

Practical 5: Advanced Sorting Algorithms

What am I doing today?

Today's practical focuses on:

1. Implementing Mergesort from pseudo-code and comparing Mergesort to Insertion Sort for increasing input sizes
2. Implementing enhanced Mergesort
3. Compare the performance between 3 sorting algorithms

Instructions

Try all the questions. Ask for help from the demonstrators if you get stuck.

Grading: Remember if you complete the practical, add the code to your GitHub repo which needs to be submitted at the end of the course **for an extra 5%**

Quick Questions

1. Mergesort guarantees to sort an array in linearithmic time, regardless of the input:
 - A. Linear time
 - B. Quadratic time
 - C. Linearithmic time
 - D. Logarithmic time

2. The main disadvantage of MergeSort is:
 - A. It is difficult to implement
 - B. It uses extra space in proportion to the size of the input**
 - C. It is an unstable sort
 - D. None of the above

3. Merge sort makes use of which common algorithm strategy?
 - A. Dynamic Programming
 - B. Branch-and-bound
 - C. Greedy approach
 - D. Divide and conquer**

4. Which sorting algorithm will take the least time when all elements of the input array are identical?
 - A. Insertion Sort
 - B. MergeSort
 - C. Selection Sort
 - D. Bogo Sort**

5. Which sorting algorithm should you use when the order of input is not known?
 - A. Mergesort**
 - B. Insertion sort
 - C. Selection sort
 - D. Shell sort

Algorithmic Development

Part 1

Let's start by implementing a version of the Merge Sort algorithm (using the pseudo-code below) that sort values in ascending order.

***please add this function to your class of sorts that you created last week.**

First implement mergeSort:

```
function mergeSort (int[] a){  
  
    N = array.length;  
  
    //base case  
    if (n == 1){  
        return array;  
    }  
  
    //create left and right sub-arrays  
    left = mergeSort(left);  
    right = mergeSort(right);  
  
    mergeArray = merge(left, right);  
  
    return mergedArray;  
}
```

Second implement the recursive merge:

```

function merge (int[] a, int[] b){

    //repeat while both arrays have elements in them
    while (a.notEmpty() && b.notEmpty()){

        //if element in 1st array is <= 1st element in 2nd array
        if (a.firstElement <= b.firstElement){
            S.insertLast(a.removeFirst());
        } else if (b.firstElement <= a.firstElement){
            S.insertLast(b.removeFirst());
        }

        //when while loop ends
        If (a.notEmpty()){
            //add remaining elements in a to S
        } else if (b.notEmpty()){
            //add remaining elements in b to S
        }

    }

    return S;
}

```

Part 2

Write a second version of MergeSort that implements the two improvements to mergesort that we covered in the lecture:

- 1) add a cutoff for small subarrays and use insertion sort (written last time) to handle them. We can improve most recursive algorithms by handling small cases differently.

Pseudo-code:

```

if (hi <= lo + CUTOFF) {
    insertionSort(dst, lo, hi);
    return;
}

```

- 2) test whether the array is already in order. We can reduce the running time to be linear for arrays that are already in order by adding a test to skip call to merge() if $a[mid]$ is less than or equal to $a[mid+1]$. **In other words, if the last element in the first sorted array is less than or equal to the first element in the second sorted array then you can just add the entire second array in without the need for comparisons.** With this change, we still do all the recursive calls, but the running time for any sorted subarray is linear.

Part 3

Compare the performance of Insertion Sort, MergeSort and MergeSortEnhanced on a range of inputs (N= 10, 1000, 10000, 100000 etc.).

Solution note: Switching to insertion sort for small subarrays will improve the running time of a typical mergesort implementation by 10 to 15 percent.

(Nano Seconds)

	InsertionSort	MergeSort	MergeSortEnhanced
N = 10	4100	14200	9600
N = 100	106400	112600	120000
N = 1000	5117000	1112500	388800
N = 10000	25522100	4618500	4480300