Modern C++ Programming

4. Basic Concepts III

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Memory

and Stack

Management: Heap

Process Address Space

addresses 0x00FFFFFF

higher memory

text

BSS and Data

Segment .bss/.data

Stack

Heap

Code

stack memory

int data[10]

new int[10]

malloc(40)

Static/Global

data

dynamic memory

int data[10]

(global scope)

lower memory addresses 0x00FF0000

new, delete

new, delete

new /new[] and delete /delete[] are C++ keywords that
perform dynamic memory allocation/deallocation, and object
construction/ destruction at runtime

malloc and free are C functions and they allocate and free
memory blocks (expressed in bytes)

Example:

```
int* array = new int[10]; // C: (int*) malloc(10 * sizeof(int))
delete[] array; // C: free(array)
```

new, delete Advantages

- Language keywords, not functions → safer
- Return type: new returns exact data type, while malloc() returns void*
- Failure: new throws an exception, while malloc() returns a NULL pointer → it cannot be ignored
- Allocated bytes: The size of the allocated memory is calculated by the compiler for new, while the user must take care of manually calculate the size for malloc()
- Initialization: new can be used to initialize an object or a set of objects

Dynamic Allocation

Allocate a single element

```
int* value = (int*) malloc(sizeof(int)); // C
int* value = new int; // C++
```

Allocate N elements

```
int* array = (int*) malloc(N * sizeof(int)); // C
int* array = new int[N]; // C++
```

Allocate and zero-initialize N elements

```
int* array = (int*) calloc(N * sizeof(int)); // C
int* array = new int[N](); // C++
```

Allocate N structures

```
MyStruct* array = (int*) malloc(N * sizeof(MyStruct)); // C
MyStruct* array = new MyStruct[N]; // C++
```

Dynamic Deallocation

Deallocate a single element

```
int* value = (int*) malloc(sizeof(int)); // C
free(value);
int* value = new int; // C++
delete value;
```

Deallocate N elements

```
int* value = (int*) malloc(N * sizeof(int)); // C
free(value);
int* value = new int[N]; // C++
delete[] value;
```

Fundamental rules:

- Each object allocated with new must be deallocated with delete
- Each object allocated with new[] must be deallocated with delete[]

Mixing new , new[] , malloc with something different from
their counterparts leads to undefined behavior

delete and delete[] applied to NULL/ nullptr pointers do
not produce errors (same as free)

Memory Leak

Memory Leak

A **memory leak** is a dynamically allocated entity in heap memory that is <u>no longer used</u> by the program, but still maintained overall its execution

Problems:

- Illegal memory accesses \rightarrow segmentation fault
- lacktriangle Undefined values ightarrow segmentation fault
- Additional memory consumption

```
int main() {
   int* array = new int[10];
   array = nullptr; // memory leak!!
} // the memory can no longer be deallocated!!
```

Note: the memory leaks are especially difficult to detect in complex code and when objects are widely used

Wild and Dangling Pointers

Wild pointer:

Dangling pointer:

```
int main() {
   int* array = new int[10];
   delete[] array; // ok -> "array" now is a dangling pointer
   delete[] array; // double free or corruption!!
   // program aborted, the value of "array" is not null
}
```

Solution:

```
int main() {
    int* array = new int[10];
    delete[] array; // ok -> "array" now is a dangling pointer
    array = nullptr; // no more dagling pointer
    delete[] array; // ok, no side effect
}
```

Stack Allocation

A local variable is either in stack memory or CPU registers

Every object which resides in the stack is not valid outside the current scope!!

```
int* wrongFunction() {
    int array[3] = {1, 2, 3};
    return array;
}
int main() {
    int* ptr = wrongFunction();
    cout << ptr[0]; // Illegal memory access!!
}</pre>
```

The organization of stack memory enables much higher performance. On the other hand, this memory space is limited!!

It is $\approx 8MB$ on linux by default

2D Memory Allocation

Easy on stack:

```
int A[3][4];
```

Dynamic Memory 2D allocation/free:

```
int** A = new int*[3];
for (int i = 0; i < 3; i++)
    A[i] = new int[4];

for (int i = 0; i < 3; i++)
    delete[] A[i];
delete[] A;</pre>
```

Dynamic memory 2D allocation/free C++11:

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Data and BSS Segment

Data/BSS (Block Started by Symbol) segments are larger than stack memory (max \approx 1GB in general) but slower

Initialization

Variable Initialization

```
C++03:
```

C++11 (Uniform Initialization):

```
int b1{2};  // direct list initialization

// Zero-initialization / Default-Initialization
int b3 = {};  // copy list initialization, default value
int b4{};  // direct list initialization, default value
```

Stack Array Initialization

One dimension:

```
int a[3] = {1, 2, 3}; // explicit size
int b[] = {1, 2, 3}; // implicit size
char c[] = "abcd"; // implicit size
int d[3] = {1, 2}; // d[2] = 0 -> zero/default value

int e[4] = {0}; // all values of D are initialized to 0
int f[3] = {}; // all values of E are initialized to 0 (C++11)
int g[3] {}; // all values of F are initialized to 0 (C++11)
```

Two dimensions:

Dynamic Initialization

```
C++03:
int* a1 = new int;  // undefined
int* a2 = new int(); // zero-initialization
int* a3 = new int(4); // allocate a single value equal to 4
int* a4 = new int[3]; // allocate 4 elements with undefined values
int* a5 = new int[4](); // allocate 4 elements zero-initialized
// int* a6 = new int[4](3); // not valid
C++11:
int* b1 = new int[4]{};  // allocate 4 elements zero-initialized
int* b2 = new int[4]{1, 2}; // set first, second, zero-initialized
```

Pointers and References

Pointers and Pointer Dereferencing

Pointer

A pointer is a value referring to a location in memory

Pointer Dereferencing

Pointer **dereferencing** means obtaining the value stored in at the location refereed to the pointer

```
int* ptr1 = new int;
*ptr1 = 4;    // deferencing (assignment)
int a = *ptr1; // deferencing (get value)
```

Common error:

```
int *ptr1, ptr2; // one pointer and one integer!!
int *ptr1, *ptr2; // ok, two pointers
```

void Pointer (Generic Pointer)

Instead of declaring different types of pointer variable it is possible to declare single pointer variable which can act as any pointer types

- A void* can be assigned to another void*
- void* can be compared for equality and inequality
- A void* can be explicitly converted to another type
- Other operations would be unsafe because the compiler cannot know what kind of object is really pointed to. Consequently, other operations result in compile-time errors

```
cout << (sizeof(void*) == sizeof(int*));  // print true

int array[] = { 2, 3, 4 };

void* ptr = array;
cout << *array;  // print 2

// cout << *ptr;  // compile error
cout << *((int*) ptr);  // print 2

// void* ptr2 = ptr + 2;  // compile error</pre>
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```

Address-of operator &

The address-of operator (&) returns the address of a variable

To not confuse with Reference syntax: T& var = ...

```
int array[4];
// &array is a pointer to an array of size 4
int size1 = (&array)[1] - array;
int size2 = *(&array + 1) - array;
cout << size1; // print 4
cout << size2; // print 4</pre>
```

$1+1 \neq 2$: Pointer Arithmetic

Pointer syntax:

```
ptr[i] is equal to *(ptr + i)
```

Pointer arithmetic rule:

```
address(ptr + i) = address(ptr) + (sizeof(T) * i)
```

where \mathtt{T} is the type of elements pointed by \mathtt{ptr}

Example:

```
char arr[3] = "abc"

value address

| 'a' | 0x0 | \leftarrow arr[0]
```

 $\begin{array}{c|cccc} \text{'b'} & 0x1 & \leftarrow \text{arr}[1] \\ \text{'c'} & 0x2 & \leftarrow \text{arr}[2] \\ \end{array}$

int arr[3] =
$$\{4,5,6\}$$

	value	address	
	4	0×0	\leftarrow arr[0]
		0×1	
		0x2	
ı		0x3	

	4	\leftarrow arr[1]
5	0×5	
5	0×6	
	0×7	

Reference

A variable **reference** (T&) is an **alias**, namely another name for an already existing variable. Both variable and variable reference can be applied to refer the value of the variable

- A pointer has its own memory address and size on the stack, reference shares the same memory address (with the original variable)
- References cab be internally implemented as pointers, but the compiler treats them in a very different way

References are safer than pointers:

- References <u>cannot have NULL</u> value. You must always be able to assume that a reference is connected to a legitimate storage
- References <u>cannot be changed</u>. Once a reference is initialized to an object, it cannot be changed to refer to another object (Pointers can be pointed to another object at any time)
- References must be <u>initialized</u> when they are created (Pointers can be initialized at any time)

Reference (Examples)

Reference syntax: T& var = ...

```
//int& d; // reference. compile error!! no initilization
int c = 2;
int& e = c; // reference. ok valid initialization
e++; // increment
cout << c; // print 3</pre>
```

Reference (Function Arguments)

Reference vs. pointer arguments:

```
void f(int* value) {} // value may be a nullptr
void g(int& value) {} // value is never a nullptr

int a = 3;
f(&a); // ok
g(a); // ok
//g(3); // compile error!! "3" is not a reference of something
```

References can be use to indicate fixed size arrays:

Reference (Arrays)

Reference:

[1] www3.ntu.edu.sg/home/ehchua/programming/cpp/cp4_PointerReference.html

Reference and struct

- The dot (.) operator is applied to local objects and references
- The arrow operator (->) is used with a pointer to an object

```
#include <iostream>
using namespace std;
struct A {
   int x = 3;
};
int main() {
   A obj;
   A* p = \&obj; // pointer
   p->x; // arrow syntax
   A& ref = obj; // reference
   cout << obj.x; // dot syntax</pre>
   cout << ref.x; // dot syntax</pre>
```

sizeof Operator

sizeof operator

sizeof

The **sizeof** is a compile-time operator that determines the size, in bytes, of a variable or data type

- sizeof returns a value of type size_t
- sizeof(incomplete type) produces compile error, e.g. void
- sizeof(bitfield member) produces compile error
- sizeof(anything) never returns 0, except for array of size 0
- sizeof(char) always returns 1
- When applied to structures, it also takes into account padding
- When applied to a reference, the result is the size of the referenced type

```
int f(int[] array) {    // dangerous!!
    cout << sizeof(array);
}
int array1[10];
int* array2 = new int[10];
cout << sizeof(array1); // print sizeof(int) * 10 = 40
cout << sizeof(array2); // print sizeof(int*) = 8 (64-bit)
f(array1); // print 8</pre>
```

```
struct B {
   int x;
   char y;
};
struct C : B { // C extends B
   short z;
};
sizeof(B); // 8 : 4 + 1 (+ 3) (padding)
sizeof(C); // 12 : sizeof(B) + 2 (+ 2) (padding)
struct A {};
sizeof(A); // 1 : sizeof never return 0 (except for arrays)
```

```
char a;
char \& b = a;
sizeof(&a); // 8 in a 64-bit OS (pointer)
sizeof(b); // 1, i.e. sizeof(char)
              // NOTE: a reference is not a pointer
struct A {};
A array1[10];
sizeof(array1); // 1 : array of empty structures
int array2[0];
sizeof(array2); // 0 : special case
```

const and constexpr

const Keyword

const keyword

The const keyword indicates objects never changing value after their initialization (they must be initialized when declared)

const variables are evaluated at compile-time value if the right
expression is also evaluated at compile-time

Constness rules:

- int* \rightarrow const int*
- const int* *→* int*

```
int f1(const int* array) { // the values of the array cannot be
                          // modified
    . . .
int f2(int* array) {}
int* ptr = new int[3];
const int* c_ptr = new int[3];
f1(ptr); // ok
f2(ptr); // ok
f1(c_ptr); // ok
// f2(c_ptr); // compile error!!
void g(const int) { // pass-by-value combined with 'const'
                    // note: it is not useful because the value
```

// is copied

- int* pointer to int
 - The value of the pointer can be modified
 - The elements refereed by the pointer can be modified
- const int* pointer to const int. Read as (const int)*
 - The value of the pointer can be modified
 - The elements refereed by the pointer cannot be modified
- int *const const pointer to int
 - The value of the pointer cannot be modified
 - The elements refereed by the pointer can be modified
- const int *const const pointer to const int
 - The value of the pointer cannot be modified
 - The elements refereed by the pointer cannot be modified

Note: const int* is equal to int const*

Tip: pointer types should be read from right to left

Common error: adding const to a pointer is <u>not</u> the same as adding const to a type alias of a pointer

```
using ptr_t = int*;
using const_ptr_t = const int*;
void f1(const int* ptr) {
// ptr[0] = 0; // not allowed: pointer to const objects
   ptr = nullptr; // allowed
void f3(const_ptr_t ptr) { // same as before
// ptr[0] = 0; // not allowed: pointer to const objects
   ptr = nullptr; // allowed
void f2(const ptr_t ptr) { // warning!!
   ptr[0] = 0;  // allowed
// ptr = nullptr; // not allowed: const pointer to
             // modifiable objects
```

constexpr (function)

C++11/C++14/C++17 guarantees compile-time evaluation of an function as long as **all** its arguments are constant

- C++11: constexpr must contain exactly one return statement and it must not contain loops or switch
- C++14: constexpr has no restrictions

constexpr (variable)

C++11/C++14/C++17 constexpr variables are evaluated at compile-time

- const guarantees the value of a variable to be fixed overall the execution of the program
- constexpr tells the compiler that the expression results is at compile-time.
 constexpr value implies const

```
const int v1 = 3;  // compile-time evaluation
const int v2 = v1 * 2; // compile-time evaluation

int     a = 3;  // "a" is dynamic
const int v3 = a;  // run-time evaluation!!

constexpr c1 = v1;  // ok
// constexpr c2 = v3; // compile error!!
```

```
constexpr int square(int value) {
    return value * value;
}

square(4); // compile-time evaluation

int a = 4; // "a" is dynamic
square(a); // run-time evaluation
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```

if constexpr

C++17 introduces **if constexpr** feature which allows conditionally compiling code based on a compile-time value

It is an if statement where the branch is chosen at compile-time (similarly to the #if preprocessor)

```
void f() {
   if constexpr (true)
      std::cout << "compile!";
   else
      std::cout << "error!"; // never compiled
}</pre>
```

constexpr example

```
constexpr int fib(int n) {
    return (n == 0 || n == 1) ? 1 : fib(n - 1) + fib(n - 2);
}
int main() {
    if constexpr (sizeof(void*) == 8)
        return fib(5);
    else
        return fib(3);
}
```

Generated assembly code (x64 OS):

```
main:
  mov eax, 8
  ret
```

Explicit Type

Conversion

Old style cast (type) value

New style cast:

 static_cast does compile-time (not run-time) checking of the types involved In many situations, this can make it the safest type of cast, as it provides the least room for accidental/unsafe conversions between various types

reinterpret_cast

```
reinterpret_cast<T*>(v) equal to (T*) v
reinterpret_cast<T&>(v) equal to *((T*) &v)
```

const_cast may be used to cast away (remove) constness or volatility

Static cast vs. old style cast:

Const cast:

```
const int a = 5;
const_cast<int>(a) = 3; // ok
```

Reinterpret cast: (bit-level conversion)

Type Punning

Pointer Aliasing

One pointer **aliases** another when they both point to the <u>same</u> memory location

Type Punning

Type punning refers to circumvent the type system of a programming language to achieve an effect that would be difficult or impossible to achieve within the bounds of the formal language

```
bool is_negativeA(float x) {
    return x < 0.0;
}
bool is_negativeB(float x) {
    unsigned int* ui = (unsigned int *) &x; // gcc warning:
    return (*ui) & 0x80000000; // -Wstrict-aliasing
} // this could lead to undefined behavior</pre>
```

Brace Initialization

C++11 The **brace initialization** can be also used to *safely* convert arithmetic types, preventing implicit *narrowing*, i.e potential value loss

The syntax is also more concise than modern casts

```
int x1 = -1;
int64 t x2 = int64 t\{-1\}; // ok
//int64_t x3 = uint64_t{-1}; // compile error!! unsafe
int64_t y1 = int64_t\{x1\}; // ok (only GCC, not clang)
//uint64 t y2 = uint64 t\{x1\}; // compile error!! unsafe
uint64 t z1 = static cast\langle uint64 t \rangle (-1); // ok
uint64_t z2 = static_cast<uint64_t>(x1); // ok
float f1 = float{10e30}; // ok
//float f2 = float{10e40}; // compile error!! unsafe
```