

## CS311H: Discrete Mathematics

### Mathematical Induction

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## Review: Induction and Strong Induction

- ▶ How does one prove something by induction?
- ▶ What is the inductive hypothesis?
- ▶ What is the difference between regular and strong induction?
- ▶ Is strong induction more powerful than standard induction?

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## Example

- ▶ For  $n \geq 1$ , prove there exist natural numbers  $a, b$  such that:

$$5^n = a^2 + b^2$$

- ▶ Hint:  $5^{n+1} = 5^2 \cdot 5^{n-1}$

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## Matchstick Example

- ▶ **The Matchstick game:** There are two piles with same number of matches initially
- ▶ Two players take turns removing any positive number of matches from one of the two piles
- ▶ Player who removes the last match wins the game
- ▶ **Prove:** Second player always has a winning strategy.

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## Matchstick Proof

- ▶  **$P(n)$ :** Player 2 has winning strategy if initially  $n$  matches in each pile
- ▶ **Base case:**
- ▶ **Induction:** Assume  $\forall j. 1 \leq j \leq k \rightarrow P(j)$ ; show  $P(k+1)$
- ▶ **Inductive hypothesis:**
- ▶ Prove Player 2 wins if each pile contains  $k+1$  matches

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## Matchstick Proof, cont.

- ▶ **Case 1:** Player 1 takes  $k+1$  matches from one of the piles.
- ▶ **What is winning strategy for player 2?**
- ▶ **Case 2:** Player 1 takes  $r$  matches from one pile, where  $1 \leq r \leq k$
- ▶
- ▶

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## Recursive Definitions

- ▶ Should be familiar with recursive functions from programming:

```
public int fact(int n) {  
    if(n <= 1) return 1;  
    return n * fact(n - 1);  
}
```

- ▶ **Recursive definitions** are also used in math for defining sets, functions, sequences etc.

## Recursive Definitions in Math

- ▶ Consider the following sequence:

$$1, 3, 9, 27, 81, \dots$$

- ▶ This sequence can be defined **recursively** as follows:

$$\begin{aligned} a_0 &= 1 \\ a_n &= 3 \cdot a_{n-1} \end{aligned}$$

- ▶ First part called **base case**; second part called **recursive step**

## Recursively Defined Functions

- ▶ Just like sequences, functions can also be defined recursively

- ▶ **Example:**

$$\begin{aligned} f(0) &= 3 \\ f(n+1) &= 2f(n) + 3 \quad (n \geq 1) \end{aligned}$$

- ▶ What is  $f(1)$ ?
- ▶ What is  $f(2)$ ?
- ▶ What is  $f(3)$ ?

## Recursive Definition Examples

- ▶ Consider  $f(n) = 2n + 1$  where  $n$  is non-negative integer
- ▶ What's a recursive definition for  $f$ ?
- ▶ Consider the sequence  $1, 4, 9, 16, \dots$
- ▶ What is a recursive definition for this sequence?
- ▶ Recursive definition of function defined as  $f(n) = \sum_{i=1}^n i$ ?

## Recursive Definitions of Important Functions

- ▶ Some important functions/sequences defined recursively

- ▶ **Factorial function:**

$$\begin{aligned} f(1) &= 1 \\ f(n) &= n \cdot f(n-1) \quad (n \geq 2) \end{aligned}$$

- ▶ **Fibonacci numbers:**  $1, 1, 2, 3, 5, 8, 13, 21, \dots$

$$\begin{aligned} a_1 &= 1 \\ a_2 &= 1 \\ a_n &= a_{n-1} + a_{n-2} \quad (n \geq 3) \end{aligned}$$

- ▶ Just like there can be multiple base cases in inductive proofs, there can be multiple base cases in recursive definitions

## Inductive Proofs for Recursively Defined Structures

- ▶ Recursive definitions and inductive proofs are very similar
- ▶ Natural to use induction to prove properties about recursively defined structures (sequences, functions etc.)
- ▶ Consider the recursive definition:
$$\begin{aligned} f(0) &= 1 \\ f(n) &= f(n-1) + 2 \end{aligned}$$
- ▶ Prove that  $f(n) = 2n + 1$

## Example

- ▶ Let  $f_n$  denote the  $n$ 'th element of the Fibonacci sequence
- ▶ Prove: For  $n \geq 3$ ,  $f_n > \alpha^{n-2}$  where  $\alpha = \frac{1+\sqrt{5}}{2}$
- ▶ Proof is by **strong induction** on  $n$  with two base cases
- ▶ **Intuition 1:** Definition of  $f_n$  has two base cases
- ▶ **Intuition 2:** Recursive step uses  $f_{n-1}, f_{n-2} \Rightarrow$  strong induction
- ▶ **Base case 1 (n=3):**  $f_3 = 2$ , and  $\alpha < 2$ , thus  $f_3 > \alpha$
- ▶ **Base case 2 (n=4):**  $f_4 = 3$  and  $\alpha^2 = \frac{(3+\sqrt{5})}{2} < 3$

## Example, cont.

Prove: For  $n \geq 3$ ,  $f_n > \alpha^{n-2}$  where  $\alpha = \frac{1+\sqrt{5}}{2}$

- ▶ **Inductive step:** Assuming property holds for  $f_i$  where  $3 \leq i \leq k$ , need to show  $f_{k+1} > \alpha^{k-1}$
- ▶ First, rewrite  $\alpha^{k-1}$  as  $\alpha^2 \alpha^{k-3}$
- ▶  $\alpha^2$  is equal to  $1 + \alpha$  because:

$$\alpha^2 = \left( \frac{1+\sqrt{5}}{2} \right)^2 = \frac{\sqrt{5}+3}{2} = \alpha + 1$$

- ▶ Thus,  $\alpha^{k-1} = (\alpha + 1)(\alpha^{k-3}) = \alpha^{k-2} + \alpha^{k-3}$

## Example, cont.

- ▶  $\alpha^{k-1} = \alpha^{k-2} + \alpha^{k-3}$
- ▶ By recursive definition, we know  $f_{k+1} = f_k + f_{k-1}$
- ▶ Furthermore, by inductive hypothesis:  
$$f_k > \alpha^{k-2} \quad f_{k-1} > \alpha^{k-3}$$
- ▶ Therefore,  $f_{k+1} > \alpha^{k-2} + \alpha^{k-3} = \alpha^{k-1}$

□

## Recursively Defined Sets and Structures

- ▶ We can also define sets and other data structures recursively
- ▶ **Example:** Consider the set  $S$  defined as:

$$3 \in S$$
$$\text{If } x \in S \text{ and } y \in S, \text{ then } x + y \in S$$

- ▶ What is the set  $S$  defined as above?
- ▶ Give a recursive definition of the set  $E$  of all even integers:
  - ▶ **Base case:**
  - ▶ **Recursive step:**

## Strings and Alphabets

- ▶ Recursive definitions play important role in study of **strings**
- ▶ Strings are defined over an **alphabet**  $\Sigma$ 
  - ▶ **Example:**  $\Sigma_1 = \{a, b\}$
- ▶ Set of all strings formed from  $\Sigma$  forms **language** called  $\Sigma^*$ 
  - ▶  $\Sigma_1^*$ :  $\{\epsilon, a, b, aa, ab, ba, bb, \dots\}$

## Recursive Definition of Strings

- ▶ The language  $\Sigma^*$  has natural recursive definition:
  - ▶ **Base case:**  $\epsilon \in \Sigma^*$  (empty string)
  - ▶ **Recursive step:** If  $w \in \Sigma^*$  and  $x \in \Sigma$ , then  $wx \in \Sigma^*$
- ▶ Since  $\epsilon$  is the empty string,  $\epsilon s = s$
- ▶ Consider the alphabet  $\Sigma = \{0, 1\}$
- ▶ How is the string "1" formed according to this definition?
- ▶ How is "10" formed?

## Recursive Definitions of String Operations

- ▶ Many operations on strings can be defined recursively.
- ▶ Consider function  $l(w)$  which yields length of string  $w$
- ▶ Example: Give recursive definition of  $l(w)$ 
  - ▶ Base case:
  - ▶ Recursive step:

## Another Example

- ▶ The **reverse** of a string  $s$  is  $s$  written backwards.
- ▶ Example: Reverse of "abc" is "bca"
- ▶ Give a recursive definition of the **reverse**( $s$ ) operation
  - ▶ Base case:
  - ▶ Recursive step:

## Palindromes

- ▶ A **palindrome** is a string that reads the same forwards and backwards
- ▶ Examples: "mom", "dad", "abba", "Madam I'm Adam", ...
- ▶ Give a recursive definition of the set  $P$  of all palindromes over the alphabet  $\Sigma = \{a, b\}$
- ▶ Base cases:
- ▶ Recursive step:

## Structural Induction

- ▶ How do we prove properties about recursively defined structures?
- ▶ **Structural induction** is a technique that allows us to apply induction on recursive definitions even if there is no integer
- ▶ Structural induction is also no more powerful than regular induction, but can make proofs much easier

## Structural Induction Overview

- ▶ Suppose we have:
  - ▶ a recursively defined structure  $S$
  - ▶ a property  $P$  we'd like to prove about  $S$
- ▶ **Structural induction** works as follows:
  1. **Base case:** Prove  $P$  about base case in recursive definition
  2. **Inductive step:** Assuming  $P$  holds for sub-structures used in the recursive step of the definition, show that  $P$  holds for the recursively constructed structure.

## Example 1

- ▶ Consider the following recursively defined set  $S$ :
  1.  $a \in S$
  2. If  $x \in S$ , then  $(x) \in S$
- ▶ Prove by **structural induction** that every element in  $S$  contains an equal number of right and left parentheses.
- ▶ **Base case:**  $a$  has 0 left and 0 right parentheses
- ▶ **Inductive step:** By the inductive hypothesis,  $x$  has equal number, say  $n$ , of right and left parentheses.
- ▶ Thus,  $(x)$  has  $n + 1$  left and  $n + 1$  right parentheses.

## Example 2

- ▶ Consider the set  $S$  defined recursively as follows:
  - ▶ **Base case:**  $3 \in S$
  - ▶ **Recursive step:** If  $x \in S$  and  $y \in S$ , then  $x + y \in S$
- ▶ Prove  $S$  is set of all positive integers that are multiples of 3
- ▶ Let  $A$  be the set of all positive integers divisible by 3
- ▶ We want to show that  $A = S$
- ▶ To do this, we need to prove  $S \subseteq A$  and  $A \subseteq S$

## Proof, Part I

Consider the set  $S$  defined recursively as follows:  $3 \in S$  and if  $x \in S$  and  $y \in S$ , then  $x + y \in S$

- ▶ Let's first prove  $S \subseteq A$ , i.e., any element in  $S$  is divisible by 3
- ▶ For this, we'll use structural induction
- ▶ **Base case:**
- ▶ **Inductive step:**

## Proof, Part II

- ▶ We showed that all integers in  $S$  are multiples of 3, but still need to show  $S$  includes **all** positive multiples of 3
- ▶ Therefore, need to prove that  $3n \in S$  for all  $n \geq 1$
- ▶ We'll prove this by strong induction on  $n$ :
  - ▶ **Base case** ( $n=1$ ):
  - ▶ **Inductive hypothesis:**
  - ▶ **Need to show:**
  - ▶
  - ▶