Software Defined Networking

SDN

Software Defined Networking

- Objective
 - A tool to enable a higher degree of control over network devices and traffic flow
- Main aspects
 - Control plane is separated from the device implementing the data plane
 - A single control plane is bale to manage multiple network devices
- Initial deployments
 - Universities: to try out radical new protocols in parallel with existing traffic
 - Busy data centres: overcome the VLAN ID tag limit (4095)
- Protocols:
 - OpenFlow (first)
 - P4 (now)

OpenFlow





- Standardised by the ONF Open Networking Foundation
- https://opennetworking.org/software-defined-standards/specifications/
- Currently on version 1.5

ABOUT BROADBAND PRIVATE 5G & EDGE OPEN RAN PROGRAMMABLE NETWORKS OTHER PROJECTS

Specifications

Technical Specifications include all standards that define a protocol, information model, functionality of components and related framework documents. It is this category of Technical Specification that is identified as such because it is a normative publication that has the ONF RAND-Z IPR policy and licensing guiding its further use.

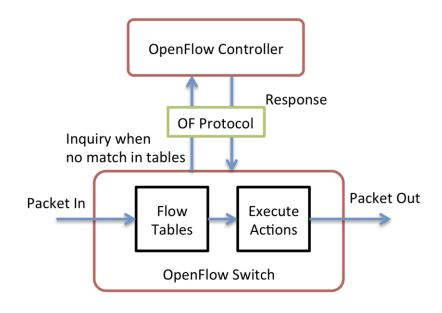
SPECIFICATIONS MODELS & APIS INFORMATIONAL

Current Versions

+ REFERENCE	+ REFERENCE DESIGNS			
+ P4 LANGUAGE & RELATED SPECIFICATIONS			Show	
OPENFLOW	OPENFLOW SPECIFICATIONS		Hide	
DATE	DOCUMENT NAME	DOCUMENT TYPE & ID	FORMAT	
06/2017	SPTN OpenFlow Protocol Extensions	TS-029	PDF	
04/2017	Optical Transport Protocol Extensions Ver. 1.0	TS-022	PDF	
04/2015	OpenFlow® Switch Specification Ver 1.5.1	TS-025	PDF	

OpenFlow

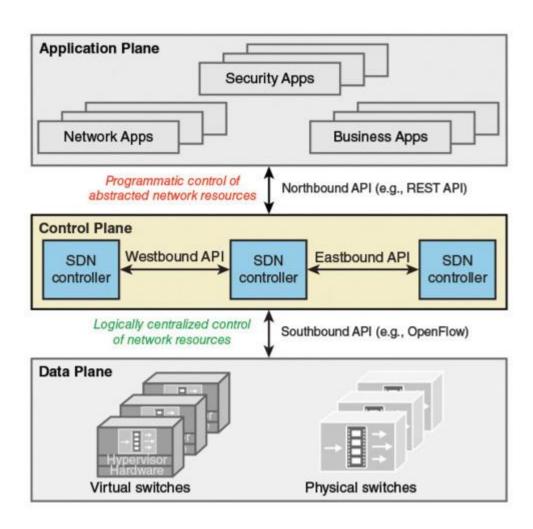
- A protocol existing between SDN switches and a new entity
 - SDN controller
- Allows the SDN controller to manage flow tables in SDN switches
- SDN switch contains
 - Openflow agent
 - Flow tables
 - Performs packet lookup and forwarding
 - Is able to communicate (securely) with the controller
- A flow table is composed of
 - Flow entries (matches properties in packets' headers)
 - Counters (for activity)
 - A set of actions to be applied (to matching packets)
 - When no actions are presente, the switch can
 - Drop the packet
 - Ask the controller what to do



Source: fibre-optic-transceiver-module.com

OpenFlow Controller

- Controller
 - Service existing in a server, which uses the OpenFlow protocol to interact with SDN switches
 - Formulates flows and programs switches
 - Is able to receive directives from external applications via northbound REST API



OpenFlow Controller

- Many existing flavours
 - ONOS Open Network Operating System
 - https://opennetworking.org/onos/
 - TeraFlow SDN
 - https://tfs.etsi.org
 - OpenDaylight
 - https://opendaylight.org
 - Floodlight (last version from 2016)
 - https://floodlight.atlassian.net
 - NOX (10 w/o maintenance)
 - POX (Python version of NOX w/ some maintenance) <u>https://noxrepo.github.io/pox-doc/html/</u>
 - Ryu (w/o maintenance since 2017)
 - https://ryu-sdn.org/
 - Trema (5y since last update)
 - https://github.com/trema/trema
 - Frenetic (last release: 2019)
 - https://github.com/frenetic-lang/frenetic











Differences between SDN Controllers

- Research vs production
- Programming language
- Performance
- Learning Curve
- User base and Support
- Focus
 - Southbound API Support
 - Northbound API
 - OpenFlow version

A Qualitative and Quantitative assessment of SDN Controllers

Pedro Bispo, Daniel Corujo, Rui L. Aguiar Instituto de Telecomunicações e Universidade de Aveiro, Portugal Email: {pedrobispo, ruilaa}@ua.pt; dcorujo@av.it.pt

Abstract—With the increasing number of connected devices, new challenges are being raised in the networking field. Software Defined Networking (SDN) enables a greater degree of dynamism and simplification for the deployment of future 5G networks. In such networks, the controller plays a major role by being able to manage forwarding entities, such as switches, through the application of flow-based rules via a southbound (SB) interface. In turn, the controller itself can be managed by means of actions and policies provided by high-level network functions, via a northbound (NB) interface.

The growth of SDN integration in new mechanisms and network architectures led to the development of different controller solutions, with a wide variety of characteristics. Despite existing studies, the most recent evaluations of SDN controllers are focused only on performance and are not up to date, since new versions of the most popular controllers are constantly being released. As such, this work provides a wider study of several open-source controllers, (namely, OpenDaylight (ODL), Open Network Operative System (ONOS), Ryu and POX), by evaluating not only their performance, but also their characteristics in a qualitative way. Taking performance as a critical issue among SDN controllers, we quantitatively evaluated several criteria by benchmarking the controllers under different operational conditions, using the Cbench tool.

Keywords—Software-Defined Networking, OpenFlow, SDN controller.

reachability optimization towards the users, and the users themselves want overall better service. Simultaneously satisfying involved actors is a highly complex task, whose harmonization is only achieved through careful planning and overprovisioning of networking resources. Nonetheless, the increase in generated data [1] and the need to dynamically adapt to changing situations in a cost-effective way, are demanding for more flexible and adaptive network control mechanisms. As a result, SDN has emerged.

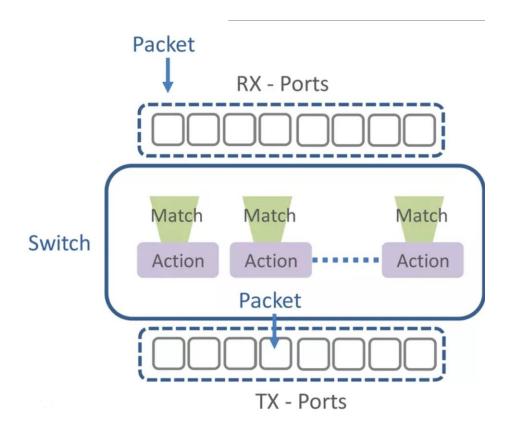
SDN provides the separation of the network control plane from the forwarding plane, allowing more control, adaptability, agility and overall cost reduction. By having a complete view of the network, an SDN controller plays an extremely important role in such networks as it can manage the network structure and services dynamically.

Several SDN controllers exist, with most of them under continuous development. This diversity, which originated from the different needs of operators and research teams that resulted in the development of their own controller versions, made comparison efforts more difficult.

This diversity is evidenced as each controller presents different Northbound (NB) and Southbound (SB) interfaces (which allow it to be interfaced by high-level entities and to control forwarding entities, respectively), development

SDN switch

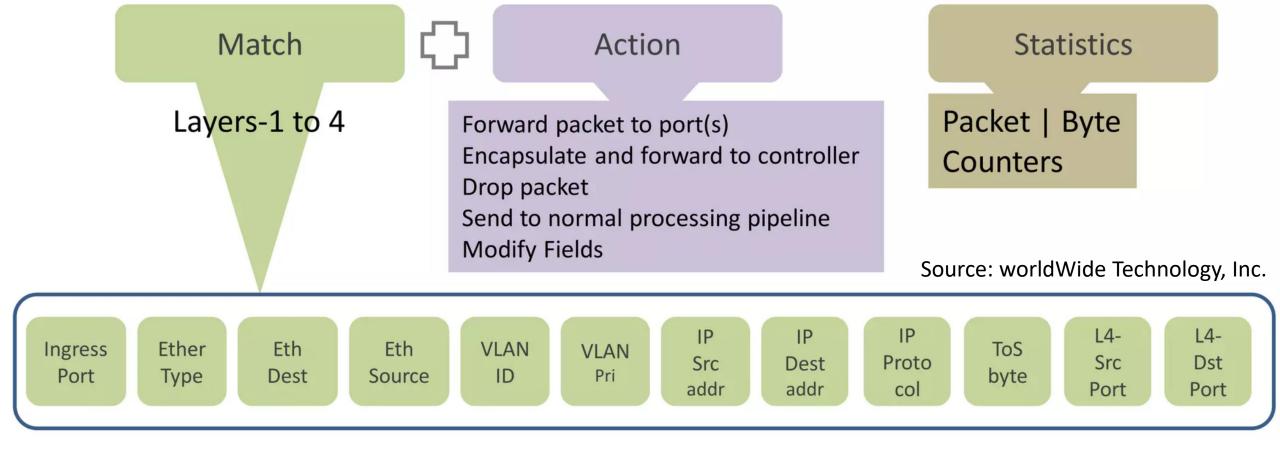
- Has an OpenFlow agent that is able to communicate with the controller
- Processes commands received by the controller
- Dataplane of a switch
 - Ports
 - Flow tables
 - Flows
 - Classifiers (Match)
 - Modifiers and actions
- Packets are matched to flows in flow tables using the match/classifiers
- Flows contain sets of modifiers and actions which are applied to each packet that it matches



Source: worldWide Technology, Inc.

SDN Switch – Flow table

• Each flow table entry contains: Match, Action and counters

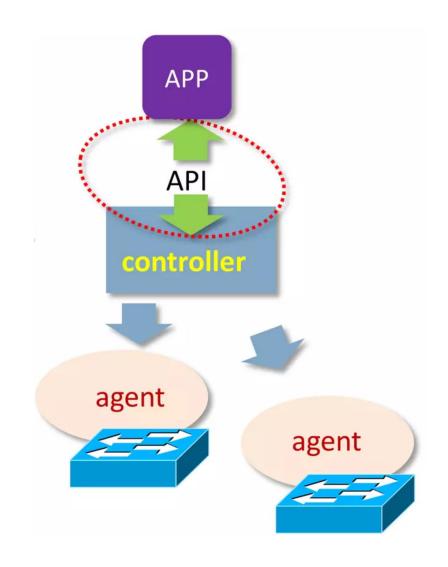


SDN Switch - Actions

- When a switch first connects to a controller, it specifies which actions are supported
 - Not all switches need to implemente all OpenFlow actions
- Examples of actions
 - Forward a packet to a set of ports
 - Drop a packet
 - Add, modify or remove VLAN ID or priority on a per-destination-port basis
 - Modify the IP DSCP (i.e., QoS)
 - Modify the destination MAC address
 - Send the packet to the OpenFlow controller (Packet In)
 - Receive the packet from the OpenFlow controller and send it to ports (Packet out)

How to "direct" the controller?

- Northbound interface allows for Northbound protocols
- Allows applications and orchestration systems to program the network and request services
- Provides a network abstraction interface to applications
- The abstraction is important when managing dissimilar network elements



Northbound Protocol Examples

- There is nothing standardized, but there are "types" of protocols used
- REST (web based) API applications which run on different machine or address space on the controller
- Web Browser
 - http://<SDN Controller IP>:8080/
- WebSockets
- OSGi framework is used for applications that will run in the same address space as the controller.
 - Open Service Gateway Initiative

Towards the next generation SDN

OpenFlow: Enabling Innovation in Campus Networks

Nick McKeown Stanford University Guru Parulkar Tom Anderson University of Washington Larry Peterson Hari Balakrishnan MIT

Stanford University Larry Pet Princeton U Jennifer Rexford Princeton University

Scott Shenker University of California, Berkeley Jonathan Turner Washington University in St. Louis

This article is an editorial note submitted to CCR. It has NOT been peer reviewed. Authors take full responsibility for this article's technical content. Comments can be posted through CCR Online.

ABSTRACT

This whitepaper proposes OpenFlow: a way for researchers to run experimental protocols in the networks they use every day. OpenFlow is based on an Ethernet switch, with an internal flow-table, and a standardized interface to add and remove flow entries. Our goal is to encourage networking vendors to add OpenFlow to their switch products for deployment in college campus backbones and wiring closets. We believe that OpenFlow is a pragmatic compromise: on one hand, it allows researchers to run experiments on heterogeneous switches in a uniform way at line-rate and with high port-density: while on the other hand, vendors do not need to expose the internal workings of their switches. In addition to allowing researchers to evaluate their ideas in real-world traffic settings, OpenFlow could serve as a useful campus component in proposed large-scale testheds like GENI. Two buildings at Stanford University will soon run OpenFlow networks, using commercial Ethernet switches and routers. We will work to encourage deployment at other schools; and We encourage you to consider deploying OpenFlow in your university network too.

Categories and Subject Descriptors

C.2 [Internetworking]: Routers

General Terms

Experimentation, Design

Keywords

Ethernet switch, virtualization, flow-based

1. THE NEED FOR PROGRAMMABLE NETWORKS

Networks have become part of the critical infrastructure of our businesses, homes and schools. This success has been both a blessing and a curse for networking researchers; their work is more relevant, but their chance of making an impact is more remote. The reduction in real-world impact of any given network innovation is because the enormous installed base of equipment and protocols, and the reluctance to experiment with production traffic, which have created an exceedingly high barrier to entry for new ideas. Today, there is almost no practical way to experiment with new network protocols (e.g., new routing protocols, or alternatives to IP) in sufficiently realistic settings (e.g., at scale carrying real traffic) to gain the confidence needed for their widespread deployment. The result is that most new ideas from the networking research community go untried and untested; hence the commonly held belief that the network infrastructure has "sossified"

Having recognized the problem, the networking commity is hard at work developing programmable networks, such as GENI [1] a proposed nationwide research facility for experimenting with new network architectures and distributed systems. These programmable networks call for programmable switches and routers that (using virtualization) can process packets for multiple isolated experimental networks simultaneously. For example, in GENI it is envisaged that a researcher will be allocated a slice of resources across the whole network, consisting of a portion of network links, packet processing elements (e.g. routers) and end-hosts; researchers program their slices to behave as they wish. A slice could extend across the backbone, into access networks, into college campuses, industrial research labs, and include wiring closets, wireless networks, and sensor networks.

Virtualized programmable networks could lower the barrier to entry for new ideas, increasing the rate of innovation in the network infrastructure. But the plans for nationwide facilities are ambitious (and costly), and it will take years for them to be deployed.

This whitepaper focuses on a shorter-term question closer to home: As researchers, how can we run experiments in our campus networks? If we can figure out how, we can start soon and extend the technique to other campuses to benefit the whole community.

To meet this challenge, several questions need answering, including: In the early days, how will college network administrators get comfortable putting experimental equipment (switches, routers, access points, etc.) into their network? How will researchers control a portion of their local network in a way that does not disrupt others who depend on it? And exactly what functionality is needed in network

P4: Programming Protocol-Independent Packet Processors

Pat Bosshartt, Dan Daly', Glen Gibb', Martin Izzardt, Nick McKeownt, Jennifer Rexfordt, Cole Schlesingertt, Dan Talaycot, Amin Vahdatt, George Vargheset, David WalkerttBarefoot Networks "Intel "Stanford University "Princeton University "Google "Microsoft Research

ABSTRACT

P4 is a high-level language for programming protocol-independent packet processors. P4 works in conjunction with SDN control protocols like OpenFlow. In its current form, OpenFlow explicitly specifies protocol headers on which it operates. This set has grown from 12 to 41 fields in a few years, increasing the complexity of the specification while still not providing the flexibility to add new headers. In this paper we propose P4 as a strawman proposal for how Open-Flow should evolve in the future. We have three goals: (1) Reconfigurability in the field: Programmers should be able to change the way switches process packets once they are deployed. (2) Protocol independence: Switches should not be tied to any specific network protocols. (3) Target independence: Programmers should be able to describe packetprocessing functionality independently of the specifics of the underlying hardware. As an example, we describe how to use P4 to configure a switch to add a new hierarchical label.

1. INTRODUCTION

Software-Defined Networking (SDN) gives operators programmatic control over their networks. In SDN, the control plane is physically separate from the forwarding plane, and one control plane controls multiple forwarding devices. While forwarding devices could be programmed in many ways, having a common, open, vendor-agnostic interface (like OpenFlow) enables a control plane to control forwarding devices from different hardware and software vendors.

Version	Date	Header Fields
OF 1.0	Dec 2009	12 fields (Ethernet, TCP/IPv4)
OF 1.1	Feb 2011	15 fields (MPLS, inter-table metadata)
OF 1.2	Dec 2011	36 fields (ARP, ICMP, IPv6, etc.)
OF 1.3	Jun 2012	40 fields
OF 1.4	Oct 2013	41 fields

Table 1: Fields recognized by the OpenFlow standard

The OpenFlow interface started simple, with the abstraction of a single table of rules that could match packets on a dozen header fields (e.g., MAC addresses, IP addresses, protocol, TCP/UDP port numbers, etc.). Over the past five years, the specification has grown increasingly more complicated (see Table 1), with many more header fields and

multiple stages of rule tables, to allow switches to expose more of their capabilities to the controller.

The proliferation of new header fields shows no signs of stopping. For example, data-center network operators intensingly want to apply new forms of packet encapsulation (e.g., NVGRE, VXLAN, and STT), for which they resort to deploying software switches that are easier to extend with new functionality. Rather than repeatedly extending the OpenFlow specification, we argue that future switches should support flexible mechanisms for parsing packets and matching header fields, allowing controller applications to leverage these capabilities through a common, open interface (i.e., a new "OpenFlow 2.0" API). Such a general, extensible approach would be simpler, more elegant, and more future-proof than today's OpenFlow 1.x standard.



Figure 1: P4 is a language to configure switches.

Recent chip designs demonstrate that such flexibility can be achieved in custom ASICs at terabit speeds [1, 2, 3]. Programming this new generation of switch chips is far from easy. Each chip has its own low-level interface, akin to microcode programming. In this paper, we sketch the design of a higher-level language for Programming Protocolindependent Packet Processors (P4). Figure 1 shows the relationship between P4—used to configure a switch, telling it how packets are to be processed—and existing APIs (such as OpenFlow) that are designed to populate the forwarding tables in fixed function switches. P4 raises the level of abstraction for programming the network, and can serve as a

ACM SIGCOMM Computer Communication Review

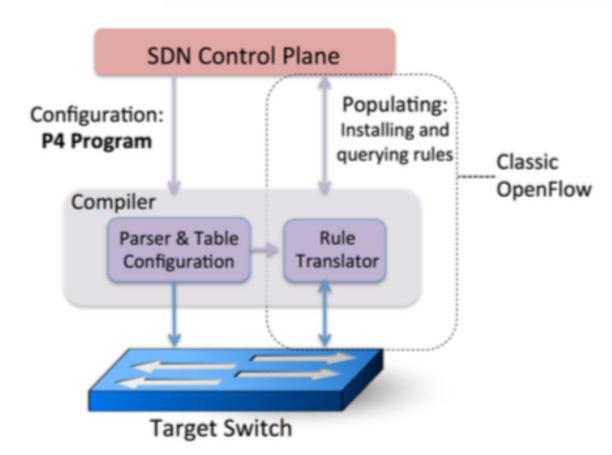
Volume 38, Number 2, April 2008

Programming Protocol-Independent Packet Processors (P4)

- Objectives
 - Reconfigurable
 - Data Plane program can be changed in the field
 - Protocol-Independence
 - No knowledge of low-level hardware organization is required
 - Compiler compiles the program for the target device
 - Architecture dependente e.g. v1model.p4, p4c-xdp.p4, psa.p4
 - Switch/vendor Independence
 - Consistent Control Plane Interface
 - Control plane APIs are automatically generated by the compiler
 - Community-driven design
 - https://p4.org

Programming Protocol-Independent Packet Processors (P4)

- P4 program is a high-level programm that configures forwarding behavior (abstract forwarding model)
- P4 compiler generates the low-level code to be executed by the desired target
- OpenFlow can still be used to install and query rules once forwarding model is defined
- Allows the definition of arbitrary headers and fields



P4 Header and Fields

- Fields have a bit width and other attributes
- Headers are collections of fields
 - Like an instantiated class in Java

```
header_type egress_metadata_t {
header type ethernet t {
                                     fields {
 fields {
                                      nhop type : 8; /* 0: L2, 1: L3, 2: tunnel */
   dstAddr : 48;
                                      encap type : 8; /* L2 Untagged; L2ST; L2DT */
   srcAddr : 48;
                                      vnid
                                                  : 24; /* gnve/vxlan vnid/gre key */
   etherType : 16;
                                                  : 8; /* vxlan; gre; nvgre; gnve*/
                                      tun type
                                      tun idx
                                                  : 8; /* tunnel index */
/* Instance of eth header */
header ethernet t inner ethernet;
                                  metadata egress metadata t egress metadata;
```

Parser

- Extracts header instances
- Selects a next "state" by returning another parser function

```
parser parse ethernet {
 extract(ethernet);
 return select(latest.etherType) {
    ETHERTYPE CPU : parse cpu header;
    ETHERTYPE VLAN : parse vlan;
    ETHERTYPE MPLS : parse mpls;
    ETHERTYPE IPV4 : parse ipv4;
    ETHERTYPE IPV6 : parse ipv6;
    ETHERTYPE ARP : parse arp rarp;
    ETHERTYPE RARP : parse arp rarp;
    ETHERTYPE NSH : parse nsh;
```

Match + Action table

Parsed representation of headers gives contexto for processing of the packets

Match semantics

An action function consists of several primitive actions

```
- exact
table acl {
                                                                               · port index : exact
     reads {
                                                                            - ternary
                                                                               · ethernet.srcAddr : ternary
          ipv4.dstAddr : ternary;
                                                                            - valid
          ipv4.srcAddr : ternary;

    vlan tag[0]: valid

          ipv4.protocol : ternary;
                                                                            - lpm
          udp.srcPort : ternary;

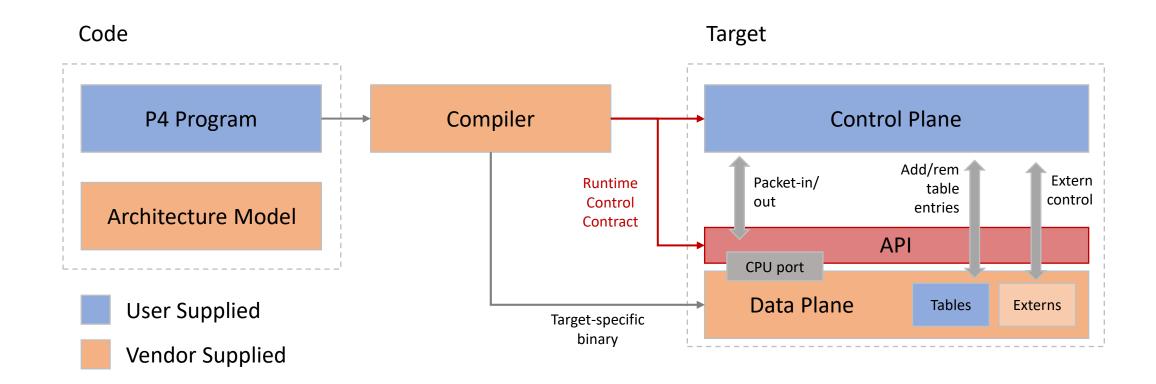
    ipv4.dstAddr:lpm

          udp.dstPort : ternary;
                                                                            - range
          ethernet.dstAddr : exact;

    udp.dstPort : range

          ethernet.srcAddr : exact;
          ethernet.etherType : ternary;
                                                                      Primitive actions
                                                                      - modify field, add to field, add,
                                                                        set field to hash index
     actions {
                                                                      - add header, remove header, copy header
                       /* permit */
          no op;
                                                                      - push, pop
          acl drop; /* reject */
                                                                      - count, meter
                                                                      - generate digest, truncate
          nhop set; /* policy-based routing */
                                                                      - resubmit, recirculate
                                                                      - clone *
                                                                      - no op, drop
```

Programming a target



P4Runtime

- Framework for runtime control of P4-defined data planes
 - Open-source API and a server implementation
- P4 program-independent
 - API doesn't change with the P4 program
- Enables field-reconfigurability
 - Ability to push new P4 program without recompiling the software stack of target switches
- API based on protobuf (serialization) and gRPC (cliente/server transport)
 - Makes it easy to implement a P4Runtime cliente/server by auto-generating code for different languages
- P4Info as a contract between control and data plane
 - Generated by P4 compiler
 - Needed by the control plane to format the body of P4Runtime messages

P4 and P4Runtime are two different things

- P4
 - Programming language used to define how a switch processes packets
 - Specifies the switch pipeline
 - Which fields does it match upon?
 - What actions does it perform on the packets?
 - In which order does it perform the matches and actions
 - Specify the behavior of an existing device
 - Specify a logical abstraction for the device

• P4Runtime

- An API used to control switches whose behavior has already been specified in the P4 language
- Works for different types of switches
 - Fixed
 - Semi-programmable
 - Fully programmable