Unified Business Logic API and Storage Layers a.k.a. "Re-Current"

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On the origin of the term "Re-Current": http://dima.ai/static/current.pdf

Think of the following out of the box:

A data storage with the APIs and the performance comparable to Redis.

- But the amount of data to work on is unbounded!
- Because, while most of the data remains in memory, most of it lives on SSDs.

• This storage is distributed, and it never loses data.

- Each mutation to the data is first replicated to N other locations.
- \circ As long as (N+1)/2 locations are available and connected to each other, the system lives on.

The storage designed for attaching eventually consistent replicas.

- The very service is a large, high-throughput, low-latency OLTP engine.
- Behind the scenes of this OLTP is the Event Store, and CDC is included in the package.

The data storage layer is not just "Redis-like", it is generic.

Redis data structures, including the MULTI statement, are just a demo.

- The idea of a strongly linearizable order of mutations is unconditionally respected.
- The actual mutations can implement arbitrary complex logic.
 - As long as it "evaluates" in a short, provably bounded time.

Strongly consistent mutations and queries go through the leader.

- Instead of Lua, a statically typed high-performance language is used.
- The storage is transactional by design, and any programmable invariant can be respected.

Eventually consistent operations are horizontally scalable.

- Mutations and strongly consistent reads have to get their sequence numbers stamped.
- Eventually consistent read or query operations can be executed against any replica.

The implementation relies heavily on:

- The fact that a single CPU can do a lot.
 - As proven by Redis.
- The fact that distributed consensus is a solved problem.
 - As proven by etcd.
- The fact that API gateways are finally merging with schema registries.
 - From Swagger to Buf.build and various gRPC proxy implementations.
- And the fact that storage is cheap these days, while "DBs" are not.
 - There is just no pure solution on the market the combines the above.
 - Hence this slide deck!

Realistically, the data is, first and foremost, an event log.

The overall solution is the Event Store.

- An immutable, append-only event log.
- And this is why CDC and data feeds come in organically.

The bodies of CQRS commands and the schema, are first-class citizens.

- They are stored alongside other data in the same storage space.
- And they are subject to the same evolution invariants to ensure correctness.

One of the major building blocks is the Total Order Broadcast (TOB) engine.

- Existing distributed consensus solutions are slow.
 - They can handle a few cross-DC mutations per second, but not more.
- We are in the throughput game, not in the latency game.
 - It is more important to persist the data than to respond quickly.
 - And distributed data replication is slow by design.
 - So, it is acceptable to have ~200ms latency for what has to be strongly consistent.
- And the fact that storage is cheap these days, while "DBs" are not.
 - There is just no pure solution on the market the combines the above.

External interfaces are part of the system itself:

- Requests enter the system through a "frontend cluster" gateway.
 - This is essential to ensure linearizability for mutations and for strongly consistent reads.
- This gateway is extensible by design.
 - In particular, HTTP+JSON and gRPC methods are organically wrapped to CQRS commands.
 - Both the strongly consistent and the eventually consistent paths are covered.
- The end-to-end system is deployed organically, as a single logical unit.
 - AWS, GCP, or really anything Kubernetes-friendly.
 - S3 is used as the cheapest storage layer on the market, if AWS is used.
 - SaaS deployments can be liberally priced based on compute and storage costs.

Behind the scenes, we are looking at a proprietary ledger.

- CQRS commands <=> orders to execute smart contracts.
 - Linearizability <=> an immutable, append-only event log.
 - Hence the Event Store paradigm, where <u>CDC</u> and data feeds come in organically.
- The CQRS queries have exclusive access to all the data.
 - If they need to be strongly consistent, they go through the leader.
 - o If they can be eventually consistency, any replica would do.
- The schema of the data is also incremental, and stored on the ledger.
 - Along with the schema, schema evolution rules are stored as needed.
 - The end users do not notice that the very schema is evolving.
 - As long as the CQRS commands in use are kept backwards- and forward-compatible.
- Same applies to the outer-world-facing interfaces.
 - TL;DR: The CQRS queries and commands are natively mapped to HTTP+JSON / gRPC calls.

Peak strongly consistent data access is never compromised.

- Asynchronous, long-running jobs can be run separately. The lifecycle is:
 - o Implement new CQRS commands and queries to "read old, write new" data.
 - Push these changes, thus activating "read-through migration".
 - The data begins to be moved, most actively used pieces first.
 - Mark the now-read-only data as immutable.
 - This is a static change. CQRS commands that try to mutate this data will fail to build.
 - Run a batch job that converts this data into new data structures in the background.
 - Logically, one could "touch" every single bit of data in some throttled way.
 - But it's better to perform this "cold migration" off the "main linearized thread".
 - Once complete, retire the old data.
- To ensure high throughput, the runtime of each command is limited
 - o In particular, no full scans and no batch mutations of the data on the leader.
- If a new "index" has to be added to the data, the above process applies.
 - So that efficient lookups can be implemented over data structures with no up front indexes.

To present the above, we solve a few "standard" problems with it. The demos include:

• A "simple" "decentralized trusted ledger" solution.

- Handle 10M++ "Alice-pays-bob" transactions per second easily.
- With guaranteed invariant validation.
- And that's a single ledger; 100M++ is easy with sub-ledgers and cross-shard txns.

• Effectively, a "large Redis".

- Implement a wide enough subset of Redis commands.
- Show these commands to run on terabytes-large datasets.

• Effectively, a "fast Kafka".

The protocol can handle more throughput than Kafka, with stronger durability guarantees.

"Standard" SysDesign interview questions

- "Design Twitter": "feed generation" and "celebrities".
- **"Design Ticketmaster":** strongly consistent transactions.
- **"Design Telegram":** 1:1 and channels and groups.

Implementation 101

The system consists of several components:

- The Gateway / Frontend Cluster.
 - o To orchestrate sending commands in parallel to N geo-distributed backends.
 - Client libraries to interface with it (incl. time skew compensation, etc.)
 - And HTTP+JSON and/or gRPC gateways for user-facing APIs.
- The internal "Front Row Orchestrator".
 - Logically, a wrapper of over some etcd, so of N clusters one is always the leader.
 - Since our leader should handle far more RPS than etcd.
 - Implementation-wise, this component is also responsible for linearizing the input stream.
 - In a way, it's "Kafka on steroids", with one topic & one partition, but millions of TPS.
 - To accomplish this, each "message" in this event bus is just a few bytes long.
 - Thus, the front row also needs to act as a distributed in-memory storage.
 - To map cooked CQRS commands' bodies into some 64-bit hashes.
- The executor of the commands.
 - This is relatively straightforward as the commands are pre-packaged.
 - Challenging sub-problems include dynamic loading of code, schema evolution, and the "ledger-like" behavior.
 - This layer includes the "infinitely accessible" "memory canvas", with paging, LRU caching, S3 backups, etc.
- The replication protocol and "hot standby"-s.
 - o In reality, there will be more than one "single topic single partition" linearized "Kafka" stream.
 - For data replication purposes, instead of executing each command on each node, just the "execution log" can be replicated.
 - It's not a commands log, but also not the logical replication log; it's a log of traces of which parts of which command should be executed.

Links

- <u>"The Pyramid" slides</u>, from 2015, the origins of <u>Current</u>.
- An earlier version of the ideas in these slides, from ~3 years ago.
- A <u>Miro board</u> with a top-level view of everything, ~2 years ago.