

ECS 20 — Lecture 7 — Fall 2013 — 18 Oct 2013

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Today: o **Sets** and
o Writing sets
o Some operations on sets
o Some important sets for math and CS

$S = \{\text{dog}, \text{cat}\}$. Order we write elements (points) in a set doesn't matter.
 $= \{\text{cat}, \text{dog}\}$. Repetitions don't matter, either:
 $= \{\text{cat}, \text{dog}, \text{cat}\}$.

$$S = \emptyset$$

Often many ways to write a set

$$\begin{aligned} A &= \{2i+1: i \in \mathbf{Z}\} \\ &= \{\dots, -5, -3, -1, 1, 3, 5, \dots\} \\ &= \{x: x \text{ is an odd integer}\} \\ &= \{n: n \in \mathbf{Z} \text{ and } \neg (\exists j \in \mathbf{Z})(2j=n)\} \end{aligned}$$

Or

Let P be the set of prime numbers.
 $P = \{n: n \text{ is a prime number}\}$
 $P = \{n \in \mathbf{N}: i \mid n \implies i \in \{-n, -1, 1, n\}\}$
 $P = \{2, 3, 5, 7, 11, \dots\}$

Can a set contain a set? **Sure.**
Can a set contain the emptyset? **Sure**
 $S = \{\mathbf{N}, \{2, 3\}, [0, 1]\}$.
 $S = \{\emptyset, \{\emptyset\}, \{\{\emptyset\}\}\}$

Naïve set theory, where we describe sets with natural language, can sometime run into trouble.

Can a set contain itself? **No**

Can a set contain “everything”? **No**

Russell's paradox: Let $R = \{x \mid x \notin x\}$ *Problem:* is $R \in R$ iff $R \notin R$?

Def: $S = T$ iff $x \in S \leftrightarrow x \in T$

Def: $S \subseteq T$ if $x \in S \rightarrow x \in T$

$\{a, b\} \subseteq \{a, b, c\}$ YES
 $\{a, b\} \subseteq \{a, b\}$ YES

$\{a, b\} \subseteq \{a, d, e\}$ NO
 $\emptyset \subseteq \{a, b, c\}$ YES (explain)
 $\{\emptyset\} \subseteq \{a, b, c\}$ NO
 $\{\emptyset\} \subseteq \{\{\emptyset\}\}$ YES
T/F: for all $S, \emptyset \subseteq S$: True

Some important sets for math and computer science
 $\mathbf{N} = \{1, 2, 3, \dots\}$ // some books include 0, some don't
 $\mathbf{R} = \{x: x \text{ is a real number}\}$
 $\mathbf{Z} = \{\dots, -2, -1, 0, 1, 2, \dots\}$
 $\mathbf{Q} = \{m/n: m, n \in \mathbf{Z}, n \neq 0\}$

$[a..b]$ integers between a and b , inclusive.
 $[a, b]$ reals between a and b , inclusive

$[1..N] = \{1, 2, \dots, N\}$
 $[N] = \mathbf{Z}_N = \{0, 1, \dots, N-1\}$

Sometimes sets come with operations on them, these operations satisfying simple algebraic properties.

Example:

Group This is a set A and an operation $*$ where:

- 1) $(x*y)*z = x*(y*z)$;
- 2) there exists an element 1 in A such that $x*1 = 1*x = x$;
- 3) for every element x there is an element y such that $x*y = 1 = y*x$

But let me emphasize that a set, all by itself, does **not** have operations defined on its elements.

- Ask questions about making \mathbf{N} , \mathbf{R} , \mathbf{Z} into groups.
- Later: ask questions about making BITS, BYTES, WORDS32 into a group, by either XOR and Modular addition operation

For computers, important sets correspond to those things that our architectures natively manipulate:

$\text{BITS} = \{0, 1\}$
 $\text{BYTES} = \{0, 1\}^8$ Signed, unsigned
 $\text{WORDS32} = \{0, 1\}^{32}$ Signed, unsigned
 $\text{WORDS64} = \{0, 1\}^{64}$ Signed, unsigned
 $\text{IEEEFLOAT64} = \{0, 1\}^{64}$ = representing exponents -1022 .. 1023 (about 16 digits of accuracy)
 Weirder than you may think

- **sign, significand** (=coefficient), **exponent** $(-1)^{\text{sign}} \cdot \text{significand} \cdot 2^{\text{exponent}}$ IEEE
- $+\infty$ and $-\infty$
- NaN (of various kinds)
- Zero can be +0 or -0



[William Kahan](#). A primary architect of the [IEEE 754](#) floating-point standard

Or particular language:

The set of all valid C programs

The set of valid URLs

The set of valid http programs

$|S|$ = the number of element in S if S is finite, ∞ otherwise

$A = \{\{a\}, b, \emptyset\}$. $|A| = 3$

$A = \{\emptyset, \{\emptyset\}, \{\{\emptyset\}\}$ $|A|=3$.

UNION

$A \cup B = \{x: x \in A \text{ or } x \in B\}$ // not really very rigorous:

$\{x: \text{blah}\}$ before the colon, the universe should be clear, with what comes after a narrowing of that. Books don't all stick to this, but that's what learned!

$\{\text{dog}, \text{cat}\} \cup \{\text{cat}, \text{fish}\}$

$\{a, b\} \cup \{\emptyset, a\} = \{a, b, \emptyset\}$

$A \cup \emptyset = A$

Can union up infinitely many things

$\bigcup_{n \in \mathbb{N}} [0, n] = \mathbb{R}$

$\bigcup_{i \in \mathbb{N}} A_i$ eg, $\bigcup_{i \in \mathbb{N}} \{2i-1\} = \text{the set of odd positive number}$

$\bigcup_{a \in \mathbb{N}} \{a^i: i \in \mathbb{N}\}$

"powers of integers"

INTERSECTION

$\{1, 2, 3\} \cap \{2, 5, 8\} = \{2\}$

$\{1, 2, 3\} \cap \{4, 5, 8\} = \emptyset$

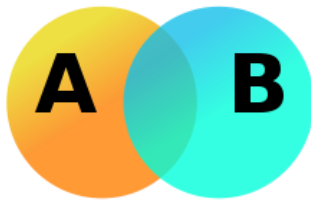
$\{1, 2, 3\} \cap \emptyset = \emptyset$ **True/False**

$S \cap \emptyset = \emptyset$ **True/False**

Can intersect infinitely many things, too:

$$\bigcap_{n \in \mathbb{N}} [0, n] = \{0\}$$

Venn Diagrams



Set Difference

$$A \setminus B \quad \text{or} \quad A - B$$

Symmetric Difference

$$A \oplus B$$

Algebra of sets

Commutative laws:

- $A \cup B = B \cup A$
- $A \cap B = B \cap A$

Associative laws:

- $(A \cup B) \cup C = A \cup (B \cup C)$
- $(A \cap B) \cap C = A \cap (B \cap C)$

Distributive laws:

- $A \cup (B \cap C) = (A \cup B) \cap (A \cup C)$
- $A \cap (B \cup C) = (A \cap B) \cup (A \cap C)$

Proof: $x \in A \cup (B \cap C)$ means

$$\begin{aligned} & (x \in A) \vee ((x \in B) \wedge (x \in C)) \quad \text{But } P \vee (Q \wedge R) = (P \vee Q) \wedge (P \vee R). \text{ So} \\ & = ((x \in A) \vee (x \in B)) \wedge ((x \in A) \vee (x \in C)) \\ & = (A \cup B) \cap (A \cup C) \end{aligned}$$

Identity laws:

- $A \cup \emptyset = A$
- $A \cap U = A$

Complement laws:

- $A \cup A^C = U$
- $A \cap A^C = \emptyset$
-

idempotent laws:

- $A \cup A = A$
- $A \cap A = A$

domination laws:

- $A \cup U = U$
- $A \cap \emptyset = \emptyset$

absorption laws:

- $A \cup (A \cap B) = A$
- $A \cap (A \cup B) = A$

double complement or involution law:

- $(A^C)^C = A$

complement laws for the universal set and the empty set:

- $\emptyset^C = U$
- $U^C = \emptyset$

De Morgan's laws:

- $(A \cup B)^C = A^C \cap B^C$
- $(A \cap B)^C = A^C \cup B^C$

Proof (of first claim): $x \in (A \cup B)^c$

iff $\neg(x \in (A \cup B))$

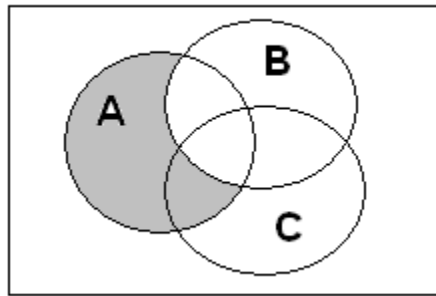
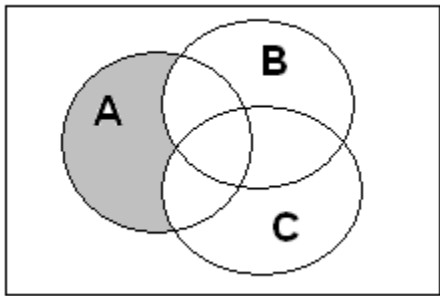
iff $\neg(x \in A \vee x \in B)$

iff $\neg(x \in A) \wedge \neg(x \in B)$

iff $x \in A^c \wedge x \in B^c$

Be careful!!

$$(A \setminus B) \setminus C \stackrel{?}{=} A \setminus (B \setminus C)$$



Cartesian Product (= Cross product) ← Did get here, continue next time

$$A \times B = \{(a,b): A \in A, B \in B\}$$

\mathbf{R}^2 points in the plane

An array of chessmen might be represented by BYTES⁶⁴

Power Set

\mathcal{P} – Power set operator, unary operator (takes 1 input). $\mathcal{P}(x)$ is the “set of all subsets of x”

$$\mathcal{P}(X) = \{A: A \subseteq X\}$$

Example: $X = \{a, b, c\}$

Example:

Variant notation: $\mathcal{P}(X) = 2^X$

Notation is suggestive of size –

For X finite, $|\mathcal{P}(X)| = 2^{|X|}$

...