

Efficient Generation of Geographically Accurate Transit Maps

Hannah Bast¹, Patrick Brosi¹ and Sabine Storandt²

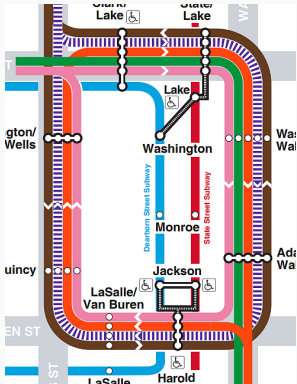
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26th ACM SIGSPATIAL - Seattle, Washington, USA

Motivation

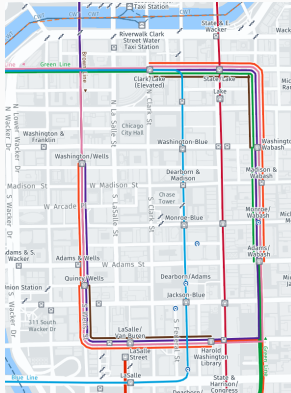
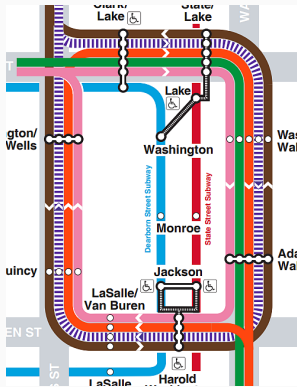
Official CTA map



Motivation

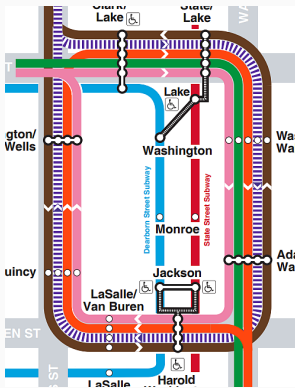
Official CTA map

HERE

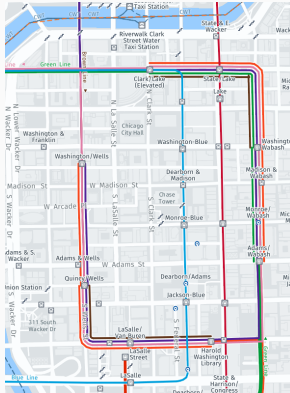


Motivation

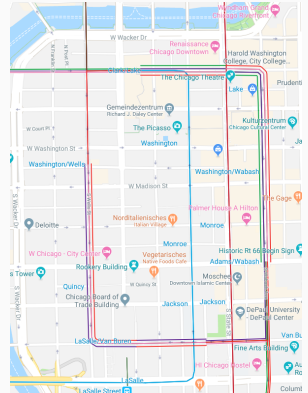
Official CTA map



HERE



Google



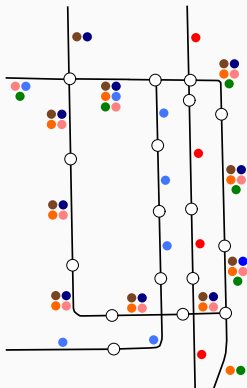
Goal

Goal: Generate these maps automatically, in high quality

"Bag of trips"

Goal

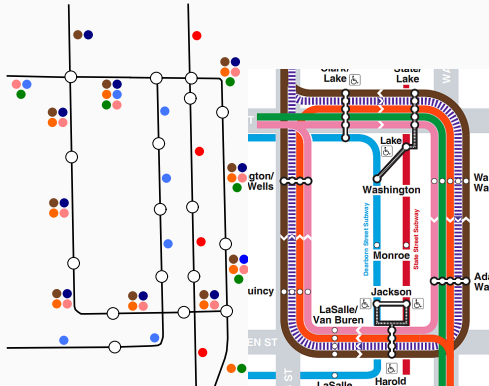
Goal: Generate these maps **automatically**, in **high quality**



"Bag of trips"

Goal

Goal: Generate these maps automatically, in high quality



"Bag of trips"

Challenges

Official



HERE



Google



I. Avoid line overlaps

Challenges

Official

HERE

Google



I. Avoid line overlaps



II. Match line orderings

Challenges

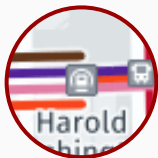
Official

HERE

Google



I. Avoid line overlaps



II. Match line orderings

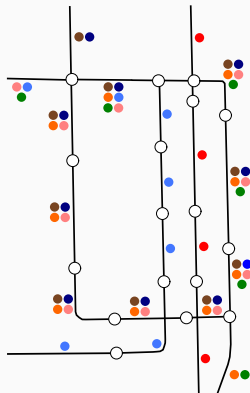


III. Clearly indicate line continuations

Line graph construction

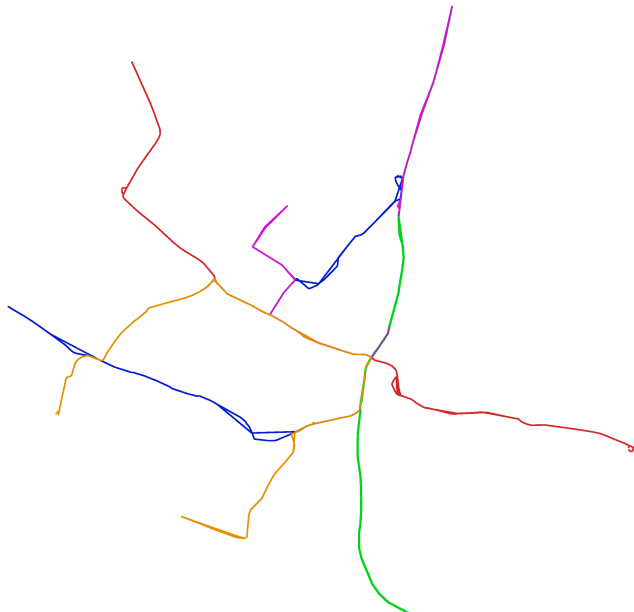
Line graph:

- Undirected labeled graph
 $G = (V, E, L)$
- Edge labels are subsets of the network lines \mathcal{L}
 $(L(e) \subseteq \mathcal{L})$
- Nodes are **usually** stations

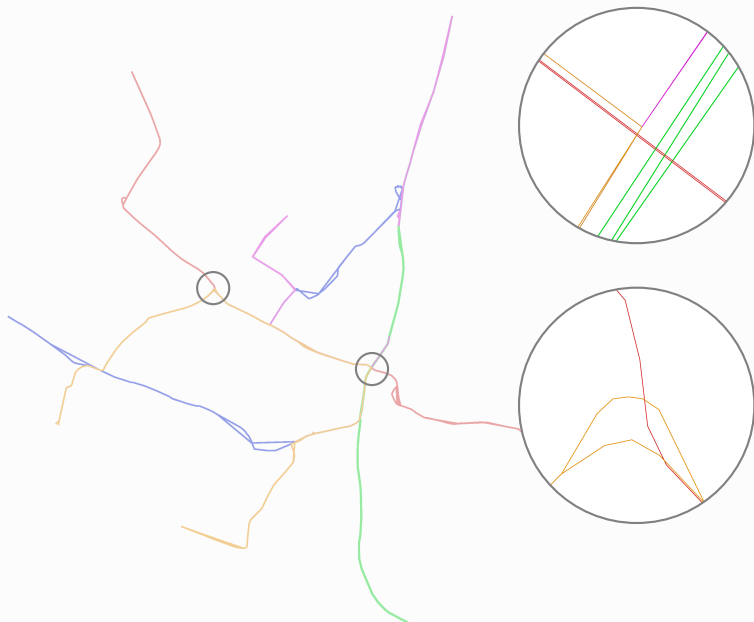


Example: $\mathcal{L} = \{\}$ and $L((a, b)) = \{\}$

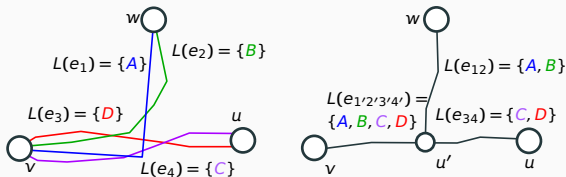
Line graph construction - Input data



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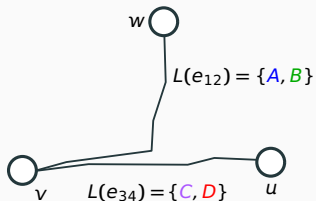


Line graph construction - Shared segment collapsing

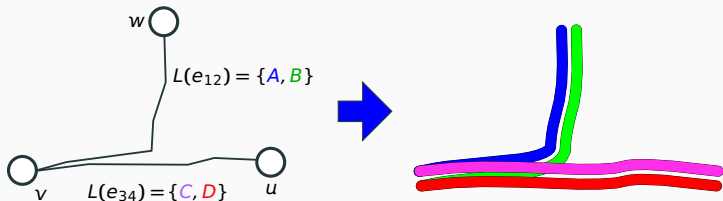


- Repeatedly collapse (segments of) two edges e and f within a distance \hat{d}
- Sweep over some edge e in steps of 10 m, measure distance d of current point on e to f
- If $d < \hat{d}$, start new segment. If not, end current (if open)
- Take average between the two "shared segments" on e and f
- Add **additional non-station nodes** at segment boundaries

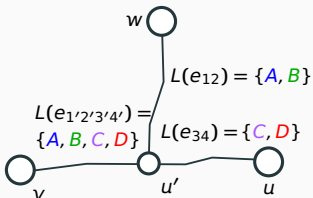
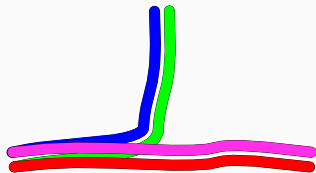
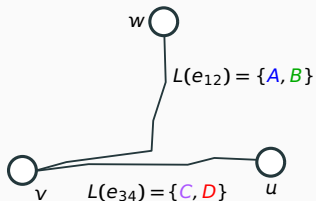
Line graph construction - Why non-station nodes?



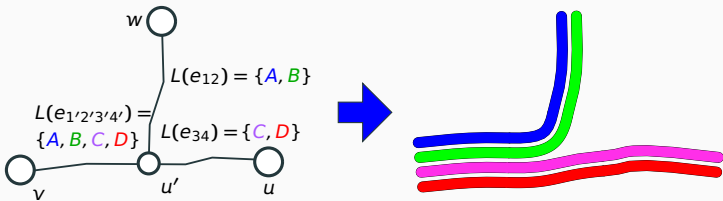
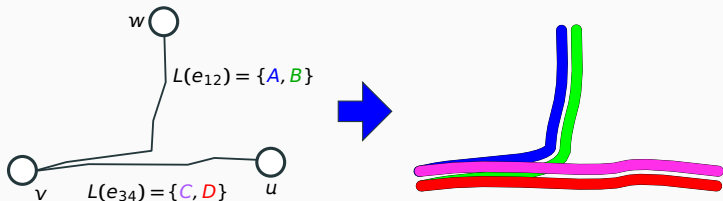
Line graph construction - Why non-station nodes?



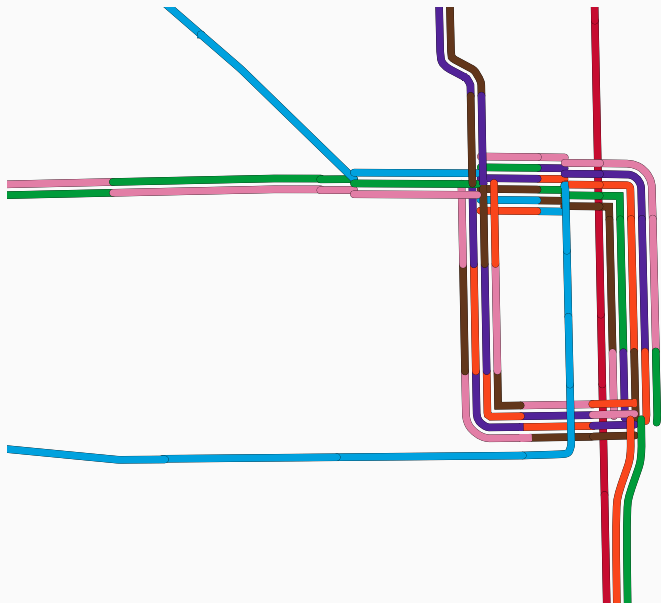
Line graph construction - Why non-station nodes?



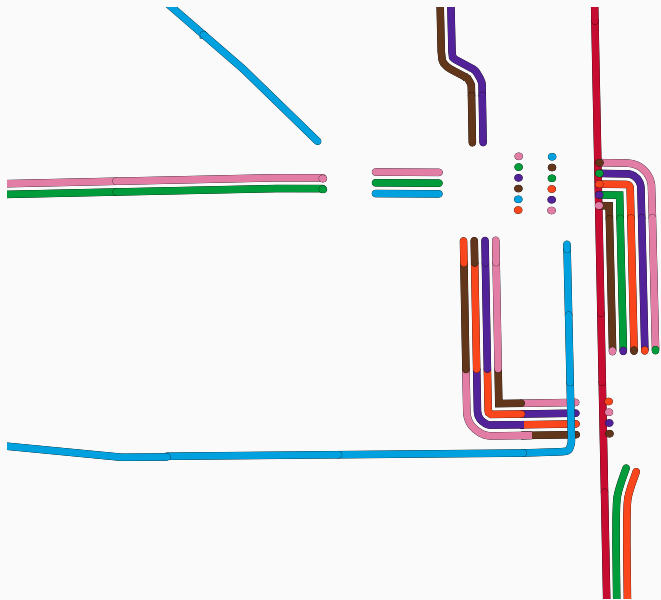
Line graph construction - Why non-station nodes?



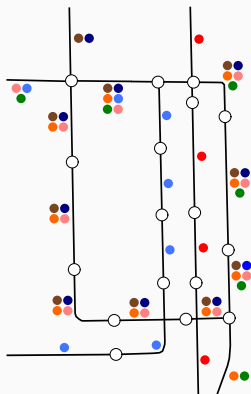
Results so far (1)



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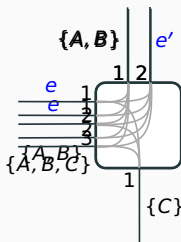
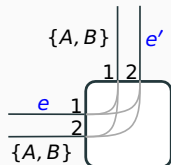
Line-ordering optimization



- 23 edges
- Each edge e has $|L(e)|!$ possible line permutations
- Possible configurations for the graph on the left: $> 2 \times 10^{17}$

⇒ **Naive exhaustive search infeasible**

Line-ordering optimization - Baseline ILP



- For each edge e , line l and position p , introduce variable $x_{elp} \in 0, 1$
- **Example:** x_{eA1} and x_{eA2} for line A
- Constraint: all x_{elp} have to sum up to 1 for a single line l on a single edge e
- **Standard crossing:** Objective variable $x_{ee'AB}$ which is 1 if $p_e(A) < p_e(B)$ and $p_{e'}(A) > p_{e'}(B)$, or else 0
- **Split crossing:** Objective variable $x_{ee'e''AB}$ which is 1 if $p_e(A) < p_e(B)$, or else 0

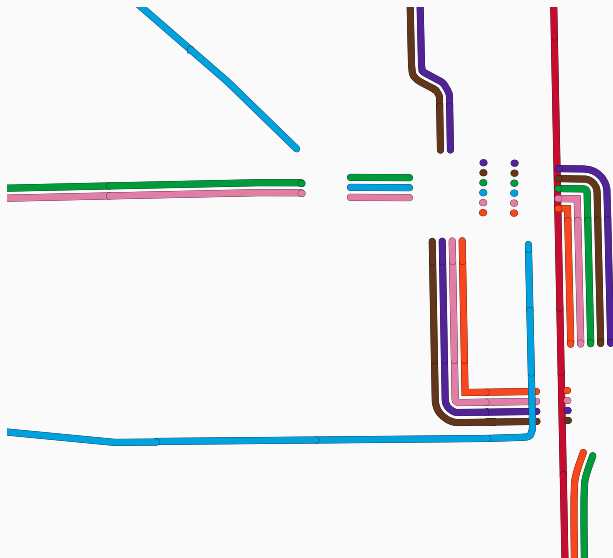
$\Rightarrow \mathcal{O}(|E|M^2)$ variables, $\mathcal{O}(|E|M^6)$ constraints

Line-ordering optimization - Improved ILP

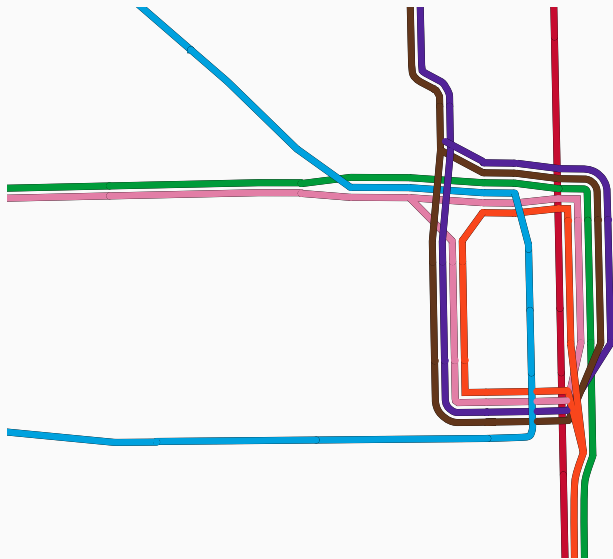
- **Observation:** we only need to check if $p_e(A) < p_e(B)$ (or vice versa) for both types of crossings
- But we explicitly enumerate all possible line positions of A and B on e
- **Basic idea:** introduce binary variables $x_{eA < B}$ and $x_{eB < A}$ which can be efficiently checked

$\Rightarrow \mathcal{O}(|E|M^2)$ variables, $\mathcal{O}(|E|M^2)$ constraints

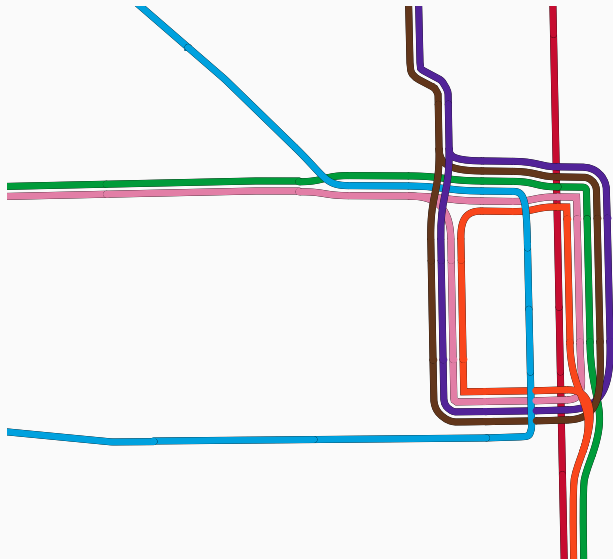
Results so far (2)



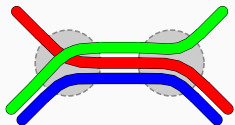
Results so far (2)



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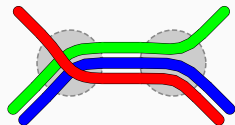
Line-ordering optimization - Line separations



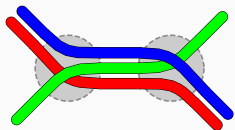
1 crossing, 1 separation

VS

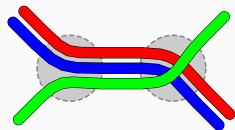
VS



2 crossings, 0 separations

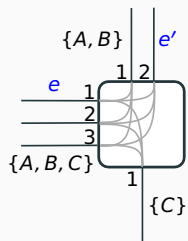


2 crossings, 1 separation



2 crossings, 0 separations

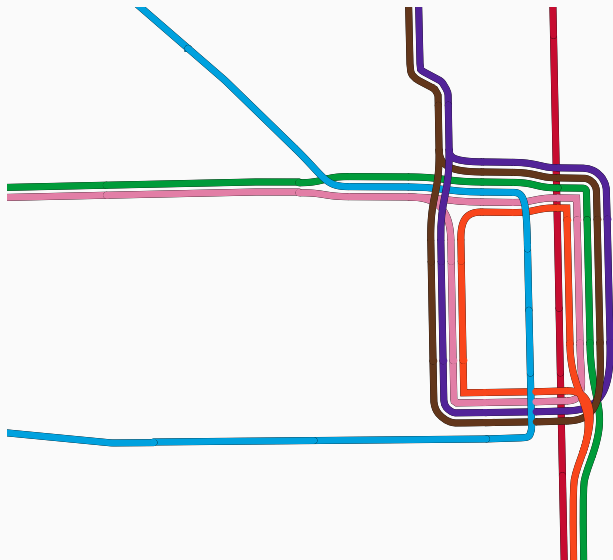
Line-ordering optimization - Line separations (ctd.)



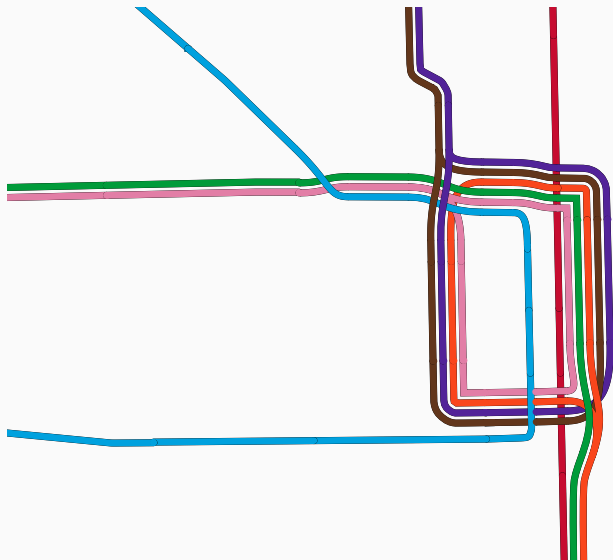
- **Idea:** If two lines A, B continue from e to e' , set a binary splitting variable $x_{ee'A||B} = 1$ if they are next to each other in e , but no in e'
- Add $x_{ee'A||B}$ to the objective function

\Rightarrow Still $\mathcal{O}(|E|M^2)$ variables, $\mathcal{O}(|E|M^2)$ constraints

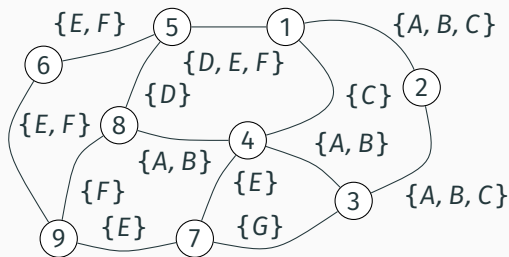
Results so far (3)



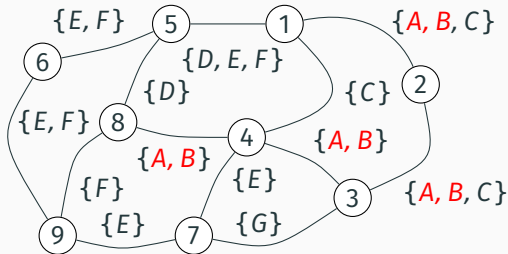
Results so far (3)



Line-ordering optimization - Core optim graph

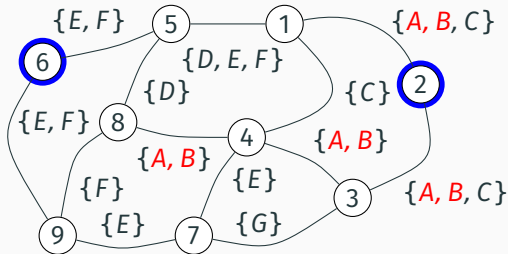


Line-ordering optimization - Core optim graph



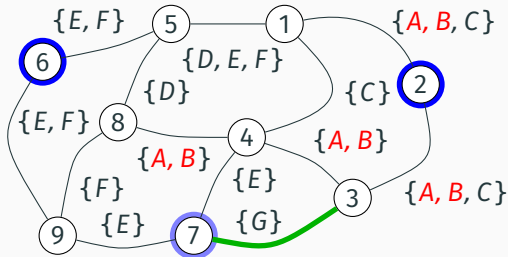
- Combine lines A, B that **always occur together** into a single new line ($A, B \rightarrow X$)

Line-ordering optimization - Core optim graph



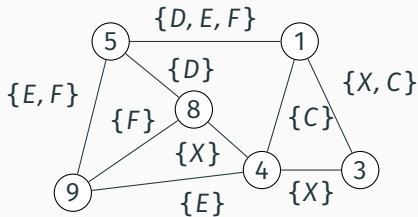
- Combine lines **A, B** that **always occur together** into a single new line ($A, B \rightarrow X$)
- Delete **nodes with degree 2** and merge their adjacent edges

Line-ordering optimization - Core optim graph



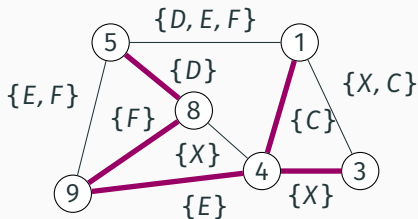
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- Remove edges (u, v) where u and v are **termini** for all $L((u, v))$

Line-ordering optimization - Core optim graph



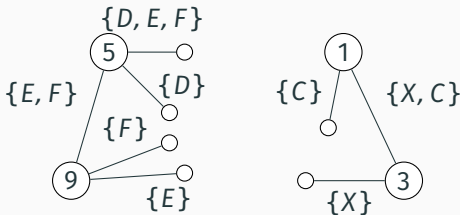
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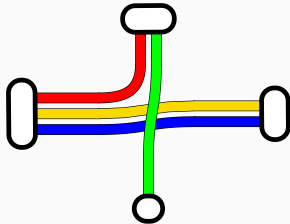
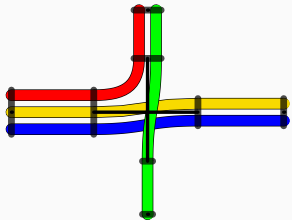
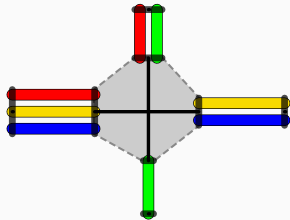
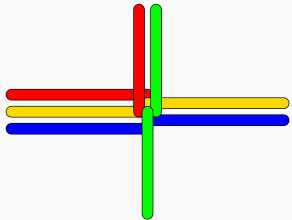
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- Cut **edges** with $|L(e)| = 1$

Line-ordering optimization - Core optim graph

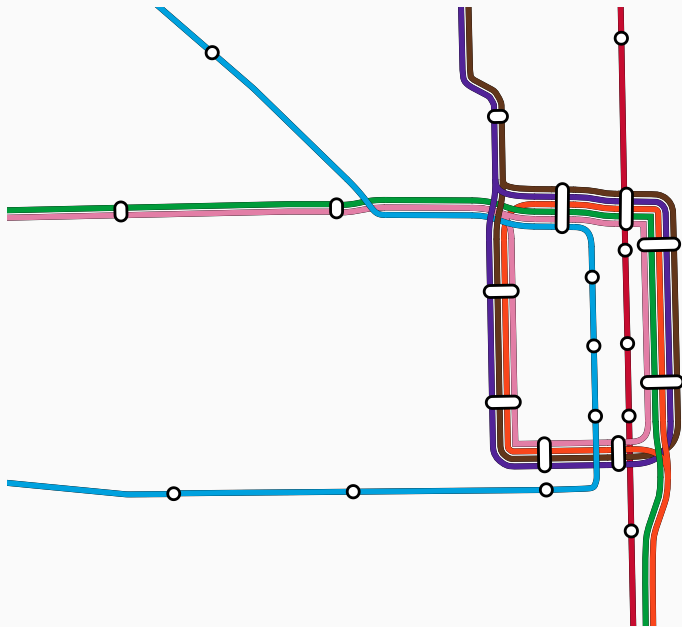


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- Delete **nodes with degree 2** and merge their adjacent edges
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- Cut **edges** with $|L(e)| = 1$

Rendering



Results so far (4)



Future work

- Additional rules for core graph reduction (work in progress)
- Faster construction times of line graph
- TODO

Thank you!