

# 5. Минимални Крайни Автомати

Def Релация на Myhill-Nerode

Нека  $A = (\Sigma, Q, I, \delta, F)$  е дет. крайн автомат

Тогаваш дет.  $\equiv_A \subseteq Q \times Q$  така:

$$(\forall p, q \in Q) [p \equiv_A q \iff (\forall w \in \Sigma^*) ((\exists f \in F)(p \xrightarrow{w} f) \iff (\exists f \in F)(q \xrightarrow{w} f))]$$

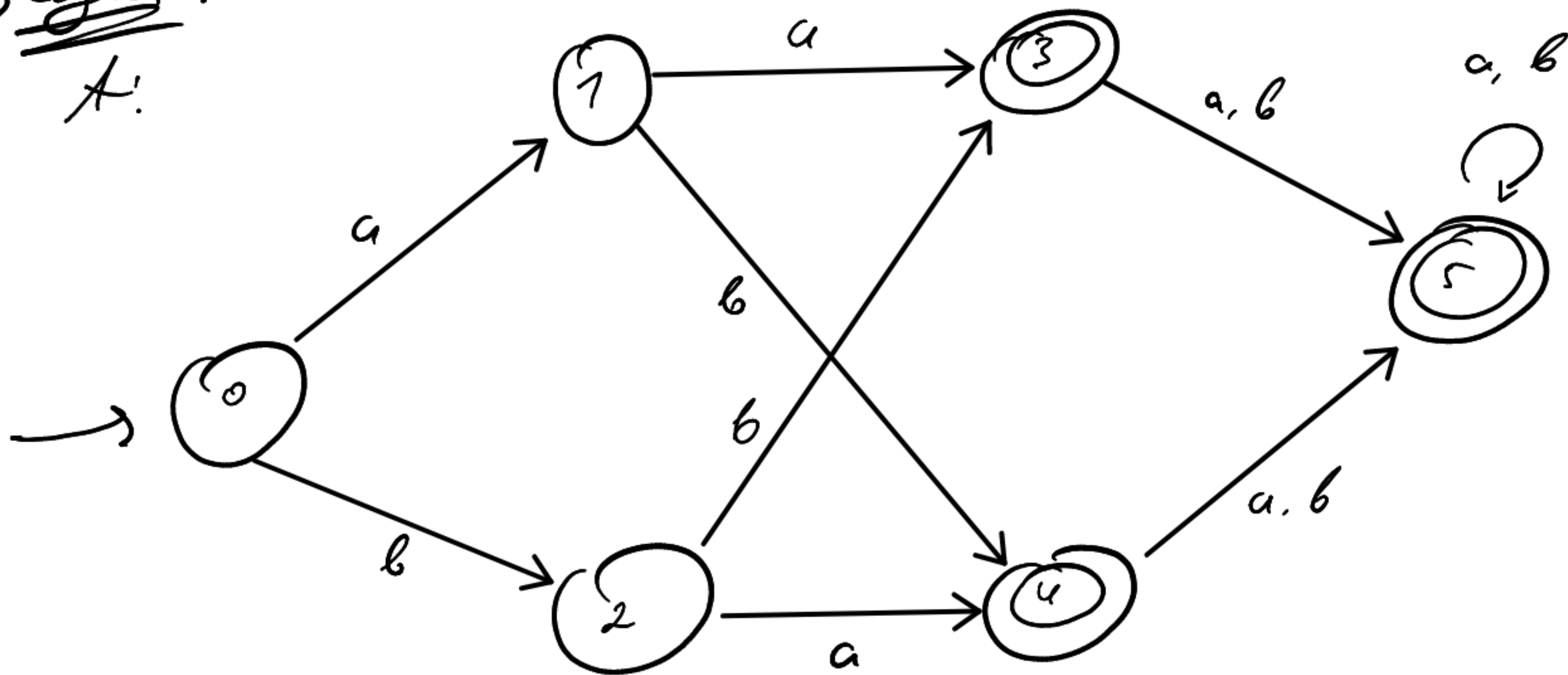
Тогаваш автоматът

$$A_m = (\Sigma, Q/\equiv_A, [q], \Delta, F/\equiv_A)$$

е мин. пот. дет. крайн. авт. с  $L(A_m) = L(A)$

Задача:

A:



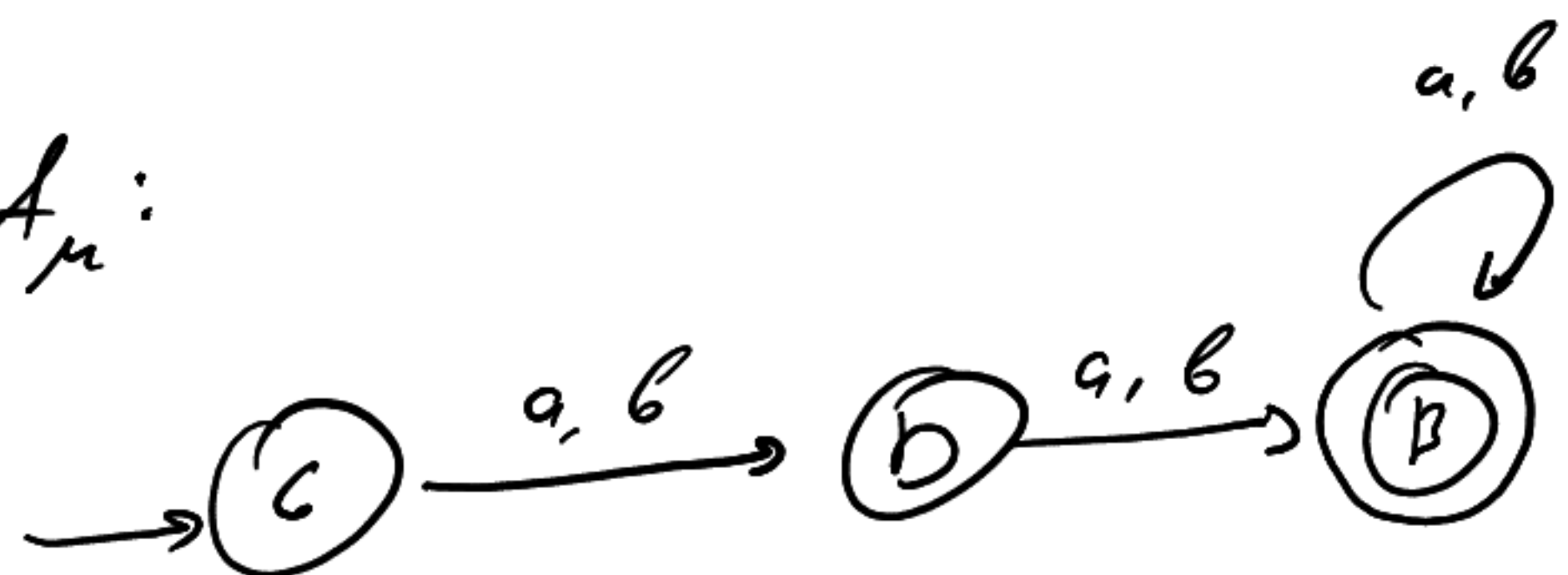
A	a	b
0	A	A
1	B	B
2	B	B

B	a	b
3	B	B
4	B	B
5	B	B

C	a	b
0	D	D

D	a	b
1	B	B
2	B	B

$A_m$ :



Задан Да се построи МДКА с език  $L$ :

$$a) \Sigma = \{a, b, c\}$$

$$L = \{w \in \Sigma^* \mid w \text{ не завършва на } abc \\ \text{и започва с } aaa\}$$

$$b) \Sigma = \{a, b, c, \dots, z\}$$

$$L = \Sigma^* \cdot \{hers, his, he\}$$

$$b) \Sigma = \{a, b, c\}$$

$$L = \Sigma^* \setminus (\{abc\}^* \cup (\{a, c\} \Sigma^* \cap \Sigma^* \{b\}))$$

Решение

$$a) L = \{w \in \Sigma^* \mid w \text{ не зав. на } 'abc'\} \cap \{w \in \Sigma^* \mid w \text{ започва с } 'aaa'\}$$

$$= \Sigma^* \setminus (\Sigma^* \cdot \{abc\}) \cap \{aaa\} \cdot \Sigma^*$$

$$A_{\Sigma^*}: L(A_{\Sigma^*}) = \Sigma^*$$



$$A_{abc}: L(A_{abc}) = \{abc\}$$

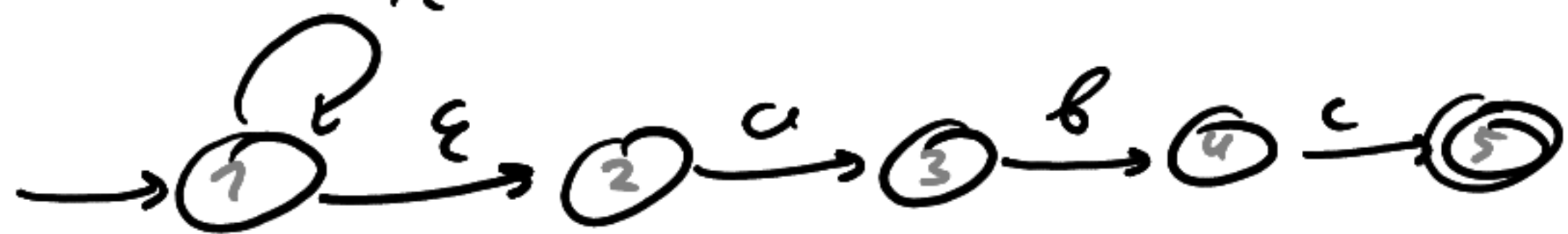


$$A_{aaa}: L(A_{aaa}) = \{aaa\}$$

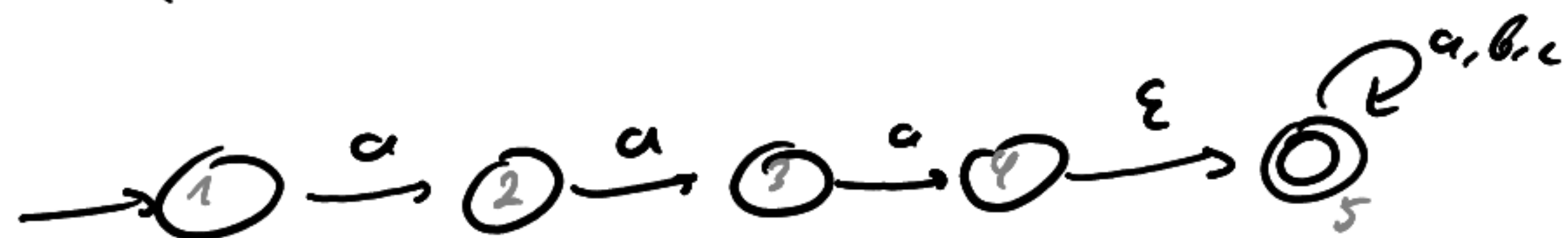




$$A_1: L(A_1) = \Sigma^* \cdot \{abc\} = L(A_{\Sigma^*}) \circ L(A_{abc})$$



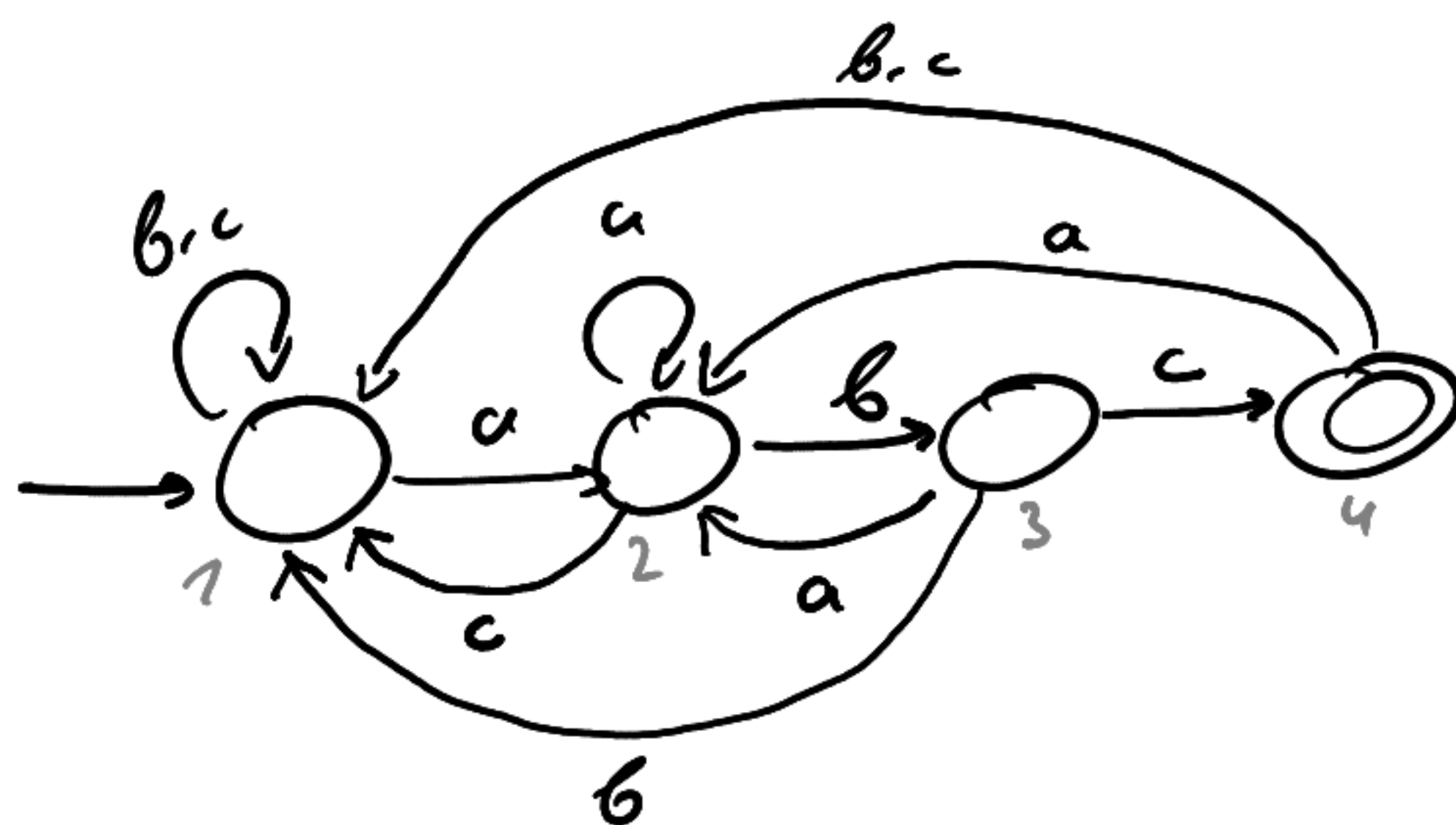
$$A_2: L(A_2) = \{aaaa\} \cdot \Sigma^* = L(A_{aaaa}) \circ L(A_{\Sigma^*})$$



Ще като конструираме за тези две регулярности, ще генерираме  $A_1$  и  $A_2$

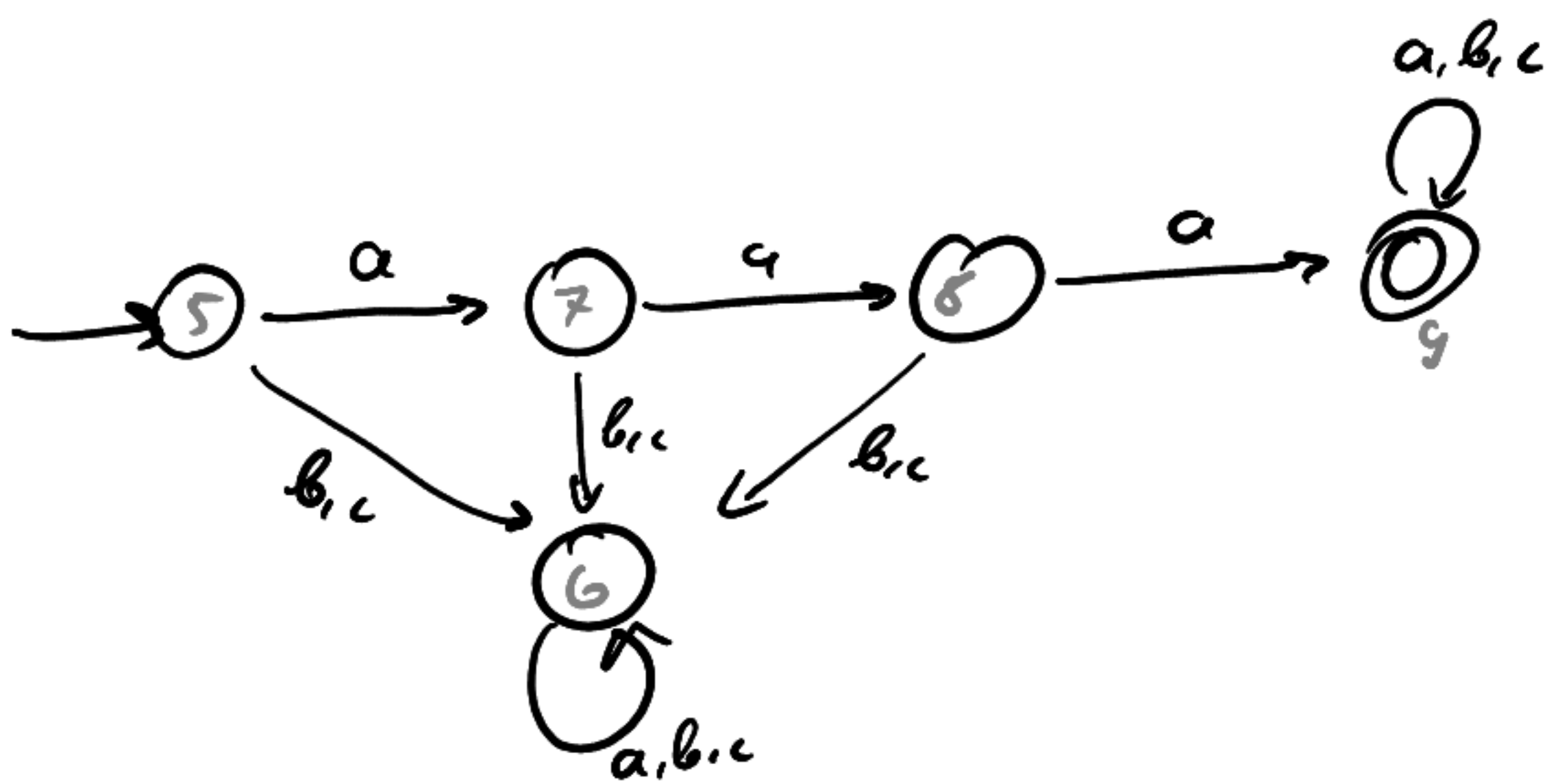
$$A_1^D:$$

		a	b	c
1:	{1, 2}	1, 3, 2 <sup>2</sup>	1, 2 <sup>1</sup>	1, 2 <sup>1</sup>
2:	{1, 2, 3}	1, 3, 2 <sup>2</sup>	1, 4, 2 <sup>3</sup>	1, 2 <sup>1</sup>
3:	{1, 2, 4}	1, 3, 2 <sup>2</sup>	1, 2 <sup>1</sup>	1, 5, 2 <sup>4</sup>
4:	{1, 2, 5}	1, 3, 2 <sup>2</sup>	1, 2 <sup>1</sup>	1, 2 <sup>1</sup>

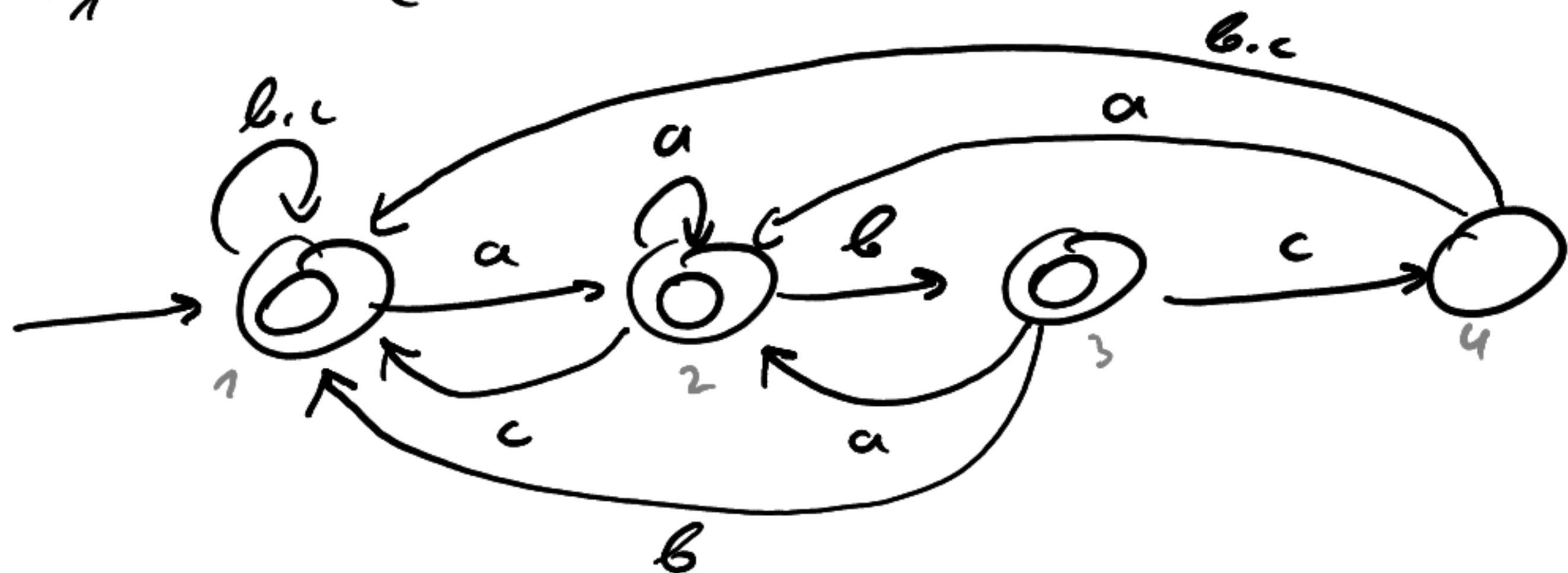


$$A_2^D:$$

		a	b	c
5:	{1}	2 <sup>7</sup>	$\emptyset^6$	$\emptyset^6$
6:	$\emptyset$	$\emptyset^6$	$\emptyset^6$	$\emptyset^6$
7:	{2}	3 <sup>6</sup>	$\emptyset^6$	$\emptyset^6$
8:	{3}	4, 5 <sup>9</sup>	$\emptyset^6$	$\emptyset^6$
9:	{4, 5}	4, 5 <sup>9</sup>	4, 5 <sup>9</sup>	4, 5 <sup>9</sup>

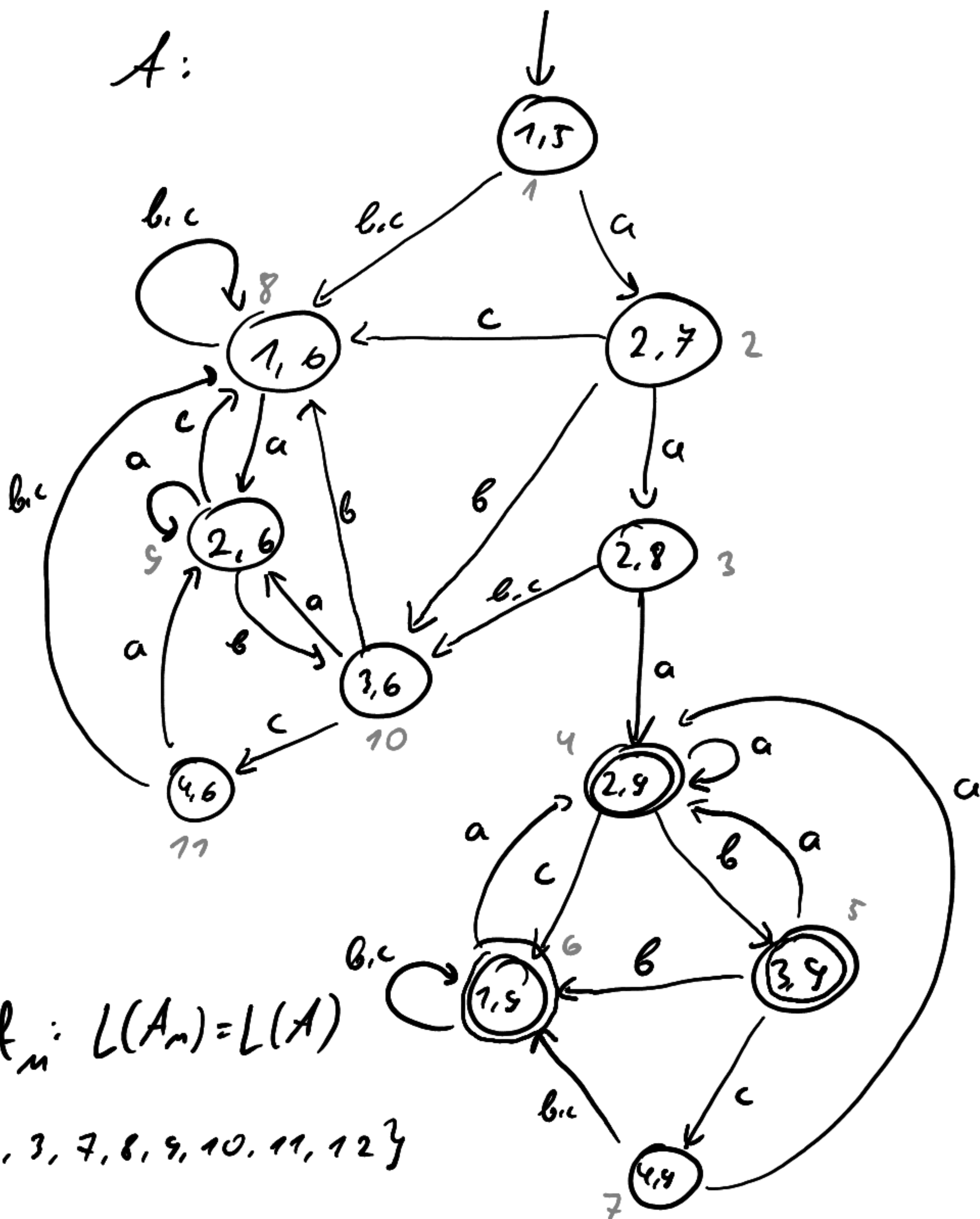


$$\overline{A_1^D}: L(\overline{A_1^D}) = \overline{L(A_1^D)} = \Sigma^* \setminus L(A_1^D)$$



Сеза остатак го покривам  $A: L(A) = L(\overline{A_1^D}) \cap L(A_2^D)$

	a	b	c
→ (1, 5)	(2, 7)	(1, 6)	(1, 6)
(1, 6)	(2, 6)	(1, 6)	(1, 6)
(2, 7)	(2, 8)	(3, 6)	(1, 6)
(2, 8)	(2, 9)	(3, 6)	(1, 6)
(3, 6)	(2, 6)	(1, 6)	(4, 6)
⊙ (2, 9)	(2, 9)	(3, 9)	(1, 9)
(2, 6)	(2, 6)	(3, 6)	(1, 6)
(4, 6)	(2, 6)	(1, 6)	(1, 6)
⊙ (3, 9)	(2, 9)	(1, 9)	(4, 9)
⊙ (1, 9)	(2, 9)	(1, 9)	(1, 9)
(4, 9)	(2, 9)	(1, 9)	(1, 9)



У сеза минимизација :  $A_m: L(A_m) = L(A)$

$$\underline{A} = \{4, 5, 6\} \quad \underline{B} = \{1, 2, 3, 7, 8, 9, 10, 11, 12\}$$

A	a	b	c
4	A	A	A
5	A	A	B
6	A	A	A

B	1	2	3	7	8	9	10	11
a	B	B	A	A	B	B	B	B
b	B	B	B	A	B	B	B	B
c	B	B	B	A	B	B	B	B

$$\underline{C} = \{4, 6\}, \underline{D} = \{5\}, \underline{E} = \{1, 2, 8, 9, 10, 11\}, \underline{F} = \{3\}, \underline{G} = \{7\}$$

C	a	b	c
4	C	D	C
6	C	C	C

E	a	b	c
1	E	E	E
2	F	E	E
8	E	E	E
9	E	E	E
10	E	E	E
11	E	E	E

$$\underline{H} = \{1, 8, 9, 10, 11\}, \underline{I} = \{2\}, \underline{J} = \{4\}, \underline{K} = \{6\}$$

H	a	b	c
1	I	K	H
8	H	K	H
9	H	K	H
10	H	K	H
11	H	K	H

$$L = \{1\}$$

$$M = \{8, 9, 10, 11\}$$

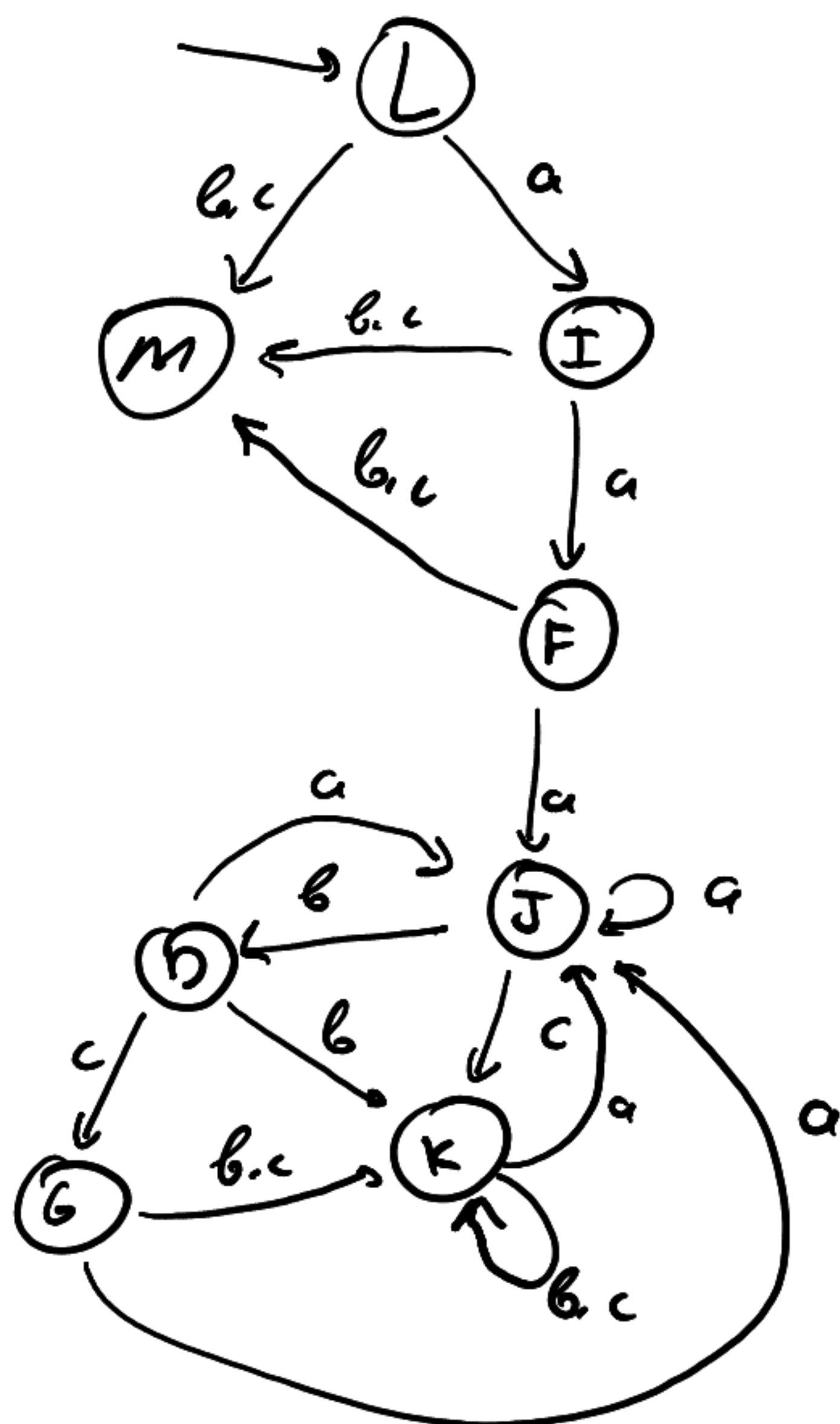
M	a	b	c
8	M	M	M
9	M	M	M
10	M	M	M
11	M	M	M

крај



		a	b	c
⊙ D		J	K	G
F		J	M	M
G		J	K	K
I		F	M	M
⊙ J		J	D	K
⊙ K		J	K	K
→ L		I	M	M
M		M	M	M

$A_m$ :



$A_m$  е минимален тотален детерминиран краен автомат с език  $L$ .

### Втори начин

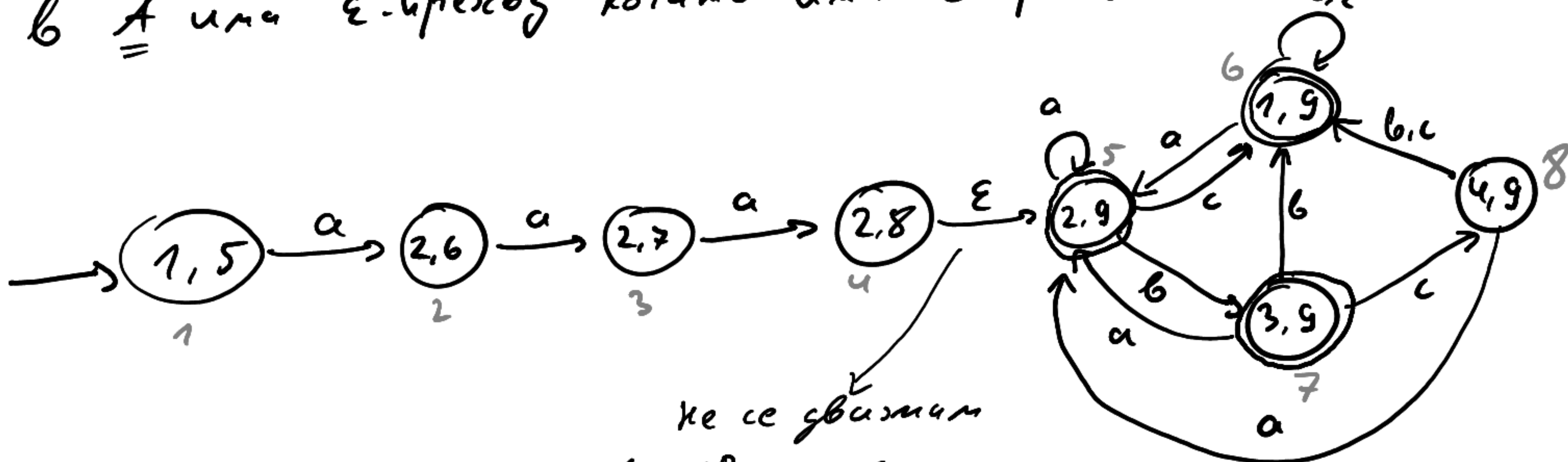
Можем да не детерминираме  $A_2$  прези сегашето (както го направихме и на дъската на упражненията)

$$A: L(A) = L(\overline{A_1^D}) \cap L(A_2)$$

Сегашето на недетерминирани автомати е гадно

!  $\underline{b \neq a}$  има преход с дъхва  $\underline{a}$  т.с.т.к. и  $\underline{b}$  дъхата автомата има преход с тази дъхва.

$\underline{b \neq a}$  има  $\epsilon$ -преход когато има  $\epsilon$ -преход в един от двата автомата



не се свързват  
в  $A_1^D$ , правим  
 $\epsilon$ -преход в  $A_2$

Детерминизираме:

$$A^D: L(A^D) = L(A)$$

	a	b	c
1 → {1}	2	∅	∅
2 {2}	3	∅	∅
3 ∅	∅	∅	∅
4 {3}	4,5	∅	∅
5 ⊙ {4,5}	5	7	6
6 ⊙ {5}	5	7	6
7 ⊙ {6}	5	6	6
8 ⊙ {7}	5	6	8
9 {8}	5	6	6

преименуваме →

	a	b	c
→ 1	2	3	3
2	4	3	3
3	3	3	3
4	5	3	3
⑤	6	8	7
⑥	6	8	7
⑦	6	7	7
⑧	6	7	9
9	6	7	7

и Максимизираме

$$A = \{1, 2, 3, 4, 9\} \quad B = \{5, 6, 7, 8\}$$

A	a	b	c
1	A	A	A
2	A	A	A
3	A	A	A
4	B	A	A
9	B	B	B

B	a	b	c
5	B	B	B
6	B	B	B
7	B	B	B
8	B	B	A

$$C = \{1, 2, 3\} \quad D = \{4\} \quad E = \{9\} \quad F = \{5, 6, 7\} \quad G = \{8\}$$

C	a	b	c
1	C	C	C
2	D	C	C
3	C	C	C

F	a	b	c
5	F	G	F
6	F	G	F
7	F	F	F

$$H = \{1, 3\}, \quad I = \{2\}, \quad J = \{5, 6\}, \quad K = \{7\}$$

H	a	b	c
1	I	H	H
3	D	H	H

J	a	b	c
5	J	G	H
6	J	G	H

$$L = \{1\} \\ M = \{3\}$$

J	a	b	c
5	J	G	K
6	J	G	K

и се рагура



Получаем следующие автоматы

$A_m$	a	b	c
D	J	M	M
E	J	K	K
G	J	K	E
I	D	M	M
J	J	G	K
K	J	K	K
L	I	M	M
M	M	M	M

