

## 1. Condition Variables Warm-up Challenge: Eat cookies fast!

Meanwhile in a Parallel Universe ...

Two threads viciously eat cookies but are blocked on a c.v. ...



```
01 int jar = 0;
02 pthread_mutex_t m = PTHREAD_MUTEX_INITIALIZER;
03 pthread_cond_t cv1 = PTHREAD_COND_INITIALIZER;
04
05 void* cookie_eater(void*arg) {
06     char* name = (char*) arg;
07     while(game_running) {
08         while(jar == 0) {
09             printf("%s nap time\n", name);
10             ?
11         }
12         jar--;
13         printf("%s eats! %d remain\n", name, jar);
14     }
15     printf("%s is exiting...", name);
16     return NULL;
17 }
```

Complete the `add_cookies` to add cookies to the cookie jar  
(Pretend cookie jar has  $\infty$  capacity)

```
18 void add_cookies(int add) {
19     assert(add > 0);
20
21
22
23
24
```

2. What must be locked before calling `p_cond_wait` ? \_\_\_\_\_

3. You wake a thread blocked inside a condition variable but it does not return from `p_cond_wait`. Why?

Another thread still \_\_\_\_\_

The blocked thread will continue when \_\_\_\_\_

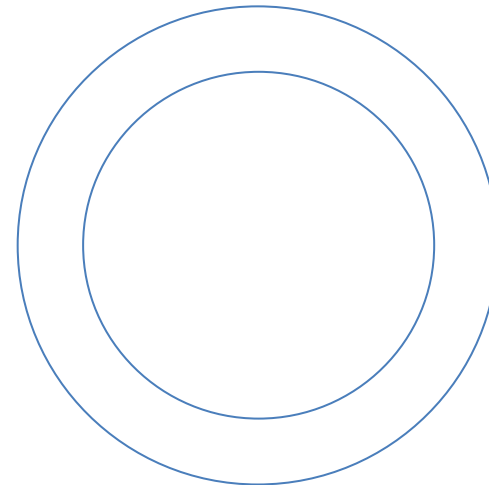
4. How do I use counting semaphores?

`sem_init`

`sem_wait`

`sem_post`

5. What is a fixed ring buffer? Why would I use it?



## 6. Producer Consumer Case Study:

Use counting semaphores to implement a fixed ring buffer

```
pthread_mutex_t m;

// (Not OSX!)

sem_t _____

void init() {
    sem_init(_____, 0, _____);
    sem_init(_____, 0, _____);
    pthread_mutex_init( &m , NULL);
}

void sync_enqueue(work_t *work) {

}

work_t* sync_dequeue() {
```

## 7. Quick quiz

i) How many threads can be executing line 8 or 14 at a time? Why?

ii) What have I made? (Missing code? + Better function names?)

```
01 pthread_mutex_t m = PTHREAD_MUTEX_INITIALIZER;
02 pthread_cond_t cv1 = PTHREAD_COND_INITIALIZER;
03 int  mystery = 5;
04
05 void A?() { // Waits if count would become -ve
06     p_m_lock(&m)
07     while(mystery == 0) p_cond_wait(&cv1, &m);
08     mystery --;
09     p_m_unlock(&m);
10 }
11 void B?() {
12     p_m_lock(&m);
13     mystery ++;
14     if( _____ ) p_cond_broadcast(&cv1);
15     p_m_unlock(&m);
16 }
```

### CRITICAL SECTION PROBLEM

Today's prizes !

Mutual Exclusion

Progress

Bounded Wait

*Candidates – remember you will be called more than once!*

Candidate #3. Change the sequence order

<b>raise A flag</b> <b>wait until B flag is down</b> <i>Critical Section code here</i> <b>lower A flag</b>	<b>raise B flag</b> <b>wait until A flag is down</b> <i>Critical Section code here</i> <b>lower B flag</b>
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Problems?

Candidate #4. Try a single turn-based shared variable.

turn=1

<b>while( turn == 2) { }</b> <i>Critical Section code here</i> <b>turn = 2</b>	<b>while( turn == 1) { }</b> <i>Critical Section code here</i> <b>turn = 1</b>
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Problems?