

The strategy we used was greedy. We had two different greedy approaches. One was a simple greedy approach which moves across each vertex assigning it the first valid coloring. The next greedy approach, known as the Welsh-Powell Algorithm, colors vertices from highest degree to lowest degree. The runtime complexity of our simple algorithm was  $O(V + E)$  while the runtime of our Welsh-Powell implementation was  $O(V \log V + E)$ . Both of these ran in polynomial time and were able to handle very large input sizes (1000 vertices and 500,000 edges).

Test cases:		
test	result	runtime
test_large_connected	passed	.6069031s
test_large_connected2	passed	1.0258840s
test_nonoptimal	passed	.0809894s
test_petersen_graph	passed	.1010330s
test_small_connected	passed	.0859973s

Our large connected 2 test for example contains 288,420 edges and our approximation code ran in only 1.02 seconds.

```
def color_highest_degree(G):
    Colors = {}
    pq = PriorityQueue()
    add all v in G to pq where key is degree

    while pq isn't empty:
        v = pq.get()
        color = assign_color(G, v, colors)
        colors[v] = color

    return len(set(values of colors))

def assign_color(G, v, colors):
    neighbor_colors = set()
    for u in G[v]:
        if u in colors:
            neighbor_colors.add(colors[u])

    color = 1
    while color in neighbor_colors:
        color += 1
    return color
```

Overall time complexity:  $O(V \log V + E)$

Above is the pseudo code for the Welsh-Powell algorithm. It runs a while loop over all the nodes which have been added to a priority queue. Each iteration it calls `pq.get()` which runs in  $O(\log n)$  time. The loop combined with the get runs in  $O(V \log V)$ . The `assign_color` function in the worst case will take in a vertex whose degree is the same as the number of edges in the graph. This runs in  $O(E)$  time. Overall this approximation is polynomial.