

CS5130 : Mathematical Tools for Theoretical Computer Science

(Scribe Lecture Notes)

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Preface

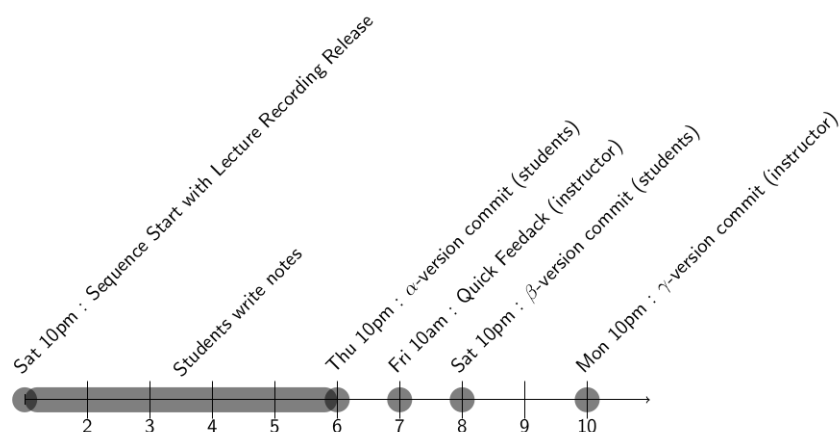
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Scribe status

Each lecture has a field called **status**. It tells which stage of the edit pipeline is the document currently. The scribe notes are due on Thursdays and Saturdays as per the following timeline.



Even after these edits, it is possible that there are still errors in the draft, which may not get noticed. If you find errors still, please report to the instructor.

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Todo list

- 1: Jayalal says: Todo - Prove that f is a bijection v
- 2: Jayalal says: Todo - Establish bijections from *Euler's* problem to *Full binary tree* problem and *hand-shaking* problem to *balanced parenthesised strings* problem x

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Lecture 1

Pigeon Hole Principle and Basic Applications

1.1 Example 1

This is a test example lecture notes.

1.2 Example 2

1.3 Example 3

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Lecture 5

Multichoosing

5.1 Introduction

Consider the definition of *set*. We know that it's a well defined collection of *distinct* objects. From a collection of n distinct symbols, the number of ways to form a *set* of length k is given by $\binom{n}{k}$. Now let's consider the definition of *multi-set*. It's similar to that of a *set*, except that it allows repetition of objects. Now it's natural ask the following question: From a collection of n distinct symbols, what is the number of ways to form a *multi-set* of length k . Multichoosing exactly answers this questions. In this lecture, we explore multichoosing in detail. We discuss several equivalent bijections to this problem and come-up with an algebraic expression for $\left(\binom{n}{k}\right)$ (spelled out as n multi-choose k).

5.2 Equivalent bijections

5.2.1 Non-negative solutions

Formally, $\left(\binom{n}{k}\right)$ is the number of ways of choosing k objects from a set of n objects where the order is not important but repetitions are allowed. For all $i = 1, 2, \dots, n$, if we denote by x_i the number of copies of i^{th} object we choose, then we have the equation

$$x_1 + x_2 + \dots + x_n = k \quad (5.1)$$

where each $x_i \geq 0$. Therefore, number of *non-negative* integral solutions to this equation gives us the required number of ways of choosing k objects from n objects with given conditions. Let's look at an equivalent problem and establish a bijection between these two.

5.2.2 Voting problem

If n candidates are contesting in an election and there are k voters, how many ways can votes of those k voters be distributed among n candidates?

If we denote by x_i , the number of votes received by i^{th} candidate and there are k voters, we have $x_1 + x_2 + \dots + x_n = k$ and thus, the number of ways of dividing votes among candidates

is the number of non-negative solutions to the equation 5.1. Formally, we can define a bijection f from set of solutions to the equation 5.1 to set of ways of dividing the votes among n candidates.

Definition: f takes the tuple $\mathbf{x} = (x_1, x_2, \dots, x_n)$ and assign x_i number of votes to i^{th} candidate where $i = 1, 2, \dots, n$.

Well defined: f is well defined because for every valid tuple $\mathbf{x} = (x_1, x_2, \dots, x_n)$, we have $x_1 + \dots + x_n = k$ and thus summing over votes received by i^{th} where $i = 1, 2, \dots, n$ will be k votes in total.

Injective: f is an injection because for every valid way of dividing the votes among candidates, there's a unique solution tuple in which x_i = number of votes received by i^{th} candidate. In other words, for any two $\mathbf{x}_1 \neq \mathbf{x}_2$, there exists an $i \in [n]$ such that $x_{1_i} \neq x_{2_i}$ and i^{th} candidate gets different votes. Thus $f(\mathbf{x}_1) \neq f(\mathbf{x}_2)$.

Surjective: f is surjective because for every way of dividing k votes among n candidates, there is a pre-image $\mathbf{x} = (x_1, \dots, x_n)$ which is a valid solution to the equation 5.1 (as there are a total of k voters, sum of number of votes received by each voter must sum up to k).

Thus f is a bijection from the set of non-negative solutions to $x_1 + \dots + x_n = k$ to the set of ways of dividing k votes among n candidates.

5.2.3 Non-decreasing subsequences

Number of non-decreasing sequences of integers between 1 and n of length k . A non-decreasing sequence is of the form $\{a_1, a_2, \dots, a_k\}$ where $1 \leq a_1 \leq a_2 \leq \dots \leq a_k \leq n$. Lets define a bijection f from set of non-negative integral solutions to Eqn. 5.1 to set of non-decreasing sequences between 1 and n of length k .

Definition: f takes $\mathbf{x} = (x_1, \dots, x_n)$ as input and writes the number i x_i times for all $i = 1, 2, \dots, n$ to obtain a sequence of length k .

Well defined: As f constructs the sequence in increasing order from 1 to n by writing i x_i times, the resulting sequence will be non-decreasing. Therefore, f is well defined.

Injective: For every $\mathbf{x}_1 \neq \mathbf{x}_2$, there exists an i such that $x_{1_i} \neq x_{2_i}$ and thus in the resulting sequences, number i is written different number of times. Therefore, f is injective.

Surjective: Every non-decreasing sequence of integers between 1 and n of length k has a pre-image $\mathbf{x} = (x_1, \dots, x_n)$ which is a valid solution to equation 5.1 (where x_i is the number of times the number i is present in the sequence and as length of sequence is k , all x_i 's where $i = 1, 2, \dots, n$ sum up to k).

Thus f is a bijection.

5.2.4 Stars and bars problem

There are k stars placed horizontally. Find the number of ways to place $n - 1$ bars in between those k stars. Lets define a bijection f from set of non-negative integral solutions to Eqn. 5.1 to set of ways of placing $n - 1$ bars among k stars.

Definition: f takes $\mathbf{x} = (x_1, \dots, x_n)$ as input and place x_i number of stars between $(i - 1)^{th}$ bar and i^{th} bar. We leave it as an exercise to prove that f is well-defined, injective and surjective.

1: Jayalal says: Todo - Prove that f is a bijection

5.3 Algebraic expression

So far in Sec. 5.2, we have established bijections between *non-negatives integral* solutions of Eq. 5.1 and various other problems and argued that number of ways of solving any particular problem is equal to the number of non-negative integral solutions to Eq. 5.1. In this section, we are interested in coming up with a concrete expression for $\binom{n}{k}$ by solving it's equivalent bijection.

Method 1 Let's solve the *stars and bars* problem defined in Sec. 5.2.4. Let's use the fact that any placement of $n - 1$ bars among k stars can be equivalently thought of as a string of length $n + k - 1$ over the alphabet $\{\star, |\}$ with k \star 's. Therefore,

$$\begin{aligned} \text{number of ways of placing } n - 1 \text{ bars among } k \text{ stars} &= \text{number of such strings} \\ &= \binom{n + k - 1}{k} \end{aligned}$$

Method 2 Let's solve the *Non-decreasing subsequences* problem defined in Sec. 5.2.1. Let's establish a bijection f from set β of non-decreasing subsequences of integers between 1 and n of length k to a set Γ of strictly increasing subsequences of integers between 1 and $n + k - 1$ of length k . A strictly increasing subsequence is of the form $1 \leq b_1 < b_2 < \dots < b_k \leq n + k - 1$

Definition: f takes as input a non-decreasing subsequence (a_1, a_2, \dots, a_k) between 1 and n and for all $i = 1, 2, \dots, k$ set $b_i = a_i + i - 1$ and output the sequence (b_1, b_2, \dots, b_k)

Well defined: For any $(a_1, a_2, \dots, a_k) \in \beta$, we have for all $i = 1, 2, \dots, k - 1$,

$$\begin{aligned} a_i &\leq a_{i+1} \\ a_i + i &\leq a_{i+1} + i \\ a_i + i - 1 &< a_{i+1} + i \\ b_i &< b_{i+1} \end{aligned}$$

Therefore, the subsequence (b_1, \dots, b_k) is strictly increasing subsequence and thus f is well defined.

Injective: For every non-decreasing subsequence (a_1, \dots, a_k) , there's a unique strictly increasing subsequence (b_1, \dots, b_k) where for all $i = 1, \dots, k$, $b_i = a_i + i - 1$. Therefore f is injective.

Surjective: For every strictly increasing subsequence (b_1, \dots, b_k) , there's a pre-image (a_1, \dots, a_k) which is non-decreasing where for all $i = 1, \dots, k$, $a_i = b_i - i + 1$

Therefore, f is a bijection. The number of ways of choosing a strictly increasing subsequence (b_1, \dots, b_k) between integers 1 and $n + k - 1$ is just choosing k integers from first $n + k - 1$ integers and arrange them in one way (increasing order). Therefore number of ways = $\binom{n+k-1}{k}$. As f is a bijection, therefore, the number of non-decreasing subsequences between 1 and n of length k are $\binom{n+k-1}{k}$

Method 3 Let's solve the *Voting* problem defined in Sec. 5.2.2. Let's ask a slightly modified question.

Question: How many ways to distribute m votes among n candidates such that each candidate gets at least one vote.

Answer 1: As every candidate gets at least one vote, let's first distribute one vote each to each of the n candidate and then distribute the remaining $m - n$ votes among n candidates. By the bijection defined in Sec. 5.2.2, the number of ways of distributing $m - n$ votes among n candidates is $\binom{n}{m-n}$

Answer 2: Let's interpret votes as \star s. Then the question essentially reduces to placing $n - 1$ bars (since there are n candidates, we divide by placing $n - 1$ bars) among m stars (since there are m voters). i^{th} candidate gets votes equal to number of stars between $(i - 1)^{th}$ | and i^{th} |. However, there are two additional constraints

1. A bar cannot be placed in the beginning or in the end (if not then either the first candidate or the last candidate gets 0 votes)
2. We cannot place two | s between same two \star s (if we place $(i - 1)^{th}$ | and i^{th} | between same two \star s, the i^{th} candidate gets 0 votes)

Hence, we have to choose $n - 1$ gaps among the $m - 1$ gaps (because we have $m + 1$ gaps and by cond. 1 we remove two) to place $n - 1$ | s without repetitions (because repeating violates cond. 2). Therefore, there are $\binom{m-1}{n-1}$ ways of doing it. Thus $\binom{n}{m-n} = \binom{m-1}{n-1}$ and by substituting $m = n + k$, we have

$$\binom{n}{k} = \binom{n+k-1}{n-1} = \binom{n+k-1}{k}$$

5.4 Identities

In this section, we discuss some identities on $\binom{n}{k}$ and argue their proofs using the idea of either double counting or bijections.

Identity 1

$$\binom{n}{k} = \binom{k+1}{n-1}$$

Proof. Let's use the bijection method to prove this. Formally, let's define sets S_1 and S_2 and count their cardinalities independently and then establish a bijection from S_1 to S_2 proving that $|S_1| = |S_2|$.

S_1 : Configuration of $k \star s$ and $n - 1 \mid s$ as described in Sec. 5.2.4. By the bijection defined in it, $|S_1| = \binom{n}{k}$

S_2 : Configuration of $n - 1 \star s$ and $k \mid s$ as described in Sec. 5.2.4. Again, by the bijection defined in it, $|S_2| = \binom{k+1}{n-1}$

Bijection: Let's define a bijection f from S_1 to S_2 . f takes a configuration from S_1 as input and interpret $\star s$ as $\mid s$ and $\mid s$ as $\star s$. Therefore it ends up with a configuration with $n - 1 \star s$ and $k \mid s$ which is a configuration in S_2 . It's easy to observe that f is a bijection.

As f is a bijection from S_1 to S_2 , we have $|S_1| = |S_2|$. This completes the proof \square

Identity 2

$$k \binom{n}{k} = n \binom{n+1}{k-1}$$

Proof. Let's use the method of double counting to prove this.

Question: In how many ways can we construct a non-decreasing sequence $1 \leq a_1 \leq a_2 \leq \dots \leq a_k \leq n$ and mark one element?

Answer 1: By the bijection established in Sec. 5.2.3 we have $\binom{n}{k}$ number of non-decreasing subsequences and for every such subsequence, we can mark any one of the k elements choose. Thus the answer is $k \binom{n}{k}$

Answer 2: Firstly, determine the value in $[n]$ which is to be marked. Let r be this value. Now, consider a non-decreasing subsequence between 1 and $n + 1$ with $k - 1$ elements. Using r and the non-decreasing sequence chosen, we construct a unique non-decreasing sequence between 1 and n of length k with r as marked in the following way:

Let $(b_1, b_2, \dots, b_{k-1})$ with $1 \leq b_1 \leq b_2 \leq \dots \leq b_{k-1} \leq n + 1$ be the chosen sequence,

- Insert marked- r in the right most position so that the resulting sequence is still sorted.
- As long as there's an $n + 1$ in the sequence, remove it and add it as r to the right of marked- r in the sequence

Therefore, number of required sequences

$$\begin{aligned} &= \text{number of ways to choose } r \times \text{number of non-decreasing sequences of length } k-1 \text{ between } 1 \text{ and } n+1 \\ &= n \times \binom{n+1}{k-1} \end{aligned}$$

This completes the proof \square

Exercise 5.5.

Prove the following by combinatorial arguments

$$\binom{\binom{n}{k}}{k} = \sum_{m=1}^n \binom{\binom{m}{k-1}}{k-1}$$

Hint: Look for bijection to number of non-decreasing subsequences

Prove the following by combinatorial arguments

$$\sum_{k=0}^m \binom{\binom{n}{k}}{k} = \binom{\binom{n+1}{m}}{m}$$

Hint: Look for bijection to Voting problem

Prove the following by combinatorial arguments

$$\binom{\binom{n}{k}}{k} = \sum_{m=0}^n \binom{\binom{n}{m}}{m} \binom{\binom{m}{k-m}}{k-m}$$

Catlan Bijections

6.1 Introduction

One of the classic examples to demonstrate the power of bijections is *Catlan numbers*. The Catlan numbers form a sequence of natural numbers that occur in various counting problems and occurs in several seemingly different contexts. Historically, *Euler* is the first person to study them. He was interested in counting the number of ways of dividing a polygon into triangles by drawing non-overlapping diagonals. Catlan numbers got their name from *Eugene Catlan* when he used them to answer the *Parenthesisation problem* which is the following: Consider a sequence $(a_1, a_2, \dots, a_{n+1})$ of $n + 1$ numbers, If we have to perform a binary operations \odot n times among them, how many number of ways are there to parenthesise (or bracket) them using n parenthesis of single type (say '()'). In this lecture, we will see a few equivalent problems to this and then arrive at an explicit expression of Catlan numbers.

6.2 Equivalent Bijections

In this section, we see a few equivalent problems of the *parenthesisation* problem and argue that answer to each of them is also the *catlan number*

Full binary trees If we observe the Parenthesisation problem carefully, we notice that every valid parentheses of those $n + 1$ numbers form a *full binary tree* (a binary tree in which every node have either two children or no children) of $n + 1$ leaves and n internal nodes where leaves represents the numbers a_1, \dots, a_{n+1} and each internal node corresponds to one operation. Therefore, there's an implicit bijection between the set of valid parenthesisations and full binary trees with n internal nodes. Therefore,

$$\text{number of valid parenthesisations of } n+1 \text{ elements} = \text{number of full binary trees with } n \text{ internal nodes} \quad (6.2)$$

Balanced parenthesised strings A balanced parenthesised string of length $2n$ is a string consists of n left brackets '(' and n right brackets ')' in which every prefix of the string has number of left brackets '(' \geq number of right brackets ')'. One can easily observe the bijection from set of balanced paranthesised string to valid parenthesisations of $n + 1$ numbers

Euler's problem Find the number of ways of triangulating a polygon with $n + 2$ edges

Handshaking problem Consider a scenario where $2n$ people are sitting around a table. How many ways they can shake hands with each other without crossing hands. We leave it as an exercise to establish bijections from *Euler's problem* to *Full binary tree problem* and *handshaking problem* to *balanced parenthesised strings problem*.

2: Jayalal says: Todo - Establish bijections from *Euler's problem* to *Full binary tree problem* and *handshaking problem* to *balanced parenthesised strings problem*

6.3 Algebraic Expression

In this section, we are interested in arriving at a concrete expression of the n^{th} *catlan number* (denoted by c_n). Let's solve another problem and then, by establishing a bijection to one of the above problems, we can arrive at an expression for c_n .

6.3.1 Monotone walk on $n \times n$ grid

Suppose we have a grid of size $n \times n$. How many ways are there to go from $(0,0)$ to (n,n) by using only downward edges or right edges. A sample path is represented in Fig. 6.1. We observe that each step can increment the value of exactly one of the co-ordinates by 1. Since we have to move from $(0,0)$ to (n,n) , we have to increase the value of both the co-ordinates by n and n and thus irrespective of the path you take, the length of a path from $(0,0)$ to (n,n) must be of length $n + n = 2n$.

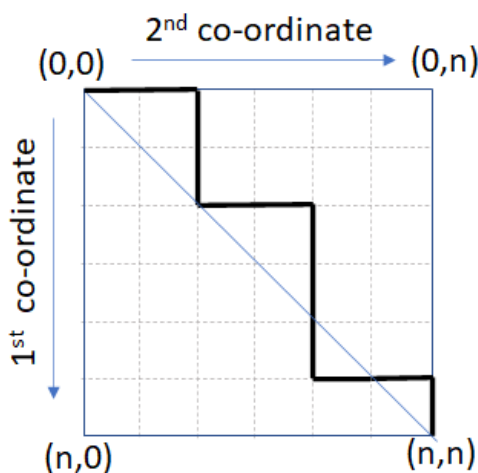


Figure 6.1: A path from $(0,0)$ to (n,n) using downward and right edges

If we represent each right move as R and each downward move as D , one can observe that there's a bijection f from the set of paths to set of strings of length $2n$ over the alphabet $\{D, R\}$ with number of D 's = number of R 's = n . Formally, if $(u_0, v_0), (u_1, v_1), \dots, (u_{2n}, v_{2n})$ represents the path where $(u_0, v_0) = (0,0)$ and $(u_{2n}, v_{2n}) = (n,n)$, and $b = b_1 b_2 \dots b_{2n}$ represents the string

where each b_i is either D or R , our bijection f takes a path as input and sets b_i as

$$b_i = \begin{cases} D & \text{if } u_i = u_{i-1} + 1 \\ R & \text{if } v_i = v_{i-1} + 1 \end{cases}$$

Well defined: As we have exactly n x co-ordinate increments and n y co-ordinate increments, we will have exactly n D 's and n R 's in our string and thus f is well defined.

Injective: Two different paths from $(0, 0)$ to (n, n) will differ in at least one (u_{i-1}, v_{i-1}) to (u_i, v_i) transition where $i = 1, 2, \dots, 2n$, their corresponding strings under f will differ in at least i^{th} position and thus f is injective.

Surjective: Every string over $\{D, R\}$ of length $2n$ with equal number of D 's and R 's has a pre-image under f which is defined by $(u_0, v_0) = (0, 0)$ and (u_i, v_i) is $(u_{i-1} + 1, v_{i-1})$ if $b_i = R$ and $(u_{i-1}, v_{i-1} + 1)$ if $b_i = D$. As there will be n D 's and n R 's, $(u_{2n}, v_{2n}) = (n, n)$ and thus f is surjective.

Thus f is bijection. As we have number of string over $\{D, R\}$ of length $2n$ with equal number of D 's and R 's equal to $\binom{2n}{n}$ (select n positions out of $2n$ available and fill them with D 's and the rest with R 's). Thus the number of paths from $(0, 0)$ to (n, n) with only downward and rightward movements is $\binom{2n}{n}$.

Lets ask a slightly question. How many ways are there to go from $(0, 0)$ to $(n + 1, n - 1)$ using only downward or right edges. Using a similar arguments as above, we can come up with a bijection to set of string over $\{D, R\}$ of length $2n$ with $n + 1$ D 's and $n - 1$ R 's. Therefore number of required paths are $\binom{2n}{n+1} = \binom{2n}{n-1}$

6.3.2 Diagonal avoiding paths and Catalan numbers

In this section we explore the connection between the above paths that we discussed and the Catalan number. Let us ask this question: How many paths are there in the grid from $(0, 0)$ to (n, n) that avoids crossing the diagonal?

We first define what *crossing the diagonal* means. The diagonal consists of the points of the form (i, i) , $i \in \{0, \dots, n\}$. A path $((u_0, v_0), \dots, (u_{2n}, v_{2n}))$ is said to be crossing the diagonal if it *intersects* through the diagonal and goes to some point below the diagonal. Mathematically, a path is a diagonal crossing path if $\exists i$ such that $u_i > v_i$. In particular, $\exists i : u_i = v_i + 1$ (refer fig. 6.2 for example. Any diagonal crossing path must necessarily pass through one of the red dots). Equivalently, in a diagonal avoiding path $\forall i \in \{0, \dots, 2n\}, v_i \geq u_i$. A sample *diagonal-avoiding path* is shown in the fig. 6.3

Before computing this number, an obvious question is what is the connection between such restricted paths and Catalan number. It turns out that the set of diagonal avoiding paths from $(0, 0)$ to (n, n) is in bijection with the set of balanced parenthesized strings of length $2n$. Hence, to count the number of balanced parenthesized strings of length $2n$, which is also the Catalan number, we only need to count the diagonal avoiding paths from $(0, 0)$ to (n, n) . Let us first establish the bijection between the two.

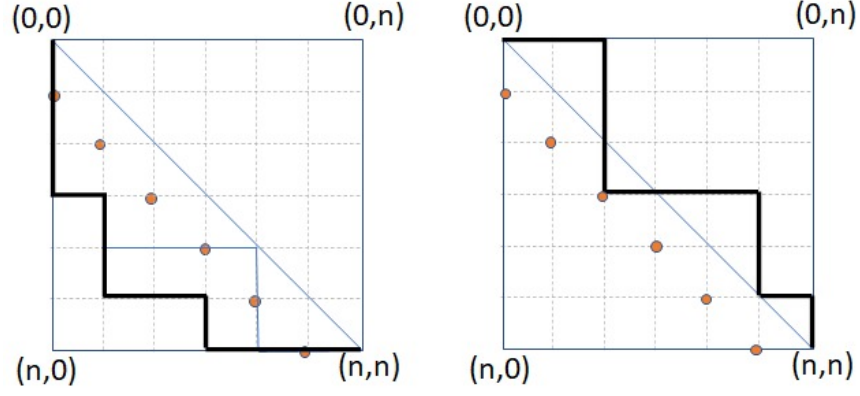


Figure 6.2: Diagonal crossing paths. Note that path in (a) is crossing the diagonal at $(0, 0)$

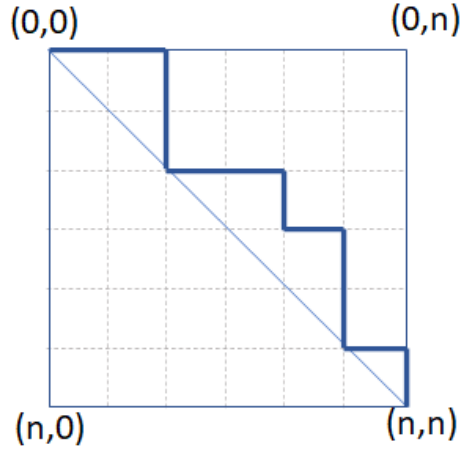


Figure 6.3: A diagonal avoiding path. Observe that it can still touch the diagonal

6.3.3 Bijection from Diagonal avoiding paths to Balanced parenthesis problem

Intuitively, the bijection can be defined as follows: for any given balanced parenthesized string $w = w_1 w_2 \dots w_{2n}$, the corresponding path from $(0, 0)$ to (n, n) is obtained by starting from position $(0, 0)$, and scanning the string from left to right. Take right move whenever '(' is encountered and a down move for ')'. Formally we define the bijection as follows:

Defining the bijection: Let P be the set of diagonal avoiding paths from $(0, 0)$ to (n, n) and B be the set of balanced paranthesized strings of length $2n$ over the alphabets $\{ (,) \}$. Define the bijection $\phi : B \rightarrow P$ as follows:

For $w = w_1 w_2 \dots w_{2n} \in B$, $\phi(w) = (u_0, v_0), (u_1, v_1), \dots, (u_i, v_i), \dots, (u_{2n}, v_{2n})$, where

1. $(u_0, v_0) = (0, 0)$
2. $\forall i \in \{1, 2, \dots, 2n\}$

$$(u_i, v_i) = \begin{cases} (u_{i-1} + 1, v_{i-1}) & \text{if } w_i =) \\ (u_{i-1}, v_{i-1} + 1) & \text{if } w_i = (\end{cases}$$

Proof of bijection

Well-defined: From the above description, given any string w , $\phi(w)$ is uniquely defined. Further, for any string $w \in B$, since the number of '(' is same as the number of ')' = n , the corresponding path has n right and n down moves and hence it ends at (n, n) . Also, since the number of left brackets is greater than or equal to the number of right brackets in any prefix of w , for all $i \in [2n]$, $v_i \geq u_i$. This shows that $\forall w \in B, \phi(w) \in P$. Hence, ϕ is well-defined.

Injective: Let w, w' be two different strings in set B . Then \exists an index $i \in [2n]$ where $w_i \neq w'_i$. Hence $\phi(w)$ and $\phi(w')$ also differ at the i th step, where one of the paths takes one step right while the other takes one step down.

Surjective: Given any path $((0, 0), (u_1, v_1), \dots, (u_{2n}, v_{2n}))$ the corresponding string $w \in B$ is defined as follows:

$\forall i \in [2n]$

$$w_i = \begin{cases} (& \text{if } (u_i, v_i) = (u_{i-1}, v_{i-1} + 1) \\) & \text{if } (u_i, v_i) = (u_{i-1} + 1, v_{i-1}) \end{cases}$$

We can verify that the string w indeed is in set B , because firstly, for any path in P , $\forall i, v_i \geq u_i$ and hence by definition, number of left brackets '(' in w is greater than or equal to number of right brackets, '(' in any prefix of w . Secondly, for any path to reach from $(0, 0)$ to (n, n) it must have n right moves (increase in 2nd coordinate) and n down moves (increase in 1st coordinate) and hence w must have n left brackets and n right brackets.

6.3.4 Counting the number of diagonal avoiding paths

Having established the bijection between Catalan number and diagonal avoiding paths, we get

$$C_n = \# \text{ of diagonal avoiding paths from } (0, 0) \text{ to } (n, n) \quad (6.3)$$

So, our next task is to count the number of diagonal avoiding paths from $(0, 0)$ to (n, n) . To count this, we take following approach. Let us call the diagonal avoiding paths as *good* paths and diagonal crossing paths as *bad* paths. Then,

$$\begin{aligned} \# \text{ of diagonal avoiding paths from } (0, 0) \text{ to } (n, n) &= \# \text{ of paths from } (0, 0) \text{ to } (n, n) - \# \text{ of diagonal crossing paths from } (0, 0) \text{ to } (n, n) \end{aligned} \quad (6.4)$$

So, now our revised goal is to count the number of diagonal crossing paths from $(0, 0)$ to (n, n) . How do we do that? Here again bijection plays an important role. The idea is to translate diagonal crossing paths into different kind of paths which are easy to count.

Let us define the following path translation: Let $\pi = (0, 0), (u_1, v_1), \dots, (u_{2n}, v_{2n})$ be a diagonal crossing path. Then there must exist i such that $u_i = v_i + 1$. There can be many such indices as the path can cross the diagonal multiple times. Choose i to be the least such index. Let $u_i = \ell$,

then the first co-ordinate after crossing the diagonal is $(\ell, \ell - 1)$. Let us call this point P (refer fig. 6.4(a)). Then to find the translated path we reflect the part of the path π after point P w.r.t. the main diagonal.

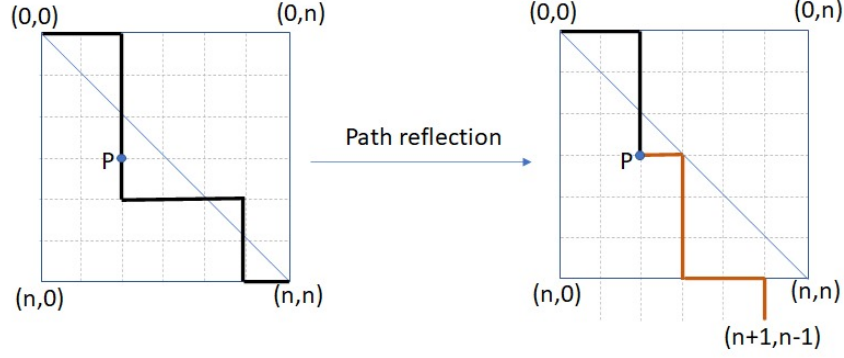


Figure 6.4: Point P in a diagonal crossing path and the reflected path after P

More precisely, we can divide the diagonal crossing path into two stretch S_1, S_2 , where S_1 is the part of the path between $(0, 0)$ to P and S_2 is the part of the path between P to (n, n) . Then to translate π into a new path, replace S_2 with S'_2 to get a new path $\pi' = S_1 S'_2$. The replacement S'_2 is defined as follows:

- replace downward edges with right edges and
- replace right edges with downward edges.

Refer fig. 6.4(b) We can observe that the new path π' described in this way is always between $(0, 0)$ to $(n + 1, n - 1)$. The argument for this goes as follows:

Originally (in S_2), $(\ell, \ell - 1)$ goes to (n, n) which means it takes $(n - \ell)$ downward moves and $(n - \ell + 1)$ right moves. Since, we are swapping the right and downward moves to get S'_2 from S_2 , there are $(n - \ell + 1)$ downward moves and $(n - \ell)$ right moves from point $P = (\ell, \ell - 1)$ in S'_2 . Thus, S'_2 goes from $(\ell, \ell - 1)$ to $(\ell + n - \ell + 1, \ell - 1 + n - \ell) = (n + 1, n - 1)$ and hence, $\pi' = S_1 S'_2$ is a path from $(0, 0)$ to $(n + 1, n - 1)$.

Thus we have established that any diagonal crossing path from $(0, 0)$ to (n, n) maps to a path from $(0, 0)$ to $(n + 1, n - 1)$ after applying the transformation described above. The converse is also true, i.e., given any path from $(0, 0)$ to $(n + 1, n - 1)$, we can translate it back to a diagonal crossing path from $(0, 0)$ to (n, n) by using the same reflection technique. Thus, we get a bijection between the set of diagonal crossing paths from $(0, 0)$ to (n, n) to the set of paths from $(0, 0)$ to $(n + 1, n - 1)$. We formally define the translation and prove that it is indeed a bijection.

Bijection: Let A be the set of diagonal crossing paths from $(0, 0)$ to (n, n) and B be the set of paths from $(0, 0)$ to $(n + 1, n - 1)$. Then the mapping $\phi : A \rightarrow B$ is formally defined as follows:

Let $\pi = (0, 0), (u_1, v_1), \dots, (u_{2n}, v_{2n})$ and (u_i, v_i) be the first point when π crosses the diagonal. Then $\phi(\pi) = \pi' = (0, 0), (u'_1, v'_1), \dots, (u'_{2n}, v'_{2n})$ is given by:

1. $\forall 1 \leq j \leq i, (u'_j, v'_j) = (u_j, v_j)$

2. $\forall i + 1 \leq j \leq 2n,$

$$(u'_j, v'_j) = \begin{cases} (u'_{j-1} + 1, v'_{j-1}) & \text{if } (u_j, v_j) = (u_{j-1}, v'_{j-1} + 1) \\ (u'_{j-1}, v'_{j-1} + 1) & \text{if } (u_j, v_j) = (u_{j-1} + 1, v'_{j-1}) \end{cases}$$

Well-defined: We already observed that any path $\pi \in A$ from $(0, 0)$ to (n, n) maps to a path $(0, 0)$ to $(n + 1, n - 1)$. Hence ϕ is well defined.

Injection: Consider two different diagonal crossing paths π_1 and π_2 . Let $\pi_1 = S_{1,1}S_{1,2}$ and $\pi_2 = S_{2,1}S_{2,2}$, where the two components $S_{i,1}$ and $S_{i,2}$ for $i \in \{1, 2\}$ are as defined before. Then following two cases are possible:

- Case1: $S_{11} \neq S_{21}$. Then $\pi'_1 \neq \pi'_2$, because the first component is copied as it is in the translation, i.e. $\pi'_1 = S_{1,1}S'_{1,2}$ and $\pi'_2 = S_{2,1}S'_{2,2}$.
- Case2: $S_{11} = S_{21}$, but $S_{12} \neq S_{22}$. In this case $S'_{12} \neq S'_{22}$ because of the way it is defined, i.e. for every right move there is a downwards move and vice-versa. Hence, $\pi'_i \neq \pi'_2$.

Surjective: Given any path π' from $(0, 0)$ to $(n + 1, n - 1)$, we can construct the corresponding path π from $(0, 0)$ to (n, n) , such that $\phi(\pi) = \pi'$, as follows.

Let $\pi' = (0, 0), (u'_1, v'_1), \dots, (u'_{2n}, v'_{2n})$. Since π' goes to $(n + 1, n - 1)$ which is below the diagonal there must exist i such that (u'_i, v'_i) is below the diagonal. Again, there can be many such indices. Take i to be the first such index. Same as before, let $\pi' = S'_1 S'_2$, where S'_1 is the path from $(0, 0)$ to (u'_i, v'_i) and S'_2 is the path from (u'_i, v'_i) to (u'_{2n}, v'_{2n}) . Then $\pi = S_1 S_2$ where S_2 is obtained from S'_2 by swapping the right and downwards moves. Mathematically, let $\pi = (0, 0), (u_1, v_1), \dots, (u_{2n}, v_{2n})$. Then

1. $\forall j \leq i, (u_j, v_j) = (u'_j, v'_j)$
2. $\forall i + 1 \leq j \leq 2n$

$$(u_j, v_j) = \begin{cases} (u_{j-1} + 1, v_{j-1}) & \text{if } (u'_j, v'_j) = (u'_{j-1}, v'_{j-1} + 1) \\ (u_{j-1}, v_{j-1} + 1) & \text{if } (u'_j, v'_j) = (u'_{j-1} + 1, v'_{j-1}) \end{cases}$$

Again by the same argument as before it can be verified that π is a diagonal crossing path from $(0, 0)$ to (n, n) . We write it here for completeness. Let $(\ell, \ell - 1)$ be the first point when π' crosses the diagonal. Then since the path from $(0, 0)$ to $(\ell, \ell - 1)$ remains as it is in π , it is a diagonal crossing path. Further since π' is path from $(0, 0)$ to $(n + 1, n - 1)$, it takes $n + 1 - \ell$ downward steps and $n - \ell$ right steps from $(\ell, \ell - 1)$. Hence, π takes $n + 1 - \ell$ right and $n - \ell$ downward steps from $(\ell, \ell - 1)$. Thus, π ends at $(\ell + n - \ell, \ell - 1 + n + 1 - \ell) = (n, n)$.

Thus, we have established a bijection between the set of diagonal crossing paths from $(0, 0)$ to (n, n) and the set of paths from $(0, 0)$ to $(n + 1, n - 1)$. Hence,

$$\begin{aligned} \# \text{of diagonal crossing paths from } (0, 0) \text{ to } (n, n) &= \# \text{of paths from } (0, 0) \text{ to } (n + 1, n - 1) \\ &= \binom{2n}{n + 1} \end{aligned}$$

Hence, from (6.3),(6.4),

$$\begin{aligned}
C_n &= \# \text{of diagonal avoiding paths from } (0,0) \text{ to } (n,n) \\
&= \# \text{ of paths from } (0,0) \text{ to } (n,n) - \# \text{ of diagonal crossing paths from } (0,0) \text{ to } (n,n) \\
&= \binom{2n}{n} - \binom{2n}{n+1} \\
&= \binom{2n}{n} - \frac{n}{n+1} \binom{2n}{n} \\
&= \frac{1}{n+1} \binom{2n}{n}
\end{aligned}$$

Here, in the second last line, we have used the identity:

$$\binom{2n}{n+1} = \frac{n}{n+1} \binom{2n}{n}.$$

Exercise 6.4.

Try to establish a bijection between the set of different possible polygon triangulation in a polygon of $n + 2$ nodes and the set of binary trees with n internal nodes.

Hint: associate each internal node with a triangle in a triangulation. Then, each internal node will have degree three, which is the case for full binary tree, except for the leaves. Leaves will correspond to those triangles whose one of the edge is the boundary of the polygon.

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Lecture 7

From Bijections to PIE

7.1 Introduction

In this lecture, we will continue with the use of bijections and use it in formally proving the two identities that we discussed in class and then see their relationship to the Principal of Inclusion and Exclusion.

7.2 The Identities

Recall that we proved following two identities in one of the discussion sessions

$$\sum_{k=0}^n (-1)^k \binom{n}{k} = 0 \quad (7.5)$$

$$\sum_{k=0}^m (-1)^k \binom{n}{k} = (-1)^m \binom{n-1}{m} \quad (7.6)$$

In this section, we will see the proofs for the above equations in detail

7.2.1 Proof for Eqn. (7.5)

$$\sum_{k=0}^n (-1)^k \binom{n}{k} = 0$$

Proof. The LHS counts the number of even sized subsets of $[n]$ with positive sign and odd size subsets with negative sign. Then we proved the result using bijection between even sized and odd sized subsets of $[n]$. Hence, we get 0 on RHS. Let us formally define the bijection here.

Let E be the set of all even sized subsets of $[n]$ and O be the set of all odd sized subsets of $[n]$. Then the bijection $\phi_i : E \rightarrow O$ is defined with respect to an element $i \in [n]$ as follows.

Let $X \subseteq [n]$, such that $|X|$ is even. Then

$$\phi_i(X) = \begin{cases} X \setminus \{i\} & \text{if } i \in X \\ X \cup \{i\} & \text{if } i \notin X \end{cases}$$

Proof of bijection:

Well-defined: Given any even sized subset X , there are two possibilities: (i) $i \in X$, (ii) $i \notin X$. In first case, i is removed from X , hence its size reduces by one and becomes odd. In the second case, i is added, hence the size of the subset increases by one and becomes odd. Hence, ϕ is well defined.

Injective: Let X and X' be two distinct subsets of $[n]$. Then $\exists j \in [n]$ such that j is present in exactly one of the two subsets. Wlog, let $j \in X$ and $j \notin X'$. Now, if $j \neq i$, then $j \in \phi(X)$ and $j \notin \phi(X')$ and hence $\phi(X) \neq \phi(X')$. On the other hand, if $j = i$, then $j \notin \phi(X)$ and $j \in \phi(X')$. Hence, $\phi(X) \neq \phi(X')$.

Surjective: Let $Y \in \mathcal{O}$ be an odd sized subset of $[n]$. From Y , we can recover X such that $\phi(X) = Y$ by the same operation as in ϕ . That is,

$$X = \begin{cases} Y \setminus \{i\} & \text{if } i \in Y \\ Y \cup \{i\} & \text{if } i \notin Y \end{cases}$$

It can easily be verified that in both the cases, X is an even sized subset of $[n]$.

This completes the proof. □

7.2.2 Proof for Eqn. (7.6)

$$\sum_{k=0}^m (-1)^k \binom{n}{k} = \binom{n-1}{m}$$

Proof. Now we look at the second identity which is even more interesting. To prove this identity we use *almost bijection* where the bijection is between a set and subset of another set.

In words, the identity to prove, can be described as

$$\begin{array}{c} \# \text{ of even sized subsets of } [n] \\ \text{of size atmost } m \end{array} - \begin{array}{c} \# \text{ of odd sized subsets of } [n] \\ \text{of size atmost } m \end{array} = (-1)^m \binom{n-1}{m}.$$

Clearly, there cannot be a bijection between the two sets (even sized subsets and odd sized subsets) in this case, since their difference is non-zero. This is where we use almost bijection.

We use following case analysis.

Case1: m is even: Then the identity to prove is:

$$\sum_{k=0}^m (-1)^k \binom{n}{k} = \binom{n-1}{m} \tag{7.7}$$

This can be interpreted as

$$\sum_{\substack{k=0, \\ k \text{ is even}}}^m \binom{n}{k} - \sum_{\substack{k=1, \\ k \text{ is odd}}}^{m-1} \binom{n}{k} = \binom{n-1}{m} \tag{7.8}$$

Let E be the set of all the even sized subsets of $[n]$ of size at most m and O be the set of odd sized subsets of $[n]$ having size at most $m - 1$. Then, Eqn. (7.8) can intuitively interpreted as follows: there is a subset $E' \subseteq E$, such that E' is in bijection with O and $|E \setminus E'| = \binom{n-1}{m}$. Thus, we have three tasks at hand

- identify the set E' , and
- define and prove the bijection between E' and O .
- prove that $|E \setminus E'| = \binom{n-1}{m}$

Defining the set E' : Set E' is the union of two sets:

$$E' = \{X \subseteq [n] : |X| \text{ is even and } |X| \leq m - 2\} \cup \{X \subseteq [n] : i \in X \text{ and } |X| = m\}$$

Defining the bijection: The bijection $\phi : E' \rightarrow B$ is defined in the same way as we defined it for first identity. That is, for $X \in E'$,

$$\phi(X) = \begin{cases} X \setminus \{i\} & \text{if } i \in X \\ X \cup \{i\} & \text{if } i \notin X \end{cases}$$

Proof of bijection

Well-defined: Let $X \in E'$, then (i) if $|X| \leq m - 2$, then $|\phi(X)|$ is odd and $|\phi(X)| \leq m - 1$, (ii) if $|X| = m$, then $i \in X$, hence $\phi(X) = X \setminus \{i\}$. This implies $|\phi(X)| = m - 1$. Thus, in both the cases $\phi(X) \in O$.

Injective: Since, the function is same as in the previous case, the same argument for injectivity works.

Surjective: Let $Y \in O$ be an odd sized subset of $[n]$. From Y , we can recover $X \in E'$ such that $\phi(X) = Y$ by the same operation as in ϕ . That is,

$$X = \begin{cases} Y \setminus \{i\} & \text{if } i \in Y \\ Y \cup \{i\} & \text{if } i \notin Y \end{cases}$$

It can easily be verified that in both the cases, $|X|$ is even. In first case, since $|Y| \leq m - 1$, $|X| \leq m - 2$, hence $X \in E'$. In second case, since $i \notin Y$ and $|Y| \leq m - 1$, $|X| \leq m$ and $i \in X$. Hence $X \in E'$, by definition.

This proves the bijection between E' and O .

Proof for: $|E \setminus E'| = \binom{n-1}{m}$

From the above definitions, $E \setminus E' = \{X \subseteq [n] : |X| = m, i \notin X\}$. This can be interpreted as $E \setminus E' = \{X \subseteq [n] \setminus \{i\} : |X| = m\}$. Hence, $|E \setminus E'| = \binom{n-1}{m}$.

Case2: m is odd: In this case the identity to prove is:

$$\sum_{k=0}^m (-1)^k \binom{n}{k} = -\binom{n-1}{m} \quad (7.9)$$

This can be interpreted as

$$\sum_{\substack{k=0, \\ k \text{ is even}}}^{m-1} \binom{n}{k} - \sum_{\substack{k=1, \\ k \text{ is odd}}}^m \binom{n}{k} = -\binom{n-1}{m} \quad (7.10)$$

Equivalently,

$$\sum_{\substack{k=1, \\ k \text{ is odd}}}^m \binom{n}{k} - \sum_{\substack{k=0, \\ k \text{ is even}}}^{m-1} \binom{n}{k} = \binom{n-1}{m} \quad (7.11)$$

This time the set of odd sized subsets of $[n]$ of size at most m is bigger than the even sized subsets of $[n]$ of size at most m . The proof is same as that for the case of even m . Let E be the set of all the even sized subsets of $[n]$ of size at most $m-1$ (since m is odd) and O be the set of odd sized subsets of $[n]$ having size at most m . Then (7.11) can be interpreted as follows: there is a subset $O' \subseteq O$, such that E is in bijection with O' and $|O \setminus O'| = \binom{n-1}{m}$.

Thus, we have two task at hand

- identify the set O' , and
- define and prove the bijection between E and O' .
- prove that $|O \setminus O'| = \binom{n-1}{m}$

Defining the set O' : Set O' to be the union of two sets:

$$O' = \{Y \subseteq [n] : |Y| \text{ is odd and } |Y| \leq m-2\} \cup \{Y \subseteq [n] : i \in Y \text{ and } |Y| = m\}$$

Defining the bijection: The bijection $\phi : E \rightarrow O'$ is defined in the same way as we defined it for first identity. That is, for $X \in E$,

$$\phi(X) = \begin{cases} X \setminus \{i\} & \text{if } i \in X \\ X \cup \{i\} & \text{if } i \notin X \end{cases}$$

Proof of bijection

Well-defined: Let $X \in E$, then $\phi(X)$ is of odd size because either an element is added or removed from X , which is of even size. Now, (i) if $i \in X$, then $\phi(X) = X \setminus \{i\}$. Hence, $|\phi(X)| \leq m-2$ (because $|X| \leq m-1$) which implies $\phi(X) \in O'$ (ii) if $i \notin X$, then, $\phi(X) = X \cup \{i\}$. This implies $|\phi(X)| \leq m$. But since, $i \in \phi(X)$, $\phi(X) \in O'$. This proves that ϕ is well-defined.

Injective: Since, the function is same as in sub section 7.2.1, the same argument for injectivity works.

Surjective: Let $Y \in O'$ be an odd sized subset of $[n]$. From Y , we can recover $X \in E$ such that $\phi(X) = Y$ by the same operation as in ϕ . That is,

$$X = \begin{cases} Y \setminus \{i\} & \text{if } i \in Y \\ Y \cup \{i\} & \text{if } i \notin Y \end{cases}$$

It can easily be verified that in both the cases, $|X|$ is even. In first case, $|Y| \leq m$ and hence $|X| \leq m - 1$. So, $X \in E$. In second case, since $i \notin Y$, $|Y| \leq m - 2$ (by definition) and hence $|X| \leq m - 1$. Hence $X \in E$.

This proves the bijection between E and O' .

Proof for: $|O \setminus O'| = \binom{n-1}{m}$

From the above definitions, $O \setminus O' = \{Y \subseteq [n] : |Y| = m, i \notin Y\}$. This can be interpreted as $O \setminus O' = \{Y \subseteq [n] \setminus \{i\} : |Y| = m\}$. Hence, $|O \setminus O'| = \binom{n-1}{m}$.

This completes the proof □

This proves both the identities.

7.3 Principle of Inclusion and Exclusion

Suppose we are given n sets $A_1, A_2, \dots, A_n \subseteq G$, where G is some ground set. We are interested in finding the size of $A = A_1 \cup A_2 \cup \dots \cup A_n$. This is very abstract scenario and we will see specific examples later, but here we are going to see classic use of the above identities in deriving this number.

So, we are interested in finding $|A| = |A_1 \cup A_2 \cup \dots \cup A_n|$.

So, here is a thought process - Clearly, we can add the size of individual sets as $|A| = |A_1| + |A_2| + \dots + |A_n|$, but this will over-count if there are some elements present in more than one sets. So, for that we need to subtract the double counting. For e.g. if $x \in A_1$ and $x \in A_2$, then it gets counted twice and to compensate for that we need to subtract $|A| = |A_1 \cap A_2|$ and we might attempt $|A| = |A_1| + |A_2| + \dots + |A_n| - \sum_{1 \leq i < j \leq n} |A_i \cap A_j|$. But then, if x is present in A_1, A_2 and A_3 , then it is under-counted (added thrice and subtracted thrice). So, again we need to compensate for that by adding $\sum_{1 \leq i < j < k \leq n} |A_i \cap A_j \cap A_k|$ in the above expression and this sequence goes on for any element being present in $k \leq n$ sets and finally we get the expression for $|A|$ as follows

$$|A| = |A_1| + \dots + |A_n| - \sum_{1 \leq i < j \leq n} |A_i \cap A_j| + \sum_{1 \leq i < j < k \leq n} |A_i \cap A_j \cap A_k| - \dots + (-1)^{n+1} |A_1 \cap A_2 \cap \dots \cap A_n| \quad (7.12)$$

For $n = 2$, the above expression gives

$$|A| = |A_1| + |A_2| - |A_1 \cap A_2|$$

which we all must have seen before and can easily prove using Venn diagram.

In this section, we will formally prove the above expression for general n using the two identities we proved in previous section.

Proof. Consider any $x \in A_1 \cup A_2 \cup \dots \cup A_n$. Let x appears in k of the A_i 's. Then let us see how x gets counted

- $|A_1| + |A_2| + \dots + |A_n|$: counts x k times (added)
- $\sum_{1 \leq i < j \leq n} |A_i \cap A_j|$: counts x $\binom{k}{2}$ times (subtracted)

- $\sum_{1 \leq i < j < k \leq n} |A_i \cap A_j \cap A_k|$: counts x $\binom{k}{3}$ times (added)
- and so on ...

Notice that in terms involving intersection of more than k sets, x never appears.

Thus,

$$\begin{aligned}
 \text{\#of times } x \text{ gets counted} &= k + \binom{k}{2} - \binom{k}{3} + \cdots + (-1)^{k+1} \binom{k}{k} \\
 &= -\binom{k}{0} + \binom{k}{1} + \binom{k}{2} - \binom{k}{3} + \cdots + (-1)^{k+1} \binom{k}{k} + \binom{k}{0} \\
 &= -\sum_{i=0}^k (-1)^i \binom{k}{i} + \binom{k}{0} \\
 &= \binom{k}{0} \quad \text{from (7.5)} \\
 &= 1
 \end{aligned}$$

Thus, irrespective of the value of k , any element $x \in A_1 \cup A_2 \cup \cdots \cup A_n$ is counted exactly once. Hence, every $x \in A_1 \cup A_2 \cup \cdots \cup A_n$ is counted exactly once in RHS in (7.12).

This proves the PIE □

Now let us look at the application of second identity that we derived. This identity is used in deriving a version of PIE which appears very naturally in several context. Let us look at one such example.

PIE says that if we want to derive $|A_1 \cup A_2 \cup \cdots \cup A_n|$, then the following expression does not give the correct count.

$$|A_1 \cup A_2 \cup \cdots \cup A_n| = |A_1| + |A_2| + \cdots + |A_n|$$

But we can ask, does this expression gives a lower or an upper bound? As we saw, this does over-counting, hence we can write

$$|A_1 \cup A_2 \cup \cdots \cup A_n| \leq |A_1| + |A_2| + \cdots + |A_n|$$

Now, suppose we include the next component, i.e.

$$|A_1| + |A_2| + \cdots + |A_n| - \sum_{1 \leq i < j \leq n} |A_i \cap A_j|$$

Again from PIE we know that this also does not give the correct count. But we ask the same question again - does it give any lower or upper bound. And as we saw that this term can do some over-subtraction and hence we can say that this expression gives the lower bound. That is,

$$|A_1 \cup A_2 \cup \cdots \cup A_n| \geq |A_1| + |A_2| + \cdots + |A_n| - \sum_{1 \leq i < j \leq n} |A_i \cap A_j|$$

Similarly,

$$|A_1 \cup A_2 \cup \dots \cup A_n| \leq |A_1| + |A_2| + \dots + |A_n| - \sum_{1 \leq i < j \leq n} |A_i \cap A_j| + \sum_{1 \leq i < j < k \leq n} |A_i \cap A_j \cap A_k|$$

and we continue like this.

Let us now formally establish this observation. We use the same technique that we used in the proof of PIE.

Let x appears in k of the sets in A_1, A_2, \dots, A_n . Suppose we cut off the PIE after $m \leq n$ sized intersections. Then

$$\begin{aligned} \text{\#of times } x \text{ gets counted} &= \binom{k}{1} - \binom{k}{2} + \dots + (-1)^{m+1} \binom{k}{m} \\ &= - \sum_{i=0}^m (-1)^i \binom{k}{i} + \binom{k}{0} \\ &= 1 + (-1)^{m+1} \binom{k-1}{m} \quad \text{from (7.6)} \end{aligned}$$

Thus, x is over counted or under counted depending on whether the second term on RHS is positive or negative. Let us analyze this for two cases.

Case1: $k \leq m$

Since, x appears in only $k \leq m$ sets and we are cutting down only after m , then this means that all possible intersections of this particular x are added and subtracted and x can not appear in any of the intersections of more than k sets. Hence, x is neither under counted nor over counted. In the expression, $\binom{k-1}{m} = 0$ Hence,

$$\text{\#of times } x \text{ is counted} = 1$$

Case2: $k > m$

In this case, x can be under counted or over counted depending upon whether m is even or odd. If m is odd then x is over counted.

If m is even then x is under counted.

Notice that either all $x \in A_1 \cup A_2 \cup \dots \cup A_n$ are correctly counted or under counted or all x are correctly counted or over counted based on the parity of m . Thus, whether a PIE cut down after m intersections gives lower bound or upper bound depends only on the parity of m . This principle is also called the *Bon Ferroni's inequality*.

REMARK 7.3.1. We used the equality in (7.7) to prove PIE. We can actually do the other way round as well, i.e. we can use PIE to prove this equality too.

This completes this lecture. In the next lecture we will look at some applications of PIE.

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Lecture 8

Live lecture discussion

Bijection from Euler's problem to Binary Trees As we have already established a bijection from set of balanced parenthesisations to set of full binary trees and established that number of full binary trees with n internal nodes is the catlan number C_n , in this section, let's establish a bijection from the *Euler's Problem* to set of full binary trees to establish that the solution to *Euler's problem* is also catlan number C_n .

Lets recall *Euler's problem* first. Consider a convex polygon with $n + 2$ edges. Euler's problem is the number of ways of triangulating it (partition the polygon into triangles) by drawing non-crossing diagonals. (Refer fig. 8.5). We know that number of non-crossing diagonals in a polygon of $n + 2$ edges is $n - 1$ (proof follows from a simple induction) and from those $n - 1$ non-crossing diagonals, we have our polygon partitioned into n triangles. Let's associate each of the triangles with a vertex (green dots in the fig. 8.5). Observe that if two triangles share an edge, it must be one of the diagonals (no two triangles can share an edge because of non-crossing diagonals). Now, let's connect the vertices whose corresponding triangles share an edge. Any edge connecting two of these vertices crosses a diagonal. Now, consider a polygon edge e . For every polygon edge

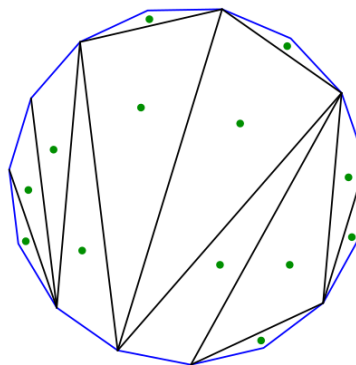


Figure 8.5: Partitioning a polygon into triangles by non-crossing diagonals. Observe that green dots in each triangle associates the triangle with a vertex

surrounding a vertex (other than e), add an open-edge originating from that vertex (see fig. 8.6). We arrive at the following claim.

CLAIM 8.0.1. *If we remove the underlying triangles (which are formed with polygon edges and diagonals), from fig. 8.6, the resulting graph obtained (see fig. 8.7) is a full binary tree with the vertices as internal nodes.*

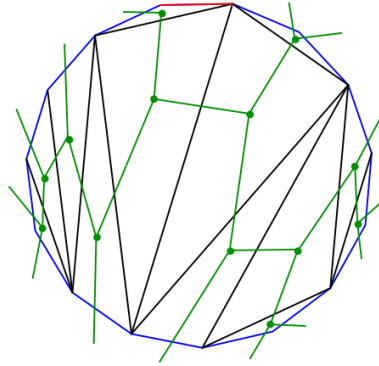


Figure 8.6: Polygon with vertices connected to form a tree

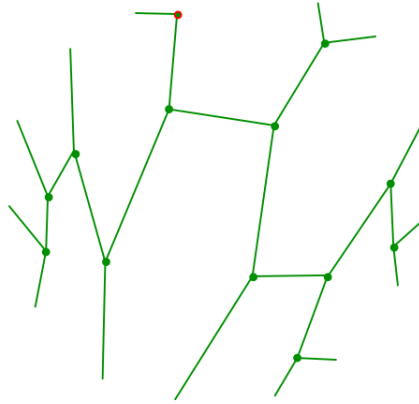


Figure 8.7: Tree formed by connecting vertices

Proof. We observe that degree of every vertex other than the vertex surrounded by edge e is 2. This vertex will act as root to our full binary tree. All other vertices have degree 3 because each vertex is surrounded by a triangle and if a side is a diagonal, it will be connected to vertex which is surrounded by triangle that shares the diagonal and if the side is a polygon edge, then there will be an open edge corresponding to it originating from the vertex. Therefore the resulting graph formed is a full binary tree with our vertices as n internal nodes and vertices corresponding to open edges are $n + 1$ leaves (because there are $n + 2$ edges and one edge is under consideration). This completes the description of bijection. \square

We leave it as an exercise to the reader to prove that the mapping defined above is indeed a bijection.

Bijection from binary trees to full binary trees In this section we are interested in connection between binary and full binary trees. Recall that a full binary tree is one in which each node has either 0 or two children. On the other hand, when we say binary tree then it only means that each node can have at most two children. We want to find a bijection between set of binary trees with n internal nodes and set of full binary trees with certain number of internal nodes.

First of all lets try to see how to convert a given binary tree into a full binary tree so that we can reverse the process, i.e. recover the original (binary) tree back from the full binary tree without ambiguity.

Here is the first attempt:

Attempt 1: First natural approach can be to add a leaf node to all non-full (internal nodes having only one child) nodes, as shown in figure 8.8

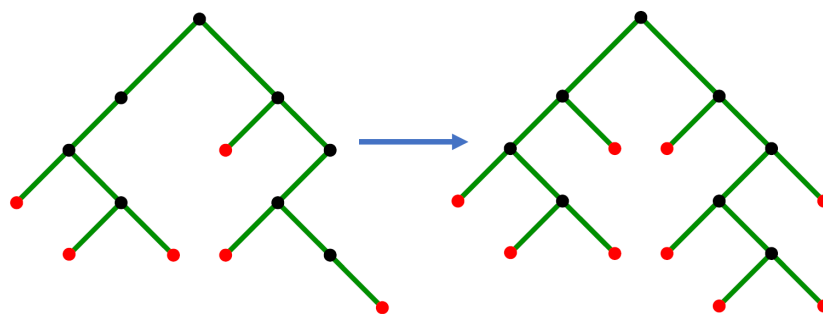


Figure 8.8: Binary to full binary tree attempt1: adding a child node to each non full node

But notice that this transformation is not injective. For example, it can be observed that both the trees in figure 8.9 map to same full binary tree.

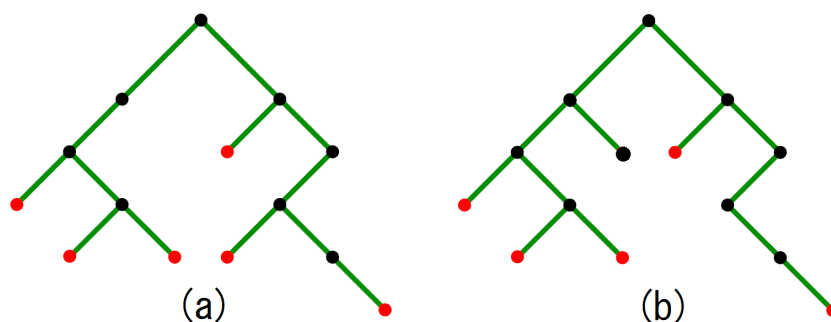


Figure 8.9: Two different binary trees that map to same full binary tree

Attempt 2(correct) Lets try a slightly different approach. Given a binary tree, do the following:

- to each leaf node, add two children
- to each internal node having only one child, add another child

Figure 8.10 shows the full binary tree constructed in this way for the same binary tree as in Figure 8.8. We can see that this solution addresses the issue in the first attempt. Intuitively because of

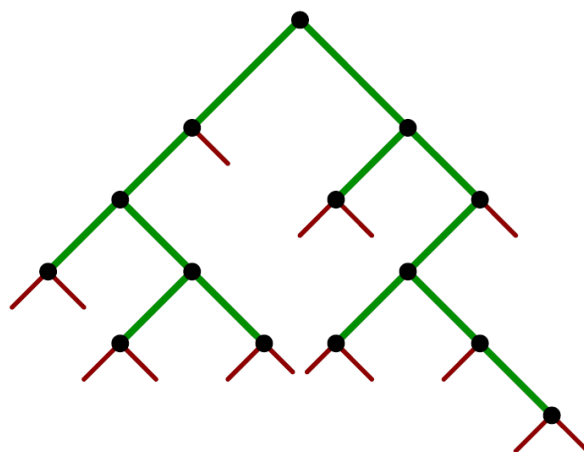


Figure 8.10: Full binary tree for the (non-full) binary tree given in fig 8.8. Notice that all the leaf nodes are added during transformation

following argument: in the previous attempt the problem was that given a full binary tree, it was hard to decide if a leaf node was originally present in the binary tree or added during transformation. Now, in the current solution, this issue does not arise, because for any leaf node originally present in the binary tree, we add two new leaves as its children. Thus, it can be observed that all the leaf nodes (and only these nodes) are added during transformation.

To see that this translation is well-defined, we can see that the transformed tree is full binary tree by construction itself. Surjectivity is also easy to prove. To recover a binary tree from any given full binary tree, simply remove all the leaf nodes. We discussed injection informally. To give a formal argument, we first need to identify how to characterize two different binary trees? One of the hint as given during the discussion is to assign address to the nodes in the form of binary string, where 0-1 represents left or right child.

Here we argued the bijection only intuitively and there are many things to be worked out formally. For example, proof for injection is not formally argued. Also, to argue surjection, we need to fix the number of nodes in full binary tree. Once we figure out this number, the argument for transformation being well-defined also need to take that into account.

Writing a complete formal proof of bijection is left as homework exercise.

Bijection between plane trees and full binary trees A plane tree is a rooted tree with an ordering among the children. A plane tree can have more than two children. Figure 8.11 shows a plane tree.

We are interested in studying the connection between plane trees and binary trees. The number of plane trees with n nodes is equal to the number of binary trees with n nodes. Thus, there is bijection between set of plane trees with n nodes and the set of binary trees with n nodes.

Here we define the bijection function.

The Bijection: Given any plane tree, do the following

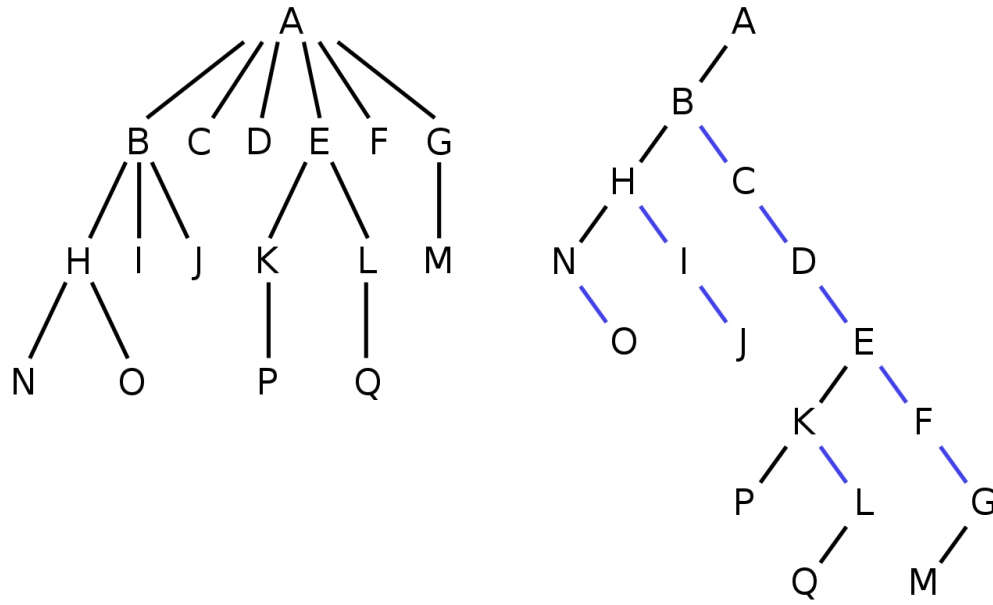


Figure 8.11: An example of plane trees and its transformation to a binary tree

- For each node in the tree,
 - add its first child in plane tree as its left child in binary tree
 - add its immediate sibling on right as its right child in binary tree.

child in the binary tree.

By following the above rule, we get a binary tree from given plane tree.

Observe that in the binary tree thus obtained, root node has only one child, while in general, in a binary tree the root can have both its children. Hence, we won't include the root as part of the binary tree.

Writing formal argument for all the properties is left as homework exercise.