Modern C++ Programming

4. Basic Concepts III

- Memory Management

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Heap and Stack

Process Address Space



lower memory addresses 0x00FF0000

Heap

BSS and Data
Segment
.bss/.data

Code
.text

Stack

stack memory

dynamic memory

Static/Global data

int data[10]

new int[10] malloc(40)

int data[10]
 (global scope)

Stack and Heap Memory Overview

Memory Organization	Contiguous	(block) Fragmented
Max size	Small (~8MB)	Whole system memory
If exceed	Program crash	Exception or nullptr
Allocation	Compile-time	Run-time

Heap

Shared among threads

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Stack

Locality High I ow Each thread in a program Thread View

has its own stack

Stack Memory

A local variable is either in stack memory or CPU registers

```
int x = 3; // not on stack
struct A {
   int k; // depends on where the instance of A is
};
int main() {
   int y = 3; // on stack
   char z[] = "abc"; // on stack
   A a; // on stack (also k)
   int* ptr = new int; // variable "ptr" is on stack
}
```

The organization of stack memory enables much higher performance. On the other hand, this memory space is limited!!

It is $\approx 8MB$ on linux by default

Stack Memory Data

Types of data stored in the stack:

Local variables Variable in a local scope

Function arguments Data passed from caller to a function

Return addresses Data passed from a function to a caller

Compiler temporaries Compiler specific instructions

Interrupt contexts

Stack Memory

Every object which resides in the stack is not valid outside his scope!!

```
int* f() {
    int array[3] = {1, 2, 3};
    return array;
}
int* ptr = f();
cout << ptr[0]; // Illegal memory access!! $\mathref{Z}$</pre>
```

new, delete

new, delete

new/new[] and delete/delete[] are C++ keywords that
perform dynamic memory allocation/deallocation, and object
construction/destruction at runtime

malloc and free are C functions and they allocate and free memory blocks (expressed in bytes)

Example:

```
int* array = new int[10]; // C: (int*) malloc(10 * sizeof(int))
delete[] array; // C: free(array)
```

new, delete Advantages

- Language keywords, not functions → safer
- Return type: new returns exact data type, while malloc() returns void*
- Failure: new throws an exception, while malloc() returns a NULL pointer → it cannot be ignored
- Allocated bytes: The size of the allocated memory is calculated by the compiler for new, while the user must take care of manually calculate the size for malloc()
- Initialization: new can be used to initialize an object or a set of objects

Dynamic Allocation

Allocate a single element

```
int* value = (int*) malloc(sizeof(int)); // C
int* value = new int; // C++
```

Allocate N elements

```
int* array = (int*) malloc(N * sizeof(int)); // C
int* array = new int[N]; // C++
```

Allocate and zero-initialize N elements

```
int* array = (int*) calloc(N * sizeof(int)); // C
int* array = new int[N](); // C++
```

Allocate N structures

```
MyStruct* array = (int*) malloc(N * sizeof(MyStruct)); // C
MyStruct* array = new MyStruct[N]; // C++
```

Dynamic Deallocation

Deallocate a single element

```
int* value = (int*) malloc(sizeof(int)); // C
free(value);
int* value = new int; // C++
delete value;
```

Deallocate N elements

```
int* value = (int*) malloc(N * sizeof(int)); // C
free(value);
int* value = new int[N]; // C++
delete[] value;
```

Dynamic Memory Notes

Fundamental rules:

- Each object allocated with new must be deallocated with delete
- Each object allocated with new[] must be deallocated with delete[]

Mixing new , new[] , malloc with something different from their counterparts leads to undefined behavior

delete and delete[] applied to NULL/ nullptr pointers do
not produce errors (same as free)

Memory Leak

Memory Leak

A **memory leak** is a dynamically allocated entity in heap memory that is <u>no longer used</u> by the program, but still maintained overall its execution

Problems:

- Illegal memory accesses → segmentation fault
- Undefined values \rightarrow segmentation fault
- Additional memory consumption

```
int main() {
   int* array = new int[10];
   array = nullptr; // memory leak!!
} // the memory can no longer be deallocated!!
```

Note: the memory leaks are especially difficult to detect in complex code and when objects are widely used

2D Memory Allocation

Easy on stack:

```
int A[3][4];
```

Dynamic Memory 2D allocation/free:

```
int** A = new int*[3];
for (int i = 0; i < 3; i++)
    A[i] = new int[4];

for (int i = 0; i < 3; i++)
    delete[] A[i];
delete[] A;</pre>
```

Dynamic memory 2D allocation/free C++11:

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Data and BSS Segment

Data/BSS (Block Started by Symbol) segments are larger than stack memory (max \approx 1GB in general) but slower

Initialization

Variable Initialization

```
C++03:
```

```
int a1;
               // default initialization (undefined value)
              // direct (or value) initialization
int a2(2):
int a3 = 2;  // copy initialization
int a4 = 2u;  // copy initialization (implicit conversion)
int();  // direct (or value) initialization (zero-init)
int(2)
            // direct (or value) initialization
int a5 = int(2); // copy initialization
int a6 = int(); // copy initialization (zero-initialization)
int a7 = {2}; // copy list initialization
// int a8();  // a8 is a function
```

Uniform Initialization

C++11 provides the **Uniform Initialization** syntax, also called brace-initialization or braced-init-list, to initialize different entities (variables, objects, structures, etc.) in a <u>consistent</u> way:

```
// direct list (or value) initialization
int b1{2};
                 // direct list (or value) initialization
int b2{}:
                 // (zero-initialization)
int b3 = {};  // copy list initialization (zero-initialization)
int{}:
               // direct list (or value) initialization (zero-init)
int{4};
               // direct list (or value) initialization
int b4 = int{}; // copy initialization
int b5 = int{4}; // copy initialization
```

Brace Initialization Advantages

The C++11 uniform initialization can be also used to *safely* convert arithmetic types, preventing implicit *narrowing*, i.e potential value loss. The syntax is also more concise than modern casts

```
int b4 = -1; // ok
int b5{-1}; // ok
unsigned b6 = -1; // ok
//unsigned b7{-1}; // compile error

float f1{10e30}; // ok
float f2 = 10e40; // ok, "inf" value
//float f3{10e40}; // compile error
```

```
struct S {
   unsigned x;
   unsigned y;
};
// C++03
S s1; // default initialization (undefined values)
S s2 = {}; // copy list initialization (zero-init)
S s3 = \{1, 2\}; // copy list initialization
// C++11
S s4{}; // direct list (or value) initialization (zero-init)
S s5{1, 2}; // direct list (or value) initialization
// S s6{1, -2}; // compile error
```

Non-Static Data Member Initialization (NSDMI):

```
struct S {
    unsigned x = 3; // equal initialization
    unsigned y = 2; // equal initialization
};
                   // also functions are allowed
struct S1 {
    unsigned x {3}; // brace initialization
    unsigned y {2}; // brace initialization
};
S s1; // call default constructor (x=3, y=2)
S s2{}; // call default constructor (x=3, y=2)
S s3{1, 4}; // set x=1, y=4
S f() {
   return {3, 2};
```

Stack Array Initialization

One dimension:

```
int a[3] = {1, 2, 3}; // explicit size
int b[] = {1, 2, 3}; // implicit size
char c[] = "abcd"; // implicit size
int d[3] = {1, 2}; // d[2] = 0 -> zero/default value

int e[4] = {0}; // all values of D are initialized to 0
int f[3] = {}; // all values of E are initialized to 0 (C++11)
int g[3] {}; // all values of F are initialized to 0 (C++11)
```

Two dimensions:

Dynamic Initialization

int* a1 = new int; // undefined

C++03:

```
int* a2 = new int();  // zero-initialization
int* a3 = new int(4);  // allocate a single value equal to 4
int* a4 = new int[4];  // allocate 4 elements with undefined values
int* a5 = new int[4]();  // allocate 4 elements zero-initialized
// int* a6 = new int[4](3);  // not valid

C++11:
int* b1 = new int[4]{};  // allocate 4 elements zero-initialized
int* b2 = new int[4]{1, 2};  // set first, second, zero-initialized
```

Pointers and References

Pointer

A pointer T* is a value referring to a location in memory

Pointer Dereferencing

Pointer **dereferencing** (*ptr) means obtaining the value stored in at the location refereed to the pointer

Subscript Operator []

The subscript operator (ptr[]) allows accessing to the pointer element at a given position

Deferencing:

```
int* ptr1 = new int;
*ptr1 = 4;    // deferencing (assignment)
int a = *ptr1; // deferencing (get value)
```

Array subscript:

```
int* ptr2 = new int[10];
ptr2[2] = 3;
int var = ptr2[4];
```

Common error:

```
int *ptr1, ptr2; // one pointer and one integer!!
int *ptr1, *ptr2; // ok, two pointers
```

Subscript operator meaning:

```
ptr[i] is equal to *(ptr + i)
```

Note: subscript operator accepts also negative values

Pointer arithmetic rule:

```
address(ptr + i) = address(ptr) + (sizeof(T) * i)
```

where T is the type of elements pointed by ptr

$$\begin{array}{c|cccc} value & address \\ \hline 'a' & 0x0 & \leftarrow arr[0] \\ 'b' & 0x1 & \leftarrow arr[1] \\ 'c' & 0x2 & \leftarrow arr[2] \\ \hline '\backslash 0' & 0x3 & \leftarrow arr[3] \\ \end{array}$$

int arr[3] =
$$\{4,5,6\}$$

value	address	
	0×0	\leftarrow arr[0]
	0×1	
4	0×2	
	0×3	
	0×4	←arr[1]
_	0×5	
5	0×6	
	0×7	
	0×8	←arr[2]
	0×9	
6	0×10	
	0×11	

Address-of operator &

The address-of operator (&) returns the address of a variable

To not confuse with Reference syntax: T& var = ...

Wild and Dangling Pointers

Wild pointer:

Dangling pointer:

```
int main() {
   int* array = new int[10];
   delete[] array; // ok -> "array" now is a dangling pointer
   delete[] array; // double free or corruption!!
   // program aborted, the value of "array" is not null
}
```

Solution:

```
int main() {
   int* array = new int[10];
   delete[] array; // ok -> "array" now is a dangling pointer
   array = nullptr; // no more dagling pointer
   delete[] array; // ok, no side effect
}
```

void Pointer (Generic Pointer)

Instead of declaring different types of pointer variable it is possible to declare single pointer variable which can act as any pointer types

- void* can be compared
- A void* can be implicitly converted to another pointer
- Other operations are unsafe because the compiler does not know what kind of object is really pointed to

```
cout << (sizeof(void*) == sizeof(int*)); // print true

int array[] = { 2, 3, 4 };
void* ptr = array; // implicit conversion
cout << *array; // print 2
// *ptr; // compile error
// ptr + 2; // compile error</pre>
```

Reference

A variable **reference T&** is an **alias**, namely another name for an already existing variable. Both variable and variable reference can be applied to refer the value of the variable

- A pointer has its own memory address and size on the stack, reference shares the same memory address (with the original variable)
- The compiler <u>can</u> internally implement references as *pointers*, but treats them in a very different way

References are safer than pointers:

- References <u>cannot have NULL</u> value. You must always be able to assume that a reference is connected to a legitimate storage
- References <u>cannot be changed</u>. Once a reference is initialized to an object, it cannot be changed to refer to another object (Pointers can be pointed to another object at any time)
- References must be <u>initialized</u> when they are created (Pointers can be initialized at any time)

Reference (Examples)

Reference syntax: T& var = ...

```
//int& a; // compile error no initilization
//int \& b = 3; // compile error "3" is not a variable
int c = 2;
int& d = c; // reference. ok valid initialization
int& e = d; // ok. the reference of a reference is a reference
d++; // increment
e++; // increment
cout << c; // print 4
int a = 3;
int* b = &a; // pointer
int* c = &a; // pointer
b++; // change the value of the pointer 'b'
*c++; // change the value of 'a' (a = 4)
int& d = a; // reference
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d++; // change the value of 'a' (a = 5)
```

Reference vs. pointer arguments:

```
void f(int* value) {} // value may be a nullptr
void g(int& value) {} // value is never a nullptr
int a = 3;
f(\&a); // ok
f(0); // dangerous but it works!! (but not with other numbers)
//f(a); // compile error "a" is not a pointer
g(a); // ok
//g(3); // compile error "3" is not a reference of something
//q(&a); // compile error "&a" is not a reference
```

References can be use to indicate fixed size arrays:

```
void f(int (&array)[3]) { // accepts only arrays of size 3
   cout << sizeof(array);</pre>
}
void g(int array[]) {
    cout << sizeof(array); // any surprise?</pre>
}
int A[3], B[4];
int* C = A;
//----
f(A); // ok
// f(B); // compile error B has size 4
// f(C); // compile error C is a pointer
g(A); // ok
g(B); // ok
g(C); // ok
```

Reference (Arrays) ★

```
int A[4];
int (&B) [4] = A; // ok, reference to array
int C[10][3]:
int (&D)[10][3] = C; // ok, reference to 2D array
auto c = new int[3][4]; // type is int (*)[4]
// read as "pointer to arrays of 4 int"
// int (&d)[3][4] = c; // compile error
// int (*e)[3] = c; // compile error
int (*f)[4] = c; // ok
int array[4];
// &array is a pointer to an array of size 4
int size1 = (&array)[1] - array;
int size2 = *(&array + 1) - array;
cout << size1; // print 4</pre>
cout << size2; // print 4</pre>
```

Reference and struct

- The dot (.) operator is applied to local objects and references
- The arrow operator (->) is used with a pointer to an object

```
struct A {
  int x = 3;
};
Aa;
A* ptr = &a; // pointer
ptr->x; // arrow syntax
A& ref = a; // reference
a.x; // dot syntax
ref.x; // dot syntax
```

const, constexpr,

consteval,

constinit

const Keyword

const keyword

The const keyword indicates objects never changing value after their initialization (they must be initialized when declared)

const variables are evaluated at compile-time value if the right
expression is also evaluated at compile-time

Constness rules:

- int* → const int*
- const int* → int*

- int* pointer to int
 - The value of the pointer can be modified
 - The elements refereed by the pointer can be modified
- const int* pointer to const int. Read as (const int)*
 - The value of the pointer can be modified
 - The elements refereed by the pointer cannot be modified
- int *const const pointer to int
 - The value of the pointer cannot be modified
 - The elements refereed by the pointer can be modified
- const int *const const pointer to const int
 - The value of the pointer cannot be modified
 - The elements refereed by the pointer cannot be modified

Note: const int* is equal to int const*

Tip: pointer types should be read from right to left

Common error: adding const to a pointer is <u>not</u> the same as adding const to a type alias of a pointer

```
using ptr_t = int*;
using const_ptr_t = const int*;
void f1(const int* ptr) {
// ptr[0] = 0; // not allowed: pointer to const objects
   ptr = nullptr; // allowed
void f3(const_ptr_t ptr) { // same as before
// ptr[0] = 0; // not allowed: pointer to const objects
   ptr = nullptr; // allowed
void f2(const ptr_t ptr) { // warning!!
   ptr[0] = 0;  // allowed
// ptr = nullptr; // not allowed: const pointer to
             // modifiable objects
```

constexpr (variable)

C++11 constexpr variables are evaluated at compile-time

- const guarantees the value of a variable to be fixed overall the execution of the program
- constexpr tells the compiler that the expression results is at compile-time.
 constexpr value implies const

constexpr (function)

C++11 constexpr guarantees compile-time evaluation of a function as long as <u>all</u> its arguments are constant

- C++11: must contain exactly one return statement and it must not contain loops or switch
- C++14: no restrictions

```
const int v1 = 3;  // compile-time evaluation
const int v2 = v1 * 2;  // compile-time evaluation

int    a = 3;  // "a" is dynamic
const int v3 = a;  // run-time evaluation!!

constexpr int c1 = v1;  // ok
// constexpr int c2 = v3; // compile error, "v3" is dynamic
```

```
constexpr int square(int value) {
    return value * value;
}

square(4); // compile-time evaluation

int a = 4; // "a" is dynamic
square(a); // run-time evaluation

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```

consteval Keyword

consteval

C++20 consteval, or immediate functions, guarantees compile-time evaluation of a function. A non-constant value produces a compilation error

```
consteval int square(int value) {
    return value * value;
}

square(4);  // compile-time evaluation

int v = 4;  // "v" is dynamic

// square(v); // compile error
```

constinit Keyword

constinit

C++20 constinit guarantees compile-time initialization of a variable. A non-constant value produces a compilation error

- The value of the variable can change during the execution
- const constinit does not imply constexpr, while the opposite is true
- constexpr requires compile-time evaluation during his entire lifetime

```
constexpr int square(int value) {
    return value * value;
}
constinit int v1 = square(4);  // compile-time evaluation
v1 = 3;  // ok, v1 can change

int a = 4;  // "v" is dynamic  45/59

// constinit int v2 = square(a); // compile error
```

if constexpr

C++17 introduces **if constexpr** feature which allows conditionally compiling code based on a compile-time value

It is an if statement where the branch is chosen at compile-time (similarly to the #if preprocessor)

```
void f() {
   if constexpr (true)
      cout << "compile!";
   else
      cout << "error!"; // never compiled
}</pre>
```

constexpr example

```
constexpr int fib(int n) {
    return (n == 0 || n == 1) ? 1 : fib(n - 1) + fib(n - 2);
}
int main() {
    if constexpr (sizeof(void*) == 8)
        return fib(5);
    else
        return fib(3);
}
```

Generated assembly code (x64 OS):

```
main:
  mov eax, 8
  ret
```

std::is_constant_evaluated

C++20 provides std::is_constant_evaluated() utility for evaluating if the current function is evaluated at compile time

```
#include <type_traits> // std::is_constant evaluated
constexpr int f(int n) {
    if (std::is_constant_evaluated())
       return 0;
   return 4;
int x = f(3): // x = 0
int v = 3;
int y = f(y); // y = 4
```

Explicit Type

Conversion

Old style cast (type) value

New style cast:

- static_cast does compile-time (not run-time) checking of the types involved In many situations, this can make it the safest type of cast, as it provides the least room for accidental/unsafe conversions between various types
- const_cast can add or cast away (remove) constness or volatility
- reinterpret_cast

```
reinterpret_cast<T*>(v) equal to (T*) v
reinterpret_cast<T&>(v) equal to *((T*) &v)
```

const_cast and reinterpret_cast do not compile to any CPU
instructions

Static cast vs. old style cast:

Const cast:

```
const int     a = 5;
const_cast<int>(a) = 3; // ok, but /\uline{undefined behavior}/
```

Reinterpret cast: (bit-level conversion)

Print the value of a pointer

Array reshaping

```
int a[3][4];
int (&b)[2][6] = reinterpret_cast<int (&)[2][6]>(a);
int (*c)[6] = reinterpret_cast<int (*)[6]>(a);
```

Pointer Aliasing

One pointer **aliases** another when they both point to the <u>same</u> memory location

Type Punning

Type punning refers to circumvent the type system of a programming language to achieve an effect that would be difficult or impossible to achieve within the bounds of the formal language

The compiler assumes that the *strict aliasing rule is never violated*. Accessing a value using a type which is different from the original one is not allowed and it is classified as *undefined behavior*

```
// slow without optimizations. The branch breaks the pipeline
float abs(float x) {
    return (x < 0.0f) ? -x : x;
}
// optimized by hand
float abs(float x) {
    unsigned uvalue = reinterpret_cast<unsigned&>(x);
    unsigned tmp = uvalue & 0x7FFFFFFF; // clear the last bit
    return reinterpret_cast<float&>(tmp);
// this is undefined behavior!!
```

GCC warning (not clang): -Wstrict-aliasing

- blog.qt.io/blog/2011/06/10/type-punning-and-strict-aliasing
- What is the Strict Aliasing Rule and Why do we care?
- Type Punning In C++17

memcpy and std::bit_cast

The right way to avoid undefined behavior is using memcpy

```
float v1 = 32.3f;
unsigned v2;
std::memcpy(&v2, &v1, sizeof(float);
// v1, v2 must be trivially copyable
```

C++20 provides std::bit_cast safe conversion for replacing reinterpret_cast

```
float v1 = 32.3f;
unsigned v2 = std::bit_cast<unsigned>(v1);
```

sizeof Operator

sizeof operator

sizeof

The **sizeof** is a compile-time operator that determines the size, in bytes, of a variable or data type

- sizeof returns a value of type size_t
- sizeof(incomplete type) produces compile error, e.g. void
- sizeof(bitfield member) produces compile error
- sizeof(anything) never returns 0, except for array of size 0
- sizeof(char) always returns 1
- When applied to structures, it also takes into account padding
- When applied to a reference, the result is the size of the referenced type

```
sizeof(int); // 4 bytes
sizeof(int*) // 8 bytes on a 64-bit OS
sizeof(void*) // 8 bytes on a 64-bit OS
sizeof(size_t) // 8 bytes on a 64-bit OS
```

sizeof (struct)

```
struct A { // "A" is aligned to the largest type: int (4 bytes)
    int x;
    char v; // offset 4
};
sizeof(A); // 8 bytes : 4 + 1 (+ 3 padding)
struct B {
    int x; // offset 0
    char y; // offset 4
    short z; // offset 6 -> 2 bytes aligned
};
sizeof(B); // 8 bytes : 4 + 1 (+ 1 padding) + 2
struct C {
    short z; // offset 0 -> 2 bytes aligned
    int x; // offset 4 -> 4 bytes aligned
    char y; // offset 8
};
                                                                 57/59
sizeof(C); // 12 bytes : 2 (+ 2 padding) + 4 + 1 + (+ 3 padding)
```

```
char a;
char \& b = a;
sizeof(&a); // 8 bytes in a 64-bit OS (pointer)
sizeof(b); // 1 byte, equal to sizeof(char)
               // NOTE: a reference is not a pointer
// SPECIAL CASES
struct A {};
sizeof(A); // 1 : sizeof never return 0 (except for arrays)
A array1[10];
sizeof(array1); // 1 : array of empty structures
int array2[0];
sizeof(array2); // 0 : special case
```

sizeof and Size of a Byte

Interesting: C++ does not explicitly define the size of a byte (see Exotic architectures the standards committees care about)