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	1	2	3	4	5
A	1,3	1	1,2	2	2
B	2	2,3	3	3,1	1
C					3

### Proof 1

Theorem: In any execution of the algorithm, if a woman receives a proposal on day<sub>i</sub>, then she receives some proposal on every subsequent day until the algorithm terminates.

*Proof by induction:*

Base case:

- On day<sub>1</sub> a woman<sub>1</sub> receives a proposal from man<sub>A</sub>.

Inductive Step:

- Given a woman<sub>1</sub> who receives at least one proposal on day<sub>i</sub> from a man.
- By part 1 of the algorithm, each man that proposed to woman<sub>1</sub> must have woman<sub>1</sub> in the front of his list.
- The second part of this algorithm states that women<sub>1</sub> must accept one of those offers, on day<sub>i</sub> let's assume that's man<sub>A</sub>.
- According to the third part of the algorithm if a man is rejected, he removes the woman from his list. Since Man<sub>A</sub>'s offers was accepted, he will not remove women<sub>1</sub> from his list, thus he will propose to her the next day as well. Hence on day<sub>i+1</sub> she will receive a proposal.
- This cycle will continue every day. QED

### Proof 2

Theorem: In any execution of the algorithm, if a woman receives no proposal on day<sub>i</sub>, then she receives no proposal on any previous day<sub>j</sub>,  $1 \leq j < i$

*Proof by Contradiction:*

Assume: In any execution of the algorithm, if a woman receives no proposal on day<sub>i</sub>, then she receives a proposal on any previous day<sub>j</sub>,  $1 \leq j < i$

The assumption stated another way: In any execution of the algorithm, if a woman receives a proposal on day<sub>j</sub>, then she will not receive a proposal on any further day<sub>i</sub>,  $1 \leq j < i$

- Given a woman<sub>1</sub> who receives a proposal on day<sub>j</sub> from a man. (1)
- By proof 1, she will receive a proposal on all later days including day<sub>i</sub>.
- This contradicts the assumption. QED

### Proof 3

Theorem: In any execution of the algorithm, there is at least one woman who only receives a single proposal.

*Direct Proof:*

- ◆ The last part of the algorithm states that the process ends when all men and women are paired, that means that there are an equal number of men and women.
- ◆ It follows that if all women receive a proposal, then they will all receive only one, and at that point the process will end.
- ◆ The last part of the algorithm therefore shows us that **only on the last day** all women will get a proposal from exactly one man, otherwise the algorithm would have ended on an earlier day.
- ◆ Following this, it comes out that on the day before the last day, (at least) one of the women must not have received an offer.
- This is because if she would have received a proposal then that would mean that every woman already received a proposal making this the last day, and that can't be considering it's the second to last day.
- ◆ That means that on the second to last day there is a woman who didn't receive an offer but on the next day she did (because all women have an offer on the last day).

◆By Proof 2, she must not have received an offer any day prior to the final day.  
-Thus, this woman will only receive an offer on one day (the last day).  
-Therefore, in any execution of the algorithm, there is at least one woman who only receives a single proposal. QED