

# Coroutine & Event Dispatcher

## Non blocking operations in LibXively

Olgierd Humeńczuk

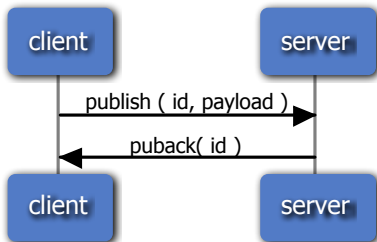
December 13, 2014

# Brief Outline

- 1 Example
  - Sample protocol
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- 2 Coroutines?
  - How does it work?
  - How it can be converted?
  - Coroutine code analysis
- 3 LibXively Coroutines
  - MQTT protocol sending analysis
- 4 Event dispatcher
  - Dispatcher Pros & Cons
  - LibXively dispatcher responsibilities and flow
  - Usage of LibXively event dispatcher on receiving example
- 5 The End
  - Thank you!

# Protocol

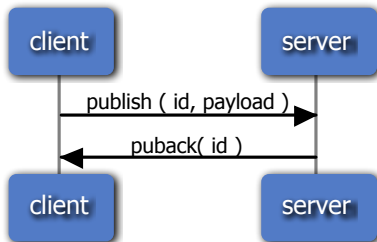
## MuP ( Micro Protocol )



- Let's create simple protocol
- Simple enough to send payloads
- But powerfull enough to make sure it's delivered at least once
- Yes it's a small fraction of MQTT

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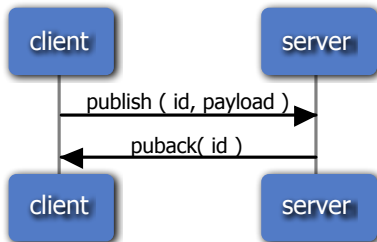
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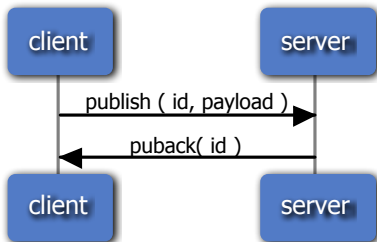
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# Definition

## Header

```
typedef enum
{
    M_UNKNOWN = 0, M_PUBLISH, M_PUBACK
} M_message_types;

typedef struct
{
    M_message_types type;
} M_header_t;
```

## Publish/Puback

```
typedef struct
{
    uint16_t id;
    uint8_t* data; uint16_t length;
}M_publish_t;

typedef struct
{
    uint16_t id;
}M_puback_t;
```

## Final Structure

```
typedef struct
{
    M_header_t header;
    union
    {
        M_publish_t publish;
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} M_protocol_t;
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# Sending

## Common

```
void handle_events( buffer_t* data, event_t e, uint32_t time_diff )
{
    static state_t state = STATE_SEND_DATA;
    static uint16_t sent = 0;
    static uint32_t timer_value = 0;
    static bool parser_not_done = false;

    timer_value += time_diff;

    switch( state )
    {
        case STATE_INIT:
        {
            state          = STATE_SEND_DATA;
            sent            = 0;
            parser_not_done = false;
        }
        break;
    }
```

# Sending

## Sending

```
case STATE_SEND_DATA:
{
    if( e == CAN_WRITE )
    {
        while( sent < data->length )
        {
            int res = write( data->data + sent, data->length - sent );

            if( res <= 0 )
            {
                // do the error handling and maybe leave
            }

            sent += res;
        }

        state      = STATE_RECV_DATA;
        timer_value = 0;
    }
}
break;
```

# Receiving

## Receiving

```
case STATE_RECV_DATA:
{
    if( recvd_msg == NULL ) { ALLOC_AT( M_protocol_t, recvd_msg ); }

    if( e == CAN_READ )
    {
        if( timer_value > TIMEOUT ) { // timeout!!! }

        while( parser_not_done == true )
        {
            ALLOC( buffer_t, recv_data );
            ALLOC_AT( recv_data->data, 32 );

            int res = read( recv_data->data, 32 );
            if( res <= 0 ) { // do the error handling }

            recv_data->length          = res;
            parser_result_t parser_result = call_parser( recv_data
                , recvd_msg
                , &parser_not_done );

            if( parser_result == PARSER_ERROR ) { // do the error handling }
            FREE( recv_data->data );
            FREE( recv_data );
        }

        state = ANALYSE;
    }
    break;
}
```

# Receiving

## Analyse

```
case ANALYSE:
{
    if( recvd_msg->header.type != M_PUBACK ) { // wrong message type }

    // take the payload

    // send it or read it

    FREE( recvd_msg );
}
break;
};
}
```

# Conclusions

- Code is long and less readable
- This is error prone
  - Memory leaks
  - Reordering sequence in wrong way
  - Level of complication
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How does it work?

# Implementation

## Coroutine macros

```
#define BEGIN_CORO( state )\
    switch( state )\
    {\
        default:\

#define YIELD( state, ret )\
    state = __LINE__; return ret; case __LINE__:\

#define YIELD_ON( state, expression, ret )\
{\
    if ( (expression) ) \
    {\
        state = __LINE__; return ret; case __LINE__;; \
    } \
};

#define YIELD_UNTIL( state, expression, ret )\
{\
    if ( (expression) ) \
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    } \
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#define CORO_END()\
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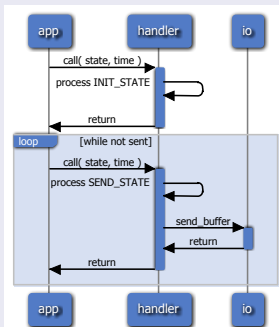
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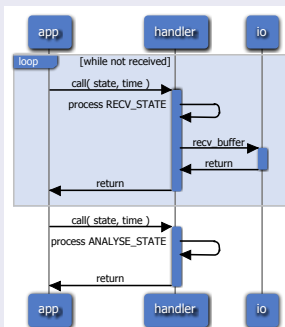
How does it work?

# Sequence of operations for sending/receiving MuP

## Init and send



## Receive and analyse



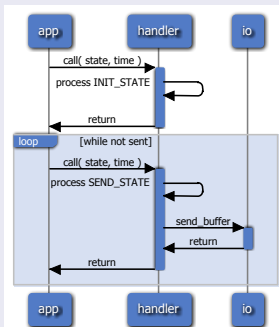
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- We prereserve the state of processing between each execution
- We enter to the proper place using switch



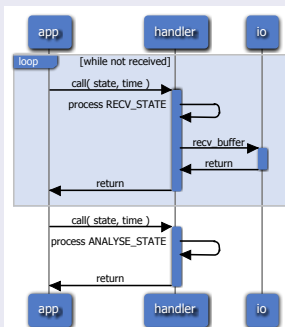
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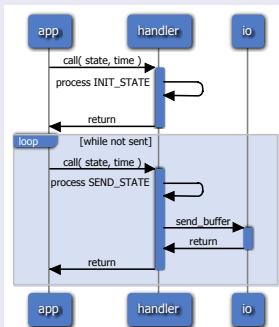


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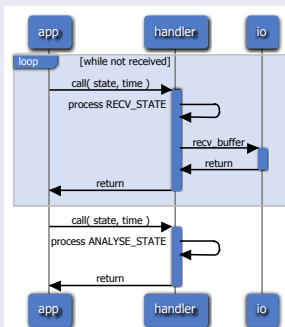
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How it can be converted?

# Solution ?

## Coroutine

```
requested_state_t process( handler_state_t state )
{
    BEGIN_CORO( state->co );

    ALLOC_AT( message_t, state->data );
    ALLOC_BUFFER_AT( message_t, state->data->data );

    do
    {
        res = write( state->data + state->offset, 32 );
        if( res <= 0 ) { // do the error handling and maybe YIELD( state->co ); }
        YIELD_UNTIL( state->co, state->sent < data->length, WANT_WRITE )
    } while( state->sent < data->length );

    FREE( state->data->data ); FREE( state->data );

    state->timeout = REGISTER_TIMER( make_handle( process, TIMEOUT_STATE ), TIMEOUT );
    YIELD( state->co, WANT_READ );

    ALLOC_AT( buffer_t, state->buffer ); ALLOC_BUFFER_AT( state->buffer->data, 32 );
    ALLOC_AT( message_t, state->recvd_msg );

    if( state == TIMEOUT ) { //timeout }
    UNREGISTER_TIMER( state->timeout );

    do
    {
        res = read( state->buffer->data, 32 );
        if( res <= 0 ) { // do the error handling and maybe YIELD( state->co ); }
        state->buffer->length = res;
        state->parser_done = parse( state->buffer, state->recvd_msg );
        YIELD_UNTIL( state->co, state->parser_done, WANT_READ );
    } while( state->parser_done );

    FREE( state->buffer->data ); FREE( state->buffer );

    if( recvd_msg->header.type != M_PUBACK ) { // wrong message type }
    // take the payload
    // send it or read it
    FREE( recvd_msg );
    END_CORO();
}
```

# Coroutine analysis - sending

## sending

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        YIELD_UNTIL( state->co
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                    , WANT_WRITE )
    } while( state->sent < data->length );
}
```

# Coroutine analysis - timeout

## timeout

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state->timeout = REGISTER_TIMER(
    make_handle( process, TIMEOUT_STATE )
    , TIMEOUT );
YIELD( state->co, WANT_READ );
```

# Coroutine analysis - receiving

## receiving

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do
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} while( state->parser_done );
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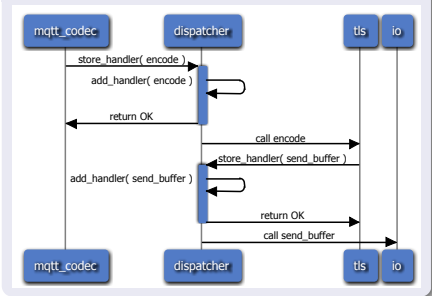
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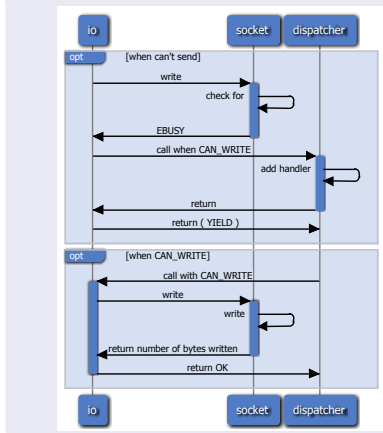
MQTT protocol sending analysis

# MQTT sending with coroutines

## protocol logic level

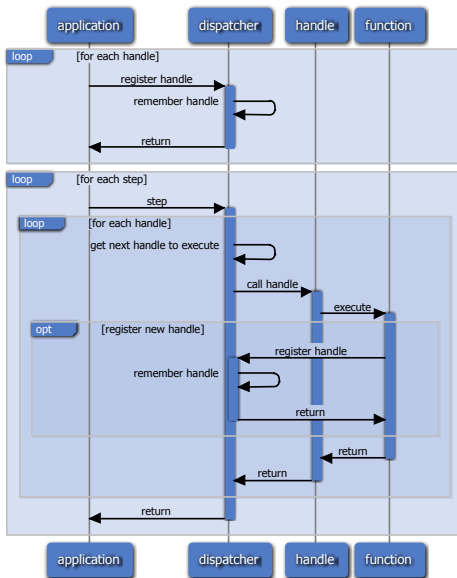


## io logic level





# Reactor pattern sequence diagram



## Dispatcher Pros &amp; Cons

## ● Pros

- Separates application specific code from the dispatcher implementation
- Increases code reusability and allows the code to be modular and well splitted
- Allows some simple coarse-grain concurrency while not adding the complexity of multiple threads to the system

## ● Cons

- The reactor pattern can be more difficult to debug due to the inverted flow of control
- By only calling request handlers synchronously, the reactor pattern limits maximum concurrency ( can be replaced with proactor that does not have that limitation )

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- Execute handle at the desired time
- Execute handle in a desired order
- Execute handle whenever desired event appear on a socket

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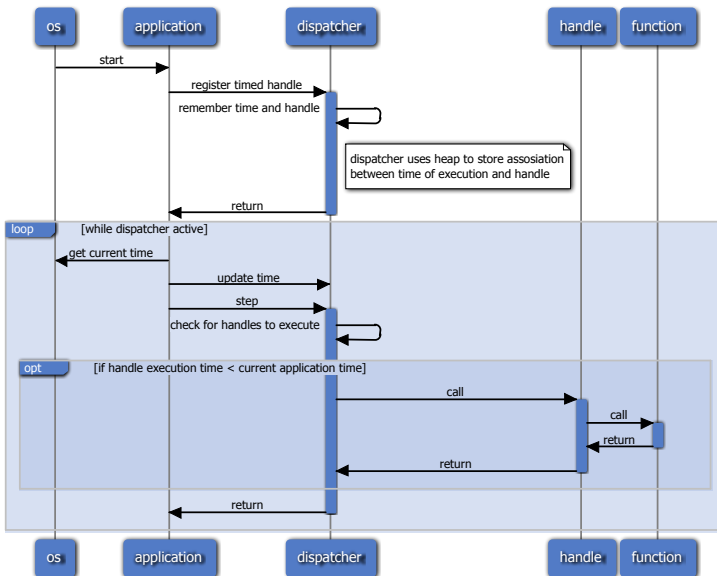
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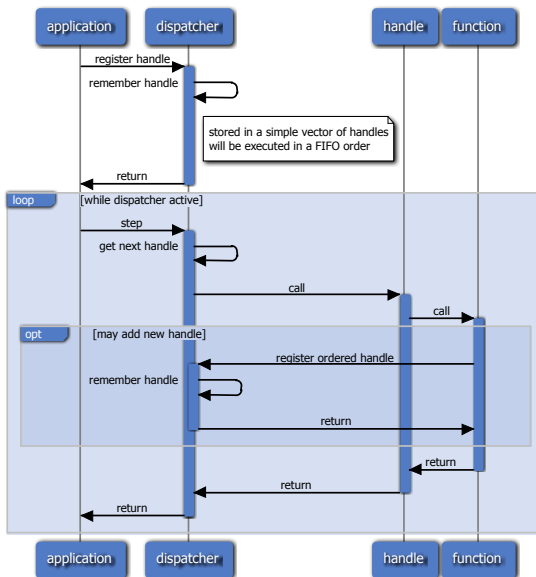
## LibXively dispatcher responsibilities and flow

## Timed events



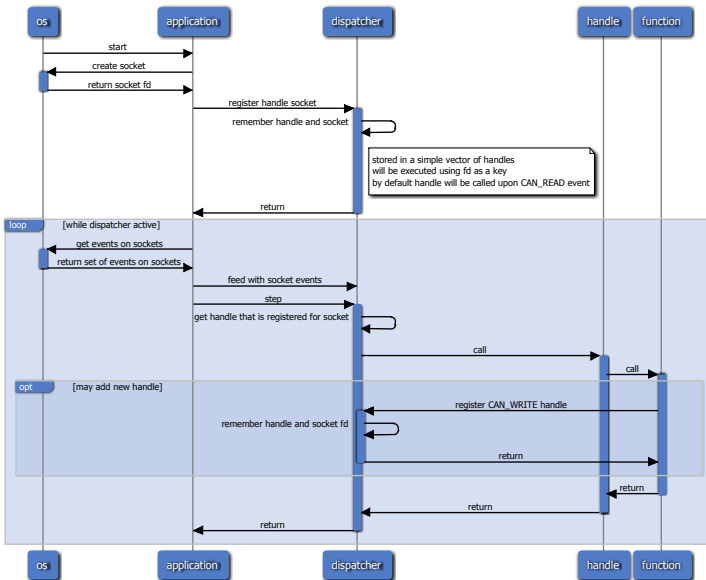
## LibXively dispatcher responsibilities and flow

## Queued events



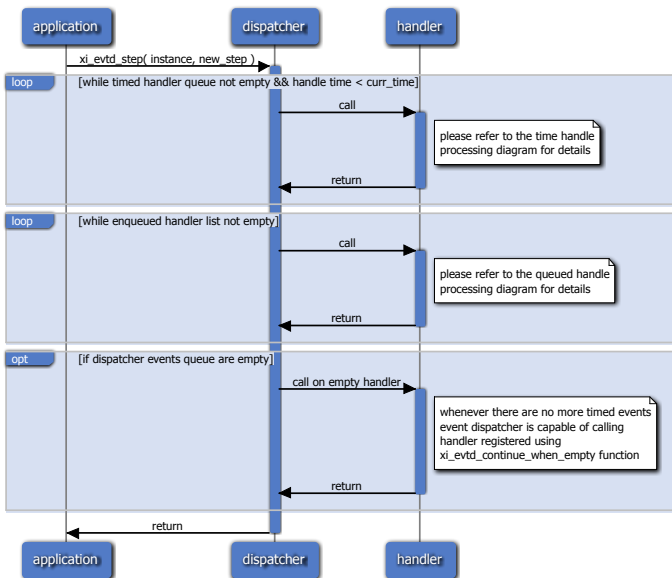
## LibXively dispatcher responsibilities and flow

## Socket events



Usage of LibXively event dispatcher on receiving example

# Processing of event



Example  
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Coroutines?  
oooooooo

LibXively Coroutines  
o

Event dispatcher  
oooooo

The End  
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Thank you!

# Questions?

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