Priority Queues



Priority Queue ADT

- A priority queue stores a collection of items
- Each item is a pair (key, value)
- Main methods of the Priority
 Queue ADT
 - add (k, x)
 inserts an item with key k
 and value x
 - remove_min()
 removes and returns the
 item with smallest key

- Additional methods
 - min()
 returns, but does not
 remove, an item with
 smallest key
 - len(P), is_empty()
- Applications:
 - Standby flyers
 - Auctions
 - Job scheduler in Operating systems

Priority Queue Example

Operation	Return Value	Priority Queue
P.add(5,A)		{(5,A)}
P.add(9,C)		{(5,A), (9,C)}
P.add(3,B)		{(3,B), (5,A), (9,C)}
P.add(7,D)		{(3,B), (5,A), (7,D), (9,C)}
P.min()	(3,B)	{(3,B), (5,A), (7,D), (9,C)}
P.remove_min()	(3,B)	{(5,A), (7,D), (9,C)}
P.remove_min()	(5,A)	{(7,D), (9,C)}
len(P)	2	{(7,D), (9,C)}
P.remove_min()	(7,D)	{(9,C)}
P.remove_min()	(9,C)	{ }
P.is_empty()	True	{ }
P.remove_min()	"error"	{ }

Total Order Relations

- Keys in a priority queue can be arbitrary objects on which an order is defined
- Two distinct
 entries in a
 priority queue can
 have the same
 key

- Mathematical conceptof total order relation ≤
 - Reflexive property:x ≤ x
 - Antisymmetric property: $x \le y \land y \le x \Rightarrow x = y$
 - Transitive property: $x \le y \land y \le z \Rightarrow x \le z$

e.g. Alphabetical order: Define x <= y if 'x' is before 'y' in alphabetical order

Composition Design Pattern

- An item in a priority queue is simply a key-value pair
- Priority queues store items to allow for efficient insertion and removal based on keys

```
class PriorityQueueBase:
      """ Abstract base class for a priority queue."""
      class _ltem:
        """Lightweight composite to store priority queue items."""
        __slots__ = '_key', '_value'
        def __init__(self, k, v):
          self.\_key = k
          self. value = v
        def __lt__(self, other):
          return self._key < other._key
                                             # compare items based on their keys
15
      def is_empty(self):
                                        # concrete method assuming abstract len
        """Return True if the priority queue is empty."""
16
        return len(self) == 0
```

Sequence-based Priority Queue

Implementation with an unsorted doubly linked list

4-5-2-3-1

Performance:

- add takes O(1) time since we can insert the item at the beginning or end of the sequence
- Remove_min and min take
 O(n) time since we have to traverse the entire sequence to find the smallest key

Implementation with a sorted doubly linked list

1-2-3-4-5

Performance:

- add takes O(n) time since we have to find the place where to insert the item
- remove_min and min take O(1) time, since the smallest key is at the beginning

Runtime of implementing PQ using Sorted and Unsorted List

Operation	Unsorted List	Sorted List
len	<i>O</i> (1)	<i>O</i> (1)
is_empty	O(1)	<i>O</i> (1)
add	<i>O</i> (1)	O(n)
min	O(n)	O(1)
remove_min	O(n)	<i>O</i> (1)

Priority Queue Sorting

- We can use a priority queue to sort a set of comparable elements
 - 1. Insert the elements one by one with a series of add operations
 - 2. Remove the elements in sorted order with a series of remove_min operations
- The running time of this sorting method depends on the priority queue implementation

```
Algorithm PQ-Sort(S, C)
    Input sequence S, comparator C for
    the elements of S
     Output sequence S sorted in
     increasing order according to C
    P \leftarrow priority queue with
         comparator C
    while ¬S.is_empty ()
         e \leftarrow S.remove\ first()
         P. add (e, \emptyset)
    while \neg P.is\_empty()
         e \leftarrow P.removeMin().key()
         S.add last(e)
```