hashs.resize(n+1, 0):
 pows[0] = 1:
 for(int i=n-1;i>=0;i--)
  hashs[i]=(1LL * hashs[i+1] * P + s[i] - 'a' +
      1) % MOD:
  pows[n-i]=(1LL * pows[n-i-1] * P) % MOD;
```

### 16 Intervals $_H$ and ling

```
map<int, int> active;

void init()
{
   active[-1] = -1;
   active[2e9] = 2e9;
   active[1] = n;
}

void add(int L, int R) //Always remove [L, R]
   before adding
{
   active[L]=R;
   ans+=R-L+1;
}

void remove(int L, int R)
{
   int removed=0;
   auto it = active.lower_bound(L);
   it--;
   if(it->second>=L)
{
    active[L] = it->second;
}
```

```
it->second = L-1:
it++:
while(it->first <= R)</pre>
 if(it->second > R)
  removed+=R + 1 - it->first;
  active[R+1] = it->second;
  removed+= it->second - it->first + 1;
 auto it2=it;
 it++;
 active.erase(it2):
ans-=removed;
//submission
for(int l=1:1<=n:1++)</pre>
 int r=1:
 while (r+1 \le n \&\& a[r+1] = = a[1])
 s.insert({1, r-l+1});
 rem.insert({-(r-l+1), l});
 l=r;
while(rem.size())
 ans++;
 auto it=*rem.begin();
 int idx=it.second;
 rem.erase(it):
 auto it2=s.lower_bound(make_pair(idx, 0));
 auto L=it2, R=it2;
 if(L!=s.begin() && R!=s.end())
 {
  L--:
  R++:
```

### 17 KMP

```
String:
vector<int> prefix_function(string &s)
{
  int n = (int)s.length();
  vector<int> pi(n);
  for (int i = 1; i < n; i++)
  {
    int j = pi[i-1];
    while (j > 0 && s[i] != s[j])
        j = pi[j-1];
    if (s[i] == s[j])
        j++;
    pi[i] = j;
  }
  return pi;
}
vector<int> find_occurences(string &text, string &pattern)
```

```
string cur=pattern + '#' + text;
int sz1=text.size(), sz2=pattern.size();
vector<int> v;
vector<int> lps=prefix_function(cur);
for(int i=sz2+1;i<=sz1+sz2;i++)</pre>
 if(lps[i]==sz2)
  v.push_back(i-2*sz2);
return v;
Vector:
vector<int> prefix_function(vector<int> &v)
int n = (int)v.size();
vector<int> pi(n);
for (int i = 1; i < n; i++)</pre>
 int j = pi[i-1];
 while (j > 0 \&\& v[i] != v[j])
  j = pi[j-1];
 if (v[i] == v[j])
  j++;
 pi[i] = j;
return pi;
vector<int> find_occurences(vector<int> &text,
    vector<int> &pattern)
vector<int> v=pattern;
v.push_back(-1);
for(auto &it:text)
 v.push_back(it);
int sz1=text.size(), sz2=pattern.size();
vector<int> lps=prefix_function(v);
vector<int> store;
for(int i=sz2+1;i<=sz1+sz2;i++)</pre>
```

```
if(lps[i]==sz2)
 store.push_back(i-sz*2);
return v;
```

```
inline bool intersect1d (double 11, double r1,
    double 12, double r2 ) {
if ( l1 > r1 ) swap ( l1, r1 );
if (12 > r2) swap (12, r2);
return max ( 11, 12 ) <= min ( r1, r2 ) + EPS ;</pre>
inline int vec ( const pt & a, const pt & b,
    const pt & c ) {
double s = (b. x - a. x) * (c. y - a. y) - (
   b. y - a. y ) * (c. x - a. x);
return abs (s) < EPS ? 0 : s > 0 ? + 1 : - 1;}
bool intersect ( const seg & a, const seg & b ) {
return intersect1d ( a. p . x , a. q . x , b. p
     x , b. q . x )
&& intersect1d ( a. p . y , a. q . y , b. p . y ,
&& vec ( a. p , a. q , b. p ) * vec ( a. p , a. q
    , b. q = 0
&& vec ( b. p , b. q , a. p ) * vec ( b. p , b. q
     . a. a ) <= 0 :}
bool operator < ( const seg & a, const seg & b )</pre>
double x = max (min (a. p. x, a. q. x),
   min (b.p.x,b.q.x));
return a. get_y ( x ) < b. get_y ( x ) - EPS ;}</pre>
#define MAX 1000
struct event
int ind; // Index of rectangle in rects
```

```
bool type; // Type of event: 0 = Lower-left ; 1 =
                                                Upper-right
                                           event() {}:
                                           event(int ind, int type) : ind(ind), type(type)
                                           point rects [MAX][12]; // Each rectangle consists
                                                of 2 points: [0] = lower-left; [1] = upper-
                                               right
LineSweep-AreaofRectangles compare_x(event a, event b) { return rects[a
                                               .ind][a.type].x<rects[b.ind][b.type].x; }
                                           bool compare_v(event a, event b) { return rects[a
                                               .ind][a.type].y<rects[b.ind][b.type].y; }
                                           int union_area(event events_v[], event events_h[],
                                               int n.int e)
                                           //n is the number of rectangles, e=2*n , e is the
                                                number of points (each rectangle has two
                                               points as described in
                                           declaration of rects)
                                           bool in set[MAX]={0}:int area=0:
                                           sort(events_v, events_v+e, compare_x); //Pre-sort
                                                of vertical edges
                                           sort(events_h, events_h+e, compare_y); // Pre-
                                               sort set of horizontal edges
                                           in_set[events_v[0].ind] = 1;
                                           for (int i=1;i<e;++i)</pre>
                                           { // Vertical sweep line
                                           event c = events_v[i];
                                           int cnt = 0; // Counter to indicate how many
                                               rectangles are currently overlapping
                                           // Delta_x: Distance between current sweep line
                                               and previous sweep line
                                           int delta_x = rects[c.ind][c.type].x - rects[
                                               events_v[i-1].ind][events_v[i-1].type].x;
                                           int begin_y;
                                           if (delta x==0){
                                           in_set[c.ind] = (c.type==0);
                                           continue;
```

```
for (int j=0;j<e;++j)</pre>
if (in_set[events_h[j].ind]==1)
//Horizontal sweep line for active rectangle
if (events_h[j].type==0)
//If it is a bottom edge of rectangle
if (cnt==0) begin_y = rects[events_h[j].ind][0].y
    ; // Block starts
//incrementing number of overlapping rectangles
else
//If it is a top edge
--cnt:
//the rectangle is no more overlapping, so remove
     it.
if (cnt==0)
//Block endsICPC Notebook
int delta_y = (rects[events_h[j].ind][13].y-
   begin_y);//length of the vertical sweep line
   cut by
rectangles
area+=delta_x * delta_v;
in_set[c.ind] = (c.type==0);//If it is a left
   edge, the rectangle is in the active set else
     not
}
return area;
```

### 19 LineSweepClosestPair

```
#define px second
#define py first
typedef pair<long long, long long> pairll;
pairll pnts [MAX];
int compare(pairll a, pairll b)
return a.px<b.px;</pre>
double closest_pair(pairll pnts[],int n)
sort(pnts,pnts+n,compare);
double best=INF;
set<pairll> box;
box.insert(pnts[0]);
int left = 0;
for (int i=1;i<n;++i)</pre>
while (left<i && pnts[i].px-pnts[left].px > best)
box.erase(pnts[left++]);
for(typeof(box.begin()) it=box.lower_bound(
    make_pair(pnts[i].py-best, pnts[i].px-best));
    it!=box.end() &&
pnts[i].py+best>=it->py;it++)
best = min(best, sqrt(pow(pnts[i].py - it->py,
    2.0)+pow(pnts[i].px - it->px, 2.0)));
box.insert(pnts[i]);
return best;
```

# $egin{array}{ccc} \mathbf{Matching}_(Hopcroft \ Karp)_i n_B i partite_G raph \end{array}$

```
//1 indexed Hopcroft-Karp Matching in O(E sqrtV)
```

```
static const int inf = 1e9:
int n;
vector<int> matchL, matchR, dist;
vector<vector<int> > g;
Hopcroft_Karp(int n) :
n(n), matchL(n+1), matchR(n+1), dist(n+1), g(n
    +1) {}
void addEdge(int u, int v)
g[u].push_back(v);
bool bfs()
 queue<int> q;
 for(int u=1;u<=n;u++)</pre>
 if(!matchL[u])
   dist[u]=0;
  q.push(u);
 else
  dist[u]=inf;
 dist[0]=inf;
 while(!q.empty())
 int u=q.front();
 q.pop();
 for(auto v:g[u])
  if(dist[matchR[v]] == inf)
    dist[matchR[v]] = dist[u] + 1;
```

```
q.push(matchR[v]);
return (dist[0]!=inf);
bool dfs(int u)
if(!u)
 return true;
for(auto v:g[u])
 if(dist[matchR[v]] == dist[u]+1 &&dfs(matchR[v])
     1))
  matchL[u]=v;
  matchR[v]=u;
  return true;
 }
dist[u]=inf;
return false;
int max_matching()
int matching=0;
while(bfs())
 for(int u=1;u<=n;u++)</pre>
  if(!matchL[u])
   if(dfs(u))
    matching++;
return matching;
```

```
};
```

### 21 $Matrix_S truct$

```
int add(int a, int b)
int res = a + b;
if(res >= MOD)
 return res - MOD;
return res;
int mult(int a, int b)
long long res = a;
res *= b:
if(res >= MOD)
 return res % MOD;
return res;
struct matrix
int arr[SZ][SZ];
void reset()
 memset(arr, 0, sizeof(arr));
void makeiden()
 reset();
 for(int i=0;i<SZ;i++)</pre>
  arr[i][i] = 1;
```

```
matrix operator + (const matrix &o) const
 matrix res;
 for(int i=0;i<SZ;i++)</pre>
  for(int j=0;j<SZ;j++)</pre>
   res.arr[i][j] = add(arr[i][j], o.arr[i][j]);
 return res;
matrix operator * (const matrix &o) const
 matrix res:
 for(int i=0;i<SZ;i++)</pre>
  for(int j=0;j<SZ;j++)</pre>
   res.arr[i][j] = 0;
   for(int k=0;k<SZ;k++)</pre>
    res.arr[i][j] = add(res.arr[i][j] , mult(arr
         [i][k] , o.arr[k][j]));
  }
 return res;
};
matrix power(matrix a, int b)
matrix res;
res.makeiden();
while(b)
 if(b & 1)
```

```
{
  res = res * a;
}
  a = a * a;
b >>= 1;
}
return res;
}
```

### 22 $\mathbf{MaxFlow}_{-P}ush_Relabel[V^3]$

```
//Push-Relabel Algorithm for Flows - Gap
   Heuristic, Complexity: O(V^3)
//To obtain the actual flow values, look at all
   edges with capacity > 0
//Zero capacity edges are residual edges
struct edge
int from, to, cap, flow, index;
 edge(int from, int to, int cap, int flow, int
    index):
 from(from), to(to), cap(cap), flow(flow), index
     (index) {}
};
struct PushRelabel
int n:
vector<vector<edge> > g;
vector<long long> excess;
vector<int> height, active, count;
 queue<int> Q;
 PushRelabel(int n):
 n(n), g(n), excess(n), height(n), active(n),
     count(2*n) {}
```

```
void addEdge(int from, int to, int cap)
g[from].push_back(edge(from, to, cap, 0, g[to].
    size())):
 if(from==to)
 g[from].back().index++;
g[to].push_back(edge(to, from, 0, 0, g[from].
    size()-1));
void enqueue(int v)
 if(!active[v] && excess[v] > 0)
  active[v]=true:
  Q.push(v);
}
void push(edge &e)
int amt=(int)min(excess[e.from], (long long)e.
     cap - e.flow);
 if(height[e.from] <= height[e.to] || amt == 0)</pre>
 return;
 e.flow += amt;
g[e.to][e.index].flow -= amt;
 excess[e.to] += amt;
 excess[e.from] -= amt;
enqueue(e.to);
void relabel(int v)
count[height[v]]--;
int d=2*n:
for(auto &it:g[v])
 if(it.cap-it.flow>0)
  d=min(d, height[it.to]+1);
```

```
height[v]=d;
 count[height[v]]++;
enqueue(v);
void gap(int k)
for(int v=0;v<n;v++)</pre>
 if (height[v] < k)</pre>
  continue;
 count[height[v]]--;
 height[v] = max(height[v], n+1);
 count[height[v]]++;
 enqueue(v);
void discharge(int v)
for(int i=0; excess[v]>0 && i<g[v].size(); i++)</pre>
 push(g[v][i]);
if(excess[v]>0)
 if(count[height[v]]==1)
  gap(height[v]);
 else
  relabel(v);
long long max_flow(int source, int dest)
count[0] = n-1:
count[n] = 1;
height[source] = n;
active[source] = active[dest] = 1;
for(auto &it:g[source])
```

```
excess[source]+=it.cap;
push(it);
}

while(!Q.empty())
{
  int v=Q.front();
  Q.pop();
  active[v]=false;
  discharge(v);
}

long long max_flow=0;
for(auto &e:g[source])
  max_flow+=e.flow;

return max_flow;
}
};
```

### 23 $\mathbf{Min}_{C}ost_{M}ax_{F}low_{-D}ijkstra$

```
//Works for negative costs, but does not work for
    negative cycles
//Complexity: O(min(E^2 *V log V, E logV * flow))
struct edge
{
    int to, flow, cap, cost, rev;
};
struct MinCostMaxFlow
{
    int nodes;
    vector<int> prio, curflow, prevedge, prevnode, q
        , pot;
    vector<vector<edge> > graph;
```

```
MinCostMaxFlow() {}
MinCostMaxFlow(int n): nodes(n), prio(n, 0),
    curflow(n, 0),
prevedge(n, 0), prevnode(n, 0), q(n, 0), pot(n,
    0), inqueue(n, 0), graph(n) {}
void addEdge(int source, int to, int capacity,
    int cost)
 edge a = {to, 0, capacity, cost, (int)graph[to
    ].size()};
edge b = {source, 0, 0, -cost, (int)graph[
    source].size()};
 graph[source].push_back(a);
graph[to].push_back(b);
void bellman_ford(int source, vector<int> &dist)
 fill(dist.begin(), dist.end(), INT_MAX);
 dist[source] = 0:
 int qt=0;
 q[qt++] = source;
 for(int qh=0;(qh-qt)%nodes!=0;qh++)
  int u = q[qh%nodes];
  inqueue[u] = false;
  for(auto &e : graph[u])
  if(e.flow >= e.cap)
   continue;
  int v = e.to;
  int newDist = dist[u] + e.cost;
  if(dist[v] > newDist)
   dist[v] = newDist;
   if(!inqueue[v])
    inqueue[v] = true;
```

```
a[at++ \% nodes] = v:
pair<int, int> minCostFlow(int source, int dest,
    int maxflow)
bellman_ford(source, pot);
int flow = 0;
int flow_cost = 0;
while(flow < maxflow)</pre>
 priority_queue<pair<int, int>, vector<pair<int</pre>
     , int> >, greater<pair<int, int> > q;
 q.push({0, source});
 fill(prio.begin(), prio.end(), INT_MAX);
 prio[source] = 0:
 curflow[source] = INT MAX:
 while(!q.empty())
  int d = q.top().first;
  int u = q.top().second;
  q.pop();
  if(d != prio[u])
   continue;
  for(int i=0;i<graph[u].size();i++)</pre>
   edge &e=graph[u][i];
   int v = e.to;
   if(e.flow >= e.cap)
    continue:
   int newPrio = prio[u] + e.cost + pot[u] -
       pot[v]:
   if(prio[v] > newPrio)
    prio[v] = newPrio;
    q.push({newPrio, v});
```

```
prevnode[v] = u;
     prevedge[v] = i;
     curflow[v] = min(curflow[u], e.cap - e.flow
  if(prio[dest] == INT_MAX)
   break:
  for(int i=0;i<nodes;i++)</pre>
   pot[i]+=prio[i];
  int df = min(curflow[dest], maxflow - flow);
  flow += df:
  for(int v=dest;v!=source;v=prevnode[v])
   edge &e = graph[prevnode[v]][prevedge[v]];
   e.flow += df:
   graph[v][e.rev].flow -= df;
   flow_cost += df * e.cost;
 return {flow, flow_cost};
};
```

### 24 Mo's<sub>A</sub>lgorithm

```
const int N=2e5+5;
const int M=1e6+5;
struct data
{
  int 1;
  int r;
  int idx;
  long long store_ans;
};
int n, q, blocksz=1000;
int a[N];
```

```
data queries[N];
long long freq[M];
long long ans=0;
bool comp(data &d1, data &d2)
 int blocka=d1.1/blocksz;
 int blockb=d2.1/blocksz;
 if(blocka<blockb)</pre>
  return true:
 else if(blocka==blockb)
  return (d1.r<d2.r)^(blocka%2);</pre>
  return false;
bool comp2(data &d1, data &d2)
 return d1.idx<d2.idx:</pre>
void update(long long k, int sign) //Sign 1 = Add
    -1 = Remove
 if(sign==1)
  ans-=freq[k]*freq[k]*k;
  freq[k]++;
  ans+=freq[k]*freq[k]*k;
 else
  ans-=freq[k]*freq[k]*k;
  freg[k]--;
  ans+=freq[k]*freq[k]*k;
void calcmo()
 int moleft=1;
 int moright=0;
 for(int i=1;i<=q;i++)</pre>
```

```
int r=queries[i].r;
int l=queries[i].l;
while(moright<r)
{
  moright++;
  update(a[moright], 1);
}
while(moright>r)
{
  update(a[moright], -1);
  moright--;
}
while(moleft<1)
{
  update(a[moleft], -1);
  moleft++;
}
while(moleft>l)
{
  update(a[moleft], 1);
}
queries[i].store_ans=ans;
}
```

### **25** Ordered<sub>S</sub> $tatistic_Tree_(PBDS)$

```
#include <bits/stdc++.h>
using namespace std;
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
#define IOS ios::sync_with_stdio(0); cin.tie(0);
    cout.tie(0);
#define endl "\n"
#define int long long
const int N=2e5+5;
```

```
#define T pair<int, int>
#define ordered_set tree<T, null_type, less<T>,
    rb_tree_tag,
    tree_order_statistics_node_update>
int getless(ordered_set &os, int R, int index)
{
    return os.order_of_key({R, index});
}
ordered_set os1;
//there are two new features it is find_by_order
    () and order_of_key().The first returns an
    iterator to the k-th
//largest element (counting from zero),the second
    the number of items in a set that are
    strictly smaller than our item
```

### **26** Random $_Generator$

```
mt19937 rng(chrono::steady_clock::now().
    time_since_epoch().count());
int getRand(int 1, int r)
{
    uniform_int_distribution<int> uid(1, r);
    return uid(rng);
}
```

### **27** Sparse $_M atrix_(RMQ)$

```
int RMQ[20][N];
void precompute()
{
   for(int i=0;(1<<i)<N;i++)
   for(int j=(1<<i);j<N && j<(1<<(i+1)); j++)
    floorlog[j]=i;
}</pre>
```

```
for(int i=n;i>=1;i--)
{
   RMQ[0][i]=dp[i];
   int mxj=floorlog[n-i+1]; //2^j <= n-i+1
   int pw=1;
   for(int j=1;j<=mxj;j++)
   {
     RMQ[j][i]=max(RMQ[j-1][i], RMQ[j-1][i+pw]);
     pw<<=1;
   }
}
int getMax(int L, int R)
{
   int k=floorlog[R-L+1]; //2^k <= R-L+1
   return max(RMQ[k][L], RMQ[k][R - (1<<k) +1]);
}</pre>
```

### 28 Strongly $Connected_{C}omponent$

```
vector<int> g[N], newg[N], rg[N], todo;
int comp[N], indeg[N];
bool vis[N];
vector<int> gr[N];

void dfs(int k)
{
  vis[k]=1;
  for(auto it:g[k])
  {
   if(!vis[it])
    dfs(it);
  }
  todo.push_back(k);
}

void dfs2(int k, int val)
```

```
comp[k]=val;
for(auto it:rg[k])
 if(comp[it]==-1)
  dfs2(it, val);
void sccAddEdge(int from, int to)
g[from].push_back(to);
rg[to].push_back(from);
void scc()
for(int i=1;i<=n;i++)</pre>
 comp[i]=-1;
for(int i=1;i<=n;i++)</pre>
 if(!vis[i])
  dfs(i);
reverse(todo.begin(), todo.end());
for(auto it:todo)
 if(comp[it]==-1)
  dfs2(it, ++grp);
```

### 29 $\mathbf{Tree}_{C}onstruction_{w}ith_{S}pecific_{\mathcal{V}}$ and $\mathbf{tice}_{\mathbf{S}G-1;i>=0;i--)}$

```
int tim=0;
int parent[LG][N];
int tin[N], tout[N], level[N], vertices[N];
vector<int> g[N], tree[N];
void dfs(int k, int par, int lvl)
tin[k]=++tim;
parent[0][k]=par;
level[k]=lvl;
for(auto it:g[k])
 if(it==par)
  continue;
 dfs(it, k, lvl+1);
tout[k]=tim:
void precompute()
for(int i=1;i<LG;i++)</pre>
 for(int j=1; j<=n; j++)</pre>
  if(parent[i-1][i])
   parent[i][j]=parent[i-1][parent[i-1][j]];
int LCA(int u, int v)
if(level[u] < level[v])</pre>
 swap(u,v);
int diff=level[u]-level[v]:
for(int i=LG-1;i>=0;i--)
 if((1<<i) & diff)</pre>
  u=parent[i][u];
}
 if(u==v)
 return u;
```

```
if(parent[i][u] && parent[i][u]!=parent[i][v])
  u=parent[i][u];
  v=parent[i][v];
return parent[0][u];
bool isancestor(int u, int v) //Check if u is an
   ancestor of v
return (tin[u] <= tin[v]) && (tout[v] <= tout[u]);</pre>
int work()
sort(vertices+1, vertices+k+1, [](int a, int b)
 return tin[a] < tin[b];</pre>
}):
int idx=k:
for(int i=1;i<idx;i++)</pre>
 vertices[++k]=LCA(vertices[i], vertices[i+1]);
 sort(vertices+1, vertices+k+1);
 k=unique(vertices+1, vertices+k+1) - vertices -
sort(vertices+1, vertices+k+1, [](int a, int b)
 return tin[a]<tin[b];</pre>
});
stack<int> s;
s.push(vertices[1]);
for(int i=2:i<=k:i++)</pre>
 while(!isancestor(s.top(), vertices[i]))
  s.pop();
 tree[s.top()].push_back(vertices[i]);
 s.push(vertices[i]);
```

```
for(int i=1;i<=k;i++)
  tree[vertices[i]].clear();
}</pre>
```

### **30** $\mathbf{Z}_{A}lgorithm$

```
//The Z-function for this string is an array of
   length n where the i-th element is equal to
   the greatest number
//of characters starting from the position i that
    coincide with the first characters of s.
vector<int> z_function(string &s)
int n=s.size();
vector<int> z(n);
for(int i=1,l=0,r=0;i<n;i++)</pre>
 if(i<=r)
  z[i]=min(r-i+1, z[i-1]);
 while(i+z[i] < n \&\& s[z[i]] == s[i+z[i]])
  z[i]++:
 if(i+z[i]-1>r)
  l=i, r=i+z[i]-1;
return z;
```

### 31 blockcuttree

```
struct graph
{
  int n;
  vector<vector<int>> adj;
  graph(int n) : n(n), adj(n) {}
  void add_edge(int u, int v)
  {
```

```
adj[u].push_back(v);
   adj[v].push_back(u);
 int add_node()
   adj.push_back({});
   return n++;
 vector<int>& operator[](int u) { return adj[u];
};
vector<vector<int>> biconnected_components(graph
   &adj)
{
 int n = adj.n;
 vector<int> num(n), low(n), art(n), stk;
 vector<vector<int>> comps;
 function<void(int, int, int&)> dfs = [&](int u,
      int p, int &t)
 {
   num[u] = low[u] = ++t:
   stk.push_back(u);
   for (int v : adj[u]) if (v != p)
     if (!num[v])
       dfs(v, u, t);
       low[u] = min(low[u], low[v]);
       if (low[v] >= num[u])
         art[u] = (num[u] > 1 || num[v] > 2);
         comps.push_back({u});
         while (comps.back().back() != v)
          comps.back().push_back(stk.back()),
              stk.pop_back();
       }
     else low[u] = min(low[u], num[v]);
```

```
};
for (int u = 0, t; u < n; ++u)
 if (!num[u]) dfs(u, -1, t = 0);
// build the block cut tree
function<graph()> build_tree = [&]()
  graph tree(0);
  vector<int> id(n);
  for (int u = 0; u < n; ++u)
   if (art[u]) id[u] = tree.add_node();
  for (auto &comp : comps)
   int node = tree.add node():
   for (int u : comp)
     if (!art[u]) id[u] = node;
     else tree.add_edge(node, id[u]);
  return tree:
}:
return comps;
```

### 32 boruvkamst

```
int V = graph - > V, E = graph - > E;
Edge * edge = graph - > edge;
struct subset * subsets = new subset[V];
int * cheapest = new int[V];
for (int v = 0; v < V; ++v) {
    subsets[v].parent = v;
    subsets[v].rank = 0;
    cheapest[v] = -1;
}
int numTrees = V, MSTweight = 0;
while (numTrees > 1) {
    for (int i = 0; i < E; i++) {</pre>
```

```
int set1 = find(subsets, edge[i].src);
   int set2 = find(subsets, edge[i].dest);
   if (set1 == set2)
       continue:
   else {
       if (cheapest[set1] == -1 ||
           edge[cheapest[set1]].weight > edge
               [i].weight)
           cheapest[set1] = i;
       if (cheapest[set1] == -1 ||
           edge[cheapest[set2]].weight > edge
               [i].weight)
           cheapest[set2] = i;
   }
for (int i = 0; i < V; i++) {</pre>
   if (cheapest[i] != -1) {
       int set1 = find(subsets, edge[cheapest
            [i]].src):
       int set2 = find(subsets, edge[cheapest
            [i]].dest):
       if (set1 == set2)
           continue:
       MSTweight += edge[cheapest[i]].weight;
       Union(subsets, set1, set2);
       numTrees--;
}
```

### 33 dpoptimization

```
// Divide and conquer optimization:
// Original Recurrence
// dp[i][j] = min(dp[i-1][k] + C[k][j]) for k < j
// Sufficient condition:
// A[i][j] <= A[i][j+1]</pre>
```

```
// where A[i][j] = smallest k that gives optimal
    answer
// How to use:
// // compute i-th row of dp from L to R. optL <=
     A[i][L] \leftarrow A[i][R] \leftarrow optR
// compute(i, L, R, optL, optR)
1. special
case L == R
//
let M = (L + R) / 2. Calculate dp[i][M] and opt[i] / /
    [M] using O(optR - optL + 1)
//
3. compute(i, L, M - 1, optL, opt[i][M])
4. compute(i, M + 1, R, opt[i][M], optR) ICPC
    Notebook
void compute(int i, int L, int R, int optL, int
    optR) {
   if (L > R) return;
   int mid = (L + R) >> 1, savek = optL;
   dp[i][mid] = inf;
   FOR(k, optL, min(mid - 1, optR)) {
       int cur = dp[i - 1][k] + getCost(k + 1,
           mid);
       if (cur < dp[i][mid]) {</pre>
           dp[i][mid] = cur;
           savek = k;
       }
    compute(i, L, mid - 1, optL, savek);
    compute(i, mid + 1, R, savek, optR);
}
Knuth Optimisation
// Original Recurrence:
// dp[i][j] = min(dp[i][k] + dp[k][j]) + C[i][j]
    for k = i+1..j-1
// Necessary & Sufficient Conditions:
// A[i][j-1] \le A[i][j] \le A[i+1][j]
```

```
// with A[i][j] = smallest k that gives optimal
    answer
 // Also applicable if the following conditions
    are met:
 // 1. C[a][c] + C[b][d] <= C[a][d] + C[b][c] (
    quadrangle inequality)
 // 2. C[b][c] <= C[a][d]
 (monotonicity)
 // for all a <= b <= c <= d
 // To use:
// Calculate dp[i][i] and A[i][i]
// FOR(len = 1..n-1)
 // FOR(i = 1..n-len) {
i = i + len
 //
FOR(k = A[i][j - 1]..A[i + 1][j])
 update(dp[i][j])
 // }
 const int MN = 2011;
int a[MN], dp[MN][MN], C[MN][MN], A[MN][MN];
 int n;
 void solve() {
    cin >> n;
    FOR(i, 1, n) {
        cin >> a[i];
        a[i] += a[i - 1];
    FOR(i, 1, n) FOR(j, i, n) C[i][j] = a[j] - a[
    FOR(i, 1, n) dp[i][i] = 0, A[i][i] = i;
    FOR(len, 1, n - 1)
    FOR(i, 1, n - len) {
        int j = i + len;
        ICPC Notebook
        dp[i][j] = 2000111000;
        FOR(k, A[i][i-1], A[i+1][i]) {
```

```
int cur = dp[i][k - 1] + dp[k][j] + C[
               i][i];
           if (cur < dp[i][j]) {</pre>
               dp[i][j] = cur;
               A[i][j] = k;
           }
       }
   }
   cout << dp[1][n] << endl;</pre>
//SOS DP - Initialize base case of nums[i] - Sum
    of all subsets
rep(i, 0, 17) rep(j, 0, N) if ((j >> i) & 1) nums
    [j] += nums[j ^ (1 << i)];
void neg(int f[], int upto) {
   for (int i = 0; i <= upto; i++)</pre>
       if (pc[i] \% 2 == 0) f[i] = -f[i];
void sos(int f[], int upto, int t) {
   for (int i = 0; i < t; ++i)
       for (int mask = 0; mask <= upto; ++mask)</pre>
           if (mask & (1 << i))</pre>
               f[mask] += f[mask ^ (1 << i)];
       }
neg(cnt, upto);
sos(cnt, upto, t);
```

### 34 euler

```
//Start with an empty stack and an empty circuit
   (eulerian path).
//- If all vertices have even degree - choose any
   of them.
//- If there are exactly 2 vertices having an odd
   degree - choose one of them.
```

```
//- Otherwise no euler circuit or path exists.
//If current vertex has no neighbors - add it to
   circuit, remove the last vertex from the
   stack and set it as the current one.
//Otherwise (in case it has neighbors) - add the
   vertex to the stack, take any of its
   neighbors, remove the edge between
//selected neighbor and that vertex, and set that
    neighbor as the current vertex.
//Repeat step 2 until the current vertex has no
   more neighbors and the stack is empty.
//Note that obtained circuit will be in reverse
   order - from end vertex to start vertex.
//Code stores the Euler Circuit path in Circuit
void EulerTour(int k) //Heirholzer's Algorithm
   int cur = k;
   stack < int > temp;
   while (true) {
       int ct = g[cur].size();
       if (!ct) {
          viscircuit[cur] = 1:
           circuit.push_back(cur);
           if (temp.size() == 0)
              break;
           cur = temp.top();
          temp.pop();
       } else {
           pair < int, int > next = * (g[cur].
              begin());
           g[cur].erase(next);
           g[next.ff].erase(mp(cur, next.ss));
           if (next.ss == 1) { in [next.ff]++;
              out[cur]++:
              ans.pb(mp(cur, next.ff));
          temp.push(cur);
           cur = next.ff;
   }
```

### 35 implicitsegtree

```
//Implicit Segment Tree :
//For assigning a unique number to every visitor,
     such that the unique number >=x, and has not
     been taken. Input:
//1 x = Tourist enters, 2 x = Tourist leaves.
struct nd {
   11 t, x;
struct segmenttree {
   int val;
   segmenttree * left, * right;
   segmenttree() {
       val = 0;
       left = NULL:
       right = NULL;
};
int n, num = 1e6;
nd queries[N];
void update(segmenttree * root, int L, int R, int
     pos, int type) {
   if (L == R) {
       if (type == 0) {
           root - > val = 0;
       elseICPC Notebook {
           root - > val = 1;
       }
       return;
   int M = (L + R) \gg 1;
   if (pos <= M)
       update(root - > left, L, M, pos, type);
   else
```

```
update(root - > right, M + 1, R, pos,
           type);
   root - > val = (root - > left) - > val * (
       root - > right) - > val;
int query(segmenttree * root, int L, int R, int
   pos) {
   if (L == R) {
       if (root - > val == 0)
           return L;
       else
           return -1;
   int M = (L + R) \gg 1;
   if (root - > left == NULL) {
       root - > left = new segmenttree();
       root - > right = new segmenttree();
   if (pos <= M) {</pre>
       if ((root - > left) - > val == 1) {
           return query(root - > right, M + 1, R,
               pos);
       } else {
           int store = query(root - > left, L, M,
               pos);
           if (store == -1) {
              return query(root - > right, M +
                  1, R, pos);
          } else {
              return store;
          }
       }
   } else {
       return query(root - > right, M + 1, R,
           pos);
   }
```

### 36 manacher

```
void Manachers() {
   int N = strlen(text);
   N = 2 * N + 1; //Position count
   int L[N]; //LPS Length Array
   L[0] = 0;
   L[1] = 1;
   int C = 1, R = 2, i = 0;
   int iMirror, maxLPSLength = 0,
       maxLPSCenterPosition = 0;
   int start = -1, end = -1, diff = -1;
   for (i = 2; i < N; i++) {</pre>
       iMirror = 2 * C - i;
       L[i] = 0;
       diff = R - i;
       if (diff > 0)
          L[i] = min(L[iMirror], diff);
       while (((i + L[i]) < N && (i - L[i]) > 0)
           (((i + L[i] + 1) \% 2 == 0) | |
              (text[(i + L[i] + 1) / 2] == text
                  [(i - L[i] - 1) / 2]))
          L[i]++;
       if (L[i] > maxLPSLength) {
          maxLPSLength = L[i];
          maxLPSCenterPosition = i;
       if (i + L[i] > R) {
          C = i:
          R = i + L[i]:
       }
   }
   start = (maxLPSCenterPosition - maxLPSLength)
        / 2;
   TCPC Notebook
   end = start + maxLPSLength - 1;
   printf("LPS of string is %s : ", text);
   for (i = start; i <= end; i++)</pre>
       printf("%c", text[i]);
```

```
printf("\n");
```

```
37
       maths
  //1) Sum of values of totient functions of all
     divisors of n is equal to n.
//4)
        Bell Number:
//In combinatorial mathematics, the Bell numbers
    count the possible partitions of a set. The
   nth of these numbers,
//Bn, counts the number of different ways to
   partition a set that has exactly n elements,
   or equivalently, the number of
//equivalence relations on it.
void processbell()
   bell[0]=1:
   bell[1]=1;
   for(int i=2;i<=5000;i++)</pre>
       for(int j=0; j<=i-1; j++)</pre>
           bell[i]+=nCr(i-1, j) * bell[j];
           bell[i]%=MOD;
   }
       Stirling number of the second kind:
//In mathematics, particularly in combinatorics,
   a Stirling number of the second kind (or
   Stirling partition number) is the
//number of ways to partition a set of n objects
   into k non-empty subsets and is denoted by S(
//Equivalently, they count the number of
   different equivalence relations with
   precisely k equivalence classes that can be
```

```
//defined on an n element set.
//Value of S(n, k) can be defined recursively as,
     S(n+1, k) = k*S(n, k) + S(n, k-1)
//with S(0, 0) = 1 and S(n, 0) = S(0, n) = 0 for
        Modular Inverse modulo N (General) :
int modInverse(int a, int m)
   int m0 = m;
   int y = 0, x = 1;
   if (m == 1)
     return 0;
   while (a > 1)
       int q = a / m;
       int t = m;
       m = a \% m, a = t;
       t = y;
       y = x - q * y;
       x = t:
   if (x < 0)
      x += m0:
   return x;
     Chinese Remainder Theorem:
int findMinX(int num[], int rem[], int k)
₹
int prod = 1;
for (int i = 0; i < k; i++)</pre>
 prod *= num[i];
 int result = 0:
 for (int i = 0; i < k; i++)</pre>
 int pp = prod / num[i];ICPC Notebook
 result += rem[i] * inv(pp, num[i]) * pp;
return result % prod;
```

```
}
//nCr%m when n.r ~ 10^18, m ~10^6
#define N 1000005
lld maxp[N];
lld extended_euclid(lld a,lld b,lld &x,lld &y) {
   11d xx = y = 0;
   lld vv = x = 1;
   while (b) {
   int q = a/b;
   int t = b; b = a%b; a = t;
   t = xx; xx = x-q*xx; x = t;
   t = yy; yy = y-q*yy; y = t;
   return a;
lld mod(lld a,lld b) {
   return ((a%b)+b)%b;
lld inversemod(lld a,lld n) {
   11d x, y;
   lld d = extended_euclid(a, n, x, y);
   if(d > 1) return -1LL:
   return mod(x,n);
plld chinese_remainder_theorem(lld x,lld a,lld y,
   11d b) {
 lld s, t;
 lld d = extended_euclid(x, y, s, t);
 if(a%d != b%d) return make_pair(OLL, -1LL);
 return make_pair(mod(s*b*x+t*a*y,x*y)/d, x*y/d)
plld chinese_remainder_theorem(const vlld &x,
    const vlld &a) { // x are the modules and a
    are remainders
   plld ret = make_pair(a[0], x[0]);
   11d tmp=x.size();
   rep(i,0,tmp) {
       ret = chinese_remainder_theorem(ret.
           second, ret.first, x[i], a[i]);
```

```
if (ret.second == -1) break:
   return ret:
lld countFact(lld n,lld p)
   11d k=0;
   while (n>=p) k+=n/p,n/=p;
   return k;
lld factorial_mod(lld n,lld m)
   lld res=1;
   while(n>0)
       for(11d i=2, m=n%MOD; i<=m; i++)</pre>
           res=(res*i) % MOD;
       if ((n/=MOD)\%2 > 0)
           res = MOD - res;
   return res:
lld nCk_get_prime_pow(lld n,lld k,lld p) {
   11d res=countFact(n,p)-countFact(k,p)-
        countFact(n-k,p);
   return res;
lld nCk_get_non_prime_part(lld n,lld k,lld p,lld
    e) {
   11d pe=powm(p,e,LLINF);
   lld r=n-k,acc=1;
   vlld fact_pe;
   fact_pe.pb(1LL);
   rep(x,1,pe) {
       if(x\%p == 0) fact_pe.pb(acc);
       else acc=(acc*x)%pe,fact_pe.pb(acc);
   }
   1ld top=1,bottom=1,is_neg=0,digits=0;
    while(n!=0) {
       if(acc!=1) {
```

```
if(digits>=e) {
              is_neg ^= (n&1);
              is_neg ^= (r&1);
              is_neg ^= (k&1);
       top=(top*fact_pe[n%pe])%pe;
       bottom=(bottom*fact_pe[r%pe])%pe;
       bottom=(bottom*fact_pe[k%pe])%pe;
       n/=p,r/=p,k/=p;
       digits+=1;
   lld res=(top*inversemod(bottom,pe))%pe;
   if(p!=2 or e<3)
       if(is_neg) res=pe-res;
   return res:
//Sum of GP in LogN
lld solve(lld x,lld n,lld m){
 if(n==0) return 1LL:
 if(n==1) return (1LL+x)%m;
 if(n\%2==0){
   lld t1=solve((x*x)%m,n/2LL-1LL,m);
   t1=(t1*(1LL+x))%m;
   t1=(t1+power(x,n,m))%m;
   return t1;
 }
 else{
   lld t1=solve((x*x)%m,n/2LL,m);
   t1=(t1*(1LL+x))%m;
   return t1;
 }
```

### 38 persistent

```
int build(int L, int R) {
   int node = ++ct;
   if (L == R) {
       return node:
   int M = (L + R) \gg 1;
   lc[node] = build(L, M);
   rc[node] = build(M + 1, R);
   return node;
int update(int onode, int L, int R, int pos) {
   int node = ++ct;
   if (L == R) {
       st[node] = st[onode] + 1;
       return node;
   }
   int M = (L + R) >> 1:
   lc[node] = lc[onode];
   rc[node] = rc[onode];
   if (pos <= M)
       lc[node] = update(lc[onode], L, M, pos);
       rc[node] = update(rc[onode], M + 1, R,
           pos);
   st[node] = st[lc[node]] + st[rc[node]];
   return node;
}
int query(int nodeu, int nodev, int L, int R, int
    pos) {
   if (L == R) {
       return L;
   }
   int M = (L + R) >> 1;
   int leftval = st[lc[nodev]] - st[lc[nodeu]];
   int rightval = st[rc[nodev]] - st[rc[nodeu]];
   if (leftval >= pos) {
```

### 39 suffixarray

```
int suffixRank[20][int(1E6)];
struct myTuple {
   int originalIndex; // stores original index
       of suffix
   int firstHalf: // store rank for first half
       of suffix
   int secondHalf:
   // store rank for second half of suffix
int cmp(myTuple a, myTuple b) {
   if (a.firstHalf == b.firstHalf) return a.
       secondHalf < b.secondHalf;</pre>
   else return a.firstHalf < b.firstHalf;</pre>
int N = s.size();
for (int i = 0; i < N; ++i)
   suffixRank[0][i] = s[i] - 'a':
mvTuple L[N]:
for (int cnt = 1, stp = 1; cnt < N; cnt *= 2, ++</pre>
    stp) {
   for (int i = 0; i < N; ++i) {</pre>
       L[i].firstHalf = suffixRank[stp - 1][i];
       L[i].secondHalf = i + cnt < N ?
           suffixRank[stp - 1][i + cnt] : -1;
       L[i].originalIndex = i;
```

```
sort(L, L + N, cmp);
   suffixRank[stp][L[0].originalIndex] = 0;
   for (int i = 1, currRank = 0; i < N; ++i) {</pre>
       if (L[i - 1].firstHalf != L[i].firstHalf
           || L[i - 1].secondHalf != L[i].
           secondHalf)
           ++currRank;
       suffixRank[stp][L[i].originalIndex] =
           currRank;
   }
//KASAI
vector < int > kasai(string txt, vector < int >
    suffixArr) {
       int n = suffixArr.size();
       vector < int > lcp(n, 0);
       vector < int > invSuff(n, 0);
       for (int i = 0; i < n; i++)</pre>
           invSuff[suffixArr[i]] = i;
       int k = 0:
       for (int i = 0; i < n; i++) {</pre>
           if (invSuff[i] == n - 1) {
              k = 0;
               continue;
           }
           int j = suffixArr[invSuff[i] + 1];
           while (i + k < n \&\& j + k < n \&\& txt[i]
                + k] == txt[i + k]
               k++;
           lcp[invSuff[i]] = k; // lcp for the
               present suffix.
           if (k > 0)
              k--:
       }
```

return lcp;