

# Data Structures - CSCI-SHU 210

**Solution**

## Final Practice Exam

<b>First Name</b>	
<b>Last Name</b>	
<b>NetID</b>	
<b>Section</b>	

### Details

- **Duration:** 75 Minutes
- Computers, calculators, phones, textbooks, or notebooks are not allowed

### Instructions

- Fill the table with your information as it appears on Brightspace and Albert.
- Show your NYU ID card to the proctor before signing the attendance sheet.
- You can use the blank space to make notes for yourself. Your notes are not evaluated.
- Turn off and place your mobile devices with your belongings at the front of the room.
- Please do not damage the exam paper.
- Do not discuss this quiz with anyone until the test is handed back.
- If you need to ask questions, please raise your hand and wait for the instructor.
- Stay seated until the instructor has collected all exams.

Please sign the statement below:

I, \_\_\_\_\_, confirm that I am a student in this class and I will complete this exam without accessing any unauthorized information during the exam. I have read and understood the exam instructions above.

# Multiple Choices Questions

After answering all knowledge questions, transfer your solution letter to the table below. Only one solution is correct for each answer option. Please mark your solution clearly:

Questions 1.1 to 1.10

	1.1	1.2	1.3	1.4	1.5	1.6	1.7	1.8	1.9	1.10
Answer	C	A	B	A	D	C	C	B	A	C
Points										

### Question 1.1. - Tree Terminology

If a binary tree has  $n$  nodes, how many edges does it have?

- ☐ A:  $\log_2 n$
- ☐ B:  $n$
- ☐ C:  $n - 1$
- ☐ D: none of the listed

### Question 1.2. - Tree Terminology

Suppose a binary search tree has  $n$  nodes. The minimum height of the binary search tree is  $\text{Ceil}(\log_2(n + 1) - 1)$

- ☐ A: The above statement is True
- ☐ B: The above statement is False

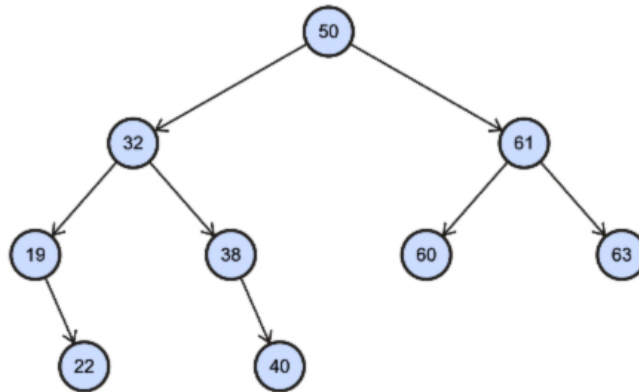
### Question 1.3. - Hashtable

The worst-case complexity for getting the length of a hashtable is  $O(n)$ , based on the Hashtable data structure from our lecture and recitation.

- ☐ A: The above statement is True
- ☐ B: The above statement is False

## Question 1.4. - Tree Traversal

Consider the following tree and select the correct statement below:



- ☐ A: The element before 32 in an in-order traversal is: 22
- ☐ B: The element before 50 in an in-order traversal is: 61
- ☐ C: The element after 19 in an in-order traversal is: 32
- ☐ D: The element after 40 in an in-order traversal is: 60

## Question 1.5. - Function Mystery

Consider the following code snippet:

```
def mystery(lis):  
    S = [0] * len(lis)  
    for i in range(len(lis)-1):  
        T = lis[0 : i + 1]  
        S[i] = sum(T) * (i + 2)  
    return S
```

Suppose  $\text{len}(\text{lis}) = n$ . The worst-case big-O analysis of this code is:

- ☐ A:  $O(\log_2 n)$
- ☐ B:  $O(\log_2 n * \log_2 n)$
- ☐ C:  $O(n)$
- ☐ D:  $O(n^2)$

## Question 1.6. - Time Complexity

What is the time complexity of the following function, assuming  $n = \text{len}(\text{string})$ :

```
def is_special(string):  
    return special_helper(string, 0)  
  
def special_helper(string, index):  
    if index >= len(string) // 2:  
        return True  
    if string[index] != string[-1 - index]:  
        return False  
    return special_helper(string, index + 1)
```

- ☐ A:  $O(n^2)$
- ☐ B:  $O(\log_2 n)$
- ☐ C:  $O(n)$
- ☐ D:  $O(1)$

## Question 1.7. - Tree Insert

The following items are inserted into a binary search tree in this order:

12, 8, 5, 3, 9, 13, 14, 52, 1, 2, 55, 4

Which node is the deepest?

- ☐ A: 9
- ☐ B: 55
- ☐ C: 2
- ☐ D: 1

### Question 1.8. - Heap Remove

What is the worst-case time complexity of removing the minimum item in a min-heap of size  $n$  implemented using a linked binary tree:

- ☐ A:  $O(n * \log_2 n)$
- ☐ B:  $O(\log_2 n)$
- ☐ C:  $O(n)$
- ☐ D:  $O(1)$

### Question 1.9. - Heap Correct Statement

Select the correct statement:

- ☐ A: The last node in a heap is the deepest rightmost leaf node in a complete tree
- ☐ B: The last node in a heap is the deepest leftmost leaf node in a complete tree
- ☐ C: The last node in a heap is the deepest rightmost non-leaf node in a full tree
- ☐ D: The last node in a heap is the deepest leftmost non-leaf node in a full tree

### Question 1.10. - Hashtable

For a given hashtable using the hash function  $h(x) = x \bmod 17$  and linear probing, as seen in the picture below:

	18	19	88	1	22	90	41	59	77	44				14	32
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15

How many probes will be required to insert 8?

- ☐ A: 4
- ☐ B: 5
- ☐ C: 3
- ☐ D: 2

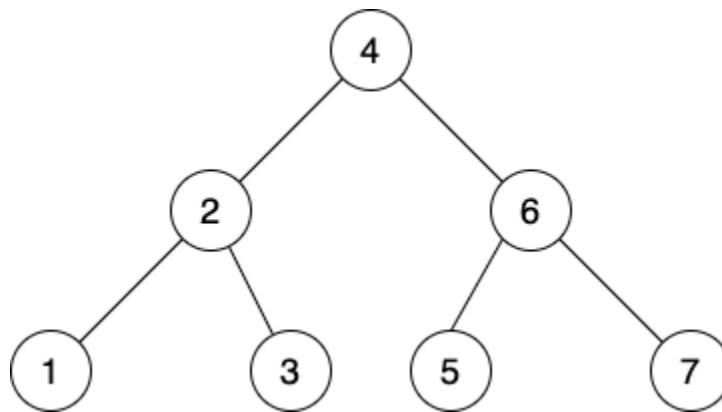
## 2. Programming Question

Please try to write your programming solution as clearly as possible.

### Question 2.1. - Search Tree from Array

Given a sorted list of numbers, convert the list to a perfect Binary Search Tree in array representation, as seen in the figure. Remember: A perfect tree is a complete tree with the maximum number of leaf nodes.

Binary Search Tree tree representation for the sorted array [1, 2, 3, 4, 5, 6, 7]:



**Example 1:**

```
def main():  
    ls = [1, 2, 3, 4, 5, 6, 7]  
    res = convert(ls)  
    print(res) # Should print: [4, 2, 6, 1, 3, 5, 7]
```

#### Requirements:

- You are not allowed to use global variables.
- You can not use any library or third-party tools.
- You are allowed to use the Math libraries log function.
- Your solution has to be  $O(n)$  time complexity.
- Your solution has to be  $O(n)$  space complexity.
- You are not allowed to create a linked binary search tree.

```

def convert(ls):

    def convert(ls):
        res = [None] * len(ls)
        convert_helper(0, ls, res, 0, len(ls) - 1)
        return res

    def convert_helper(result_idx, input, output, left, right):
        if left > right:
            return

        mid = (left + right) // 2
        output[result_idx] = input[mid]

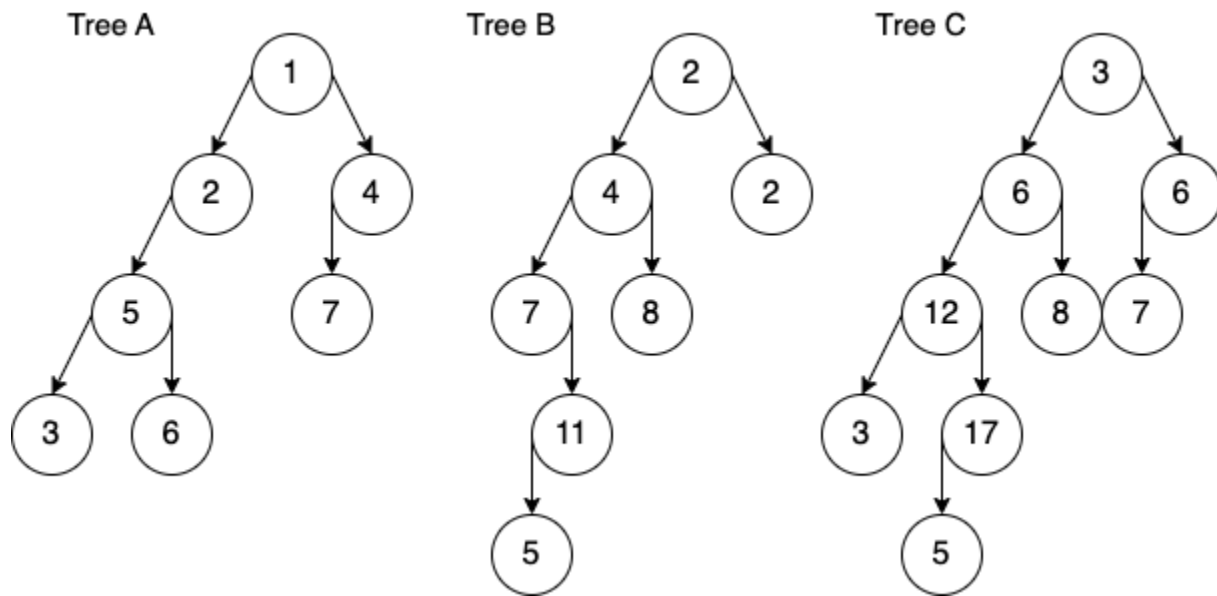
        convert_helper(result_idx*2+1, input, output, left, mid - 1)
        convert_helper(result_idx*2+2, input, output, mid + 1, right)

```



## Question 2.2. - Merging Trees

Given two binary trees A and B, of size  $n$  and  $m$ , respectively, write a function to merge them into a tree C. The function should perform the sum of the elements in two nodes if the corresponding node exists in both trees (e.g., the root node in tree A and tree B in the figure). Your function should merge all nodes in the output tree C. Your function has to create new nodes and return the root node of a new tree C. You cannot modify the values or structures in the input trees.



The above picture shows the merge of trees A and B, with the resulting tree C.

### Requirements:

- You are not allowed to use global variables.
- You cannot use any library or third-party tools.
- Your solution has to be in  $O(n+m)$  time complexity
- Your solution has to be  $O(n+m)$  space complex
- You are not allowed to modify the input trees.
- Your function has to return the root node of a new tree.

**Example 1:**

```
treeA = TreeWithoutParent(1)
treeA._left = TreeWithoutParent(2)
treeA._right = TreeWithoutParent(4)
treeA._left._left = TreeWithoutParent(5)
treeA._right._left = TreeWithoutParent(7)
treeA._left._left._left = TreeWithoutParent(3)
treeA._left._left._right = TreeWithoutParent(6)

treeB = TreeWithoutParent(2)
treeB._left = TreeWithoutParent(4)
treeB._right = TreeWithoutParent(2)
treeB._left._left = TreeWithoutParent(7)
treeB._left._right = TreeWithoutParent(8)
treeB._left._left._right = TreeWithoutParent(11)
treeB._left._left._right._left = TreeWithoutParent(5)

treeC = treeA + treeB
```

```

def __add__(self, other):

    def _add_helper(self, t1, t2):

        # case 1: the nodes exist in tree A and tree B
        if t1 and t2:
            t3 = TreeWithoutParent(t1._element + t2._element)
            t3._left = self._add_helper(t1._left, t2._left)
            t3._right = self._add_helper(t1._right, t2._right)

        # case 2: the node only exists in tree A
        elif t1:
            t3 = TreeWithoutParent(t1._element)
            t3._left = self._add_helper(t1._left, None)
            t3._right = self._add_helper(t1._right, None)

        # case 3: the node only exists in tree B
        elif t2:
            t3 = TreeWithoutParent(t2._element)
            t3._left = self._add_helper(None, t2._left)
            t3._right = self._add_helper(None, t2._right)

        # base case: there is no node
        else:
            return None

    return t3

def __add__(self, other):
    return self._add_helper(self, other)

```

## Question 2.3.- Binary Tree Layers

Given an array representation of a binary tree, return a list of all node values at a given depth. For example, if depth  $d$  is equal to 0, you should return a list containing only the element of the root node. If the depth is equal to 1, you should return a list containing the elements of the direct children of the root node, and so on. Skip any non-existent nodes. Notice that the value  $d$  can be any non-negative integer.

### Example:

```
def main():
    lissy = [4, 2, 6, None, 1, 3, 5, 7]

    print(get_layer(lissy, 0)) # Expect: [4]
    print(get_layer(lissy, 1)) # Expect: [2, 6]
    print(get_layer(lissy, 2)) # Expect: [1, 3, 5]
    print(get_layer(lissy, 3)) # Expect: [7]
    print(get_layer(lissy, 4)) # Expect: []
```

### Requirements:

- You are not allowed to use global variables.
- You are allowed to use helper functions.
- Your solution has to be recursive. Loops are not allowed.
- Do not use any library or third-party tools.
- Your solution must have a running time of at most  $O(2^d)$ .
- You can not create a tree of nodes to solve this problem.
- Your output list can not contain *None* values.

```
def get_layer(lissy, depth):  
    return get_layer_helper(lissy, depth, 0, 0)  
  
def get_layer_helper(lissy, d, cur_root_idx, i):  
    if cur_root_idx >= len(lissy):  
        return []  
    if i == d:  
        value = lissy[cur_root_idx]  
        return [] if value is None else [value]  
    return get_layer_helper(lissy, d, cur_root_idx*2+1, i+1) +  
        get_layer_helper(lissy, d, cur_root_idx*2+2, i+1)
```

## Appendix 1 - TreeWithoutParent Implementation

```
class TreeWithoutParent:
    def __init__(self, element, left=None, right=None):
        self._element = element
        self._left = left
        self._right = right

    def __add__(self, other):
        # You need to implement this function.
        # Here, this function is to merge self & other into a new tree
        # Return type: an instance of TreeWithoutParent
        pass
```

*You can tear off this page and use it as scratch paper.*