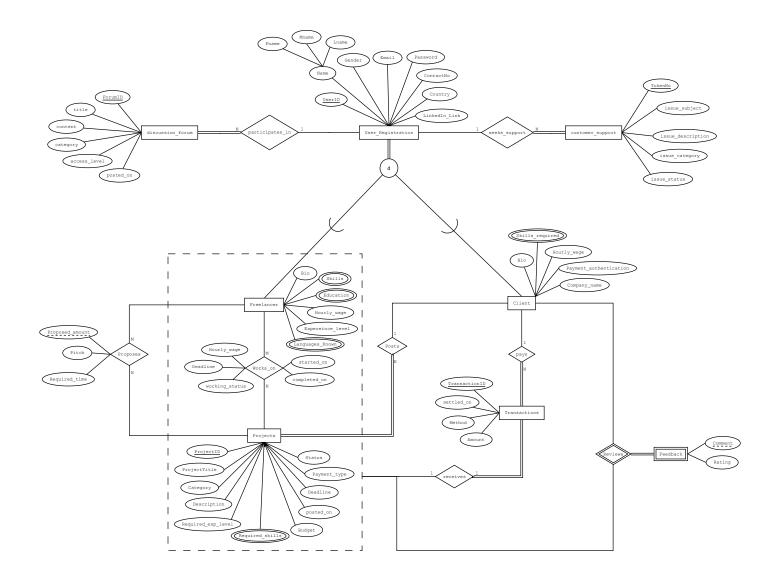
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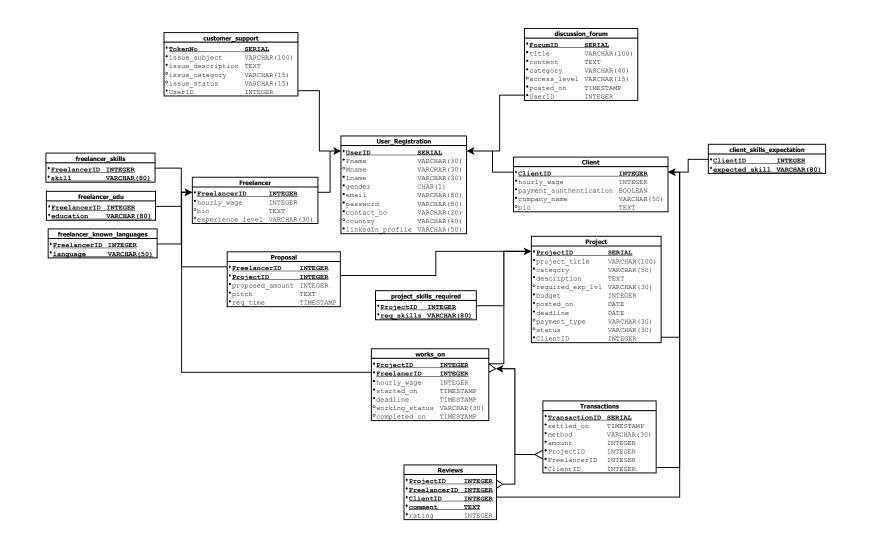


IT214 – Database Management System <u>Database of Freelancing Website</u>

T601

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MINIMAL FD SET OF THE RELATION

UserID → **Fname**

UserID → Mname

UserID → **Lname**

UserID → **Gender**

UserID → **Email**

UserID → **Password**

UserID → **Contact_No**

UserID → **Country**

UserID → **linkedin_profile**

Email → UserID

Contact_No → UserID

linkedin_profile → UserID

FreelancerID → Bio

FreelancerID → hourly_wage

FreelancerID → experience_level

ClientID → hourly_wage

ClientID → payment_authentication

ClientID → company name

ClientID → Bio

ProjectID → **project_title**

ProjectID → category

ProjectID → **description**

ProjectID → required_exp_level

```
ProjectID → budget
ProjectID → posted on
ProjectID → deadline
ProjectID → payment type
ProjectID → status
ProjectID → ClientID
{ProjectID, FreelancerID} → proposed amount
{ProjectID, FreelancerID} → pitch
{ProjectID, FreelancerID} → req_time
{ProjectID, FreelancerID} → hourly_wage
{ProjectID, FreelancerID} → started on
{ProjectID, FreelancerID} → deadline
{ProjectID, FreelancerID} → working status
{ProjectID, FreelancerID} → completed on
TranscationID → settled on
TranscationID → method
TranscationID → amount
TranscationID → ProjectID
TranscationID → FreelancerID
TranscationID → ClientID
{ClientID, FreelancerID, ProjectID, comment}→ rating
TokenNo → issue subject
TokenNo → issue description
TokenNo → issue category
TokenNo → issue status
```

TokenNo → UserID

 $\textbf{ForumNo} \rightarrow \textbf{title}$

ForumNo → content

ForumNo → category

 $\textbf{ForumNo} \rightarrow \textbf{access_level}$

 $\textbf{ForumNo} \rightarrow \textbf{posted_on}$

 $ForumNo \rightarrow UserID$

Proof that relations are in BCNF:

1. 'User_Registration' relation:

Attributes (UserID, Fname, Mname, Lname, gender, email, password, contact_no, country, linkedin_profile)

Functional Dependencies:

UserID → Fname

UserID → Mname

UserID → Lname

UserID → Gender

UserID → Email

UserID → Password

 $UserID \rightarrow Contact_No$

UserID → Country

UserID → linkedin profile

Email → UserID

Contact No → UserID

linkedin profile → UserID

Let X = UserID

Clouse of X:

X⁺ = {UserID, Fname, Mname, Lname, gender, email, password, contact no, country, linkedin profile}

Let Y = email

Clouse of Y:

Y⁺ = {UserID, Fname, Mname, Lname, gender, email, password, contact_no, country, linkedin_profile}

Let **Z = linkedin_profile**

Clouse of Z:

Z⁺ = {UserID, Fname, Mname, Lname, gender, email, password, contact no, country, linkedin profile}

Let W = contact no

Clouse of W:

W⁺ = {UserID, Fname, Mname, Lname, gender, email, password, contact_no, country, linkedin_profile}

Hence, **UserID**, **email**, **linkedin_profile**, **contact_no** are candidate keys of the relation.

The left side of FDs in the minimal FD set for the relation 'User_Registration' is either UserID, email, linkedin_profile or contact_no, which are the candidate keys of this relation. Hence, relation 'User_Registration' is in BCNF.

2. 'Freelancer' relation

Attributes (FreelancerID, hourly_wage, bio, experience_level)

Functional Dependencies:

FreelancerID → Bio

FreelancerID → hourly wage

FreelancerID \rightarrow experience_level

Let X = FreelancerID

Clouse of X:

X⁺ = {FreelancerID, hourly wage, bio, experience level}

Since closure of **FreelancerID** has all the attributes of **'Freelancer'** relation, **FreelancerID** is the candidate key of this relation.

The left side of all the minimal FDs of the relation '**Freelancer**' is **FreelancerID**, which is the key of this relation, so '**Freelancer**' is in BCNF.

3. 'freelancer skills' relation

Attributes (FreelancerID, skills)

Here, Candidate Key = {FreelancerID, skills}

In this relation, all attributes create a candidate key.

According to the theorem, such relations are always in BCNF.

So, 'freelancer_skills' is in BCNF.

4. 'freelancer_edu' relation

Attributes (FreelancerID, education)

Here, Candidate Key = {FreelancerID, education}

All attributes are a part of the candidate key. So, 'freelancer_edu' is in BCNF.

5. 'freelancer_known_languages' relation

Attributes (FreelancerID, language)

Here, **Candidate Key** = {FreelancerID, language}

According to the theorem, if all the attributes of a relation make a candidate key, then the relation is always in BCNF.

So, 'freelancer_known_languages' is in BCNF.

6. 'Client' relation

Attributes(ClientID, hourly_wage, payment_authentication, company_name, bio)

Functional Dependencies:

ClientID → hourly_wage

 $ClientID \rightarrow payment_authentication$

ClientID → company_name

ClientID → Bio

Let X = ClientID

Clouse of X:

X⁺ = {ClientID, hourly_wage, payment_authentication, company_name, bio}

ClientID is the candidate key of a minimal FD set.

The left side of all the minimal FDs in the relation 'Client' is ClientID, which is the key of this relation, so 'Client' is in BCNF.

7. 'client skills expectation' relation

Attributes(ClientID,expected_skills)

Here Candidate Key = {ClientID, expected_skills}

According to the theorem, if all the attributes of a relation make a candidate key, then the relation is always in BCNF.

So, 'client skills expectation' is in BCNF.

8. 'Project' relation

Attributes (ProjectID, project_title, category, description, required_exp_lvl, budget, posted_on, deadline, payment_type, status, ClientID)

Functional Dependencies:

ProjectID → project_title

 $ProjectID \rightarrow category$

ProjectID → description

 $ProjectID \rightarrow required_exp_level$

 $ProjectID \rightarrow budget$

 $ProjectID \to posted_on$

ProjectID → deadline

ProjectID → payment_type

```
ProjectID → status
ProjectID → ClientID
```

Let X = ProjectID

Clouse of X:

X⁺ = {ProjectID,project_title,category,description,required_exp_lvl,budget, posted_on,deadline,payment_type,status,ClientID}

Since closure of **ProjectID** has all the attributes of '**Project**' relation, **ProjectID** is the candidate key.

The left side of all the minimal FDs in the relation '**Project**' is **ProjectID**, which is the key of this relation, so '**Project**' is in BCNF.

9. 'project_skills_required' relation

```
Attributes(ProjectID,required_skills)

Here Candidate Key = {ProjectID,required_skills}\

All attributes are a part of the candidate key.

So, 'project_skills_required' is in BCNF.
```

10. 'Proposal' relation

Attributes (FreelancerID, ProjectID, proposed_amount, pitch, req_time)

Functional Dependencies:

```
\{ProjectID, FreelancerID\} \rightarrow proposed\_amount \\ \{ProjectID, FreelancerID\} \rightarrow pitch \\ \{ProjectID, FreelancerID\} \rightarrow req\_time
```

```
Let X = {ProjectID, FreelancerID}
Clouse of X:
```

 X^+ = {FreelancerID, ProjectID, proposed_amount, pitch, req_time} {**ProjectID, FreelancerID**} is the candidate key of a minimal FD set.

The left side of all the FDs in the minimal set of FDs for the relation '**Proposal**' is **{ProjectID, FreelancerID}**, which is the key of this relation, so '**Proposal**' is in BCNF.

11. 'works_on' relation

Attributes (ProjectID, FreelancerID, hourly_wage, started_on, deadline, working_status, completed_on)

Functional Dependencies:

```
{ProjectID, FreelancerID} → hourly_wage

{ProjectID, FreelancerID} → started_on

{ProjectID, FreelancerID} → deadline

{ProjectID, FreelancerID} → working_status

{ProjectID, FreelancerID} → completed on
```

Let X = {ProjectID, FreelancerID}

Clouse of X:

X⁺ = {ProjectID,FreelancerID,hourly_wage,started_on,deadline, Working status, completed on}

{ProjectID, FreelancerID} is the candidate key of a minimal FD set.

The left side of all the minimal FDs of the relation 'works_on' is {ProjectID, FreelancerID}, which is the key of this relation, so 'works_on' is in BCNF.

12. 'Transactions' relation

Attributes (TransactionID, settled_on, method, amount, ProjectID, FreelancerID, ClientID)

Functional Dependencies:

TranscationID → settled_on

```
TranscationID \rightarrow method
```

TranscationID → amount

TranscationID → ProjectID

TranscationID → FreelancerID

TranscationID → ClientID

{ProjectID, FreelancerID, ClientID} → TransactionID

Let X = {TranscationID}

Clouse of X:

X⁺ = {TransactionID , settled_on,method,amount,ProjectID,FreelancerID, ClientID}

Let Y = {ProjectID, FreelancerID, ClientID}

Clouse of Y:

Y⁺ = {TransactionID , settled_on,method,amount,ProjectID,FreelancerID, ClientID}

Thus, **TranscationID**, **{ProjectID**, **FreelancerID**, **ClientID}** is the candidate key of a minimal FD set.

The left side of all the minimal FDs of the relation '**Transactions**' is candidate key, so '**Transactions**' is in BCNF.

13. 'Reviews' relation

Attributes (ProjectID, FreelancerID, ClientID, comment, rating)

Functional Dependencies:

{ClientID, FreelancerID, ProjectID, comment} \rightarrow rating

Let X = {ClientID, FreelancerID, ProjectID, comment}

Clouse of X:

X⁺ = {ProjectID,FreelancerID,ClientID,comment,rating}

(ClientID, FreelancerID, ProjectID, comment) is the candidate key of a minimal FD set.

The left side of all the minimal FDs of the relation 'Reviews' is {ClientID, FreelancerID, ProjectID, comment} which is the key of this relation, so 'Reviews' is in BCNF.

14. 'customer_support' relation

Attributes (TokenNo, issue_subject, issue_description, issue_category, issue_status, UserID)

Functional Dependencies:

```
TokenNo → issue_subject
TokenNo → issue_description
TokenNo → issue_category
TokenNo → issue_status
TokenNo → UserID
```

Let X = TokenNo

Clouse of X:

```
X* = {TokenNo,issue_subject,issue_description,issue_category,
    issue_status,UserID}
```

Since closure of **TokenNo** has all the attributes of **'customer_support'** relation, **TokenNo** is the candidate key.

The left side of all the minimal FDs of the relation 'customer_support' is **TokenNo**, which is the key of this relation, so 'customer_support' is in BCNF.

15. 'discussion_forum' relation

Attributes(ForumID,title,content,category,access_level,posted_on,UserID)

Functional Dependencies:

```
ForumNo → title
ForumNo → content
ForumNo → category
ForumNo → access_level
ForumNo → posted_on
```

ForumNo \rightarrow UserID

Let X = ForumID

Clouse of X:

X⁺ = {ForumID,title,content,category,access_level,posted_on,UserID}

Since closure of **ForumID** has all the attributes of 'discussion_forum' relation, **ForumID** is the candidate key.

The left side of all the minimal FDs of the relation 'discussion_forum' is ForumID, which is the key of this relation, so 'discussion_forum' is in BCNF.