

Term Work - 04

Problem statement

Write a PROLOG for a MENU-DRIVEN PROGRAM for
member, concatenation, add, delete and permutation functions.

Theory

In Prolog, a menu-driven program can be implemented using a combination of facts, rules, and predicates. Here's general outline of how you can structure a menu-driven program in Prolog:

1. Define the menu option:
Create facts to represent each menu option, where the fact structure represents the option's name & any associated parameter as argument.
2. Implement menu handling rules:
 - Write rules that define the behaviour for menu option
 - Each rule should have a head that matches the selected menu option & a body that specifies the action to be performed.
3. Implement the main menu loop:
 - Write a predicate that displays the menu options & prompts the user for input.
 - Repeat the main menu loop until the user chooses to exit the program.

PROGRAM

% Hyper predicates

concatenate(T, L, R).

concatenate([H|T], L, [H|R]) :- concatenate(T, L, R).

add_element(X, L, [X|L]).

delete_element(_, [], []).

delete_element(X, [X|T], T).

delete_element(X, [H|T], [H|R]) :- X \= H, delete_element(X, T, R).

permute(T, [T]).

permute([H|T], R) :- permute(T, X), select(H, R, X).

list_member(X, [X|_]).

list_member(X, [_|T]) :- list_member(X, T).

% menu-driven function

menu:-

write('MENU'), nl.

write('1. Concatenate list'), nl,

write('2. Add element'), nl,

write('3. Delete element'), nl,

write('4. Permute list'), nl,

write('5. Check member'), nl,

write('6. Exit'), nl,

write('Enter the number of your choice:')

read(choice),

process(choice).

Process 1:-

```
write('Enter first list: '),  
read(L1),  
write('Enter second list: '),  
read(L2),  
concatenate(L1, L2, Result),  
write('Concatenated list: '),  
write(Result), nl.
```

menu.

Process 2:-

```
write('Enter an element: '),  
read(X),  
write('Enter a list: '),  
read(L),  
add_element(X, L, Result),  
write('List after adding element: '),  
write(Result), nl.
```

menu.

Process 3:-

```
write('Enter an element: '),  
read(X),  
write('Enter a list: '),  
read(L),  
delete_element(X, L, Result),  
write('List after deleting element: '),  
write(Result), nl.
```

menu.

Process 4:-

```
write('Enter a list: '),  
read(L),
```

```
final (X, Permuted (L, X), Result),  
write('Permutations: '), n!, maplist(writeLn, Result),  
menu.
```

Process 5):-

```
write('Goodbye!'), n!,
```

% Main Predicate

main :-

menu.

Process 6):-

```
write('Enter list').
```

```
read(L)
```

```
write('Enter element to be searched / checked').
```

```
read(L1)
```

```
list_member (L, L1, Result).
```

```
write (Result).
```

menu.

Process 6):-

```
write('Goodbye!'), n!
```

% Main Predicate

main :-

menu

Output

Menu

1. Concatenate lists
2. Add element
3. Delete element
4. Print list
5. Check member
6. Quit.

Enter the number of your choice : 1

Enter the first list : [1, 2, 3, 4]

Enter the second list : [5, 6, 7, 8]

Concatenated list : [1, 2, 3, 4, 5, 6, 7, 8].

Enter the first number of your choice : 2

Enter an element : 2

Enter a list : 4, 5, 6

List after adding element : [4, 5, 6, 2]

Enter the number of your choice : 3

Enter an element : 5

Enter a list : 4, 5, 6, 2

List after deleting element : [4, 6, 2]

Termwork

PROBLEM STATEMENT - Design an algorithm for TO implement depth first search and develop a PROLOG program for the same.

Theory

Depth - first search or DFS algorithm is a recursive algorithm that uses the backtracking principle. It entails conducting exhaustive searches of all nodes by moving forward if possible & backtracking, if necessary. To visit the next node, pop the top node from the stack & push all of its nearby nodes into a stack. Topological sorting, Scheduling problems, graph cycle detection, & solving puzzles with just one solution, such as a maze or a Sudokus puzzle.

A standard DFS implementation puts each vertex of the graph into one of two categories:

1. visited.
2. not visited.

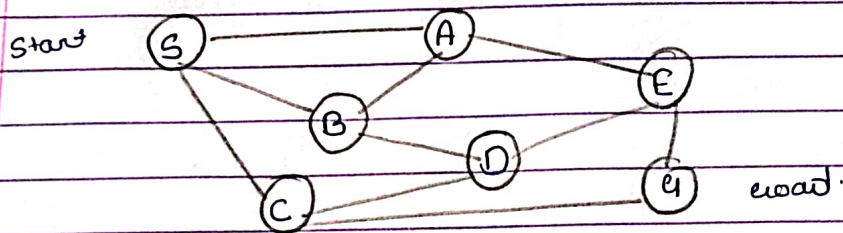
The purpose of the algorithm is to mark each vertex as visited while avoiding cycles.

The DFS algorithm works as follows.

1. Start by putting any one of the graph's vertices on top of a stack.

2. Take the top item of the stack & add it to the visited list.
3. Create a list of that vertex's adjacent nodes. Add the ones which aren't in the visited list to the top of the stack.
4. Keep repeating steps 2 & 3 until the stack is empty.

Example



open	closed
(S, nil)	()
(A, S), (B, S), (C, S)	(S, nil)
(E, A), (D, B), (C, S)	(A, S), (C, S, nil)
(D, E), (E, E), (B, S)	(E, A), (A, S), (S, nil)
(C, S)	
(E, E), (B, S), (C, S)	(D, E), (E, A), (A, S), (S, nil)

path is $u \rightarrow E, E \rightarrow A, A \rightarrow S$
 $u \rightarrow E \rightarrow A \rightarrow S$
 $S \rightarrow A \rightarrow E \rightarrow S$

Algorithm

DepthFirstSearch()

open \leftarrow start NIL

closed \leftarrow {}

while not null(open)

do nodepair \leftarrow head(open)

node \leftarrow head(nodepair)

if isGoal(node) = TRUE

then return ReconstructPath(nodepair, closed)

else closed \leftarrow cons(nodepair, closed)

children \leftarrow children(node)

noLoops \leftarrow RemoveSeen(children, open, closed)

new \leftarrow MakePairs(noLoops, node)

open \leftarrow Append(new, Tail(open))

return "No solution found"

RemoveSeen(nodepair, openList, closedList)

if Null(NodeList)

then return {}

else n \leftarrow head(NodeList)

if (GoalTest(n, openList) OR GoalTest(n, closedList))

then return RemoveSeen(Tail(NodeList), openList, closedList)

else return cons(n, RemoveSeen(Tail(NodeList), openList, closedList))

GoalTest(node, listofPairs)

if Null(listofPairs)

then return FALSE

else if $n = \text{Head}(\text{Head}(\text{list of Pairs}))$

then return TRUE.

else return recurIn(nodes, Tail(list of Pairs))

MakePairs(list, parent)

if Null(list)

then return ()

else return cons(MakeList(Head(list), parent),
MakePairs(Tail(list), parent)).

PROGRAM

child (S, E)

child (S, A)

child (S, B)

child (B, A)

child (A, E)

child (C, D)

child (D, E)

child (E, e)

path (A, e, [A|Z]);- /* to find the path from root to leaf */

childnode (A, e, z).

childnode (A, e, [e]);- /* to determine whether a node is child of
other */

child (A, e).

childnode (A, e, [x|L]);-

child (A, x).

childnode (x, e, L).