COMP2010 Data Structures and Algorithms

Lecture 16: Depth-First Search

Department of Computer Science & Technology
United International College



Depth-First Search (DFS)

- DFS is another popular graph search strategy
 - Idea is similar to pre-order traversal (visit node, then visit children recursively)

- DFS can provide certain information about the graph that BFS cannot
 - ◆ It can tell whether we have encountered a cycle or not

DFS Algorithm

- DFS will continue to visit neighbors in a recursive pattern
 - Whenever we visit v from u, we recursively visit all unvisited neighbors of v. Then we backtrack (return) to u.

Note: it is possible that w2 was unvisited when we recursively visit w1, but became visited by the time we return from the recursive call.

DFS Algorithm

Algorithm DFS(s)

- 1. **for** each vertex v
- 2. **do** flag[v] := false;
- 3. RDFS(s);

Flag all vertices as not visited

Algorithm RDFS(v)

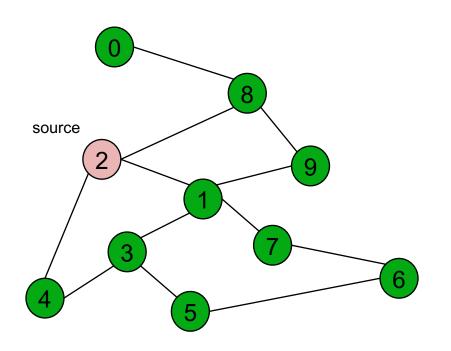
- 1. flag[v] := true;
- 2. **for** each neighbor w of v
- 3. **do if** flag[w] = false
- 4. then RDFS(w);

Flag yourself as visited

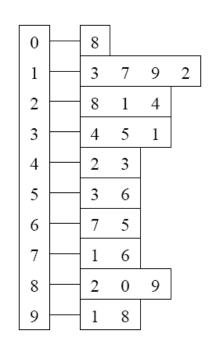
For unvisited neighbors, call RDFS(w) recursively

We can also record the paths using pred[].

Example



Adjacency List

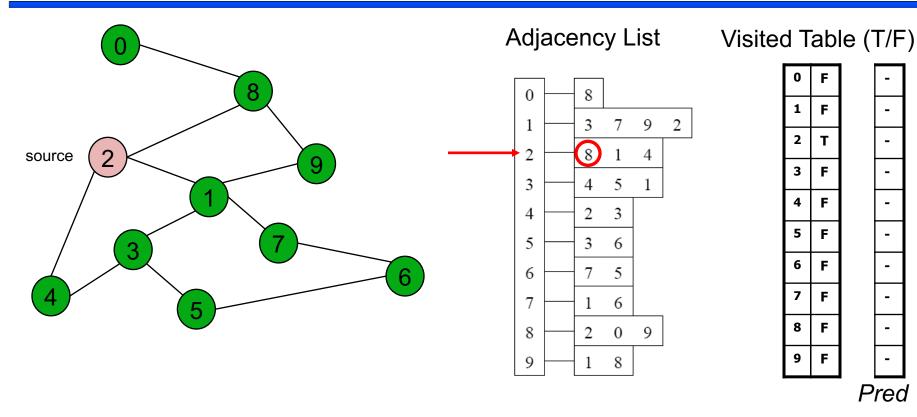


Visited Table (T/F)

0	F		-
1	F		-
2	F		-
3	F		-
4	F		-
5	F		-
6	F		-
7	F		-
8	F		-
9	F		-
		P	rea

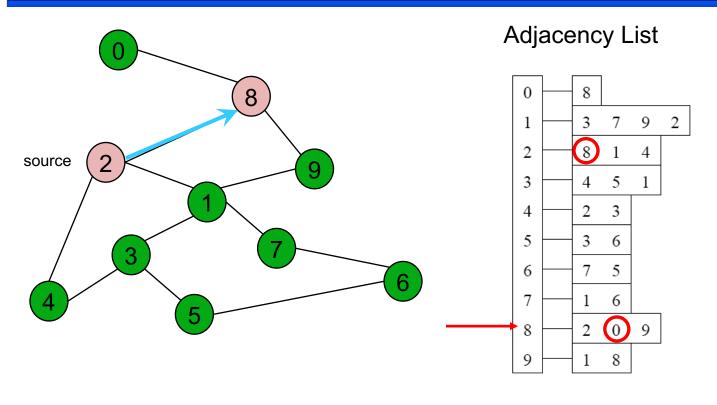
Initialize visited table (all False)

Initialize Pred to -1



Mark 2 as visited

RDFS(2)
Now visit RDFS(8)



Visited Table (T/F)

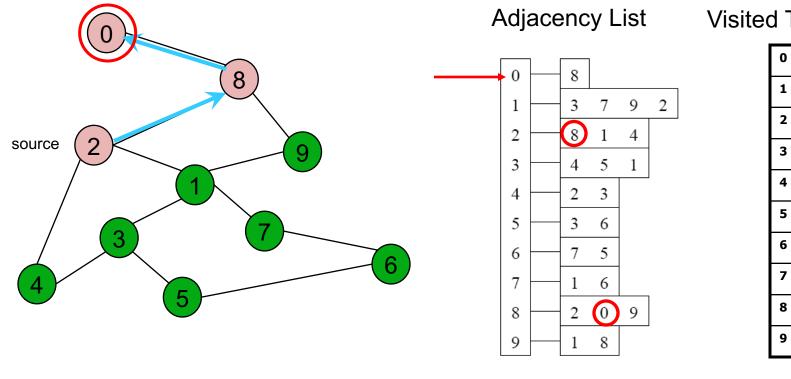
0	F		-
1	F		-
2	T		-
3	F		-
4	F		-
5	F		-
6	F		-
7	F		-
8	T		2
9	F		-
Pred			

Mark 8 as visited

mark Pred[8]

Recursive RDFS(2) calls RDFS(8)

2 is already visited, so visit RDFS(0)



Visited Table (T/F)

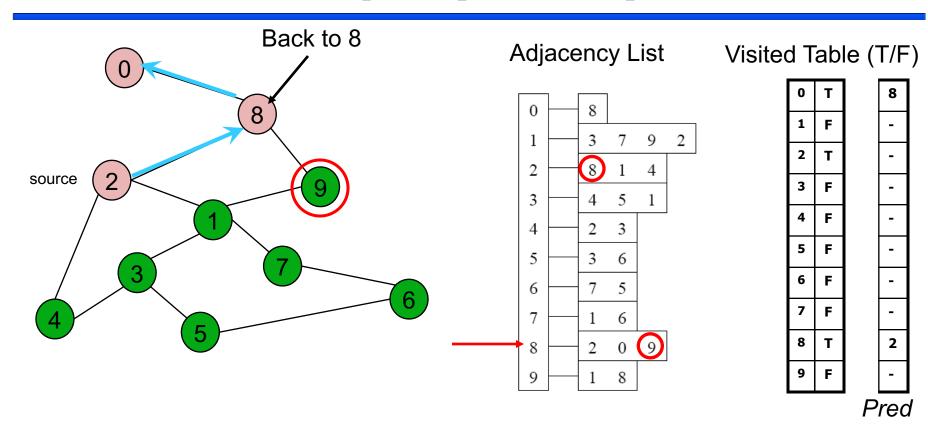
0	T		8
1	F		-
2	T		-
3	F		-
4	F		-
5	F		-
6	F		•
7	F		-
8	T		2
9	F		-
		P	rec

Mark 0 as visited

Mark Pred[0]

Recursive RDFS(2) calls RDFS(8)

RDFS(0) -> no unvisited neighbors, return to call RDFS(8)



Recursive RDFS(2)
calls RDFS(8)
Now visit 9 -> RDFS(9)

RDFS(2)

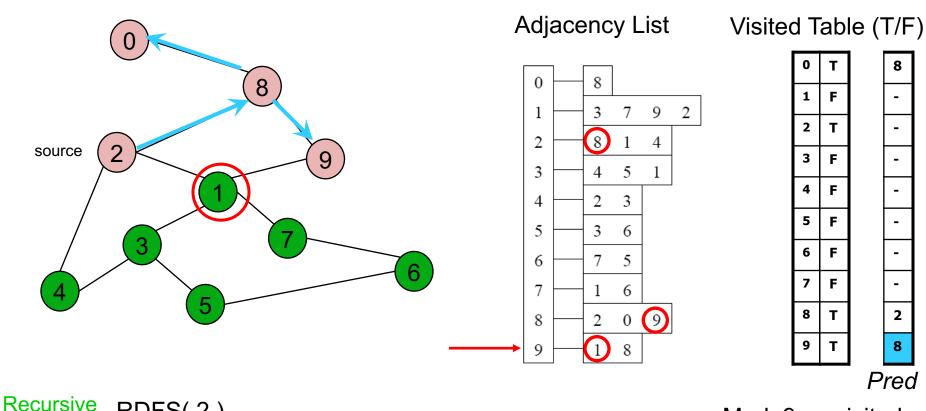
RDFS(8)

RDFS(9)

-> visit 1, RDFS(1)

calls

Example (Cont'd)



Mark 9 as visited

2

8

Pred

Mark Pred[9]

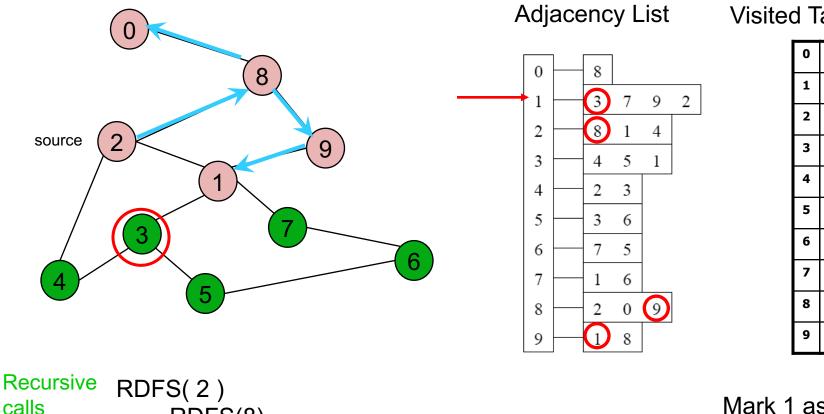
RDFS(8)

RDFS(9)

RDFS(1)

visit RDFS(3)

Example (Cont'd)

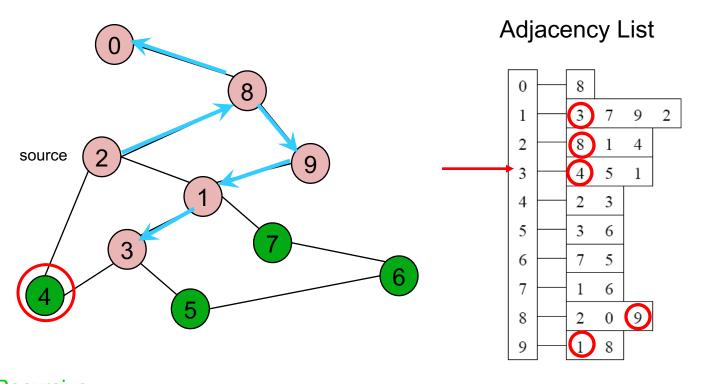


Visited Table (T/F)

0	Т		8
1	T		9
2	T		
3	F		-
4	F		
5	F		-
6	F		•
7	F		
8	T		2
9	T		8
		P	rec

Mark 1 as visited

Mark Pred[1]



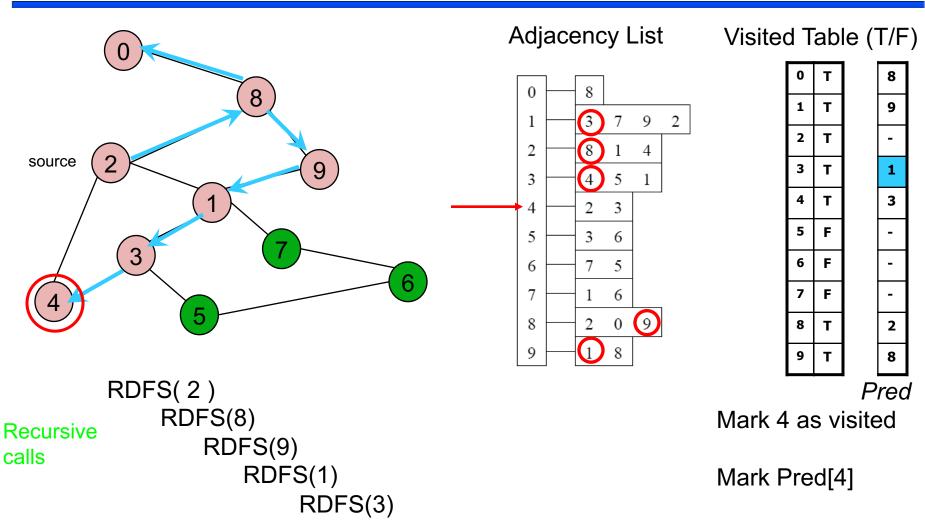
Visited Table (T/F)

0	T		8
1	Т		9
2	T		-
3	T		1
4	F		-
5	F		-
6	F		-
7	F		-
8	T		2
9	T		8
Pred			

Mark 3 as visited

Mark Pred[3]

Recursive RDFS(2)
calls RDFS(8)
RDFS(9)
RDFS(1)
RDFS(3)
visit RDFS(4)



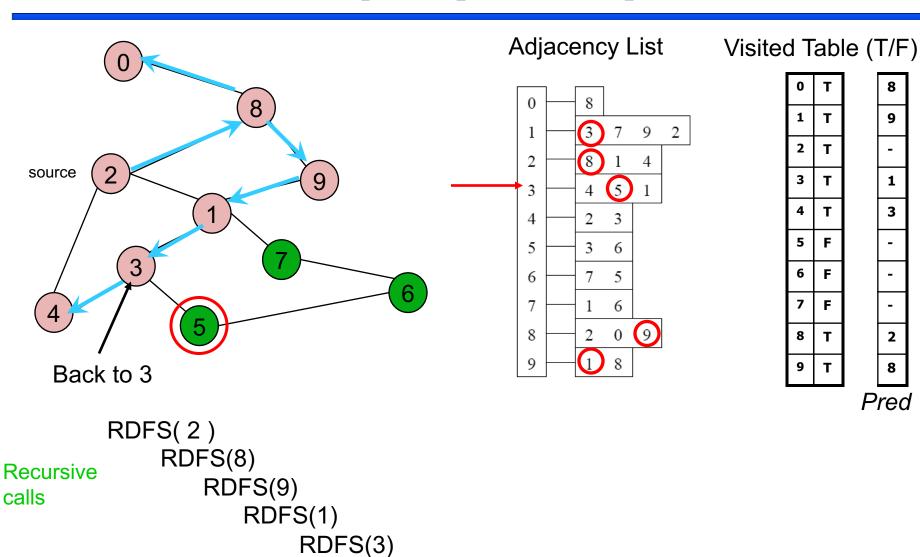
RDFS(4) → STOP all of 4's neighbors have been visited return back to call RDFS(3)

9

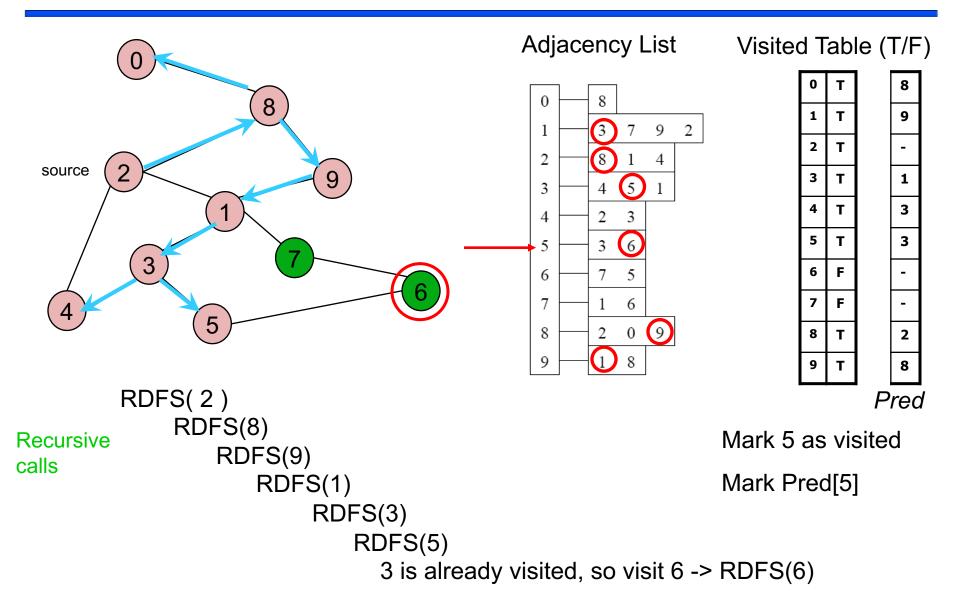
1

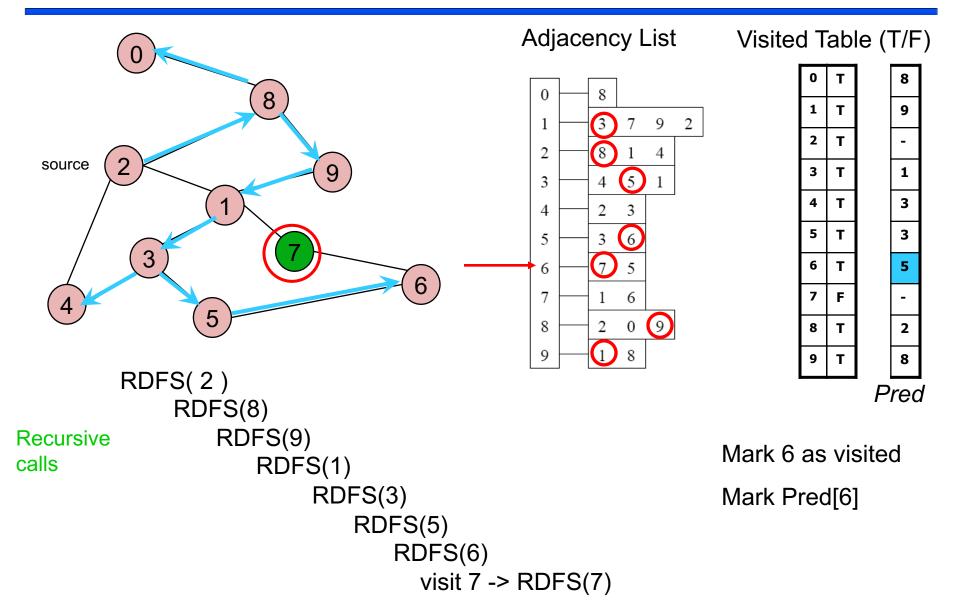
3

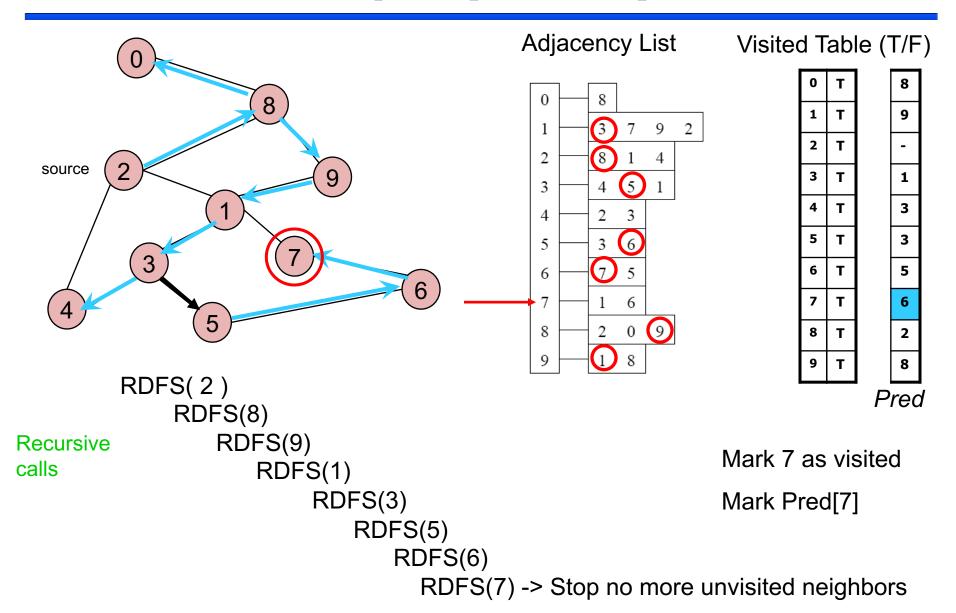
2

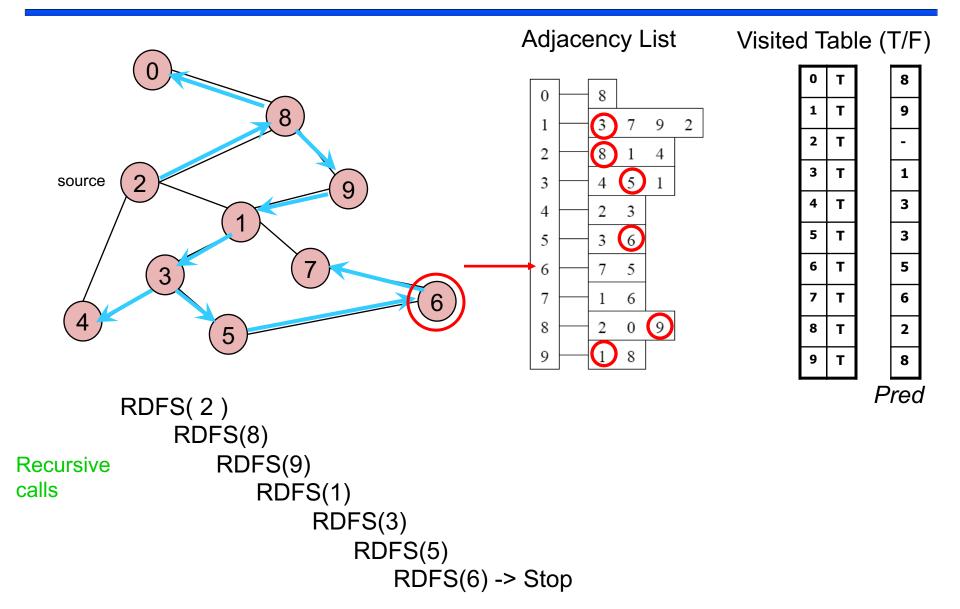


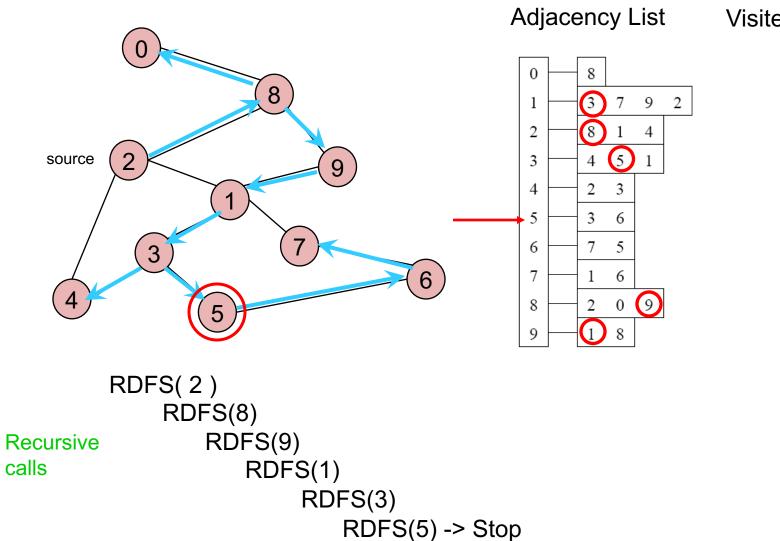
visit 5 -> RDFS(5)



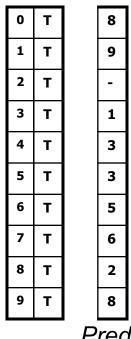




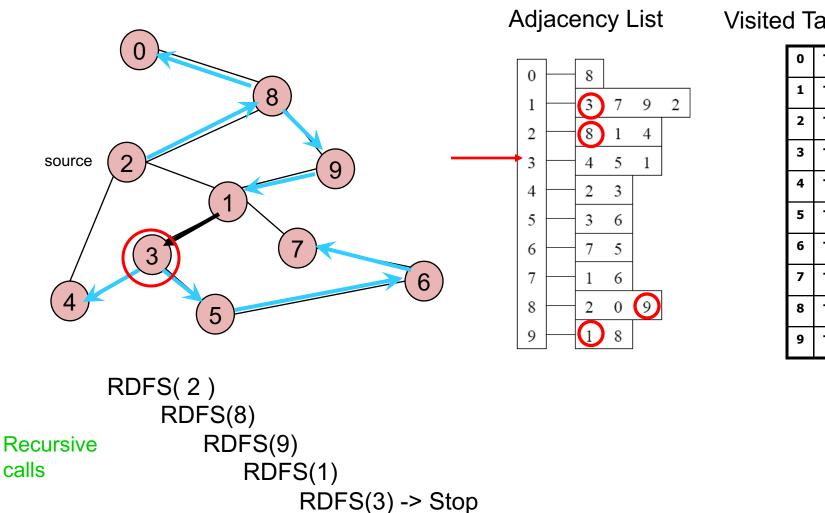




Visited Table (T/F)

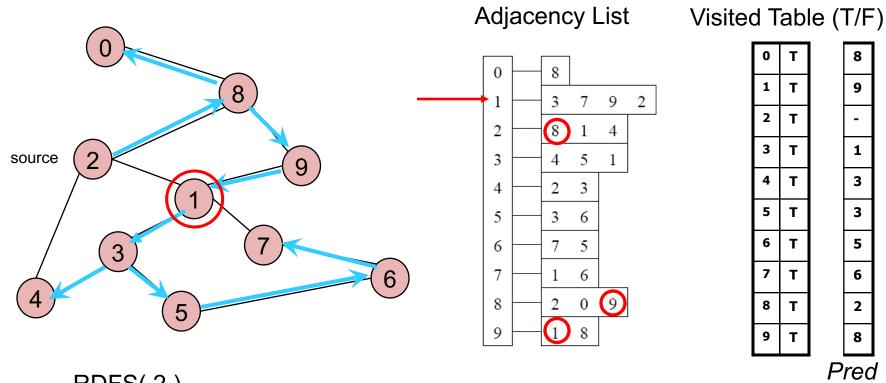


Pred



Visited Table (T/F)

		_	
0	Т		8
1	T		9
2	T		-
3	Т		1
4	T		3
5	T		3
6	T		5
7	T		6
8	T		2
9	T		8
		P	rec



9

1

3

3

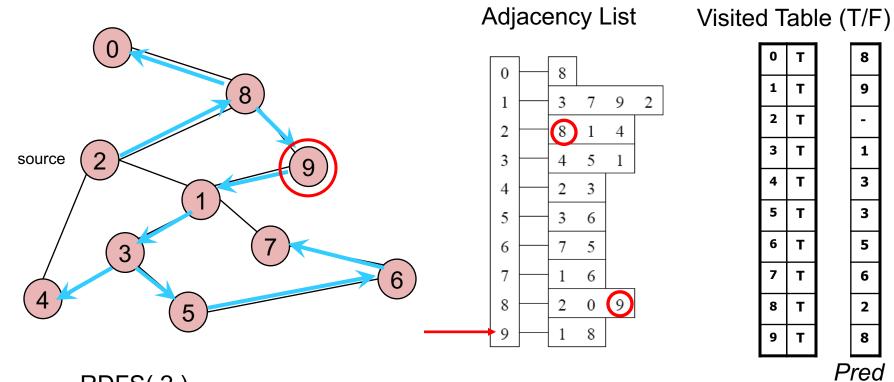
5

6

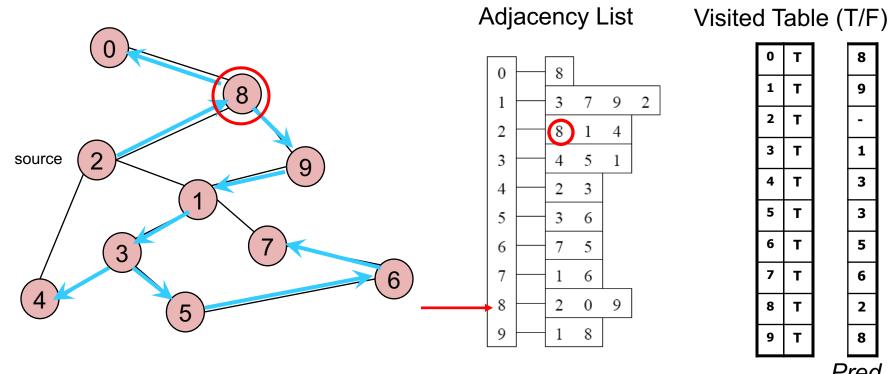
2

8

RDFS(2) RDFS(8) RDFS(9) Recursive RDFS(1) -> Stop calls



RDFS(2)
RDFS(8)
Recursive RDFS(9) -> Stop calls

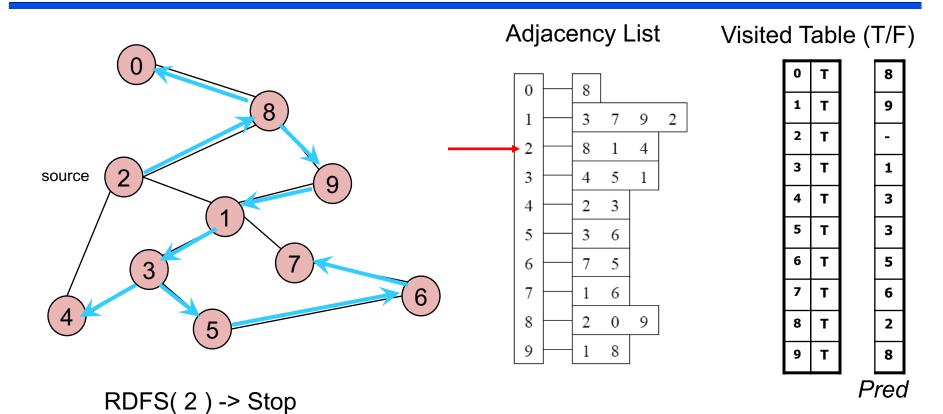


		_	
0	Т		8
1	T		9
2	T		-
3	Т		1
4	T		3
5	Т		3
6	Т		5
7	Т		6
8	Т		2
9	T		8
		P	rec

RDFS(2) RDFS(8) -> Stop

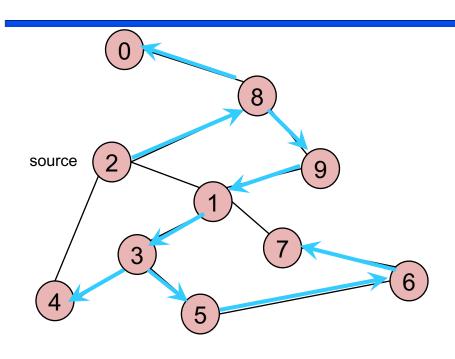
Recursive calls

Example Finished



Recursive calls finished

DFS Path Tracking



Adjacency List

Visited Table (T/F)

0	Т		8
1	T		9
2	T		-
3	T		1
4	T		3
5	T		3
6	Т		5
7	Т		6
8	T		2
9	T		8
		P	rec

Pred

DFS find out path too

Algorithm Path(w)

- 1. if $pred[w] \neq -1$
- 2. then
- 3. Path(pred[w]);
- 4. output w

Try some examples.

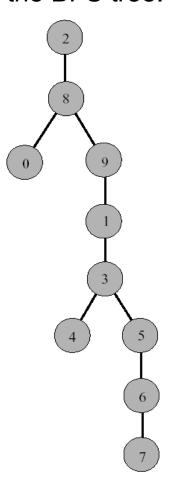
- Path(0) ->
- Path(6) ->
- Path(7) ->

DFS Tree

Resulting DFS-tree.

Notice it is much "deeper" than the BFS tree.





Captures the structure of the recursive calls

- when we visit a neighbor w of v, we add w as child of v
- whenever DFS returns from a vertex v, we climb up in the tree from v to its parent

Time Complexity of DFS

(Using adjacency list)

- We never visited a vertex more than once
- We had to examine all edges of the vertices
 - We know $\Sigma_{\text{vertex } \nu} \text{degree}(\nu) = 2m$ where m is the number of edges
- So, the running time of DFS is proportional to the number of edges and number of vertices (same as BFS)
 - \bullet O(n + m)
- You will also see this written as:
 - O(|v|+|e|) |v| = number of vertices (n) |e| = number of edges (m)